

# 1 Organization and Introduction

- The Art of managing complexity
  - Abstraction: Hiding details when they are not important
  - Discipline: Intentionally restricting your design choices to that you can work more productively at higher abstraction levels
  - The three -Y's
    - \* Hierarchy: A system is divided into modules of smaller complexity
    - \* Modularity: Having well defined functions and interfaces
    - \* Regularity: Encouraging uniformity, so modules can be easily re-used
- Bit: **B**inary **d**igit

# 2 Binary Numbers

- Powers of two:
 

|            |             |                  |
|------------|-------------|------------------|
| $2^0 = 1$  | $2^5 = 32$  | $2^{10} = 1024$  |
| $2^1 = 2$  | $2^6 = 64$  | $2^{11} = 2048$  |
| $2^2 = 4$  | $2^7 = 128$ | $2^{12} = 4096$  |
| $2^3 = 8$  | $2^8 = 256$ | $2^{13} = 8192$  |
| $2^4 = 16$ | $2^9 = 512$ | $2^{14} = 16384$ |
- Binary to decimal conversion
 
$$\begin{aligned}
 10011_2 &= 2^4 \times 1 + 2^3 \times 0 + 2^2 \times 0 + 2^1 \times 1 + 2^0 \times 1 \\
 &= 16 \times 1 + 8 \times 0 + 4 \times 0 + 2 \times 1 + 1 \times 1 \\
 &= 16 + 0 + 0 + 2 + 1 = 19_{10}
 \end{aligned}$$
- Convert decimal to binary (roughly). Example with  $47_{10}$  to binary
 

|            |                   |     |   |                 |
|------------|-------------------|-----|---|-----------------|
| $2^6 = 64$ | is $64 \leq 47$ ? | no  | 0 | do nothing      |
| $2^5 = 32$ | is $32 \leq 47$ ? | yes | 1 | $47-32=15$      |
| $2^4 = 16$ | is $16 \leq 15$ ? | no  | 0 | do nothing      |
| $2^3 = 8$  | is $8 \leq 15$ ?  | yes | 1 | $15-8=7$        |
| $2^2 = 4$  | is $4 \leq 7$ ?   | yes | 1 | $7-4=3$         |
| $2^1 = 2$  | is $2 \leq 3$ ?   | yes | 1 | $3-2=1$         |
| $2^0 = 1$  | is $1 \leq 1$ ?   | yes | 1 | $1-1=0$ ; done! |

$\Rightarrow 47_{10}$  to binary is  $0101111_2$
- Binary values and range
  - $N$ -digit decimal number
    - \* How many values:  $10^N$
    - \* Range:  $[0, 10^N - 1]$
    - \* Example (3-digit number):  $10^3 = 1000$  possible values, range:  $[0, 999]$
  - $N$ -bit binary number
    - \* How many values:  $2^N$

- \* Range:  $[0, 2^N - 1]$
- \* Example (3-digit number):  $2^3 = 8$  possible values, range:  $[0, 7] = [000_2 \text{ to } 111_2]$

- Hexadecimal (Base-16) Numbers

| Decimal | Hexadecimal | Binary |  | Decimal | Hexadecimal | Binary |
|---------|-------------|--------|--|---------|-------------|--------|
| 0       | 0           | 0000   |  | 8       | 8           | 1000   |
| 1       | 1           | 0001   |  | 9       | 9           | 1001   |
| 2       | 2           | 0010   |  | 10      | A           | 1010   |
| 3       | 3           | 0011   |  | 11      | B           | 1011   |
| 4       | 4           | 0100   |  | 12      | C           | 1100   |
| 5       | 5           | 0101   |  | 13      | D           | 1101   |
| 6       | 6           | 0110   |  | 14      | E           | 1110   |
| 7       | 7           | 0111   |  | 15      | F           | 1111   |

- Bits, Bytes, Nibbles...



