

## The Code

The following code is the implementation for both problem 1, 2

```

"""
    The coefficient of finite difference of 2D 5 points laplacian without
    the h^2 multiplier and assume equally spaced gridpoints.
"""
function Laplacian5Stencil()
    stencil = Vector{Float64}()
    push!(stencil, (0, 0, -4))
    push!(stencil, (-1, 0, 1))
    push!(stencil, (0, -1, 1))
    push!(stencil, (1, 0, 1))
    push!(stencil, (0, 1, 1))

    return stencil end

"""
    The coefficient of finite difference of 2D 9 points laplacian without
    the h^2 multiplier and assume equally spaced gridpoints.
"""
function Laplacian9Stencil()
    stencil = Vector{Float64}()
    push!(stencil, (-1, -1, 1/6))
    push!(stencil, (-1, 1, 1/6))
    push!(stencil, (1, -1, 1/6))
    push!(stencil, (1, 1, 1/6))
    push!(stencil, (0, -1, 4/6))
    push!(stencil, (0, 1, 4/6))
    push!(stencil, (1, 0, 4/6))
    push!(stencil, (-1, 0, 4/6))
    push!(stencil, (0, 0, -20/6))

    return stencil end

"""
    Make a Linear system representin the Poisson Problem.
    Grid point:
        Equally space on both x and y direction on the domain of [0, 1] x [0, 1]
    Boundary Conditions:
        4 pices of boundary conditions are keyword paremter, should be passed in
        as a scalar function.
    Keyworld Parameters:
        * u_left, u_right, u_top, u_bottom are the function modeling the bouncary
        conditions.
        * f is the righthand side function.
        * stencil: It's a stencil generative function.
"""
function MakeLaplacianSystem(
    m::Int64;
    f=nothing,          # keyword argument
    u_left=nothing,
    u_right=nothing,
    u_top=nothing,
    u_bottom=nothing,
    stencil::Function=Laplacian5Stencil
)
    @assert m > 2 "m should be larger than 2. "
    zeroFunc(x, y) = 0
    returnsMatrixOnly = isnothing(u_top) &&
        isnothing(u_bottom) &&
        isnothing(u_left) &&
        isnothing(u_right) &&

```

```

isnothing(f)

# if every keyword argument are nothing, then we just need the
# laplacian matrix

u_left = isnothing(u_left) ? (x) -> 0 : u_left
u_right = isnothing(u_right) ? (x) -> 0 : u_right
u_top = isnothing(u_top) ? (x) -> 0 : u_top
u_bottom = isnothing(u_bottom) ? (x) -> 0 : u_bottom
f = isnothing(f) ? zeroFunc : f

Tl(i, j) = i + (j - 1)*m
# make the matrix
h = 1/(m + 1)
rhsMod = zeros(m^2)
RowIndex = Vector{Float64}()
ColIdx = Vector{Float64}()
Vals = Vector{Float64}()

function AddCoeff(TupleIdx1, TupleIdx2, coef)
    i1, j1 = TupleIdx1
    i2, j2 = TupleIdx2
    # interior
    if i2 >= 1 && i2 <= m && j2 >= 1 && j2 <= m
        push!(RowIndex, Tl(i1, j1))
        push!(ColIdx, Tl(i2, j2))
        push!(Vals, coef)
    # boundary.
    elseif i2 == 0 || i2 == m + 1 # left or right B.C
        if i2 == 0 # left
            rhsMod[Tl(i1, j1)] -= u_left(j2*h)*coef
        else # right
            rhsMod[Tl(i1, j1)] -= u_right(j2*h)*coef
        end
    elseif j2 == 0 || j2 == m + 1 # top or bottom
        if j2 == 0 # bottom
            rhsMod[Tl(i1, j1)] -= u_bottom(i2*h)*coef
        else # top
            rhsMod[Tl(i1, j1)] -= u_top(i2*h)*coef
        end
    else
        error("this should not happen, I don't expect this")
    end
end
return end

# Construct using Stencil
stencil = stencil()
for i in 1:m, j in 1:m
    for (k, l, v) in stencil
        AddCoeff((i, j), (i + k, j + l), v/h^2)
    end
    rhsMod[Tl(i, j)] += f(i*h, j*h)
end

if returnsMatrixOnly
    return sparse(RowIdx, ColIdx, Vals)
end

return sparse(RowIdx, ColIdx, Vals), rhsMod end

"""
Return the interior points of a partitioning of the grid points
in a vector of tuples of x, y if the parameter u is not given.
if it's given then it will just put these tuples of values into the
given u function, values are returned as a vector of natural ordering.

parameter:

