

YSC2229: Introductory Data Structures and Algorithms

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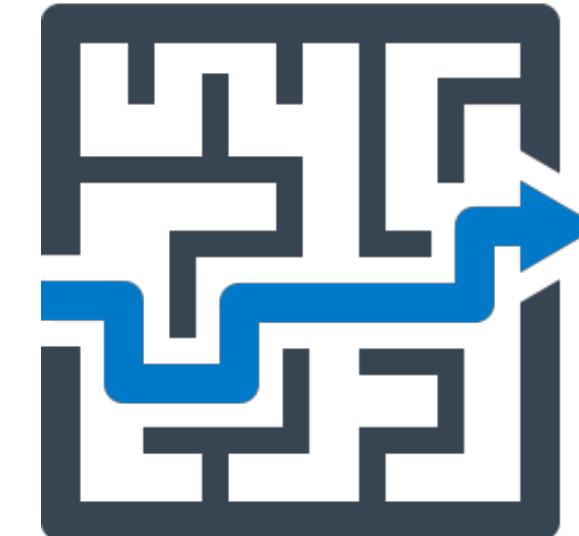
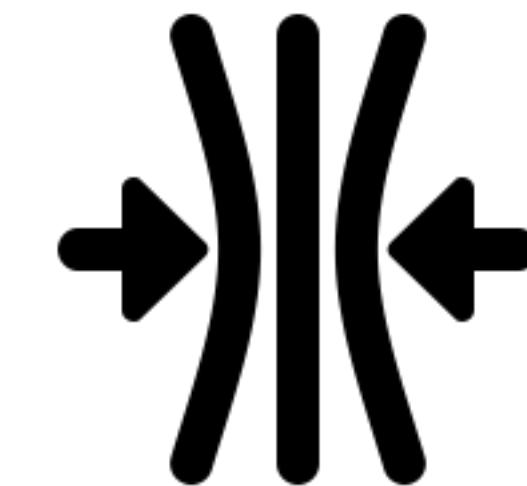
Some Terminology

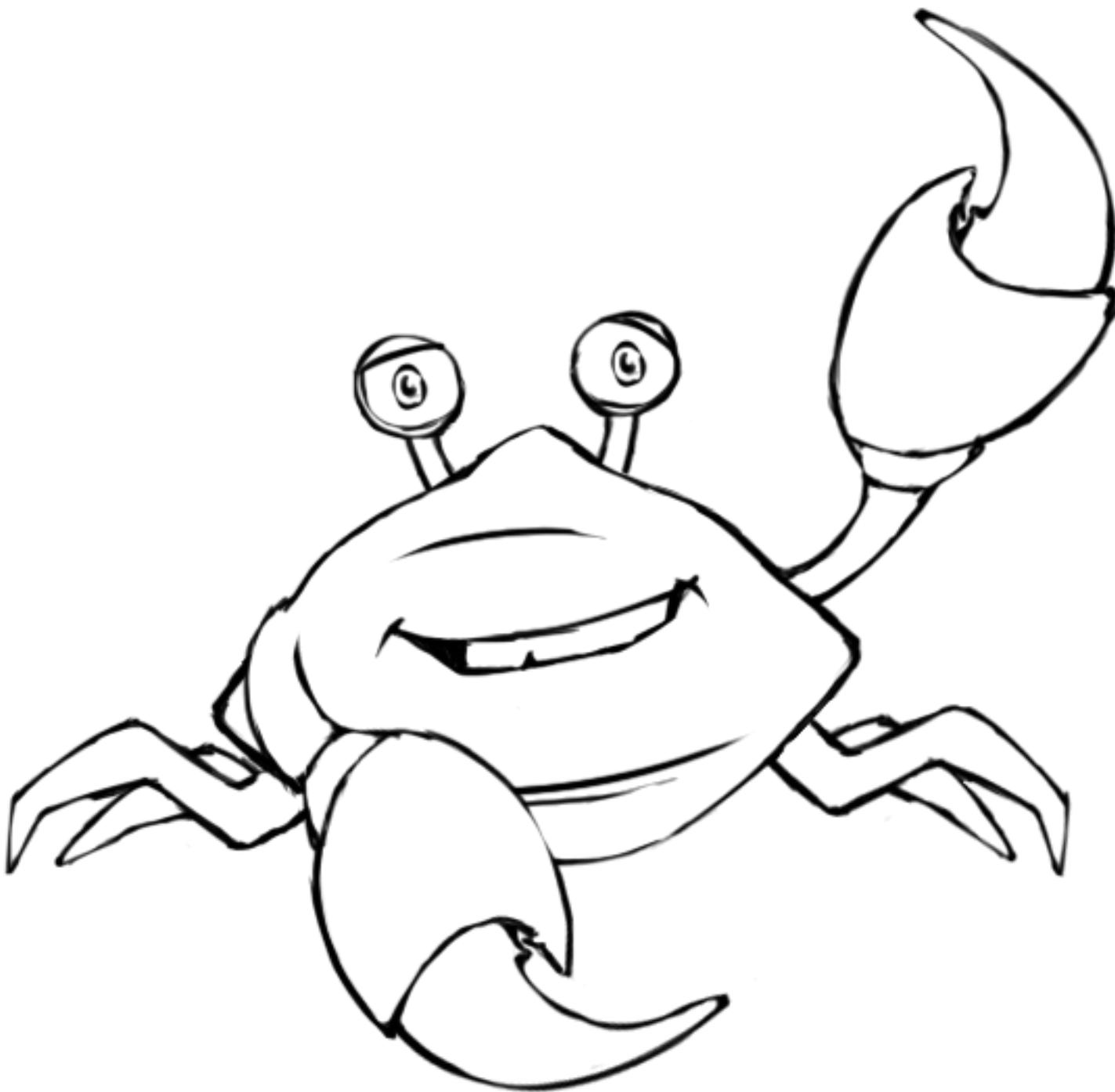
- *Data* represents *information*
- *Computations* represent *data processing*
- An *algorithm* is a sequence of computational steps that transform the *input* data (given) into the *output* data (wanted).

What this course is about?

Solving computational problems

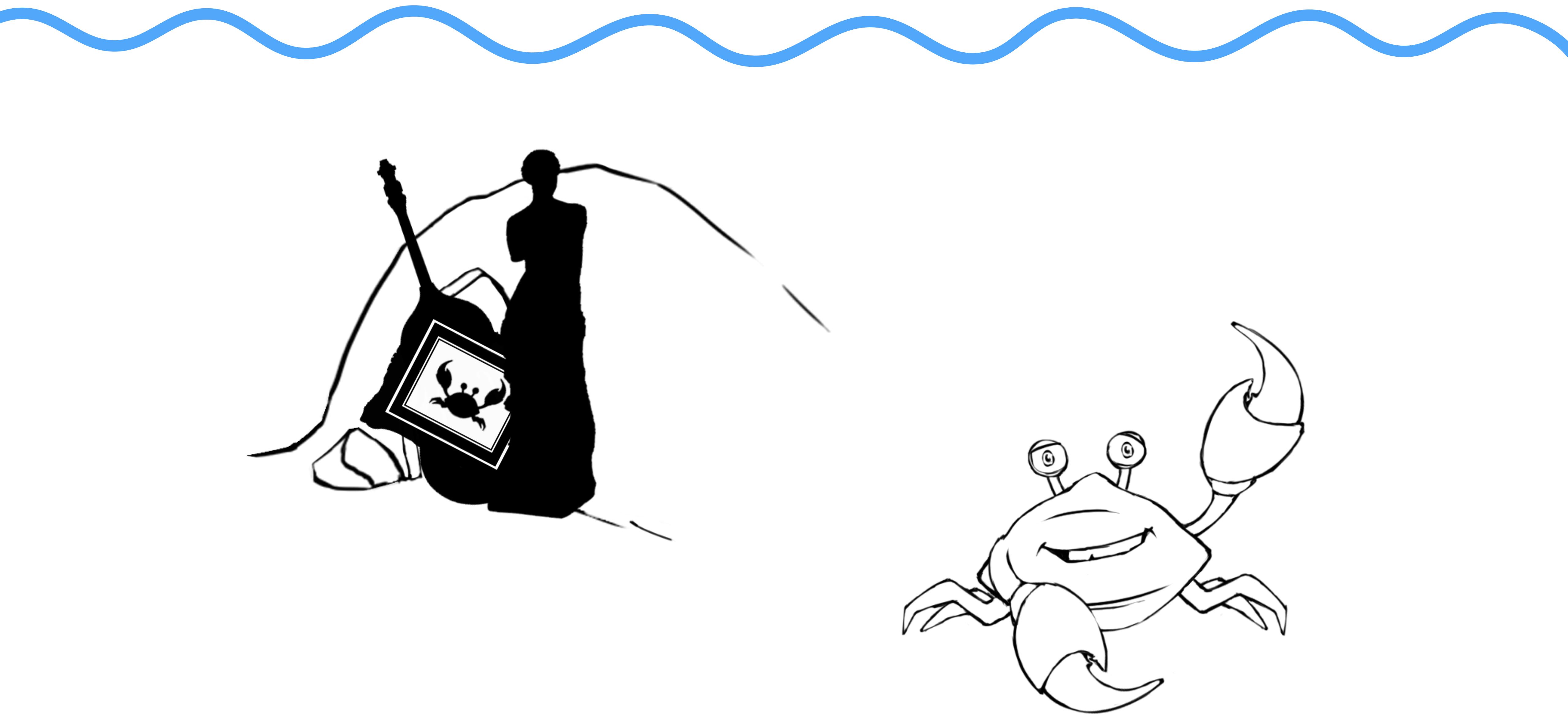
- **Searching:**
finding a word in a text or an article to buy on Amazon
- **Storing and retrieving data:**
representing files in you computer
- **Data compression/decompression:**
transferring files on the internet
- **Path finding:**
getting from a point A to point B in the most efficient way
- **Geometric problems:**
finding the closes fuel station, shape intersection



