

# Notes for Tropical Geometry

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## § *Entry 1*

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### 0 Introduction/Motivation

Tropical geometry is the study of discrete structures appearing in limits of polynomial equations.

Course outline:

(1) Hypersurface amoebas, their skeleta, and tropical limits

(2)

### 1 Hypersurface amoebas, their skeleta, and tropical limits

#### 1.1 Laurent polynomial ring

$\mathbb{C}[z_1^{\pm 1}, \dots, z_n^{\pm 1}]$ . Each such Laurent polynomial defines a holomorphic (algebraic) map  $f : (\mathbb{C}^\times)^n \rightarrow \mathbb{C}$  whose zero locus  $V(f) \subseteq (\mathbb{C}^\times)^n$   $f \neq 0$  is a **complex hypersurface**. The ring  $\mathbb{C}[z_1^{\pm 1}, \dots, z_n^{\pm 1}]$  is a unique factorization domain which implies  $f = f_1^{\alpha_1} \cdots f_m^{\alpha_m}$  where the  $f_i$  are irreducible, pairwise different, and hence  $Z(f) = Z(f_1) \cup \dots \cup Z(f_m)$ . This locus is *always* a complex submanifold, even in the case of the nodal cubic for instance, of  $\dim_{\mathbb{C}} = n - 1$  outside of a real codimension 2 subset  $Z(f) \cap Z(\partial_1 f) \cap \dots \cap Z(\partial_n f)$ .

#### Example 1.1.

(a)  $V(z + w) \subseteq (\mathbb{C}^\times)^2$  is isomorphic as a  $\mathbb{C}$ -manifold or as an algebraic variety to  $\mathbb{C}^\times$ . The map  $\mathbb{C}^\times \mapsto V(z + w)$  given  $u \mapsto (u, -u)$  parameterizes this curve.

(b)  $V(z + w + 1) \subseteq (\mathbb{C}^\times)^2$  is isomorphic to  $\mathbb{C}^\times \setminus \{0, 1\}$  via the map  $u \mapsto (u, 1 - u)$ .

#### 1.2 The Log Map

Forget phases and use logarithmic coordinates.

$$\text{Log} : (\mathbb{C}^\times)^n \xrightarrow{1.1} \mathbb{R}_{>0}^n \xrightarrow{\log} \mathbb{R}^n$$

given by

$$(z_1, \dots, z_n) \mapsto (|z_1|, \dots, |z_n|) \mapsto (\log |z_1|, \dots, \log |z_n|).$$

**Definition 1.2.** The **Hypersurface amoeba** of  $f \in \mathbb{C}[z_1^{\pm 1}, \dots, z_n^{\pm 1}] \setminus \{0\}$  is

$$\mathcal{A}_f = \text{Log}(V(f)) \subseteq \mathbb{R}^n$$

(Gelfand, Vapranov, Zelevabsky)

#### Example 1.3.

(a)  $f = z + w$

(b)  $f = z + w + 1$

(c)  $f = 1 + 5zw + w^2 - z^2 + 3z^2w - z^2w^2$

(add pictures later) careful to draw these such that the complements of the amoeba are all convex.

### Observations:

- connected cusps of  $\mathbb{R}^n \setminus \mathbb{C}_f$  are convex in  $\dim = 2$ .  $\mathcal{A}_f$  looks like a thickened graph. We'll sketch a proof of a more general result.

Recall:  $\mathcal{U} \subseteq \mathbb{C}$ ,  $f : \mathcal{U} \setminus \{p_1, \dots, p_r\} \rightarrow \mathbb{C}$  are meromorphic with  $m$  poles  $(p_1, \dots, p_r)$  and  $s$  zeros with multiplicity. This implies

$$s - r = \frac{1}{2\pi i} \int_C \frac{f'(z)}{f(z)} dz.$$

This is the argument principle from complex analysis. Appears in the derivative of  $\frac{1}{2\pi i} \int_{S^1} \log |f| dz$ . This appears in the Jensen formula:  $\mathcal{U} \subseteq \mathbb{C}$  an open subset and assume it contains a closed disk of radius  $r$   $\{z \mid |z| \leq r\} = D$ . Important that it includes the boundary. Then if we have a holomorphic function  $f : \mathcal{U} \rightarrow \mathbb{C}$  with zeros of  $f$  in  $D$   $a_1, \dots, a_k$  such that  $0 < |a_1| \leq |a_2| \leq \dots \leq |a_k|$  (with multiplicity) then we have

$$\frac{1}{2\pi i} \int_0^{2\pi} \log |f(re^{i\theta})| d\theta = \log |f(0)| + \sum_{j=1}^k \log \frac{r}{|a_j|}.$$

This is the Jensen formula.