Where MSB=Most significant Bit and LSB=Least significant Bit

- Addition in base two works exactly the same as in base 10, using carries
- Overflow
  - Digital systems operate on a fixed number of bits
  - Addition overflows when the result is too big to fit in the available number of bits
- Signed Binary Numbers
  - Sign/Magnitude Numbers
    - \* 1 sign bit,  $N - 1$  magnitude bits
    - \* Sign bit is the most significant (left-most) bit
    - \* Example: 4-bit sign/mag repr. of  $\pm 6$ :
      - $+6 = \mathbf{0110}$
      - $-6 = \mathbf{1110}$
    - \* Range of an  $N$ -bit sign/magnitude number:  $[-(2^{N-1} - 1), 2^{N-1} - 1]$
    - \* Problems:
      - Addition doesn't work
      - Two representations of 0 ( $\pm 0$ ): 1000 and 0000
      - Introduces complexity in the processor design
  - One's Complement Numbers

- \* A negative number is formed by reversing the bits of the positive number (MSB still indicates the sign of the integer)

| $2^7$ | $2^6$ | $2^5$ | $2^4$ | $2^3$ | $2^2$ | $2^1$ | $2^0$ |     | One's Compl. | Unsigned |
|-------|-------|-------|-------|-------|-------|-------|-------|-----|--------------|----------|
| 0     | 0     | 0     | 0     | 0     | 0     | 0     | 0     | =   | 0            | 0        |
| 0     | 0     | 0     | 0     | 0     | 0     | 0     | 1     | =   | 1            | 1        |
| 0     | 0     | 0     | 0     | 0     | 0     | 1     | 0     | =   | 2            | 2        |
| ...   | ...   | ...   | ...   | ...   | ...   | ...   | ...   | ... | ...          | ...      |
| 0     | 1     | 1     | 1     | 1     | 1     | 1     | 1     | =   | 127          | 127      |
| 1     | 0     | 0     | 0     | 0     | 0     | 0     | 0     | =   | -127         | 128      |
| 1     | 0     | 0     | 0     | 0     | 0     | 0     | 1     | =   | -126         | 129      |
| ...   | ...   | ...   | ...   | ...   | ...   | ...   | ...   | ... | ...          | ...      |
| 1     | 1     | 1     | 1     | 1     | 1     | 0     | 1     | =   | -2           | 253      |
| 1     | 1     | 1     | 1     | 1     | 1     | 1     | 0     | =   | -1           | 254      |
| 1     | 1     | 1     | 1     | 1     | 1     | 1     | 1     | =   | -0           | 255      |

- \* Range of  $n$ -bit number:  $[-2^{n-1} - 1, 2^{n-1} - 1]$ , 8 bits:  $[-127, 127]$
- \* Addition: Done using binary addition with end-around carry. If there is a carry out of the MSB of the sum, this bit must be added to the LSB of the sum

#### – Two's Complement Numbers

- \* Don't have same problems as sign/magnitude numbers:
  - addition works
  - Single representation for 0
- \* Has advantages over one's complement:
  - Has a single 0 representation
  - Eliminates the end-around carry operation required in one's complement addition.
- \* A negative number is formed by reversing the bits of the positive number (MSB still indicates the sign of the integer) and adding 1:

| $2^7$ | $2^6$ | $2^5$ | $2^4$ | $2^3$ | $2^2$ | $2^1$ | $2^0$ |     | Two's Compl. | Unsigned |
|-------|-------|-------|-------|-------|-------|-------|-------|-----|--------------|----------|
| 0     | 0     | 0     | 0     | 0     | 0     | 0     | 0     | =   | 0            | 0        |
| 0     | 0     | 0     | 0     | 0     | 0     | 0     | 1     | =   | 1            | 1        |
| 0     | 0     | 0     | 0     | 0     | 0     | 1     | 0     | =   | 2            | 2        |
| ...   | ...   | ...   | ...   | ...   | ...   | ...   | ...   | ... | ...          | ...      |
| 0     | 1     | 1     | 1     | 1     | 1     | 1     | 1     | =   | 127          | 127      |
| 1     | 0     | 0     | 0     | 0     | 0     | 0     | 0     | =   | -128         | 128      |
| 1     | 0     | 0     | 0     | 0     | 0     | 0     | 1     | =   | -127         | 129      |
| ...   | ...   | ...   | ...   | ...   | ...   | ...   | ...   | ... | ...          | ...      |
| 1     | 1     | 1     | 1     | 1     | 1     | 0     | 1     | =   | -3           | 253      |
| 1     | 1     | 1     | 1     | 1     | 1     | 1     | 0     | =   | -2           | 254      |
| 1     | 1     | 1     | 1     | 1     | 1     | 1     | 1     | =   | -1           | 255      |

