

Amath 515, March 8<sup>th</sup>

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## Recap:

$$\min_x g(x) + h(Ax - b) - \langle c, x \rangle$$

$$\max_z -\langle z, b \rangle - h^*(z) - g^*(-A^T z + c)$$

Problem without red additions will lead us to the Chambolle-pock algorithm for  $\min_x g(x) + h(Ax)$

→ match the algo in their paper.

\* Add red terms, will get a slightly modified version we need for Lasso on hw.

Opt conditions:

$$\begin{bmatrix} 0 \\ \lambda_0 \end{bmatrix} \in \left( \begin{bmatrix} \partial g & 0 \\ 0 & \partial h^* \end{bmatrix}^+ + \begin{bmatrix} 0 & A^T \\ -A & 0 \end{bmatrix} \right) \begin{bmatrix} \bar{x} \\ \bar{z} \end{bmatrix}$$

Iterations:

\* ① Basic:  $w^+ = (I + T)^{-1} w$

\*\* ② 'Tricky':  $w^+ = (I + P^\top T)^{-1} w$   
Chambolle-Pock iteration

Shown  $T$  is monotone

$$\langle x - y, u - v \rangle \geq 0$$

$x, y \in T_x$

$v \in T_y$

Plan:  $\langle x - y, T_x - T_y \rangle \geq 0$

① Show  $(I + T)^{-1}$  is FNE

② Pf sketch for convergence of

\*, \*\*

③ Understand 'Trick', and  $P$

Thm: T monotone  $\Rightarrow$  S =  $(I+T)^{-1}$  FNE

PF: pick  $x, u \in S_x, y, v \in S_y$

$u \in S_x$  means  $u \in (I+T)^{-1}x$

multiply by  $I+T$ , get

$$x \in \underline{(I+T)u}$$

$$u + Tu$$

$$\underline{x - u \in Tu}$$

$v \in S_y$  means  $y - v \in Tv$

T monotone means flat

$$\langle u - v, (x - u) - (y - v) \rangle \geq 0 \quad \checkmark$$

$$\begin{array}{ccc} \uparrow & & \uparrow \\ Tu & & Tv \\ x - S_x & & y - S_y \end{array} \quad \star$$

$$\langle S_x - S_y, (I-s)x - (I-s)y \rangle \geq 0$$

Choice of  $x, y$  arbitrary, so concl.  
holds for all  $x, y$ .

\* holding for all  $x, y$  means

$s$  is FNE by how 3b.

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Thm: Krasnosel'skii - Mann

Informal: Any Lip- $\frac{1}{\epsilon}$  (NE) operator  
can be tweaked so repeated application  
of it converges to its fixed pts,

Cor: If  $s$  is FNE, then don't  
met 'tweak' so

$w^+ = \lim w$  conv to fixed points

of  $s$ .

## Formal KM statement:

Let  $R$  be NE,  $\|R_x - R_y\| \leq \|x - y\|$ .

Suppose  $\text{Fix } R = \{x : R_x = x\} \neq \emptyset$ .

For  $\lambda \in (0, 1)$ , define 'twist'.

$$R_\lambda w = (1-\lambda)w + \lambda R w.$$

①  $\text{Fix } R = \text{Fix } R_\lambda$  (same fix pfs)

②  $w^+ = R_\lambda w$  conv. to  $\text{Fix } R$   
at a sublinear rate.

Pf details: by hw 2B, we have

for any  $\bar{w} \in \text{Fix } R$ ,

$$\|R_\lambda w - \bar{w}\|^2 \leq \|w - \bar{w}\|^2 - \lambda(1-\lambda) \|R w - \bar{w}\|^2$$

$w^{k+1} = R_\lambda w^k$  our iteration, can  
get

$$\lambda(1-\lambda) \underbrace{\|w^* - R w^k\|^2}_{=} \underbrace{\|w^k - \bar{w}\|^2 - \|w^{k+1} - \bar{w}\|^2}_{\rightarrow}$$

Telescopes

add up across  $i = 0, \dots, K$ , get

$$\min_{0 \leq i \leq K} \|R w^i - w^i\|^2 \leq \frac{1}{K} \sum_{i=1}^K \|R w^i - \bar{w}\|^2$$

$$\frac{\|w_0 - \bar{w}\|^2}{K(\lambda)(1-\lambda)}$$

Big picture: some work we do not show here.

- ①  $\{w^k\}$  stay bounded
- ②  $\{w^k\}$  have unique limit point

- ③ Our detail above gives a sublinear conv' rate to this limit point.

Cor: If  $s$  is PNE, RM simplifies to  $w^+ = \underline{s} w$ .

