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Mata kuliah : Kriptografi

KSA (Key Scheduling Algorithm)

Inisialisasi : $S_0 = S_1 \dots S_{255} = 255$

key = Saputra \rightarrow length key = 8

Iterasi ke-0

$i = 0 \quad j = 0 \quad s = 115$

$$\begin{aligned} j &= (j + s[i] + k[(i \bmod \text{len}(k))]) \bmod 256 \\ &= (0 + 0 + k[0 \bmod 8]) \bmod 256 \\ &= (0 + k[0]) \bmod 256 \\ &= (0 + 115) \bmod 256 \\ &= 115 \bmod 256 \end{aligned}$$

$j = 115$

swap = $s[i], s[j] = s[0], s[115]$

$S = 115, 2, 3, 11, 5, 6, 7, \dots, 114, 0, 116, \dots, 255$

Iterasi ke-1

$i = 1 \quad j = 213 \quad P = 112$

$$\begin{aligned} j &= (j + s[i] + k[(i \bmod \text{len}(k))]) \bmod 256 \\ &= (213 + 2 + k[2 \bmod 8]) \bmod 256 \\ &= (215 + k[2]) \bmod 256 \\ &= (215 + 112) \bmod 256 \\ &= (327 \bmod 256) \Rightarrow j = 71 \end{aligned}$$

swap = $s[i], s[j] = s[1], s[71]$

$S = 115, 213, 71, 3, 4, 5, \dots, 70, 2, 72, \dots, 114, 0, 116, \dots, 212, 1, 114, \dots, 255$

Iterasi ke-2

$i = 2 \quad j = 191 \quad t = 116$

$$\begin{aligned} j &= (j + s[i] + k[(i \bmod \text{len}(k))]) \bmod 256 \\ &= (191 + 3 + k[3 \bmod 8]) \bmod 256 \\ &= (194 + k[3]) \bmod 256 \\ &= (194 + 116) \bmod 256 \\ &= 310 \bmod 256 \end{aligned}$$

$j = 55$

swap = $s[i], s[j] = s[2], s[55]$

$S = 115, 213, 71, 191, 55, 5, \dots, 54, 4, 56, \dots, 70, 2, 72, \dots, 114, 0, 116, \dots, 190, 3, 192, \dots, 212, 1, 214, \dots, 255$

$S = 115, 213, 71, 191, 55, 5, \dots, 54, 4, 56, \dots, 70, 2, 72, \dots, 114, 0, 116, \dots, 190, 3, 192, \dots, 212, 1, 214, \dots, 255$

Iterasi ke-3

$i = 3 \quad j = 115 \quad a = 97$

$$\begin{aligned} j &= (j + s[i] + k[(i \bmod \text{len}(k))]) \bmod 256 \\ &= (115 + 4 + k[4 \bmod 8]) \bmod 256 \\ &= (119 + k[4]) \bmod 256 \\ &= (119 + 97) \bmod 256 \\ &= 216 \bmod 256 \end{aligned}$$

$j = 213$

swap = $s[i], s[j] = s[3], s[213]$

$S = 115, 213, 3, 4, 5, \dots, 114, 0, 116, \dots, 212, 1, 214, \dots, 225$

Iterasi ke-4

$i = 4 \quad j = 71 \quad u = 117$

$$\begin{aligned} j &= (j + s[i] + k[(i \bmod \text{len}(k))]) \bmod 256 \\ &= (71 + 5 + k[5 \bmod 8]) \bmod 256 \\ &= (76 + k[5]) \bmod 256 \\ &= (76 + 117) \bmod 256 \\ &= 193 \bmod 256 \end{aligned}$$

$j = 191$

swap = $s[i], s[j] = s[4], s[191]$

$S = 115, 213, 71, 191, 4, 5, \dots, 70, 2, 73, \dots, 114, 0, 116, \dots, 192, 212, 1, 214, \dots, 255$

Algoritma : Pseudo-random Generation Algorithm (PRGA)

Array s : [115, 213, 71, 191, 55, 124, 21, 77, 8, ..., 19, 20, 6, 22, 23
 53, 54, 4, 56, 56, ..., 69, 70, 2, 72, 73, 74, 75, 76, 77, 78, ..., 113
 114, 0, 116, 117, ..., 172, 173, 5, 175, 176, ..., 189, 190, 3, 192
 193, 1, ..., 211, 212, 1, 214, 215, ..., 250, 251, 252, 253, 254, 255]

Plaintext = "2090"

• Liturasi pertama - $idx = 0$

$$i = 0$$

$$j = 0$$

$$\begin{aligned} \Rightarrow i &= (i+1) \% 256 \\ &= (0+1) \% 256 \\ &= 1 \% 256 \end{aligned}$$

$$= 1$$

$$\begin{aligned} \Rightarrow j &= (j + s[i]) \% 256 \\ &= (0 + s[1]) \% 256 \\ &= (0 + 213) \% 256 \\ &= 213 \end{aligned}$$

Swap (s[i], s[j])

Swap (s[i], s[213])