```

```

    """
    m:: An integer larger than 2.
    """
function Grid2Vec(m::Int64, u=nothing)
    v = Vector{Float64}{}
    h = 1/(1 + m)
    if isnothing(u)
        for i in 1:m, j in 1:m
            push!(v, (i*h, j*h))
        end
        return v
    end

    if typeof(u) <: Function
        for i in 1:m, j in 1:m
            push!(v, u(i*h, j*h))
        end
        return v
    end

    if typeof(u) <: AbstractMatrix
        for i in 1:m, j in 1:m
            push!(v, u[j, i])
        end
        return v
    end
end

return v end

    """
    convert the naturally ordered vector into a grid, including the
    boundary conditions.

    * Returns a matrix that is literally the function values over the grid points.
    the (i,j) element is (j*h, i*h), so it's the 2d quadrant rotated to the right
    by 180

    * Plotting of the matrix using heatmap will recover the function's plot
    over [0, 1] x [0, 1] like, exactly as what you expected because
    julia plots it upside down.
    """

function Vec2Grid(
    m, v;
    bc_left=nothing,
    bc_right=nothing,
    bc_top=nothing,
    bc_bottom=nothing
)
    g = zeros(m, m)
    for i in 1:m, j in 1:m
        g[j, i] = v[(j - 1)*m + i]
    end
    if isnothing(bc_left) && isnothing(bc_right) && isnothing(bc_top) && isnothing(bc_bottom)
        return g
    end
    @error("Not yet implemented.")
    gWithBC = zeros(m + 2, m + 2)
    gWithBC[2:end - 1, 2:end - 1] = g

return gWithBC end

    """
    Performs conjugate gradient to solve the sparse linear effectively. The implementation of
    this method
    is faster than writting out an FFT based Poisson solver.
    """

