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# Introduction: The Quest for Quantum Advantage

**Context.** Quantum Information Science seeks tasks where quantum systems fundamentally outperform classical ones.

**This Work's Focus.** A physically natural problem involving **permutations of particles**.

- No Oracles.
- No Computational Promises.
- **Sharp, Non-asymptotic separation.**

# The Problem: Parity of a Hidden Permutation



**Problem Definition.** Given a set of  $n$  particles subject to a hidden permutation  $\sigma \in S_n$ , determine the parity:

- Is  $\sigma$  **Even**? (e.g., identity, 3-cycles)
- Is  $\sigma$  **Odd**? (e.g., single swap)

**Goal:** Identify parity with **Certainty** ( $P_s = 1$ ).

# The Classical Limit

Case 1: Distinguishable ( $d = n$ )



Unique Parity ✓

Case 2: Indistinguishable ( $d < n$ )



Ambiguous Parity ?  
(Even or Odd?)

## Full Distinguishability ( $n$ labels)

- perfectly reconstruct the permutation
- determine its parity
- $P_s = 1$

## Insufficient Labels ( $< n$ labels)

- perfect identification becomes impossible
- indistinguishable opposite-parity counterparts
- limited to random guessing
- $P_s = 1/2$

## Theorem: The Quadratic Quantum Advantage

**Theorem (Main Result).** Perfect parity identification ( $P_s = 1$ ) is achievable if

$$d \geq d_{\min} := \lceil \sqrt{n} \rceil$$

Below this threshold, assuming all permutations are equally likely, parity remains indistinguishable and random guessing is the best one can do ( $P_s = 1/2$ ).

Below  $d = \lceil \sqrt{n} \rceil$ , even Quantum Mechanics cannot help ( $P_s = 1/2$ ).

## Example: 4 Qubits ( $n = 4, d = 2$ )

**Physical Setup.** Consider a system of **four spin-1/2 particles**. Alice acts by reshuffling their positions *without* disturbing the spin degrees of freedom.

### Schur-Weyl Decomposition.

$$(\mathbb{C}^2)^{\otimes 4} \cong (\mathcal{K}_2 \otimes \mathcal{H}_2) \oplus (\mathcal{K}_1 \otimes \mathcal{H}_1) \oplus (\mathcal{K}_0 \otimes \mathcal{H}_0)$$

- $\mathcal{H}_j$ :  $SU(2)$ -invariant subspaces (labeled by total spin  $j$ ).
- $SU(2)$  acts *only* on  $\mathcal{H}_j$  and trivially on  $\mathcal{K}_j$ .

### Multiplicity Spaces ( $\mathcal{K}_j$ ).

- $j = 2$ :  $\mathcal{K}_2 \cong \mathbb{C}$  (1 copy)
- $j = 1$ :  $\mathcal{K}_1 \cong \mathbb{C}^3$  (3 copies)
- $j = 0$ :  $\mathcal{K}_0 \cong \mathbb{C}^2$  (2 copies)

# Symmetry and Representation: The $S_4$ Perspective

## Permutational Symmetry ( $S_4$ )

- Reshuffling 4 qubits  $\longleftrightarrow$  “Rotation” within subspaces  $K_j$ .
- These spaces carry **irreducible representations (irreps)** of  $S_4$ .
- Note: No proper subspace is invariant under  $\sigma|\psi\rangle$ .

## Schur-Weyl Duality Structure

- Agent's operations are restricted to permutations.
- **Consequence:** Action is **trivial** on spin sectors  $H_j$ .
- **Interpretation:**
  - $K_j$ : Irrep spaces (Dynamic part).
  - $H_j$ : Multiplicity spaces (Static part).

# Restriction to Even Permutations ( $A_4 \subset S_4$ )

**Group Restriction:** Restrict agent actions to *even permutations* ( $A_4$ ).

- $\mathcal{K}_1$  and  $\mathcal{K}_2$ : Remain irreducible.
- $\mathcal{K}_0$ : **Decomposes** into two orthogonal 1D invariant subspaces:

$$\mathcal{K}_0 \xrightarrow{A_4} \mathcal{K}_{0a} \oplus \mathcal{K}_{0b}$$

## Action of Odd Permutations

( $\sigma_{\text{odd}} \in S_4 \setminus A_4$ )

Any odd permutation swaps the subspaces:

$$\sigma_{\text{odd}} : \mathcal{K}_{0a} \longleftrightarrow \mathcal{K}_{0b}$$

Note: This explains why  $\mathcal{K}_{0a/b}$  are not invariant under the full group  $S_4$ .

