

# Error Bound Can Give Near Optimal Convergence Rate for Inexact Accelerated Proximal Gradient Method

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## Abstract

This is still a draft. [5].

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## 1 Introduction

**Notations.** Let  $g : \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$ , we denote  $g^*$  to be the Fenchel conjugate.  $I : \mathbb{R}^n \rightarrow \mathbb{R}^n$  denotes the identity operator. For a multivalued mapping  $T : \mathbb{R}^n \rightarrow 2^{\mathbb{R}^n}$ ,  $\text{gra } T$  denotes the graph of the operator, defined as  $\{(x, y) \in \mathbb{R}^n \times \mathbb{R}^n : y \in Tx\}$ .

### 1.1 Epsilon subgradient and inexact proximal point

{def:esp-subgrad}

**Definition 1.1 ( $\epsilon$ -subgradient)** *Let  $g : \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$  be proper, lsc. Let  $\epsilon \geq 0$ . Then the  $\epsilon$ -subgradient of  $g$  at some  $\bar{x} \in \text{dom } g$  is given by:*

$$\partial g_\epsilon(\bar{x}) := \{v \in \mathbb{R}^n \mid \langle v, x - \bar{x} \rangle \leq g(x) - g(\bar{x}) + \epsilon \forall x \in \mathbb{R}^n\}.$$

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When  $\bar{x} \notin \text{dom } g$ , it has  $\partial g_\epsilon(\bar{x}) = \emptyset$ .

**Remark 1.2**  $\partial_\epsilon g$  is a multivalued operator and, it's not monotone, unless  $\epsilon = 0$ , which makes it equivalent to Fenchel subgradient  $\partial g$ .

If we assume lsc, proper and convex  $g$ , we will now introduce results in the literatures that we will use.

**Fact 1.3 ( $\epsilon$ -Fenchel inequality)** *Let  $\epsilon \geq 0$ , then:*

$$x^* \in \partial_\epsilon f(\bar{x}) \iff f^*(x^*) + f(\bar{x}) \leq \langle x^*, \bar{x} \rangle + \epsilon \implies \bar{x} \in \partial_\epsilon f^*(x^*).$$

*They are all equivalent if  $f^{**}(\bar{x}) = f(\bar{x})$ .*

**Remark 1.4** The above fact is taken from Zalinascu [4, Theorem 2.4.2].

We will now define inexact proximal point based on  $\epsilon$ -subgradient

**Definition 1.5 (inexact proximal point)** *For all  $x \in \mathbb{R}^n, \epsilon \geq 0, \lambda > 0$ ,  $\tilde{x}$  is an inexact evaluation of proximal point at  $x$ , if and only if it satisfies:*

$$\lambda^{-1}(x - \tilde{x}) \in \partial_\epsilon g(\tilde{x}).$$

*We denote it by  $\tilde{x} \approx_\epsilon \text{prox}_{\lambda g}(x)$ .*

**Remark 1.6** This definition is nothing new, for example see Villa et al. [3, Definition 2.1]

**Fact 1.7 (the resolvent identity)** *Let  $T : \mathbb{R}^n \rightarrow 2^{\mathbb{R}^n}$ , then it has:*

$$(I + T)^{-1} = (I - (I + T^{-1})^{-1}).$$

**Theorem 1.8 (inexact Moreau decomposition)** *Let  $g : \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$  be a closed, convex and proper function. It has the equivalence*

$$\tilde{y} \approx_\epsilon \text{prox}_{\lambda^{-1}g^*}(\lambda^{-1}y) \iff y - \lambda\tilde{y} \approx_\epsilon \text{prox}_{\lambda g}(y).$$

*Proof.* Consider  $\tilde{y} \approx_\epsilon \text{prox}_{\lambda^{-1}g^\star}(\lambda^{-1}y)$ , then it has:

$$\begin{aligned}
& \tilde{y} \in (I + \lambda^{-1}\partial_\epsilon g^\star)^{-1}(\lambda^{-1}y) \\
& \iff (\lambda^{-1}y, \tilde{y}) \in \text{gra}(I + \lambda^{-1}\partial_\epsilon g^\star)^{-1} \\
& \stackrel{(1)}{\iff} (\lambda^{-1}y, \tilde{y}) \in \text{gra}(I - (I + \partial_\epsilon g \circ (\lambda I))^{-1}) \\
& \iff (\lambda^{-1}y, \lambda^{-1}y - \tilde{y}) \in \text{gra}(I + \partial_\epsilon g \circ (\lambda I))^{-1} \\
& \iff (\lambda^{-1}y - \tilde{y}, \lambda^{-1}y) \in \text{gra}(I + \partial_\epsilon g \circ (\lambda I)) \\
& \iff (y - \lambda\tilde{y}, \lambda^{-1}y) \in \text{gra}(\lambda^{-1}I + \partial_\epsilon g) \\
& \iff (y - \lambda\tilde{y}, y) \in \text{gra}(I + \lambda\partial_\epsilon g) \\
& \iff y - \lambda\tilde{y} \in (I + \lambda\partial_\epsilon g)^{-1}y \\
& \iff y - \lambda\tilde{y} \approx_\epsilon \text{prox}_{\lambda g}(y).
\end{aligned}$$

At (1) we can use Fact 1.7, and it has  $(\lambda^{-1}\partial_\epsilon g^\star)^{-1} = \partial_\epsilon g \circ (\lambda I)$  by Fact 1.3 and the assumption that  $g$  is closed, convex and proper. ■

## 1.2 Inexact proximal gradient inequality

{ass:for-inxt-pg-ineq}

**Assumption 1.9 (for inexact proximal gradient)** The assumption is about  $(F, f, g, L)$ . We assume that

- (i)  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  is a convex,  $L$  Lipschitz smooth function (i.e:  $\nabla f$  is  $L$  a Lipschitz continuous mapping).
- (ii)  $g : \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$  is a convex, proper, and lsc function which we do not have its exact proximal operator.
- (iii) The over all objective is  $F = f + g$ .

No, we develop the theory based on the use of epsilon subgradient as in Definition 1.1. Let  $\rho > 0$ , the exact proximal gradient operator defined for  $(f, g, L)$  satisfying Assumption 1.9 has

$$\begin{aligned}
T_\rho(x) &= \underset{z \in \mathbb{R}^n}{\text{argmin}} \left\{ g(z) + \langle \nabla f(x), z \rangle + \frac{\rho}{2} \|z - x\|^2 \right\} \\
&= \text{prox}_{\rho^{-1}g} \left( x - \rho^{-1} \nabla f(x) \right).
\end{aligned}$$

The following definition extends the proximal gradient operator to the inexact case using the concept of  $\epsilon$ -subgradient as given by Definition 1.1.

{def:inxt-pg}

**Definition 1.10 (inexact proximal gradient)** Let  $(f, g, L)$  satisfies Assumption 1.9. Let  $\epsilon \geq 0, \rho > 0$ . Then,  $\tilde{x} \approx_\epsilon T_\rho(x)$  is an inexact proximal gradient if it satisfies variational

inequality:

$$\mathbf{0} \in \nabla f(x) + \rho(x - \tilde{x}) + \partial_\epsilon g(\tilde{x}).$$

**Remark 1.11** We assumed that we can get exact evaluation of  $\nabla f$  at any points  $x \in \mathbb{R}^n$ .

{lemma:other-repr-inxt-pg}

**Lemma 1.12 (other representations of inexact proximal gradient)**

Let  $(f, g, L)$  satisfies Assumption 1.9,  $\epsilon \geq 0, \rho > 0$ , then for all  $\tilde{x} \approx_\epsilon T_\rho(x)$ , it has the following equivalent representations:

$$\begin{aligned} (x - \rho^{-1}\nabla f(x)) - \tilde{x} &\in \rho^{-1}\partial_\epsilon g(\tilde{x}) \\ \iff \tilde{x} &\in (I + \rho^{-1}\partial_\epsilon g(\tilde{x}))^{-1}(x - \rho^{-1}\nabla f(x)) \\ \iff x &\approx_\epsilon \text{prox}_{\rho^{-1}g} (x - \rho^{-1}\nabla f(x)) \end{aligned}$$

*Proof.* It's direct. ■

{thm:inxt-pg-ineq}

**Theorem 1.13 (inexact over-regularized proximal gradient inequality)**

Let  $(F, f, g, L)$  satisfies Assumption 1.9,  $\epsilon \geq 0, B \geq 0, \rho > 0$ . Consider  $\tilde{x} \approx_\epsilon T_{B+\rho}(x)$ . Denote  $F = f + g$ . If in addition,  $\tilde{x}, B$  satisfies the line search condition  $D_f(\tilde{x}, x) \leq B/2\|x - \tilde{x}\|^2$ , then it has  $\forall z \in \mathbb{R}^n$ :

$$-\epsilon \leq F(z) - F(\tilde{x}) + \frac{B+\rho}{2}\|x - z\|^2 - \frac{B+\rho}{2}\|z - \tilde{x}\|^2 - \frac{\rho}{2}\|\tilde{x} - x\|^2.$$

*Proof.* By Definition 1.10 write the variational inequality that describes  $\tilde{x} \approx_\epsilon T_B(x)$ , and the definition of epsilon subgradient (Definition 1.1) it has for all  $z \in \mathbb{R}^n$ :

$$\begin{aligned} -\epsilon &\leq g(z) - g(\tilde{x}) - \langle (B + \rho)(\tilde{x} - x) - \nabla f(x), z - \tilde{x} \rangle \\ &= g(z) - g(\tilde{x}) - (B + \rho)\langle \tilde{x} - x, z - \tilde{x} \rangle + \langle \nabla f(x), z - \tilde{x} \rangle \\ &\stackrel{(1)}{\leq} g(z) + f(z) - g(\tilde{x}) - f(\tilde{x}) - (B + \rho)\langle \tilde{x} - x, z - \tilde{x} \rangle - D_f(z, x) + D_f(\tilde{x}, x) \\ &\stackrel{(2)}{\leq} F(z) - F(\tilde{x}) - (B + \rho)\langle \tilde{x} - x, z - \tilde{x} \rangle + \frac{B}{2}\|\tilde{x} - x\|^2 \\ &= F(z) - F(\tilde{x}) + \frac{B+\rho}{2}(\|x - z\|^2 - \|\tilde{x} - x\|^2 - \|z - \tilde{x}\|^2) + \frac{B}{2}\|\tilde{x} - x\|^2 \\ &= F(z) - F(\tilde{x}) + \frac{B+\rho}{2}\|x - z\|^2 - \frac{B+\rho}{2}\|z - \tilde{x}\|^2 - \frac{\rho}{2}\|\tilde{x} - x\|^2. \end{aligned}$$

