### E303: Communication Systems

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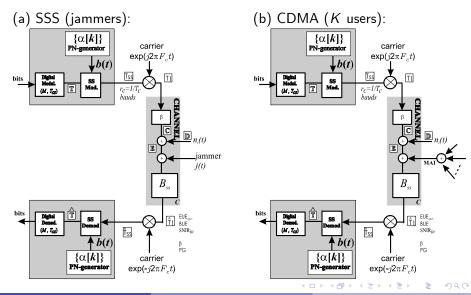
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An Overview of Fundamentals: PN-codes/signals & Spread Spectrum (Part B)

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### Introduction



# Modelling of PN-signals in SSS

- Consider  $-\begin{cases} \{\alpha[n]\} = \text{a sequence of } \pm 1/s \\ c(t) = \text{an energy signal of duration } T_c, \text{ e.g. } c(t) = \text{rect } \left\{\frac{t}{T_c}\right\} \end{cases}$ We have seen that the PN signal b(t) can be modelled as follows:
  - DS-SSS (Examples: DS-BPSK, DS-QPSK):

$$b(t) = \sum_{n} \alpha[n].c(t - nT_c)$$
 (1)

FH-SSS (Examples: FH-FSK)

$$b(t) = \sum_{n} \exp \{ j(2\pi k[n]F_1t + \phi[n]) \} .c(t - nT_c)$$
 (2)

where  $\{k[n]\}$  is a sequence of integers such that

$$\{\alpha[n]\} \mapsto \{\mathsf{k}[n]\}\tag{3}$$

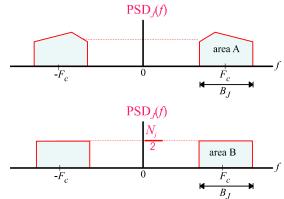
and with  $\phi[n]=\mathsf{random}\colon \mathsf{pdf}_{\phi[n]}=rac{1}{2\pi}\mathsf{rect}\Big\{rac{\phi}{2\pi}\Big\}$ 

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# Equivalent Energy Utilisation Efficiency (EUE)

#### Remember:

- ★ Jamming source, or, simply Jammer = intentional interference
- ★ Interfering source = unintentional interference



- \* With area-B = area-A we can find  $N_j$
- \*  $P_j = 2 \times \underline{\text{area} \mathbf{A}} = 2 \times \underline{\text{area} \mathbf{B}} = N_j B_j \Rightarrow N_j = \frac{P_j}{B_i}$

if

$$B_J = qB_{ss}; \ 0 < q \le 1 \tag{4}$$

then

$$EUE_J = \frac{E_b}{N_J} = \frac{P_s.B_J}{P_J.r_b} = \frac{P_s.q.B_{ss}}{P_J.B} = PG \times SJR_{in} \times q \quad (5)$$

$$EUE_{equ} = \frac{E_b}{N_0 + N_J} \tag{6}$$

$$= \underbrace{\mathsf{PG} \times \mathsf{SJR}_{in} \times q}_{\mathsf{EUE}_j} \times \left(\frac{\mathsf{N}_0}{\mathsf{N}_j} + 1\right)^{-1} \tag{7}$$

where

$$SJR_{in} \triangleq \frac{P_s}{P_I} \tag{8}$$

### Comments

- EUE<sub>equ</sub> (or EUE<sub>J</sub>): very important since bit error probabilities are defined as function of EUE<sub>equ</sub> (or of EUE<sub>J</sub>)
- For a specified performance
  - $\left\{ \begin{array}{l} \text{the smaller the SIR}_{in} \Rightarrow \text{the better for the signal} \\ \text{the larger the SIR}_{in} \ \Rightarrow \text{the better for the jammer} \end{array} \right.$
- Jammer limits the performance of the communication system i.e. effects of channel noise can be ignored i.e. Jammer Power $\gg P_n \Rightarrow \mathsf{EUE}_{equ} = \frac{E_b}{N_0 + N_J} \simeq \frac{E_b}{N_J} = \mathsf{EUE}_J$



N.B. exception to this:

- i) non-uniform fading channels
- ii) multiple access channels
- ∃ ∞ number of possible jamming waveforms
- There is no single jamming waveform that is worst for all SSSs
- There is no single SSS that is best against all jamming waveforms.