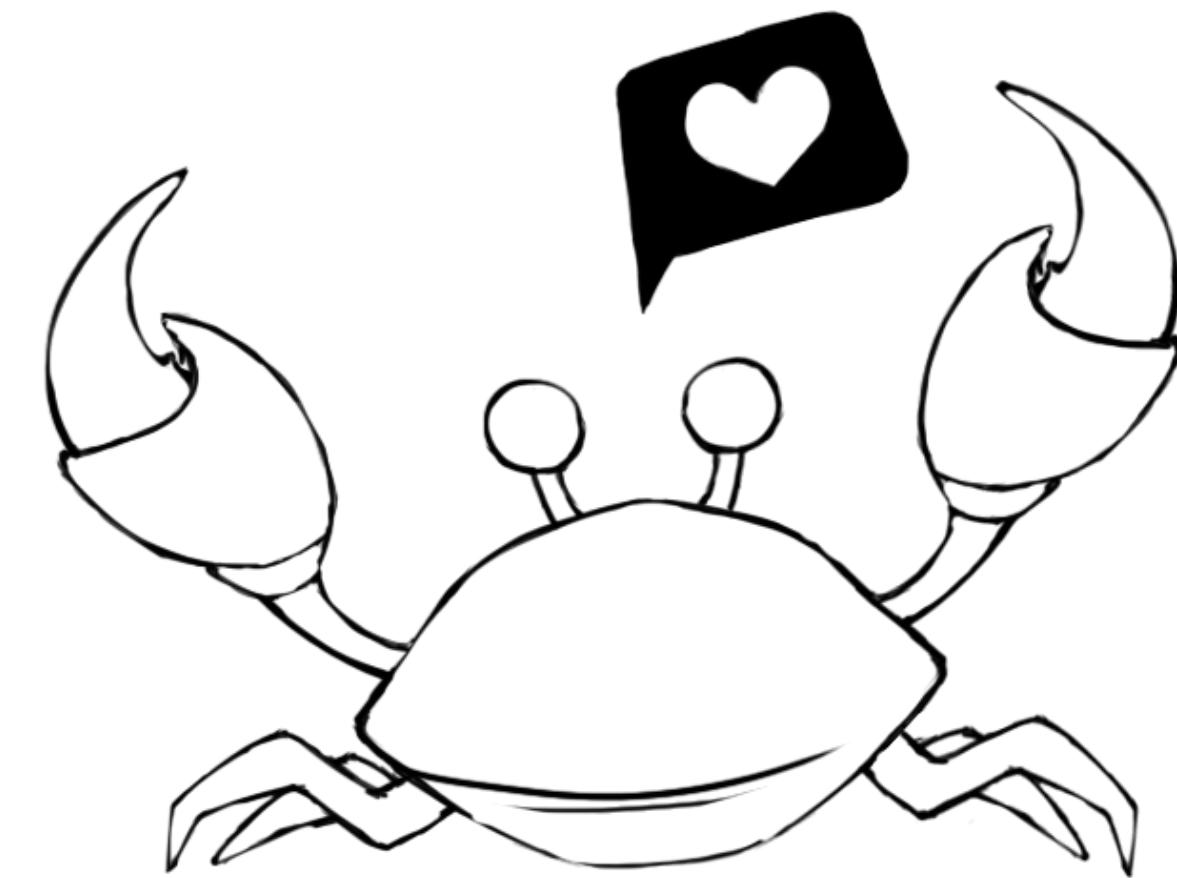


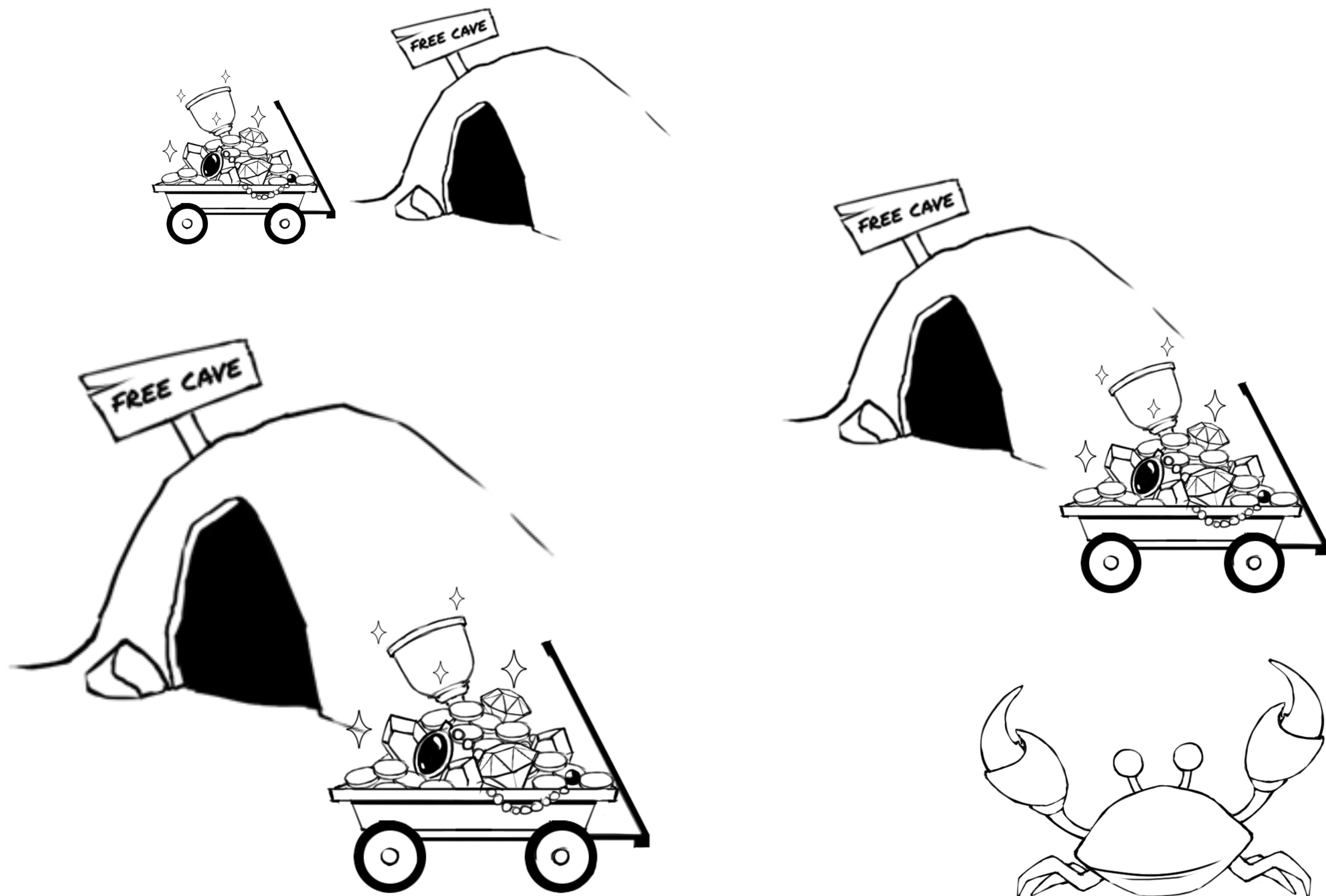
Torpe

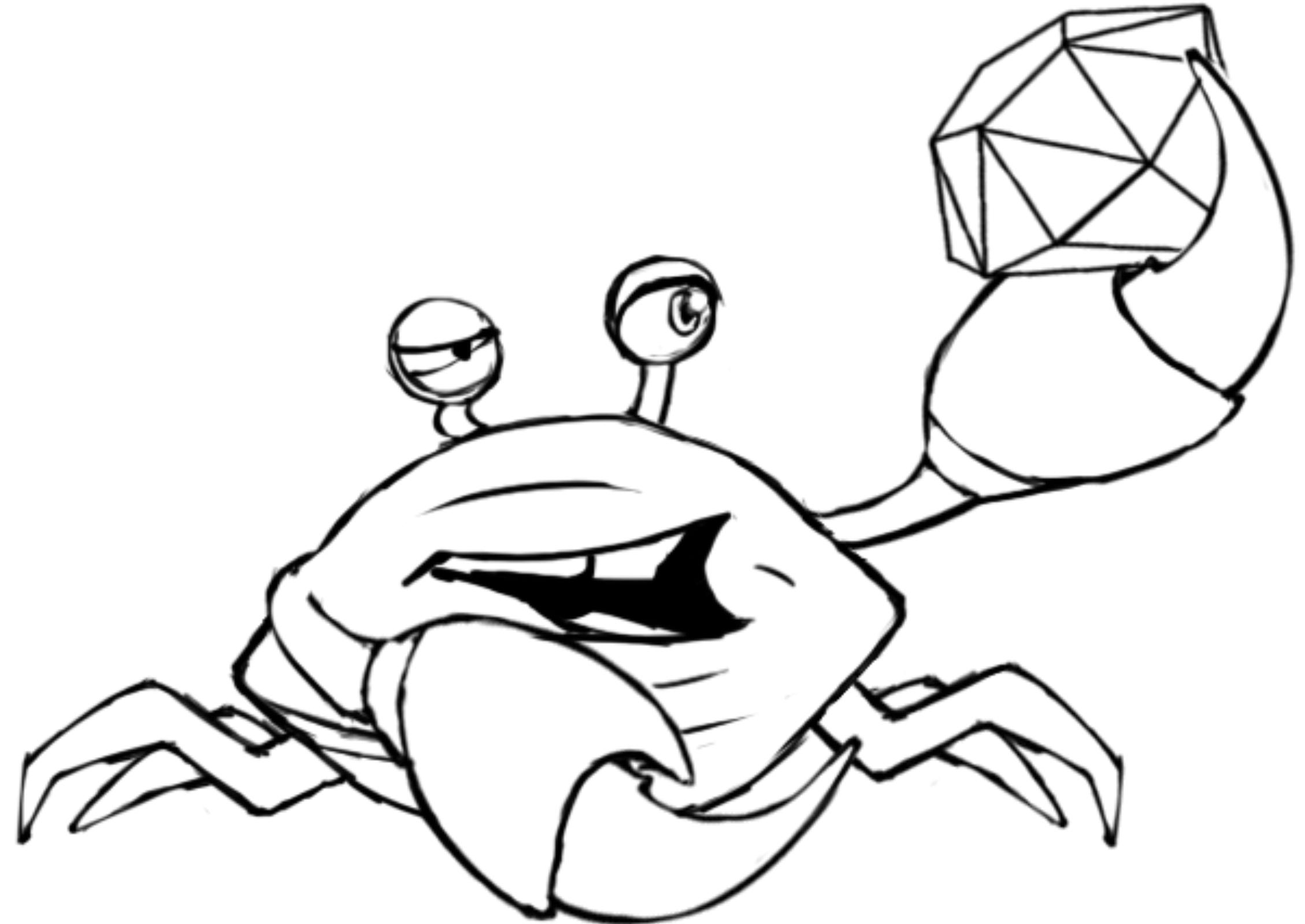
(the prosperous crab)