## Parity Identification Protocol

Prepare initial state  $|\psi_e\rangle \in \mathcal{K}_{0a} \otimes \mathcal{H}_0$ .  
The odd/even sectors become orthogonal:

$$\sigma_{\text{odd}}|\psi_e\rangle \perp \sigma_{\text{even}}|\psi_e\rangle \quad (1)$$

**Result:** Parity can be identified with certainty via projective measurement.

# Explicit Construction: The M4 State

## Construction via Projection

We construct  $|\psi_e\rangle$  by projecting the basis state  $|0011\rangle$  (with magnetic quantum number  $m = 0$ ) onto the invariant subspace  $\mathcal{K}_{0a} \otimes \mathcal{H}_0$ :

$$|\psi_e\rangle = \mathcal{P}_{0a}|0011\rangle, \quad \text{where } \mathcal{P}_{0a} = \frac{1}{|A_4|} \sum_{\sigma \in A_4} \bar{\chi}_{0a}(\sigma) \sigma$$

## Computational Basis Expansion:

$$|\psi_e\rangle = (|0011\rangle + |1100\rangle) + \zeta_3 (|0101\rangle + |1010\rangle) + \zeta_3^2 (|0110\rangle + |1001\rangle)$$

- Also known as the **M4 (or D4) state**.
- A Fourier-type combination of **Dicke states**.

$\zeta_3 = e^{i2\pi/3}$  (Cubic root of unity).

# Labeling Invariant Subspaces: Young Diagrams



**Partitions & Young Diagrams (YDs).** Subspaces are labeled by partitions  $\lambda \vdash n$ , visualized as YDs with  $n$  boxes.

$$\lambda = [\lambda_1, \lambda_2, \dots]$$

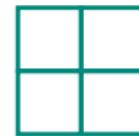
$$(\lambda_1 \geq \lambda_2 \geq \dots \geq 0, \quad \sum \lambda_p = n)$$

- $S_n$  **Irreps** ( $\mathcal{K}_\lambda$ ): Correspond to any distinct  $\lambda \vdash n$ .
- $SU(d)$  **Irreps** ( $\mathcal{H}_\lambda$ ): Restricted by dimension  $d$ :

$$\ell(\lambda) \leq d$$

(The YD cannot have more than  $d$  rows)

# Labeling Subspaces: Young Diagrams (YDs)



**Partitions of  $n = 4$**

$[4], [3, 1], [2, 2], [2, 1, 1], [1^4]$

**Crucial Subspaces for  $n = 4$**

For  $d \geq 4$ , all 5 irreps appear.

- $\mathcal{K}_{[4]}$ : **Fully Symmetric** (Trivial, +1).
- $\mathcal{K}_{[1^4]}$ : **Fully Antisymmetric** (Sign,  $\pm 1$ ).

*Mixing these enables perfect parity identification.*

$$|\psi_e\rangle = |\psi_{\text{sym}}\rangle + |\psi_{\text{ant}}\rangle$$

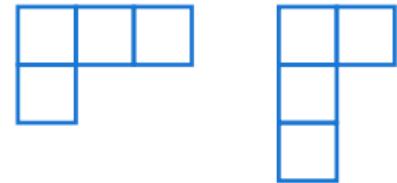
$$|\psi_{\text{sym}}\rangle \in \mathcal{K}_{[4]} \otimes \mathcal{H}_{[4]}$$
$$|\psi_{\text{ant}}\rangle \in \mathcal{K}_{[1^4]} \otimes \mathcal{H}_{[1^4]}$$

## Case 2: Restricted Dimensions ( $d = 3$ )

### The Constraint Problem ( $d = 3$ )

Fully antisymmetric partition  $[1^4]$  is **forbidden**.

$$\ell([1^4]) = 4 > d = 3 \implies |\psi_{\text{ant}}\rangle \text{ does not exist.}$$



**Solution: Conjugate Pairs.** We use conjugate partitions  $\lambda$  and  $\lambda^T$  (reflecting YD along diagonal). They satisfy a crucial representation property:

$$D_{\lambda^T}(\sigma) = \text{sign}(\sigma) D_\lambda(\sigma)$$

**New Construction:** Combine partitions  $\lambda = [3, 1]$  and  $\lambda^T = [2, 1^2]$ :

$$|\psi_e\rangle = |\psi_{[3,1]}\rangle + |\psi_{[2,1^2]}\rangle$$

Even though we lack the trivial/sign singlets, these higher-dimensional irreps mimic the sign-flipping behavior required for orthogonality.

## Case 3: Self-Conjugate Irreps & The General Rule

### Definition: Self-Conjugate Partitions

Its Young Diagram is symmetric along the main diagonal ( $\lambda = \lambda^T$ ).

- Example ( $n = 4$ ):  $\lambda = [2, 2]$  (The square diagram).

