YSC3248: Parallel, Concurrent and Distributed Programming

Concurrent Stacks and Elimination

Last and This Lectures

- Queue
 - Bounded, blocking, lock-based
 - Unbounded, non-blocking, lock-free
- Stack
 - Unbounded, non-blocking lock-free
 - Elimination-backoff algorithm

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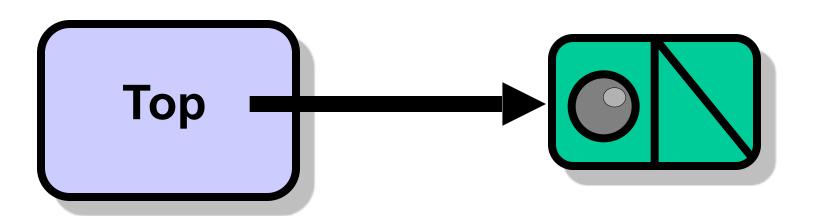
Demo:

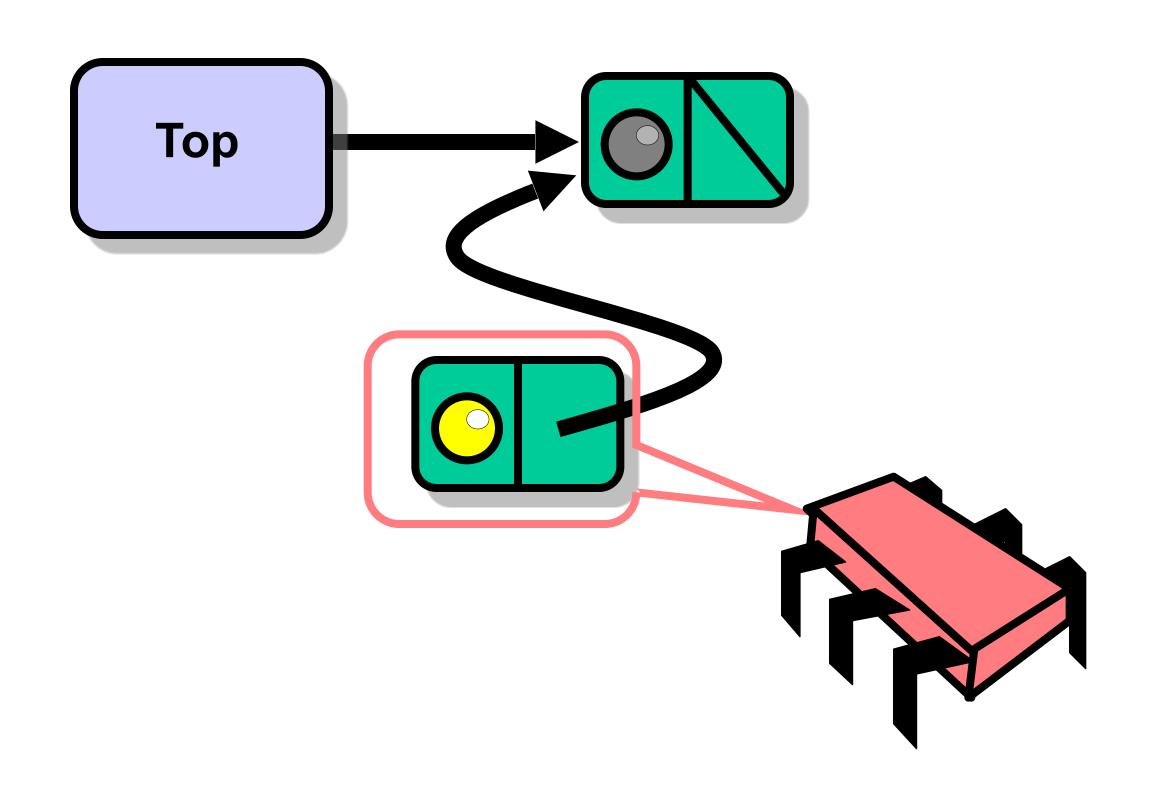
Implementing and Testing Lock-Based Stack

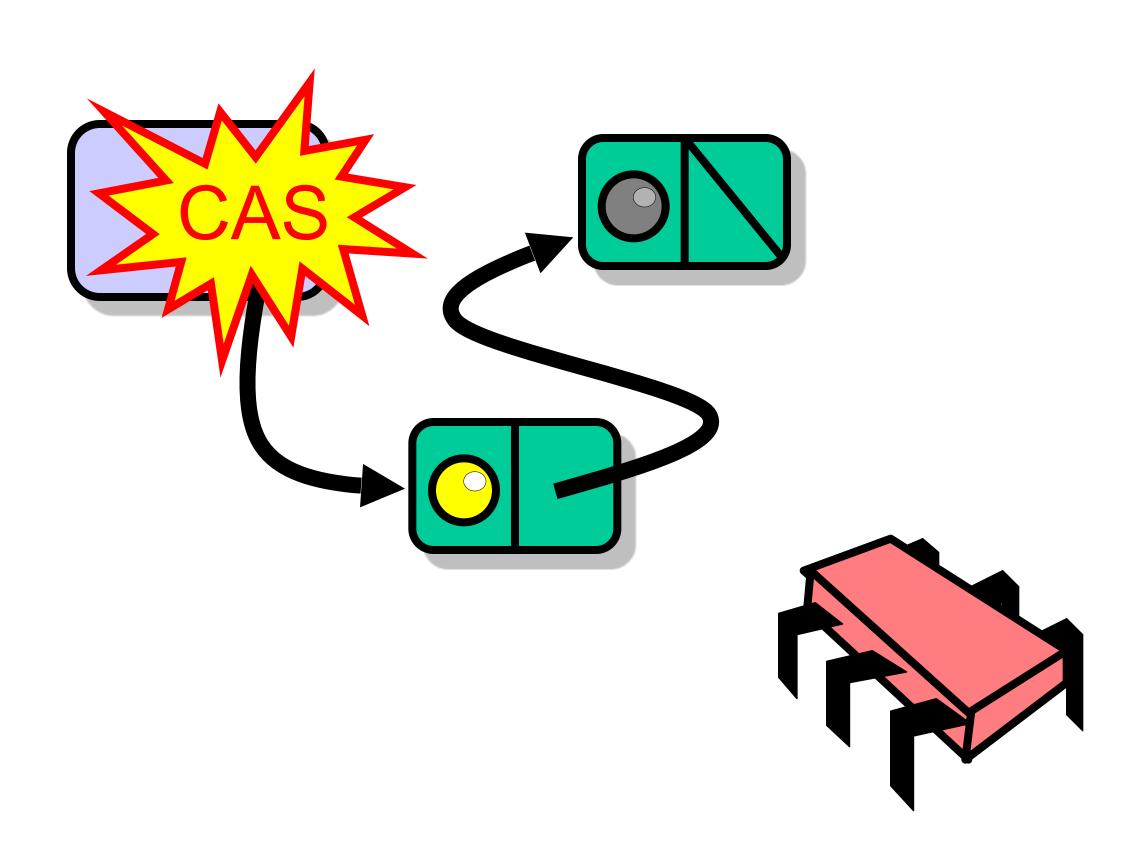
Concurrent Stack

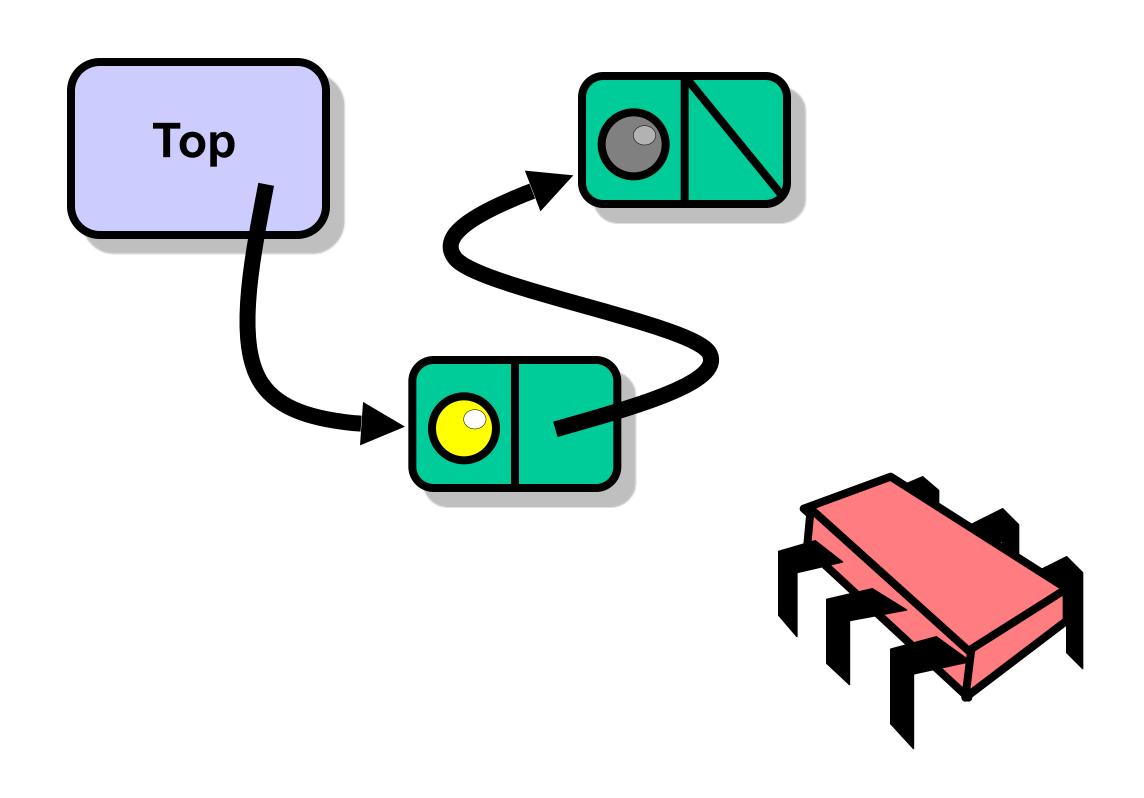
- Methods
 - push(x)
 - pop()
- Last-in, First-out (LIFO) order
- Lock-Free!