- \* Same as unsigned binary, but the most significant bit (MSB) has value of  $-2^{N-1}$ 
  - Most positive 4-bit number: 0111
  - Most negative 4-bit number: 1000
- \* The most significant bit still indicates the sign (1=neg., 0=pos.)
- \* Range of an  $N$ -bit two's comp. number:  $[-2^{N-1}, 2^{N-1} - 1]$ , 8 bits:  $[-128, 127]$

- Increasing bit width (assume from  $N$  to  $M$ , with  $M > N$ ):

- Sign-extension

- \* Sign bit is copied into MSB
- \* Number value remains the same
- \* Give correct result for two's compl. numbers
- \* Example 1:
  - 4-bit representation of  $3 = 0011$
  - 8-bit sign-extended value: **00000011**
- \* Example 2:
  - 4-bit representation of  $-5 = 1011$
  - 8-bit sign-extended value: **11111011**

- Zero-extension

- \* Zeros are copied into MSB
- \* Value will change for negative numbers
- \* Example 1:
  - 4-bit value:  $0011_2 = 3_{10}$
  - 8-bit zero-extended value: **00000011**<sub>2</sub> =  $3_{10}$
- \* Example 2:
  - 4-bit value:  $1011_2 = -5_{10}$
  - 8-bit zero-extended value: **00001011**<sub>2</sub> =  $11_{10}$

### 3 Short Introduction to Electrical Engineering (EE Perspective)

- The goal of circuit design is to optimize:
  - Area: Net circuit area is proportional to the cost of the device
  - Speed/Throughput: We want circuits that work faster, or do more
  - Power/Energy
    - \* Mobile devices need to work with a limited power supply
    - \* High performance devices dissipate more than  $100\text{W}/\text{cm}^2$
  - Design time
    - \* Designers are expensive
    - \* The competition will not wait for you
- (Frank's) Principles for engineering
  - Good engineers are lazy: They do not want to work unnecessarily, be creative
  - They know how to ask the question “why”? : take nothing for granted
  - Engineering is not a religion: Use what works best for you
  - Keep it simple and stupid: Engineers' job is to manage complexity

- Building blocks for microchips
  - Conductors: Metals (Aluminium, Copper)
  - Insulators: Glass ( $\text{SiO}_2$ ), Air
  - Semiconductors: Silicon (Si), Germanium (Ge)
- N-type Doping: Add extra electron (negatively charged), zone becomes negatively charged
- P-type Doping: Remove electron, zone becomes positively charged
- Semiconductors:
  - You can “Engineer” its properties, i.e.
    - \* Make it P type by injecting type-III elements (b, Ga, In)
    - \* Make it N type by injecting elements from type-V (P, As)
  - You can combine P and N regions to each other, from a pure semiconductor
  - Allows you to make interesting electrical devices (Diodes, Transistors, Thyristors)
- pMOS is a P type transistor, nMOS an N type transistors; combined they are a CMOS
- CMOS (Properties)
  - No input current: Capacitive input, no resistive path from the input
  - No current when output is at logic levels: Little static power, current is needed only when switching
  - Electrical properties determined directly by geometry: A transistor that is 2 times larger drives twice the current
  - Very simple to manufacture: pMOS and nMOS can be manufactured on the same substrate
- CMOS Gate Structure
  - The general form used to construct any inverting logic, such as: NOT, NAND, NOR
    - \* The networks may consist of transistors in series or parallel
    - \* When transistors are in parallel, the network is ON if either transistor is ON
    - \* When transistors are in series, the network is ON only if all transistors are ON
  - In a proper logic gate: One of the networks should be ON and the other OFF at any given time
  - Use the rule of conduction complements:
    - \* When nMOS transistors are in series, the pMOS transistor must be in parallel

