

# Static analyzer for the Nix Expression Language

Статический анализатор для Nix Expression Language

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*Source code:* [github.com/ilyakooo0/tix](https://github.com/ilyakooo0/tix)

*Nix is a powerful package manager that makes package management reliable and reproducible.*

– **[nixos.org](https://nixos.org)**

# Nix

- Package manager
- Build system\*
- Ad hoc development environments
- Easy (cross-community) dependency management
- Reproducible builds
- *And much more: NixOS, NixOps, Build Caching, Docker ...*

# Nix Expression Language

- Purely functional
- Lazy
- **Strict, but not static typing**

```
let pkgs = import ./nixpkgs {};  
in derivation {  
    name = "simple";  
    builder = "${pkgs.bash}/bin/bash";  
    args = [ ./simple_builder.sh ];  
    gcc = pkgs.gcc;  
    coreutils = pkgs.coreutils;  
    src = ./simple.c;  
    system = builtins.currentSystem;  
}
```

# Terms and definitions

**Type system** – a logical system comprising a set of rules that assigns a property called a type to the various constructs of a computer program

**Type checking** – the process of verifying and enforcing the constraints of types.

**Hindley-Milner type system** – a classical type system for the lambda calculus with parametric polymorphism.

**Parametric polymorphism** – a type of polymorphism where types are specified by abstract symbols that can represent any type.

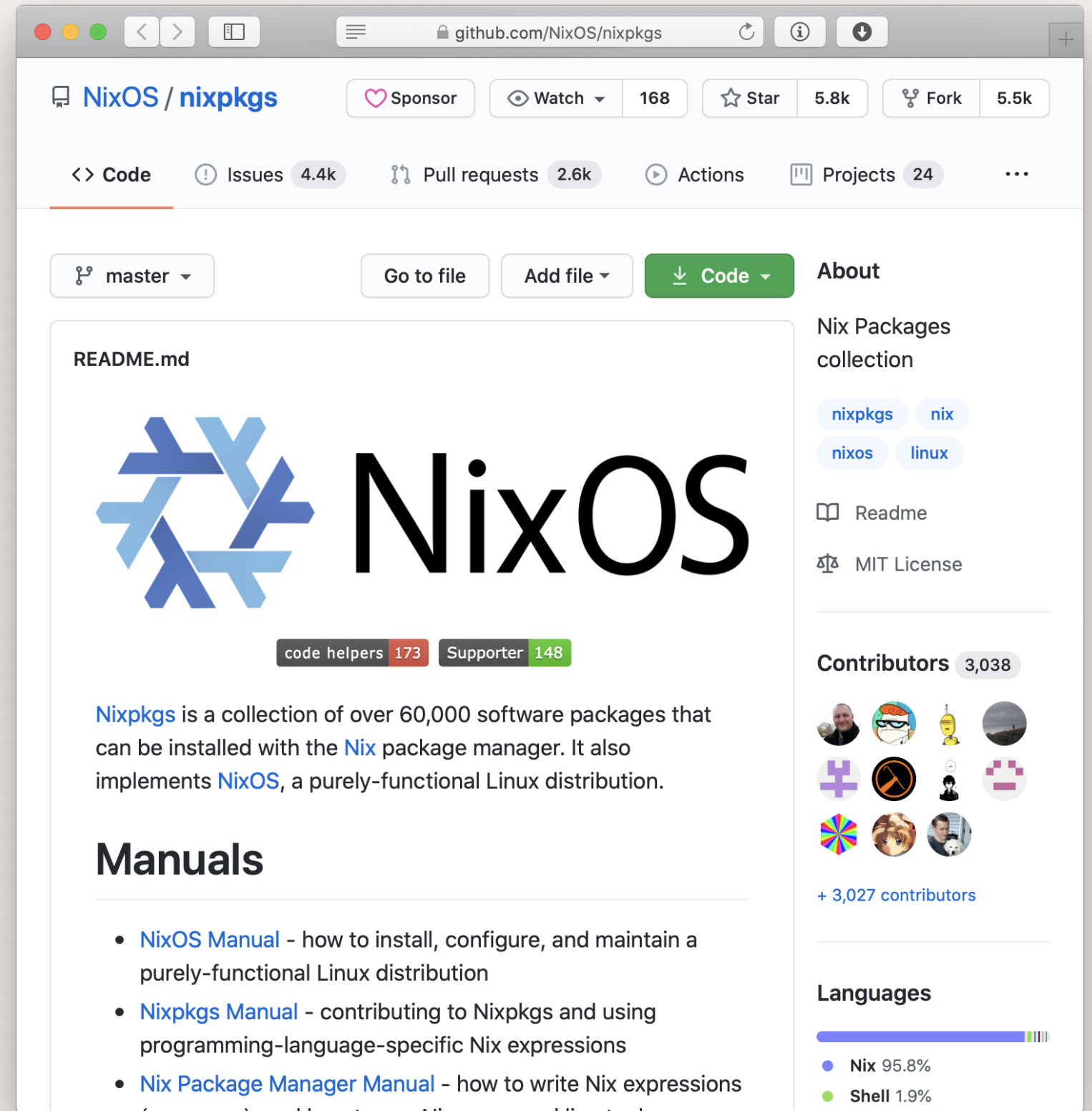
**Row polymorphism** – a kind of polymorphism that allows one to write programs that are polymorphic on record field types.