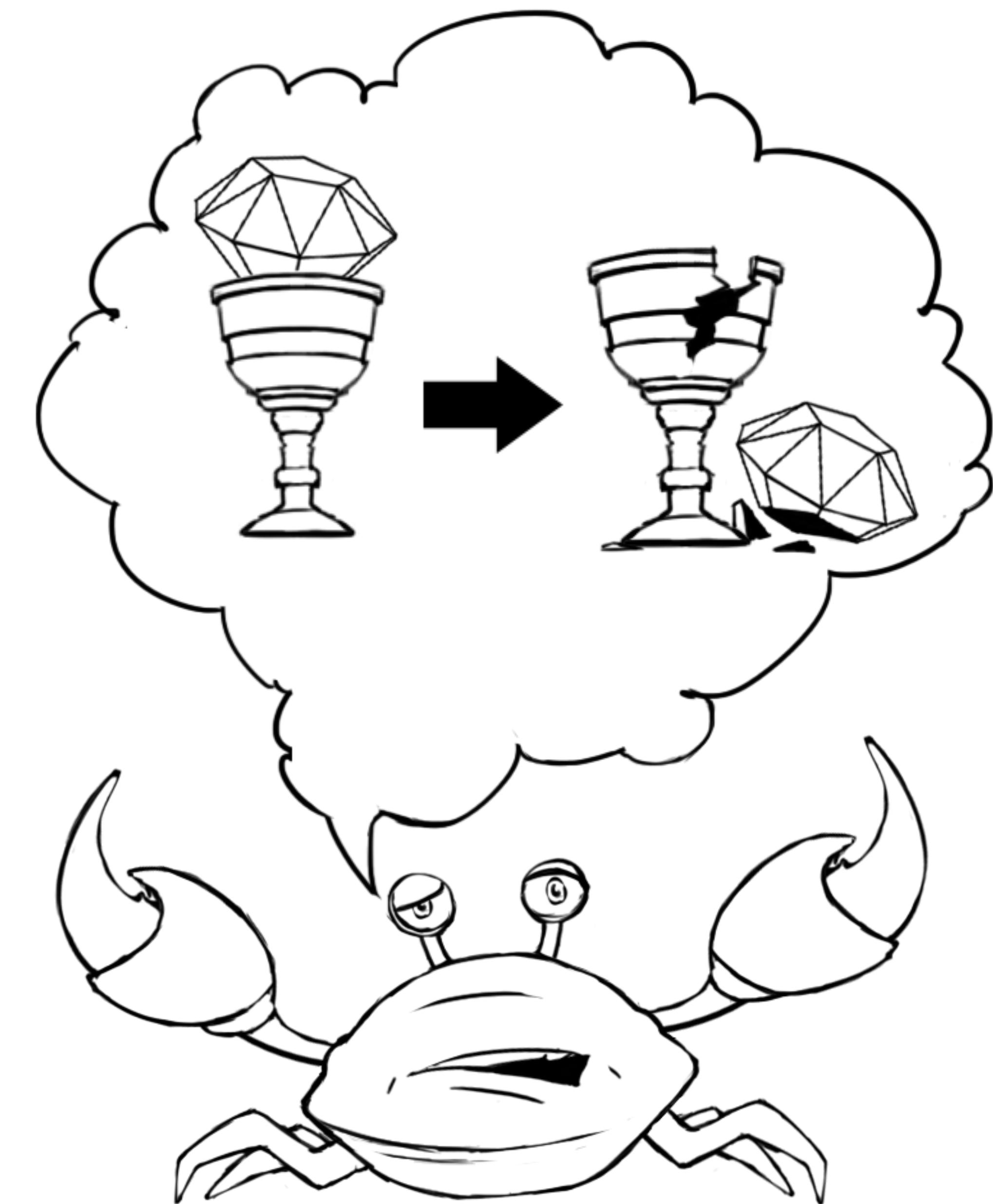


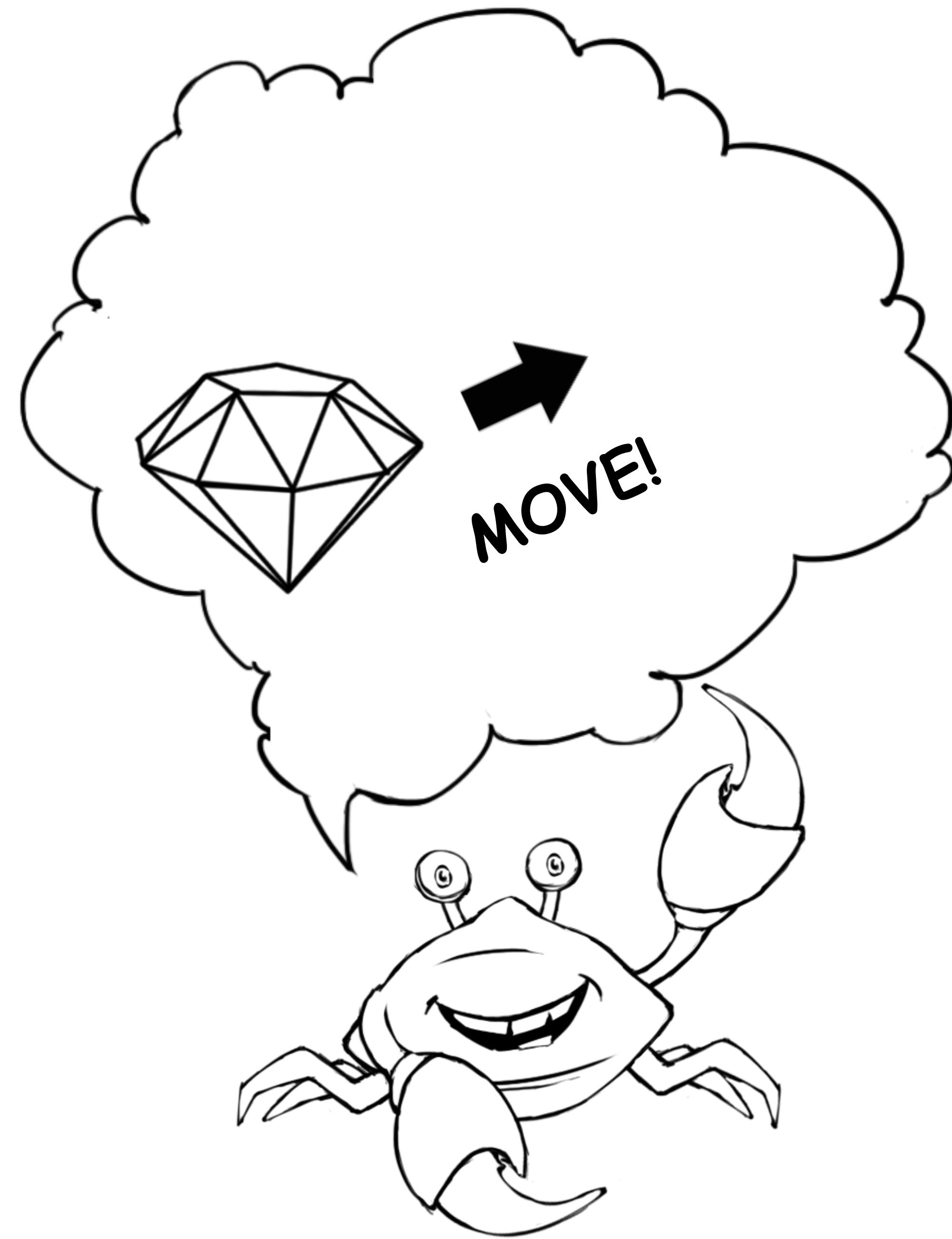


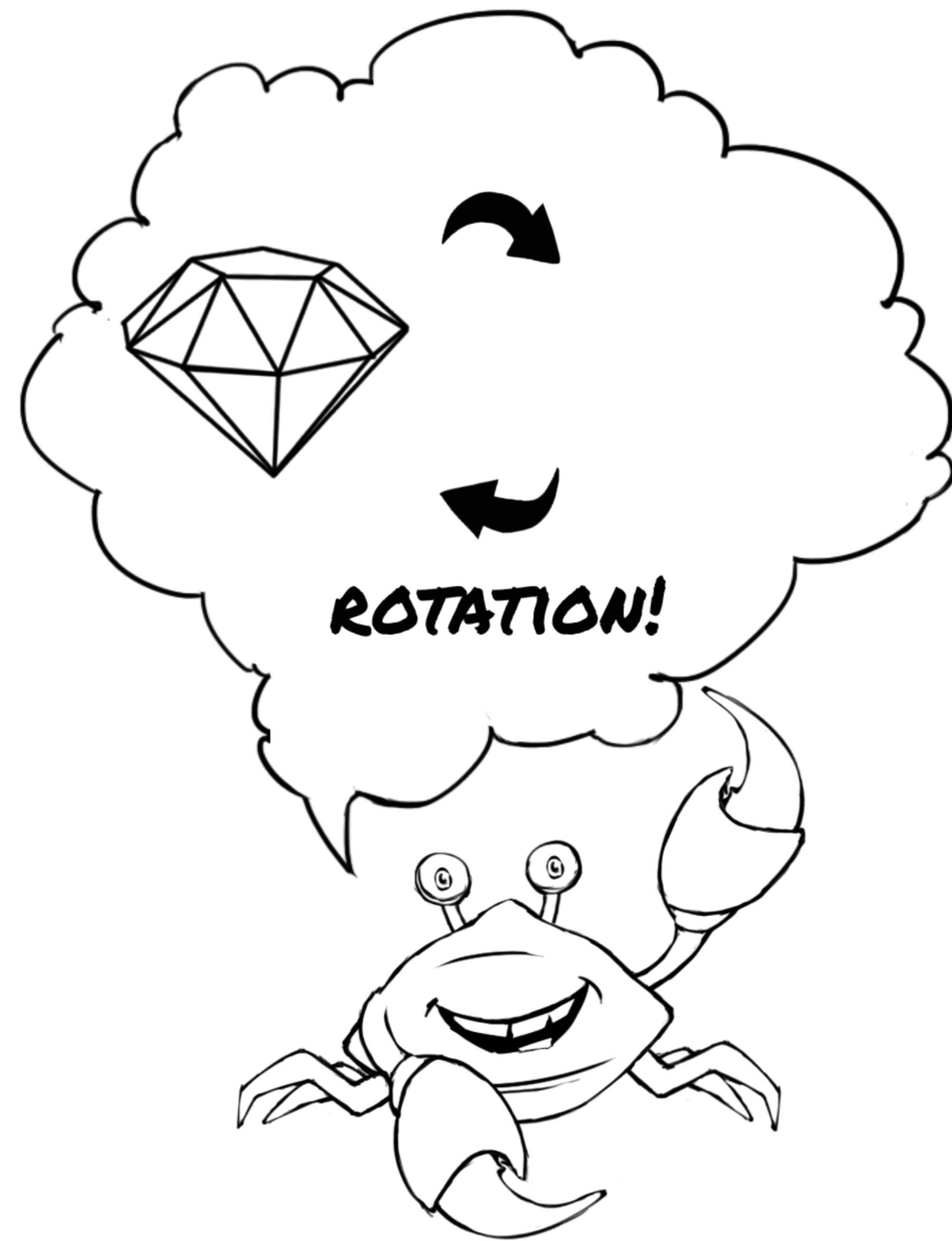


