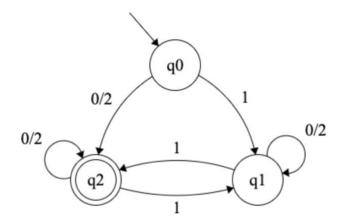
Exercise 3

Submit by Wednesday 14/04/21

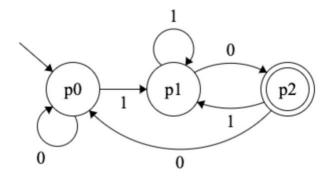
Question 1 (20 pts)

Define **DFA** for the following languages. It is enough to draw a state diagram for each language. For items c,d,e use the product automaton and specify the accepting states.

a. $L_1 = \{w \in \{0,1,2\}^+ \mid \text{ the number of } w \text{ } i \text{n basis 3 is divisible by 2, leading 0's are permitted} \}$

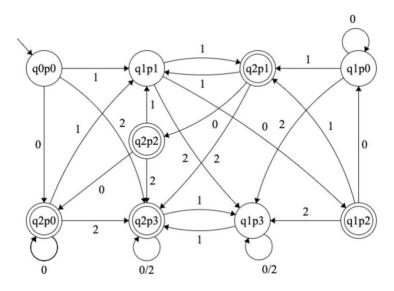


b.
$$L_2 = \{w10 | w \in \{0,1\}^*\}$$

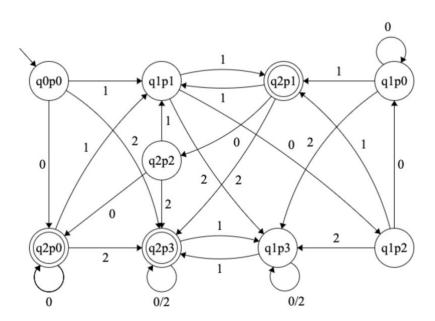


c.
$$L_3 = \{ w \mid w \in L_1 \text{ or } w \in L_2 \}$$

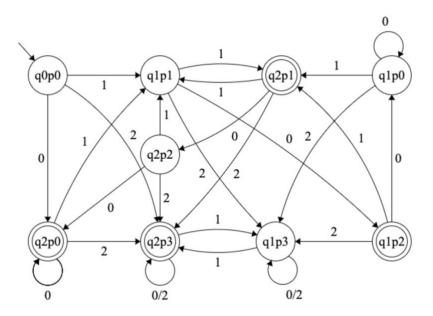
We add a state to L_2 , call it p3, which handles {2} and the general product automaton without the unreachable states is:



d. $L_4 = \{ w \mid w \in L_1 \ and \ w \notin L_2 \}$



e.
$$L_5 = \{ w \mid w \in L_1 - L_2 \text{ or } w \in L_2 - L_1 \}$$

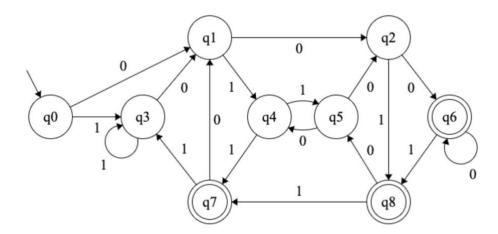


Question 2 (80 pts)

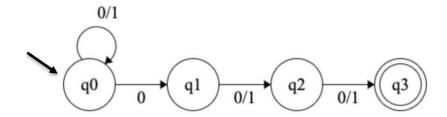
Define NFAs (non-deterministic automaton) for the following languages. It is enough to draw a state diagram for each language. Full credits will be given only for the NFA with the minimal possible states. For items a,b,c draw also the DFAs for these languages.

a. $L = \{w \in \{0,1\}^* \mid \text{the 3rd letter from the end is zero } \}$

DFA:

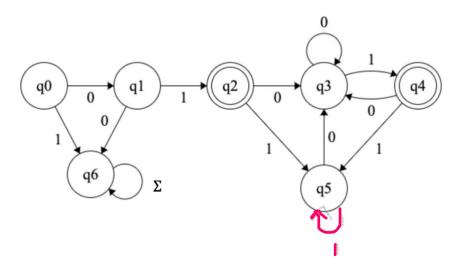


NFA:

