Efficient Problem Solving: Using the Right Data Structures and the Right Algorithms



Part 1: Looking Back

- Why do we study data structures and algorithms?
- ■Why do we study asymptotic analysis?

Problem Solving: Good/Better/Best Algorithms.

- **☐** How to verify your implementation?
- ☐ Implement algorithms the correct/wrong way?

- **Example:** Using the right algorithm
- If the goal is to sort an almost sorted array with efficient worst-case time complexity, use insertion sort.
- This is decided by the inversion number of the input sequence.

- Example: Using the right algorithm
- If the goal is to sort a random array with efficient worst-case time complexity, use merge-sort.
- If the goal is to sort a random array with efficient average-case time complexity, use quicksort.

- Example: Using the right data structure
- If the goal is to support insertion/search/deletion with good worst-case time complexity, use the red-black tree.
- **■** What about a linked list?
- **■** What about a binary search tree?
- What about a hash table?

- Example: Using the right data structure
- If the goal is to support insertion/deletemin/decreasekey with good worst-case time complexity, use red-black tree or min-heap.
- **■** Min-heap is very easy to implement.
- RBT is much harder to implement.

Why Study Asymptotic Analysis?

- ☐ It abstracts the running time in a simple way!
- □O(log(n)) running time
- □O(n) running time
- O(n log(n)) running time
- O(n²) running time
- O(n³) running time

□ O(2ⁿ) running time

The Goal Is to Solve a Problem Efficiently

- Suppose you have several algorithms for solving the same problem. The algorithms have worst-cast time complexities given in the following:
- □ A: O(log(n)) running time
- ☐ B: O(n) running time
- C: O(n log(n)) running time
- D: O(n²) running time
- E: O(n³) running time
- F: O(2ⁿ) running time
- Which algorithm is more appealing?

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- E: O(n³) running time
- ☐ F: O(2ⁿ) running time
- Algorithm A is more appealing.



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Part 2: Efficient Problem Solving

- **□**Why do we study data structures and algorithms?
- **□** Why do we study asymptotic analysis?

Problem Solving: Good/Better/Best Algorithms.

- **☐** How to verify your implementation?
- Implement algorithms the correct/wrong way?

- We are given an unsorted array A of size n. Let k be an integer between 1 and n. We need to find the k-th smallest element of A.
- ■What algorithms can you come up with?
- What are the corresponding worst-case time complexities?

- We are given an unsorted array A of size n. Let k be an integer between 1 and n. We need to find the k-th smallest element of A.
- ■What algorithms can you come up with?
- Approach A: Find the smallest element and delete it. Repeat this process k times.
- Time complexity is O(k n).

- We are given an unsorted array A of size n. Let k be an integer between 1 and n. We need to find the k-th smallest element of A.
- ■What algorithms can you come up with?
- \square Approach B: Sort A using mergesort. Return A[k].
- Time complexity is $O(n \log n)$.
- This is better than Approach A.

- We are given an unsorted array A of size n. Let k be an integer between 1 and n. We need to find the k-th smallest element of A.
- What algorithms can you come up with?
- Approach C: Build a min-heap. Delete-Min k times.
- Time complexity is $O(n) + O(k \log n) = O(n + k \log n)$.
- This is better than Approach B.

- We are given an unsorted array A of size n. Let k be an integer between 1 and n. We need to find the k-th smallest element of A.
- ■What algorithms can you come up with?
- Approach D: Use linear time selection.
- \square Time complexity is O(n).
- This is better than Approach C.

- We are given an unsorted array A of size n. Let k be an integer between 1 and n. We need to find the k-th smallest element of A.
- ■What algorithms can you come up with?
- ☐ Can we do better than Approach D?
- No. There are n elements in unsorted order. Any correct algorithm needs to check all n elements. This requires a time complexity of $\Omega(n)$.
- ☐ D is the best approach (asymptotically).

- We are given an unsorted array A of size n. Let k be an integer between 1 and n. We need to return the k smallest elements of A in sorted order.
- ■What algorithms can you come up with?
- What are the corresponding worst-case time complexities?

- We are given an unsorted array A of size n. Let k be an integer between 1 and n. We need to return the k smallest elements of A in sorted order.
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- Approach A: Find the smallest element and delete it. Repeat this process k times.
- Time complexity is O(k n).