```

```
function SimpleConjugateGradient(A::AbstractMatrix, b::AbstractVector; epsilon=1e-10)
    error("Not Yet implemented yet. ")
return end
```

```
### Hands on testing part -----
```

```
using SparseArrays, Plots, LinearAlgebra, Statistics
```

```
"""
```

```
    Basic test.
```

```
"""
```

```
function BasicTests()
    f = (x, y) -> x^2 + y^2
    m = 3
    M, b = MakeLaplacianSystem(
        3,
        u_right=(x)->1, u_top=(x) -> 1, u_bottom =(x)-> 1, u_left=(x) ->1,
        stencil=Laplacian9Stencil
    )
    @info "Carrying out some basic testing: "
    (M*6)/(m + 1)^2 |> display
    b*6/(m + 1)^2 |> display
    Vec2Grid(m, Grid2Vec(m, (x, y)-> x + y)) |> heatmap |> display

    @info "Testing whether the 9 points Laplacian works"
    m = 63
    A, b = MakeLaplacianSystem(
        m, f=f,
        u_right=(x)->1, u_top=(x) -> 1, u_bottom =(x)-> 1, u_left=(x) ->1,
        stencil=Laplacian9Stencil
    )
    println("M matrix: ")
    Vec2Grid(
        m,
        A\b,
    ) |> heatmap |> display
return end
```

```
function Problem1()
    # super fine grid point solution come first.
    M = 2^10 - 1
    f(x, y) = x^2 + y^2
    u_bc(x) = 1
    A, b = MakeLaplacianSystem(
        M,
        f=f,
        u_left=u_bc,
        u_right=u_bc,
        u_top=u_bc,
        u_bottom=u_bc
    )

    @info "The system we are solving is: "
    A |> display
    b |> display
    uVeryFine = Vec2Grid(M, A\b)
    uVeryFine |> heatmap |> display

    # estimating the error using super fine grid as a reference
    Errors = Vector{Float64}() # L2 error over the grid
    gridWidth = Vector{Float64}()
    for m in 2 .^ (collect(2:8)) .- 1
        h = 1/(m + 1)
        A, b = MakeLaplacianSystem(
```

```

        m,
        f=f,
        u_left=u_bc, u_right=u_bc, u_top=u_bc, u_bottom=u_bc
    )
    uCoarse = Vec2Grid(m, A\b)
    skip = convert(Int64, (M + 1)/(m + 1) |> floor)
    push!(Errors,
        h*norm(
            (uVeryFine[skip:skip:end, skip:skip:end] - uCoarse)[:])
        )
    )
    push!(gridWidth, h)
end
# Printing things out:
@info "These are a print out for the error vectors and grid width. "
Errors |> display
gridWidth |> display
gridWidthLogged = log2.(gridWidth); ErrorsLogged = log2.(Errors)
fig = plot(
    gridWidthLogged,
    ErrorsLogged,
    title="Log2 vs Log2 Error",
    label="5 p stencil",
    legend=:bottomright
)
plot!(fig, gridWidthLogged, gridWidth.^2 .|> log2, label="reference h^2")
xlabel!(fig, "log2(h)")
ylabel!(fig, "log2(E)")
fig |> display
savefig(fig, "p1_fig.png")

# Estimating rate of convergence.
println("The estimated rate of convergence for the first problem is: ")
(ErrorsLogged[1:end - 1] - ErrorsLogged[2:end])./
(gridWidthLogged[1:end - 1] - gridWidthLogged[2:end])|> mean |> display

return end

function Problem2()
    function SolveWithDeferredCorrection(m::Int64)
        # super fine grid point solution come first.
        f(x, y) = x^2 + y^2
        u_bc(x) = 1
        A, b = MakeLaplacianSystem(
            m,
            f=f,
            u_left=u_bc,
            u_right=u_bc, u_top=u_bc, u_bottom=u_bc,
            stencil=Laplacian9Stencil
        )
        b .+= 1/(3*(m + 1)^2)
    end
    return Vec2Grid(m, A\b)
end
uVeryFine = SolveWithDeferredCorrection(2^10 - 1)
@info "uVeryFine 9 points stencils solution looks like: "
uVeryFine |> display

# estimating the error using super fine grid as a reference
Errors = Vector{Float64}() # L2 error over the grid
gridWidth = Vector{Float64}()
M = 2^10 - 1
for m in 2 .^ (collect(2:8)) .- 1
    h = 1/(m + 1)
    uCoarse = SolveWithDeferredCorrection(m)
    skip = convert(Int64, (M + 1)/(m + 1) |> floor)
    push!(Errors,
        h*norm(
            (uVeryFine[skip:skip:end, skip:skip:end] - uCoarse)[:])
    )
end

```

```

    )
    push!(gridWidth, h)
end

# Printing things out:
@info "These are a print out for the error vectors and grid width. "
Errors |> display
gridWidth |> display
gridWidthLogged = gridWidth .|> log2
ErrorsLogged = Errors .|> log2
fig = plot(
    gridWidthLogged,
    ErrorsLogged,
    title="Log2 vs Log2 Error",
    label="9 p stencil",
    legend=:bottomright
)
plot!(
    fig,
    gridWidthLogged,
    gridWidth.^3 .|> log2,
    label="reference h^3"
)
xlabel!(fig, "log2(h)")
ylabel!(fig, "log2(E)")
fig |> display
savefig(fig, "p2_fig.png")
# estimate the convergence numerically.
println("The estimated rate of convergence for the second problem is: ")
(ErrosLogged[1:end - 1] - ErrorsLogged[2:end])./(
    gridWidthLogged[1:end - 1] - gridWidthLogged[2:end])|> mean |> display
end

BasicTests()
Problem1()
Problem2()

```

## Problem 1

### Theory that goes with the code

Natural ordering of the element refers to the mapping of tuple of indices from the grid to the actual point in space, given as  $u_{i,j} = u(ih, jh)$ . Here, we use equally space point in both xy direction on the unit square:  $[0, 1] \times [0, 1]$ . And we use  $m$  number of interior points, and including the boundary we will have  $m + 2$  boundary point.

The implementation focuses on building up the system of matrices using the stencil defined on each point indexed grid point  $(i, j)$ . For each point, we convert the coordinate indices  $(i, j)$  to linear indices for row and column location in the sparse matrix, the coefficient on the stencil is then added to the matrix or the RHS of the system. This is done for  $(i, j) \in \{0, 1, \dots, m\}^2$ . If any of the point are laying outside, we move the coefficient to the RHS vector.

Alternatively, one can also use Kronecker Product to construct the structural matrice, but it will lose the elegence of on the code and less generic. Algebraically, one can handle the boundary conditions by thinking about the stencils on the boundary. I only did it for the 5 point stencil.