# Relevance

The Nix ecosystem is slowly becoming popular.

The central repository contains more than 60000 packages.

There is a need to make writing Nix expressions less error-prone.



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Nix Expression Language: **Strict, but not static typing.**

Adding static typing would:

- Increase reliability of code
- Lower costs of development (resolving issue earlier is cheaper)

*The typechecker should help find issues in existing code*

# The goal

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1. Develop a type static system for the existing Nix expression language

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2. Develop a type checker for this type system

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- [github.com/haskell-nix/hnix](https://github.com/haskell-nix/hnix) – has a module with a type checker. Is very experimental and is not used.



# Existing approaches

Type systems for functional languages is an actively researched field.

The Hindley–Milner type system[1] is a simple, but well-studied type system for a very simple language. It is fully decidable, does not require annotations. It is the basis for many real-world type systems.

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- The system should report the use of undefined terms.
- The system should report the resulting types of expressions to the user.

# Methods and algorithms

The type system will be heavily influenced by the Hindley–Milner[1] type system.

The type system will implement parametric polymorphism to support code that does not require a specific type.

The type system will implement arbitrary-ranked polymorphism[3] with deep instantiation to disambiguate types in existing untyped code.

The type system will implement row polymorphism to work with attribute sets.



# Used tools

1. Haskell
2. Visual Studio Code
3. Haskell Language Server
4. Haskell Stack

# Implementation

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1. Inferring the types of terms and collecting constraints
2. Solving the collected constraints

# Current results

A typechecker executable that successfully typechecks *most* of the selected files from *nixpkgs*.

The typechecker currently has some support for all language constructs.

```
test/golden/importless/purpose.nix: OK
test/golden/importless/mptcp.nix: OK
test/golden/importless/ocserv.nix: OK (0.02s)
test/golden/importless/gaps.nix: OK (0.01s)
test/golden/importless/magic-wormhole-mailbox-server.nix: OK (0.03s)
test/golden/importless/mullvad-vpn.nix: OK (0.03s)
test/golden/importless/fakeroute.nix: FAIL (0.01s)
  ([],[],[CanNotUnify (NAtomic Bool) (NBruijn (DeBruijn 1 0)),KeysDontMatch (fromL
  CallStack (from HasCallStack):
    error, called at test/src/Main.hs:35:88 in main:Main
test/golden/importless/trusted.nix: FAIL (0.02s)
  ([],[],[KeysDontMatch (fromList ["name"])],[])
  CallStack (from HasCallStack):
    error, called at test/src/Main.hs:35:88 in main:Main
test/golden/importless/strongswan.nix: FAIL (0.03s)
  ([([SrcSpan {spanBegin = SourcePos {sourceName = "<string>", sourceLine = Pos 29,
  29, sourceColumn = Pos 13}},UndefinedVariable "builtins"),(SrcSpan {spanBegin = Sou
  SourcePos {sourceName = "<string>", sourceLine = Pos 5, sourceColumn = Pos 21}},Und
  ,"strongswan"])),CanNotUnify (NAttrSet (fromList [("strongswan",[] :=> NAttrSet (from
  NAtomic String)])))) (List (NAtomic String)),CanNotUnify (NAttrSet (fromList []))
  atch (fromList ["example"])),CanNotUnify (NAttrSet (fromList [])) (NBruijn (DeBruijn
  CallStack (from HasCallStack):
    error, called at test/src/Main.hs:35:88 in main:Main
test/golden/mine/polyApply.nix: OK
test/golden/mine/dhall-packages.nix: OK
test/golden/mine/with.nix: OK
test/golden/mine/inftype.nix: FAIL
  ([],[],[InfinityType (NAttrSet (fromList [("a",[] :=> NAtomic Integer),("f",[] :
  CallStack (from HasCallStack):
    error, called at test/src/Main.hs:35:88 in main:Main
test/golden/mine/other_row.nix: OK
test/golden/mine/agda-packages.nix: OK
test/golden/mine/polyArgs.nix: OK
test/golden/mine/z3.nix: OK
test/golden/mine/other.nix: OK (0.01s)
test/golden/mine/org-generated.nix: OK
test/golden/mine/row.nix: FAIL
  ([],[],[CanNotUnify (NAtomic Integer) (NAtomic String)],[])
  CallStack (from HasCallStack):
    error, called at test/src/Main.hs:35:88 in main:Main
test/golden/mine/unionfs.nix: FAIL
  ([],[],[KeysDontMatch (fromList ["boot","system"])],[])
  CallStack (from HasCallStack):
    error, called at test/src/Main.hs:35:88 in main:Main
test/golden/mine/0.9.nix: OK
test/golden/mine/curry.nix: OK
test/golden/mine/nuget.nix: OK

1244 out of 2944 tests failed (15.67s)
```