### Classification of Jammers in SSS

#### BROADBAND NOISE JAMMER



Spreads Gaussian Noise of total power P, evently over the total frequency range of the spread spectrum bandwidth  $B_{aa}$ 

$$EUE_J = \frac{E_b}{N_j}$$

The only knowldge available to (and exploited by) the jammer is the bandwidth  $B_{ij}$  of the SSS

 $D_{\cdot\cdot}$ : same as that with additive white Gaussian noise of PSD\_(f)=N/2

The performance with this type of jammer is known as

BASELINE PERFORMANCE

# PARTIAL NOISE



Spreads noise of total power P, evently over some frequency range  $B_i$  with  $\hat{B}_{i,j} \ni \hat{B}_i$ 

$$o = \frac{B_j}{B_{ss}} \le 1$$

#### CWJAMMER



$$j(t) = \sqrt{2P_j} \cos(2\pi F_1 t + \vartheta)$$

#### MULTITONE JAMMER



$$j(t) = \sum_{\mathbf{I}} \sqrt{\frac{2P_j}{N_j}} \cos(2\pi F_l t + \vartheta_l)$$

#### PULSE JAMMER

The jammer transmits with power

$$P_{j, peak} = \frac{P_j}{\rho}$$

for a fraction ρ of the time and nothing for the ramaining 1-ρ of the time The jammer first estimates some parameters of the SSSignal and

#### REPEAT-BACK **JAMMER**

frequ.following iammer

then transmits jamming signals which use this information.

Effective against FH-SSS with slow hop-rate enough for the jammer to respond within the hop duration.

Can be neutralized by increasing hop-rate.

#### ARBITRARY JAMMER POWER DISTRIBUTION

Spread Gaussian noise with arbitrary PSD(f) of total power P over the total frequency range (or some frequency range) of the SSSignal bandwidth  $B_{ss}$ 

# Direct Sequence SSS

#### Introductory Concepts & Mathematical Modelling

 If BPSK digital modulator is employed then the BPSK-signal can be modelled as:

**BPSK** 
$$s(t) = A_c \cdot \sin(2\pi F_c t + m(t) \cdot \frac{\pi}{2})$$
 (9)

where the data waveform can be modelled as follows:

$$m(t) \equiv \sum_{n} a[n] \cdot \underbrace{c_1(t - n \cdot T_{cs})}_{rect\left\{\frac{t - n \cdot T_{cs}}{T_{cs}}\right\}}; \quad nT_{cs} \leq t < (n+1) \cdot T_{cs}$$
 (10)

with  $\{a[n]\}=$  sequ. of independent data (message) bits  $(\pm 1$ 's)

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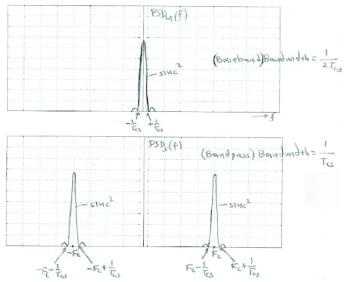
• Equation (9) can be rewritten as follows:

$$\mathsf{BPSK} \mid s(t) = A_c \cdot m(t) \cdot \cos(2\pi F_c t)$$

$$\therefore \mathsf{BPSK} \mathsf{ can be considered as } \left\{ \begin{array}{c} \mathsf{PM} \\ \mathsf{AM} \end{array} \right. \tag{11}$$

remember:

### • The PSD(f)'s of m(t) and s(t) are shown below





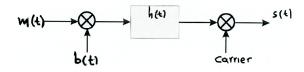
 If a DS/BPSK modulator is employed, then the SS-transmitted signal is

DS/BPSK: 
$$s(t) = A_c \sin \left( 2\pi F_c t + \underbrace{m(t)b(t)}^{\pm 1} \frac{\pi}{2} \right)$$
 (12)  
=  $A_c m(t)b(t) \cos(2\pi F_c t)$  (13)

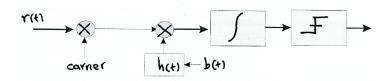
$$\text{where} \left\{ \begin{array}{l} m(t) \equiv \sum\limits_{n} \mathbf{a}[n] \cdot \underbrace{c_1(t-nT_{cs})}_{\uparrow}; \quad nT_{cs} \leq t < (n+1)T_{cs} \\ \text{rect}\{\frac{t-nT_{cs}}{T_{cs}}\} \\ b(t) = \sum\limits_{k} \alpha[k] \cdot \underbrace{c_2(t-kT_{cs})}_{\uparrow}; \quad kT_c \leq t < (k+1)\overset{\downarrow}{T_c} \\ \text{rect}\{\frac{t-kT_c}{T_c}\} \end{array} \right.$$

# BPSK DS/CDMA Transmitter & Receiver

• TX (see also Appendix-B):



• Rx (see also Appendix-B):



- The PN-sequence  $\{\alpha[I]\}$  (whose elements have values  $\pm 1$ )
- is M times faster than the data sequence  $\{a[n]\}$ .