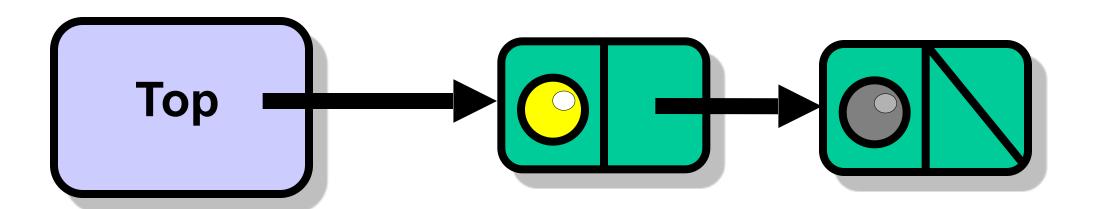
Empty Stack

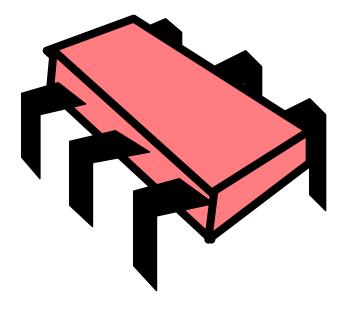


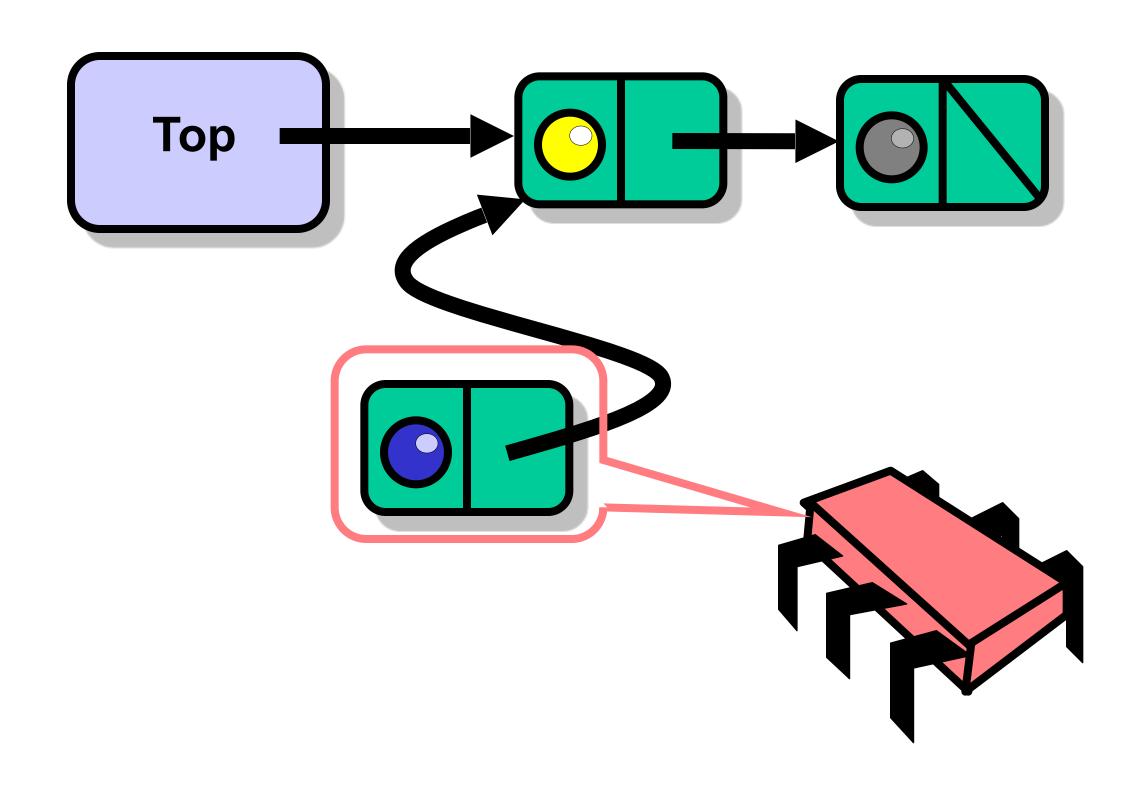


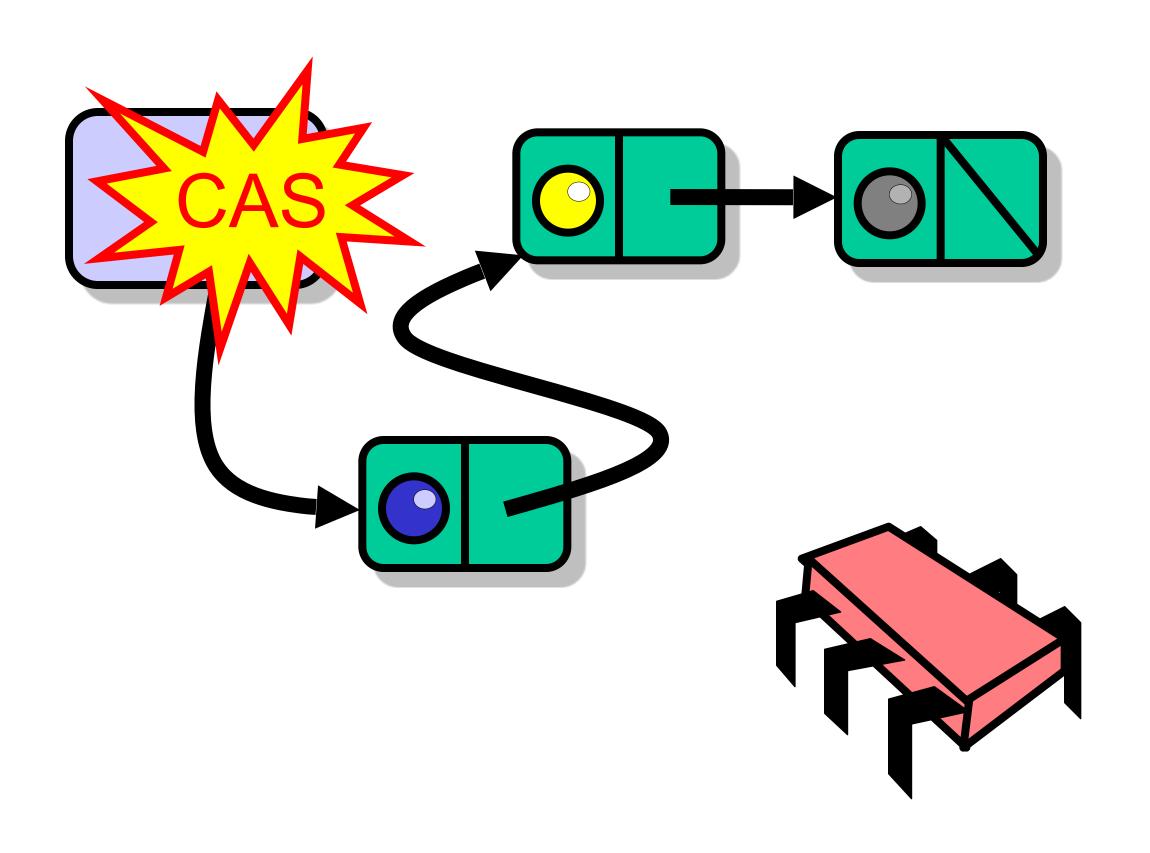


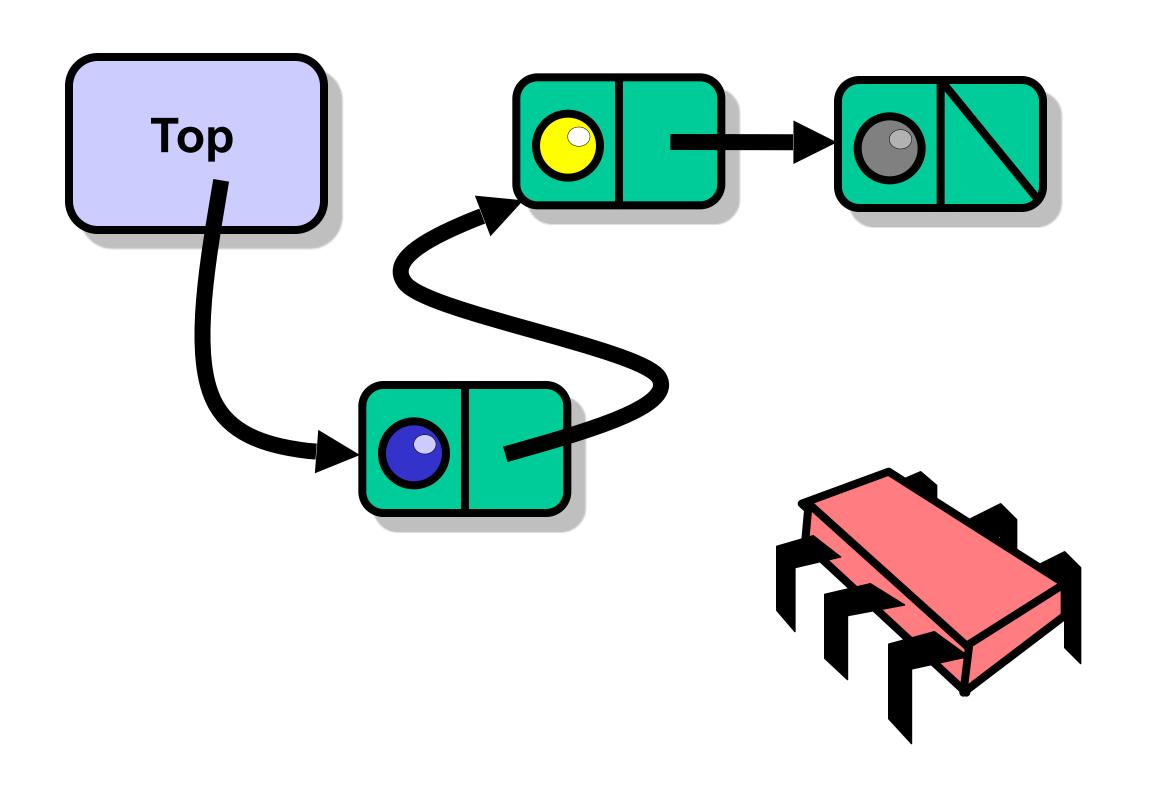


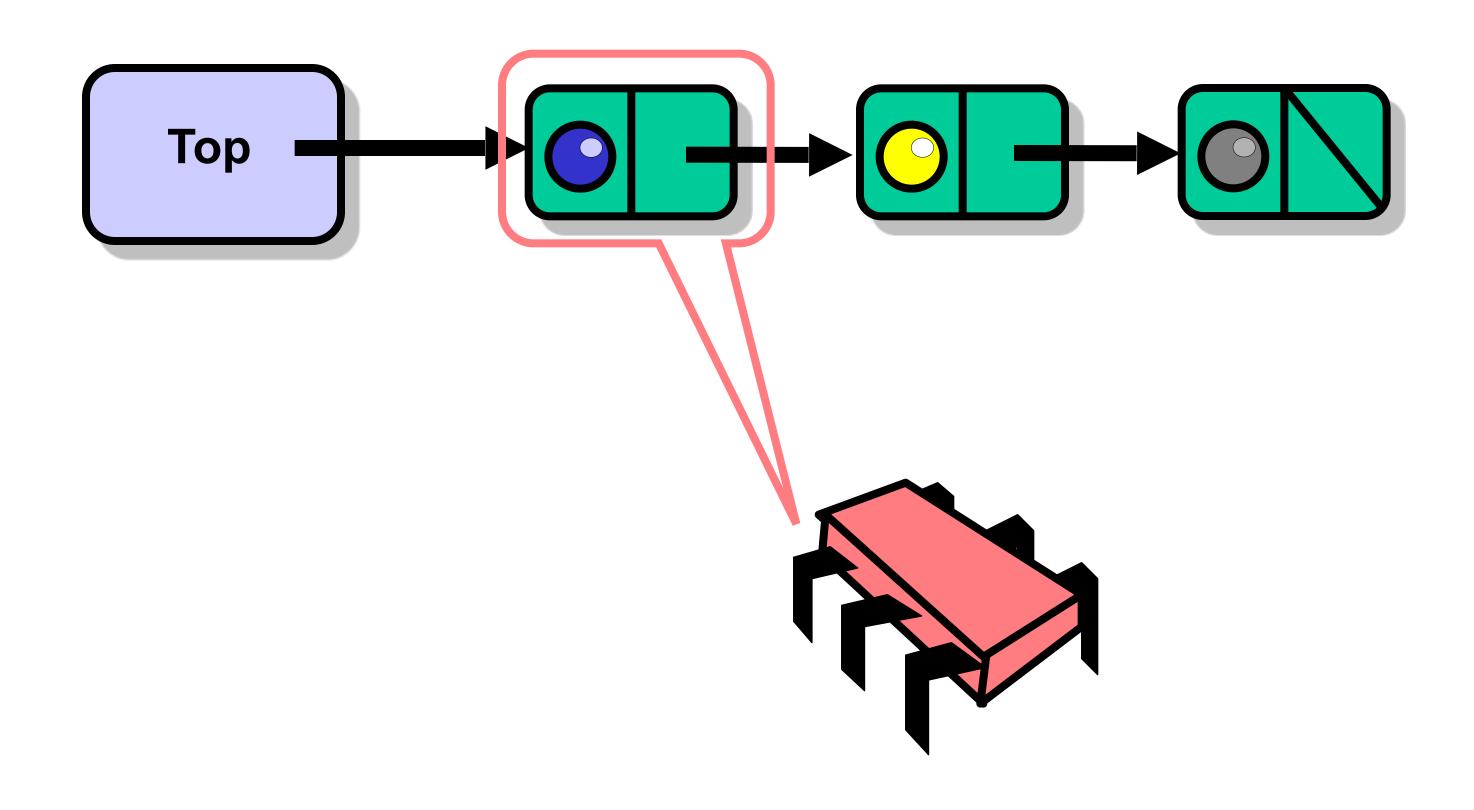


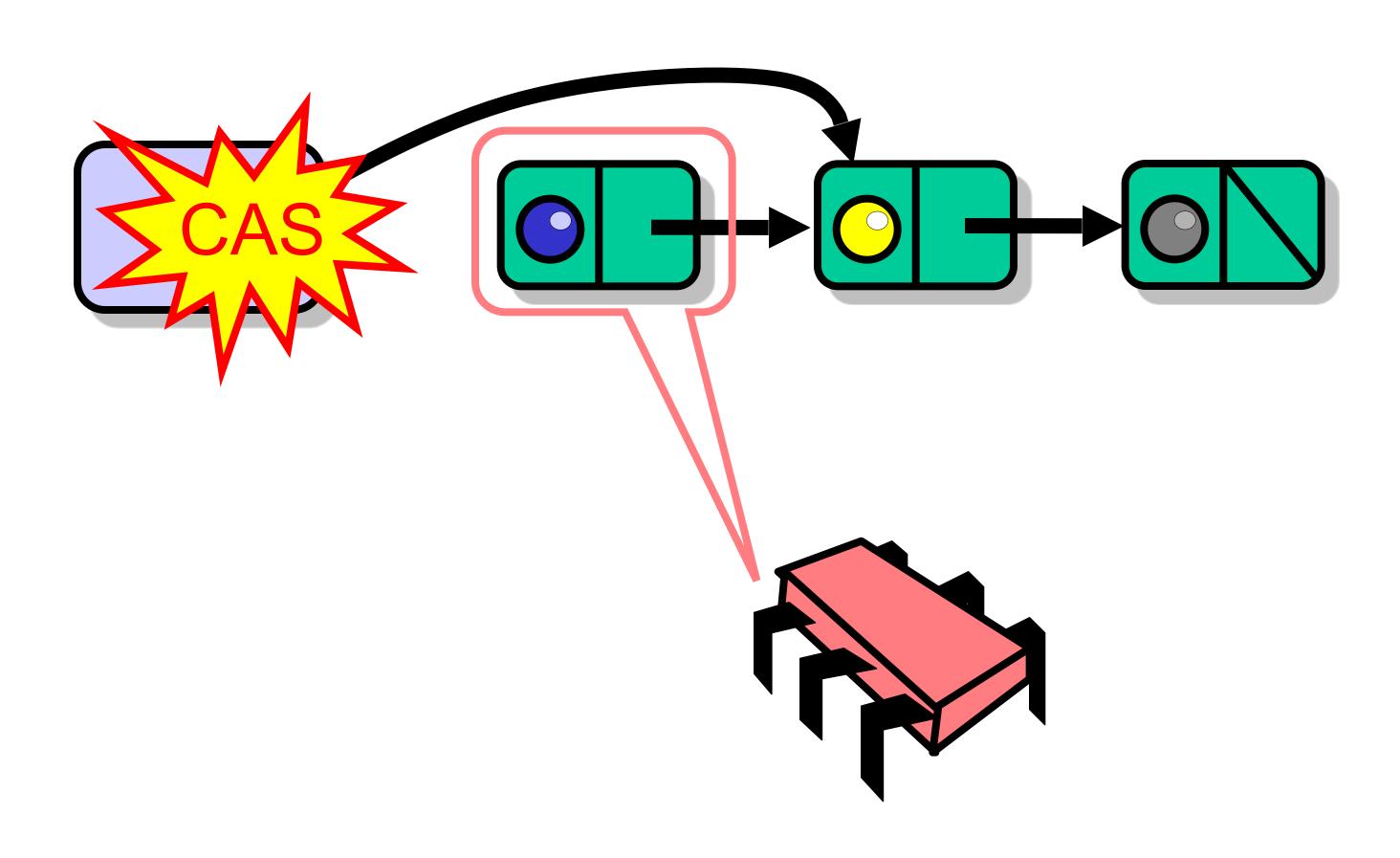


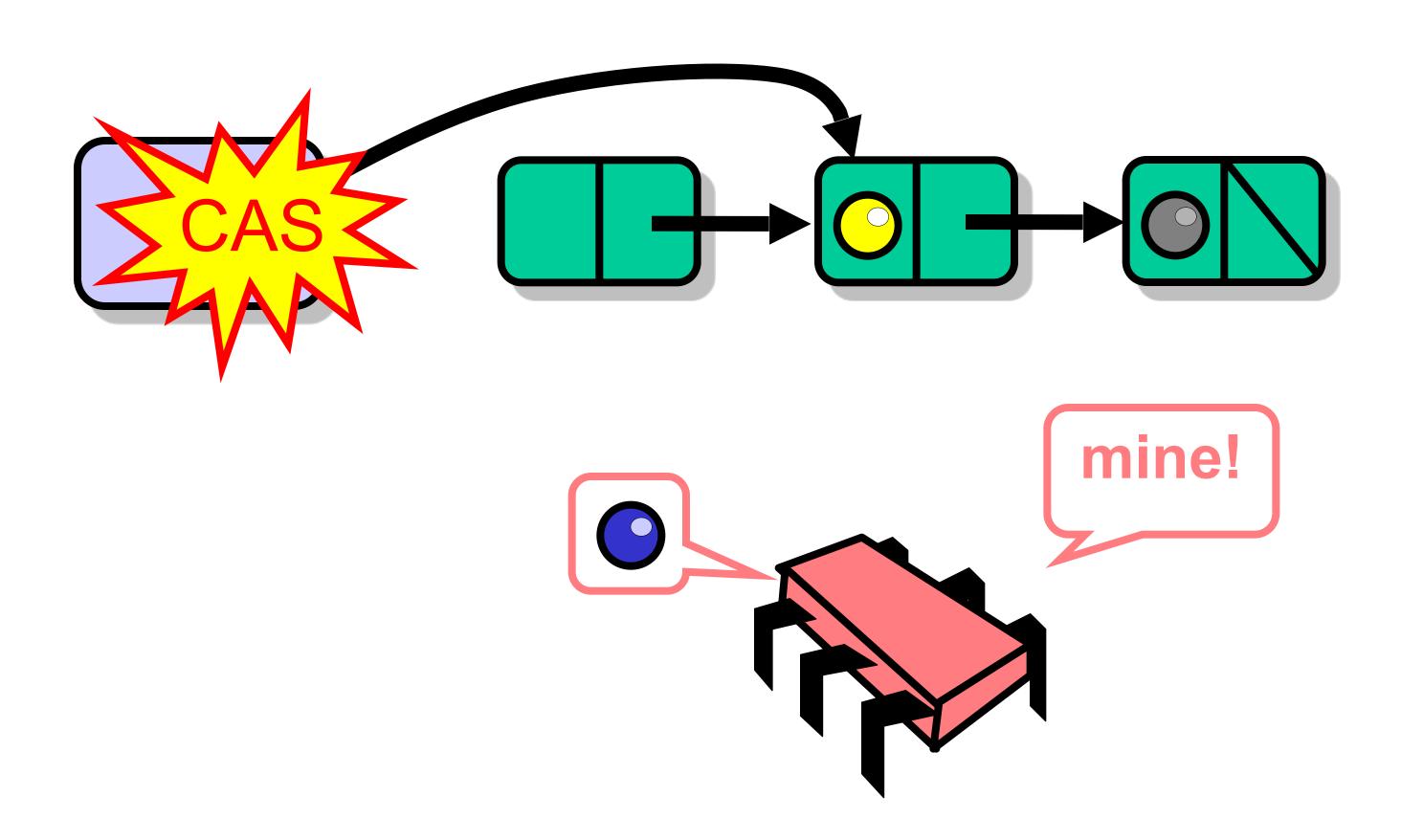


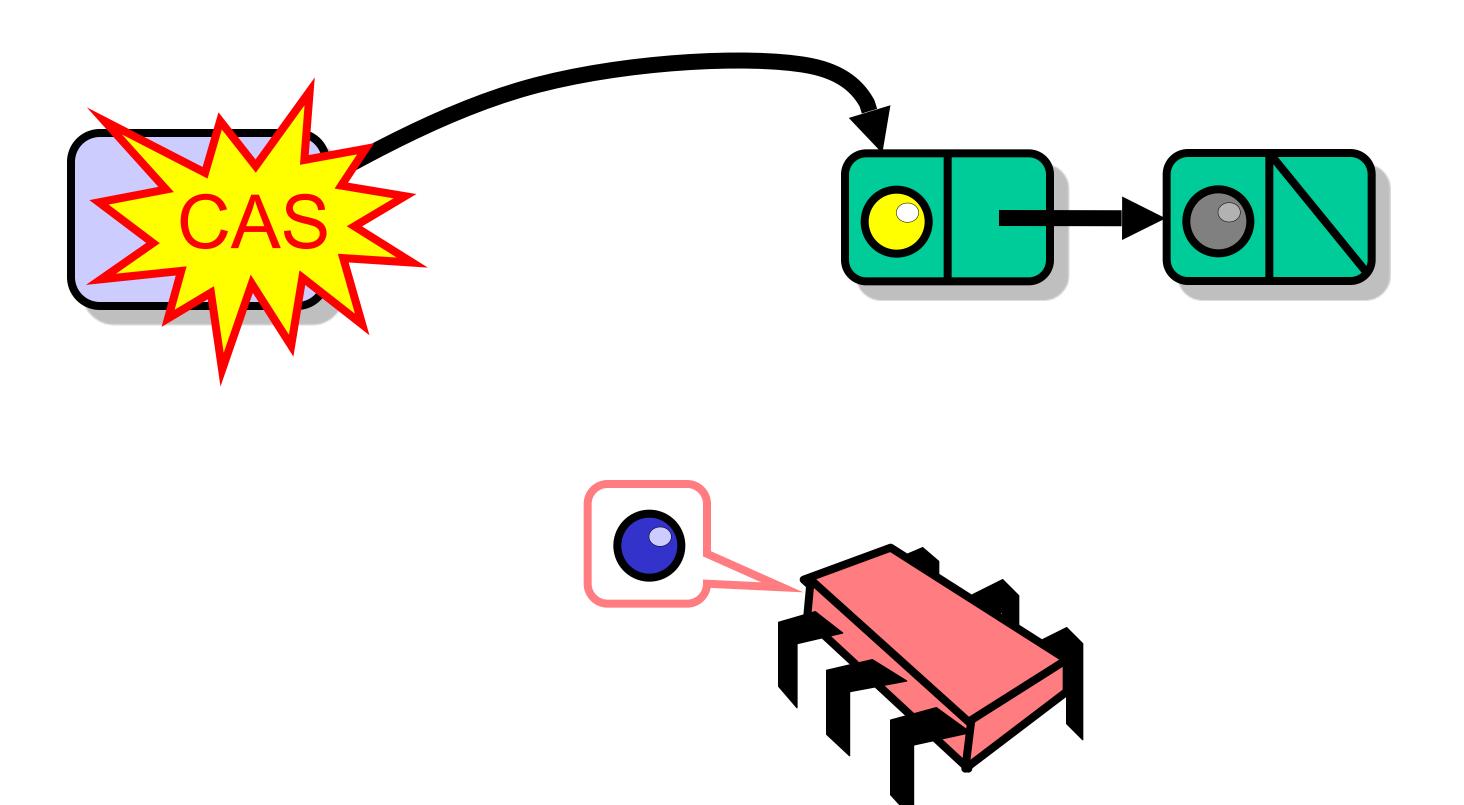


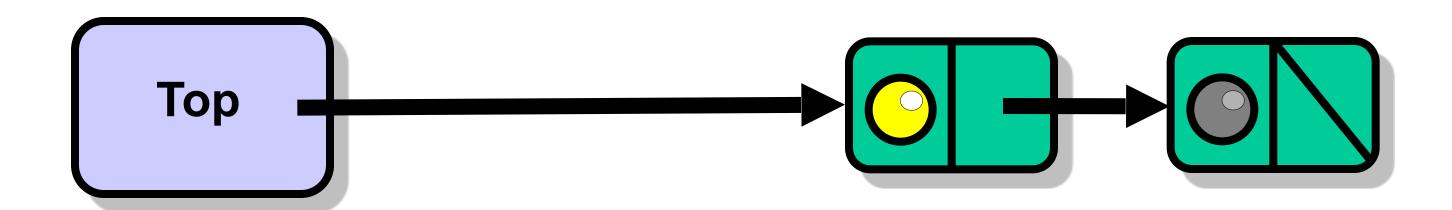


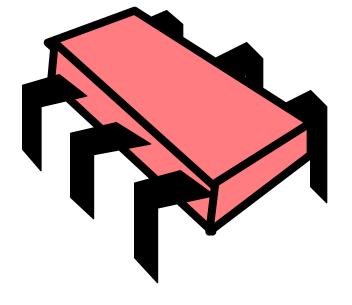












```
class LockFreeStack[T] extends ConcurrentStack[T] {
 val top = new AtomicReference[Node] (null)
 protected def tryPush(node: Node): Boolean = {
   val oldTop = top.get()
    node.next = oldTop
    top.compareAndSet(oldTop, node)
  override def push(value: T): Unit = {
   val node = new Node(value)
   while (true) {
      if (tryPush(node))
        return
      else backoff.backoff()
```

```
class LockFreeStack[T] extends ConcurrentStack[T] {
 wal ton = new AtomicReference[Nodel(null)
 protected def tryPush(node: Node): Boolean
   val oldTop = top.get(
    node.next = oldTop
    top.compareAndSet(oldTop,
 override def push (value: T):
   val node = new Node(value)
   while (true) {
     if (tryPush (node))
        return
     el tryPush attempts to push a node
```

```
class LockFreeStack[T] extends ConcurrentStack[T] {
 val top = new AtomicReference[Node] (null)
 protected def tryPush(node: Node): Boolean = {
   val oldTop = top.get()
    node.next = oldTop
    top.compareAndSet(oldTop, node)
 override def push(value T): Unit = {
   val node = new Node (value)
    while (true) {
      if (tryPush(node))
        return
     else backoff. Read top value
```

```
class LockFreeStack[T] extends ConcurrentStack[T] {
 val top = new AtomicReference[Node] (null)
 protected def tryPush(node: Node): Boolean = {
   val oldTop = top.get()
   node.next = oldTop
    top.compareAndSet(oldTop, node)
 override def push(value: T): Unit = {
   val node = new Node(value)
    while (true) {
      if (tryPush (node)
        return
      else backoff.backofi()
     current top will be new node's successor
```

```
class LockFreeStack[T] extends ConcurrentStack[T] {
