

## Exercise 1

Submit by Wednesday 17/03/21

### Question 1 (30 pts)

Write the elements of the following sets:

- a.  $\{a, b, c\} \times \{2, 3, 4\} = \{(a, 2), (a, 3), (a, 4), (b, 2), (b, 3), (b, 4), (c, 2), (c, 3), (c, 4)\}$
- b.  $\{b, c, \{a\}, \{a, b\}\} - \{\{b, a\}, a, b, d\} = \{c, \{a\}\}$
- c.  $2^{\{1,2\}} \times (\{a, \{b, c\}\} \cup 2^\phi) = \{\phi, \{1\}, \{2\}, \{1,2\}\} * (\{a, \{b, c\}\} \cup \{\phi\}) =$   
 $\{(\phi, a), (\phi, \{b, c\}), (\phi, \phi), (\{1\}, a), (\{1\}, \{b, c\}), (\{1\}, \phi), (\{2\}, a),$
- d.  $(\{2\}, \{b, c\}), (\{2\}, \phi), (\{1,2\}, a), (\{1,2\}, \{b, c\}), (\{1,2\}, \phi)\}$
- e.  $\{\{a, b, c, d\}\} - \{a, b, c, d\} = \{\{a, b, c, d\}\}$
- f.  $2^{\{a,b,c\}} - 2^{\{c,d\}} = \{\phi, (a), (b), (c), (a, b), (a, c), (b, c), (a, b, c)\} -$   
 $\{\phi, (c), (d), (c, d)\} = \{(a), (b), (a, b), (a, c), (b, c), (a, b, c)\}$
- g.  $2^\phi = \{\phi\}$
- h.  $2^{2^\phi} = \{\phi, \{\phi\}\}$
- i.  $\{a, b, \phi\}^2 = \{a, b, \phi\} * \{a, b, \phi\} =$   
 $\{(a, a), (a, b), (a, \phi), (b, a), (b, b), (b, \phi), (\phi, a), (\phi, b), (\phi, \phi)\}$
- j.  $\{1\} \times \phi \times \{1,2\} = \phi$
- k.  $\{a, b\} \times \{\phi\} = \{(a, \phi), (b, \phi)\}$

### Question 2 (10 pts)

The following theorem is proven by a bi-directional inclusion. See below a proof for the 2nd direction. Provide the 1st direction –

**Prove:**  $(L_1 \cup L_2)^* = (L_1^* \circ L_2^*)^*$

**Proof:** We prove by a bi-directional inclusion

**2<sup>nd</sup> direction:**  $(L_1^* \circ L_2^*)^* \subseteq (L_1 \cup L_2)^*$ :

Let  $w \in (L_1^* \circ L_2^*)^*$ .

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This gives that  $w$  can be decomposed into  $w = w_1 w_2 \dots w_k$

Where each  $w_i \in L_1^* \circ L_2^*$ . But each  $w_i$  can be decomposed again into  $w_i = u_i v_i$ ,  $u_i \in L_1^*$ ,  $v_i \in L_2^*$

That gives that  $\forall i \exists i_k \geq 0$  s.t.  $u_i = u_{i_1} u_{i_2} \dots u_{i_k}$  where each  $u_{i_j} \in L_1$

Similarly,  $\forall i \exists i_l \geq 0$  s.t.  $v_i = v_{i_1} v_{i_2} \dots v_{i_l}$  where each  $v_{i_j} \in L_2$

Therefore,  $\exists n$  s.t.  $w = z_1 z_2 z_3 \dots z_n$  where each  $z_i \in L_1$  or  $z_i \in L_2$

This gives (by the definition of Kleen Star) that  $w \in (L_1 \cup L_2)^*$  q.e.d.

1<sup>st</sup> direction:  $(L_1 \cup L_2)^* \subseteq (L_1^* \circ L_2^*)^*$  : **provide a proof.**