At (1), we used considered the following:

$$\begin{aligned} \langle \nabla f(x), z - x \rangle &= \langle \nabla f(x), z - x + x - \tilde{x} \rangle \\ &= \langle \nabla f(x), z - x \rangle + \langle \nabla f(x), x - \tilde{x} \rangle \\ &= -D_f(z, x) + f(z) - f(x) + D_f(\tilde{x}, x) - f(\tilde{x}) + f(x) \\ &= -D_f(z, x) + f(z) + D_f(\tilde{x}, x) - f(\tilde{x}). \end{aligned}$$

At (2), we used the fact that  $f$  is convex hence  $-D_f(z, x) \leq 0$  always, and in the statement hypothesis we assumed that  $B$  has  $D_f(\tilde{x}, x) \leq B/2\|\tilde{x} - x\|^2$ . We also used  $F = f + g$ . ■

**Remark 1.14** When  $\epsilon = 0, \rho = 0$ , this reduces to proximal gradient inequality in the exact case. In this inequality, observe that the parameter  $\epsilon$  controls the inexactness of the proximal gradient evaluation. More specifically,  $\epsilon_k$  controls the absolute perturbations of the proximal gradient inequality compared to its exact counterpart.  $\rho$  on the other hand, it is the over-relaxation of proximal gradient operator, and it compensates the perturbations caused by  $\approx_\epsilon$  relative to the term  $\|\tilde{x} - x\|^2$ .

### 1.3 Optimizing the inexact proximal point problem

{sec:optz-inxt-pp-problem}

In this section we will present the optimization problem that obtains a  $\tilde{x}$  such that  $\tilde{x} \approx_\epsilon \text{prox}_{\lambda g}(z)$ . Eventually we want to evaluate  $T_\rho(x)$  of some  $F = f + g$  inexactly using Lemma 1.12. To do that one would need to evaluate  $\text{prox}_{\rho^{-1}g}$  inexactly which is defined in Definition 1.5.

Most of these results that will follow are from the literature. To start, we must assume the following about a function  $g : \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$ , with  $g$  closed, convex and proper.

{ass:for-inxt-prox} **Assumption 1.15 (linear composite of convex nonsmooth function)**

This assumption is about  $(g, \omega, A)$ . Let  $m \in \mathbb{N}, n \in \mathbb{R}^n$ , we assume that

- (i)  $A \in \mathbb{R}^{m \times n}$  is a matrix.
- (ii)  $\omega : \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$  is a closed and convex function such that it admits proximal operator  $\text{prox}_{\lambda\omega}$  and, its conjugate  $\omega^*$  is known.
- (iii)  $g := \omega(Ax)$  such that  $\text{rng } A \cap \text{ri dom } g \neq \emptyset$ .

Now, we are ready to discuss how to choose  $\tilde{x} \approx_\epsilon \text{prox}_{\lambda g}(x)$ . Fix  $y \in \mathbb{R}^n, \lambda > 0$ , we are ultimately interested in minimizing:

$$\{\text{eqn:primal-pp}\} \quad \Phi_\lambda(u) := \omega(Au) + \frac{1}{2\lambda}\|u - y\|^2 \quad (1.1)$$

This problem admits dual objective in  $\mathbb{R}^m$ :

$$\{\text{eqn:dual-pp}\} \quad \Psi_\lambda(v) := \frac{1}{2\lambda}\|\lambda A^\top v - y\|^2 + \omega^*(v) - \frac{1}{2\lambda}\|y\|^2. \quad (1.2)$$

We define the duality gap

$$\{\text{eqn:duality-gap-pp}\} \quad \mathbf{G}_\lambda(u, v) := \Phi_\lambda(u) + \Psi_\lambda(v). \quad (1.3)$$

If strong duality holds, it exists  $(\hat{u}, \hat{v})$  such that we have the following:

$$\mathbf{G}_\lambda(\hat{u}, \hat{v}) = 0 = \min_u \Phi_\lambda(u) + \min_v \Psi_\lambda(v)$$

{thm:primal-dual-trans} The following theorem quantifies a sufficient conditions for  $\tilde{x} \approx_\epsilon \text{prox}_{\lambda g}(x)$ . The theorem below is from [3, Proposition 2.2].

**Theorem 1.16 (primal translate to dual [3, Proposition 2.2])** *Let  $(g, \omega, A)$  satisfies assumption 1.15,  $\epsilon \geq 0$ , then*

$$(\forall z \approx_\epsilon \text{prox}_{\lambda g}(y)) (\exists v \in \text{dom } \omega^\star) : z = y - \lambda A^\top v.$$

This theorem that follows is from Villa et al. [3, Proposition 2.3].

{thm:dltty-gap-inxt-pp} **Theorem 1.17 (duality gap of inexact proximal problem [3, Proposition 2.3])** *Let  $(g, \omega, A)$  satisfies Assumption 1.15, for all  $\epsilon \geq 0$ ,  $v \in \mathbb{R}^n$  consider the following conditions:*

- (i)  $\mathbf{G}_\lambda(y - \lambda A^\top v, v) \leq \epsilon$ .
- (ii)  $A^\top v \approx_\epsilon \text{prox}_{\lambda^{-1}g^\star}(\lambda^{-1}y)$ .
- (iii)  $y - \lambda A^\top v \approx_\epsilon \text{prox}_{\lambda g}(y)$ .

*They have (a)  $\implies$  (b)  $\iff$  (c). If in addition  $\omega^\star(v) = g^\star(A^\top v)$ , then all three conditions are equivalent.*

*Proof.* The proof of (a)  $\implies$  (b), and the case when (a)  $\iff$  (b), we refer readers to Villa et al. [3, Proposition 2.3], and to show (b)  $\iff$  (c) use Theorem 1.8. ■

{fact:minimizing-dual-pp} The following fact from the literature indicates that it's sufficient to minimize the dual problem  $\Psi_\lambda$  to obtain an element of the inexact proximal point operator. The following fact is Proposition is from Villa et al. [3, Theorem 5.1].

**Fact 1.18 (minimizing dual of the proximal problem [3, Theorem 5.1])** *Let  $\bar{v}$  be a solution of  $\Psi_\lambda$ . Suppose that  $(v_n)_{n \geq 0}$  is a minimizing sequence for  $\Psi_\lambda$ . Let  $z_n = y - \lambda A^\top v_n$ , and  $\bar{z} = y - \lambda A^\top \bar{v}$ . Then the following are true:*

- (i) *The sequence  $z_n \rightarrow \bar{z}$  as well.*
- (ii) *There exists a constant  $K$  that depends on the sequence  $z_n$ , and  $\Phi_\lambda$  such that for all  $k \geq 0$  the inequality:*

$$\Phi_\lambda(z_n) - \Phi_\lambda(\bar{z}) \leq K \|z_n - \bar{z}\| \leq K \sqrt{2\lambda} (\Psi_\lambda(v_n) - \Psi_\lambda(\bar{v}))^{1/2}.$$

We remark that the above fact translates any algorithm that optimizes the dual problem  $\Psi_\lambda$  into optimizing the duality gap  $\mathbf{G}(z_n, v_n)$  because it shows the primal is bounded by the dual through a Lipschitz-like inequality.

With the proof back in Villa et al. [3, Theorem 5.1], pay close attention to the constant  $L_\Phi$  exists by the virtue of  $\Psi$  being convex, and the sequence  $v_n \rightarrow \bar{v}$  which allows the sequence  $z_n \rightarrow \bar{z}_n$  to be bounded inside the domain of  $\Phi_\lambda$ . By convexity,  $\Phi_\lambda$  is locally Lipschitz on  $\text{ri dom } \Phi_\lambda$ , by convergence of  $z_n$  the sequence is in a compact set in  $\text{ri dom } \Phi_\lambda$  hence a Lipschitz bound with constant  $K$  exists.

For this reason, the number of iterations of the inner loop required to achieve  $\mathbf{G}(z_n, v_n) < \epsilon$  for a given  $\epsilon$  is related to the convergence rate of the algorithms used to optimize  $\Psi_\lambda(v_n)$ . With the theorem derived above, and using Theorem 1.17 it implies that any algorithm which can optimize function value  $\Psi_\lambda$  will produce iterates sufficient to achieve  $\approx_\epsilon \text{prox}_{\lambda g}(y)$ .

## 1.4 Literature reviews

## 1.5 Our contributions

# 2 The inexact accelerated proximal gradient with controlled errors

In this section, we present an accelerated algorithm with controlled error using Definition 1.10, and show that it can have a convergence rate under certain error conditions.

{def:inxt-apg}

### Definition 2.1 (our inexact accelerated proximal gradient)

Suppose that  $(F, f, g, L)$  and, sequences  $(\alpha_k, B_k, \rho_k, \epsilon_k)_{k \geq 0}$  satisfies the following

- (i)  $(\alpha_k)_{k \geq 0}$  is a sequence such that  $\alpha \in (0, 1]$  for all  $k \geq 0$ .
- (ii)  $(B_k)_{k \geq 0}$  has  $B_k > 0 \forall k$ , it characterizes any potential line search, back tracking routine.
- (iii)  $(\rho_k)_{k \geq 0}$  be a sequence such that  $\rho_k \geq 0$ , characterizing the over-relaxation of the proximal gradient operator.
- (iv)  $(\epsilon_k)_{k \geq 0}$  has  $\epsilon_k > 0$  for all  $k \geq 0$ , it characterizes the errors of inexact proximal evaluation.
- (v)  $(f, g, L)$  satisfies Assumption 1.9, and let  $F = f + g$ .