 val top = new AtomicReference[Node] (null)
 protected def tryPush(node: Node): Boolean = {
   val oldTop = top.get()
   node.next = oldTop
   top.compareAndSet(oldTop, node)
 override def push (value: T
   val node = new Node (valu
    while (true) {
      if (tryPush (node)
        return
 Try to swing top, return success or failure
```

```
class LockFreeStack[T] extends ConcurrentStack[T] {
 val top = new AtomicReference[Node] (null)
 protected def tryPush(node: Node): Boolean = {
    val oldTop = top.get()
    node.next = oldTop
    top.compareAndSet(oldTop, node)
  override def push (value: T): Unit = {
    val node = new Node(value)
    while (true)
     if (tryPush(node))
        return
      else backoff.backoff()
```

Push calls tryPush

```
class LockFreeStack[T] extends ConcurrentStack[T] {
 val top = new AtomicReference[Node] (null)
 protected def tryPush(node: Node): Boolean = {
   val oldTop = top.get()
    node.next = oldTop
    top.compareAndSet(oldTop, node)
  override def push (value: T): Unit = {
   val node = new Node(value)
      if (tryPush(no
        return
      else backoff.backoff()
                 Create new node
```

```
class LockFreeStack[T] extends ConcurrentStack[T] {
 val top = new AtomicReference[Node] (null)
 protected def
                      If tryPush() fails,
   val oldTop = t
   node.next = back off before retrying
    top.compareAndSet(OldLop, node)
 override def push (value: T): Unit = {
   val node = new Node (value)
   while (true) {
     if (tryPush(node))
        return
     else backoff.backoff()
```

- Good
 - No locking
- Bad
 - Without GC, fear ABA
 - Without backoff, huge contention at top
 - In any case, no parallelism

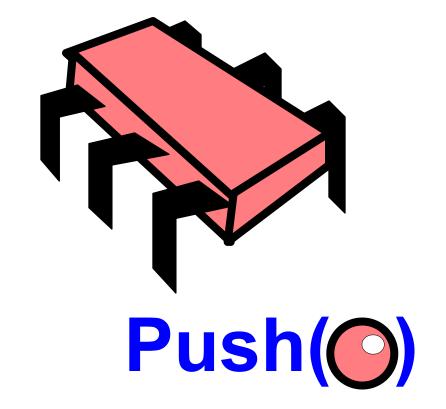
Big Question

- Are stacks inherently sequential?
- Reasons why
 - Every pop() call fights for top item
- Reasons why not
 - Stay tuned …

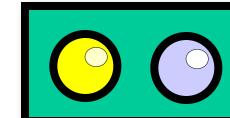
Elimination-Backoff Stack

- How to
 - "turn contention into parallelism"
- Replace familiar
 - exponential backoff
- With alternative
 - elimination-backoff

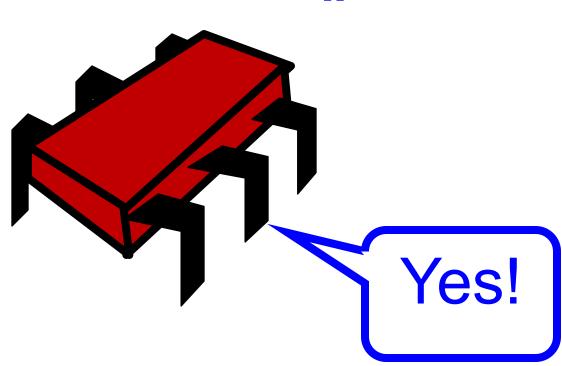
Observation



linearizable stack

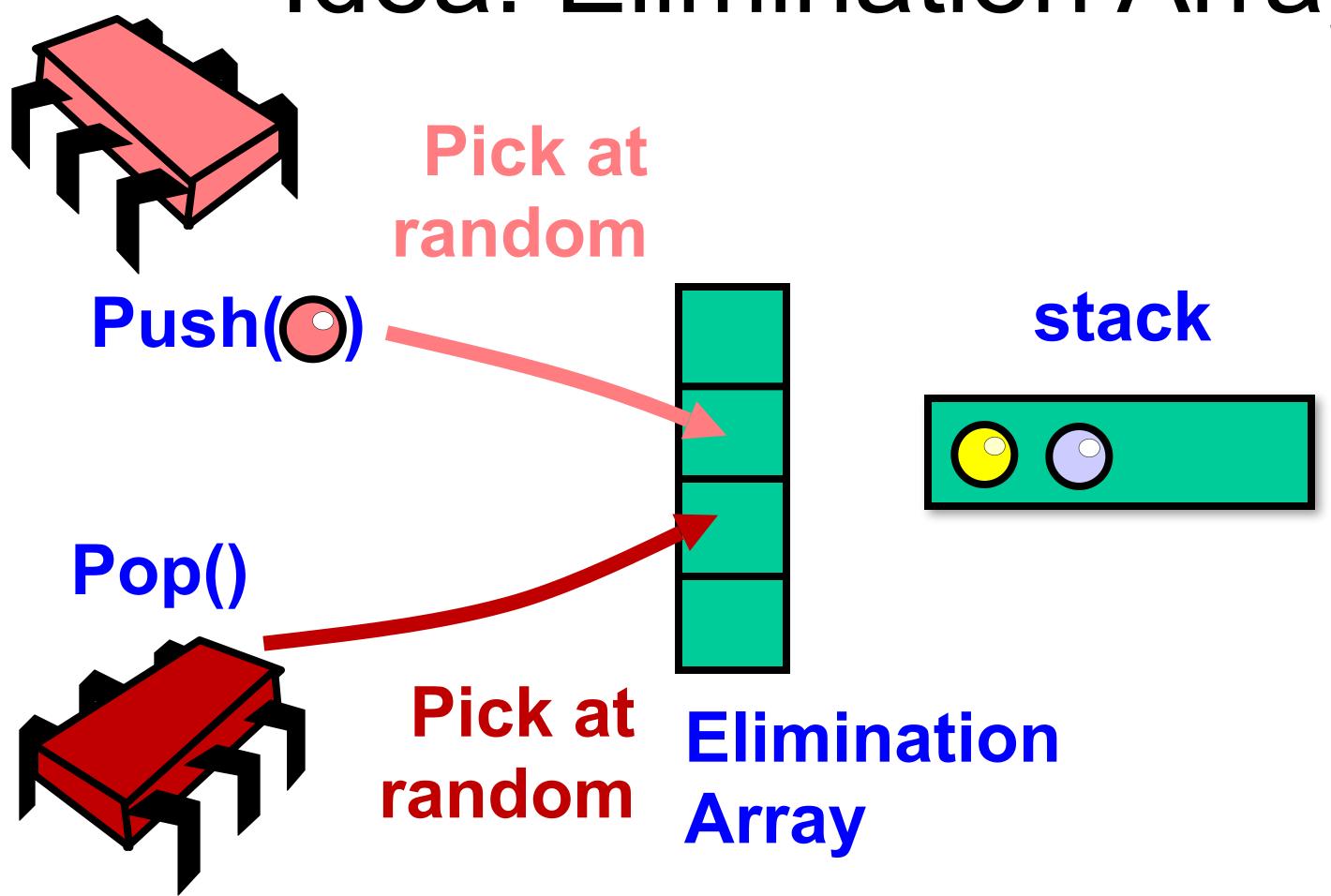


Pop()

