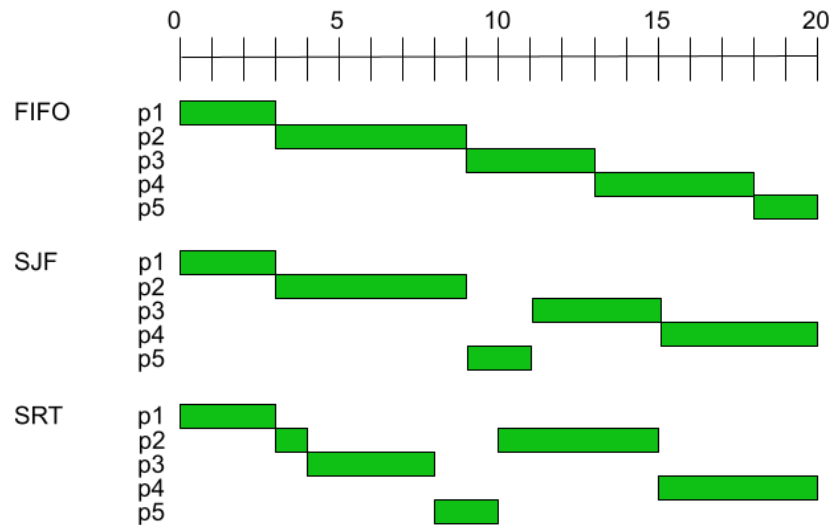


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3.2.1

(a)

For the 5 processes described below, draw a timing diagram showing when each process executes under FIFO, SJF, and SRT.



(b)

Determine the ATT for each scheduling algorithm for the 5 processes.

Process	p1	p2	p3	p4	p5	ATT
FIFO: Turnaround times	3	7	9	12	12	8.60
SJF: Turnaround times	3	7	11	14	3	7.60
SRT: Turnaround times	3	13	4	14	2	7.20

3.2.2 (a)

Starting at time 0, a new process p of length 3 arrives every 4 time units.
 Starting at time 1, a new process q of length 1 arrives every 4 time units.
 Determine the ATT under FIFO, SJF, and SRT.

	p	q	ATT
	waiting+running	waiting+running	
FIFO, SJF	0+3	2+1	$6/2 = 3$
SRT	1+3	0+1	$5/2 = 2.5$

3.2.3

(a)

i	0	1	2	3	4	5
T_i	7	5	6	15	15	15
S_i	7	7	5.4	5.88	13.176	14.635

(b)

i	0	1	2	3	4	5
T_i	7	5	6	15	15	15
S_i	7	7	6	6	10.5	12.75

3.2.4

(a)

$$\begin{aligned}
 S_5 &= 0.5 * T_4 + 0.5 * S_4 && \text{Thus } T_4 \text{ contributes 50\% to } S_5 \\
 &= 0.5 * T_4 + 0.5 * (0.5 * T_3 + 0.5 * S_3) \\
 &= 0.5 * T_4 + 0.25 * T_3 + 0.25 * S_3 && \text{Thus } T_3 \text{ contributes 25\% to } S_5 \\
 &= 0.5 * T_4 + 0.25 * T_3 + 0.25(0.5 * T_2 + 0.5 * S_2) \\
 &= 0.5 * T_4 + 0.25 * T_3 + 0.125 * T_2 + 0.125 * S_2 && \text{Thus } T_2 \text{ contributes 12.5\% to } S_5 \\
 &= 0.5 * T_4 + 0.25 * T_3 + 0.125 * T_2 + 0.125(0.5 * T_1 + 0.5 * S_1) \\
 &= 0.5 * T_4 + 0.25 * T_3 + 0.125 * T_2 + 0.0625 * T_1 + 0.0625 * S_1 && \text{Thus } T_1 \text{ contributes 6.25\% to } S_5.
 \end{aligned}$$

3.3.1

(a) Determine the average turnaround time, ATT, when the quantum is $Q = 1$ time unit.