• i.e.

$$T_{cs} = M \cdot T_c \tag{15}$$

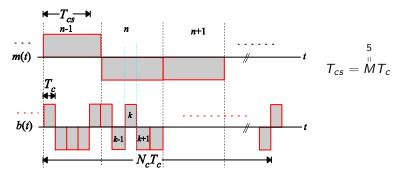
i.e.

$$PN-signal Bandwidth \geqslant \boxed{data-Bandwidth}$$
 (16)

Systems which have coincident data and SS code clocks

are often said to have a "data privacy feature"

such systems are easy to build and can be combined in single units



• Note:  $chip = T_c = smallest time increment$ 

• If the above "data privacy feature" is taken into account

then

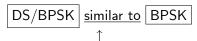
$$\begin{cases} m(t) \equiv \sum_{n} a[n].c_1(t - n \cdot T_{cs}) & nT_{cs} < t < (n+1) T_{cs} \\ b(t) = \sum_{k} \alpha[k].c_2(t - kT_c) & kT_c < t < (k+1) T_c \end{cases}$$
with  $T_{cs} = MT_c$ ;  $\left\lfloor \frac{k}{M} \right\rfloor = n$ ; (17)

where

$$n \cdot T_{cs} + k' \cdot T_c \le t < n \cdot T_{cs} + (k'+1) \cdot T_c$$
  
 $\forall k' = 0, 1, \dots, M-1 = k \mod M$ 

#### Conclusion

•



except that the apparent data rate is M times faster



signal spectrum is M times wider

Therefore

$$\mathsf{PG} = \frac{B_{\mathsf{ss}}}{B} = M \tag{18}$$

- Note:
  - message cannot be recovered without knowledge of PN-sequence i.e. PRIVACY
  - 2 typical:
    - **★** PN-chip-rate → several **M bits/sec**
    - ★ data rate → few bits/sec



# PSD of DS/BPSK/SS Transmitted Signal

Tx signal

$$\begin{array}{rcl} s(t) & = & m(t) \cdot b(t) \cdot A_c \cdot \cos(2\pi F_c t) \\ & & \Downarrow \\ \operatorname{PSD}_s(f) & = & \operatorname{PSD}_m(f) * \operatorname{PSD}_b(f) * \operatorname{PSD}_{A_c \cos(2\pi F_c t)}(f) \\ & & \Downarrow \\ \operatorname{PSD}_s(f) & = & \operatorname{PSD}_m(f) * \operatorname{PSD}_b(f) * \operatorname{PSD}_{A_c \cos(2\pi F_c t)}(f) \end{array}$$

remember

$$\mathsf{PSD}_{A_c \cos(2\pi F_c t)}(f) = \frac{A_c^2}{4} \cdot \left(\delta \left(f - F_c\right) + \delta \left(f + F_c\right)\right)$$

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• 
$$m(t) \equiv \sum_{n} a[n].c(t - n \cdot T_{cs}) \Rightarrow (\text{see "line-codes"})$$

$$\boxed{ \mathsf{PSD}_m(f) = \frac{|\mathsf{FT}(c(t))|^2}{T_{cs}} \cdot \left[ R_{\mathsf{aa}}[0] + \sum_{\substack{k = -\infty \\ k \neq 0}}^{+\infty} R_{\mathsf{aa}}[k] \exp(-j2\pi f k T_{cs}) \right] }$$
 (19)

• Note that if statistical indep. we have:

$$R_{aa}[k] = \begin{cases} \mathcal{E}\left\{a_n^2\right\} \text{ for } k = 0\\ \mathcal{E}\left\{a_n\right\} \cdot \mathcal{E}\left\{a_{n+k}\right\} \text{ for } k \neq 0 \end{cases}$$

i.e. 
$$R_{\rm aa}[k]=\left\{ egin{array}{l} \mu_{\rm a}^2+\sigma_{\rm a}^2 \ {\rm for} \ k=0 \ {\rm where} \ \mu_{\rm a}={\rm mean \ and} \ \sigma_{\rm a}={\rm std} \ \mu_{\rm a}^2 \ {\rm for} \ k\neq 0 \end{array} 
ight.$$

then

$$\boxed{ \mathsf{PSD}_m(f) = \underbrace{\sigma_{\mathsf{a}}^2 \frac{\big|\mathsf{FT}(c(t))\big|^2}{T_{cs}}}_{\mathsf{Continuous Spectrum}} + \underbrace{\frac{\mu_{\mathsf{a}}^2}{T_{cs}^2} \cdot \mathsf{comb}_{\frac{1}{T_{cs}}} \big(\big|\mathsf{FT}(c(t))\big|^2\big)}_{\mathsf{Discrete Spectrum}} }$$

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# PSD(f) of a Random Pulse Signal

• For a random pulse signal  $m(t) \equiv \sum\limits_{n} a[n].c(t-n\cdot T_{cs})$ 