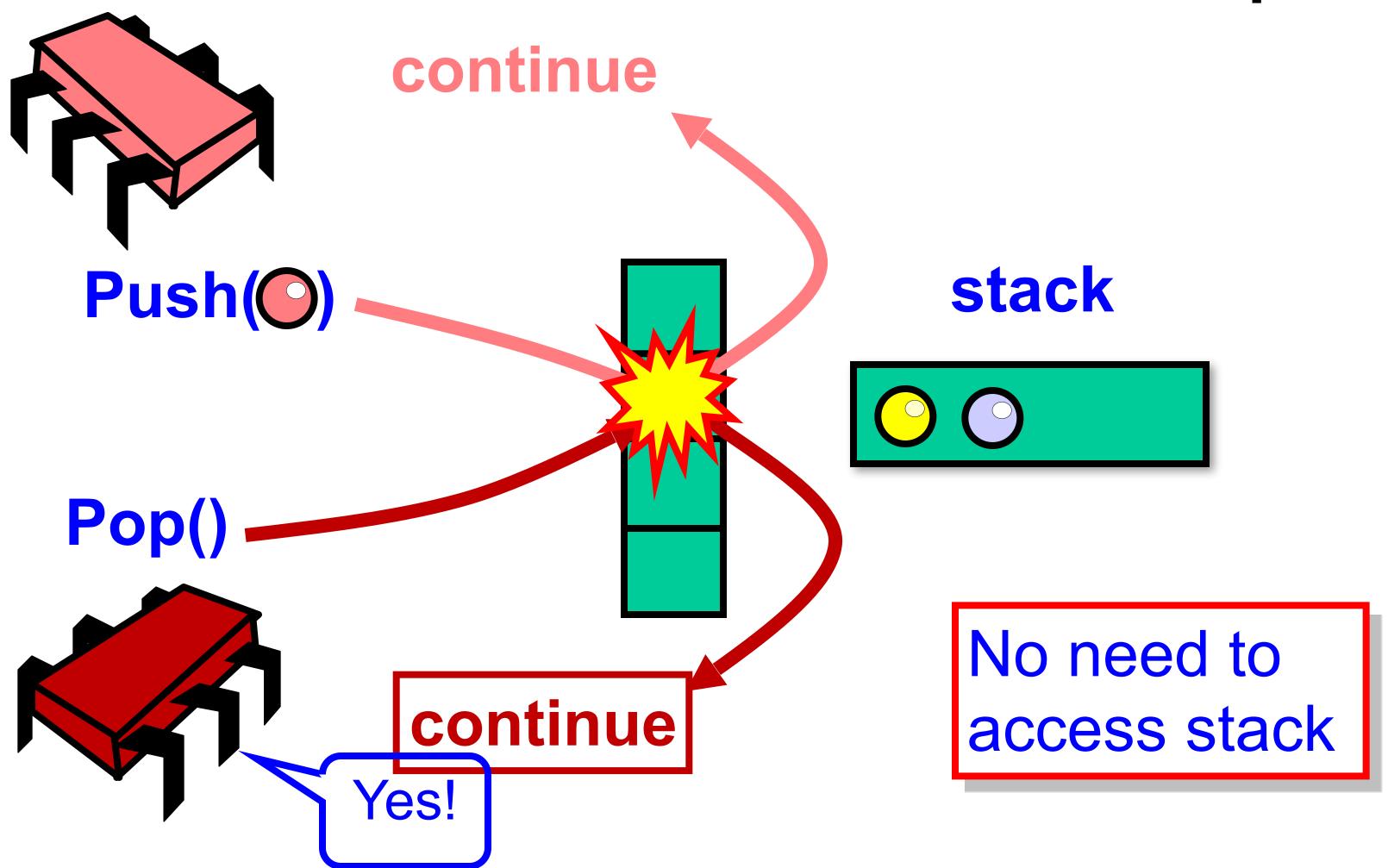


After an equal number of pushes and pops, stack stays the same

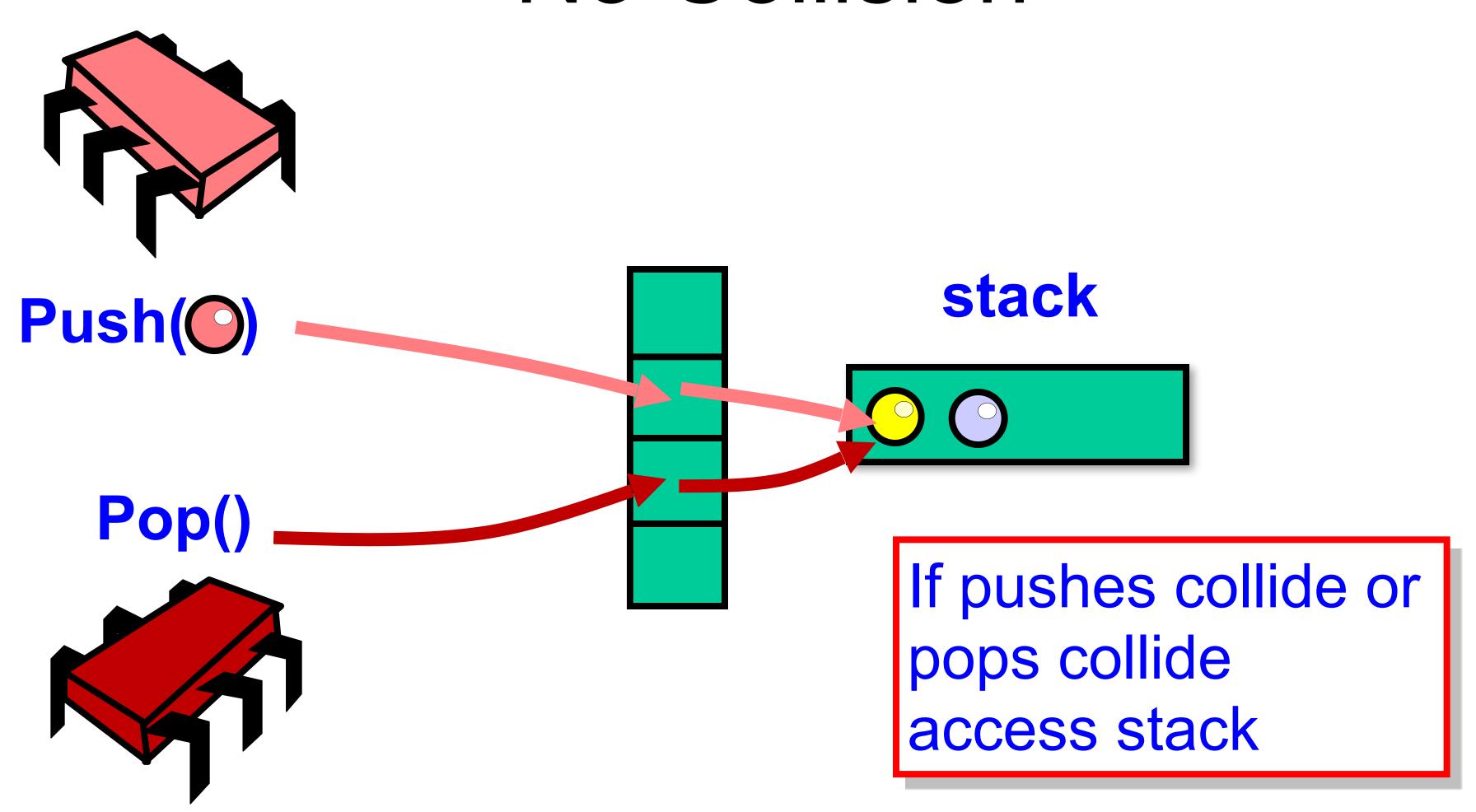
Idea: Elimination Array



Push Collides With Pop



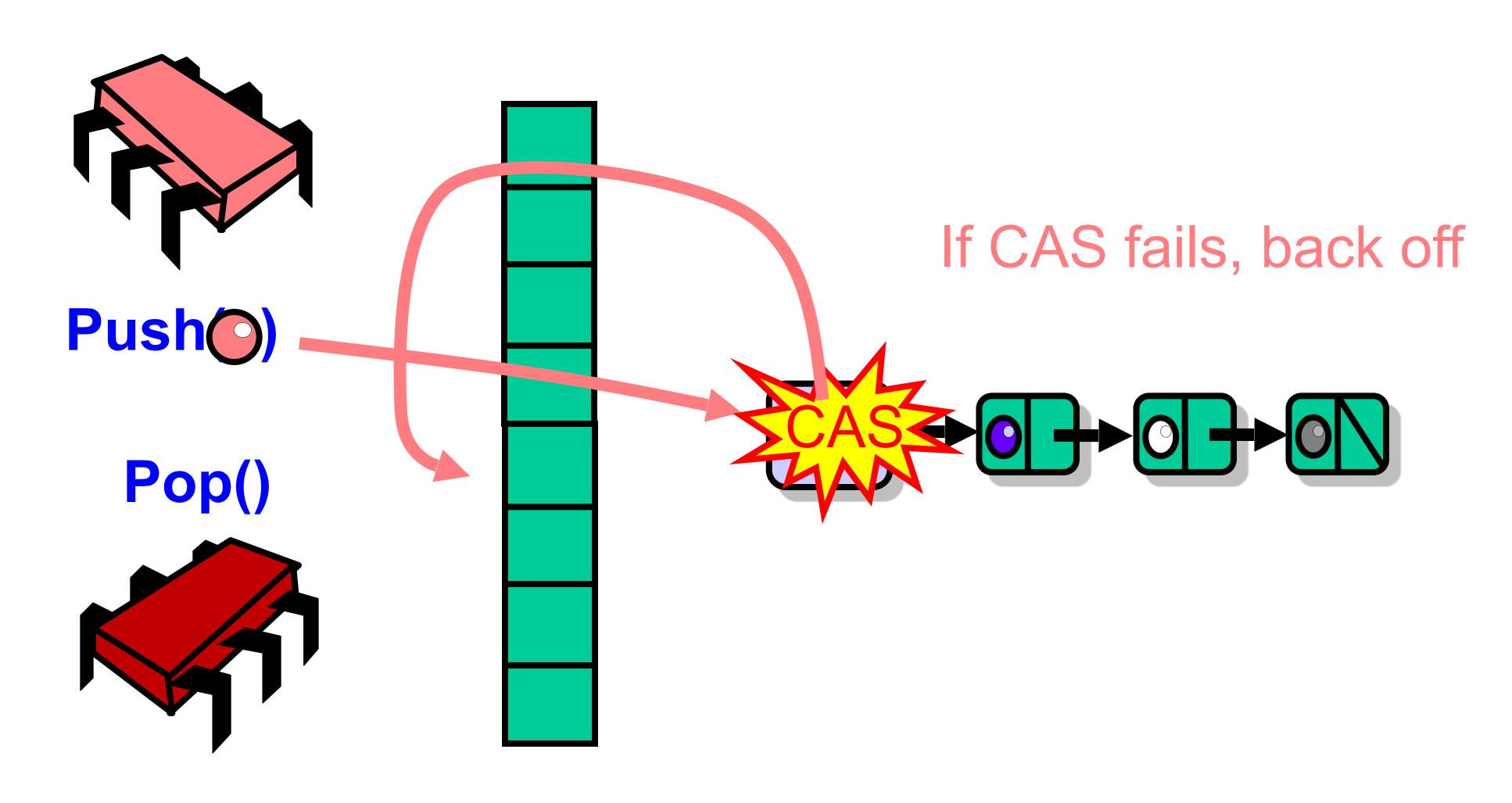
No Collision



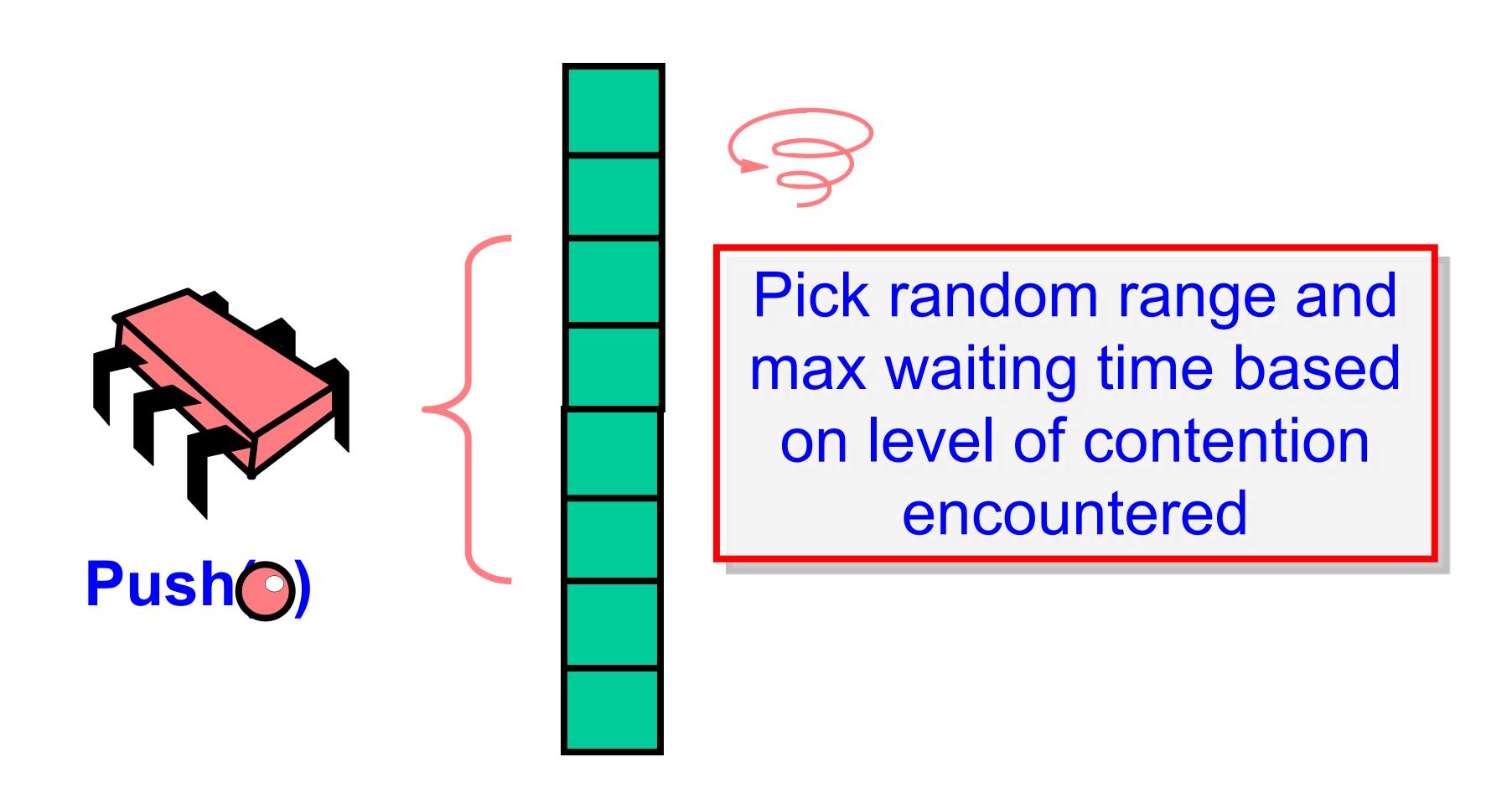
Elimination-Backoff Stack

- Lock-free stack + elimination array
- Access Lock-free stack,
 - If uncontended, apply operation
 - if contended, back off to elimination array and attempt elimination