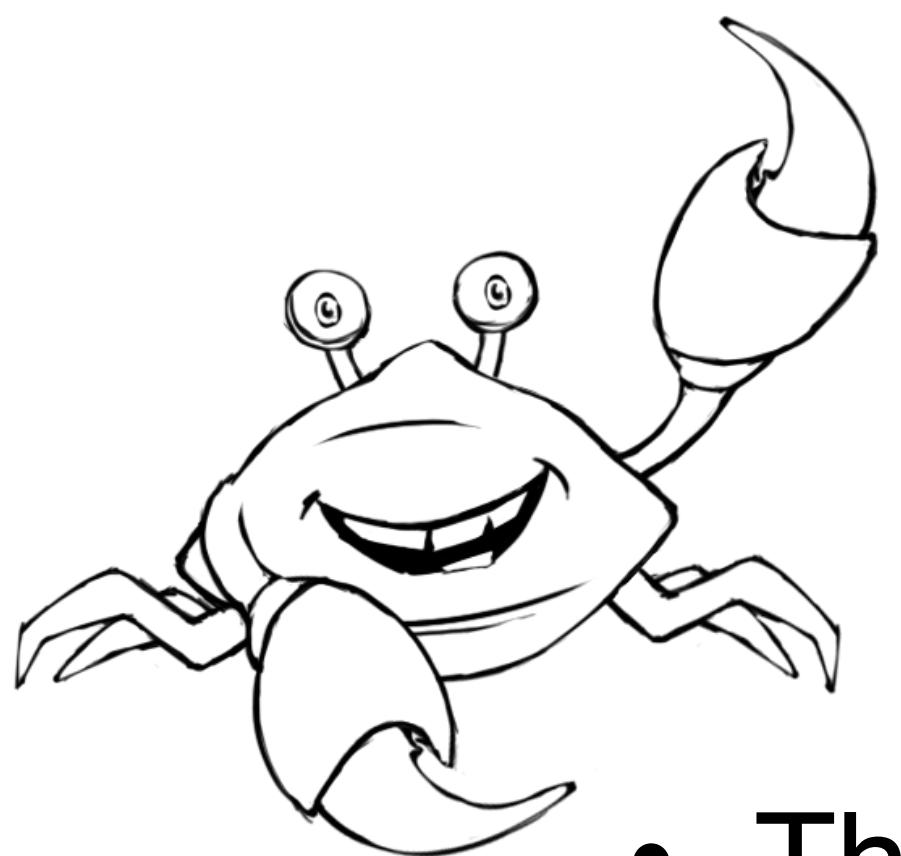










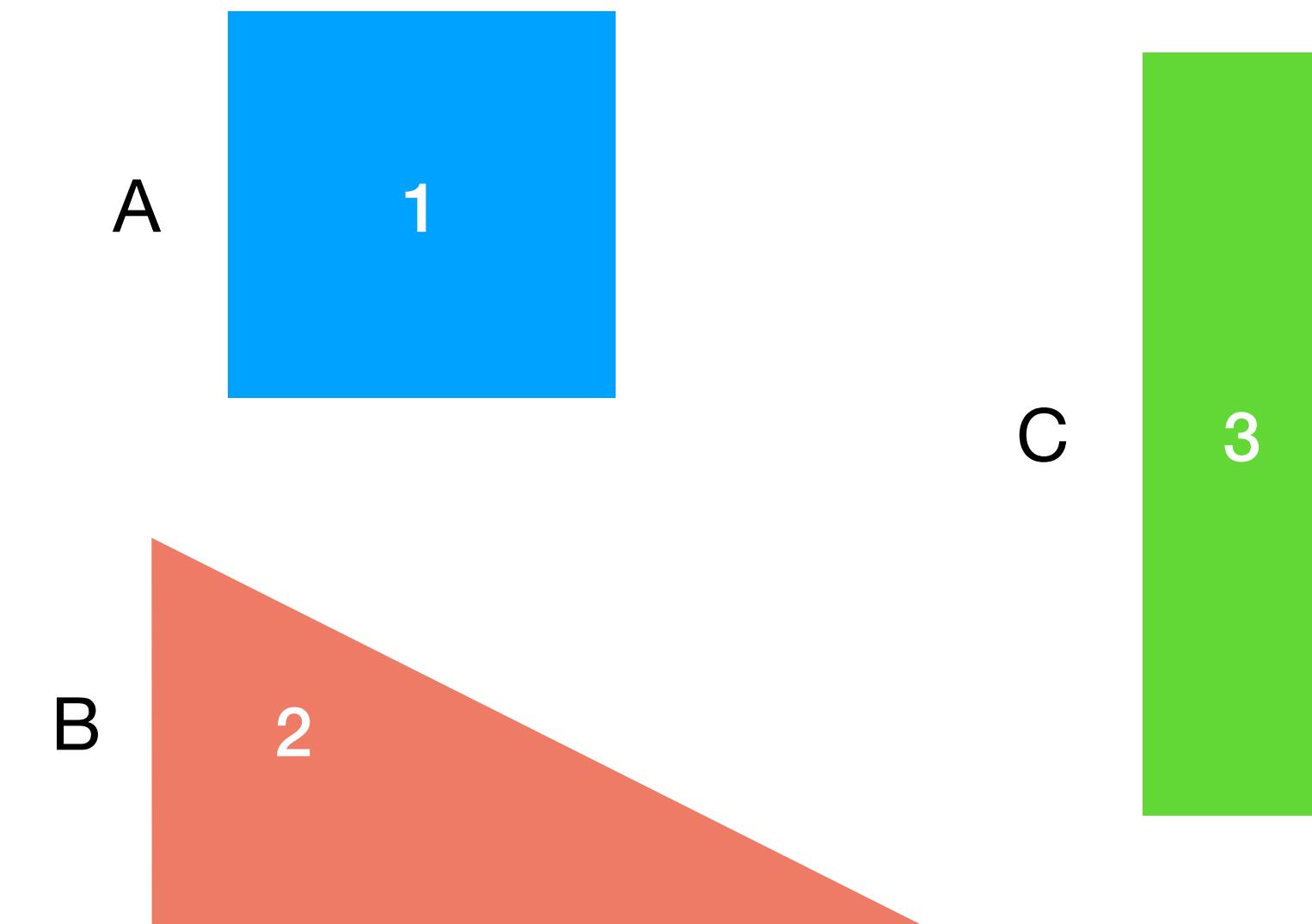
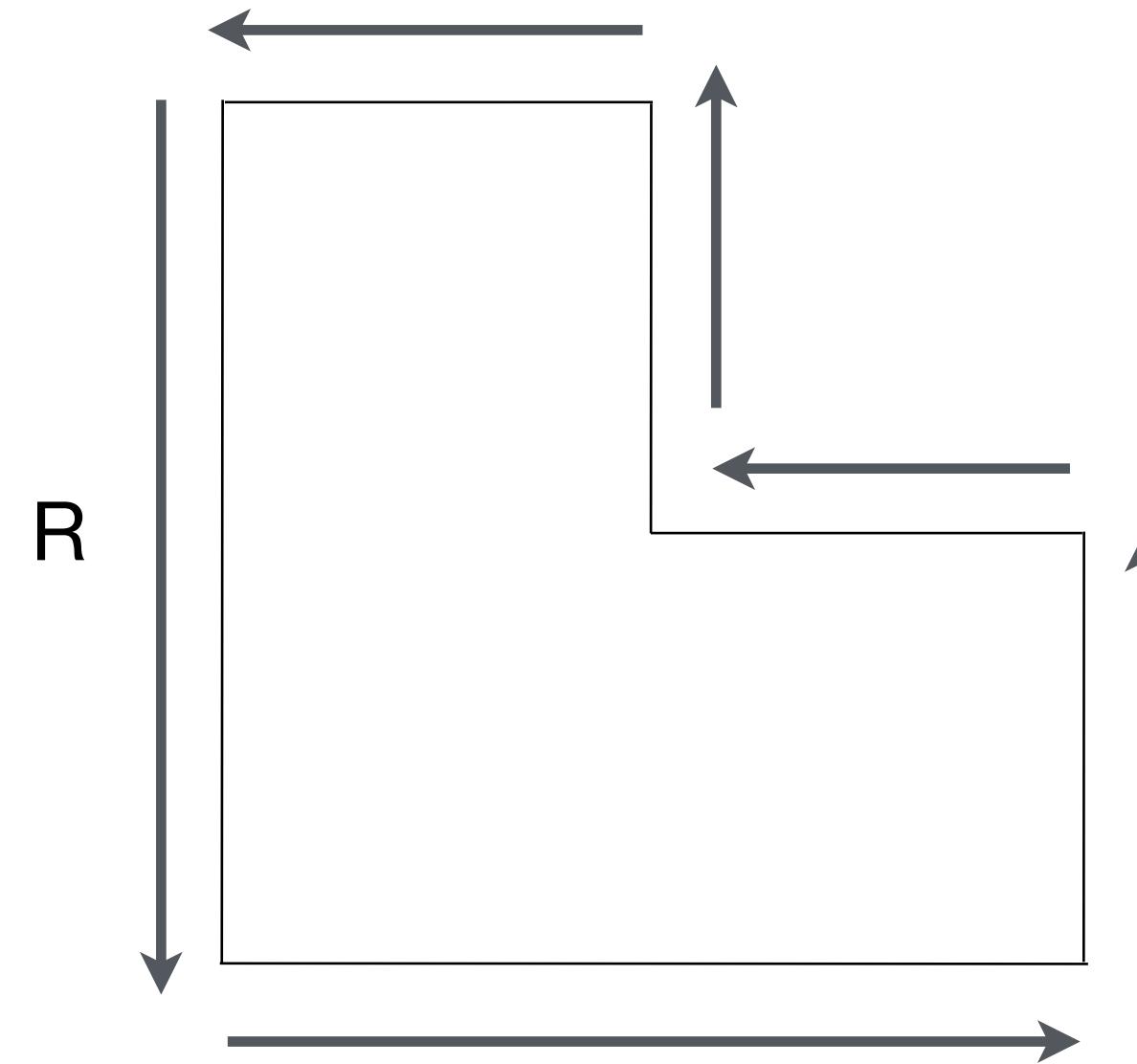


