# Algorytmy metaheurystyczne 2

Paweł Cegieła, Wojciech Sęk 5 maja 2022

## 1 Teoretyczna złożoność

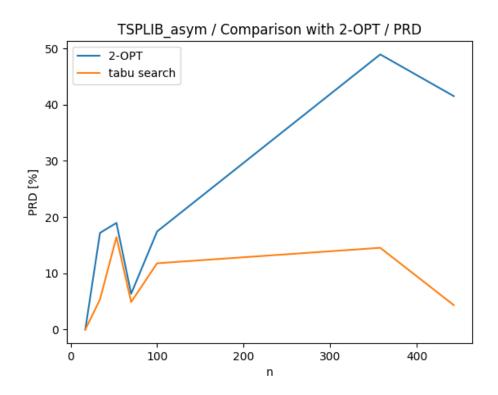
Warunkiem wyjścia w naszym algorytmie było przekroczenie 15n iteracji, gdzie n = |V| lub n ruchów bez zmiany na lepsze rozwiązanie. Rozważmy najgorszy możliwy przypadek, gdzie wykonujemy 15n iteracji. Niech k to długość listy tabu. Implementacja listy tabu za pomocą VecDeq pozwala na dostęp do i-tego elementu w czasie stałym, a usuwanie i dodawanie elementów w czasie liniowym.

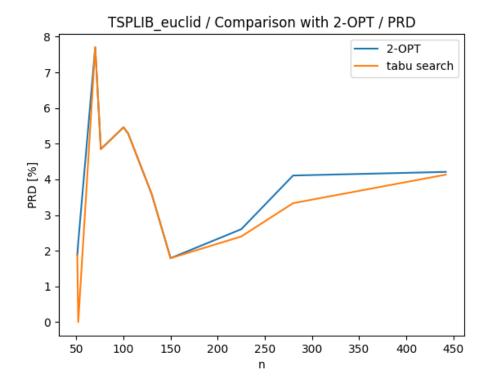
W każdym kroku algorytmu przeglądamy wszystkich  $\frac{n(n-1)}{2}$  sąsiadów danego rozwiązania i dla każdego sprawdzamy z O(k) czy jest na liście tabu. Sprawdzenie o ile sąsiad zmienia wartość permutacji jest stały (dla *invert* liczymy wcześniej z  $O(n^2)$  pomocnicze tablice. Niech O(l) to złożoność przybliżenia początkowego (dla 2-opt  $O(n^3)$ ). Wybieramy najlepszego z nich. Ostatecznie mamy (dla k stałego i 2-opta):

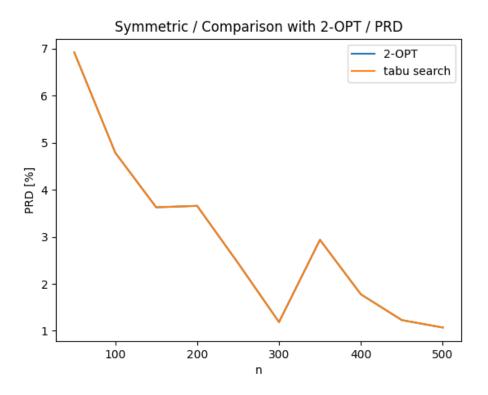
$$O\left(15n \cdot \frac{n(n-1)}{2} \cdot k + l\right) = O(n^3)$$