Maybe add a definition or a better explanation


- \* When nMOS transistors are in parallel, the pMOS transistors must be in series

Add picture on slide 34, 03 - EEPerspective

- Logic Gates


- Perform logic functions: Inversion (NOT), AND, OR, NAND, NOR, etc.
- Single input: NOT gate, buffer
- Two-input: AND, OR, XOR, NAND, NOR, XNOR

Buffer




| A | Z |
|---|---|
| 0 | 0 |
| 1 | 1 |

AND




| A | B | Z |
|---|---|---|
| 0 | 0 | 0 |
| 0 | 1 | 0 |
| 1 | 0 | 0 |
| 1 | 1 | 1 |

OR




| A | B | Z |
|---|---|---|
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 1 |

XOR




| A | B | Z |
|---|---|---|
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 0 |

Inverter




| A | Z |
|---|---|
| 0 | 1 |
| 1 | 0 |

NAND




| A | B | Z |
|---|---|---|
| 0 | 0 | 1 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 0 |

NOR



| A | B | Z |
|---|---|---|
| 0 | 0 | 1 |
| 0 | 1 | 0 |
| 1 | 0 | 0 |
| 1 | 1 | 0 |

XNOR



| A | B | Z |
|---|---|---|
| 0 | 0 | 1 |
| 0 | 1 | 0 |
| 1 | 0 | 0 |
| 1 | 1 | 1 |

- Multiple-Input:
  - \* 3, 4, or even more input AND, OR, XOR gates
  - \* Compound gates
    - AND-OR
    - OR-AND
    - AND-OR-INVERT
    - OR-AND-INVERT
  - \* Other cells: Multiplexers and Adders

- Logic Levels

- Define ranges of discrete voltages to represent 1 and 0 (i.e. 0 for ground and 1 for 5V ( $V_{DD}$ )) and allow for noise.

- Noise: Is anything that degrades the signal (i.e. resistance, power supply noise, etc.)

- Moore's Law

- “Number of transistors that can be manufactured doubles roughly every 18 months.” - Gordon Moore, 1965

- How do we keep Moore's Law:

- Manufacturing smaller structures: some structures are already a few atoms in size
- Developing materials with better properties
- Optimizing the manufacturing steps
- New technologies
- Power consumption
  - Power = Energy consumed per unit time
  - Two types of power consumption:
    1. Dynamic power consumption: Power to charge transistor gate capacitances
 
$$P_{\text{dynamic}} = \frac{1}{2}CV_{DD}^2f$$
    2. Static power consumption: Power consumed when no gates are switching, caused by the leakage current
 
$$P_{\text{static}} = I_{DD}V_{DD}$$

## 4 Combinational Circuits: Theory

- Circuit elements. A circuit consists of:
  - Inputs
  - Outputs
  - Nodes (wires): Connections between I/O and circuit elements. To count them, look at
    - \* Outputs of every circuit elements
    - \* Inputs to the entire circuit
  - Circuit elements
- Types of Logic Circuits
  - Combinational Logic
    - \* Memoryless
    - \* Outputs determined by current values of inputs
    - \* In some books called Combinatorial Logic
  - Sequential Logic
    - \* Has Memory
    - \* Outputs determined by previous and current values of inputs
- Rules of Combinational Composition
  - Every circuit element is itself combinational
  - Every node of the circuit is either
    - \* Designated as an input to the circuit
    - \* Connects to exactly one output terminal of a circuit element

- The circuit contains no cyclic paths: Every path through the circuit visits each node at most once
- Boolean Equations<sup>1</sup>

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<sup>1</sup>For a more in depth look, use the material from Diskrete Mathematik