Let  $\mathcal{I}(i, j) = (j - 1)m + i \quad \forall (i, j) \in [1, \dots, m] \times [1, \dots, m]$  be a mapping from the coordinate indices of the grid point to linear index. Next, we consider partitioning the vector  $u$  and  $b$  the RHS vector by rows on

the grid using natural ordering.

$$\begin{aligned}
u &= \begin{bmatrix} u^{[1]} \\ u^{[2]} \\ \vdots \\ u^{[m]} \end{bmatrix} \quad \vec{b} = \begin{bmatrix} b^{[1]} \\ b^{[2]} \\ \vdots \\ b^{[m]} \end{bmatrix} \quad \vec{f} = \begin{bmatrix} f^{[1]} \\ f^{[2]} \\ \vdots \\ f^{[m]} \end{bmatrix} \\
u^{[i]} &= \begin{bmatrix} u_{1,i} \\ u_{2,i} \\ \vdots \\ u_{m,i} \end{bmatrix}
\end{aligned} \tag{1.1}$$

For the boundary conditions, we consider them row by row. Consider the first row:

$$\begin{aligned}
\vec{f}_{\mathcal{I}(i,1)} &= \frac{1}{h^2} \text{sum} \left( \begin{bmatrix} & u_{i,2} & \\ u_{i-1,1} & -4u_{i,1} & u_{i+1,1} \\ & u_{i,0} & \end{bmatrix} \right) \quad \forall 1 \leq i \leq m \\
\vec{b}^{[1]} &= -\frac{1}{h^2} \begin{bmatrix} u_{1,0} \\ u_{2,0} \\ \vdots \\ u_{m,0} \end{bmatrix} - \frac{u_{0,1} \mathbf{e}_1^{(m)}}{h^2} - \frac{u_{m+1,1} \mathbf{e}_m^{(m)}}{h^2} + f^{[1]}
\end{aligned} \tag{1.2}$$

Next, we consider all the rows in the middle rows:

$$\begin{aligned}
\vec{f}_{\mathcal{I}(i,j)} &= \frac{1}{h^2} \text{sum} \left( \begin{bmatrix} & u_{i,j+1} & \\ u_{i-1,j} & -4u_{i,j} & u_{i+1,j} \\ & u_{i,j-1} & \end{bmatrix} \right) \quad \forall 1 \leq i \leq m, 2 \leq j \leq m-1 \\
\vec{b}^{[j]} &= -\frac{u_{0,j} \mathbf{e}_1^{(m)}}{h^2} - \frac{u_{m+1,j} \mathbf{e}_m^{(m)}}{h^2} + f^{[j]}
\end{aligned} \tag{1.3}$$

Finally we consider the last row on the grid point:

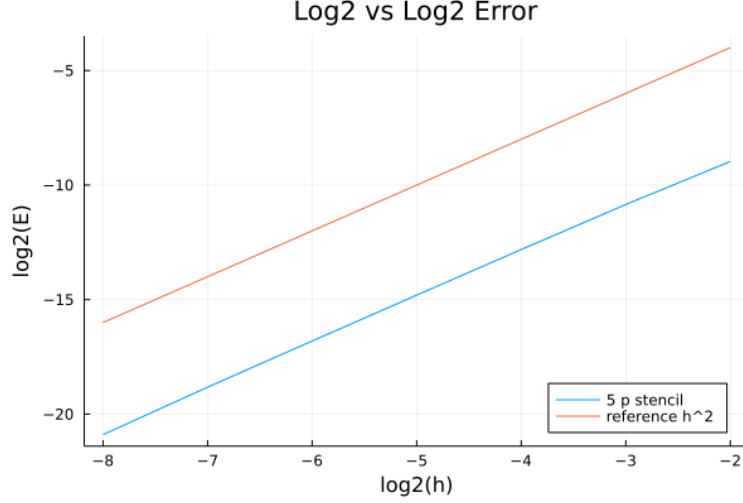
$$\begin{aligned}
\vec{f}_{\mathcal{I}(i,m)} &= \frac{1}{h^2} \text{sum} \left( \begin{bmatrix} & u_{i,m+1} & \\ u_{i-1,m} & -4u_{i,m} & u_{i+1,m} \\ & u_{i,m-1} & \end{bmatrix} \right) \quad \forall 1 \leq i \leq m \\
\vec{b}^{[m]} &= -\frac{1}{h^2} \begin{bmatrix} u_{1,m+1} \\ u_{2,m+1} \\ \vdots \\ u_{m,m+1} \end{bmatrix} - \frac{u_{0,m} \mathbf{e}_1^{(m)}}{h^2} - \frac{u_{m+1,m} \mathbf{e}_m^{(m)}}{h^2} + f^{[m]}
\end{aligned} \tag{1.4}$$

This completes the analysis for the RHS vector that includes the boundary. Take notice that a very different approach has been coded. This part is just for analysis.