**Proof.** (Rudin, "Real and complex analysis")

- (1) Assume  $f$  has no zeros and hence that  $\log |f|$  harmonic. Using the mean value property for harmonic functions (go review Analysis) yields the Jensen Formula.

- (2) For the general case, suppose we have  $|a_1|, \dots, |a_n| < r$ , and then that  $|a_{m+1}|, \dots, |a_k| = r$ . Consider  $g(z) = f(z) \cdot \prod_{j=1}^m \frac{r^2 - \bar{a}_j z}{r(a_j - z)} \prod_{j=m+1}^k \frac{a_j}{a_j - z}$  with no zeros in  $|z| \leq r$ . This implies

$$g(0) = f(0) \cdot \prod_{j=1}^m \frac{r}{a_j}$$

by our first case.

- (3)  $|z| = r$ , so on the boundary, we have

$$\left| \frac{r^2 - a_j z}{r(a_j - z)} \right| = \frac{1}{r} \left| \frac{r^2 \bar{z} - a_j |z|^2}{r(a_j - z)} \right| = \frac{r}{r} = 1$$

$$\implies \log |g(re^{i\theta})| = \log |f(re^{i\theta})| - \sum_{j=m+1}^k \log \overbrace{|1 - e^{i(\theta - \theta_j)}|}^{a_j = re^{i\theta_j}}$$

(4) Lemma:  $\int_0^{2\pi} \log(1 - e^{i\theta}) d\theta = 0$ . These four things together prove the Jensen formula.

□

For  $n > 1$  we define something called the Ronkin function. We have  $f \in \mathcal{O}(\text{Log}^{-1}(\Omega))$ ,  $\Omega \subseteq \mathbb{R}^n$  a (convex) open set. Then the **Ronkin Function** is defined

$$N_f(x) = \left(\frac{1}{2\pi i}\right)^n \int_{\text{Log}^{-1}(x)} \text{Log} |f(z_1, \dots, z_n)| \frac{dz_1}{z_1} \vee \dots \vee \frac{dz_n}{z_n}$$

**Theorem 1.4.** (a)  $N_f$  is a convex  $\mathcal{C}^0$ -function

(b)  $\mathcal{A}_f = \text{Log}(V(f)) \subseteq \Omega$  an Amoeba. For all  $\mathcal{U} \subseteq \Omega$  open, connected  $\mathcal{U} \cap \mathcal{A}_f = \emptyset \iff N_f|_{\mathcal{U}}$  affine linear.

(c)  $x \in \Omega \setminus \mathcal{A}_f \implies \text{grad } N_f(x) = (v_1, \dots, v_n)$ ,

$$v_j = \frac{1}{(2\pi i)^n} \int_{\text{Log}^{-1}(x)} \frac{z_j \partial_j f}{f} \frac{dz_1}{z_1} \vee \dots \vee \frac{dz_n}{z_n}.$$

Picture:  $N_f(x) = \langle \alpha_1, x \rangle + c_1$

**Proof.** (sketch)

(a)  $\log |f|$  is plurisubharmonic (i.e. is subharmonic (i.e. somehow less than harmonic functions on a circle) on each each holomorphic image of a disk). We have the following fact: if  $h : \mathcal{U} \rightarrow \mathbb{R}$  is subharmonic,  $\mathcal{U} \subseteq \mathbb{C}$  a domain containing  $\{|z| \leq R\}$ , then  $\varphi(r) = \int_{|z|=r=\exp(s)} h(x) dz$  is a convex function in  $\log r = s$ . Found this proof in a book of Ronkin called “Introduction to the theory of entire functions,” page 84.