### Mechanism: Splitting under $A_n$

Unlike generic irreps, self-conjugate ones split when restricted to even permutations ( $A_n$ ):

$$\mathcal{K}_\lambda \xrightarrow{A_n} \mathcal{K}_{\lambda a} \oplus \mathcal{K}_{\lambda b}$$



Action of Odd Permutations:

$$\mathcal{K}_{\lambda a} \xrightarrow{\sigma_{\text{odd}}} \mathcal{K}_{\lambda b}$$

## Example: Five Qutrits ( $d = 3$ ) & Non-Trivial Coefficients

### Setup: Self-Conjugate Irrep $\mathcal{K}_{[3,1^2]}$

Consider  $(\mathbb{C}^3)^{\otimes 5}$ . The partition  $\lambda = [3, 1^2]$  is self-conjugate with dimension  $d_{[3,1^2]} = 6$ .

- Parity can be identified with certainty using any state  $|\psi_e\rangle$  in the invariant subspace  $\mathcal{K}_{[3,1^2]a} \otimes \mathcal{H}_{[3,1^2]}$ .

### Mechanism: Projector Construction

The projector  $\mathcal{P}_{[3,1^2]a}$  onto the subspace (analogous to  $\mathcal{P}_{0a}$ ) is given by:

$$\mathcal{P}_{\lambda a} = \frac{d_\lambda/2}{|A_n|} \sum_{\sigma \in A_n} \bar{\chi}_{\lambda a}(\sigma) \sigma$$

$$|\psi_e\rangle = 3(|00012\rangle - |00021\rangle) - |00102\rangle + |00120\rangle + |00201\rangle - |00210\rangle \\ - |01002\rangle + |01020\rangle - |10002\rangle + |10020\rangle + \dots \\ + \sqrt{5}(|01200\rangle - |02100\rangle - |10200\rangle + |12000\rangle + \dots)$$

# Proof (I): Hypothesis Testing & Schur-Weyl Decomposition

**Binary Discrimination Problem.** Distinguish between even ( $\rho_0$ ) and odd ( $\rho_1$ ) mixtures of permutations.

$$\rho_0 = \frac{1}{|A_n|} \sum_{\sigma \in A_n} \sigma |\psi_e\rangle \langle \psi_e| \sigma^{-1}$$

$$\rho_1 = (12)\rho_0(12)$$

- Perfect discrimination requires orthogonal supports:  
 $\text{Tr}(\rho_0\rho_1) = 0$ .
- If  $\rho_0 = \rho_1$ , then  $P_s = 1/2$  (random guessing).

## Schur-Weyl Decomposition.

The Hilbert space decomposes into sectors labeled by partitions  $\lambda \vdash n$ :

$$(\mathbb{C}^d)^{\otimes n} \cong \bigoplus_{\lambda \vdash n, \ell(\lambda) \leq d} \mathcal{K}_\lambda \otimes \mathcal{H}_\lambda$$

## Partition $\lambda$ Into 3 Sectors.

- 1  $\lambda = \lambda^T$  (Self-conjugate).
- 2  $\lambda \neq \lambda^T$  and  $\ell(\lambda^T) \leq d$  (Conjugate pairs).
- 3  $\ell(\lambda^T) > d$  (Restricted sector).

## Proof (II): Action of $A_n$ and Block Structure

**Lemma (Block-diagonal Form).** Under the restriction to the Alternating Group  $A_n$ , the irreps split differently in each sector. By Schur's lemma,  $\rho_0$  takes the block-diagonal form:

$$\rho_0 = \rho_0^{(i)} + \rho_0^{(ii)} + \rho_0^{(iii)}$$

- 1  $[\lambda = \lambda^T]$ : Irreps split  $\mathcal{K}_\lambda \downarrow_{A_n} = \mathcal{K}_{\lambda a} \oplus \mathcal{K}_{\lambda b}$ .

$$\rho_0^{(i)} = \sum_{\lambda \in R_i} \left( \frac{\mathbb{1}_{\lambda a}}{d_\lambda/2} \otimes \Phi^{\lambda a} + \frac{\mathbb{1}_{\lambda b}}{d_\lambda/2} \otimes \Phi^{\lambda b} \right)$$

If either  $\Phi^{\lambda a} = 0$  or  $\Phi^{\lambda b} = 0$ , then  $\rho_0^{(1)}$  and  $\rho_1^{(1)}$  have orthogonal supports;

## Proof (II): Action of $A_n$ and Block Structure

2  $[\lambda \neq \lambda^T]$ : Irreps appear in conjugate pairs,  $(\mathcal{K}_\lambda \oplus \mathcal{K}_{\lambda^T}) \downarrow_{A_n} = \mathcal{K}_\lambda \otimes \mathbb{C}^2$ .