- We are given an unsorted array A of size n. Let k be an integer between 1 and n. We need to return the k smallest elements of A in sorted order.
- ■What algorithms can you come up with?
- \square Approach B: Sort A using mergesort. Return A[1] through A[k].
- Time complexity is $O(n \log n + k)$.
- ☐ This is better than Approach A.

- We are given an unsorted array A of size n. Let k be an integer between 1 and n. We need to return the k smallest elements of A in sorted order.
- ■What algorithms can you come up with?
- Approach C: Build a min-heap. Delete-Min k times.
- Time complexity is $O(n)+O(k \log (n))=O(n+k \log n)$.
- This is better than Approach B.

- We are given an unsorted array A of size n. Let k be an integer between 1 and n. We need to return the k smallest elements of A in sorted order.
- What algorithms can you come up with?
- Approach D: Use linear time selection to return the k smallest elements. Then sort them using mergesort.
- Time complexity is $O(n)+O(k \log (k))=O(n+k \log k)$.
- This is better than Approach C.

- We are given an unsorted array A of size n. Let k be an integer between 1 and n. We need to return the k smallest elements of A in sorted order.
- Can we do better than Approach D?
- No. There are n elements in unsorted order. Any correct algorithm needs to check all n elements. This requires a time complexity of $\Omega(n)$. Any algorithm needs to sort at least k elements. This requires a time complexity of $\Omega(k \log k)$. Therefore any algorithm requires $\Omega(n + k \log k)$ time.
- ☐ D is the best approach (asymptotically).



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Part 3: Algorithm Implementation

- ☐ Why do we study data structures and algorithms?
- **■** Why do we study asymptotic analysis?

□ Problem Solving: Good/Better/Best Algorithms.

- How to verify your implementation?
- Implement algorithms the correct/wrong way?

How to Verify Your Implementation

- Suppose you are implementing an algorithm for solving a given problem. How do you know that you have implemented your algorithm correctly?
- You can test your algorithm with carefully designed test cases (instances).

- □ Suppose that you have implemented an algorithm with worst-case time complexity O(n).
- You can run your algorithm on test cases with n=1, 2, 4, 8, 16, 32, 64, ..., and measure the running time for each test case.

How to Verify Your Implementation

- You can run your algorithm on test cases with n=1, 2, 4, 8, 16, 32, 64, ..., and measure the running time for each test case.
- You can compute the ratio T(2n)/T(n) for each of these values of n, where T(n) is the measured running time for the particular test case.
- ■What do you expect to see T(2n)/T(n) as n goes to infinity?
- ☐ Suppose it is bounded by 2, what do you learn?
- ☐ Suppose it approaches 3 what do you learn?

How to Verify Your Implementation

■What happens if the algorithm has worst-case time complexity O(n²)?

■What happens if the algorithm has worst-case time complexity O(log(n))?

■What happens if the algorithm has worst-case time complexity O(n log(n))?

How to Implement Algorithms Wrongly?

- Suppose that you are implementing Dijkstra's shortest path algorithm using a binary heap. The worst-case running time should be O(m+n)log(n)).
- ■You can generate test cases with m=5n, and n doubles each time.

☐ If the running time doubles each time, is it OK?

☐ If the running time quadruples each time, is it OK?

How to Implement Algorithms Wrongly?

- Suppose that you are implementing Dijkstra's shortest path algorithm using a binary heap. The worst-case running time should be O(m+n)log(n)).
- How can you implement it with a worst-case time complexity O(n(m+n))?
- If the running time quadruples each time you double the size of the instance, what do you learn?
- ☐ How can this happen???

More on Data Structures

- You need to write a program to manipulate polynomials: evaluation, addition, subtraction, multiplication.
- **■** Which data structure should you use?
 - Array
 - Singularly linked list
 - What happens if the polynomial is $x^{103} + 3$?

More on Algorithms

- You are given a sequence of integers. Your task is to find the distinct integers in the sequence, in sorted order.
- ☐ Can you use algorithms learned in class to solve this problem?

100, 2, 24, 2, 34, 24, 60, 50, 60



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