

Hamming code – one error detection and one correction

$$A = [1 \ 0 \ 1 \ 0]$$

m – represents the number of information bits (in our case = 4)

k – represents the number of control bits (in our case = 3)

$$2^k \geq m + k + 1$$

$$k = 3$$

$2^0, 2^1, \dots, 2^{k-1} \Rightarrow 1, 2, 4$ control bits position

$V = [c_1 \ c_2 \ a_3 \ c_4 \ a_5 \ a_6 \ a_7]$ – vector to be transmitted

$a_3 = \text{first information bit} (= 1)$

$a_5 = \text{second information bit} (= 0)$

$a_6 = \text{third information bit} (= 1)$

$a_7 = \text{fourth information bit} (= 0)$

The values of C_1, C_2, C_4 are calculated from:

$$H \cdot V^T = 0$$

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where: H – Hamming matrix (is generated by us)

it has n columns ($m + k$) and k rows

$$H = \begin{bmatrix} 0 & 0 & 0 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 & 0 & 1 & 1 \\ 1 & 0 & 1 & 0 & 1 & 0 & 1 \end{bmatrix}$$

$$\begin{bmatrix} 0 & 0 & 0 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 & 0 & 1 & 1 \\ 1 & 0 & 1 & 0 & 1 & 0 & 1 \end{bmatrix} \cdot \begin{bmatrix} c_1 \\ c_2 \\ a_3 \\ c_4 \\ a_5 \\ a_6 \\ a_7 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix} \Rightarrow \begin{aligned} c_4 \oplus a_5 \oplus a_6 \oplus a_7 &= 0 \\ c_2 \oplus a_3 \oplus a_6 \oplus a_7 &= 0; \quad c_4 = 1; \quad c_2 = 0; \quad c_1 = 1; \\ c_1 \oplus a_3 \oplus a_5 \oplus a_7 &= 0 \end{aligned}$$

The sending station (Tx) will send the V vector/word, and the receiving station will receive the E vector/word.

To determine if the transmission was successful, we will be verifying as follows:

$$H \cdot E^T = Z$$

Z – error vector

If $Z = 0$, then the transmission was successful

If $Z \neq 0$, the transmission had an error at the position given by Z . (transform from binary to decimal the Z vector, and you will get the erroneous position on the E vector).

Examples of error detection and correction:

a. $E = [0 \ 0 \ 1 \ 1 \ 0 \ 1 \ 0]$

$$H \cdot E^T = Z$$

$$Z = \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix} \Rightarrow E_{\text{corr}} = [1 \ 0 \ 1 \ 1 \ 0 \ 1 \ 0]$$

b. $E = [0 \ 1 \ 1 \ 1 \ 0 \ 1 \ 0]$

$$H \cdot E^T = Z$$

$$Z = \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix} \Rightarrow E_{\text{corr}} = [0 \ 1 \ 0 \ 1 \ 0 \ 1 \ 0]$$

correcting the wrong bit

c. $E = [1 \ 0 \ 1 \ 1 \ 0 \ 1 \ 0]$

$$H \cdot E^T = Z$$

$$Z = \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix} \Rightarrow \text{successful transmission, no bit correction}$$

Hamming code – two error detection and one error correction

$$A = [1 \ 0 \ 1 \ 0]$$

m – represents the number of information bits (in our case = 4)

k - represents the number of control bits (in our case = 3)

$$2^k \geq m + k + 1$$

$$k = 3$$

$2^0, 2^1, \dots, 2^{k-1} \Rightarrow 1, 2, 4$ control bits position

$$V = [c_1 \ c_2 \ a_3 \ c_4 \ a_5 \ a_6 \ a_7] - \text{vector to be transmitted}$$