(b) Prove this next time

(c)  $x \in \Omega \setminus \mathcal{A}_f$ . Note:

$$\frac{\partial}{\partial x_j} \log |f| = \frac{1}{2} \frac{\partial}{\partial x_j} \log(f\bar{f}) = \text{Re} \left( z_j \frac{\partial}{\partial z_j} \log f\bar{f} \right) = \text{Re} \left( \frac{z_j \partial_j f}{f} \right).$$

$x \in \Omega \setminus \mathcal{A}_f$  implies that

$$\frac{\partial}{\partial x_j} N_f(x) = \text{Re} \left( \frac{1}{2\pi i} \int_{\text{Log}^{-1}} \frac{z_j \partial_j f}{f} \frac{dz_1}{z_1} \wedge \dots \wedge \frac{dz_n}{z_n} \right).$$

Note: for all  $j$ , we have

$$\gamma_j = \frac{1}{(2\pi i)^n} \int_{\text{Log}^{-1}(x)} \frac{z_j \partial_j f}{f} \frac{dz_1}{z_1} \wedge \dots \wedge \frac{dz_n}{z_n}.$$

This is a locally constant  $n$ -form on  $\mathcal{U} \setminus \mathcal{A}_f$  and is not defined on  $\mathcal{A}_f$  since  $f$  is zero on  $\mathcal{A}_f$ . In fact,

$\gamma_j \in \mathbb{Z} : \frac{1}{2\pi i} \int_{|z_j|=e^{x_j}} \frac{\partial_j f(z)}{f(z)} dz_j \in \mathbb{Z}$  by the argument principle.

Look at Passare, Rullgard “Amoebas, Monge – Ampere, measures and triangulations DMJ 2004”

□

## § Entry 2

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Recall that last time we had  $V(f) \subseteq (\mathbb{C}^\times)^n \xrightarrow{\text{Log}} \mathbb{R}^n$ , and we took  $f \in \mathbb{C}[z_1^\pm, \dots, z_n^\pm]$ . This map has image in  $\mathcal{A}_f \subseteq \mathbb{R}^n$ . Recall also that the complement of the amoeba decomposes as the following union of connected components.

$$\mathbb{R}^n \setminus \mathcal{A}_f = \Omega_1 \cup \dots \cup \Omega_k.$$

These connected components correspond to integral points of the Newton polyhedron  $\text{conv}\{I \mid a_I \neq 0\}$  where  $f = \sum_{\text{finite}} a_I z^I$ . Ronkin function is

$$N_f(x) = \frac{1}{(2\pi i)^n} \int_{\text{Log}^{-1}(x)} \text{Log} |f(x)| \frac{dz_1}{z_1} \wedge \dots \wedge \frac{dz_n}{z_n}$$

is convex on  $\mathbb{R}^n$  and is **affine linear on each**  $\Omega_i$  which then implies that each  $\Omega_i$  is convex.

Note:  $\mathcal{U} = \text{Log}^{-1}(\Omega)$ , where  $\Omega$  is open, connected is a **circular domain**, i.e. change the argument of an element in the set and you're still in the set. These are called **Reinhardt domains**.

It is a fact that  $\mathcal{U}$  is a domain of holomorphy if and only if  $\Omega$  is convex. Laurent series converge on  $\text{Log}^{-1}(\Omega)$  since  $\Omega$  is convex.

**Corollary 1.5.**  $\text{Log}^{-1}(\Omega_i)$  are the domains of convergence of the Laurent series expansions of  $f$ .

### 1.3 The spine of a hypersurface amoeba

Let  $\varphi_i = N_f|_{\text{Log}^{-1}(\Omega_i)} = \langle \alpha_i, \cdot \rangle + c_i$  with  $\alpha_i \in (\mathbb{R}^n)^*$  and  $c_i \in \mathbb{R}$  be the piecewise affine approximation of  $N_f$ . Define

$$\varphi = \max\{\varphi_i\}.$$

Note that whenever  $N_f$  is convex we get that  $\varphi \leq N_f$ . **CHECK THIS, SWAPPED FROM MIN TO MAX, CHECK THIS INEQUALITY REMAINS SAME**

**Definition 1.6.**

$$\begin{aligned} \varphi_f &:= \{x \in \mathbb{R}^n \mid \varphi \text{ not affine linear near } x\} \\ &= \{x \in \mathbb{R}^n \mid \varphi \text{ not differentiable at } x\} \\ &= \{x \in \mathbb{R}^n \mid \exists i \neq j \text{ s.t. } \varphi_i(x) = \varphi_j(x) = \max_k \{\varphi_k(x)\}\} \end{aligned}$$

is called the **spine** of  $\mathcal{A}_f$ .