Denote  $L_k = B_k + \rho_k$  for short. Given any initial condition  $v_{-1}, x_{-1} \in \mathbb{R}^n$ , the algorithm

generates the sequences  $(y_k, x_k, v_k)_{k \geq 0}$  such that they satisfy for all  $k \geq 0$ :

$$\begin{aligned} y_k &= \alpha_k v_{k-1} + (1 - \alpha_k) x_{k-1}, \\ x_k &\approx_{\epsilon_k} T_{L_k}(y_k), \\ D_f(x_k, y_k) &\leq \frac{B_k}{2} \|x_k - y_k\|^2, \\ v_k &= x_{k-1} + \alpha_k^{-1} (x_k - x_{k-1}). \end{aligned}$$

{lemma:inxt-apg-cnvg-prep1}

**Lemma 2.2 (inexact accelerated proximal gradient preparation stage I)**

Let  $(F, f, g, L)$ , and  $(\alpha_k, B_k, \rho_k, \epsilon_k)_{k \geq 0}$ , be given by Definition 2.1. Denote  $L_k = B_k + \rho_k$ . Then, for any  $\bar{x} \in \mathbb{R}^n$ , the sequences  $(y_k, x_k, v_k)_{k \geq 0}$  generated satisfy for all  $k \geq 1$  the inequality:

$$\begin{aligned} &\frac{\rho_k}{2} \|x_k - y_k\|^2 - \epsilon_k \\ &\leq (1 - \alpha_k)(F(x_{k-1}) - F(\bar{x})) + F(\bar{x}) - F(x_k) \\ &+ \max \left( 1 - \alpha_k, \frac{\alpha_k^2 L_k}{\alpha_{k-1}^2 L_{k-1}} \right) \frac{\alpha_{k-1}^2 L_{k-1}}{2} \|\bar{x} - v_{k-1}\|^2 - \frac{\alpha_k^2 L_k}{2} \|\bar{x} - v_k\|^2. \end{aligned}$$

When,  $k = 1$  it instead has:

$$\begin{aligned} &\frac{\rho_0}{2} \|x_0 - y_0\|^2 - \epsilon_0 \\ &\leq (1 - \alpha_0)(F(x_{-1}) - F(\bar{x})) + F(\bar{x}) - F(x_0) + \frac{\alpha_0^2 L_0}{2} \|\bar{x} - v_{-1}\|^2 - \frac{\alpha_0^2 L_0}{2} \|\bar{x} - v_0\|^2. \end{aligned}$$

*Proof.* Two intermediate results are in order before we can prove the inequality. Define  $z_k := \alpha_k \bar{x} + (1 - \alpha_k) x_{k-1}$  for short. It has for all  $k \geq 1$  the equality:

$$\begin{aligned} z_k - x_k &= \alpha_k \bar{x} + (1 - \alpha_k) x_{k-1} - x_k \\ &= \alpha_k x^+ + (x_{k-1} - x_k) - \alpha_k x_{k-1} \\ &= \alpha_k \bar{x} - \alpha_k v_k. \end{aligned} \tag{a}$$

It also has for all  $k \geq 1$  the equality:

$$\begin{aligned} z_k - y_k &= \alpha_k \bar{x} + (1 - \alpha_k) x_{k-1} - y_k \\ &= \alpha_k \bar{x} - \alpha_k v_{k-1}. \end{aligned} \tag{b}$$



Let's denote  $L_k = B_k + \rho_k$  for short. Recall that  $(f, g, L)$  satisfies Assumption 1.9, if we choose  $x = y_k$  so  $\tilde{x} = x_k \approx_{\epsilon_k} T_{L_k}(y_k)$ , and set  $z = z_k, \epsilon = \epsilon_k$  then Theorem 1.13 has:

$$\begin{aligned}
& \frac{\rho_k}{2} \|x_k - y_k\|^2 - \epsilon_k \\
& \leq F(z_k) - F(x_k) + \frac{L_k}{2} \|y_k - z_k\|^2 - \frac{L_k}{2} \|z_k - x_k\|^2 \\
& \stackrel{(1)}{\leq} \alpha_k F(\bar{x}) + (1 - \alpha_k) F(x_{k-1}) - F(x_k) + \frac{L_k}{2} \|y_k - z_k\|^2 - \frac{L_k}{2} \|z_k - x_k\|^2 \\
& \stackrel{(2)}{=} (1 - \alpha_k)(F(x_{k-1}) - F(\bar{x})) + F(\bar{x}) - F(x_k) + \frac{\alpha_k^2 L_k}{2} \|\bar{x} - v_{k-1}\|^2 - \frac{\alpha_k^2 L_k}{2} \|\bar{x} - v_k\|^2 \\
& \leq (1 - \alpha_k)(F(x_{k-1}) - F(\bar{x})) + F(\bar{x}) - F(x_k) \\
& \quad + \max \left( 1 - \alpha_k, \frac{\alpha_k^2 L_k}{\alpha_{k-1}^2 L_{k-1}} \right) \frac{\alpha_{k-1}^2 L_{k-1}}{2} \|\bar{x} - v_{k-1}\|^2 - \frac{\alpha_k^2 L_k}{2} \|\bar{x} - v_k\|^2.
\end{aligned}$$

At (1) we used the fact that  $F = f + g$  hence  $F$  is convex. At (2) we used (a), (b). Finally, if  $k = 0$ , then take the RHS of  $\stackrel{(1)}{=}$  then:

$$\begin{aligned}
& \frac{\rho_0}{2} \|x_0 - y_0\|^2 - \epsilon_0 \\
& \leq (1 - \alpha_0)(F(x_{-1}) - F(\bar{x})) + F(\bar{x}) - F(x_0) + \frac{\alpha_0^2 L_0}{2} \|\bar{x} - v_{-1}\|^2 - \frac{\alpha_0^2 L_0}{2} \|\bar{x} - v_0\|^2.
\end{aligned}$$

■

The following assumption encapsulate assumptions on the errors such that a near optimal convergence rate is still attainable by an algorithm that satisfies Definition 2.1.

**Assumption 2.3 (valid error schedule)** The following assumption is about an algorithm satisfying Definition 2.1, its parameters  $(\alpha_k, B_k, \rho_k, \epsilon_k)_{k \geq 0}$  in relation to its iterates  $(y_k, x_k, v_k)_{k \geq 0}$  and, some additional parameters  $(\beta_k)_{k \geq 0}, \mathcal{E}_0$  and  $p$ . Let

- (i)  $(\alpha_k, B_k, \rho_k, \epsilon_k)_{k \geq 0}, (F, f, g, L)$  and  $(y_k, x_k, v_k)_{k \geq 0}$  be given by Definition 2.1.
- (ii)  $\mathcal{E}_0 \geq 0$  be arbitrary;
- (iii) the sequence  $(\beta_k)_{k \geq 0}$  be defined as  $\beta_k := \prod_{i=1}^k \max \left( 1 - \alpha_i, \frac{\alpha_i^2 L_i}{\alpha_{i-1}^2 L_{i-1}} \right)$  for all  $k \geq 1$ , with the base case being  $\beta_0 = 1$ ;
- (iv)  $p \geq 1$  is some constant which will bound the error  $\epsilon_k$  relative to  $\rho_k \|x_k - y_k\|^2, \beta_k$  and,  $k$ .

In addition, we assume that the error parameter  $\epsilon_k \geq 0$  and over-relaxation parameter  $\rho_k$ , iterates  $x_k, y_k$  and  $\beta_k$  together satisfies for all  $k \geq 0$  the relations:

$$\frac{-\mathcal{E}_0 \beta_k}{k^p} \leq \frac{\rho_k}{2} \|x_k - y_k\|^2 - \epsilon_k.$$

The following proposition is a prototype of the convergence rate together with the error schedule that delivers convergence of algorithms satisfying Definition 2.1.

{prop:inxt-apg-cnvg-generic}

**Proposition 2.4 (generic convergence rate under valid error schedule)**

Let  $(F, f, g, L)$ ,  $(\alpha_k, B_k, \rho_k, \epsilon_k)_{k \geq 0}$ ,  $(\beta_k)_{k \geq 0}$ ,  $\mathcal{E}_0, p$  as assumed in Assumption 2.3. Fix any  $\bar{x} \in \mathbb{R}^n$  for all  $k \geq 0$  and assume that  $\alpha_0 = 1$ . Then for the iterates generated  $(y_k, x_k, v_k)_{k \geq 0}$  by the algorithm, for all  $k \geq 0$  they will satisfy:

$$F(x_k) - F(\bar{x}) + \frac{\alpha_k^2 L_k}{2} \|\bar{x} - v_k\|^2 \leq \beta_k \left( \frac{L_0}{2} \|\bar{x} - v_{-1}\|^2 + \mathcal{E}_0 \sum_{n=1}^k \frac{1}{n^p} \right).$$

*Proof.* Consider results from Lemma 2.2 has  $\forall k \geq 1$ :

$$\begin{aligned} & \frac{\rho_k}{2} \|x_k - y_k\|^2 - \epsilon_k \\ & \leq (1 - \alpha_k)(F(x_{k-1}) - F(\bar{x})) + F(\bar{x}) - F(x_k) \\ & + \max \left( 1 - \alpha_k, \frac{\alpha_k^2 L_k}{\alpha_{k-1}^2 L_{k-1}} \right) \frac{\alpha_{k-1}^2 L_{k-1}}{2} \|\bar{x} - v_{k-1}\|^2 - \frac{\alpha_k^2 L_k}{2} \|\bar{x} - v_k\|^2. \\ & \leq \max \left( 1 - \alpha_k, \frac{\alpha_k^2 L_k}{\alpha_{k-1}^2 L_{k-1}} \right) \left( F(x_{k-1}) - F(\bar{x}) + \frac{\alpha_{k-1}^2 L_{k-1}}{2} \|\bar{x} - v_{k-1}\|^2 \right) \\ & + F(\bar{x}) - F(x_k) - \frac{\alpha_k^2 L_k}{2} \|\bar{x} - v_k\|^2 \end{aligned}$$

For notation brevity, we introduce  $\beta_k, \Lambda_k$ :

$$\begin{aligned} \beta_0 &= 1, \\ \beta_k &:= \prod_{i=1}^k \max \left( 1 - \alpha_i, \frac{\alpha_i^2 L_i}{\alpha_{i-1}^2 L_{i-1}} \right), \\ \Lambda_k &:= -F(\bar{x}) + F(x_k) + \frac{\alpha_k^2 L_k}{2} \|\bar{x} - v_k\|^2. \end{aligned}$$

Now, suppose that in addition there is a non-negative sequence  $(\mathcal{E}_k)_{k \geq 0}$  such that

- (i) For all  $k \geq 0$ , it has  $\frac{-\mathcal{E}_k}{k^p} \leq \frac{\rho_k}{2} \|x_k - y_k\|^2 - \epsilon_k$  where  $p \geq 1$ ,
- (ii) For all  $k \geq 1$ , it has  $\mathcal{E}_k = \frac{\beta_k}{\beta_{k-1}} \mathcal{E}_{k-1}$ , with  $\mathcal{E}_0 \geq 0$ .