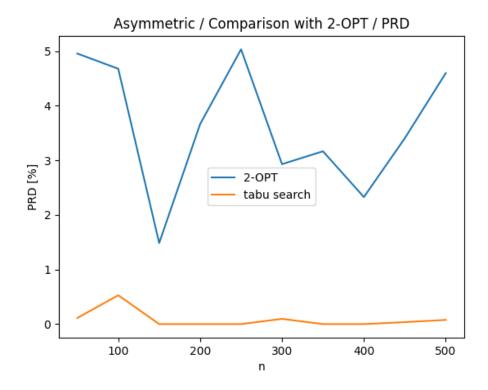
## 2 Porównanie Tabu Search z 2-OPT

#### 2.1 Dane z TSPLIB

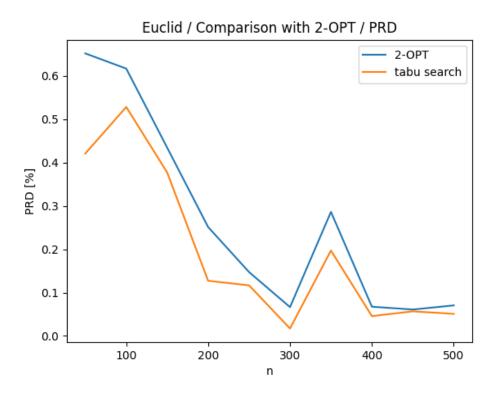








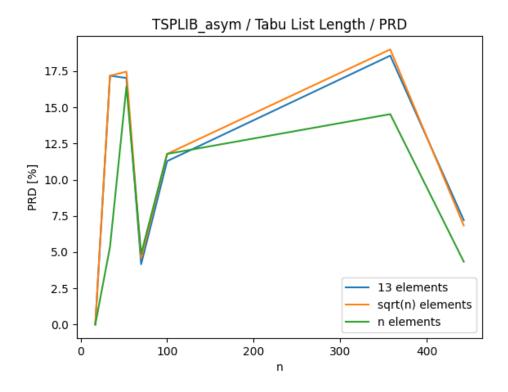
## 2.2.3 Grafy euklidesowe

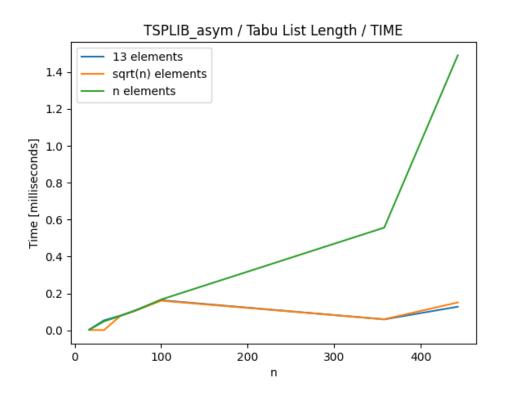


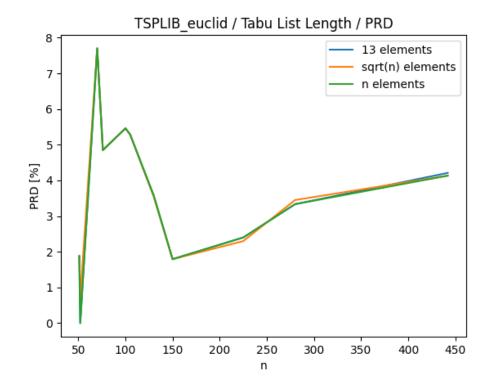
### 2.4 Tabele

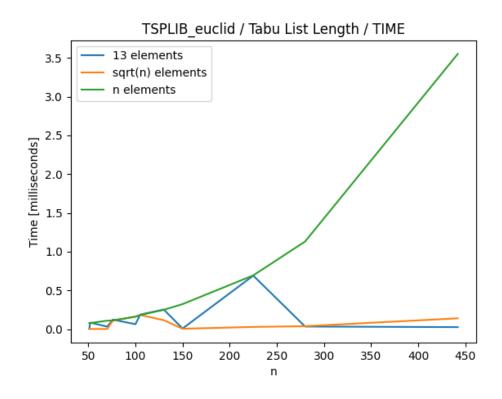
## 3 Porównanie różnych długości listy tabu