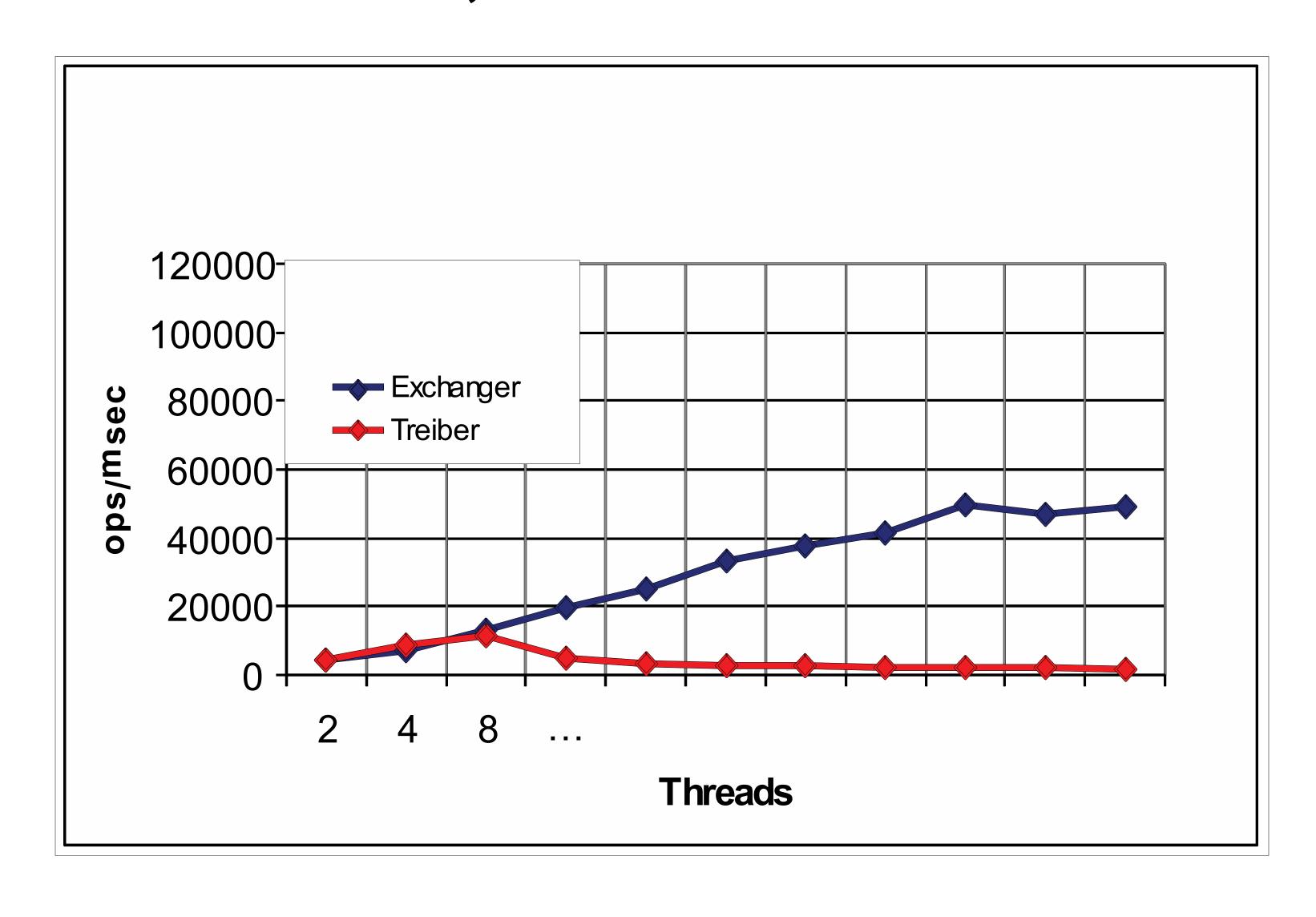
Elimination-Backoff Stack



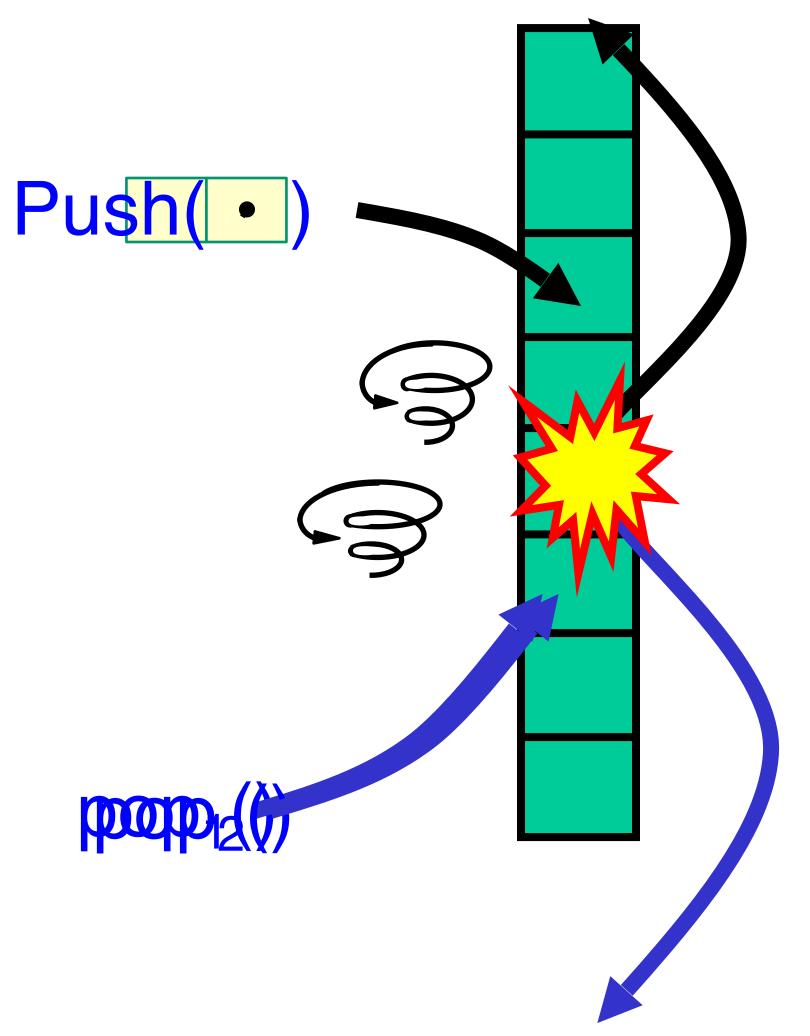
Dynamic Range and Delay



50-50, Random Slots

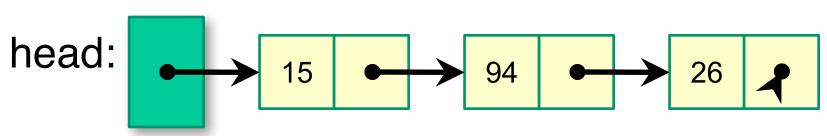


Asymmetric Rendevous

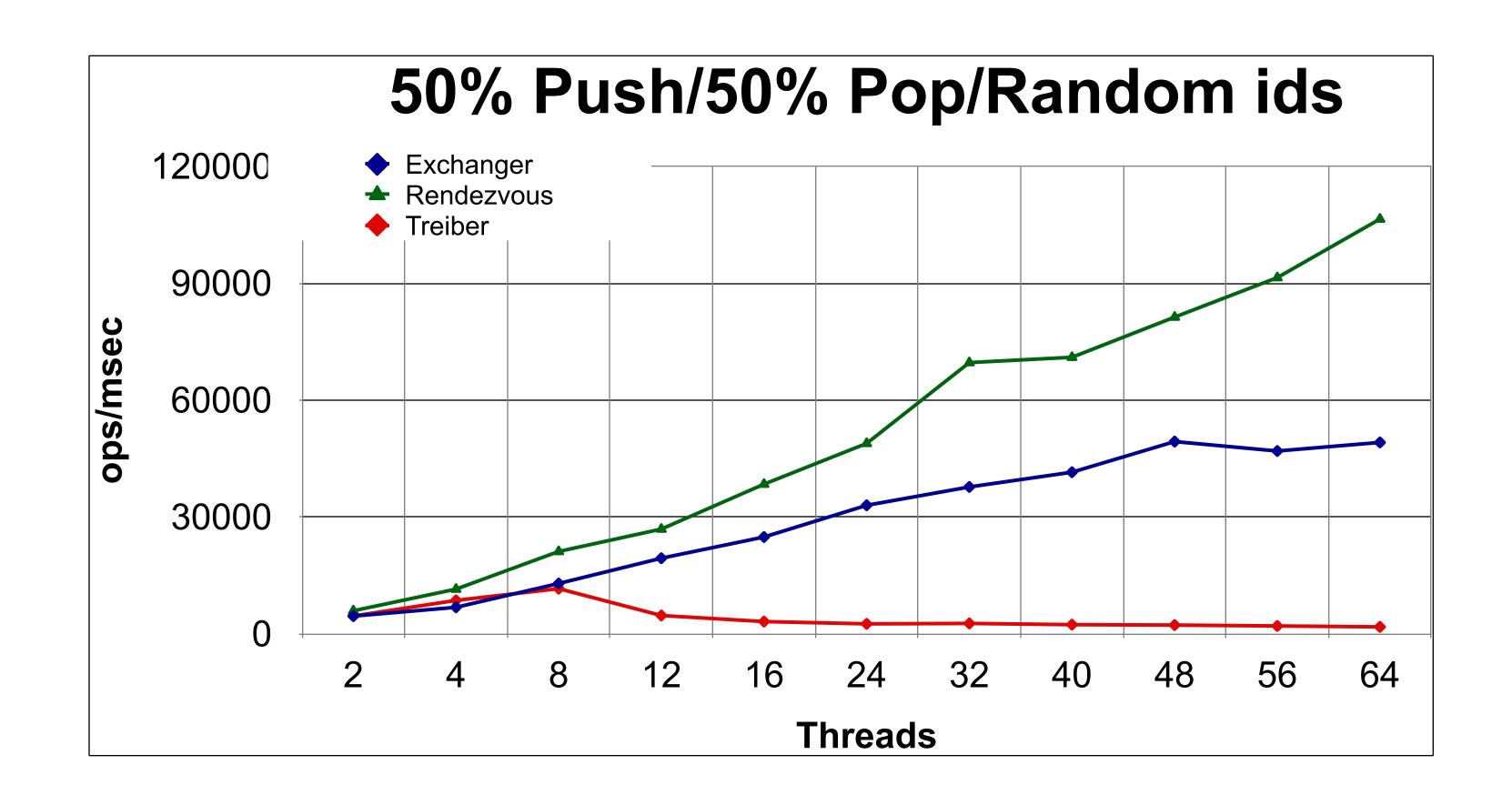


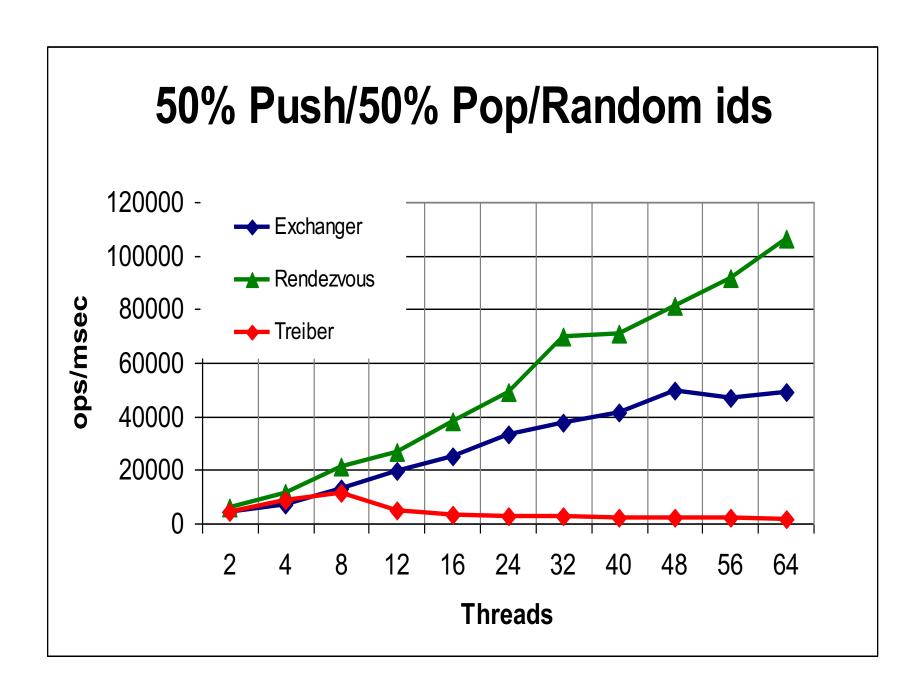
Pops find first vacant slot and spin.

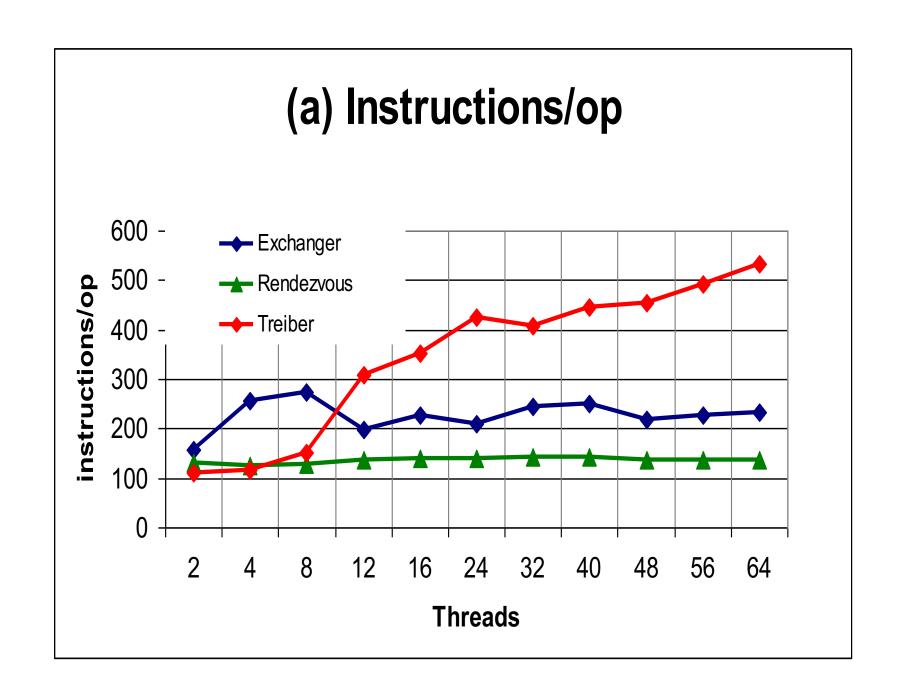
Pushes hunt for pops.