These conditions are equivalent to the assumption that  $\frac{-\mathcal{E}_0 \beta_k}{k^p} \leq \frac{\rho_k}{2} \|x_k - y_k\|^2 - \epsilon_k$  (which was stated in Assumption 2.3). One can show that by unrolling recurrence on  $\mathcal{E}_k$ . Then (2.1) implies  $\forall k \geq 1$ :

{ineq:inxt-apg-cnvg-generic-pitem-1}

$$\frac{-\mathcal{E}_k}{k^p} \leq \frac{\beta_k}{\beta_{k-1}} \Lambda_{k-1} - \Lambda_k \iff \Lambda_k \leq \frac{\beta_k}{\beta_{k-1}} \Lambda_{k-1} + \frac{\mathcal{E}_k}{k^p}. \quad (2.1)$$

Now, we show the convergence of  $\Lambda_k$ , using the relations of  $\mathcal{E}_k, \Lambda_k, \beta_k$  above.

$$\begin{aligned}
\Lambda_k &\leq \frac{\beta_k}{\beta_{k-1}} \Lambda_{k-1} + \frac{\mathcal{E}_k}{k^p} \\
&\leq \frac{\beta_k}{\beta_{k-1}} \Lambda_{k-1} + \frac{\beta_k}{\beta_{k-1}} \frac{\mathcal{E}_{k-1}}{k^p} \\
&= \frac{\beta_k}{\beta_{k-1}} \left( \Lambda_{k-1} + \frac{\mathcal{E}_{k-1}}{k^p} \right) \\
&\leq \frac{\beta_k}{\beta_{k-1}} \left( \frac{\beta_{k-1}}{\beta_{k-2}} \Lambda_{k-2} + \frac{\mathcal{E}_{k-1}}{(k-1)^p} + \frac{\mathcal{E}_{k-1}}{k^p} \right) \\
&= \frac{\beta_k}{\beta_{k-2}} \left( \Lambda_{k-2} + \frac{\mathcal{E}_{k-2}}{(k-1)^p} + \frac{\mathcal{E}_{k-2}}{k^p} \right) \\
&\dots \\
&\leq \frac{\beta_k}{\beta_1} \left( \Lambda_1 + \mathcal{E}_1 \sum_{n=2}^k \frac{1}{n^p} \right) \\
&\leq \frac{\beta_k}{\beta_1} \left( \frac{\beta_1}{\beta_0} \Lambda_0 + \mathcal{E}_1 \sum_{n=1}^k \frac{1}{n^p} \right) \\
&= \frac{\beta_k}{\beta_0} \left( \Lambda_0 + \mathcal{E}_0 \sum_{n=1}^k \frac{1}{n^p} \right).
\end{aligned}$$

Therefore, it points to the following inequality:

$$\begin{aligned}
&F(x_k) - F(\bar{x}) + \frac{\alpha_k^2 L_k}{2} \|\bar{x} - v_k\|^2 \\
&\leq \beta_k \left( F(x_0) - F(\bar{x}) + \frac{\alpha_0^2 L_0}{2} \|\bar{x} - v_0\|^2 + \mathcal{E}_0 \sum_{n=1}^k \frac{1}{n^p} \right).
\end{aligned}$$

Finally, when  $\alpha_0 = 1$ , then the results from 2.2 with  $k = 0$  simplifies the above inequality and give:

$$F(x_k) - F(\bar{x}) + \frac{\alpha_k^2 L_k}{2} \|\bar{x} - v_k\|^2 \leq \beta_k \left( \frac{L_0}{2} \|\bar{x} - v_{-1}\|^2 + \mathcal{E}_0 \sum_{n=1}^k \frac{1}{n^p} \right).$$

■

Now, it only remains to determine the sequence  $\alpha_k$  to derive a type of convergence rate for the algorithm because from the above theorem, we have the convergence rate  $\beta_k$  and, the error parameters  $\epsilon_k, \rho_k$  both controlled by the sequence  $(\alpha_k)_{k \geq 0}$ .

## 2.1 Convergence results of the outer loop

This section will give specific instances of the error control sequence  $(\epsilon_k)_{k \geq 0}$ ,  $(\rho_k)_{k \geq 0}$  and, momentum sequence  $(\alpha_k)_{k \geq 0}$  such that an optimal convergence rate of  $\mathcal{O}(1/k^2)$  can be achieved.

{ass:opt-mmntm-seq}

### Assumption 2.5 (the optimal momentum sequence)

Keeping everything assumed in Assumption 2.3 about  $(\alpha_k, B_k, \rho_k, \epsilon_k)_{k \geq 0}$ ,  $(F, f, g, L)$ ,  $(y_k, x_k, v_k)_{k \geq 0}$ ,  $(\beta_k)_{k \geq 0}$ ,  $\mathcal{E}_0$  and  $p$ . We assume in addition that the sequence  $(\alpha_k)_{k \geq 0}$  satisfies for all  $k \geq 0$  the equality:  $(1 - \alpha_k) = \alpha_k^2 L_k \alpha_{k-1}^{-2} L_{k-1}^{-1}$  and,  $p > 1$ .

{lemma:opt-mmntm-seq}

### Lemma 2.6 (the optimal momentum sequence is indeed valid and optimal)

Let  $(\alpha_k)_{k \geq 0}, (\beta_k)_{k \geq 0}$  be given by Assumption 2.5. If we choose  $\alpha_0 \in (0, 1]$  then for all  $k \geq 1$  it has:

{eqn:lemma:opt-mmntm-seq-result1}

$$\alpha_k = \frac{L_{k-1}}{2L_k} \left( -\alpha_{k-1}^2 + \left( \alpha_{k-1}^4 + 4\alpha_{k-1} \frac{L_k}{L_{k-1}} \right)^{1/2} \right) \in (0, 1) \quad (2.2)$$

For the sequence  $(\beta_k)_{k \geq 0}$  it has  $\forall k \geq 1$

{ineq:lemma:opt-mmntm-seq-result2}

$$\left( 1 + \alpha_0 \sqrt{L_0} \sum_{i=1}^k \sqrt{L_i^{-1}} \right)^{-2} \leq \beta_k \leq \left( 1 + \frac{\alpha_0 \sqrt{L_0}}{2} \sum_{i=1}^k \sqrt{L_i^{-1}} \right)^{-2}. \quad (2.3)$$

*Proof.* Firstly, we will show (2.2). We will prove using induction. Assume inductively that  $\alpha_k \in (0, 1]$ . We can solve for  $\alpha_k$  using the recursive equality from Assumption 2.5 which always has one of the strictly positive root given by:

$$\begin{aligned} \alpha_k &= \frac{L_{k-1}}{2L_k} \left( -\alpha_{k-1}^2 + \left( \alpha_{k-1}^4 + 4\alpha_{k-1} \frac{L_k}{L_{k-1}} \right)^{1/2} \right) \\ &= \frac{\alpha_{k-1}^2 L_{k-1}}{2L_k} \left( -1 + \left( 1 + \frac{4L_k}{L_{k-1} \alpha_{k-1}^2} \right)^{1/2} \right) \\ &\stackrel{(1)}{<} \frac{\alpha_{k-1}^2 L_{k-1}}{2L_k} \left( -1 + 1 + 2\alpha_{k-1}^{-2} \frac{L_k}{L_{k-1}} \right) \\ &= 1 \end{aligned}$$

At (1) we completed a square inside the radical, and use the assumption that  $\alpha_k > 0$  and,  $L_k > 0, L_{k-1} > 0$  because we had  $B_k > 0$ :

$$\begin{aligned} 1 + 4L_k L_{k-1}^{-1} \alpha_{k-1}^{-2} &= 1 + 4L_k L_{k-1}^{-1} \alpha_{k-1}^{-2} + 2L_k^2 L_{k-1}^{-2} \alpha_{k-1}^{-4} - 2L_k^2 L_{k-1}^{-2} \alpha_{k-1}^{-4} \\ &= (1 + 2L_k L_{k-1}^{-1} \alpha_{k-1}^{-2})^2 - 2\alpha_{k-1}^{-4} L_k^2 L_{k-1}^{-2} \\ &< (1 + 2L_k L_{k-1}^{-1} \alpha_{k-1}^{-2})^2. \end{aligned}$$

To see that  $\alpha_k > 0$ , recall the same fact that  $L_k > 0$ , and use the inductive hypothesis that  $\alpha_{k-1} \in (0, 1]$  then  $\alpha_k > 0$  because:

$$\begin{aligned}\alpha_k &= \frac{L_{k-1}}{2L_k} \left( -\alpha_{k-1}^2 + \left( \alpha_{k-1}^4 + 4\alpha_{k-1} \frac{L_k}{L_{k-1}} \right)^{1/2} \right) \\ &\geq \frac{L_{k-1}}{2L_k} \left( -\alpha_{k-1}^2 + (\alpha_{k-1}^4)^{1/2} \right) \\ &> 0.\end{aligned}$$

Therefore, inductively it holds that  $\alpha_k \in (0, 1)$  too.