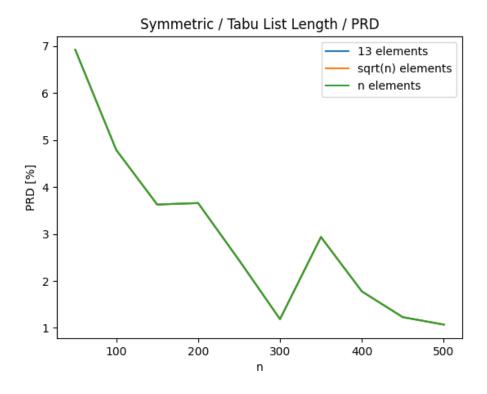
## 3.1 Dane z TSPLIB

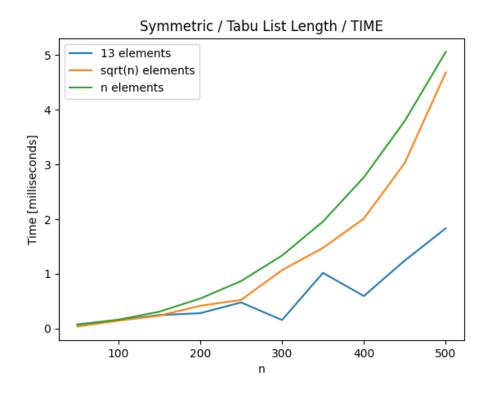


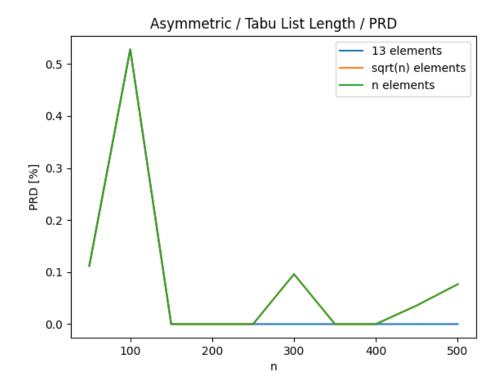


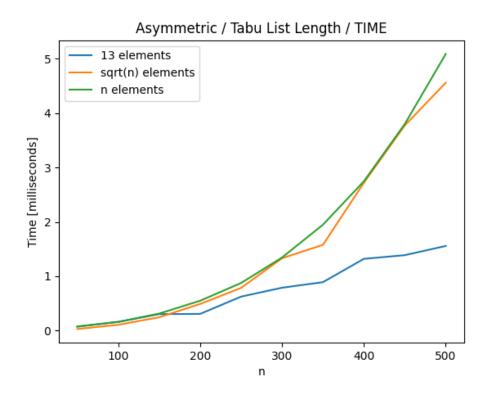


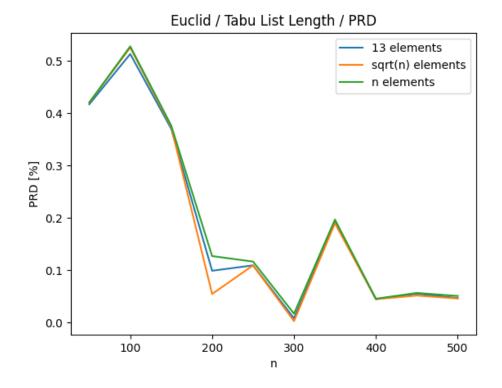


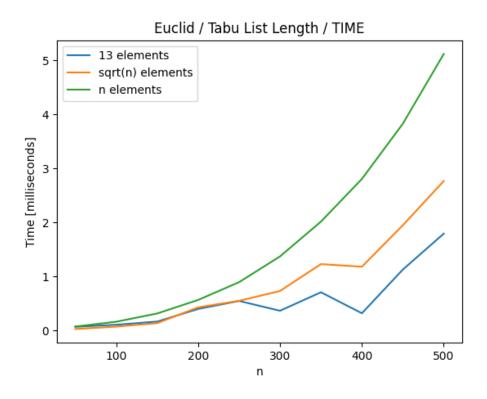








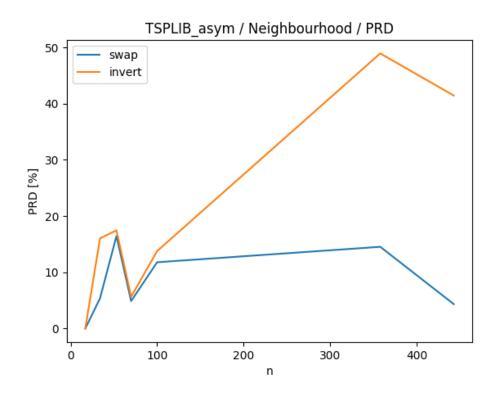