Time	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
p1	x			x			x			x		x		x	x	x
p2		x			x			x								
p3			x			x			x		x		x			

ATT = $(16 + 8 + 13)/3 = 12.33$

(b) Determine the average turnaround time, ATT, when the quantum is $Q = 3$ time units.

Time	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
p1	x	x	x							x	x	x			x	x
p2				x	x	x										
p3							x	x	x				x	x		

ATT = $(16 + 6 + 14)/3 = 12$

3.3.2 (a) How long will the execution take on a machine with n CPUs?

T ms

(b) How long will the execution take on a single CPU machine when the context switch overhead is zero?

$n \cdot t$ ms

(c) How long will the execution take on a single CPU machine when:

- the length of the time quantum is Q ms
- the time to perform each context switch is S ms

of context switches: $n \cdot T / Q$, total execution: $n \cdot T + (n \cdot T / Q) \cdot S$ ms

(d) Repeat the previous calculation using $n=5$, $T=10,000$, $Q=100$, $S=10$.

$5 \cdot 10,000 + (5 \cdot 10,000 / 100) \cdot 10 = 55,000$ ms.

3.3.3

(a) What should be the quantum size Q such that the gap between the end of one quantum and the start of the next quantum of any process does not exceed M ms?

Between the end of one quantum and the start of the next quantum, $n - 1$ processes will each execute one quantum: $(n - 1)Q$.

Additionally, n context switches will be needed for the interrupted process to restart: nS .

The sum of the two times must not exceed the limit M , i.e., $(n - 1)Q + nS \leq M$. The largest Q that satisfies this condition is $Q = (M - nS) / (n - 1)$.

(b) For $n = 5$, $S = 10$, and $M = 450$, $M = 90$, $M = 50$, determine:

- The corresponding values of Q
- The percentage of CPU time wasted on context switching

M	Q	% wasted
450	100	$100 \cdot 10 / (100 + 10) = 9.09\%$
90	10	$100 \cdot 10 / (10 + 10) = 50\%$
50	0	$100 \cdot 10 / (0 + 10) = 100\%$

3.3.4

(a) $T < Q$

When $T < Q$, the quantum never expires. The execution repeats the timing sequence: T, S, T, S,

Thus for every T ms of execution, T + S ms of CPU time are needed. The fraction of wasted CPU time is $S/(T + S)$.

(b) $T \gg Q$

When $T \gg Q$, each execution of T needs many quanta to complete. Each quantum is followed by a context switch. The fraction of wasted CPU time is $S/(Q + S)$.

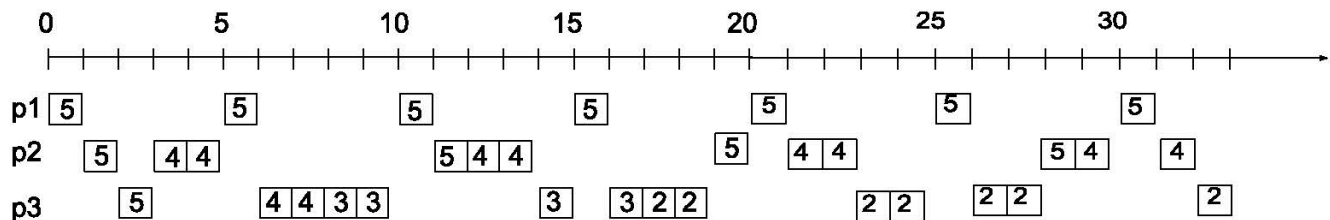
(c) Q approaches 0

When Q approaches 0, the value of $S/(Q + S)$ approaches 1 (overhead approaches 100%).

(d) Under what condition will the wasted fraction of CPU time be 50%?

50% overhead means $S/(Q + S) = 0.5$, which is true when $Q = S$.

3.3.5 (a) Draw a timing diagram for the first 33 ms. On each of the 3 lines (one per process) show when the process is running and at which priority level.



(b) Determine the ATT for each process.

P1 – ATT = 1

P2 – ATT = 4

P3 – ATT = 32

3.4.1 (a) Determine if a feasible schedule exists.