We will now show (2.3). From the assumption that  $(\alpha_k)_{k \geq 0}$  has  $(1 - \alpha_k) = \alpha_k^2 L_k \alpha_{k-1}^{-2} L_{k-1}^{-1}$  for all  $k \geq 0$  and, the definition of  $\beta_k$ , it yields the following equalities:

$$\beta_k = \prod_{i=1}^k \max \left( 1 - \alpha_i, \frac{\alpha_i^2 L}{\alpha_{i-1}^2 L_{i-1}} \right) = \prod_{i=1}^k (1 - \alpha_i) = \prod_{i=1}^k \frac{\alpha_i^2 L}{\alpha_0^2 L_0} = \frac{\alpha_k^2 L_k}{\alpha_0^2 L_0}.$$

With the above relation, and the definitions of the sequences  $(\alpha_k)_{k \geq 0}, (\beta_k)_{k \geq 0}$  it satisfies for all  $k \geq 1$  the properties:

- (a)  $\beta_k$  is monotone decreasing and  $\beta_k > 0$  for all  $k \geq 0$  because  $\beta_k = \prod_{i=1}^k (1 - \alpha_i)$  and,  $\alpha_k \in (0, 1]$ .
- (b) It has the equalities  $\beta_k / \beta_{k-1} = (1 - \alpha_k) = \frac{\alpha_k^2 L_k}{\alpha_{k-1}^2 L_{k-1}}$  for all  $k \geq 1$ .

Using the above observations, we can show the chain of equalities  $\alpha_k^2 = (1 - \beta_k / \beta_{k-1})^2 = \beta_k L_0 \alpha_0^2 L_k^{-1}$  for all  $k \geq 0$ . This is true by first considering the relations  $\prod_{i=1}^k (1 - \alpha_i) = \beta_k$ :

$$\begin{aligned}\{eqn:opt-mmntm-seq-pitem1\} \quad & (1 - \alpha_k) = \beta_k / \beta_{k-1} \\ & \iff \alpha_k = 1 - \beta_k / \beta_{k-1} \\ & \implies \alpha_k^2 = (1 - \beta_k / \beta_{k-1})^2.\end{aligned} \tag{2.4}$$

Next, the recursive relation of  $(\alpha_k)_{k \geq 0}$  gives

$$\begin{aligned}\{eqn:opt-mmntm-seq-pitem2\} \quad & \alpha_k^2 = (1 - \alpha_k) \alpha_{k-1}^2 L_{k-1} L_k^{-1} \\ & = (1 - \alpha_k) \frac{\alpha_{k-1}^2 L_{k-1}}{\alpha_0^2 L_0} \frac{L_{k-1} \alpha_0^2 L_0}{L_k} \\ & = (\beta_k \beta_{k-1}^{-1}) \beta_{k-1} L_0 \alpha_0^2 L_k^{-1} \\ & = \beta_k L_0 \alpha_0^2 L_k^{-1}.\end{aligned} \tag{2.5}$$

Combining (2.4), (2.5) and the fact that  $\beta_k > 0 \forall k \geq 0$ , it would mean for all  $i \geq 1$  it has:

$$\begin{aligned} L_0 \alpha_0^2 L_i^{-1} &= \beta_i^{-1} \left( 1 - \frac{\beta_k}{\beta_{k-1}} \right)^2 \\ &= \beta_i \left( \beta_i^{-1} - \beta_{i-1}^{-1} \right)^2 \\ &= \beta_i \left( \beta_i^{-1/2} - \beta_{i-1}^{-1/2} \right)^2 \left( \beta_i^{-1/2} + \beta_{i-1}^{-1/2} \right)^2 \\ &= \left( \beta_i^{-1/2} - \beta_{i-1}^{-1/2} \right)^2 \left( 1 + \beta_i^{1/2} \beta_{i-1}^{-1/2} \right)^2. \end{aligned}$$

Since  $\beta_i$  is monotone decreasing, it has  $0 < \beta_i^{-1/2} \beta_{i-1}^{-1/2} \leq 1$ , this gives:

$$\beta_i^{-1/2} - \beta_{i-1}^{-1/2} \leq \alpha_0 \sqrt{\frac{L_0}{L_i}} \leq 2 \left( \beta_i^{-1/2} - \beta_{i-1}^{-1/2} \right).$$

Performing a telescoping sum, use the fact that  $\beta_0 = 1$  will yield the desired results after some algebraic manipulations.  $\blacksquare$

{prop:opt-cnvg-outr-loop}

**Proposition 2.7 ( $\mathcal{O}(1/k^2)$  optimal convergence rate of the outer loop)**

Let  $(f, g, L)$ ,  $(\alpha_k, B_k, \rho_k, \epsilon_k)_{k \geq 0}$ ,  $(\beta_k)_{k \geq 0}$ ,  $\mathcal{E}_0, p$  be given by Assumption 2.5. Assume in addition that

- (i) there exists  $\bar{x} \in \mathbb{R}^n$  that is a minimizer of  $F = f + g$ ;
- (ii) the sequence  $L_k := B_k + \rho_k$  is bounded, and there exists an  $L_{\max}$  such that for all  $k \geq 0$  it has  $L_{\max} \geq \max_{k \geq i \geq 0} L_i$ .

Then, it has  $\forall k \geq 0$ :

$$F(x_k) - F(\bar{x}) + \frac{\alpha_k^2 L_k}{2} \|\bar{x} - v_k\|^2 \leq \left( 1 + \frac{k \alpha_0 \sqrt{L_0}}{2 \sqrt{L_{\max}}} \right)^{-2} \left( \frac{L_0}{2} \|\bar{x} - v_{-1}\|^2 + \mathcal{E}_0 \sum_{n=1}^k \frac{1}{n^p} \right).$$

Since,  $p > 0$  the sum is convergent and hence the above inequality claims an overall convergence rate  $\mathcal{O}(1/k^2)$ .

*Proof.* We use the results from Lemma 2.6 because of the same assumption on  $(\alpha_k)_{k \geq 0}$ , then using the fact that  $L_k$  is bounded:

$$\beta_k \leq \left( 1 + \frac{\alpha_0 \sqrt{L_0}}{2} \sum_{i=1}^k \sqrt{L_i^{-1}} \right)^{-2} \leq \left( 1 + \frac{k \alpha_0 \sqrt{L_0}}{2 \sqrt{L_{\max}}} \right)^{-2}.$$

Then, apply Theorem 2.4 to obtain the desired results.  $\blacksquare$

**Remark 2.8** In this remark, we assert the fact that all assumptions made in the theorem can be satisfied on practice, and we will also bring attention to the current, and future roles played by some parameters used in the inexact algorithm.

Pay attention that  $\alpha_k$  had been determined in the above theorem (as seemed in Lemma 2.6), and  $(B_k)_{k \leq 0}$  is reserved for potential line search routine, the only parameter left undetermined in Definition 2.1 is the over-relaxation sequence  $(\rho_k)_{k \geq 0}$ . Since we only need the entire sequence  $L_k = \rho_k + B_k$  to be bounded above, we give the freedom to the practitioners to choose  $(\rho_k)_{k \geq 0}$ . However, this sequence  $\rho_k$  is not useless because it counters  $\epsilon_k$  in the proximal gradient inequality, this has huge impact in the earlier stage (the first few iterations) of the algorithm  $\|x_k - y_k\|$  is large. Of course, the parameter  $\mathcal{E}_0$  is also free to choose.

Finally, observe that from the above proof, in case when  $p = 1$ , the convergence rate would be  $\mathcal{O}(\log(k)/k^2)$ .

In the next subsection, we will show that under the assumption of the above theorem, there exists an error sequence  $\epsilon_k$  such that it can never approach 0 at an arbitrarily fast rate.

## 2.2 The fastest rate of which the error schedule can shrink

To have overall convergence claim of the algorithm, it's necessary to prevent the error schedule  $(\epsilon_k)_{k \geq 0}$  from crashing into zero too quickly. Following Assumption 2.3, in this section, we will provide the lower bound results for  $\epsilon_k$  in Theorem 2.7 to show that in the worst case it cannot approach zero arbitrarily fast, if we choose the largest possible  $\epsilon_k$  using  $\beta_k$ .

{lemma:err-schedule-lbnd}

### Lemma 2.9 (error schedule lower bound)

Let,  $(\alpha_k, B_k, \rho_k, \epsilon_k)_{k \geq 0}$ ,  $(\beta_k)_{k \geq 0}$ ,  $\mathcal{E}_0, p$  be given by Assumption 2.5,  $L_k := \rho_k + B_k$ . Let  $(\epsilon_k)_{k \geq 0}$  be given by  $\epsilon_k := \frac{\mathcal{E}_0 \beta_k}{k^p} + \rho_k \frac{\|x_k - y_k\|^2}{2}$ , then it will be a valid error sequence and so that it satisfies the assumption. If in addition, there exists  $L_{\min}$  such that for all  $k \geq 0$  such that it has  $L_{\min} \leq \min_{1 \leq i \leq k} L_i$  and, we assume  $\mathcal{E} > 0$ , then  $\epsilon_k$  admits the non-trivial lower bound:

$$\epsilon_k \geq \frac{\mathcal{E}_0}{k^p} \left( 1 + k\alpha_0 \sqrt{L_0} \sqrt{L_{\min}^{-1}} \right)^{-2} \geq \mathcal{O}(k^{-2-p}).$$

*Proof.* Recall Assumption 2.3, the largest valid error schedule is  $\epsilon_k = \frac{\mathcal{E}_0 \beta_k}{k^p} + \rho_k \frac{\|x_k - y_k\|^2}{2}$ . Then it has

$$\begin{aligned} \epsilon_k &\geq \frac{\mathcal{E}_0 \beta_k}{k^p} \\ &\stackrel{(1)}{\geq} \left(1 + \alpha_0 \sqrt{L_0} \sum_{i=1}^k \sqrt{L_i^{-1}}\right)^{-2} \frac{\mathcal{E}_0}{k^p} \\ &\stackrel{(2)}{\geq} \frac{\mathcal{E}_0}{k^p} \left(1 + k \alpha_0 \sqrt{L_0} \sqrt{L_{\min}^{-1}}\right)^{-2} \\ &\geq \mathcal{O}(k^{-2-p}). \end{aligned}$$

At (1), we used Lemma 2.6. At (2), we used that  $L_{\min} \leq L_i$  for all  $i = 0, 1, 2, \dots$  ■

### 3 Linear convergence for the proximal problem in the inner loop

In this section, we continue the discussion from Section 1.3. As an important reminder, we will fix the vector  $y \in \mathbb{R}^n$ , which is in the inexact proximal problem as a constant in this entire section.