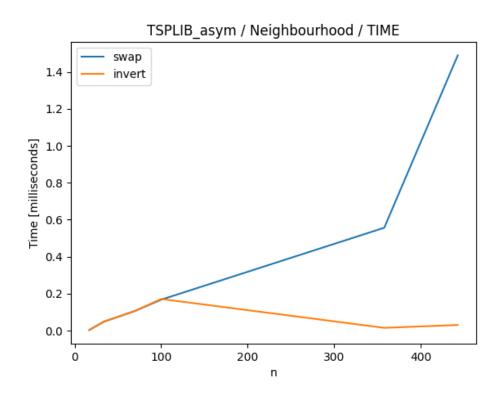


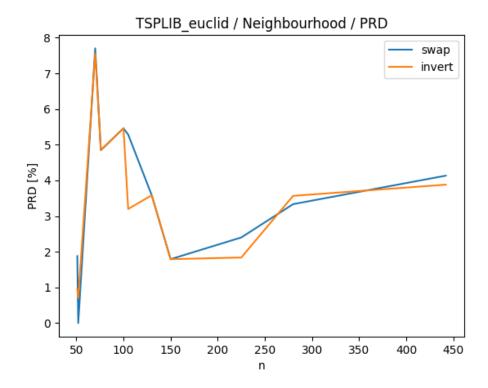
## 3.4 Tabele

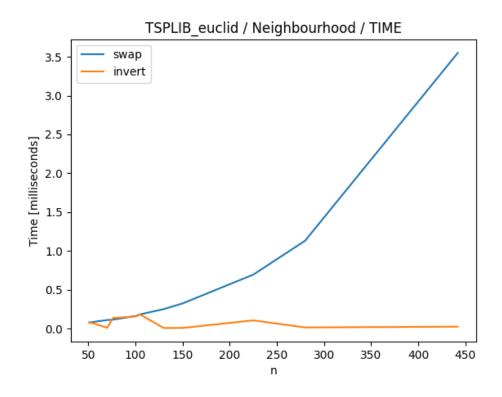
# 4 Porównanie sąsiedztwa insert i swap

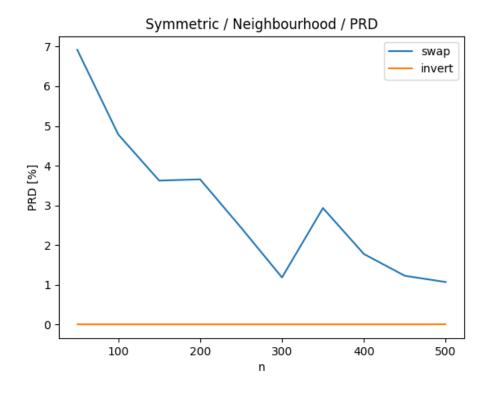
## 4.1 Dane z TSPLIB

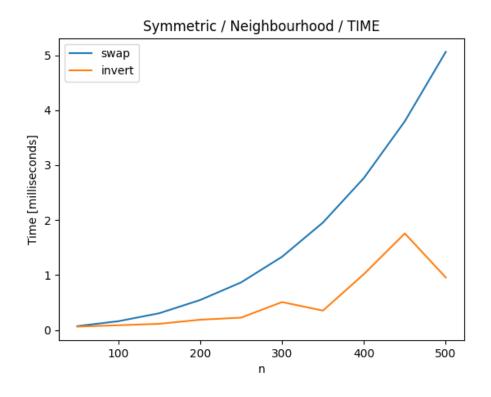


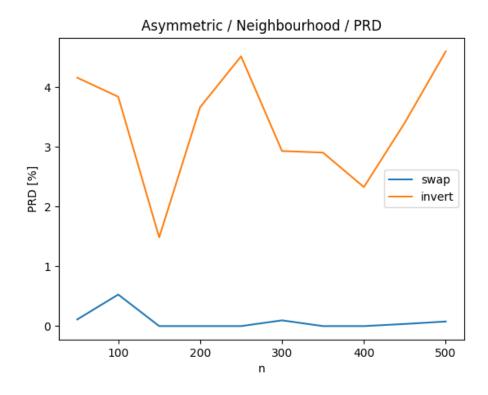


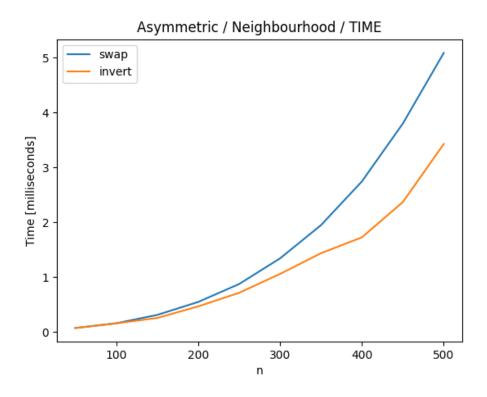


