

**Discrete Mathematics Quiz 1**  
2008.04.08

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**Instructions.** This is a 2-hr close book quiz. Please manage your time well. No dictionary, calculator, PDA, or any other electronic device. Any dishonorable cheating behavior will give you a miserable future. (Totally 115 points)

1. [10%] Let  $d_{\min}$  denote the minimum degree of a vertex in an undirected graph  $G = (V, E)$ , where  $V$  and  $E$  denote the set of vertices and edges. Show that the graph contains a simple path containing at least  $d_{\min}$  edges. (hint: proof by contradiction, assume the longest simple path  $v_0 - v_1 - \dots - v_k$  contains  $k < d_{\min}$  edges, and then argue why it can NOT be the longest path.)

**Ans:** Suppose not, which means the longest simple path  $v_0 - v_1 - \dots - v_k$  contains  $k < d_{\min}$  edges. Take vertex  $v_k$  for example, there exist at least  $d_{\min}$  adjacent vertices to  $v_k$ , which means at least one vertex in  $N - \{v_0, v_1, \dots, v_{k-1}\}$  is adjacent to  $v_k$ , and thus it implies the simple path  $v_0 - v_1 - \dots - v_k$  can be at least one edge longer which contradicts the fact that it is already the longest.

2. [8%] In  $Z_7$ , which of the following congruence classes are equal?  
[2], [7], [10], [16], [39], [45], [-1], [-3], [-6], [-17], [-23]

**Ans:**

[2]=[16], [10]=[45], [39]=[-3]=[-17]

3. [62%] The following table lists a number of tasks that must be completed in order for a crew of workers to finish a project.

Task	A	B	C	D	E	F	G	H	I	J
Time (in days)	2	2	3	1	1	2	3	4	3	3

Suppose the project starts at Apr. 9. Task A must be carried out before any other tasks can start. Task B must precede tasks E and F, and both E and F must be completed before H can begin. Tasks C and D must precede task G, which in turn must precede I. Task J must be carried out last. It is assumed that there are enough workers to carry out any number of tasks simultaneously.

- (a) [5%] Draw a PERT diagram showing the relations of these tasks.
- (b) [5%] Draw the critical path.
- (c) [5%] We say a "conflict of tasks" occurs when two tasks violate their precedent relations. Give a suitable sequence for conducting all the tasks so that no conflict of tasks occurs in the entire process.

- (d) [5%] What is the fewest number of days needed to make this product?
- (e) [5%] What is the earliest date for this project to be done?
- (f) [5%] If we start to do task D in Apr.14, will it affect your answer in (e)? If yes, affect how much? If no, explain why not.
- (g) [5%] What is the latest date for task H to begin, in order not to affect your answer in (e)?
- (h) [9%] For any two tasks  $x$  and  $y$ , suppose we define  $x\mathbf{R}y$  to be true if either task  $x$  equals to task  $y$  or task  $x$  can not be started until task  $y$  is completed. Explain why  $\mathbf{R}$  is reflexive, antisymmetric, and transitive.
- (i) [6%] Following (h), what are  $BRJ$ ,  $IRC$ ,  $HRD$ ? (i.e. True or False?)
- (j) [4%] Following (h), let  $S$  be set of all tasks and  $\mathbf{R}$  be a relation on  $S$ . List all the minimal and maximal elements of  $S$  with respect to  $\mathbf{R}$ .
- (k) [4%] Following (h), what are the infimum and supremum of  $E$  and  $F$ ? what are the infimum and supremum of  $G$  and  $F$ ?
- (l) [4%] Following (h), is  $(S, \mathbf{R})$  a lattice? Explain your answer. (an answer without explanation gets at most 1 point)