Pf:  $s$  FNE  $\Leftrightarrow$   $2s - I$  NE by HW 3C

Apply RM to  $2s - I$ ,  $\lambda = \frac{1}{2}$

$$\underbrace{(2s - I)}_{R \times w} \underbrace{\frac{1}{2}w}_{= s w} = \frac{1}{2}w + \frac{1}{2}(2s - I)w$$

Summary:  $T$  monotone, then

$w^+ = (I + T)^{-1}w$  will converge to fixed pts of  $(I + T)^{-1}$ , which satisfy  $0 \in T\bar{w}$

Asider.  $\min_x f(x)$ ,  $f$  convex, closed,  
proper, --

$$0 \in \frac{\partial f}{T}(x)$$

Definition:  $x^+ \in (I + \partial f)^{-1} x$   
what is this??

$$x \in (I + \partial f)x^+$$

$$x \in x^+ + \partial f(x^+)$$

$$\frac{x - x^+ \in \partial f(x^+)}{\text{opt or prox}}$$

$$x^+ = \text{prox}_f(x)$$

\* proximal point algorithm

Why is this not as satisfying  
as it seems?

$$f(x) = g(x) + h(Ax)$$

$$x^+ = (\text{IFT})^{-1} x$$

$$x^+ = \underbrace{\text{prox}_{g+h(A \cdot)}(x)}_{\substack{\downarrow \\ \partial f}}$$

$$\begin{aligned} &= \underset{y}{\operatorname{argmin}} \quad \frac{1}{2} \|x-y\|^2 + g(y) \\ &\qquad\qquad\qquad \xrightarrow{\hspace{10em}} + h(Ay) \end{aligned}$$

Our  $T$  was

$$\left[ \begin{array}{c} \alpha^* \\ 2\alpha^* \end{array} \right] + \left[ \begin{array}{cc} 0 & -A^T \\ -A & 0 \end{array} \right]$$

If we try  $(I+T)^{-1}$ , we will find  
we can't easily implement it.

Trick: want to use a 'preconditioner'  $P$ .

Let  $P$  be any pos def matrix,  
we can define a new inner product:

$$\langle x, y \rangle_P : \underline{x^T P y}$$

Fact:  $P^{-1}T$  is monotone with respect  
to  $\langle x, y \rangle_P \iff T$  is monotone.

Allows us to use iteration

\*  $w^+ = (I + P^{-1}T)^{-1} w$  for

any  $P$  pos. def.  
 $\hookrightarrow w \in (I + P^{-1}T) w^+$

$$w - w^+ \in P^{-1}T w^+$$

\*  $P(w - w^+) \in Tw^+$

Our optimality conditions:

$$\begin{bmatrix} 0 \\ 0 \end{bmatrix} \in \begin{bmatrix} A^T z + \partial g(x) \\ -Ax + \partial L^*(z) \end{bmatrix}, w = \begin{bmatrix} x \\ z \end{bmatrix}$$

$T_w$

Clever P:

$$\begin{bmatrix} \frac{1}{\delta} I & -A^T \\ -A & \frac{1}{\delta} I \end{bmatrix}$$

① Can pick  $\delta$  to make sure

P is pos def

② P was designed to make the iteration  $w_t = (I + P^{-1})^{-1} w$  implementable.

Let's now compute  $w^+$  via

$$P(w - w^+) \in T w^+$$

$$\underbrace{\begin{bmatrix} \frac{1}{\gamma} I & -A^T \\ -A & \frac{1}{\gamma} I \end{bmatrix}}_P \underbrace{\begin{bmatrix} x - x^+ \\ z - z^+ \end{bmatrix}}_{w - w^+} \in \underbrace{\begin{bmatrix} A^T z^+ + \partial g(x^+) - c \\ -A x^+ + \partial h^*(z^+) + b \end{bmatrix}}_{T w^+}$$

Row 1:

$$\frac{1}{\gamma} (x - x^+) - A^T (z - z^+) \in \cancel{A^T z^+ + \partial g(x^+)} - c$$

$$\frac{1}{\gamma} (x - x^+) - A^T z \in \partial g(x^+)$$

$$\frac{1}{\gamma} (x - \gamma A^T z - x^+) \in \partial g(x^+)$$

$$\underbrace{x^+ = \text{prox}_{\gamma g}(x - \gamma A^T z)}$$

$$\text{prox}_{\gamma g}(x - \gamma(A^T z - c))$$

Row 2:

$$-A(x - x^+) + \frac{1}{\gamma} (z - z^+) \leftarrow -Ax^+ + \cancel{\gamma h^*(z^+)}$$

$$\frac{1}{\gamma} (z - z^+) - \cancel{Ax^+ + 2Ax^+} \leftarrow \cancel{-Ax^+ + \gamma h^*(z^+)}$$

$$\frac{1}{\gamma} (z - \gamma A(x - 2x^+) - z^+) \leftarrow \cancel{\gamma h^*(z^+)}$$

$$z^+ = \text{prox}_{\gamma h^*} (z - \gamma A(x - 2x^+))$$

$$z^+ = \text{prox}_{\gamma h^*} (z - \gamma A(x - 2x^+) - \gamma b)$$

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Derived the channels Ue - Pack  
iteration (see their paper)

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Try to make things faster in how

$$r_x = b - Ax, \quad r_{x^+} = b - Ax^+$$

$$r_z = c - A^T z$$

Iteration in the code:

$$x^+ = \rho^{rx} \gamma_g (x - \alpha(r_z))$$

$$z^+ = \rho^{rz} \gamma_h (z + \alpha(zr_x - r_x))$$

$z$   $\hookrightarrow$   $v$  in code

in lecture

$g$   $\hookrightarrow$   $\underline{k}$  in code