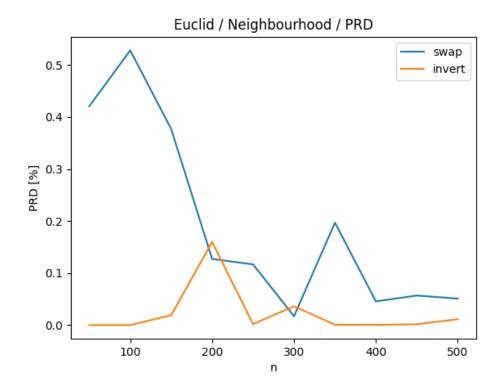


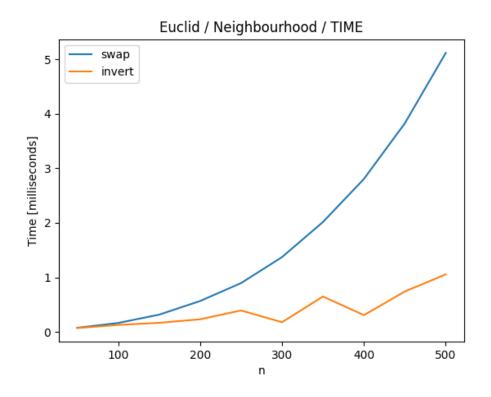








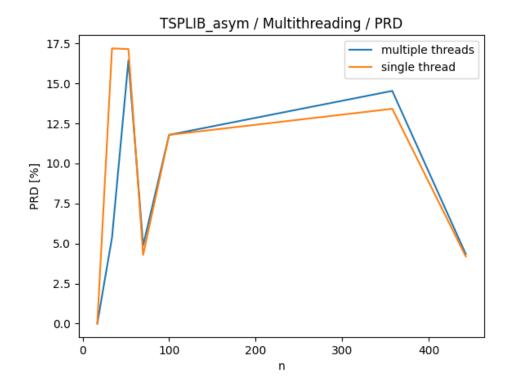


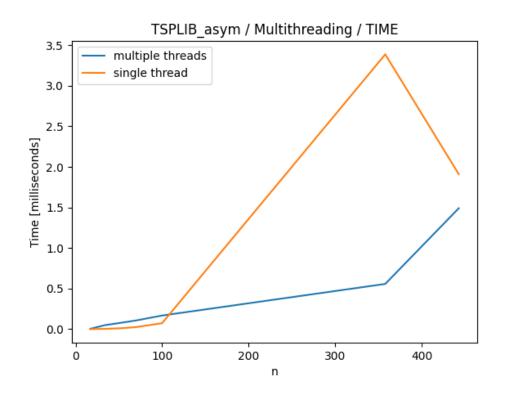


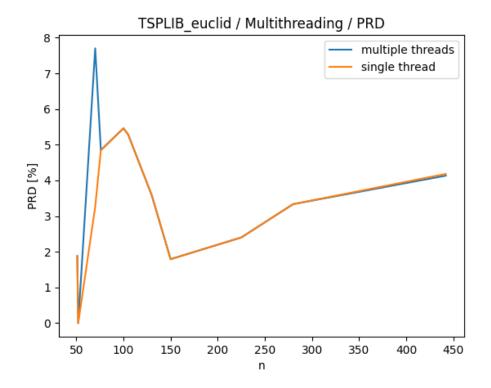
### 4.4 Tabele

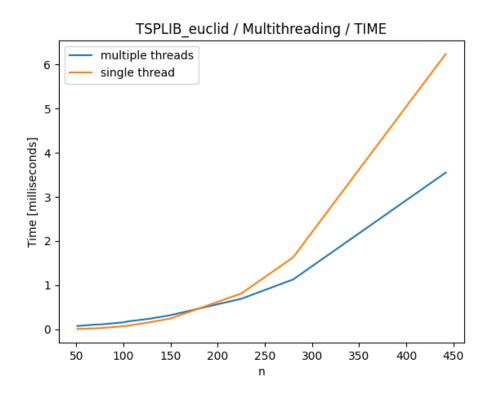
# 5 Porównanie wersji wielowątkowej i wersji jednowątkowej