(i.e. a sequence of pulses where there is an invariant average time of separation  $T_{cs}$  between pulses) with all pulses of the same form but with

- random amplitudes a[n] with mean= $\mu_a = \mathcal{E}\{a[n]\}$  and std= $\sigma_a^2 = \mathcal{E}\{(a[n] \mu_a)^2\}$
- statistical independent random time of occurrence ,

Then:

$$\mathsf{PSD}_m(f) = \underbrace{\sigma_a^2 \frac{\left|\mathsf{FT}(c(t))\right|^2}{T_{cs}}}_{\mathsf{Continuous Spectrum}} + \underbrace{\frac{\mu_a^2}{T_{cs}^2} \cdot \mathsf{comb}_{\frac{1}{T_{cs}}} (\left|\mathsf{FT}(c(t))\right|^2)}_{\mathsf{Discrete Spectrum}}$$

Note that if  $\mu_a=0$  (this is a very common case ) then

$$| \mathsf{PSD}_m(f) = \mathcal{E}\{\mathsf{a}[n]^2\} \frac{|\mathsf{FT}(\mathsf{single pulse})|^2}{T_{cs}}$$

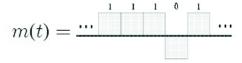
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# Example: PSD of a random BINARY signal m(t)

 Consider a random binary sequence of 0's and 1's. This binary sequence is transmitted as random signal m(t) with 1's and 0's being sent using the pulses shown below.



For instance a random binary sequence/waveform could be



• If 1's and 0's are statistically independent with Pr(0) = Pr(1) = 0.5, the PSD of the transmitted signal can be estimated as follows:

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$$PSD_{m}(f) = \mathcal{E}\left\{a[n]^{2}\right\} \cdot \frac{\left\|FT\left(\underbrace{T_{cs}}\right)^{1}\right\|^{2}}{T_{cs}}$$

$$= \mathcal{E}\left\{a[n]^{2}\right\} \cdot \frac{T_{cs}^{2} \cdot \operatorname{sinc}^{2}(fT_{cs})}{T_{cs}}$$

$$= \mathcal{E}\left\{a[n]^{2}\right\} \cdot \frac{T_{cs} \cdot \operatorname{sinc}^{2}(fT_{cs})}{T_{cs}}$$

$$= \mathcal{E}\left\{a[n]^{2}\right\} \cdot \frac{T_{cs} \cdot \operatorname{sinc}^{2}(fT_{cs})}{T_{cs}}$$

$$= \mathcal{A}^{2}\left\{a[n]^{2}\right\} \cdot \frac{T_{cs} \cdot \operatorname{sinc}^{2}(fT_{cs})}{T_{cs}}$$

$$= \mathcal{A}^{2}\left\{T_{cs}\operatorname{sinc}^{2}\left\{fT_{cs}\right\}\right\}$$

# PSD(f) of a PN-Signal b(t) in DS-SSSs

• Autocorrelation function:  $R_{bb}(\tau)$ 

$$R_{bb}(\tau) = \frac{N_c + 1}{N_c} \operatorname{rep}_{N_c T_c} \left\{ \Lambda \left( \frac{\tau}{T_c} \right) \right\} - \frac{1}{N_c}$$

$$N_c^{1} \qquad N_c T_c \qquad (20)$$

• Using the FT tables the  $PSD_b(f) = FT\{R_{bb}(\tau)\}$  of the signal b(t) is:

$$PSD_{b}(f) = \frac{N_{c}+1}{N_{c}^{2}} comb_{\frac{1}{N_{c}T_{c}}} \left\{ sinc^{2} \left\{ f \cdot T_{c} \right\} \right\} - \frac{1}{N_{c}} \delta(f)$$

$$\frac{sinc^{2}}{\sqrt{N_{c}T_{c}}} \left\{ \frac{A_{c}^{2}}{\sqrt{N_{c}T_{c}}} \right\}$$

$$\frac{A_{c}^{2}}{\sqrt{N_{c}T_{c}}} \left\{ \frac{A_{c}^{2}}{\sqrt{N_{c}T_{c}}} \right\}$$

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# PSD(f) of DS/BPSK Spread Spectrum Tx Signal s(t)

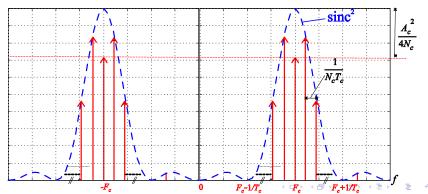
$$s(t) = m(t) \cdot b(t) \cdot A_c \cdot \cos(2\pi F_c t)$$

$$\Rightarrow \mathsf{PSD}_s(f) = \mathsf{PSD}_m(f) * \underbrace{\mathsf{PSD}_b(f) * \frac{A_c^2}{4} \left(\delta(f - F_c) + \delta(f + F_c)\right)}_{\text{term 1}}$$