Array s = [115, 1, 71, 191, 55, 174, 21, 77, 8, ..., 19, 20, 6, 22, 23, ...,
 53, 54, 4, 56, 57, ..., 69, 70, 2, 72, 73, 74, 75, 76, 77, 78, ...,
 113, 114, 0, 116, 117, ..., 172, 173, 5, 175, 176, 189, 190, 3, 192,
 193, ..., 212, 213, 214, ..., 250, 251, 253, 254, 255]

$$\begin{aligned} \Rightarrow f &= (s[i] + s[j]) \% 256 \\ &= (s[1] + s[213]) \% 256 \\ &= (1 + 213) \% 256 \\ &= 214 \end{aligned}$$

$$\begin{aligned} \Rightarrow &= s[f] \\ &= s[214] = 214 \Rightarrow \text{binner } 214 = 11010110 \end{aligned}$$

$$\begin{aligned} \Rightarrow &= 4 \oplus p[idx] \\ &= 4 \oplus p[0] \\ &= 4 \oplus "2" \Rightarrow \text{binner "2"} = 110010 \\ &= \begin{array}{r} 11010110 \\ 00110010 \\ \hline 11100100 \end{array} \end{aligned}$$

c = "4" didekripsi/can menjadi 223

Iterasi kedua $\rightarrow idx = 1$

$$i = 1$$

$$j = 213$$

$$\Rightarrow i = (1 + 1) / 256$$

$$= (1 + 1) / 256$$

$$= 2$$

$$\Rightarrow j = (j + s[i] * 256)$$

$$= (213 + s[2] * 256)$$

$$= (213 + 71) / 256$$

$$= 284 / 256$$

$$= 28$$

$$\text{Swap} = (s[i], s[j])$$

$$\text{Swap} = s[2], s[28]$$

Array $s = [115, 1, 20, 191, 55, 174, 21, 77, 8, \dots, 19, 20, 6, 22, 23, \dots, 26, 27, 71, 29, 30, \dots, 53, 54, 4, 56, 57, \dots, 69, 70, 71, 72, 73, 74, 75, 76, 77, 78, \dots, 113, 114, 0, 116, 117, \dots, 172, 173, 5, 175, 176, \dots, 189, 190, 3, 192, 193, \dots, 212, 213, 214, 215, \dots, 250, 251, 252, 253, 254, 255]$

$$\Rightarrow t = (s[i] + s[j] * 256)$$

$$= (s[2] + s[28] * 256)$$

$$= (20 + 71) * 256$$

$$= 94 * 256$$

$$= 94$$

$$\Rightarrow u = s[t]$$

$$= s[94]$$

$$= 99 \Rightarrow \text{biner } 99 = 1100011$$

$$\Rightarrow c = u \oplus p[idx]$$

$$= u \oplus p[1]$$

$$= u \oplus "0" = \text{biner "0"} = 110000$$

$$= 1100011$$

$$= 110000$$

$$= 110011$$

$$c = "S" \text{ desimal } = 83$$

$$\Rightarrow t = (s[i] + s[j] * 256)$$

$$= (s[3] + s[219] * 256)$$

$$= (219 + 191) * 256$$

$$= 410 * 256$$

$$= 104$$

$$\Rightarrow u = s[t]$$

$$= s[104]$$

$$= 154 \text{ biner } 154 =$$

$$\Rightarrow c = u \oplus p[idx]$$

$$= u \oplus p[2]$$

$$= u \oplus "3" = \text{biner "3"} = 110011$$

$$= 10011010$$

$$= 110011$$

$$= 10101001$$

$$c = "C" \text{ desimal } 169$$

Iterasi ketiga $= idx = 2$

$$i = 2, j = 28$$

$$i = (1 + 1) * 256$$

$$= (2 + 1) / 256$$

$$= 3$$

$$\text{Swap} (s[i], s[j])$$

$$\text{Swap} (s[3], s[219])$$

Array $s = [115, 1, 20, 214, 55, 174, 21, 77, 8, \dots, 19, 20, 6, 22, 23, 26, 27, 71, 29, 30, 53, 54, 4, 56, 57, \dots, 69, 70, 71, 72, 73, 74, 75, 76, 77, 78, 79, 113, 114, 0, 116, 117, \dots, 172, 173, 5, 175, 176, \dots, 189, 190, 3, 192, 193, \dots, 212, 213, 214, 215, 216, 217, 218, 191, 220, \dots, 253, 254, 255]$

• Herasi keempat $\Rightarrow \text{idx} = 3$

$$i = 3, j = 219$$

$$\begin{aligned}\Rightarrow i &= (i + 1) \% 256 \\ &= (3 + 1) \% 256 \\ &= 4\end{aligned}$$

$$\begin{aligned}\Rightarrow j &= (j + s[i]) \% 256 \\ &= (219 + s[4]) \% 256 \\ &= (219 + 55) \% 256 \\ &= 274 \% 256 \\ &= 18\end{aligned}$$

Swap (s[i], s[j])

Swap (s[4], s[18])

Array s = [115, 1, 28, 219, 18, 174, 21, 77, 8, ..., 16,
17, 55, 114, 20, 6, 22, 23, 24, 25, 26, 27,
71, 29, 30, ..., 53, 54, 4, 56, 57, 69, 70, 71,
72, 73, 74, 75, 76, 77, 78, 79, ..., 113, 114, 0,
116, 117, ..., 172, 173, 5, 175, 176, ..., 189, 190,
3, 191, 193, ..., 212, 213, 214, 215, 216, 217, 218,
191, 220, ..., 253, 254, 255]

$$\begin{aligned}\Rightarrow t &= (s[i] + s[j]) \% 256 \\ &= (s[4] + s[18]) \% 256 \\ &= 18 + 55 \% 256 \\ &= 73\end{aligned}$$

$$\begin{aligned}\Rightarrow u &= s[4] \\ &= s[73] \\ &= 73 \Rightarrow \text{Biner } 73 = 1001001\end{aligned}$$

$$\begin{aligned}\Rightarrow c &= 4 \oplus p[\text{idx}] \\ &= 4 \oplus p[3] \\ &= 4 \oplus "0" \Rightarrow \text{biner "0"} = 110000 \\ &= 1001001 \\ &\quad \underline{110000} \quad 0 \\ &\quad 1111001 \\ c &= "4" \text{ desimal} = 121\end{aligned}$$