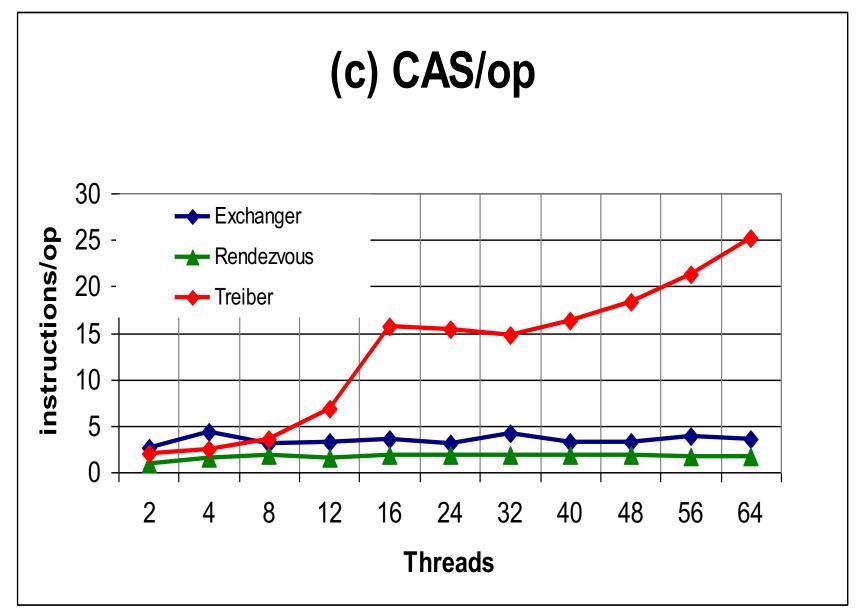


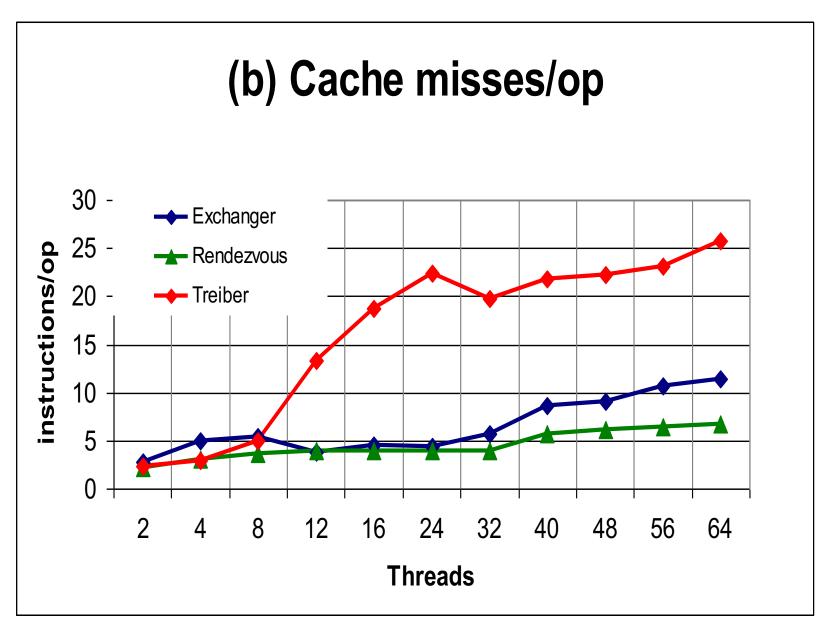
Asymmetric vs. Symmetric



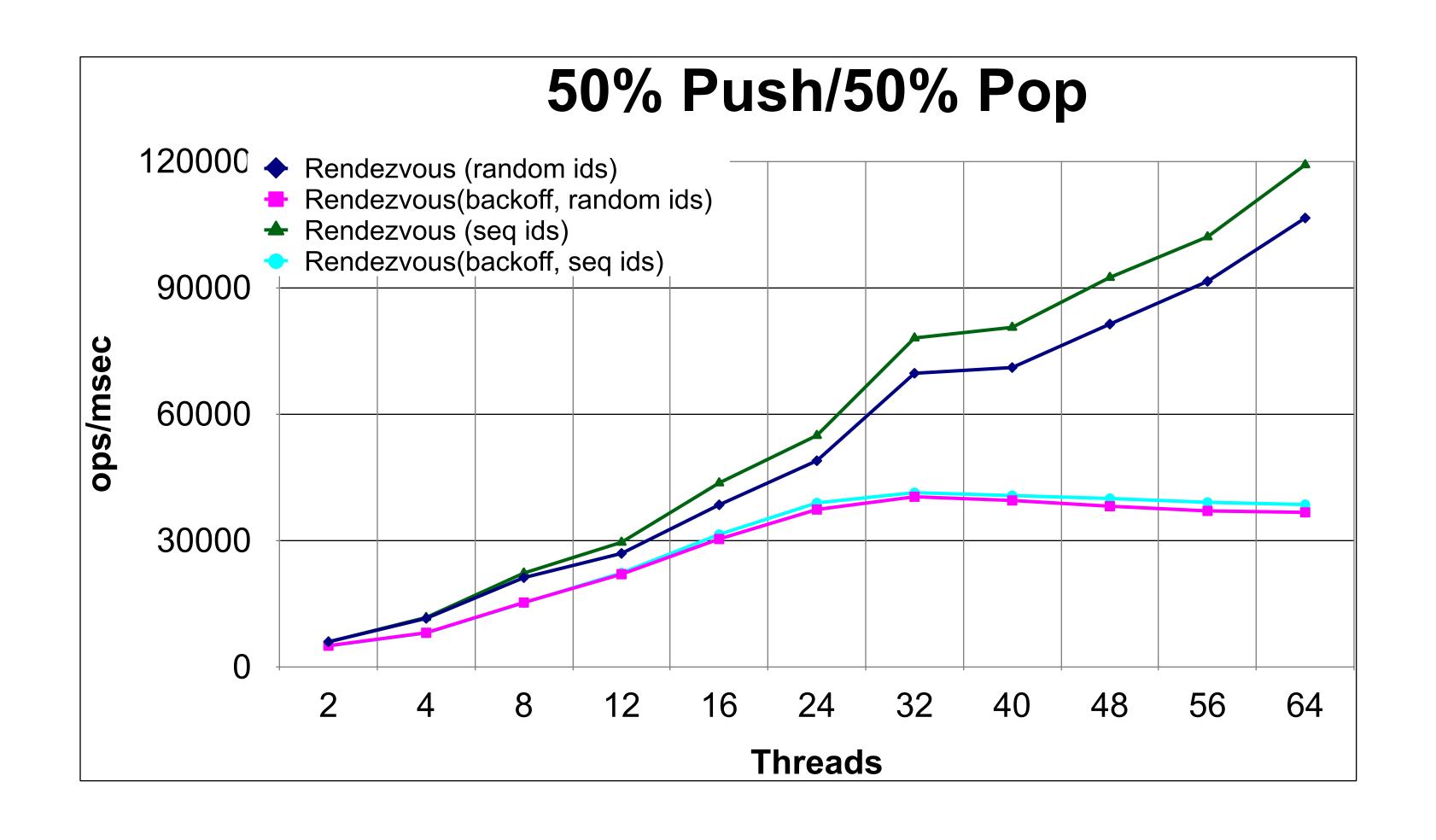






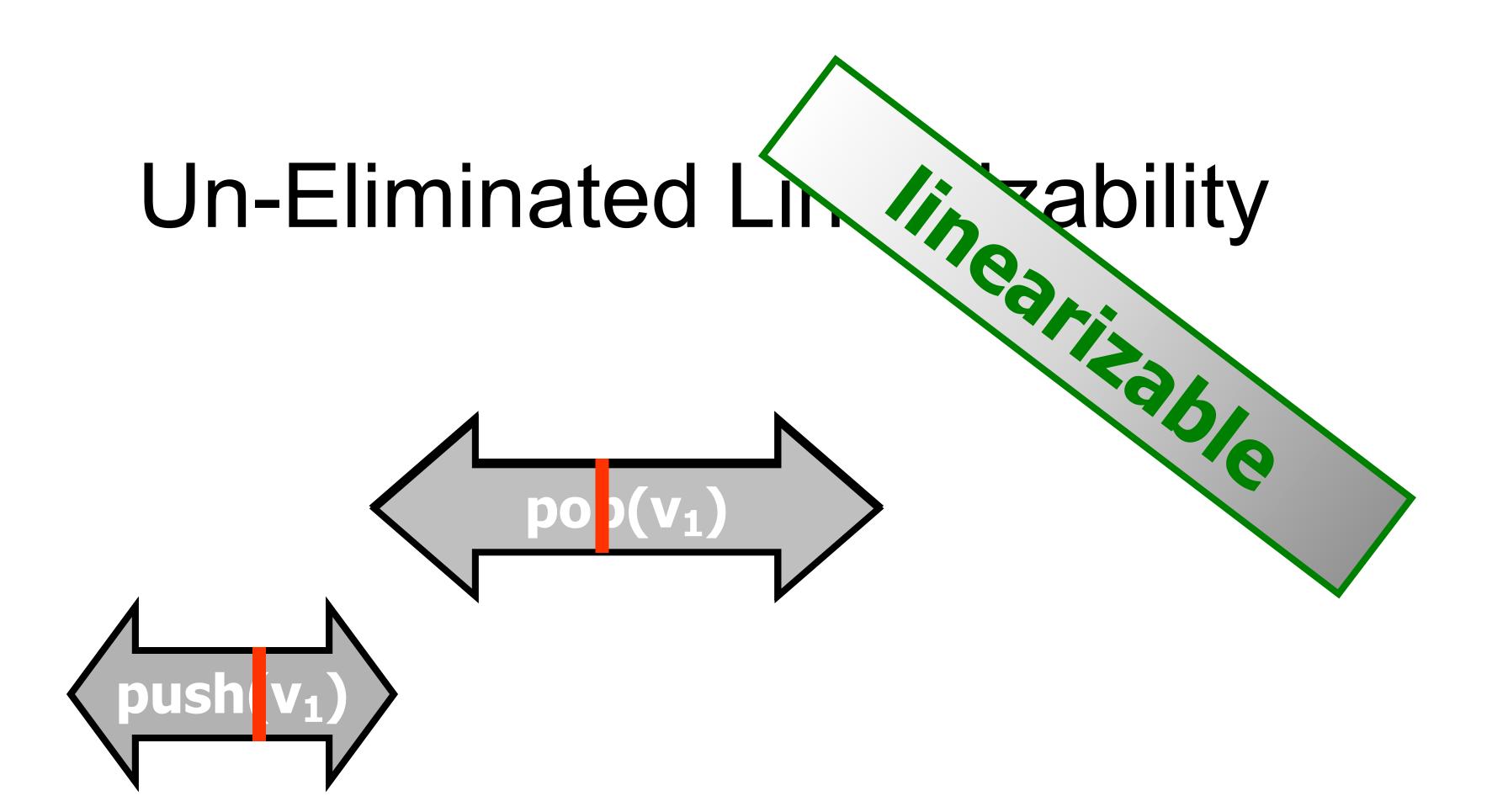


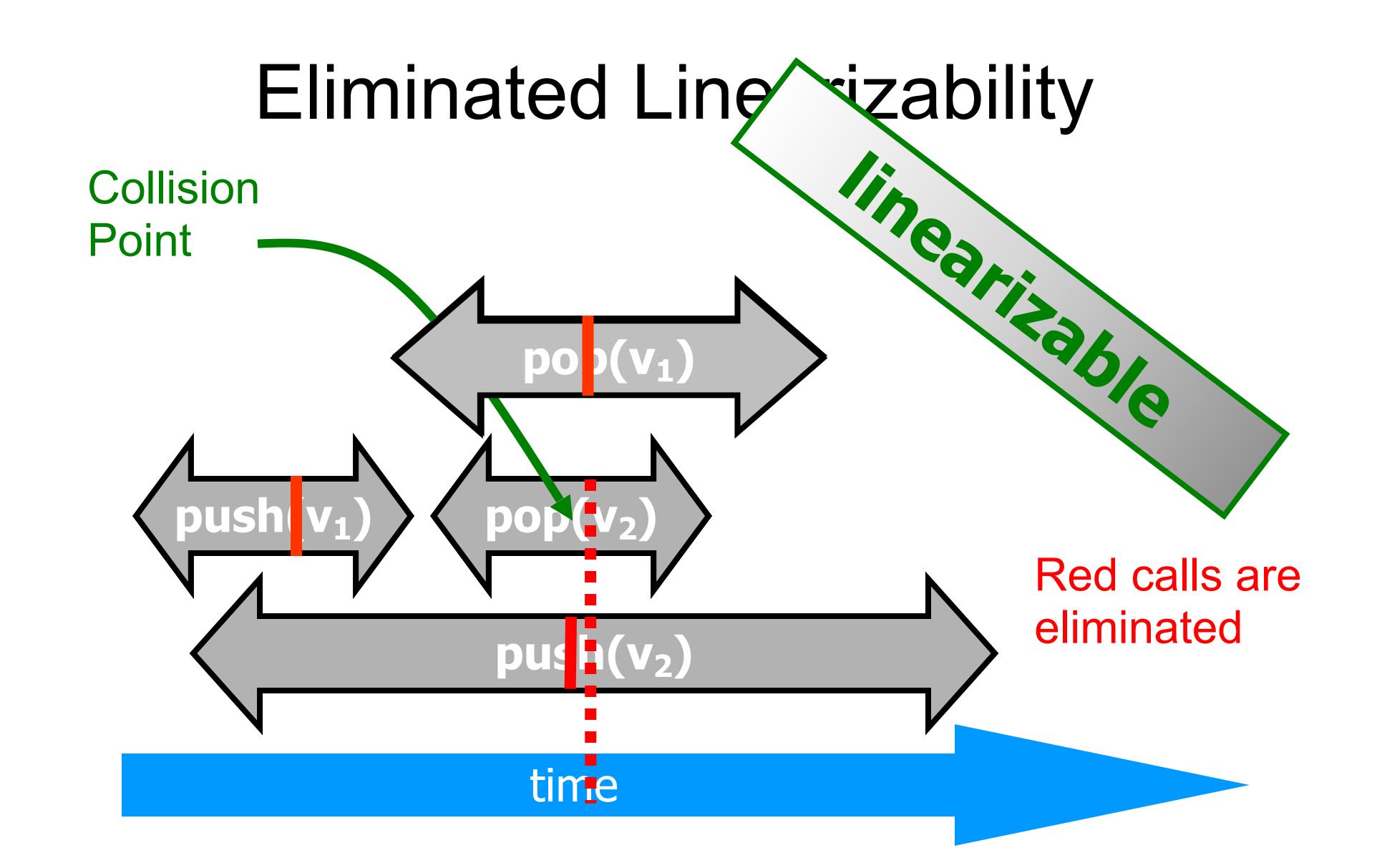
Effect of Backoff and Slot Choice



Linearizability

- Un-eliminated calls
 - linearized as before
- Eliminated calls:
 - linearize pop() immediately after matching push()
- Combination is a linearizable stack





Backoff Has Dual Effect

- Elimination introduces parallelism
- Backoff to array cuts contention on lock-free stack
- Elimination in array cuts down number of threads accessing lock-free stack

Elimination Array

```
class EliminationArray[T: ClassTag] {
  val duration = 10
 private val size = ...
 val exchangers = Array.fill(size)(new Exchanger[T])
 val random = new Random()
 def visit (value: T, range: Int): T = {
    val slot = random.nextInt(range)
    exchangers (slot).exchange (value, duration)
```

Elimination Array

```
class EliminationArray[T: ClassTag] {
 val duration = 10
 private val size = ...
 val exchangers = Array.fill(size)(new Exchanger[T])
                      An array of Exchangers
 def visit (value: T,
   val slot = random.nextInt(range)
    exchangers (slot).exchange (value, duration)
```

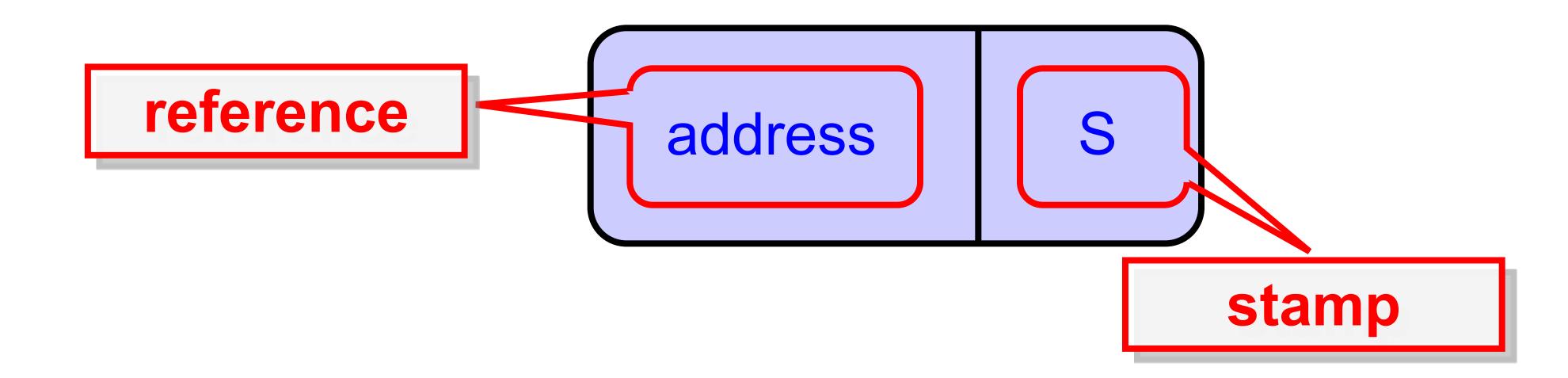
Digression: A Lock-Free Exchanger

A Lock-Free Exchanger

```
class Exchanger[T] {
 val slot = new AtomicStampedReference[T] (null, 0)
  def exchange (mylte
                timeout: Long,
                Atomically modifiable
                  reference + status
```