The inner loop of the algorithm is another algorithm that evaluates  $x_k \approx_\epsilon T_{(B_k + \rho_k)}(y_k)$  for a given value of  $\epsilon, B + \rho$  and at the point  $y_k$ . To accomplish, let  $\lambda = (B_k + \rho_k)^{-1}$ , the algorithm needs to resolve the following equivalent inexact proximal point problem:

$$x_k \approx_\epsilon \text{prox}_{\lambda g}(y_k - \lambda \nabla f(y_k)).$$

Unfortunately recall that  $g = \omega \circ A$  in the context of the outer loop hence it's impossible to directly evaluate the proximal operator of  $g$  and hence we optimize the function  $\Phi_\lambda$  as given by (1.1).

In this section, we will show that there exists an algorithm generating the sequences  $z_n, v_n$  such that  $\mathbf{G}_\lambda(z_n, v_n)$  converges linearly if  $\Psi_\lambda$  satisfies the error bound conditions. Using results available in the literature, we will characterize the exact scenarios of  $\omega \circ A$  when this is possible to achieve. To start, the following assumption is the general error bound condition of a convex with additive composite structure.

{ass:pg-eb} **Assumption 3.1 (gradient mapping error bound)**  
The following assumption is about  $(F, f, g, L, S, \gamma)$ . Assume that

- (i)  $(f, g, L)$  satisfies Assumption 1.9,



- (ii) let  $\tau > 0$  be the step size inverse, let  $T_\tau$  be the proximal operator of  $f + g$  as given by  $T_\tau(x) := \text{prox}_{\tau^{-1}g}(x - \tau^{-1}\nabla f(x))$ ,
- (iii)  $S = \underset{x}{\text{argmin}} f(x) + g(x) \neq \emptyset$ ,
- (iv) the objective function is given by  $F = f + g$ .

Let the gradient mapping  $\mathcal{G}_\tau$  be defined as:  $\mathcal{G}_\tau(x) := \tau(x - T_\tau(x))$  In addition, assume that the optimization problem  $F$  satisfies the error bound condition if it has for all  $\tau \geq L, x \in \mathbb{R}^n$  there exists  $\gamma > 0$ :

$$\{\text{def:ista}\} \quad \|\mathcal{G}_\tau(x)\| \geq \gamma \text{dist}(x|S).$$

**Definition 3.2 (proximal gradient method)** Suppose that  $(f, g, L)$  satisfies Assumption 1.9. Let  $\tau \geq L$ , and  $x_0 \in \mathbb{R}^n$ . Then an algorithm is a proximal gradient method if it generates iterates  $(x_k)_{k \geq 0}$  such that they satisfy for all  $k \geq 1$ :

$$x_{k+1} = \text{prox}_{\tau^{-1}g}(x_k + \tau^{-1}\nabla f(x_k)).$$

### 3.1 Error bound and linear convergence of ISTA

The following theorem characterizes linear convergence of the proximal gradient method under gradient mapping error bound condition.

$\{\text{thm:lin-cnv-g-ista-eb}\}$  **Theorem 3.3 (linear convergence under gradient mapping error bound)** Assume that  $(F, f, g, L, S, \gamma)$  is given by Assumption 3.1. Under this assumption, the iterates  $(x_k)_{k \geq 0}$  given by Definition 3.2 satisfies for all  $k \geq 0, \bar{x} \in S$  and  $\tau \geq L$  the inequality:

$$F(x_{k+1}) - F(\bar{x}) \leq \left(1 - \frac{\gamma}{2\tau}\right) (F(x_k) - F(\bar{x})).$$

Hence, the algorithm generates  $F(x_k) - F(\bar{x}) \leq \mathcal{O}((1 - \gamma/(2\tau))^k)$ .

*Proof.* Two important immediate results will be presented first. Consider the proximal gradient inequality from Theorem 1.13, but with  $\rho = 0, \epsilon = 0, B = \tau$ , then for all  $x$  such

that  $\|\mathcal{G}_\tau(x)\| > 0$  it has for  $\tilde{x} = T_\tau(x), z \in \mathbb{R}^n$  the inequality

$$\begin{aligned} F(\tilde{x}) - F(z) &\leq \frac{\tau}{2}\|x - z\|^2 - \frac{\tau}{2}\|z - \tilde{x}\|^2 \\ &= -\frac{\tau}{2}\|x - \tilde{x}\|^2 + \tau\langle x - z, x - \tilde{x} \rangle \\ &= -\frac{1}{2\tau}\|\mathcal{G}_\tau(x)\|^2 + \langle x - z, \mathcal{G}_\tau(x) \rangle \\ &\leq -\frac{1}{2\tau}\|\mathcal{G}_\tau(x)\|^2 + \|x - z\|\|\mathcal{G}_\tau(x)\| \\ &= \|\mathcal{G}_\tau(x)\|^2 \left( \frac{\|x - z\|}{\|\mathcal{G}_\tau(x)\|} - \frac{1}{2\tau} \right). \end{aligned}$$

Now, for all  $z = \bar{x} \in S$ , from Assumption 3.4  $\exists \gamma > 0$  such that:

$$\frac{\|x - z\|}{\|\mathcal{G}_\tau(x)\|} \leq \frac{\|x - z\|}{\gamma \text{dist}(x|S)} \leq \frac{1}{\gamma}.$$

Hence, for all  $\bar{x} \in S$  it has

$$\{ \text{ineq:lin-cnvg-ista-eb-pitem1} \} \quad 0 \leq F(\tilde{x}) - F(\bar{x}) \leq \|\mathcal{G}_\tau(x)\|^2 \left( \frac{1}{\gamma} - \frac{1}{2\tau} \right). \quad (3.1)$$

Obviously it has  $\gamma^{-1} - (1/2)\tau^{-1} > 0$ . When  $z = x$ , we have the inequality:

$$\{ \text{ineq:lin-cnvg-ista-eb-pitem2} \} \quad F(\tilde{x}) - F(x) \leq -\frac{1}{2\tau}\|\mathcal{G}_\tau(x)\|^2. \quad (3.2)$$

To derive the linear convergence, we use (3.1) with  $x = x_k, \tilde{x} = x_{k+1}$ :

$$\begin{aligned} 0 &\leq \|\mathcal{G}_\tau(x_k)\|^2 \left( \frac{1}{\gamma} - \frac{1}{2\tau} \right) - F(x_{k+1}) + F(\bar{x}) \\ &= \frac{1}{2\tau}\|\mathcal{G}_\tau(x_k)\|^2 \left( \frac{2\tau}{\gamma} - 1 \right) - F(x_{k+1}) + F(\bar{x}) \\ &\stackrel{(1)}{\leq} \left( \frac{2\tau}{\gamma} - 1 \right) (F(x_k) - F(x_{k+1})) - F(x_{k+1}) + F(\bar{x}) \\ &= \left( \frac{2\tau}{\gamma} - 1 \right) (F(x_k) - F(\bar{x}) + F(\bar{x}) - F(x_{k+1})) - F(x_{k+1}) + F(\bar{x}) \\ &= \frac{2\tau}{\gamma}(F(\bar{x}) - F(x_{k+1})) + \left( \frac{2\tau}{\gamma} - 1 \right) (F(x_k) - F(\bar{x})). \end{aligned}$$

At (1) we used (3.2). Multiple both side by  $\frac{\gamma}{2\tau}$  then we are done. ■

{sec:conds-lin-cnvg-pp}

### 3.2 Conditions for linear convergence of the proximal problem

In this section, we will focus on the sufficient characterization of the proximal problem which allows proximal gradient method to achieve linear convergence rate. The following assumption characterizes a set of sufficient conditions of the proximal problem such that linear convergence rate of applying ISTA to dual proximal objective  $\Psi_\lambda$  can be achieved.

{ass:lin-cnvg-for-pp}

#### Assumption 3.4 (conditions for linear convergence of proximal problem)

This assumption is about  $(g, \omega, A, y, \Phi_\lambda, \Psi_\lambda, L_\Phi^\lambda, \gamma_\lambda)$ . Here are the assumptions

- (i)  $y \in \mathbb{R}^n$  is the vector of which the proximal problem is anchored at, fixed it to be an arbitrary vector.
- (ii) Assume  $(g, \omega, A)$  satisfies Assumption 1.15. The primal objective  $\Phi_\lambda$  is given by (1.1), and dual  $\Psi_\lambda$  by (1.2). This means if we let  $h(v) := \frac{1}{2\lambda} \|\lambda A^\top v - y\|^2 - \frac{1}{2\lambda} \|y\|^2$ , then  $\Psi_\lambda(v) = h(v) + \omega^*(v)$ .
- (iii) Next, assume  $\Psi_\lambda = h + \omega^*$  satisfies error bound condition (Assumption 3.1) with  $\Psi_\lambda = F$  and  $\gamma = \gamma_\lambda$  and  $f = h$ . Note that we can do this because  $h$  is quadratic hence obviously Lipschitz continuous and Lipschitz smooth.
- (iv) We assume in addition, the primal objective  $\Phi_\lambda$  of proximal problem is  $L_\Phi^\lambda$  Lipschitz continuous on its domain  $\text{dom } \Phi_\lambda = \text{dom}(g \circ A)$ .

The following definition specifies the algorithm that can achieve linear convergence rate with the assumptions above.