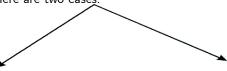
• **Ignore** (for the time being) the effects of m(t)

$$\begin{split} \mathsf{PSD}_{\mathsf{term1}}(f) &= \frac{A_c^2}{4} \cdot \mathsf{PSD}_b(f) * \left( \delta(f - F_c) + \delta(f + F_c) \right) \\ &= \frac{A_c^2}{4} \cdot \left\{ \frac{N_c + 1}{N_c^2} \cdot \mathsf{comb}_{\frac{1}{N_c T_c}} \left\{ \mathsf{sinc}^2(f \cdot T_c) \right\} - \frac{1}{N_c} \cdot \delta(f) \right\} \\ &* \left( \delta(f - F_c) + \delta(f + F_c) \right) \\ &= \frac{A_c^2}{4} \cdot \frac{N_c + 1}{N_c^2} \cdot \left( \mathsf{comb}_{\frac{1}{N_c T_c}} \left\{ \mathsf{sinc}^2 \left\{ (f - F_c) T_c \right\} \right\} \right) \\ &+ \mathsf{comb}_{\frac{1}{N_c T_c}} \left\{ \mathsf{sinc}^2 \left\{ (f + F_c) T_c \right\} \right\} \right) \\ &- \frac{A_c^2}{4} \frac{1}{N_c} \left( \delta(f - F_c) + \delta(f + F_c) \right) \end{split}$$

$$PSD_{term1}(f) = \frac{A_c^2}{4} \cdot \frac{N_c + 1}{N_c^2} \cdot \left( comb_{\frac{1}{N_c T_c}} \left\{ sinc^2 \left\{ (f - F_c) T_c \right\} \right\} + comb_{\frac{1}{N_c T_c}} \left\{ sinc^2 \left\{ (f + F_c) T_c \right\} \right\} \right) - \frac{A_c^2}{4} \frac{1}{N_c} \left( \delta(f - F_c) + \delta(f + F_c) \right)$$
(22)



- If m(t) is used, then each discrete frequency in Equation (22) becomes a sinc<sup>2</sup> function.
- ► There are two cases:



#### • CASE-1:

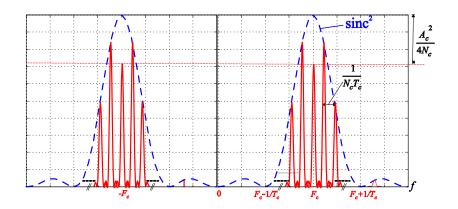
$$\frac{1}{T_{cs}} < \frac{1}{2N_c T_c}$$

#### • CASE-2:

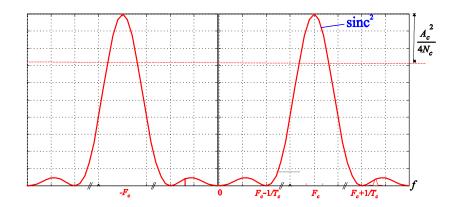
$$\frac{1}{T_{cs}} \geq \frac{1}{N_c T_c}$$

the peaks will merge into a continuous smooth spectrum

#### • CASE-1:



#### • CASE-2:



#### N.B.:

► If

$$b(t) = random$$

then

$$s(t) = A_c m(t) b(t) \cos(2\pi F_c t)$$

▶ If the effects of m(t) are ignored then

$$PSD_{s}(f) = A_{c}^{2} T_{c}^{2} sinc \left\{ fT_{c} \right\} * \frac{1}{4} \left( \delta(f - F_{c}) + \delta(f + F_{c}) \right)$$

$$\downarrow \qquad \qquad \downarrow$$

$$PSD_{s}(f) = \frac{T_{c}}{4} A_{c}^{2} \left( sinc^{2} \left\{ (f - F_{c}) T_{c} \right\} + sinc^{2} \left\{ (f + F_{c}) T_{c} \right\} \right)$$

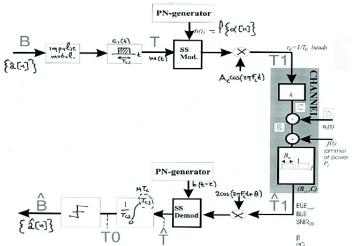
i.e. PSD is similar to 'CASE-2' above



### DS-BPSK Spread Spectrum:

#### Output SNIR

 Consider the block diagram of a SS-Communication System which employs a BPSK digital modulator:



• N.B.: 
$$B_{ss} = \frac{1}{T_c}$$
;  $B = \frac{1}{T_{cs}}$ ;  $PG = \frac{B_{ss}}{B} = \frac{T_{cs}}{T_c}$ 