Atomic Stamped Reference

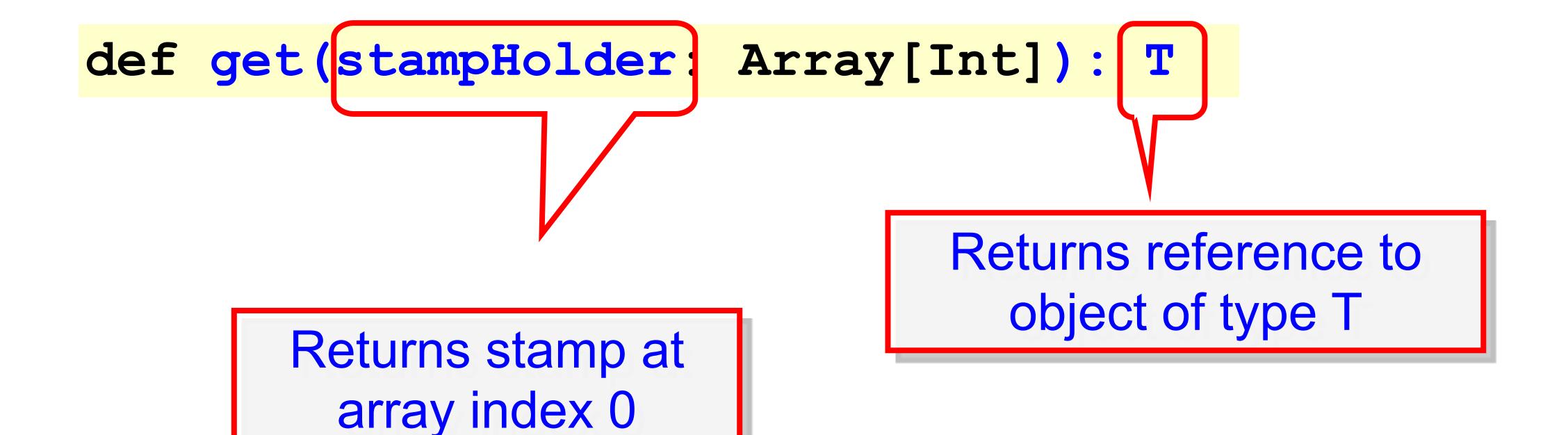
- AtomicStampedReference class
 - Java.util.concurrent.atomic package
- In C or C++:



Extracting Reference & Stamp

```
def get(stampHolder: Array[Int]): T
```

Extracting Reference & Stamp



Exchanger Status

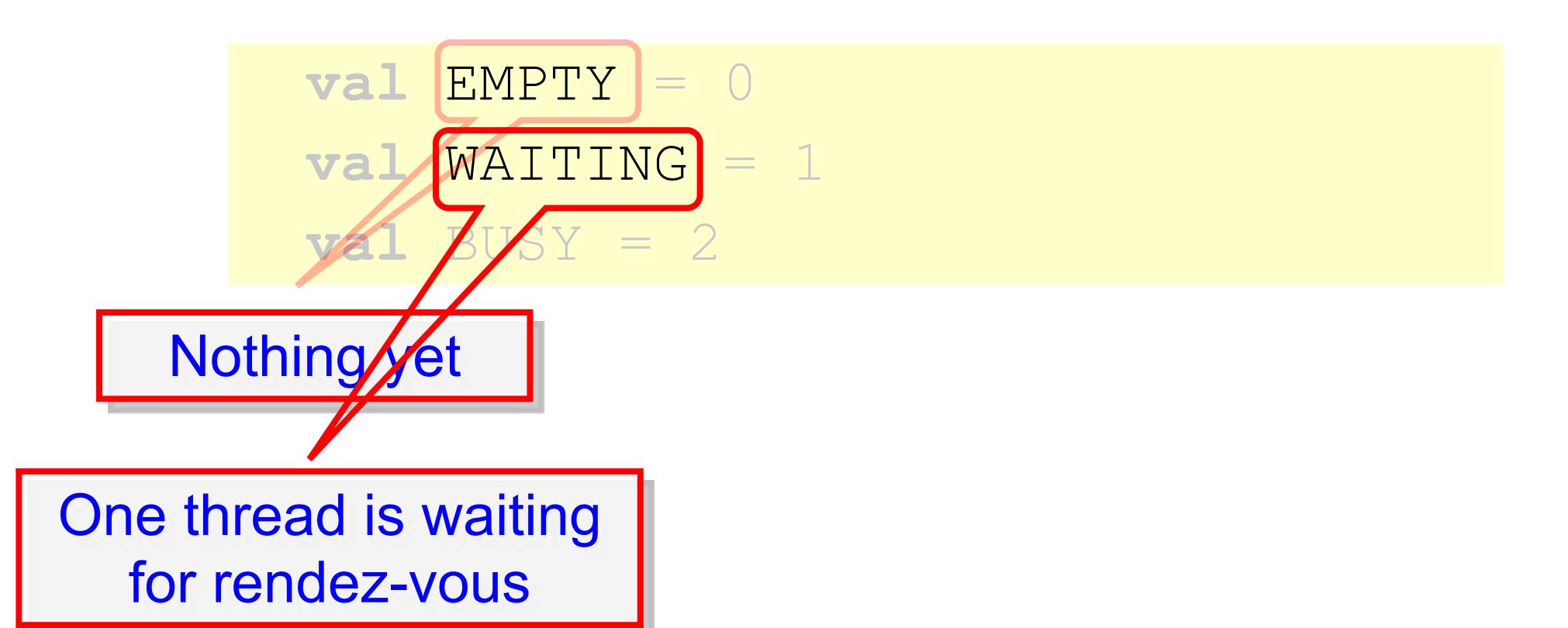
```
val EMPTY = 0
val WAITING = 1
val BUSY = 2
```

Exchanger Status

```
val EMPTY = 0
val MAITING = 1
val BUSY = 2
```

Nothing yet

Exchange Status



Exchange Status

```
val EMPTY
          val WAITING = 1
              BUSY
     Nothing yet
 One thread is waiting
   for rendez-vous
Other threads busy
with rendez-vous
```

```
def exchange (myItem: T, timeout: Long): T = {
 val timeBound = System.nanoTime() + timeout
 val stapmholder = Array(EMPTY)
 while (true) {
    if (System.nanoTime() > timeBound)
      throw new TimeoutException
   var yrItem = slot.get(stapmholder)
    stapmholder(0) match {
     case EMPTY => ... // slot is free
      case WAITING => ... // someone waiting for me
     case BUSY => ... // others exchanging
```

```
def exchange myItem: T, timeout: Long: T = {
  val timeBound = System.nanoTime() + timeout
  val stapmholder = Array(EMPTY
  while (true) {
   if (System.nano Item and timeout
      throw new TimeoutException
    var yrItem = slot.get(stapmholder)
    stapmholder(0) match {
      case EMPTY => ... // slot is free
      case WAITING => ... // someone waiting for me
      case BUSY => ... // others exchanging
```

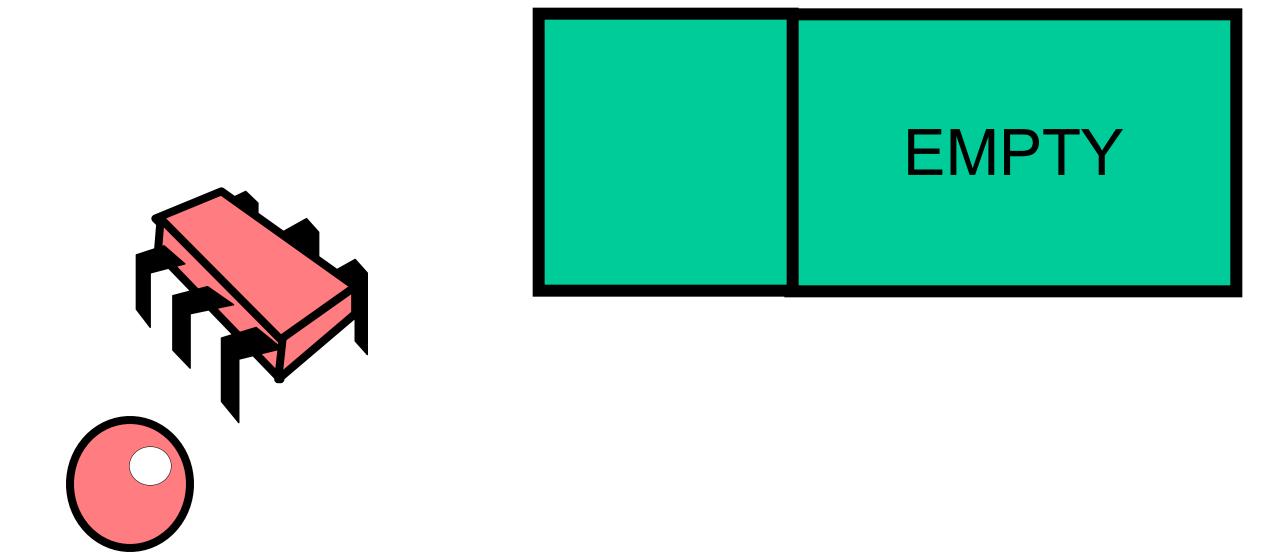
```
def exchange (myItem: T, timeout: Long): T = {
  val timeBound = System.nanoTime() + timeout
 val stapmholder = Array(EMPTY)
    if (System.nanoTime() > timeBound)
      throw new TimeoutException
    var yrItem = s
stapmholder(0)
Array holds status
      case EMPTY => ... // slot is free
      case WAITING => ... // someone waiting for me
      case BUSY => ... // others exchanging
```

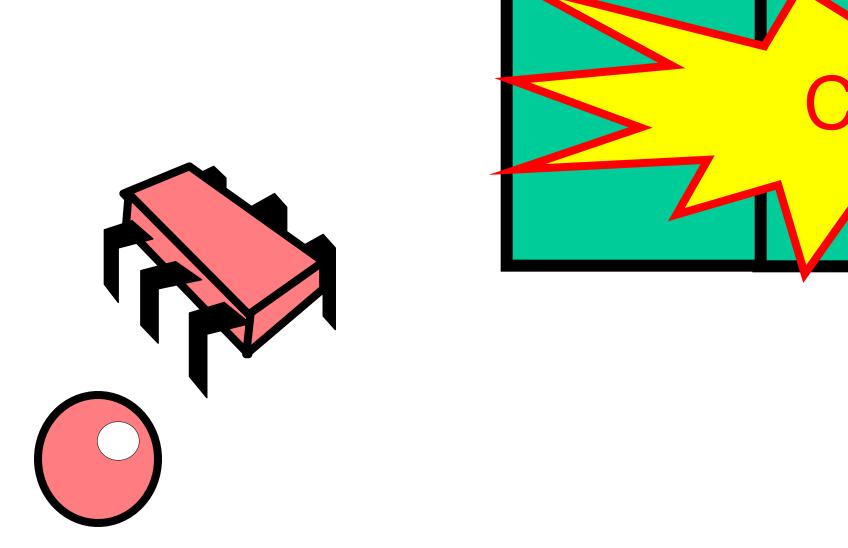
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def exchange (myItem: T, timeout: Long): T = {
  val timeBound = System.nanoTime() + timeout
  val stapmholder = Array(EMPTY)
  while (true)
   if (System.nanoTime() > timeBound)
      throw new TimeoutException
    var yrItem = slot.get(stapmholder)
    stapmholder(0) match {
      case EMPTX => ... // slot is free
      case WAITING => ... // someone waiting for me
      case BUSY => ... // others exchanging
             Loop until timeout
```

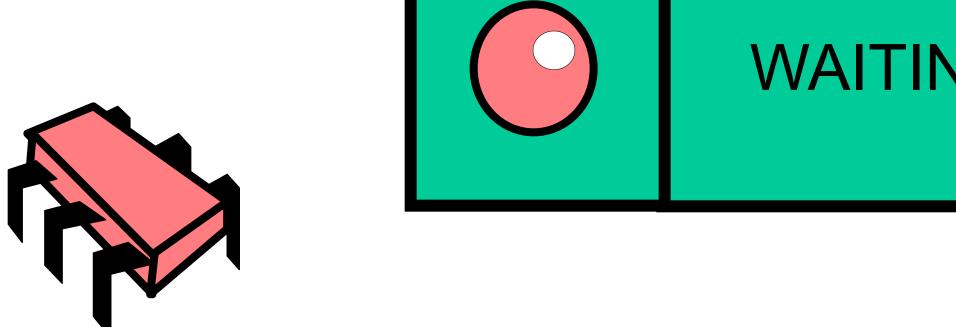
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def exchange (myItem: T, timeout: Long): T = {
 val timeBound = System.nanoTime() + timeout
 val stapmholder = Array(EMPTY)
 while (true) {
   if (System.nanoTime() > timeBound)
     throw new TimeoutException
   var yrItem = slot.get(stapmholder)
   stapmholder(0) match {
     case WAITING => ... // someone waiting for me
     case BUSY => ... // others exchanging
                Get other's item and status
```

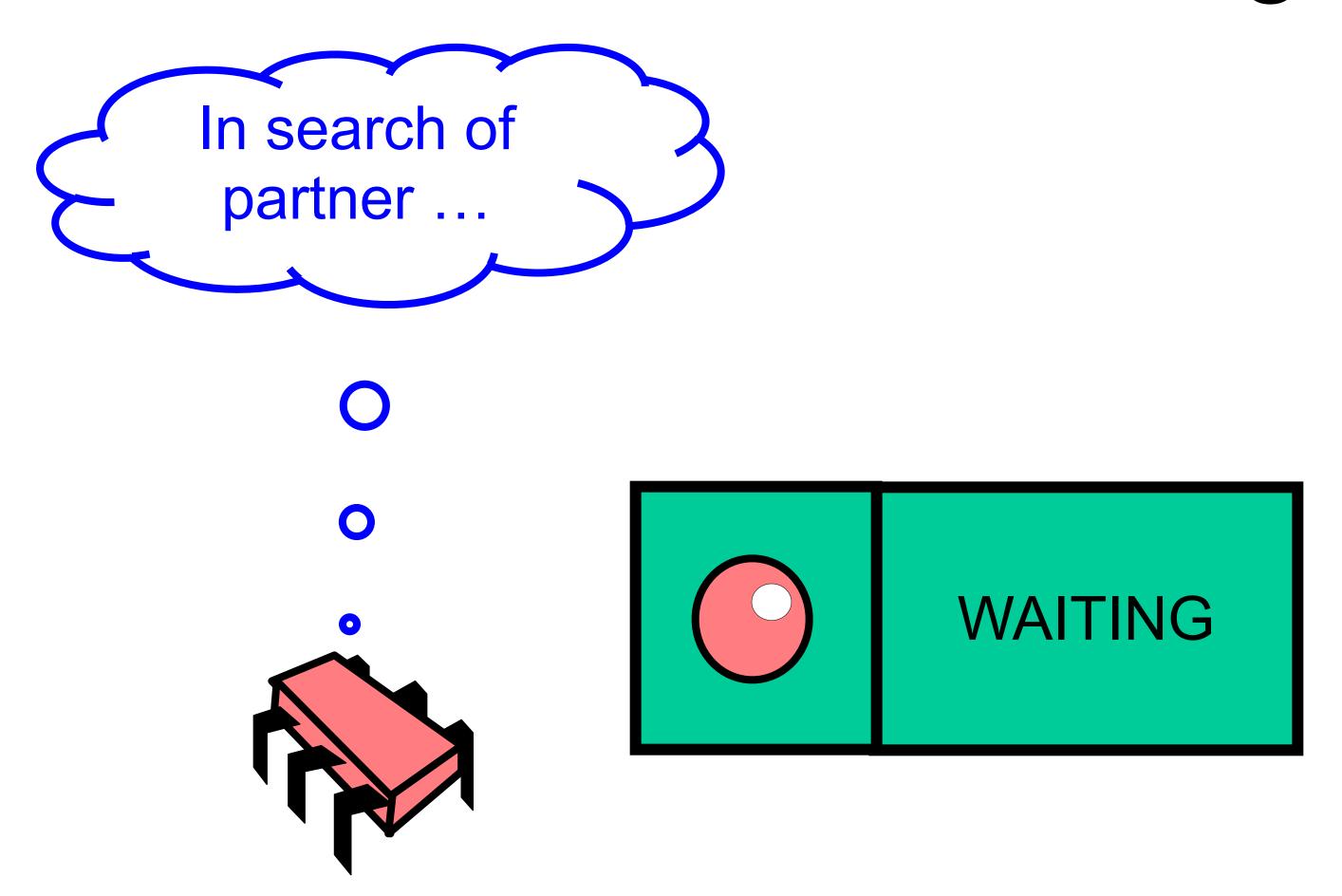
An Exchanger has three possible states

```
val timeBound = System.nanoTime() + timeout
val stapmholder = Array(EMPTY)
while (true)
 if (System.nanoTime() > timeBound)
    throw new TimeoutException
  var vrItem = slot.get(stapmholder)
  stapmholder(0) match {
   case EMPTY => ... // slot is free
   case WAITING => ... // someone waiting for me
   case BUSY => ... // others exchanging
```