**Theorem 1.7.** [(Passare, Rullgard)]

- (a)  $\varphi_f$  is the  $(n-1)$ -skeleton of a face-fitting decomposition of  $\mathbb{R}^n$  into convex (with integrally defined facets) polyhedra.
- (b)  $\mathcal{A}_f$  deformation retracts onto  $\varphi_f$ .

This notation is slightly confusing to me –  $\varphi_f$  is a subset of the graph of  $\varphi_f$ , it is not itself a function.

## 1.4 Tropical Limits and Maslov “dequantization”

$(\mathbb{R}_{>0}, +, \cdot) \xrightarrow{h \cdot \log = \log_t} (\mathbb{R}, \oplus_h, \odot_h)$  is a semiring isomorphism. The inverse is  $(\mathbb{R}_{>0}, +, \cdot) \xleftarrow{\exp(x/h) \leftarrow x} (\mathbb{R}, \oplus_h, \odot_h)$  with

$$\begin{aligned} x \oplus_h y &= h \cdot \log \left( \exp \left( \frac{x}{h} \right) + \exp \left( \frac{y}{h} \right) \right) \xrightarrow{h \rightarrow 0} \max\{x, y\} \\ x \odot_h y &= h \cdot \log \left( \exp \left( \frac{x}{h} \right) \cdot \exp \left( \frac{y}{h} \right) \right) = x + y. \end{aligned}$$

Now consider  $f_h \in \mathbb{C}(h)[z_1^\pm, \dots, z_n^\pm]$  e.g.  $\frac{h^2+1}{h}z_1^2 + (h^3 - h^2)z_1z_2^{-1}$ . For all  $h$  we have that

$$\mathcal{A}_n(f_h) = \text{Log}_t(V(f_h)) = h \cdot \mathcal{A}(f_h) \subseteq \mathbb{R}^n$$

are the amoeba for the rescaled Log-map  $\text{Log}_t = h \text{Log}$ . Here’s a theorem from a paper prior to tropical geometry truly kicking off.

**Theorem 1.8.**  $\mathcal{A}_h(f_h)$  converges for  $h \rightarrow 0$  in the Hausdorff distance to the tropical hypersurface  $V(\text{trop}(f_h))$ .

$$f_h = \alpha_1 z^{u_1} + \dots + \alpha_r z^{u_r}, \quad \alpha_i \in \mathbb{C}(h)$$

then

$$\text{trop } f_h = \max\{\langle u_1, - \rangle + c_1, \dots, \langle u_r, - \rangle + c_r\}$$

where  $c_i = \text{val}_0(\alpha_i)$ , order of  $\alpha_i(h)$  at  $h = 0$ .

$$\text{val}_0\left(\frac{h^2+1}{h}\right) = -1, \text{val}_0(h^3 - h^2) = 2.$$

**INCLUDE BOARD WITH HAUSDORFF DISTANCE**

## 2 Tropical Arithmetic

### 2.1 Tropical semiring

**Definition 2.1.**  $(\mathbb{R} \cup \{\infty\}, \oplus, \odot)$  is the tropical semiring or the min-plus algebra. We set

- $x \oplus y := \min\{x, y\}$
- $x \odot y := x + y$ .

Both operations are commutative, associative, and are together distributive.

We have the following identities:

- $x \odot (y \oplus z) = x \odot y \oplus x \odot z$
- $x \oplus \infty = x$

- $x \oplus 0 = \begin{cases} 0 & x \geq 0 \\ x & x < 0 \end{cases}$
- $x \odot 0 = x$
- $x \odot \infty := \infty$

Explanation:

$$\begin{aligned}
 (x \oplus y)^3 &= (x \oplus y) \odot (x \oplus y) \odot (x \oplus y) \\
 &= 3 \min\{x, y\} \\
 &= \min\{3x, 3y\} = x^3 \oplus y^3 \\
 &= \min\{3x, 2x + y, x + 2y, 3y\} = x^3 \oplus x^2y \oplus xy^2 \oplus y^3
 \end{aligned}$$

Noting that  $x^3 = 0 \odot x^3$ ,  $x^2y = 0 \odot x^2y$ , etc. we see that these are the coefficients of Pascal's triangle in tropical land, and that the coefficients are all 0. Hence the tropical Pascal triangle is just a bunch of 0's.