Rules of the Game

- The problem: pack Torpe's belongings into a cave (2D)
- Requirements:
 - No overlapping, all within the room, at least 30% covered
 - Try to find the best (maximal cost)
- Available actions:
 - Moving the furniture
 - Rotating the furniture

Data Structures

- Representing the cave
- Representing the furniture items
- Encoding the item costs
- Encoding the solutions



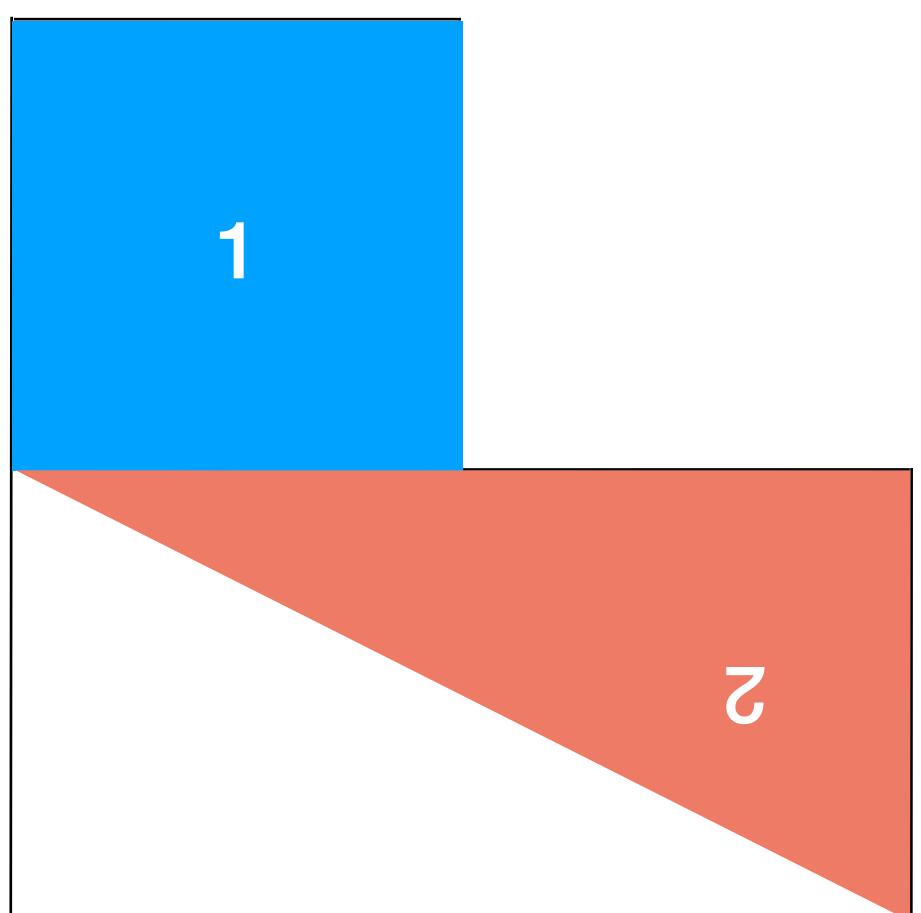
```
1: (0,0), (2,0), (2,1), (1,1), (1,2), (0,2) # 1:(0,0), (1,0), (1,1), (0,1); 2:(0,0), (2,0), (0,1); 3:(0,0), (0.5,0), (0.5,2), (0,2)
```