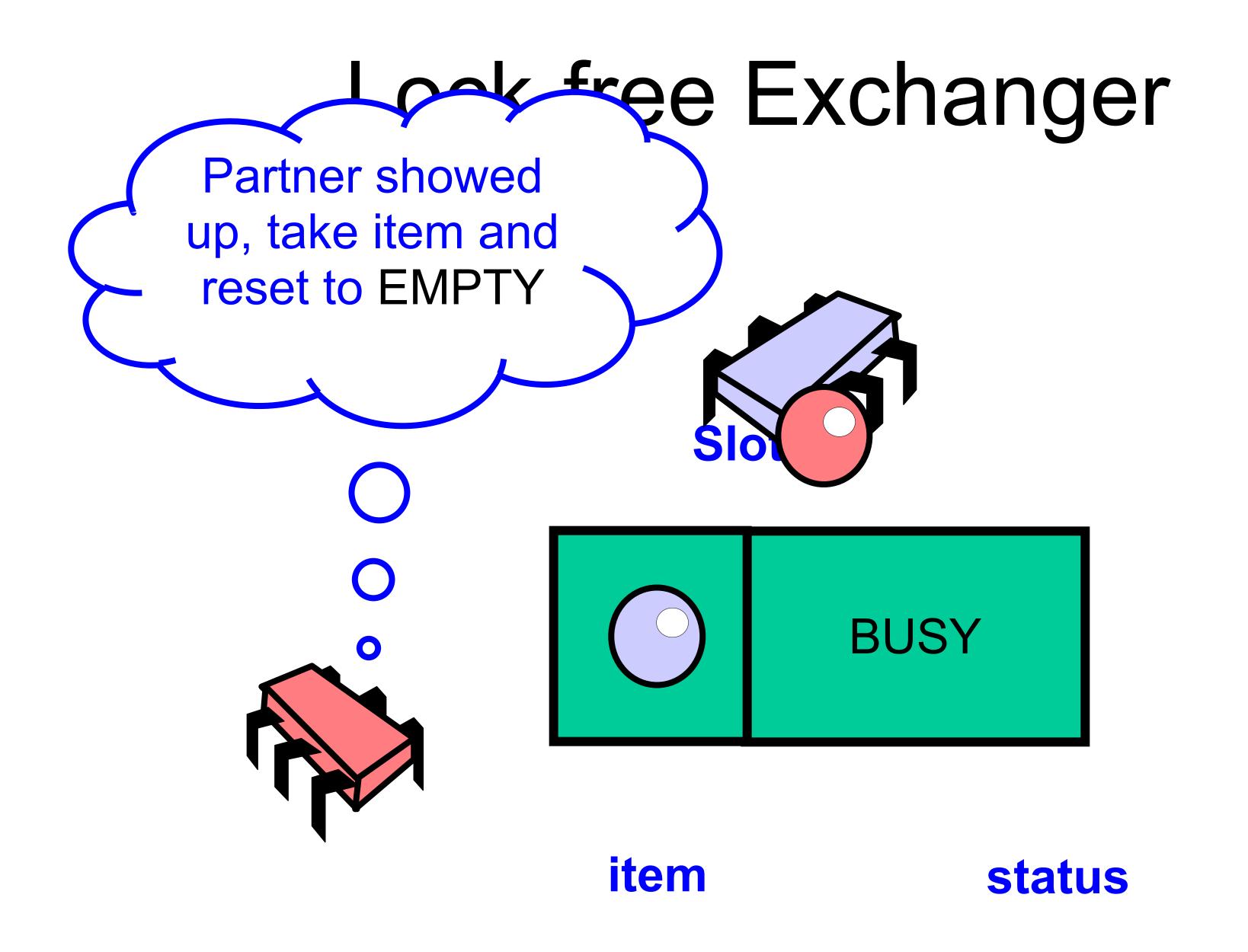


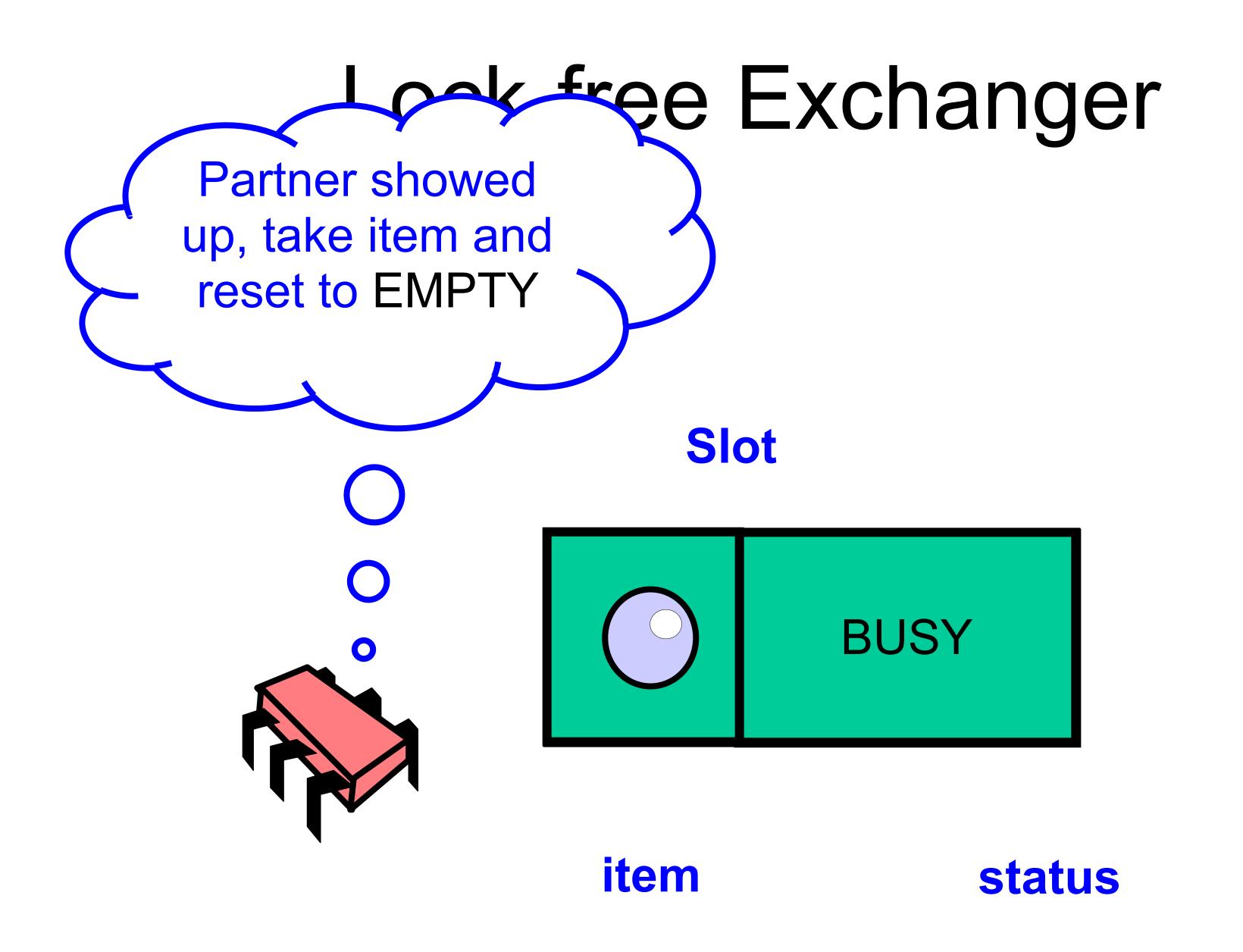












Can start skipping from here if running out of time

Exchanger State EMPTY

```
case EMPTY =>
  if (slot.compareAndSet(yrItem, myItem, EMPTY, WAITING)) {
    while (System.nanoTime() < timeBound) {</pre>
      yrItem = slot.get(stapmholder)
      if (stapmholder(0) == BUSY) {
         slot.set(null, EMPTY)
         return yrItem
    if (slot.compareAndSet(myItem, null, WAITING, EMPTY)) {
      throw new TimeoutException
    } else {
      yrItem = slot.get(stapmholder)
      slot.set(null, EMPTY)
      return yrItem
```

Exchanger State EMPTY

```
case EMPTY =>
     slot.compareAndSet(yrItem, myItem, EMPTY, WAITING)
         (System.nanoTime()
     yrItem = slot.get(stapmholder)
     if (stapmholder(0) == BUSY) {
        slot.set(null, EMPTY
        Try to insert myltem and
                  change state to WAITING
   if (slot.compareAndSet(myItem, null, WAITING, EMPTY))
      throw new TimeoutException
    } else {
     yrItem = slot.get(stapmholder)
     slot.set(null, EMPTY)
     return yrItem
```

```
case EMPTY =>
  if (slot.compareAndSet(vrItem, myItem, EMPTY, WAITING)) {
   while (System.nanoTime() < timeBound)</pre>
      yrItem = slot.get(stapmholder)
      if (stapmholder(0) == BUSY)
         slot.set(null, EMPTY)
         return yrItem
                      Spin until either
    if (slot.c
                                               ING, EMPTY)) {
      throw new myltem is taken or timeout
     else {
      yrItem = slot.get(stapmholder)
      slot.set(null, EMPTY)
      return yrItem
```

```
case EMPTY =>
  if (slot.compareAndSet(yrItem, myItem, EMPTY, WAITING)) {
    while (System.nanoTime() < timeBound) {</pre>
      yrItem = slot.get(stapmholder)
      if (stapmholder(0) == BUSY) {
         slot.set(null, EMPTY)
         return yrItem
    if (slot.compareAndSet(myItem, null, WAITING, EMPTY)) {
      throw new TimeoutException
      else
               myltem was taken,
      yrIte
                 so return yrltem
      slot.
      retur
             that was put in its place
```

```
m, EMPTY, WAITING)) {
nd) {
case EMPTY =>
      Otherwise we ran out of time,
      try to reset status to EMPTY
               and time out
         return vrItem
    if (slot.compareAndSet(myItem, null, WAITING, EMPTY)) {
      throw new TimeoutException
      else
      yrItem = slot.get(stapmholder)
      slot.set(null, EMPTY)
      return yrItem
```

```
case EMPTY =>
 if (slot.compareAndSet(yrItem, myItem, EMPTY, WAITING)) {
    while (System.nanoTime() < timeBound)
      yrItem
                       If reset failed,
      if (sta
              someone showed up after all,
         retu
                     so take that item
               mpareAndSet(myItem, null, WAITING, EMPTY)) {
      else {
      yrItem = slot.get(stapmholder)
      slot.set(null, EMPTY)
      return yrItem
```

```
case EMPTY =>
  if (slot.compareAndSet(yrItem, myItem, EMPTY, WAITING)) {
    while (System.nanoTime() < timeBound) {</pre>
      yrItem = slot.get(stapmholder)
      if (stapmholder(0) == BUSY) {
             Clear slot and take that item
               mpareAndSet(myItem, null, WAITING, EMPTY)) {
      throw new TimeoutException
      else
                slot get (stapmholder)
      slot.set(null, EMPTY)
      return yrItem
```

```
case EMPTY =>
 if (slot.compareAndSet(yrItem, myItem, EMPTY, WAITING)) {
    while (System.nanoTime() < timeBound) {</pre>
      yrItem = slot.get(stapmholder)
                   If initial CAS failed,
          then someone else changed status
               from EMPTY to WAITING,
                                                   EMPTY))
                    so retry from start
      else
             = slot.get(stapmholder)
      slot/set(null, EMPTY)
```

```
case WAITING =>
  if (slot.compareAndSet(yrItem, myItem, WAITING, BUSY)) {
    return yrItem
  }
case BUSY =>
case x =>
  throw new Exception("Cannot happen")
```

```
case WAITING =>
if (slot.compareAndSet(yrItem, myItem, WAITING, BUSY)) {
   return yrItem
}
case BUSY =>
case x =>
throw new E;
   someone is waiting to exchange,
        so try to CAS my item in
        and change state to BUSY
```

```
case WAITING =>
  if (slot.compareAndSet(yrItem, myItem, WAITING, BUSY)) {
    return yrItem
  }
case BUSY =>
  case x =>
  throw new Exception("Cannot happen")
```

If successful, return other's item, otherwise someone else took it, so try again from start

The Exchanger Slot

- Exchanger is lock-free
- Because the only way an exchange can fail is if others repeatedly succeeded or no-one showed up
- The slot we need does not require symmetric exchange

Stop skipping

Back to the Stack: the Elimination Array

```
class EliminationArray[T: ClassTag] {
    ...
    def visit(value: T, range: Int): T = {
       val slot = random.nextInt(range)
       exchanger(slot).exchange(value, duration)
    }
}
```

Elimination Array

```
class EliminationArray[T: ClassTag] {
    def visit(value: T, range: Int): T = {
       val slot = random.NextInt(range)
       exchanger(slot).exchange(value, duration)
    }
}
```

visit the elimination array with fixed value and range

Elimination Array

```
class EliminationArray[T: ClassTag] {
    def visit(value: T, range: Int): T = {
      val slot = random.nextInt(range)
      exchanger(slot).exchange(value, duration)
    }
}
```

Pick a random array entry

Elimination Array

```
override def push(value: T): Unit = {
 while (true) {
    if (tryPush(node))
      return
    else try {
      val otherValue =
        eliminationArray.visit(value, rangePolicy.getRange)
      if (otherValue == null)
        return
```

```
override def push(value: T): Unit = {
  while (true)
    if (tryPush(node))
      return
    else try
      val otherValue
        eliminationArray.wisit(value, rangePolicy.getRange)
      if (otherValue ==
        return
                        First, try to push
```

```
override def push (value: T): Unit = {
  while (true) {
    if (tryPush (node)) If I failed, backoff & try to eliminate
      return
    else try
      val otherValue =
        eliminationArray.visit(value, rangePolicy.getRange)
      if (otherValue == null)
        return
```

```
override def push (value: T): Unit =
                           Value pushed and range to try
  while (true) {
    if (tryPush(node))
      return
    else try {
      val otherValue =
        eliminationArray.visit (value, rangePolicy.getRange)
      if (otherValue == null)
        return
```

```
override def
                Only pop () leaves null,
 while (true so elimination was successful
    if (tryP
      return
    else try
      val cherValue =
                    rray.visit(value, rangePolicy.getRange)
      if (otherValue == null)
        return
```

```
override def push (value: T): Unit =
        Otherwise, retry push () on lock-free stack
  while
    if (tryPush (node))
      return
    else t
      val/otherValue =
         liminationArray.visit(value, rangePolicy.getRange)
         (otherValue == null)
        return
```

Elimination Stack Pop

```
override def pop(): T = {
  while (true) {
    val returnNode = tryPop()
    if (returnNode != null) {
      return returnNode.value
    } else try {
      val otherValue =
        eliminationArray.visit(null, rangePolicy.getRange)
      if (otherValue != null) {
        return otherValue
```

Elimination Stack Pop

```
override def pop(): T = {
   If value not null, other thread is a push (),
             so elimination succeeded
      else try
      val otherValue =
                ionArray.visit(null, rangePolicy.getRange)
     if (otherValue != null) {
        return otherValue
```

Demo: Benchmarking Stacks

Summary

- We saw both lock-based and lock-free implementations of
- queues and stacks
- Don't be quick to declare a data structure inherently sequential
 - Linearizable stack is not inherently sequential (though it is in the worst case)
- ABA is a real problem, pay attention



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