## $\mathcal{O}(h^2)$ Convergence Rate

Error computations is made by comparing solution on sparser grid point with the solution obtained on a fine grid point. I don't happen to know the solution of the PDE besides the solution expressed using the Green's function with the given density function.

For the implementation of this part, refers to the "Problem1()" function. The finer grid is taken with  $m = 2^{10} - 1$ , and then  $m$  follows the pattern of  $2^k - 1$  so that the discretized grid-point of the sparser grids is a subset of the finer grid. This is a plot of the error measured by  $L2$  norm in 2D in log log plot, referenced with  $h^2$  convergence:

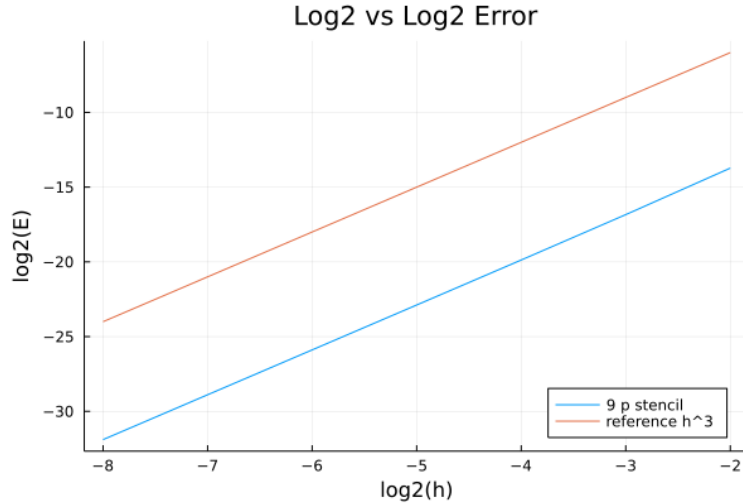


Numerically, I take the slope of the adjacent points on the log log plot and take the averaging, producing an estimate of the rate of convergence for the 5 points stencil method. The estimated rate is: 1.986773401321141.

## Problem 2

The same routine is run for the 9-point stencils but with a different definition of the stencil. Deferred correction is implemented by the end, right before the calculation of solution.

The estimated rate of convergence of error is 3.0268502045280115, The error is  $\mathcal{O}(n^3)$ , which differs from the theory. Here is a plot of the error under log log plot:



Observe that this is a mismatch from what is predicted by the theory where it's stated that with deferred correction, the error should have convergence rate of  $\mathcal{O}(h^4)$ . The cause of the problem boundary conditions. More specifically, the  $\nabla^4[f(x, y)]$  at the boundary point is not defined. Numerically, we have:

$$\begin{aligned}\nabla^4[f(x, y)] &= (\partial_x^4 + \partial_y^4 + 2\partial_x^2\partial_y^2)[x^2 + y^2] \\ &= 2 + 2 + 0 = 4\end{aligned}\tag{2.1}$$

However, at the boundary point if we assume that the function is  $\equiv 1$ , then  $\nabla^4[u(x, y)]_{(x, y)=\text{bd}([0, 1]^2)} \equiv 0$ , making the Laplacian of the RHS function discontinuous at the boundary of the grid, making  $\nabla^4[f(x, y)] \neq \nabla^4[u(x, y)]$  at the boundary, failing the assumptions made by the Deferred Correction Method, hence the accuracy of  $\mathcal{O}(h^4)$  is not guaranteed to exist for this particular case.

I am not sure why it's  $\mathcal{O}(h^3)$  specifically instead of  $\mathcal{O}(h^2)$  or anything in between.