**Ans:**

- (a) see Fig.1
- (b) A-C-G-I-J
- (c) ABCDEFGHIJ (there are many other permutations, as long as it follows topological ordering)
- (d) 14 days
- (e) Apr.22, since  $\text{Apr.9} + 14 - 1 = \text{Apr.22}$
- (f) Yes, the earliest date for G to start was supposed to be Apr.14, determined by C. If D starts at Apr. 14, the earliest date for G to start becomes Apr.15 which will make the tight schedule 1 day longer.
- (g) Apr.16. Since the earliest time for J to start is Apr.20 ( $\text{Apr.22} - 3 + 1$ ), which means the latest date for H to finish is Apr.19, and thus  $\text{Apr.19} - 4 + 1 = \text{Apr.16}$  is the latest date for H to start, without affecting the tightest schedule in (e).
- (h)  $\mathbf{R}$  is obviously reflexive since a task  $x$  equals to itself.  
 Suppose  $x \neq y$ . If  $x\mathbf{R}y$  is true, task  $x$  is not started until  $y$  is completed, and thus  $y$  must have been started before  $x$  is completed, which means  $y\mathbf{R}x$  is false. This makes the only chance to have both  $x\mathbf{R}y$  and  $y\mathbf{R}x$  be true is to let  $x = y$ , which is true since  $\mathbf{R}$  is reflexive. This shows  $\mathbf{R}$  is antisymmetric.
- Suppose  $x \neq y$ . If both  $x\mathbf{R}y$  and  $y\mathbf{R}z$  are true, which means  $y$  is completed before  $x$  is started, and  $z$  is completed before  $y$  is started, thus  $z$  is completed before  $x$  is started and  $x\mathbf{R}z$  becomes true. So  $\mathbf{R}$  is transitive.
- (i)  $BRJ = \text{False}$ ,  $IRC = \text{True}$ ,  $HRD = \text{False}$
- (j) minimal elements: J; maximal element: A
- (k) infimum of E,F: H, supremum of E,F: B; infimum of G,F: J, supremum of G,F: A
- (l) yes, since all pairs of tasks uniquely meet or join each other.

4. [16%] Let  $X$  denote the set of real numbers. Suppose we have a function  $f : X \rightarrow X$ ,  $f(x) = \frac{1}{|x|+1}$
- (a) [4%] Is  $f$  well-defined? why?
  - (b) [2%]  $\text{rng } f = ?$
  - (c) [6%] Is  $f$  injective? why? Is  $f$  surjective? why?
  - (d) [4%] Derive  $f^{-1}$ , if one exists; or explain why it does not exist.

**Ans:**

(a) yes. since for each  $x \in X$ ,  $f(x)$  is uniquely defined in  $X$  (i.e. for a given  $x \in X$ , there will be only one  $f(x)$ )

(b) Let  $y = \frac{1}{|x|+1}$ .

If  $x \geq 0$ ,  $y = \frac{1}{x+1}$ ,  $x = \frac{1-y}{y} = f^{-1}(y)$  with  $y \in X - \{0\}$

If  $x < 0$ ,  $y = \frac{1}{-x+1}$ ,  $x = \frac{y-1}{y} = f^{-1}(y)$  with  $y \in X - \{0\}$

It is easy to show  $y \in (0, 1]$ , since  $|x| \geq 0$ ,  $1 + |x| \geq 1$ , which means  $\frac{1}{|x|+1} \leq 1$ ; also,  $\lim_{x \rightarrow \infty} \frac{1}{|x|+1} = 0$ .

Thus  $\text{rng } f = (0, 1]$ .

(c)  $f$  is NOT injective since if  $x_1 = -x_2 \neq 0$ ,  $f(x_1) = f(x_2)$

$f$  is NOT surjective since  $\text{rng } f = (0, 1] \neq X = \mathbb{R}$ .

(d) by (c), we know  $f^{-1}$  does not exist, since it is not one-to-one correspondence.

5. [19%] If  $X$  has  $m$  elements and  $Y$  has  $n$  elements,
- (a) [2%] Is it possible that  $m > n$ , if there is a one-to-one function with domain  $X$  and codomain  $Y$ ? Explain your answer.
  - (b) [5%] How many one-to-one functions are there with domain  $X$  and codomain  $Y$ ?
  - (c) [2%] Is it possible that  $m > n$ , if there is an onto function with domain  $X$  and codomain  $Y$ ? Explain your answer.
  - (d) [10%] How many onto functions are there with domain  $X$  and codomain  $Y$ ?

**Ans:**

(a) No, since such a function will make some element in  $X$  not mappable to  $Y$ .

(b) When  $m \leq n$ , totally there are  $P_m^n = n \cdot (n-1) \cdot \dots \cdot (n-m+1)$  possible mapping relations (i.e. functions).

(c) Yes, as long as each element in  $X$  is mapped to an element in  $Y$ , and each element in  $X$  is mapped.

(d) It is not possible that  $m < n$ , if there is an onto function with domain  $X$  and codomain  $Y$ , since it will force some element in  $X$  maps to more than one element in  $Y$ , contradicts the definition of functions. Thus for an onto function with domain  $X$  and codomain  $Y$ ,  $m \geq n$ . In this case, the total possibilities can be calculated by the total number of functions -  $\sum_{i=1}^{n-1}$  (the total number of functions without mapping  $i$  elements in  $Y$ ) which equals to  $n^m - \sum_{i=1}^{n-1} (-1)^{i+1} C_i^n (n-i)^m$ .