Custom Computing: Assessed Coursework

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Question 1

Recurring engin eering costs are the costs that will occur in a repeating fashion during the production, usually involving fabriction. These costs are usually descriped in a per unit form.

Non-recurring engineering cost is the one-time up-front cost for research, design, testing and development of a new product.

As we can see below, the minimum number of units that need to be sold for the ASIC implementation to be cost-effective is 1 million units.

$C_{FPGA} > C_{ASIC} \Rightarrow £2 \times N_{units} > £10^6 + £1 \times N_{units} \Rightarrow N_{units} > 10^6$

Question 2

(a) Diagramatic and symbolic Simulation

Diagram of circuit Q1

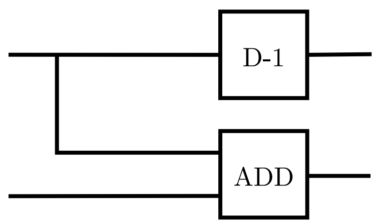

Figure 1: the circuit as derrived from Q1

Diagram of circuit P1

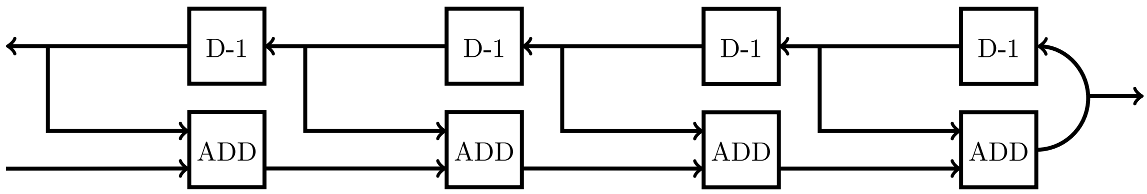

Figure 2: the circuit as derrived from P1

Simulation

The source code of the simulation (uninitialized delay) is the following:

```
INCLUDE "prelude.rby".
P1 n = Q1^n; fork^~1 .
Q1 = snd fork; rsh; [add,D^~1].
current = P1 4.
```

The result after executing `re "a;b;c"`

```
0 - <a,?> ~ ((((a + ?) + ?) + ?) + ?)
1 - <b,?> ~ ((((b + ?) + ?) + ?) + ((((a + ?) + ?) + ?) + ?))
2 - <c,?> ~ ((((c + ?) + ?) + ((((a + ?) + ?) + ?) + ?)) + ((((b + ?) + ?) + ?)
          + ((((a + ?) + ?) + ?) + ?)))
```

The source code of the simulation (initialized delay with 0) is the following:

```
INCLUDE "prelude.rby".
P1 n = Q1^n; fork^~1 .
Q1 = snd fork; rsh; [add,DI 0^~1].
current = P1 4.
```

The result after executing `re "a;b;c"`

```
0 - <a,0> ~ ((((a + 0) + 0) + 0) + 0)
1 - <b,0> ~ ((((b + 0) + 0) + 0) + ((((a + 0) + 0) + 0) + 0))
2 - <c,0> ~ ((((c + 0) + 0) + ((((a + 0) + 0) + 0) + 0)) + ((
            ((b + 0) + 0) + 0) + ((((a + 0) + 0) + 0) + 0)))
```

The result after executing `re -s 4 1"`

```
0 - <1,0> ~ 1
1 - <1,0> ~ 2
2 - <1,0> ~ 4
3 - <1,0> ~ 8
```

The simulation output (for 4 cycles) can be found in the included zip file.

Question 3

(a) Proof by induction

In order to show that $[P, Q]^n; R = R; Q^n$ for $n > 0$, we first have to show that it is $True$ for n = 1.

Base case: $[P, Q]^1; R = R; Q^1$

This is intuitively shown to be true by the given assumption $[P, Q]^n; R = R; Q$ which is equivalent.

Assuming that it is also true for $n = k > 0$

Inductive Hypothesis: $[P, Q]^k; R = R; Q^k$

We need to show that the same is true for $n = k + 1$ and $[P, Q]^{k+1}; R = R; Q^{k+1}$

$[P, Q]^{k+1}; R$ **LHS**

$= [P, Q]^k; [P, Q]; R$ (by sequential expansion of $[P, Q]^{k+1}$)

$= [P, Q]^k; R; Q$ (since $[P, Q]; R = R; Q$ given)

$= R; Q^k; Q$ (by the i.h. $[P, Q]^k; R = R; Q^k$)

$= R; Q^{k+1}$ (by sequential contraction of $Q^k; Q$) **RHS**

So by induction we have **proved** that if $[P, Q]; R = R; Q$ is given to be $True$, for $n > 0$:

$[P, Q]^n; R = R; Q^n$ is also $True$

(b) Inductive Definitions

Right-reduction

$rdr_1 = fst\,[-]^{-1}; R.$

$rdr_{n+1} = fst\,apl_n^{-1}; lsh; snd(rdr_n\ R); R.$

Delta (triangle)

$\Delta_0 = [\,].$

$\Delta_{n+1} = [\Delta_n, R^n]\backslash apr_n.$

(c) Horner's Rule

Left-hand side

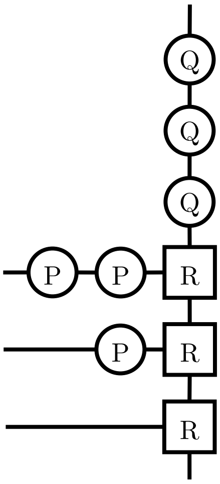

Figure 3: LHS of the rule for n = 3

Right-hand side

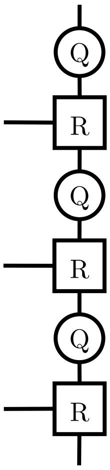

Figure 4: RHS of the rule for n = 3

(d) Polynomial Evaluation

R stands for the add operation (addition), P and Q both stand for multiplication by a constant (let this constant be x). For the given coefficients a_0, a_1, a_2, a_3, the circuit will be the following.

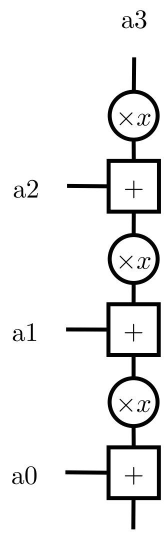

Figure 5: The optimised circuit as adjusted for polynomial evaluation

and the simulation written in Ruby:

```
INCLUDE "prelude.rby".
multc n = pi1^~1;snd n;mult.
Q = multc 'x'.
R = add.
POL n = rdr n (snd Q; R).
current = POL 3.
```

run with `re "a_0 a_1 a_2 a_3"` produces the following output:

```
Simulation start :

    0 - <<a_0,a_1,a_2>,a_3> ~ (a_0 + ((a_1 + ((a_2 + (a_3 * x)) * x)) * x))

Simulation end :
```

Question 4