### Problem 3

Our goal here is to investigate the format of the System of Linear Equations ( $A\vec{c} = \vec{b}$ ) needed for inhomogenous Poisson Equation on a 2D unit square  $\Omega = [0, 1] \times [0, 1]$  with boundary conditions  $g(x, y)$  at the boundary. For simplicity of the discussion, we assume the use of gridpoint (Not necessarily equally spaced) and piecewise bilinear approximation for the problem in  $\Omega$ . This is the Boundary Value Problem:

$$\begin{cases} \nabla^2 \cdot u = f & \forall x \in \Omega \\ u = g & \forall x \in \text{bd}(\Omega) \end{cases} \quad (3.1)$$

Here we use  $\text{bd}(\Omega)$  to denote the boundary of the Domain.

Next we consider the weak form and using  $S$  as the basis function set:

$$\begin{aligned} \forall \hat{v} \in S : \langle \mathcal{L}[\hat{u}], \hat{v} \rangle &= \langle f, \hat{v} \rangle \\ \forall \hat{v} \in \mathcal{S} : \langle \mathcal{L}[\hat{u}], \hat{v} \rangle &= \iint_{\Omega} (\hat{u}_{xx} + \hat{u}_{yy}) \hat{v} dx dy \\ &= - \iint_{\Omega} (\hat{u}_x \hat{v}_x + \hat{u}_y \hat{v}_y) dx dy + \int_{\text{bd}(\Omega)} (\nabla \hat{u}) \hat{v} d\hat{\mathbf{n}} \\ &= \int_{\text{bd}(\Omega)} (\nabla u) \hat{v} d\hat{\mathbf{n}} - \iint_{\Omega} (\nabla \hat{u}) \cdot (\nabla \hat{v}) dx dy \end{aligned} \quad (3.2)$$

When the set  $S$  is finite, we can create a system of equations to assert the conditions presented by the weak form formulation. The definition of inner product is implicitly defined here.

We assume that the grid partitioning  $\Omega$  is in the size of  $(n_x + 2) \times (n_y + 2)$ , including the grid point.  $n_x, n_y$  each denotes the number interior discretized grid points on the x, y direction. We use the following notations for mapping coordinate indices on  $\Omega$  to linear indices:

$$\lfloor i, j \rfloor = i + 1 + j(n_y + 2)$$

Please observe that  $0 \leq i \leq n_x, 0 \leq j \leq n_y$ . Now, we are ready to define some set of indices to simplify things:

$$\begin{aligned} \overline{G} &:= \{0, 1, \dots, n_x + 1\} \times \{0, 1, \dots, n_y + 1\} \\ G &:= \{1, \dots, n_x\} \times \{1, \dots, n_y\} \\ \overline{\mathbb{B}} &:= \overline{G} \setminus G \\ \mathcal{N}(i, j) &:= \{(k, l) \in \overline{G} : |i - k| + |j - l| \leq 2\} \end{aligned} \quad (3.3)$$

The set  $\overline{G}$  denotes all the nodes including the nodes at the boundary. And the set  $G$  denotes the nodes that are only inside of the grid, and the set  $\overline{\mathbb{B}}$  denotes that set of all the points at the boundary. Finally, the function  $\mathcal{N}(i, j)$  denotes that set of indices for the nodes that are neighbouring the node at  $(i, j)$ , including the node that is at  $(i, j)$ .

In our case, we present the use of Bi-linear Functions on the grid point. Which is parameterized by 5 points and separated into 5 cases:

$$\varphi_{\lfloor i, j \rfloor}(x, y) = \begin{cases} \frac{(x - x_{i-1})(y - y_{j-1})}{(x_i - x_{i-1})(y_j - y_{j-1})} & \text{in } [x_{i-1}, x_i] \times [y_{j-1}, y_j] \\ \frac{(x - x_{i+1})(y - y_{j-1})}{(x_i - x_{i+1})(y_j - y_{j-1})} & \text{in } [x_i, x_{i+1}] \times [y_{j-1}, y_j] \\ \frac{(x - x_{i-1})(y - y_{j+1})}{(x_i - x_{i-1})(y_j - y_{j+1})} & \text{in } [x_{i-1}, x_i] \times [y_j, y_{j+1}] \\ \frac{(x - x_{i+1})(y - y_{j+1})}{(x_i - x_{i+1})(y_j - y_{j+1})} & \text{in } [x_i, x_{i+1}] \times [y_j, y_{j+1}] \\ 0 & \text{elsewhere} \end{cases} \quad (3.4)$$

Take notice that, when the function  $\varphi$  is at the boundary, for example  $\varphi_{\lfloor 0, 0 \rfloor}$ , then it will only have 2 cases instead of 5, because only the top left corner of the node is defined inside of the grid. By a similar line of

reasoning, one can define the bi-linear function for all node with indices in  $\overline{\mathbb{B}}$ . Now, we are ready to define that the set of basis functions we use is:  $S := \{\varphi_{[i,j]} : (i,j) \in \overline{G}\}$ .