## 2.2 Linear algebra

The usual operations (formally) make sense over  $(\mathbb{R} \cup \{\infty\}, \oplus, \odot)$ , e.g.

$$\begin{aligned}
 (u_1, u_2, u_3) \cdot (v_1, v_2, v_3)^T &= u_1 \odot v_1 \oplus u_2 \odot v_2 \oplus u_3 \odot v_3 \\
 &= \min\{u_1 + v_1, u_2 + v_2, u_3 + v_3\}.
 \end{aligned}$$

$$(u_1, u_2, u_3)^T \odot (v_1, v_2, v_3) = \begin{pmatrix} u_1 \odot v_1 & u_1 \odot v_2 & \dots \\ u_2 \odot v_2 & \dots & \\ & & u_3 \odot v_3 \end{pmatrix}$$

**Definition 2.2.** Matrices that can be written as  $u^t \odot v$  have **tropical rank 1**.

**Definition 2.3.** The Barviahok rank of  $A \in M(m \times n, \mathbb{R})$  is  $\min\{k \mid \exists u_1, \dots, u_k, v_1, \dots, v_k, A = u_1^T \odot v_1 \oplus \dots \oplus u_k^T \odot v_k\}$ .

There are other notions of rank: Kapronov rank, tropical rank [MLS, S.5.3].

Looking at **tropical linear systems**  $A \odot x = b$  has applications in engineering, dynamic programming (optimization via recursive structures, e.g. Find a shortest (weighted) path through a directed graph) etc. More on this in section 3.

## 2.3 Tropical Polynomials

**Definition 2.4.** A **Tropical polynomial** is a Laurent polynomial over  $x_1, \dots, x_n$ , i.e. is a function on  $\mathbb{R}, \oplus, \odot)^n$ . A monomial is

$$x_1^{u_1} \odot x_2^{u_2} \cdot \dots \cdot x_n^{u_n}$$

### § Entry 3



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Recall that a tropical polynomial  $f = a_1 \odot x^{u_1} \oplus \dots \oplus a_n \odot x^{u_n}$  is a concave piecewise affine function

$$p(x) = \min\{\langle u_1 \rangle + a_1, \dots, \langle u_n \rangle + a_n\}.$$

**Example 2.5.**  $p = a \odot x^3 \oplus b \odot x^2 \oplus c \odot x \oplus d = \min\{3x + a, 2x + b, x + c, d\}$ . We say that the linear breaks of this graph are the vanishing points of  $p$ .

**Lemma 2.6.** For any concave, piecewise affine function with  $\mathbb{Z}$ -derivatives  $p : \mathbb{R}^n \rightarrow \mathbb{R}$  there exists a tropical polynomial  $f$  with  $p(x) = (x \mapsto f(x) \text{ in } (\mathbb{R}, \oplus, \odot))$ .

Note:  $f = \oplus a_I \odot x^I$  is only unique if we assume that for each  $I$  with  $a_I \neq \infty$  we have that the map  $x \mapsto \langle I, x \rangle + a_I$  agrees with  $p$  in a neighborhood of  $x \in \mathbb{R}^n$ .

**Exercise 2.7** (Tropical Fundamental Theorem of Algebra). Every PA function  $p : \mathbb{R} \rightarrow \mathbb{R}$  with integral derivatives (constant derivatives which are integers) can be written uniquely as a minimal product of tropical linear functions  $a \odot x$ .

**Example 2.8** (Example of Tropical FTA Decomposition). Take  $f = x^2 \oplus 17 \odot x \oplus 2$ . We then have

$$\begin{aligned} f &= x^2 \oplus 17 \odot x \oplus 2 \\ &= \min\{2x, x + 17, 2\} \\ &= \min\{2x, x + 1, 2\} \\ &= (x \oplus 1) \odot (x \oplus 1) \end{aligned}$$

Unique factorization fails for  $n > 1$ .