R

A

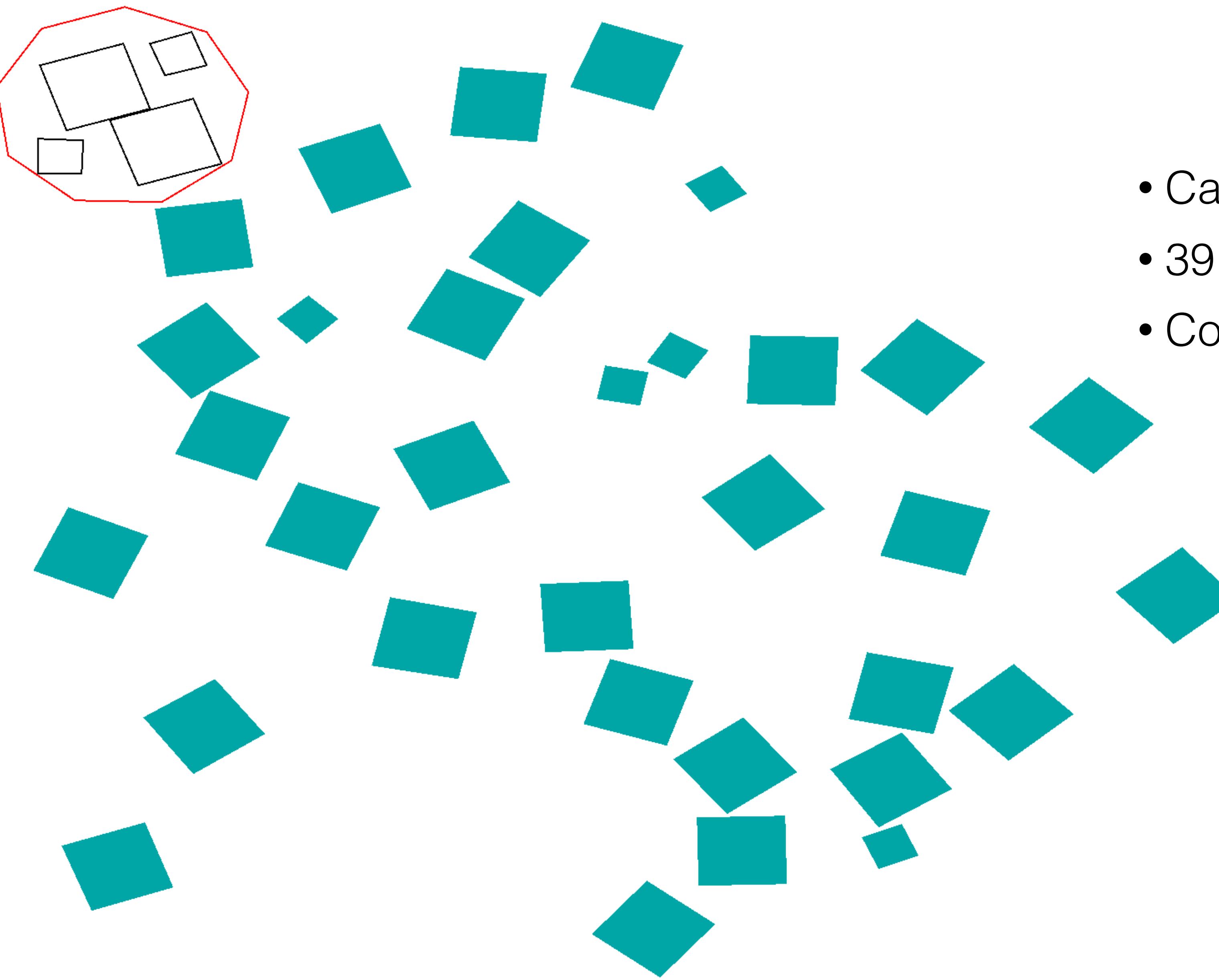
B

C

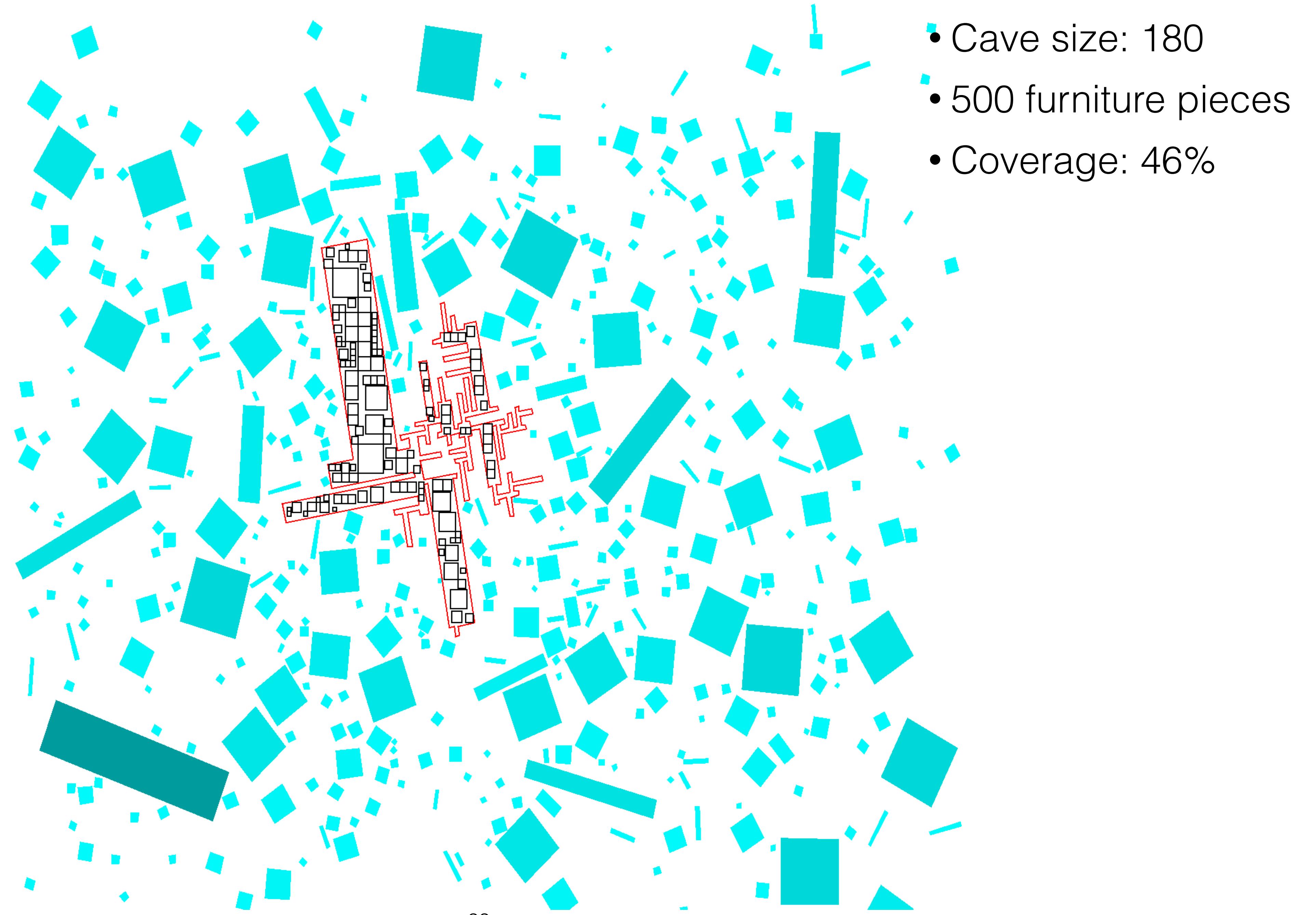


Checking a Solution

- What is an acceptable solution?
- How to check it using the data types we already have?



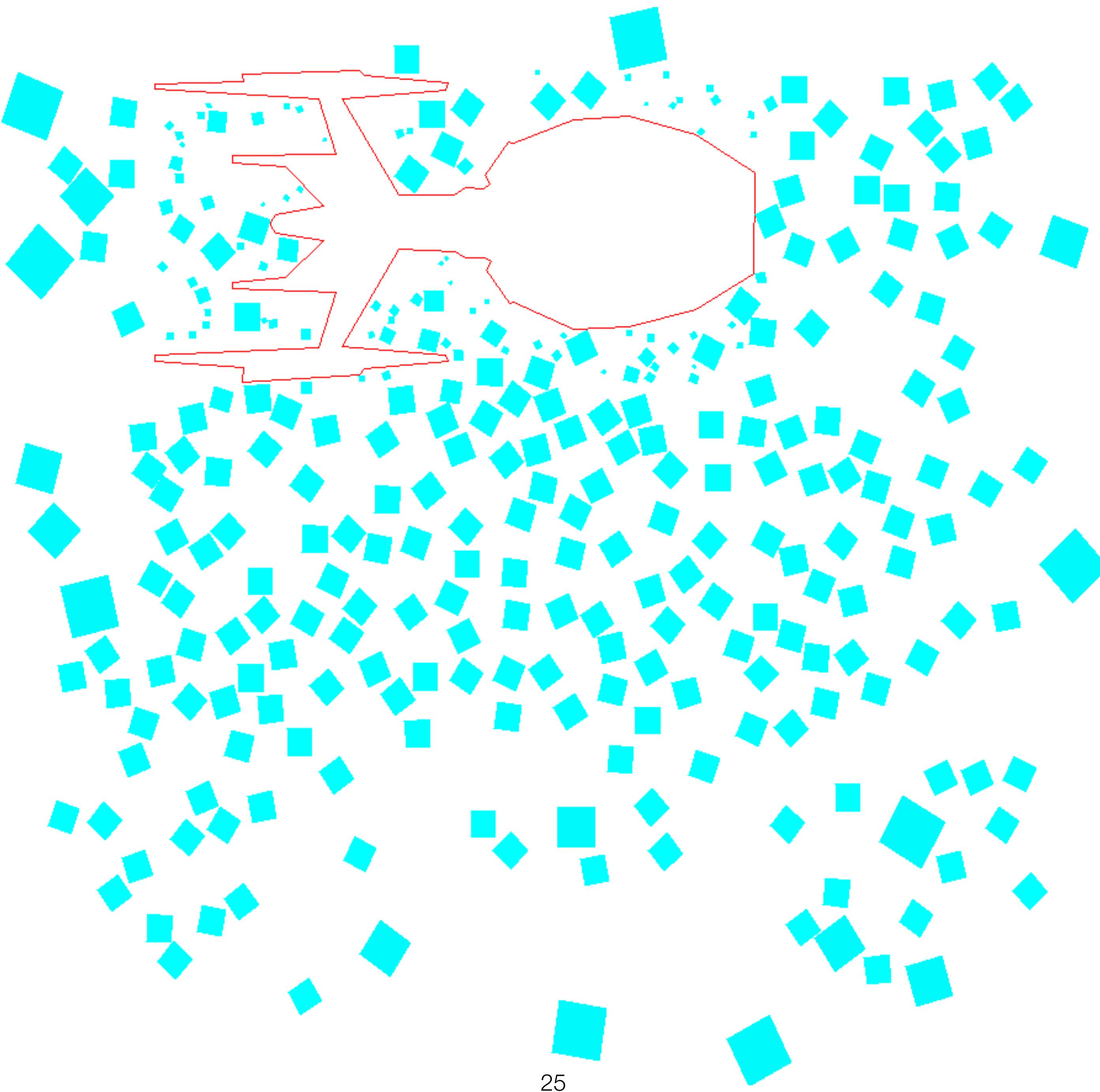
- Cave size: 9
- 39 furniture pieces
- Coverage: 40%



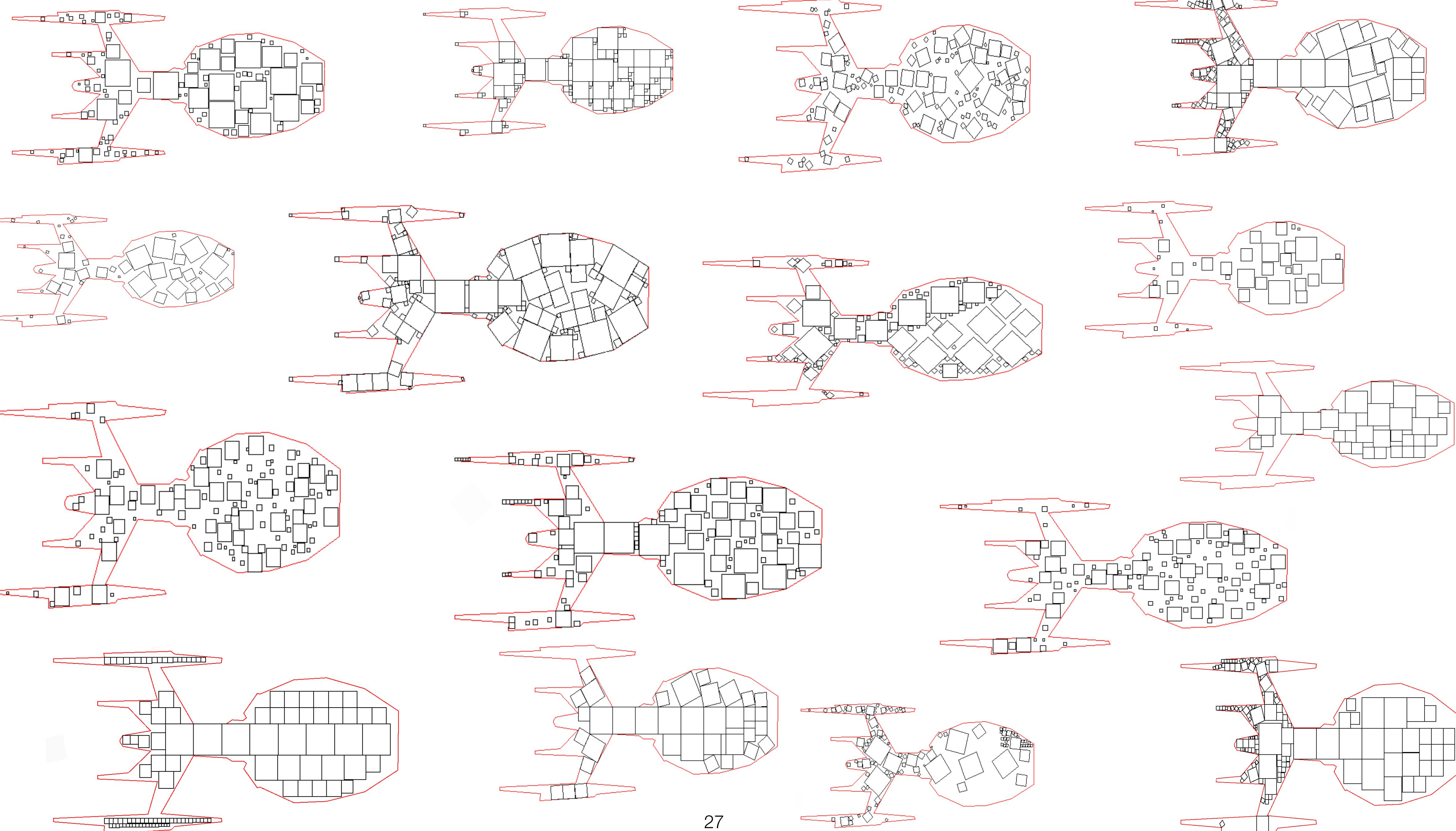
Towards an Algorithm

- What are the main steps?
- How to produce an acceptable solution?
- When should we stop?