$a_3 = \text{first information bit (= 1)}$

$a_5 = \text{second information bit (= 0)}$

$a_6 = \text{third information bit (= 1)}$

$a_7 = \text{fourth information bit (= 0)}$

c_0 – parity bit

The values of C_0, C_1, C_2, C_4 are calculated from:

$$H_2 \cdot V_2^T = 0$$

where: H_2 – Hamming matrix (is generated by us)

it has n columns ($m + k + 1$) and ($k + 1$) rows

$$H_2 = \begin{bmatrix} 0 & H \\ \vdots & \\ 1 & 1 \end{bmatrix}$$

$$H_2 = \begin{bmatrix} 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 & 0 & 1 & 0 & 1 \\ 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \end{bmatrix}$$

$$H_2 \cdot V_2^T = 0$$

$$\begin{bmatrix} 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 & 0 & 1 & 0 & 1 \\ 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \end{bmatrix} \cdot \begin{bmatrix} c_0 \\ c_1 \\ c_2 \\ a_3 \\ c_4 \\ a_5 \\ a_6 \\ a_7 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

$$c_4 \oplus a_5 \oplus a_6 \oplus a_7 = 0$$

$$c_2 \oplus a_3 \oplus a_6 \oplus a_7 = 0$$

$$c_1 \oplus a_3 \oplus a_5 \oplus a_7 = 0$$

$$; c_4 = 1; c_2 = 0; c_1 = 1; c_0 = 0$$

$$\Rightarrow c_0 \oplus c_1 \oplus c_2 \oplus a_3 \oplus c_4 \oplus a_5 \oplus a_6 \oplus a_7 = 0;$$

The sending station (Tx) will send the V_2 vector/word, and the receiving station will receive the E_2 vector/word.

To determine if the transmission was successful, we will be verifying as follows:

$$H_2 \cdot E_2^T = Z_2$$

Z_2 – error vector $\rightarrow (k + 1)$ rows and one column

$$Z_2 = \begin{bmatrix} Z_c \\ Z_p \end{bmatrix}; Z_c \rightarrow k \text{ rows and one column}; Z_p \rightarrow 1 \text{ row and 1 column}$$

If $Z_c = 0$ and $Z_p = 0 \Rightarrow \nexists \text{ error}$

If $Z_c = 0$ and $Z_p = 1 \Rightarrow \exists \text{ error and it is the parity bit} \Rightarrow E_{corr}$

If $Z_c < > 0$ and $Z_p = 1 \Rightarrow \exists \text{ error and } Z_c \text{ will give the error bit position} \Rightarrow E_{corr}$

If $Z_c < > 0$ and $Z_p = 0 \Rightarrow \exists \text{ two errors and it can't be corrected}$

Examples of error detection and correction:

$$E_2 = [0 \ 1 \ 0 \ 1 \ 1 \ 0 \ 1 \ 0]$$

$$H_2 \cdot E_2^T = Z_2$$

a.

$$Z_2 = \begin{bmatrix} 0 \\ 0 \\ 0 \\ 0 \end{bmatrix}; \text{no errors [successful transmission]}$$

$$E_2 = [1 \ 0 \ 0 \ 1 \ 1 \ 0 \ 1 \ 0]$$

$$H_2 \cdot E_2^T = Z_2$$

b.

$$Z_2 = \begin{bmatrix} 0 \\ 0 \\ 1 \\ 0 \end{bmatrix}; Z_p = 0; Z_c < > 0. \exists \text{two errors}$$

[unsuccessful transmission; too many errors and it can't be corrected]

$$E_2 = [1 \ 1 \ 0 \ 1 \ 1 \ 0 \ 1 \ 0]$$

$$H_2 \cdot E_2^T = Z_2$$

c.

$$Z_2 = \begin{bmatrix} 0 \\ 0 \\ 0 \\ 1 \end{bmatrix}; Z_p = 1; Z_c = 0. \exists \text{one error} \Rightarrow E_{2\text{corr}} = [0 \ 1 \ 0 \ 1 \ 1 \ 0 \ 1 \ 0]$$

[unsuccessful transmission; one error to be corrected]

$$E_2 = [0 \ 1 \ 1 \ 1 \ 0 \ 1 \ 0]$$

$$H_2 \cdot E_2^T = Z_2$$

d.

$$Z_2 = \begin{bmatrix} 0 \\ 1 \\ 0 \\ 1 \end{bmatrix}; Z_p = 1; Z_c < > 0. \exists \text{one error} \Rightarrow Z_c \text{ gives the error position;}$$

$$E_{2\text{corr}} = [0 \ 1 \ 0 \ 1 \ 1 \ 0 \ 1 \ 0]$$

[unsuccessful transmission; one error to be corrected]