**Example 2.9.** Take  $f(x, y) = (x \oplus 0) \odot (y \oplus 0) \odot (x \odot y \oplus 0) = (x \odot y \oplus x \oplus 0) \odot (x \odot y \oplus y \oplus 0)$ . The Newton polygon of  $p = a_1 x^{u_1} \oplus \dots \oplus a_r x^{u_r}$  gives us

$$\text{Newt}(p) = \text{conv}\{\underline{u}_i \mid a_i \neq \infty\}$$

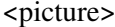
Here:  $f = x^2 y^2 \oplus x^2 y \oplus x y^2 \oplus x y \oplus x \oplus y \oplus 0$ .

### 3 Dynamic Programming

#### 3.1 Shortest paths in graphs

If  $G$  is a directed graph with  $n$  nodes  $1, 2, \dots, n$  and directed edges  $(i, j)$  have a weight  $d_{ij} \in \mathbb{R}_{\geq 0}$  with  $d_{ii} = 0$ . We say  $d_{ij} = \infty$  if there is no edge from  $i$  to  $j$ . We can conveniently present these distances in an  $n \times n$  **adjacency matrix** in the extended reals, i.e

$$D_G = (d_{ij})_{i,j=1,\dots,n} \in M(n \times n, \mathbb{R} \cup \{\infty\}).$$

**Example 3.1.** 

$$D_G = \begin{pmatrix} 0 & 3 & 1 & \infty \\ 1 & 0 & \infty & 3 \\ 1 & 2 & 0 & 0 \\ \infty & 1 & 1 & 0 \end{pmatrix}.$$

**Proposition 3.2.** The shortest length of a path from  $i$  to  $j$  is

$$(ij)\text{-entry of } D_G^{\otimes(n-1)} = \overbrace{D_G \odot \dots \odot D_G}^{(n-1)\text{-times}}.$$

*Proof:* We have that

$$d_{ij}^r := \min\{(\text{weighted}) \text{ length of a path from } i \text{ to } j \text{ with } \leq r \text{ edges}\}.$$

We have that  $d_{ij}^{(1)} = d_{ij}$ . If  $d_{ij} \geq 0$ , then a shortest path in the number of edges runs through each node at most once (otherwise, reverse the loop from  $i$  to  $i$  to arrive at a shorter path).

This implies that  $d_{ij}^{(n-1)}$  = length of shortest weighted path from  $i$  to  $j$ . Recursively this gives

$$\begin{aligned} d_{ij}^{(r)} &= \min\{d_{ik}^{(r-1)} + d_{kj} \mid k \in \{1, \dots, n\}\} \\ &= d_{i1}^{(r-1)} \odot d_{1j} \oplus \dots \oplus d_{in}^{(r-1)} \odot d_{nj} \\ &= \begin{pmatrix} d_{i1}^{(r-1)} & \dots & d_{in}^{(r-1)} \end{pmatrix} \odot \begin{pmatrix} d_{1j} \\ \vdots \\ d_{nj} \end{pmatrix} \\ &= d_{ij}^{(r)} = (i, j)\text{-th entry of } D_G^{\odot r}. \end{aligned}$$

□

This can also be viewed as a limit of a quantum computation (Maslov's dequantization). Replace  $D_G$  with a matrix  $A_G(\epsilon)$  where  $A_G(\epsilon)_{ij} = \epsilon^{D_G(i,j)}$ .

### 3.2 Integer Linear Programming

Given  $A = (a_{ij}) \in M(d \times n, \mathbb{N})$  with  $w \in \mathbb{R}^n$  and  $b \in \mathbb{N}^d$ . We'd like to solve the optimization problem  $w \cdot u$  for  $u \in \mathbb{N}^n$  subject to  $Au \leq b$  or  $Au = b$ .

We can simplify this in the following way. For all  $j$ , take  $\sum_i a_{ij} = \alpha$ . Column sums are equal. We then have  $b_1 + \dots + b_d = m\alpha$ , for  $m \in \mathbb{N}$ .

Then:  $Au = b \implies u_1 + \dots + u_n = m$ . Indeed,  $m\alpha = b_1 + \dots + b_d = \sum_{i,j} a_{ij} u_j = \sum_j (\sum_i a_{ij}) u_j = \alpha(u_1 + \dots + u_n)$ .

**Proposition 3.3.**

$$\min\{w \cdot u \mid Au = b\} = \text{coeff of } x_1^{b_1} \oplus \dots \oplus x_d^{b_d}$$

$$\text{in } (w \odot x_1^{a_{12}} \odot)$$

### 3.3 The assignment problem and the tropical determinant