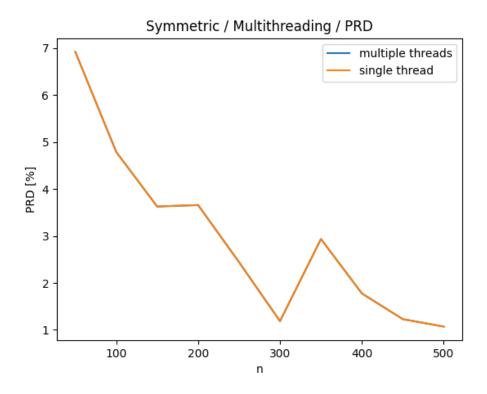
## 5.1 Dane z TSPLIB

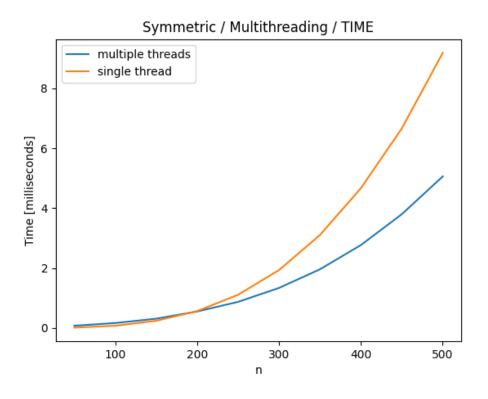


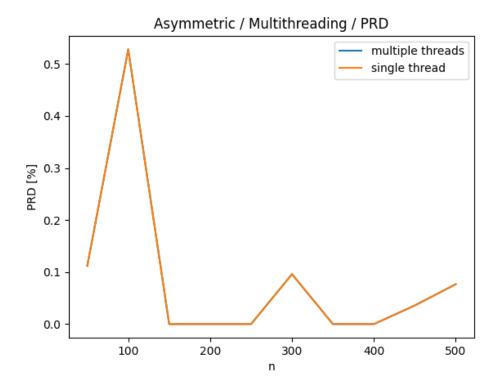


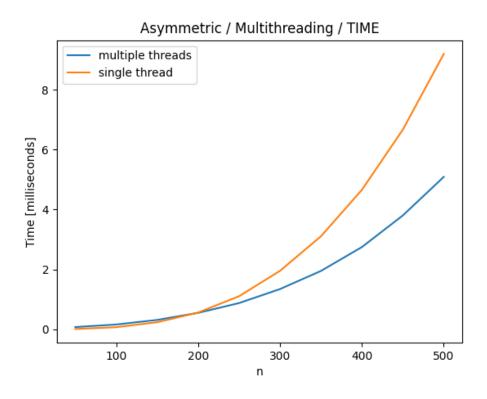


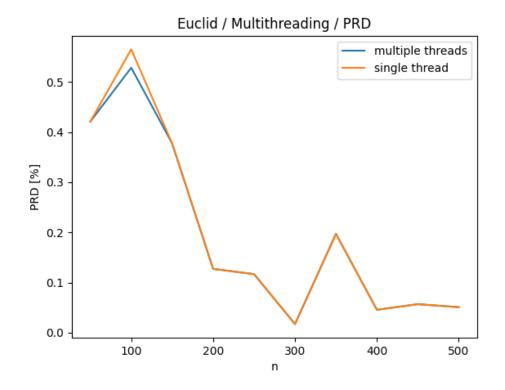


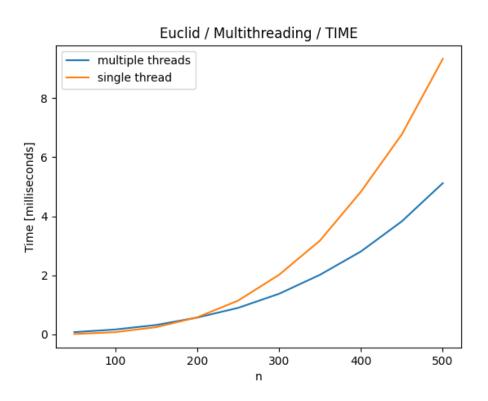












## 5.4 Tabele

## 6 Wnioski