{def:ista-inner-lp}

#### Definition 3.5 (the ISTA inner loop algorithm)

Let  $\lambda > 0, \epsilon > 0$ , and  $(g, \omega, A, y, \Phi_\lambda, \Psi_\lambda, L_\Phi^\lambda, \gamma_\lambda)$  satisfies Assumption 3.4.

- (i) Let  $v_0 \in \text{dom } \omega^*$  be a feasible initial guess of  $\Psi_\lambda$ , and let  $\tau > \|AA^\top\|$  be the inverse stepsize.
- (ii) Define  $z_0 = y - \lambda A^\top v_0$  and smooth part of  $\Psi_\lambda$  as  $h := v \mapsto \frac{1}{2\lambda} \|\lambda A^\top v - y\|^2$ .

The algorithm that solves the proximal problem generates the primal dual sequences  $(z_n, v_n)$  such that for all  $n = 0, 1, 2, \dots$ , by the rules:

$$\begin{aligned} v_{n+1} &= \text{prox}_{\tau^{-1}\omega^*}(v_n - \tau^{-1}\nabla h(v_n)), \\ z_{n+1} &= y - \lambda A^\top v_{n+1}. \end{aligned}$$

Terminates if  $\mathbf{G}_\lambda(z_n, v_n) \leq \epsilon$  where  $\mathbf{G}_\lambda$  is given by (1.3).

**Remark 3.6** The value of  $\mathbf{G}_\lambda(z_n, v_n)$  is easy to compute, it only requires access to matrix  $A, A^\top$ , and the function  $\omega$ . In case when the proximal operator for  $\omega^*$  is nontrivial, we can use the Moreau identity and the proximal operator of  $\omega$  instead. The gradient  $\nabla f(v)$  is easy to compute, and it is:  $AA^\top v - Ay$ .

The following propositions precisely show that the linear convergence is achievable when Assumption 3.4 holds.

{prop:inn-loop-lin-cnvg}

**Proposition 3.7 (sufficient conditions of linear convergence of the inner loop)**

Let the parameters  $(g, \omega, A, y, \Phi_\lambda, \Psi_\lambda, L_\Phi^\lambda, \gamma_\lambda)$  of a proximal problem satisfy Assumption 3.4. Let  $\tau, v_0, \epsilon$  be given by an algorithm that satisfies Definition 3.5 and, it generates iterates  $(z_n, v_n)_{n \geq 0}$ . Let  $\bar{v}$  be a minimizer of  $\Psi_\lambda$ , then the followings are true:

{prop:inn-looplin-cnvg-item1}

(i) The sequence  $\Psi_\lambda(v_n) - \Psi_\lambda(\bar{v})$  converges linearly to zero.

{prop:inn-looplin-cnvg-item2}

(ii) If we let  $\bar{z} := y - \lambda A^\top \bar{v}$ , then  $\Phi_\lambda(z_n) - \Phi_\lambda(\bar{z})$  converges linearly to zero.

{prop:inn-looplin-cnvg-item3}

(iii) The duality gap has convergence  $\mathbf{G}_\lambda(z_n, v_n) \leq \mathcal{O}\left(\left(1 - \frac{\gamma_\lambda}{2\tau}\right)^{n/2}\right)$ .

{prop:inn-looplin-cnvg-item4}

(iv) If  $\mathbf{G}_\lambda(z_n, v_n) \leq \epsilon$ , then  $z_n \approx_\epsilon \text{prox}_{\lambda g}(y)$ .

*Proof.* Note that  $\Psi_\lambda$  satisfies error bound condition in Assumption 3.4, and Definition 3.5 has  $\tau \geq \|A^\top A\|$  which satisfies Definition 3.2 therefore the linear convergence results of Theorem 3.3 applies to  $\Psi_\lambda(v_n) - \Psi_\lambda(\bar{v})$ . Therefore, we claim:

{ineq:inn-loop-lin-cnvg-pitem1}

$$\Psi_\lambda(v_n) - \Psi_\lambda(\bar{v}) \leq \left(1 - \frac{\gamma_\lambda}{2\tau}\right)^n (\Psi_\lambda(v_0) - \Psi_\lambda(\bar{v})). \quad (3.3)$$

We will now proceed to proving (ii). Recall that the iterates  $z_n = y - \lambda A^\top v_n$ , and  $\bar{z} = y - \lambda A^\top \bar{v}$ , and we have Lipschitz continuity of  $\Phi_\lambda$ , hence Theorem 1.18 applies, in addition since  $\Phi_\lambda$  were assumed to be globally Lipschitz on  $\text{dom } \Phi_\lambda$  in Assumption 3.4, this means that we can replace  $K$  regardless of  $\bar{v}$ , and convergent primal sequence  $z_n \rightarrow \bar{z}$  with  $L_\Phi^\lambda$  which gives:

{ineq:inn-loop-lin-cnvg-pitem2}

$$\begin{aligned} \Phi_\lambda(z_n) - \Phi_\lambda(\bar{z}) &\leq L_\Phi^\lambda \sqrt{2\lambda} (\Psi_\lambda(v_n) - \Psi_\lambda(\bar{v}))^{1/2} \\ &\stackrel{(1)}{\leq} L_\Phi^\lambda \sqrt{2\lambda} \left(1 - \frac{\gamma_\lambda}{2\tau}\right)^{n/2} (\Psi_\lambda(v_0) - \Psi_\lambda(\bar{v}))^{1/2}. \end{aligned} \quad (3.4)$$

At (1), we used (3.3).

To prove (iii), by definition of duality gap in (1.3) we have:

$$\begin{aligned}
& \mathbf{G}_\lambda(z_n, v_n) \\
& \stackrel{(1)}{=} \Phi_\lambda(z_n) + \Psi_\lambda(v_n) + 0 \\
& = \Phi_\lambda(z_n) + \Psi_\lambda(v_n) - \Phi_\lambda(\bar{z}) - \Psi_\lambda(\bar{v}) \\
& = \Phi_\lambda(z_n) - \Phi_\lambda(\bar{z}) + \Psi_\lambda(v_n) - \Psi_\lambda(\bar{v}) \\
& \stackrel{(2)}{\leq} L_\Phi^\lambda \sqrt{2\lambda} \left(1 - \frac{\gamma_\lambda}{2\tau}\right)^{n/2} (\Psi_\lambda(v_0) - \Psi_\lambda(\bar{v}))^{1/2} + \left(1 - \frac{\gamma_\lambda}{2\tau}\right)^n (\Psi_\lambda(v_0) - \Psi_\lambda(\bar{v})) \\
& = \left(1 - \frac{\gamma_\lambda}{2\tau}\right)^{n/2} \left( L_\Phi^\lambda \sqrt{2\lambda} (\Psi_\lambda(v_0) - \Psi_\lambda(\bar{v}))^{1/2} + \left(1 - \frac{\gamma_\lambda}{2\tau}\right)^{n/2} (\Psi_\lambda(v_0) - \Psi_\lambda(\bar{v})) \right) \\
& \leq \left(1 - \frac{\gamma_\lambda}{2\tau}\right)^{n/2} \left( L_\Phi^\lambda \sqrt{2\lambda} + (\Psi_\lambda(v_0) - \Psi_\lambda(\bar{v}))^{1/2} \right) (\Psi_\lambda(v_0) - \Psi_\lambda(\bar{v}))^{1/2}.
\end{aligned}$$

At (1) we used strong duality because constraint qualification is in Assumption 1.15. At (2) we used (3.3), (3.4). To see (iv), apply Theorem 1.17. ■

**Remark 3.8** In practice, use the proximal operator of  $\omega^\star$  to choose a feasible  $v_0$ , equivalently we can translate any primal feasible solution into a dual feasible solution using Theorem 1.16.

As a prelude, to have a global convergence rate it requires an upper bound of the initial optimally gap of  $\Psi_\lambda$  of for all iterations of the outer loop. Suppose for each outer iteration  $k$ , the inner loop is initialized to start on  $v_0^{(k)}$ , then there must be an upper bound for all  $\Psi_\lambda(v_0^{(k)}) - \Psi_\lambda(\bar{v})$ .

{prop:1nrm-prox-problem}

The next proposition will characterize a precise case where Assumption 3.4 is true, and it's a case widely available in applications.

**Proposition 3.9 (1-norm problem)** *Let  $(g, \omega, A)$  satisfies Assumption 1.15. In addition, if  $g := \|\cdot\|_1$ , then the function  $\Psi_\lambda, \Phi_\lambda$  satisfies Assumption 3.4.*

*Proof.* To verify, it requires that

- (i) The function  $\Psi_\lambda$  satisfies the error bound condition as specified in Assumption 3.1;
- (ii)  $\Psi_\lambda$  is also a Lipschitz continuous function.

Take note that the dual has closed form  $g^\star(z) = \delta(z|\{x : \|x\|_1 \leq 1\})$ . The  $g^\star$  is an indicator of a polytope constraint. The objective function  $\Psi_\lambda(v) = \frac{1}{2\lambda} \|\lambda A^\top v - y\|^2 + \omega^\star(v) - \frac{1}{2\lambda} \|y\|^2$  so it's Lipschitz continuous on its domain, and it must have a set of minimizers  $S$  because its domain is compact. In addition,  $\Psi_\lambda$  fits the assumption of Necoara et al. [1, Theorem 8]. Therefore,  $\Psi_\lambda$  is quasi-strongly convex then by [1, Theorem 4], and [1, Theorem 7], it satisfies error bound condition as given in Assumption 3.1, hence Assumption 3.4 also.