Determine if a feasible schedule exists.

$$5/20 + 10/100 + 42/120 = 0.25 + 0.1 + 0.35 = 0.7$$

$0.7 < 1$ and thus a feasible schedule exists.

(b)

$$5/20 + 10/100 + 42/120 = 0.25 + 0.1 + 0.35 = 0.7$$

$0.7 < 1$ and thus a feasible schedule exists.

(c)

The CPU fraction used by the 3 processes p1 through p3 is already 0.7.

Adding any more processes to the mix would likely result in an infeasible schedule.

3.4.2

(a)

For each case, determine if a feasible schedule is likely to be generated by:

- RM
- EDF

Case 1: Both RM & EDF. $U = 3/50 + 70/1000 + 5/40 = 0.255$

Case 2: Neither RM Nor EDF. $U = 5/50 + 5/10 + 1/4 = 1.05$

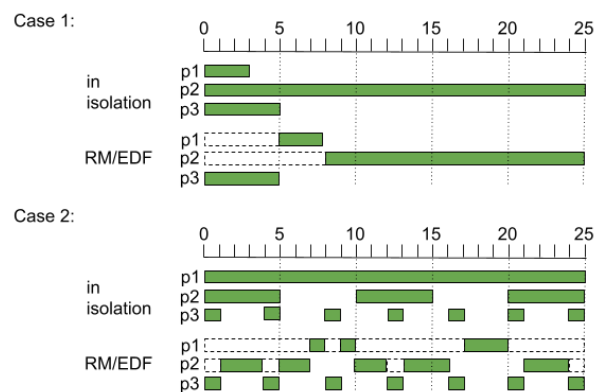
Case 3: EDF will produce feasible schedule. RM will most likely not produce a feasible schedule.

$$U = 5/20 + 7/10 + 4/100 = 0.99$$

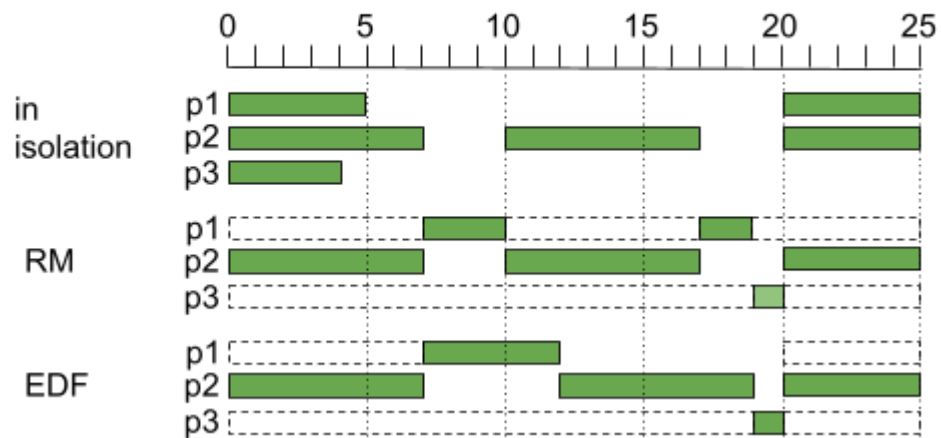
(b) For each case, determine if a feasible schedule is likely to be generated by:

- RM

- EDF



Case 3:



3.5.1

(a)

Show at which priority level the process p is executing during each of the 18Q time units.

Running interval	Priority levels	Explanation
1st 3Q	4, 3, 3	p starts at level 4, runs for 1Q and drops to level 3 for the remaining 2Q units.
Next Q	2	p drops to level 2 and runs for 1Q to execute the blocking operation.
Next 7Q	3, 3, 2, 2, 2, 2, 1	p re-enters at level 3 and runs for 2Q, then for 4Q at level 2, then drops to level 1 for 1Q.
Next Q	1	p continues running at level 1 for 1Q to execute the blocking operation.
Next 3Q	2, 2, 2	p re-enters at level 2 and completes all 3Q at the same level 2.
Next Q	2	p continues running at level 2 for 1Q to execute the blocking operation.
Next 2Q	3, 3	p re-enters at level 3 and completes both units at the same level 3.