• Then, at point | T1 |, we have

$$s(t) = A_c \cdot m(t) \cdot b(t) \cdot \cos(2\pi F_c t) \tag{23}$$

and at point  $\left|\frac{\wedge}{\mathsf{T}1}\right|$ :

$$s(t) + n(t) + j(t)$$
 (for  $k = 1$ ) (24)

 At the input of the receiver the Signal-to-Noise-plus-Interference Ratio (SNIR<sub>in</sub>) is:

$$SNIR_{in} = \frac{EUE}{PG \cdot (1 + JNR_{in})} = \frac{EUE_{equ}}{PG}$$
 (25)

• at point  $\stackrel{\wedge}{\mathsf{T}}$ :

$$(s(t) + n(t) + j(t)) \cdot b(t - \tau) \cdot 2\cos(2\pi F_c t + \theta)$$
 (26)

• at point T0:

$$P_{\text{unwanted}} = P_{n_{\text{out}}} + P_{\text{code-noise}} + P_{j_{\text{out}}}$$
 (27)

• if  $\tau = 0$  and  $\theta = 0$  (i.e. the system is synchronized) then:

$$P_{\mathsf{code} ext{-noise}} = 0$$

and

$$SNIR_{out-max} = 2EUE_{equ}$$
 (28)

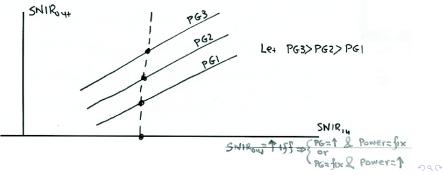


However,

$$SNIR_{in} = \frac{EUE}{PG \cdot (1 + JNR_{in})} = \frac{EUE_{equ}}{PG}$$
 (29)

Therefore,

$$SNIR_{out-max} = 2EUE_{equ} = 2.PG.SNIR_{in}$$
 (30)



# Bit Error Probability with Jamming

#### A. CONSTANT POWER BROADBAND JAMMER:

• From the "Detection Theory" topic we know that the bit-error-probability  $p_e$  for a Binary Phase-Shift Key (BPSK) communication system is given by:

$$p_{e} = T \left\{ \underbrace{\sqrt{2 \cdot EUE}}_{SNR_{out, matched filter}} \right\} \quad \text{where EUE} = \frac{E_{b}}{N_{0}}$$
 (31)

• Consider a DS/BPSK SSS which operates in the presence of a constant amplitude broadband jammer with double sided power spectral density

$$\mathsf{PSD}_j(f) = \frac{N_j}{2} \tag{32}$$
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Then.

$$p_{\mathsf{e}} = T \left\{ \sqrt{2 \cdot \mathsf{EUE}_{\mathsf{equ}}} 
ight\}$$

where 
$$\mathsf{EUE}_{\mathsf{equ}} = \frac{E_b}{N_0 + N_j}$$
 with  $N_j = \frac{P_J}{B_{\mathsf{ss}}}$ 

• If we make the assumption that  $N_i \gg N_0$  then

$$p_{e} = T\{\sqrt{2 \cdot EUE_{J}}\}$$
 (33)

where 
$$EUE_J = \frac{E_b}{N_j}$$

 This is known as the BASELINE PERFORMANCE of a DS/BPSK SSS

#### **B. PULSE JAMMER:**

- Consider a DS/BPSK SSS which operates in the presence of a jammer which transmits "broadband noise" with large power but only a fraction of the time.
- The double-sided power spectral density of the jammer is given by:

$$PSD_j(f) = \frac{N_j}{2\rho} \tag{34}$$

- where
  - $ho \equiv$  the fraction of time the jammer is "on".
  - $ightharpoonup P_i$  = average jamming power
  - $ightharpoonup rac{P_j}{
    ho}=$  actual power during a jamming pulse duration

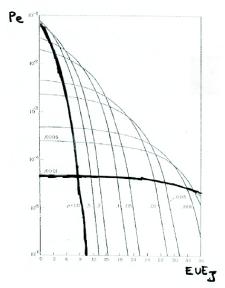
- Let the jammer pulse duration be greater than  $T_{cs}$  (data bit time).
  - Then  $\begin{cases} & \text{Pr(jammer} = \text{"on"}) = \rho \\ & \text{Pr(jammer} = \text{"off"}) = 1 \rho \end{cases}$  and the bit-error-probability is given by:

$$p_{e} = \underbrace{(1 - \rho)T\left\{\sqrt{2\frac{E_{b}}{N_{0}}}\right\}}_{\simeq 0 \text{ (very small)}} + \rho T\left\{\sqrt{2\frac{E_{b}}{N_{0} + \frac{N_{j}}{\rho}}}\right\}$$
(35)

which can be simplified to

$$p_e = \rho \cdot T \left\{ \sqrt{2\rho \cdot EUE_J} \right\} \quad \text{where } EUE_J = \frac{E_b}{N_i}$$
 (36)