# Current results

## Paper status

- 1. Nix Expression Language (80%)
- 2. Existing typecheckers (100%)
- 3. Existing approaches (100%)
- 4. Architecture (20%)
- 5. Implementation (0%)
- 6. Conclusion (60%)

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- A lot of minor big fixes in type inference
- Add type information about built-in functions

# Potential future developments

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# Potential future developments

- Add support for user-supplied type annotations.
  - Fork parser
  - Bidirectional type inference[4]
- Type holes
- Language Server Integration

# References

1. Damas L., Milner R. Principal type-schemes for functional programs //Proceedings of the 9th ACM SIGPLAN-SIGACT symposium on Principles of programming languages. – 1982. – C. 207-212.
2. Jones S. P. et al. Practical type inference for arbitrary-rank types //Journal of functional programming. – 2007. – T. 17. – N°. 1. – C. 1-82. MLA
3. Stuckey P. J., Sulzmann M., Wazny J. Improved inference for checking type annotations //arXiv preprint cs/0507036. – 2005. MLA

# Demonstration

# Parametric polymorphism

let  $f = x: y: \{ a = x; b = y + 1; \};$  in  $f$

$\forall \alpha. \alpha \rightarrow \text{Integer} \rightarrow \{a = \alpha; b = \text{Integer};\}$

# Parametric polymorphism

```
let f = x: y: { a = x; b = y + 1; }; in f
```

$\forall \alpha. \alpha \rightarrow \text{Integer} \rightarrow \{a = \alpha; b = \text{Integer};\}$

```
let f = x: y: { a = x; b = y + 1; }; in f "hello" 42
```

$\{a = \text{String}; b = \text{Integer};\}$

# Row-polymorphism with *attribute sets*

let  $f = x: x.a;$  in  $f$

$\forall \alpha \beta. (\alpha.a = \beta) \Rightarrow \alpha \rightarrow \beta$

# Row-polymorphism with *attribute sets*

```
let f = x: x.a; in f
```

$$\forall \alpha \beta. (\alpha.a = \beta) \Rightarrow \alpha \rightarrow \beta$$

```
let f = x: x.a; in f {a = 1; b = "hello";}
```

Integer

# Row-polymorphism with *attribute sets*

```
let f = x: x.a; in f
```

$$\forall \alpha \beta. (\alpha.a = \beta) \Rightarrow \alpha \rightarrow \beta$$

```
let f = x: x.a; in f {a = "1"; b = "hello";}
```

String



# Row-polymorphism with *attribute sets*

```
let f = x: x.a; in f
```

$$\forall \alpha \beta. (\alpha.a = \beta) \Rightarrow \alpha \rightarrow \beta$$

```
let f = x: x.a; in f {b = "hello";}
```

Error: key "a" NotPresent in { b = String; }

# Row-polymorphism with *attribute sets*

```
let f = x: x // {a = 69;}; in f
```

$$\forall \alpha \beta. (\alpha // \{a = \text{Integer};\} \sim \beta) \Rightarrow \alpha \rightarrow \beta$$

# Row-polymorphism with *attribute sets*

```
let f = x: x // {a = 69;}; in f
```

$$\forall \alpha \beta. (\alpha // \{a = \text{Integer};\} \sim \beta) \Rightarrow \alpha \rightarrow \beta$$

```
let f = x: x // {a = 69;}; in f {a = "hello"; b = "world";}  
{a = Integer; b = String;}
```

# Row-polymorphism with *attribute sets*

```
let f = x: x // {a = 69;}; in f
```

$$\forall \alpha \beta. (\alpha // \{a = \text{Integer};\} \sim \beta) \Rightarrow \alpha \rightarrow \beta$$

```
let f = x: x // {a = 69;}; in f {}  
  
{a = Integer;}
```