Let  $\mu_f = \|A\|^2/\kappa_f$ , let  $\kappa_f = \theta^{-2}(A, C)$  be the Hoffman Constant as presented by Necoara et al. in [1, Section 4] where  $C$  is the constraint matrix of the inequality set  $\delta(z|\{x : \|x\|_1 \leq 1\})$ , then the error bound constant can be satisfied with

$$\gamma_\lambda = \frac{\kappa_f}{1 + \lambda\mu_f + \sqrt{1 + \lambda\mu_f}}.$$

■

**Remark 3.10** It is very difficult to obtain a lower estimate for  $\gamma$  in practice.

**Definition 3.11 (piecewise linear-quadratic function [2, Definition 10.20])** *A function  $f : \mathbb{R}^n \rightarrow \overline{\mathbb{R}}$  is piecewise linear-quadratic when  $\text{dom } f$  is the union of finitely many piecewise polyhedral set, such that on each polyhedral partition of  $f$  there exists  $\alpha \in \mathbb{R}, a \in \mathbb{R}^n, C \in \mathbb{R}^{n \times n}$  as a symmetric matrix is where  $f$  has the representation  $x \rightarrow \langle x, Cx \rangle + \langle a, x \rangle + c$ .*

The following proposition characterizes another case where error bound conditions of problem holds.

**Proposition 3.12 (convex PLQ proximal problem)**

**Proposition 3.13 (scenario II, convex PLQ linear composite)**

**Example 3.14 (inner loop linear convergence rate examples)**

## 4 Total complexity of the algorithm

This section puts results regarding the total complexity of the proposed inexact proximal gradient algorithm.

### 4.1 Inner loop complexity

We remind that the parameters in the inner loop will change depending on the outer loop's iteration. To discuss the convergence we unfortunately have to re-introduce the proximal problem in the context of the accelerated proximal gradient method of the outer loop.

Let  $(F, f, g, L)$  and sequence  $(\alpha_k, B_k, \rho_k, \epsilon_k)_{k \geq 0}$  satisfies Definition 2.1. Fix any  $k \geq 0$  to be the iteration counter of the outer loop. Let  $(g, \omega, A)$  satisfies Assumption 1.15. The inner

loop is an algorithm that solves the inexact proximal gradient problem  $x \approx_{\epsilon_k} T_{B_k + \rho_k}(y_k)$  which is equivalent to evaluating:

$$x_k \approx_{\epsilon_k} \text{prox}_{L_k^{-1}g}(y_k - L_k^{-1}\nabla f(y_k)).$$

Where,  $L_k := B_k + \rho_k$ . Let  $\lambda^{(k)} := L_k^{-1}$ ,  $\tilde{y}_k := y_k - L_k^{-1}\nabla f(y_k)$ , the proximal problem boils down to optimizing the following function  $\Phi_\lambda^{(k)}$  as defined by:

$$\{\text{eqn:primal-pp-k}\} \quad \Phi_\lambda^{(k)}(u) := \frac{1}{2\lambda^{(k)}} \|u - \tilde{y}_k\|^2 + \omega(Au). \quad (4.1)$$

And its dual which is

$$\{\text{eqn:dual-pp-k}\} \quad \Psi_\lambda^{(k)}(v) := \frac{1}{2\lambda} \|\lambda^{(k)} A^\top v - \tilde{y}_k\|^2 + \omega^*(v) - \frac{1}{2\lambda^{(k)}} \|\tilde{y}_k\|^2. \quad (4.2)$$

The primal dual gap is  $\mathbf{G}_\lambda^{(k)}(u, v) = \Phi_\lambda^{(k)}(u) + \Psi_\lambda^{(k)}(v)$ . The above are identical to proximal problem defined in Section 1.3 except for the introduction of iteration counter  $k$  from the outer loop, and we had specified  $\tilde{y}_k$  in relation to the outer loop. Finally, the inner loop algorithm is responsible for optimizing the optimality gap, it satisfies that  $\mathbf{G}_\lambda^{(k)}(u, v) \leq \epsilon_k$  and all results from Section 3.2 applies.

The following assumption specifies conditions where the complexity of the inner using ISTA can be globally bounded for all iteration of the accelerated proximal gradient algorithm of the outer loop. This is crucial for analyzing the global convergence rate of the algorithm.

$\{\text{ass:bounded-inn-cmplx}\}$

**Assumption 4.1 (globally bounded inner loop complexity)**

For any integer  $k \geq 0$ , this assumption is about parameters of the proximal problem  $(g, \omega, A, \tilde{y}_k, \Phi_\lambda^{(k)}, \Psi_\lambda^{(k)}, L_\Phi^{\lambda, (k)}, \gamma_\lambda^{(k)})$  introduced at the beginning of this section, iterates  $(z_n^{(k)}, v_n^{(k)})_{n \geq 0}$ , the primal dual solutions  $(\bar{z}^{(k)}, \bar{v}^{(k)})$  and additional parameters  $\bar{\lambda}, \underline{\lambda}, \gamma^\circ$ . The assumptions now follow.

- (i) As specified in this section, the parameter  $\lambda^{(k)} = (B_k + \rho_k)^{-1}$  are from the outer loop. We denote  $L_k = B_k + \rho_k$  for short. We assume in addition, there exists  $\lambda^{\min}, \lambda^{\max}$  such that:

$$-\infty < \lambda^{\min} \leq \inf_{k \in \mathbb{N} \cup \{0\}} \lambda^{(k)} \leq \sup_{k \in \mathbb{N} \cup \{0\}} \lambda^{(k)} \leq \lambda^{\max} < \infty.$$

- (ii) All the initial guesses given by an algorithm in the outer loop  $v_0^{(k)}$  and minimizers  $\bar{v}^{(k)}$  must have  $\sup_{k \in \mathbb{N} \cup \{0\}} \|v_0^{(k)} - \bar{v}^{(k)}\| < \infty$ .
- (iii) There exists smallest error bound constant  $\gamma^{\min}$  such that it lower bound all the constants for the proximal problem, i.e:  $\exists \gamma_{\min} : 0 < \gamma^{\min} \leq \inf_{k \in \mathbb{N} \cup \{0\}} \gamma_\lambda^{(k)}$ .

- (iv)  $\tilde{y}_k, \Phi_\lambda^{(k)}, \Psi_\lambda^{(k)}, L_\Phi^{\lambda, (k)}, \gamma_\lambda^{(k)}$  all satisfies error bound conditions (Assumption 3.4) with  $y = \tilde{y}_k, \Phi_\lambda = \Phi_\lambda^{(k)}, \Psi_\lambda = \Psi_\lambda^{(k)}, L_\Phi^\lambda = L_\Phi^{\lambda, (k)}$  and  $\gamma_\lambda = \gamma_\lambda^{(k)}$ .
- (v) The iterates of the inner loop  $(z_n^{(k)}, v_n^{(k)})_{k \geq 0}$  are produced by an algorithm that satisfies Definition 3.5.

The proposition that follows characterize the bare minimum requirements so that the inner loop's complexity depends only on required accuracies  $\epsilon$ , and parameter  $\lambda$  for all initial guess of the outer loop.

{prop:inner-lp-cmplx}

**Proposition 4.2 (inner loop complexity can be bounded globally)**

Let  $(g, \omega, A, \tilde{y}_k, \Phi_\lambda^{(k)}, \Psi_\lambda^{(k)}, L_\Phi^{\lambda, (k)}, \gamma_\lambda^{(k)})$ ,  $(z_n^{(k)}, v_n^{(k)})_{n \geq 0}$ ,  $(\bar{z}^{(k)}, \bar{v}^{(k)})$  and  $C_1$  satisfies Assumption 4.1.

{prop:inner-lp-cmplx-result1}

{prop:inner-lp-cmplx-result2}

- (i) If  $\lambda_1 > \lambda_2$  then  $\Psi_{\lambda_1}(v) \geq \Psi_{\lambda_2}(v)$  for all  $v \in \text{dom } \omega^*$ .
- (ii) As a consequence of the former, the sequence  $L_\Phi^{\lambda, (k)}$  as specified back in Fact 1.18 is bounded above so  $\sup_{k \in \mathbb{N} \cup \{0\}} L_\Phi^{\lambda, (k)} < \infty$ .
- (iii) There exists a constant  $C_1$  that bounds the optimality gap of all initial guesses:

$$\infty > C_1 \geq \sup_{k \in \mathbb{N} \cup \{0\}} \left\{ \Psi_\lambda^{(k)}(v_0^{(k)}) - \Psi_\lambda^{(k)}(\bar{v}^{(k)}) \right\}.$$

{prop:inner-lp-cmplx-result3}

{prop:inner-lp-cmplx-result4}

- (iv) There exists a constant  $\infty > C_2 \geq \sqrt{C_1} \sup_{k \in \mathbb{N} \cup \{0\}} \left\{ L_\Phi^{\lambda, (k)} \sqrt{2\lambda^{(k)}} + \sqrt{C_1} \right\}$ .
- (v) Denote  $N_\epsilon$  to be the total number of iteration of inner loop needed to achieve  $\mathbf{G}_\lambda^{(k)}(z_n, v_n) \leq \epsilon$ , then it's sufficient to have

$$N_\epsilon = \left\lceil 2 \frac{\log(\epsilon/C_2)}{\log\left(1 - \frac{\gamma_{\min}^{\min}}{2\lambda^{\max}}\right)} \right\rceil.$$

*Proof.* We will start with (i). After some algebra it can be shown that  $\Psi_\lambda^{(k)}(v) = \frac{\lambda}{2} \|A^\top v\|^2 - \langle A^\top v, \tilde{y}_k \rangle + \omega^*(v)$ . From then on it's direct and it has  $\Psi_{\lambda_1}(v) - \Psi_{\lambda_2}(v) = \frac{\lambda_1 - \lambda_2}{2} \|A^\top v\|^2 \geq 0$ .

To see (ii), observe that  $\Phi_\lambda^{(k)}$  is  $L_\Phi^{\lambda, (k)}$  Lipschitz continuous from Assumption 3.4, 4.1. By its definition in (4.1),  $\forall k \geq 0$  it satisfies  $L_\Phi^{\lambda, (k)}$  is the largest when  $\lambda = \lambda^{\min}$ . And since  $\lambda^{(k)}$  is finite, it has  $\sup_{k \in \mathbb{N} \cup \{0\}} L_\Phi^{\lambda, (k)} < \infty$  too.

To see (iii), we use (i) and it's sufficient to have  $C_1 = \Psi_{\lambda^{\max}}^{(k)}$  ■



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