To make use of the weak Form of Finite Element, we represent the approximated solution  $\tilde{u}$  using the set of basis function from  $S$ :

$$\tilde{u}(x, y) = \sum_{(i,j) \in \overline{G}} c_{[i,j]} \varphi_{[i,j]}(x, y) \quad (3.5)$$

Before we start, Observe that all the basis functions at the boundary  $\text{bd}(\Omega)$  must satisfies the boundary conditions, which immediately gave us  $4 + 2(n_x + n_y)$  constraints to work with:

$$\begin{aligned} c_{[i,j]} &= g(x_i, y_j) \quad \forall (i, j) \in \overline{\mathbb{B}} \\ \implies A_{[i,j],[i,j]} &= 1 \quad \forall (i, j) \in \overline{\mathbb{B}} \\ \vec{b}_{[i,j]} &= g(x_i, y_j) \quad \forall (i, j) \in \overline{\mathbb{B}} \end{aligned} \quad (3.6)$$

Reconsider the weakform formulations with the given form of approximated solution and the basis set  $S$ , this will give use constraint for each of the node with indices in  $G$ , a total of  $n_x n_y$  constraints for the system.

$$\begin{aligned} \langle \mathcal{L}\tilde{u}, \varphi_{[i,j]} \rangle &= \langle f, \varphi_{[i,j]} \rangle \quad \forall (i, j) \in G \\ \left\langle \mathcal{L} \left[ \sum_{(k,l) \in G} c_{[k,l]} \varphi_{[k,l]} \right], \varphi_{[i,j]} \right\rangle &= \langle f, \varphi_{[i,j]} \rangle \\ \left\langle \mathcal{L} \left[ \sum_{(k,l) \in G} c_{[k,l]} \varphi_{[k,l]} \right], \varphi_{[i,j]} \right\rangle &= \langle f, \varphi_{[i,j]} \rangle \\ \forall \\ \sum_{(k,l) \in \mathcal{N}(i,j)} c_{[k,l]} \langle \mathcal{L}[\varphi_{[k,l]}], \varphi_{[i,j]} \rangle &= \langle f, \varphi_{[i,j]} \rangle \end{aligned} \quad (3.7)$$

For each of the node with indices in the interior, denoted by  $G$ , it will interact with all its 8 neighbours, including itself via the dot product defined, and that leads from the second last line to the last line. This gives at most 9 non-zeros element for  $n_x n_y$  number of row of the matrix  $A$ . At this point, we are finally able to present the  $(n_x + 2) \times (n_y + 2)$  linear system in a visually pleasant way:

$$\begin{aligned} \forall (i, j), (k, l) \in \overline{G} : \\ A_{[i,j],[k,l]} &= \begin{cases} \langle \mathcal{L}[\varphi_{[k,l]}], \varphi_{[i,j]} \rangle & (k, l) \in \mathcal{N}(i, j) \\ 1 & (k, l), (i, j) \in \overline{\mathbb{B}} \\ 0 & \text{else} \end{cases} \\ \vec{b}_{[i,j]} &= \begin{cases} g(x_i, y_j) & (i, j) \in \overline{\mathbb{B}} \\ \langle f, \varphi_{[i,j]} \rangle & \text{else} \end{cases} \end{aligned} \quad (3.8)$$

This way of capturing all  $(n_x + 2)(n_y + 2)$  number of variables simplifies the right handside of the equation comparing to setting up the  $4 + 2(n_x + n_y)$  number of variables for  $c_{[i,j]}$  at the boundary  $\overline{\mathbb{B}}$  and then create a system of the size  $n_x n_y$  for all the node in  $G$ . What remains to be evaluated is the inner product between basis function at neighbouring nodes and itself. And recall from (3.2), when the one of the nodes at  $(i, j)$  or  $(k, l)$  is in  $\overline{\mathbb{B}}$ , we will have to evaluate an extra integral over the boundary  $\text{bd}(\Omega)$ .