Let  $w \in (L_1 \cup L_2)^*$ , we will show that  $w \in (L_1^* \circ L_2^*)^*$ . We decompose  $w$  into  $w = k_1 k_2 \dots k_i$ , and each  $w_i \in L_1 \cup L_2$ . As we get that  $w_i \in L_1$  or  $w_i \in L_2$  (this follows from the definition of union and closure of Kleen). In the case where  $w_i \in L_2$ , we have  $w_i \in L_2^*$ , by the definition of Kleen Star. Since we know that we have  $\varepsilon \in L_1^*$  (by the definition), we see that  $w_i = \varepsilon w_i \in L_1^* \circ L_2^*$ . Similarly, we know that if  $w_i \in L_1$ , then  $w_i \in L_1^*$ , by definition. And from the definition of a Kleen star we also have that  $\varepsilon \in L_2^*$ , and we therefore know that  $w_i = w_i \varepsilon \in L_1^* \circ L_2^*$ . Therefore,  $\exists k$  s.t.  $w = w_1 w_2 w_3 \dots w_k$  where each  $w_i \in L_1^* \circ L_2^*$ . So finally we get by the closure of Kleen Star that  $w \in (L_1^* \circ L_2^*)^*$

### Question 3 (20 pts)

Prove (by a formal and complete proof) or disprove (showing a counter example) the following claims:

a. If  $L_1 \subseteq L_2$  then  $L_1^* \subseteq L_2^*$

True. To show this, let  $L_1 \subseteq L_2$ , and let  $w$  be a word in  $L_1^*$ . From the definition of closure of Kleen, we get that  $w = w_1 w_2 \dots w_k$  where we have  $w_i \in L_1$  for  $1 \leq i \leq n$ . Now for  $w_i \in L_2$  for  $1 \leq i \leq n$ , by the same definition of Kleen closure, we get that  $w \in L_2^*$ . Therefore:  $L_1^* \subseteq L_2^*$  by definition as required.

b.  $(L_1^+ \circ L_2^+) \subseteq (L_1 \circ L_2)^+$

False.

Let  $L_1 = \{a\}$ ,  $L_2 = \{b\}$ . Let  $w = aabbbb$ , then we have that  $w \in L_1^+ \circ L_2^+$ , but we can see that  $w \notin (L_1 \circ L_2)^+$ , since  $aa \notin L_1$ . Therefore,  $(L_1^+ \circ L_2^+) \not\subseteq (L_1 \circ L_2)^+$ .

c.  $L_1 \cap (L_2 \circ L_3) = (L_1 \cap L_2) \circ L_3$

False: Let  $L_1 = \{ab\}$ ,  $L_2 = \{a\}$ ,  $L_3 = \{b\}$ . We can see that we have  $L_1 \cap (L_2 \circ L_3) = \{ab\} \cap \{ab\} = \{ab\}$ , however, we can also see that  $(L_1 \cap L_2) \circ L_3 = (\{ab\} \cap \{a\}) \circ \{b\} = \emptyset \circ \{b\} = \emptyset$ . As  $\emptyset \neq \{ab\}$  trivially, we get that in fact  $L_1 \cap (L_2 \circ L_3) \neq (L_1 \cap L_2) \circ L_3$ .

d.  $(L_1 \circ L_2)^R = L_2^R \circ L_1^R$

True, by the definition of reverse words.

For the first direction: let  $w \in (L_1 \circ L_2)^R$ . We can decompose  $w$  into  $w = (a \circ b)^R$ , where  $a \in L_1$  and  $b \in L_2$ . We can further decompose  $a = a_1 a_2 \dots a_k$ ,  $b = b_1 b_2 \dots b_l$ .

Where  $a \circ b = a_1 a_2 \dots a_k b_1 b_2 \dots b_l$ , therefore  $w = b_l \dots b_2 b_1 \circ a_k \dots a_2 a_1 = b^R \circ a^R$

Therefore we showed that  $(L_1 \circ L_2)^R \subseteq L_2^R \circ L_1^R$

For the second direction: let  $w \in L_2^R \circ L_1^R$ . We can decompose  $w = b^R \circ a^R$  where  $a \in L_1$  and  $b \in L_2$ . Therefore we have that  $w = b_l \dots b_2 b_1 a_k \dots a_2 a_1 = b^R \circ a^R$ . If  $a$  and  $b$  are decomposed as above then we have that  $a = a_1 a_2 \dots a_k$ ,  $b = b_1 b_2 \dots b_l$ . As  $a \circ b = a_1 a_2 \dots a_k b_1 b_2 \dots b_l$ , trivially  $(a \circ b)^R = b_l \dots b_2 b_1 \circ a_k \dots a_2 a_1$ . As we have that  $b_l \dots b_2 b_1 \circ a_k \dots a_2 a_1 = b_l \dots b_2 b_1 a_k \dots a_2 a_1$  and  $w = b_l \dots b_2 b_1 a_k \dots a_2 a_1$ , then  $w = b_l \dots b_2 b_1 \circ a_k \dots a_2 a_1$  so  $w \in (L_1 \circ L_2)^R$  we can see that  $L_2^R \circ L_1^R \subseteq (L_1 \circ L_2)^R$ .