 $\bullet$  By plotting the above equation for different values of  $\rho$  we get:



#### Note that

- lacktriangle the value of which maximizes  $p_e$  decreases with increasing values of  ${\sf EUE}_j$
- there is a value of  $\rho$  which maximizes the probability of error  $p_e$ . This value can be found by differentiating Equation-36 with respect to  $\rho$ . That is

$$\frac{dp_e}{d\rho} = 0 \Rightarrow \rho^* = \begin{cases} \frac{0.709}{\text{EUE}_j} & \text{if EUE}_j > 0.709\\ 1 & \text{if EUE}_j \le 0.709 \end{cases}$$
 (37)

Therefore

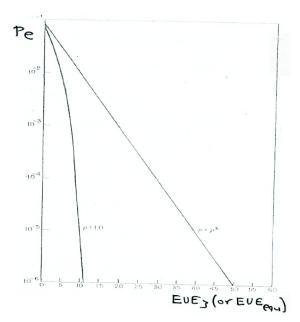
$$p_{e_{\text{max}}} = \max_{\rho} \left\{ \rho \cdot T \left\{ \sqrt{2\rho \text{EUE}_{j}} \right\} \right\}$$

$$p_{e_{\text{max}}} = \rho^{*} \cdot T \left\{ \sqrt{2\rho^{*} \text{EUE}_{j}} \right\}$$
(38)

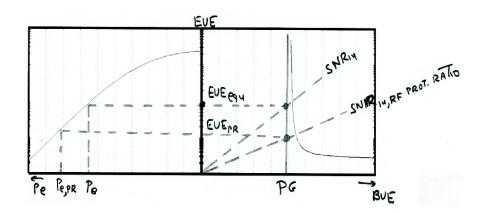
$$\implies p_{\mathsf{e}_{\mathsf{max}}} = \begin{cases} \frac{0.083}{\mathsf{EUE}_j} & \text{if } \mathsf{EUE}_j > 0.709 \\ \mathsf{T} \left\{ \sqrt{2\mathsf{EUE}_j} \right\} & \text{if } \mathsf{EUE}_j \le 0.709 \end{cases} \tag{39}$$

#### N.B.:

- when jammer pulse length is shorter than a data bit time  $T_{cs}$  then the above expression is not valid.
- ▶ However, Equation-39 represents an UPPER BOUND on the bit-error-probability  $p_e$ .
- The next graph illustrates
  - ▶ the bit-error-prob. plotted against the EUE; for a baseline jammer (i.e.  $\rho=1$ ) and
  - the worst case jammer (i.e.  $\rho = \rho^*$ ) for a DS/BPSK SSS.
- Note the huge difference between the two curves.



### DS-SSS on the (pe, EUE, BUE)-parameter plane



## Anti-Jam Margin

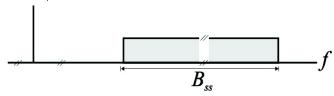
 An important parameter of SSS is the ANTIJAM MARGIN which is defined as follows:

$$\label{eq:db} \begin{split} \mathsf{dB}(\mathsf{AJM}) &\equiv \mathsf{dB}(\mathsf{EUE}_{\mathsf{equ}}) - \mathsf{dB}(\mathsf{EUE} \; \mathsf{which} \; \mathsf{corresponds} \; \mathsf{to} \; \mathsf{the} \; \rho_{e,PR}) \\ \Rightarrow & \boxed{\mathsf{dB}(\mathsf{AJM}) \equiv 10 * \mathsf{log}(\mathsf{EUE}_{\mathsf{equ}}) - 10 * \mathsf{log}(\mathsf{EUE}_{\rho_e,PR})} \end{split}$$

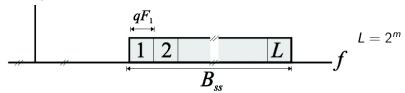
 AJM represents a safety margin against jammer (or against jammer plus noise).