Some solutions

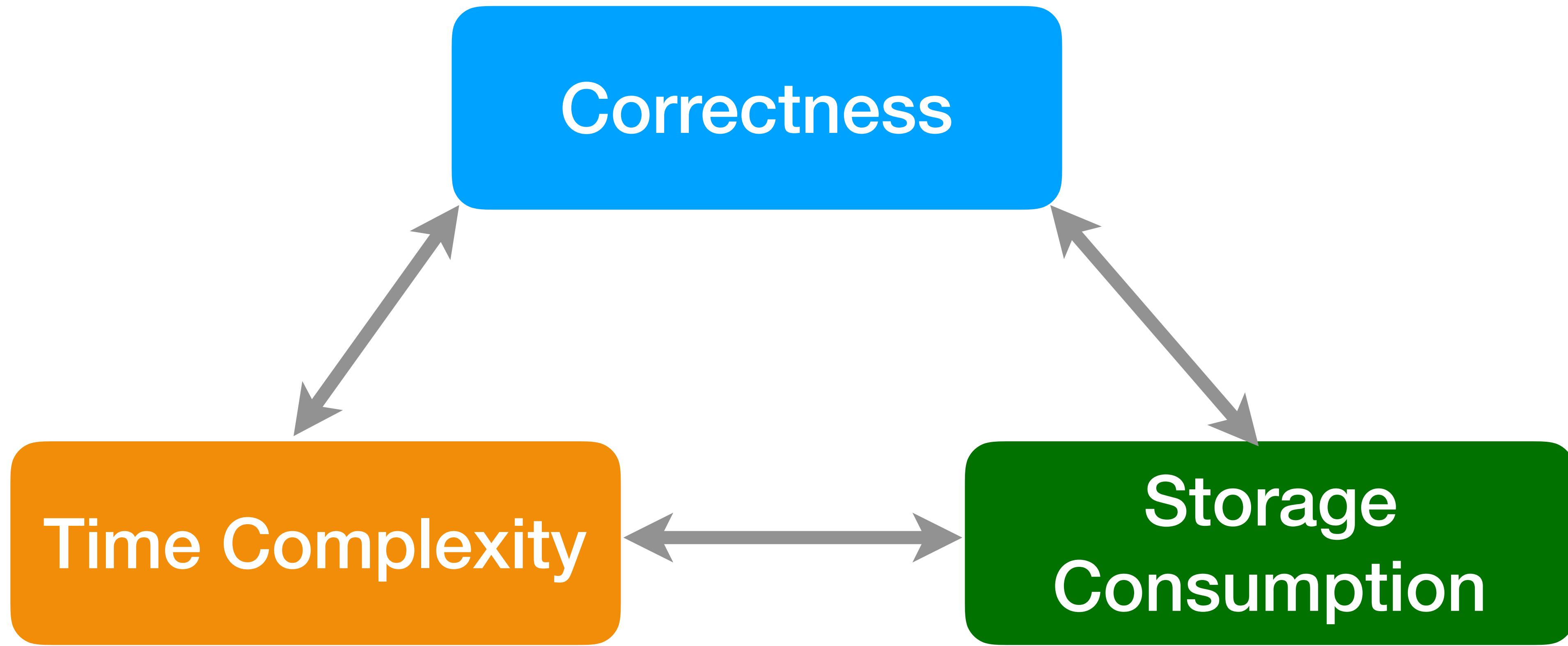






What else this course is about

Analysis of Algorithms



Algorithmic problems and Time Complexity

- **tractable problems** — admit solutions that run in “reasonable” time (e.g., sorting, searching, compression/decompression)
- **possibly intractable** — probably don’t have reasonable-time algorithmic solutions (e.g., SAT, graph isomorphism)
- **practically intractable** — definitely don’t have such solutions (e.g. the Towers of Hanoi)
- **non-computable** — can’t be solved algorithmically at all (e.g., the halting problem)

Aspects that we will study

- Algorithm Correctness
- Algorithm Termination
- Time complexity
 - Worst case
 - Average case
 - Best/Worst case

Aspects that we will study

- Algorithm Correctness — *Does my algorithm really do what it's supposed to do?*
- Algorithm Termination — *Does my algorithm always complete its work?*
- Time complexity — *How slow is my algorithm...*
 - Worst case — ... *in the worst possible case?*
 - Average case — ... *in an average case?*
 - Best/Worst case — ... *if I do my best to optimise it?*

Example: Determinant of a matrix

Laplace expansion: $|M| = \sum_{i=1}^n (-1)^{i-1} M_{1,i} \overbrace{|M^{1,i}|}^{\substack{M's \text{ element} \\ \text{at row 1,} \\ \text{column } i}} \quad (1,i)\text{-minor of } M$

For a 3x3 matrix:
$$\begin{vmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{vmatrix} =$$

$$a_{11} \begin{vmatrix} a_{22} & a_{23} \\ a_{32} & a_{33} \end{vmatrix} - a_{12} \begin{vmatrix} a_{21} & a_{23} \\ a_{31} & a_{33} \end{vmatrix} + a_{13} \begin{vmatrix} a_{21} & a_{22} \\ a_{31} & a_{32} \end{vmatrix} =$$

$$a_{11}(a_{22} \cdot a_{33} - a_{23} \cdot a_{32}) - a_{12}(a_{21} \cdot a_{33} - a_{23} \cdot a_{31}) + a_{13}(a_{21} \cdot a_{32} - a_{22} \cdot a_{31})$$

Example: Determinant of a matrix

Laplace expansion: $|M| = \sum_{i=1}^n (-1)^{i-1} M_{1,i} |M^{1,i}|$

(in Haskell)

```
detLaplace :: Num a => Matrix a -> a

detLaplace m
| size m == 1 = m ! (1,1)
| otherwise    =
  sum [ (-1)^(i-1) * m ! (1,i) * detLaplace (minorMatrix 1 i m) |
        i <- [1 .. ncols m] ]
```

(demo)

Triangular matrices

$$L = \begin{pmatrix} a_{1,1} & 0 & \cdots & 0 \\ a_{2,1} & a_{2,2} & 0 & 0 \\ \vdots & \vdots & \ddots & \vdots \\ a_{n,1} & a_{n,2} & \cdots & a_{n,n} \end{pmatrix} \quad U = \begin{pmatrix} a_{1,1} & a_{1,2} & \cdots & a_{1,n} \\ 0 & a_{2,2} & \cdots & a_{2,n} \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \cdots & a_{n,n} \end{pmatrix}$$

Determinant of a triangular matrix
is a *product* of its diagonal elements.

For a 3x3 matrix:

$$\begin{vmatrix} a_{11} & a_{12} & a_{13} \\ 0 & a_{22} & a_{23} \\ 0 & 0 & a_{33} \end{vmatrix} =$$

$$a_{11}(a_{22} \cdot a_{33} - \cancel{a_{23} \cdot 0}) - a_{12}(\cancel{0 \cdot a_{33}} - \cancel{a_{23} \cdot 0}) + a_{13}(\cancel{0 \cdot a_{32}} - \cancel{a_{22} \cdot 0}) = a_{11} \cdot a_{22} \cdot a_{31}$$

Determinants via LU-decomposition

LU-decomposition: any square matrix M , such that its top-left element is non-zero can be represented in a form

$$M = LU$$

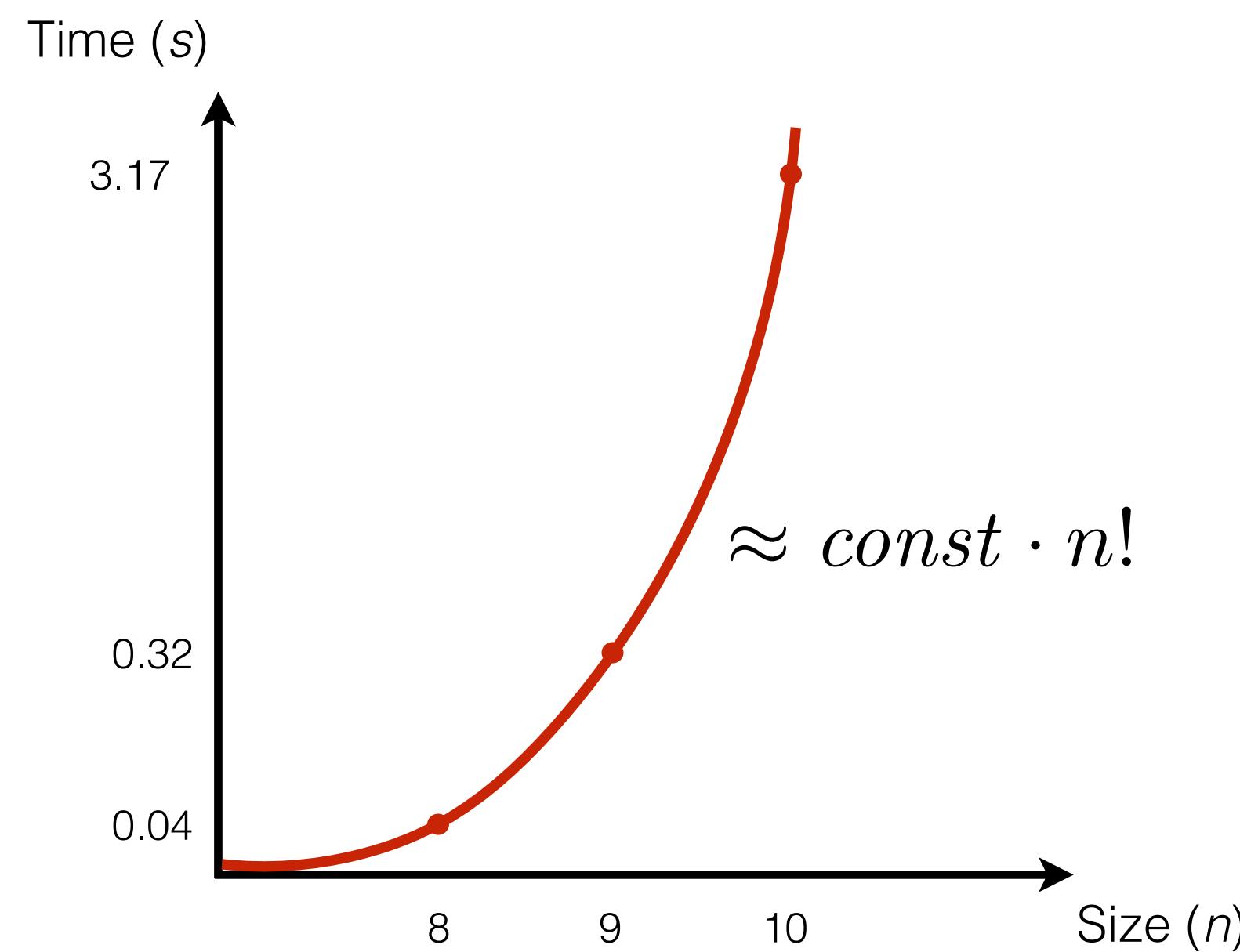
where L and U are lower- and upper-triangular matrices.