$$\rho_0^{(ii)} = \sum_{\lambda \in R_{ii}} \frac{\mathbf{1}_\lambda}{d_\lambda} \otimes \sum_{k=1}^{d_\lambda} |\phi_k^\lambda\rangle \langle \phi_k^\lambda|, \quad |\phi_k^\lambda\rangle \in \mathbb{C}^2 \otimes \mathcal{H}_\lambda$$

$$|\phi_k^\lambda\rangle = \sum_{p=\pm} |s_p\rangle |\phi_{k,p}^\lambda\rangle \in \mathbb{C}^2 \otimes \mathcal{H}_\lambda.$$

$$|s_\pm\rangle \xrightarrow{\sigma} \pm |s_\pm\rangle \quad \text{for } \sigma \in S_n \setminus A_n$$

$$\text{Tr} \left( \rho_0^{(ii)} \rho_1^{(ii)} \right) = \sum_{\lambda \in R_{ii}} \sum_{k,l=1}^{d_\lambda} \frac{\left| \langle \phi_{k,+}^\lambda | \phi_{l,+}^\lambda \rangle - \langle \phi_{k,-}^\lambda | \phi_{l,-}^\lambda \rangle \right|^2}{d_\lambda}$$

$\implies$  Supports are orthogonal if  $\langle \phi_{k,+}^\lambda | \phi_{l,+}^\lambda \rangle = \langle \phi_{k,-}^\lambda | \phi_{l,-}^\lambda \rangle$  for all  $\lambda, k, l$ .

## Proof (II): Action of $A_n$ and Block Structure

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Restriction has no effect ( $\mathcal{K}_\lambda \downarrow_{A_n} = \mathcal{K}_\lambda$ ).

$$\rho_0^{(iii)} = \sum_{\lambda \in R_{iii}} \frac{\mathbb{1}_\lambda}{d_\lambda} \otimes \Phi^\lambda$$

$\rho_0^{(iii)}$  is invariant under all of  $S_n$ , so  $\rho_0^{(iii)} = \rho_1^{(iii)}$ .

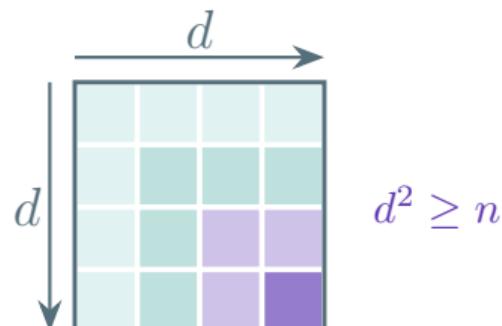
## Proof (III): Distinguishability Conditions

$d \geq d_{\min} := \lceil \sqrt{n} \rceil$ . ① and/or ② hold, enabling perfect distinguishability.

$d \leq d_{\min} := \lceil \sqrt{n} \rceil$ . ① and ② fail, only ③ remains

$$\rho_0 = \rho_1 \text{ and } P_s = 1/2.$$

This geometrically corresponds to the impossibility of fitting certain Young Diagrams within a  $d \times d$  box when  $d$  is too small.



# Entanglement Requirement for Parity Identification

**Question.** How entangled must a parity-detecting state be?

**Metric: geometric measure of entanglement (GME)**

$$E(|\psi\rangle) = 1 - \max_{|\phi\rangle \in \text{SEP}} |\langle \phi | \psi \rangle|^2$$

**Minimum Required Entanglement:**

$$E_{\lambda a} = 1 - \max_{|\phi\rangle \in \text{SEP}} \langle \phi | \mathcal{P}_{\lambda a} | \phi \rangle$$

This lower bounds the entanglement of any parity-detecting state supported on this subspace.

**Table:** Minimum required ( $E_{\lambda a}$ ) vs. known maximal ( $E_{\max}$ ) entanglement.

$n$	$E_{\lambda a}$	$E_{\max}$
3	5/9	5/9 <sup>†</sup>
4	7/9	7/9 <sup>†</sup>
5	17/20	$\approx 0.96$

<sup>†</sup> indicates proven maxima.

**Observation:**

maximal for  $n = 3, 4$ , and close to the maximum for  $n = 5$ .

# Summary and Outlook

## A Sharp Quantum Advantage

Achievable with  $d = \lceil \sqrt{n} \rceil$ .

Relies purely on permutation symmetry without invoking specific dynamics.

Requires **no ancillas** (increasing multiplicities does not help).

Demands entanglement close to the theoretical maximum.

## Practical Implementations:

- State preparation via experimentally accessible Hamiltonians.
- Designing specific quantum circuits and measurements.

## Theoretical Extensions:

- Assessing robustness of the  $\sqrt{n}$  advantage under noise.
- Exploring other symmetry groups.