## Frequency Hopping SSSs

 Consider that the following part of the spectrum has been allocated to a FH/SSS:



• Let us partition the above spectrum onto L different frequency slots of bandwidth  $F_1$  (or of bandwidth  $qF_1$  where q is a constant). Then  $B_{ss} = q \cdot F_1 \cdot L$ 



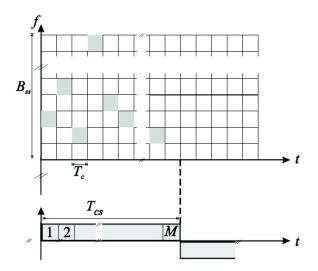
• Define the following symbols:

$$T_c = ext{hop duration (i.e. hop rate } r_{hop} = rac{1}{T_c})$$
 $T_{cs} = ext{message bit duration (i.e. bit rate } r_b = rac{1}{T_{cs}})$ 
 $M = ext{number of hops per message bit (i.e. } T_{cs} = M \cdot T_c)$ 

$$ullet$$
 FH/SSS: 
$$\left\{ egin{array}{ll} {\sf Fast hop} & 
ightarrow r_{hop} > r_b \ {\sf slow hop} & 
ightarrow r_{hop} < r_b \ {\sf balance hop} & 
ightarrow r_{hop} = r_b \end{array} 
ight.$$

• The frequency slot is constant in each time chip  $T_c$ , BUT changes from chip-to-chip.

This can be represented by the following diagram:



• In general, the No. of different frequency slots *L* over which the signal may hop, is a power of 2.

•  $F_1$  is, in general, equal to  $\frac{1}{T_c}$ ,

i.e.  $F_1 = \frac{1}{T_c}$  (but this is not a necessary requirement).

- several FH signals occupy a common RF channel
- FH model of b(t) (complex representation):

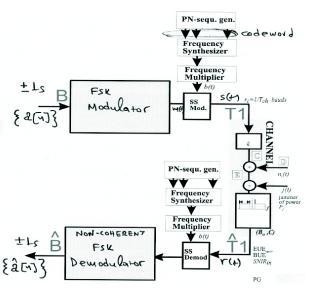
$$b(t) = \sum_{n} \exp \left\{ j(2\pi.k[n].F_1t + \phi_n) \right\} \cdot \text{rect} \left\{ \frac{t - nT_c}{T_c} \right\}$$

where

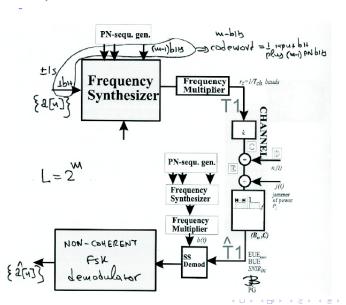
- $k[n] = \mathbf{f}_{\{PN-\text{sequ } \{\alpha[n]\}\}}$
- $\blacktriangleright$  k[n] is an integer that is formed by a codeword which is formed by one or more m-sequences



Transmitter-Receiver path:



## A Different Implementation:

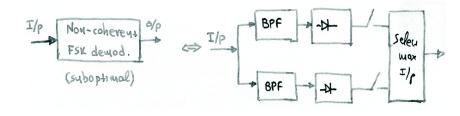


• L frequencies are produced by the digital Frequ. Synthesizer, separated by  $F_1$ 

$$B_s \simeq q \cdot F_1 \cdot L \tag{40}$$

- Note:
  - ▶ if reception= coherent:
    - \* more difficult to achieve
    - ★ places constraints on the transmitted signal and transmitted medium
  - ▶ if reception= non-coherent :
    - ★ PN-gen. can run at a considerably slower rate in this type of system than in a DS system

FH: non-coherent ⇒ poorer performance against thermal noise.



- Performance:
  - Coherent FSK (CFSK)

$$p_{e,CFSK} = T \left\{ \sqrt{\mathsf{EUE}} \right\}$$

Non-Coherent FSK (NFSK)

$$p_{e,NFSK} = \frac{1}{2} \exp\left(-\frac{\mathsf{EUE}}{2}\right)$$



very strong signals at receiver swapping out the effects of weaker signal 

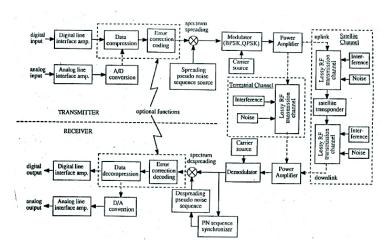
"near-far" problem

- A serious problem is the
  - ▶ DS: severe problem
  - ▶ FH: much more susceptible
- acquisition: much faster in FH than in DS
- $PG = \frac{B_{ss}}{B} = it$  is not very good criterion for FH

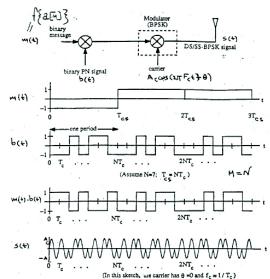
## **Appendices**

#### Appendix A. Block Diagram of a Typical SSS

(terrestrial & satellite comm. systems)



# Appendix B. BPSK/DS/SS Transmitter and Receiver BPSK/DS/SS Transmitter



# BPSK/DS/SS Receiver