Therefore, $|M| = |L| \cdot |U|$

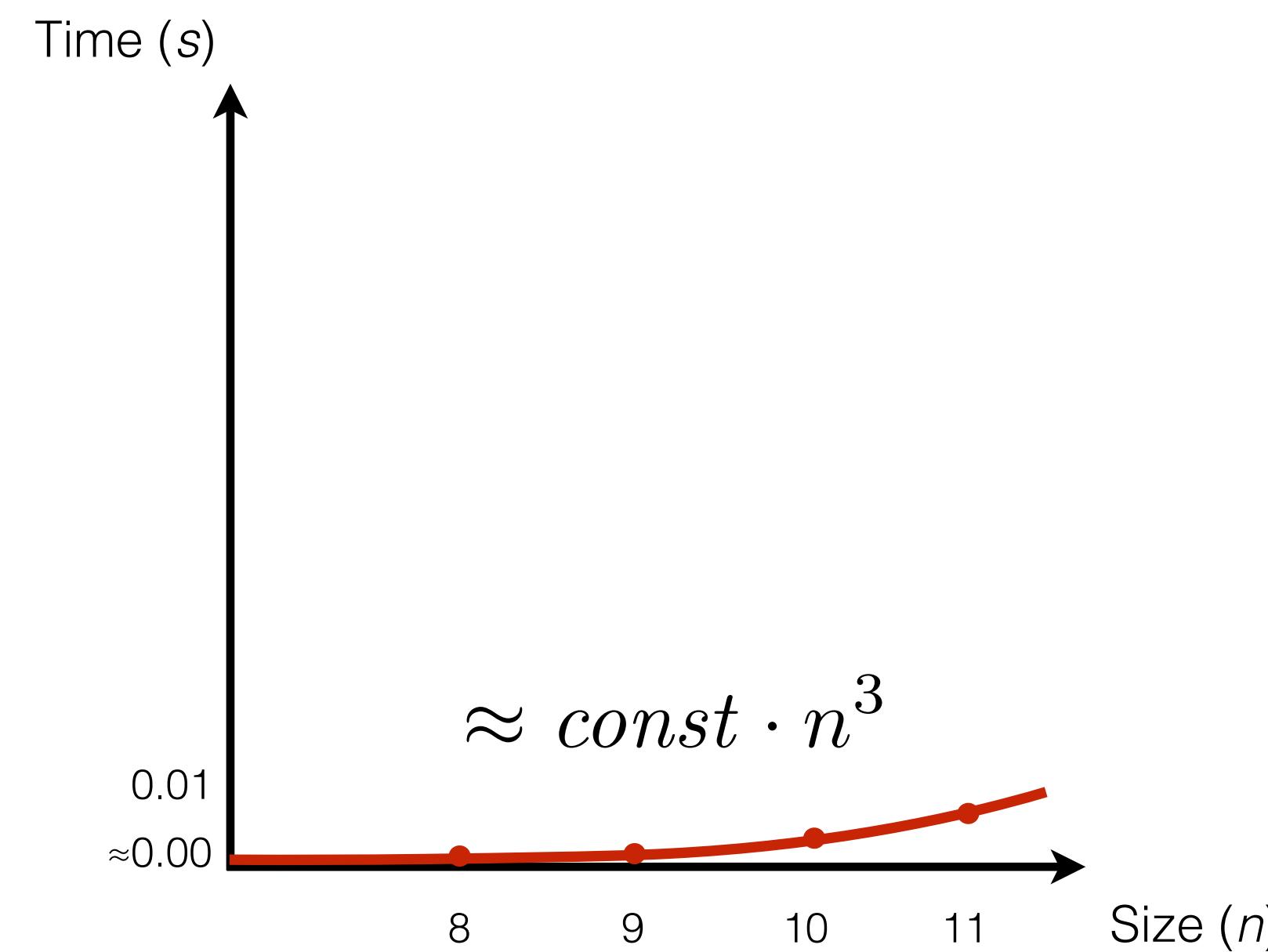
```
detLU :: Num a => Matrix a -> a
detLU m = case luDecomp m of
    (l, u) -> diagProd l * diagProd u
```

(demo)

Running time as a function of size



Determinant via
Laplace expansion



Determinant via
LU-decomposition

Time demand depends on problem size

Function	Problem size			
	10	10^2	10^3	10^4
$\log_2 n$	3.3	6.6	10	13.3
n	10	100	1000	10^4
$n \log_2 n$	33	700	10^4	1.3×10^5
n^2	100	10^4	10^6	10^8
n^3	1000	10^6	10^9	10^{12}
2^n	1024	1.3×10^{30}	$> 10^{100}$	$> 10^{100}$
$n!$	3×10^6	$> 10^{100}$	$> 10^{100}$	$> 10^{100}$

“Sizes” of different problems

Problem	Input size, n
sorting	number of items to be sorted
searching	size of the set to query
determinant calculation	number of rows and columns in the matrix
finding a shortest path	number of “checkpoints” to choose from

Two ways to analyse algorithms

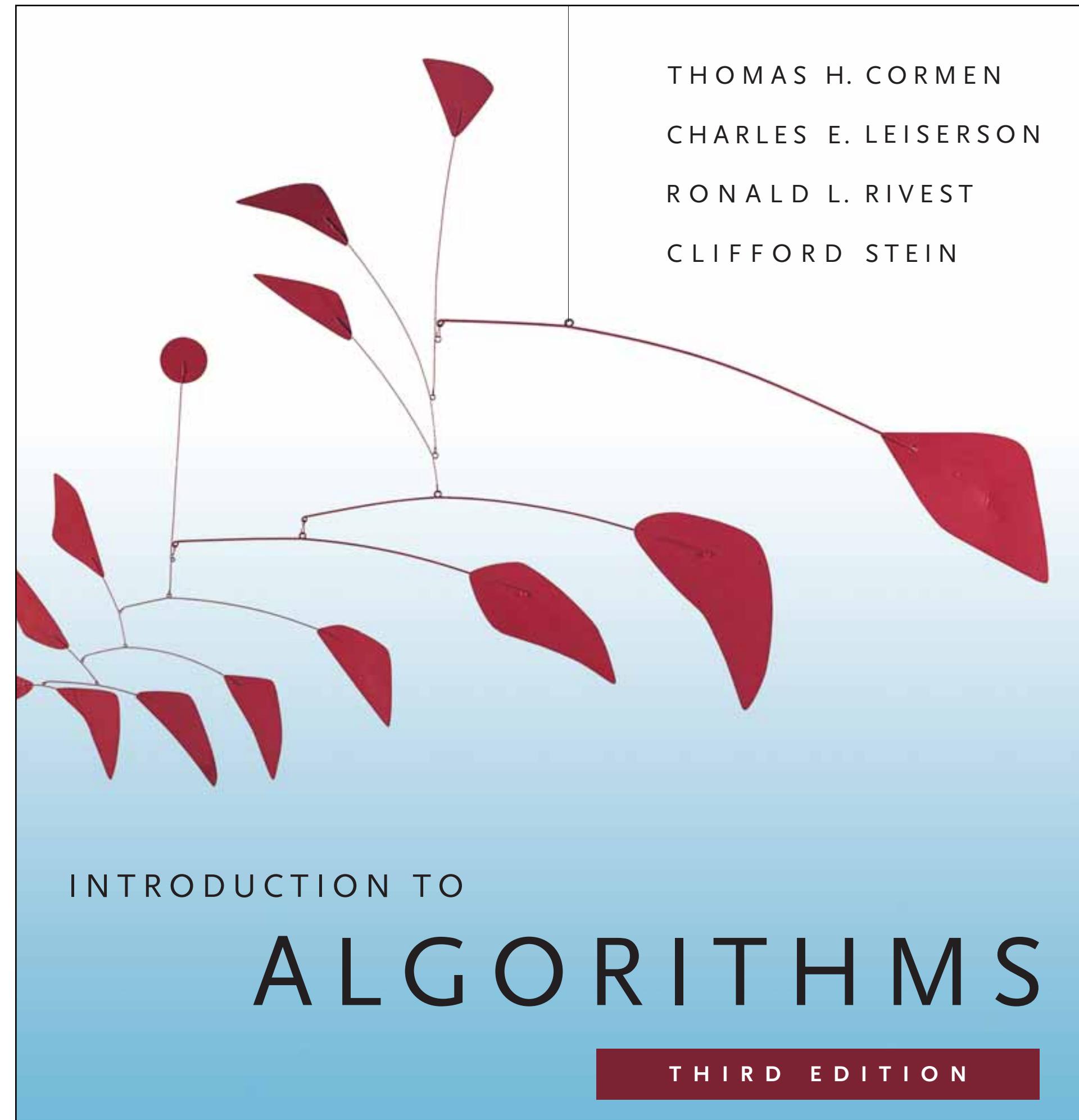
- **Empirical** — repeatedly run algorithm with different inputs to get some idea of behaviour on different inputs
 - was our selection of inputs representative?
 - this consumes the very resource (time) we are trying to conserve!
- **Theoretical** — analysis of a “paper” version of the algorithm
 - can deal with all cases (even impractically large input instances);
 - machine-independent.

Working Tools

- OCaml
- Emacs/Aquamacs
- Toolbox
 - Tuareg mode for syntax highlighting and REPL
 - Merlin mode for type information
 - Company mode auto-completion and types
 - **ocp-indent** for smart indentation
 - <https://github.com/ocaml/merlin/wiki/emacs-from-scratch>



The Textbook



Lecture Notes (WIP)

ilyasergey.net/YSC2229

Assessment

- 30% homework exercises
- 30% mid-term project
- 35% final project
- 5% class participation

Homework

- Done in groups of 3-4 people (3 is optimal)
- Deliverables:
 - an OCaml file with the solutions
 - a PDF with explanations of what has been done
- Graded out of 20
- No extra points for recommended exercises (sorry)

Before the Break

- Tell a bit about yourself:
 - Your name (optionally)
 - What is your programming background?
 - Why are you interested in Computer Science?
 - Which computing problems did you deal with?
 - What are your expectations from this course?