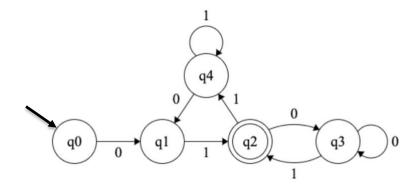


b. $L = \{w \in \{0,1\}^* \mid w \text{ starts with } 01 \text{ and ends with } 01\}$

DFA:

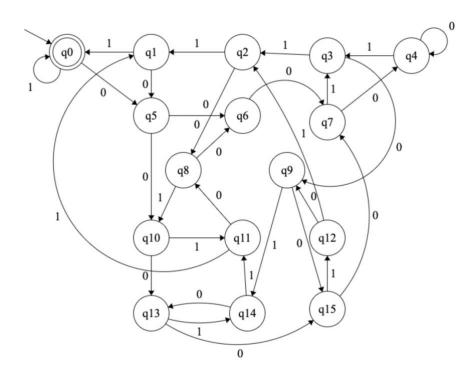


NFA:

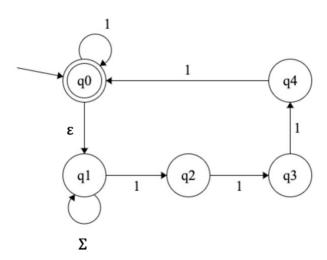


c. $L = \{w \in \{0,1\}^* \mid \text{none of the last 4 letters is 0}\}$

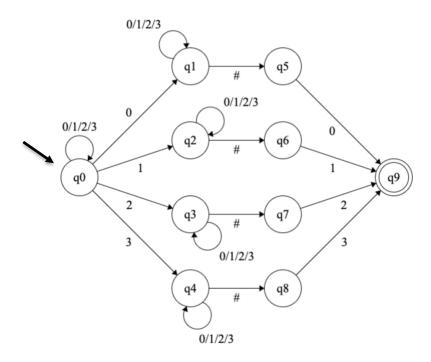
DFA:



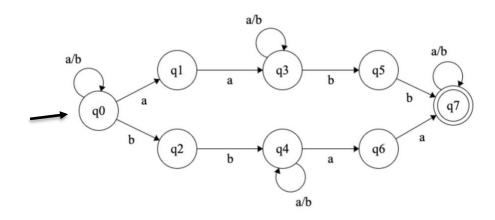
NFA:



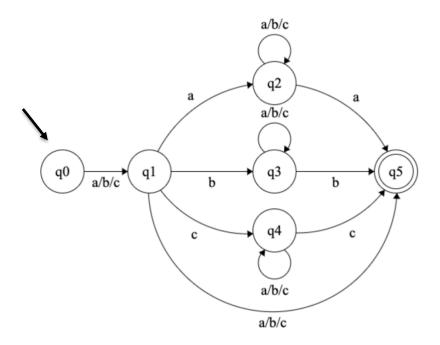
d. $L = \{ w \# \sigma \mid where \ \sigma \ \text{occurs in} \ w, w \in \Sigma^*, \sigma \in \Sigma, \ \Sigma = \{0,1,2,3\} \ \}$, for example, 13031#3 is in L, but not 13031#2.



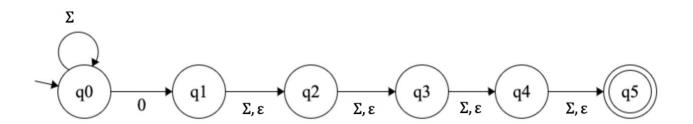
e. $L = \{w \in \{a, b\}^* \mid aa \text{ and } bb \text{ are substrings in } w\}$



f. $L = \{w \in \{a, b, c\}^* \mid \text{where the second letter in } w \text{ is identical to the last letter } \}$



g. $L = \{w \in \{0,1\}^* \mid \text{at least one of the last 5 letters is 0} \}$



h. $L = \{\sigma_1 u \sigma_2 v \sigma_3 | \sigma_1, \sigma_2, \sigma_3 \in \Sigma, u, v \in \Sigma^*, |u| = |v|, \sigma_2 = \sigma_1 \text{ or } \sigma_2 = \sigma_3 \text{ but not both, } \Sigma = \{0,1\} \}$ i.e. the middle letter is similar to the first letter or the last letter, but not both.