e.  $(L_1 \cup L_2)^* = (L_1^* \circ L_2)^* \cup (L_1 \circ L_2^*)^*$

Counterexample - we can define the following:

$$\Sigma = \{0,1\} \quad L_1 = \{0\} \quad L_2 = \{1\} \quad w = 10$$

We have on the left-hand side that  $w \in (L_1 \cup L_2)^* = \{0,1\}^*$ , but according to the right-hand side,  $w$  cannot be contained in  $(L_1^* \circ L_2)^* \cup (L_1 \circ L_2^*)^*$  since one expression cannot end with 0, and one cannot start with 1, so therefore there is no such equality and  $w$  does not belong.

#### Question 4 (20 pts)

Show concrete examples for the following languages:

a. Languages  $L_1$  and  $L_2$  over  $\{a,b\}$  where  $(L_1 \cup L_2)^* \neq L_1^* \cup L_2^*$

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Let  $L_1 = \{a\}$ ,  $L_2 = \{b\}$ , for the word  $w = abab$ ,  $w \in (L_1 \cup L_2)^*$  but we have that  $w \notin L_1^* \cup L_2^*$ .

b. Two infinite languages  $L_1$  and  $L_2$  where  $L_1 \not\subseteq L_1 \circ L_2$  and  $L_2 \not\subseteq L_1 \circ L_2$

Let  $L_1 = \{a\}^+$ ,  $L_2 = \{b\}^+$ , we define the words  $w_1 = aa$  and  $w_2 = bb$ ,  $w_1 \in L_1$

however, we can see that  $w_1 \notin L_1 \circ L_2$ , which means that trivially  $L_1 \not\subseteq L_1 \circ L_2$ , by definition, similarly we can see that  $w_2 \in L_2$  but that  $w_2 \notin L_1 \circ L_2$ , which means that we have  $L_2 \not\subseteq L_1 \circ L_2$ , by definition.

c. Languages  $L_1$  and  $L_2$  over  $\{a,b\}$  where  $\overline{L_1 \cup L_2} \neq \overline{L_1} \cup \overline{L_2}$

Let  $L_1 = \{a\}$ , and  $L_2 = \{b\}$ , and let our word be  $w = b$ , then we have that  $\overline{L_1 \cup L_2} = \overline{\{a\} \cup \{b\}} = \overline{\{a,b\}} = \{a,b\}^* - \{a,b\}$ , and we can see that  $\overline{L_1} \cup \overline{L_2} = \overline{\{a\}} \cup \overline{\{b\}} = \{a,b\}^*$ , so we get that  $w \in \overline{L_1} \cup \overline{L_2}$  but  $w \notin \overline{L_1 \cup L_2}$ .

d. Language  $L$  over  $\{a,b\}$  where  $L \neq L \circ L$ , but  $L^* = L \circ L$

Let  $L = \{\varepsilon, a, aaa \dots\}$ , then  $L \circ L = \{\varepsilon, a, aa, aaa, \dots\} = L^*$ . But we can see then that  $L \neq L \circ L = L^*$

### Question 5 (20 pts)

Given are the following languages:

$$L_1 = \{\varepsilon\} ; L_2 = \phi ; L_3 = \{ab, b, a\} ; L_4 = \{aa, \varepsilon\} ; L_5 = \{\varepsilon, bb, aa, b\}$$

Define the languages below. No need to explain or prove but give the simplest expression (for example, if  $L = \{1,10\}$  and  $L' = \{0\}$  and  $L'' = (L \cup L')^*$  then  $L'' = \{0,1\}^*$  is a simpler expression than  $L'' = \{0,1,10\}^*$  )

a.  $L_1 \circ L_5 = \{\varepsilon\} \circ \{\varepsilon, bb, aa, b\} = \{\varepsilon, bb, aa, b\}$

b.  $L_2 \circ L_3 = \phi \circ \{ab, b, a\} = \phi$

c.  $L_2^* \circ L_5 = \{\varepsilon\} \circ \{\varepsilon, bb, aa, b\} = \{\varepsilon, bb, aa, b\}$

d.  $L_3^* - L_5 = \{a, b\}^+ - \{bb, aa, b\}$

e.  $(L_3 \cup L_5)^R = \{\varepsilon, a, b, ab, aa, bb\}^R = \{\varepsilon, a, b, ab, aa, bb\}$