YSC4231: Parallel, Concurrent and Distributed Programming

Theory Assignment 3

In this theory assignment, some of the problems refer to an accompanying repository.[1] Feel free to clone it and play with the code, modifying it in any way you like. For your solution, though, you only have to submit a PDF with explanations and proofs — no code is required.

Problem 1. The `AtomicInteger` Java class (in the `java.util.concurrent.atomic` package) is a container for an integer value. One of its methods is

```
def compareAndSet(expect: Int, update: Int): Boolean.
```

This method compares the object's current value to expect. If the values are equal, then it atomically replaces the object's value with update and returns `true`. Otherwise, it leaves the object's value unchanged, and returns `false`. This class also provides `get() : Int` which returns the object's actual value.

Consider the FIFO queue implementation given by the class `IQueue`. It stores its items in an array items. It has two `AtomicInteger` fields: tail is the index of the next slot from which to remove an item, and head is the index of the next slot in which to place an item. Give an example showing that this implementation is *not* linearizable. Use the tests to help your reasoning.

Problem 2. Consider the class `ReaderWriter`. According to what you have been told about the Java memory model, will the reader method *divide by zero* if

(a) `writer()` and `reader()` are invoked by two concurred threads (a writer thread and a reader thread, correspondingly) just once each?
(b) `write()` and `reader()` are repeatedly invoked by two concurred threads (writer/reader) multiple times?

Please, explain your reasoning and supply examples of executions, if necessary. Feel free to play with the class `RWTest`, changing it to emulate both these scenarios (hint: you can create and start new threads within loops).

Problem 3. For each of the histories shown in Figures 1 and 2, are they sequentially consistent? Linearizable? Justify your answer.

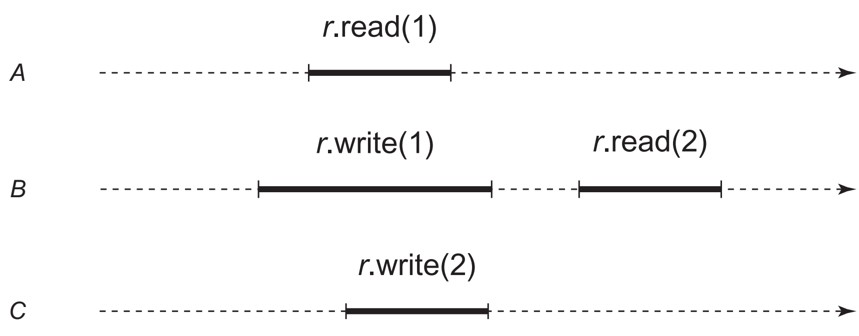

Figure 1: First history for Problem 3.

Problem 4. This problem examines a stack implementation `AGMStack` whose `push()` method does not have a single fixed linearization point *in the code.*

The stack stores its items in an items array. The tail field is an `AtomicInteger`, initially zero. The `push()` method reserves a slot by incrementing tail, and then stores the item at that location. Note that these two steps are not atomic: there is an interval after tail has been incremented but before the item has been stored in the array.

[1] https://github.com/ysc4231/hw05-linearizability-examples

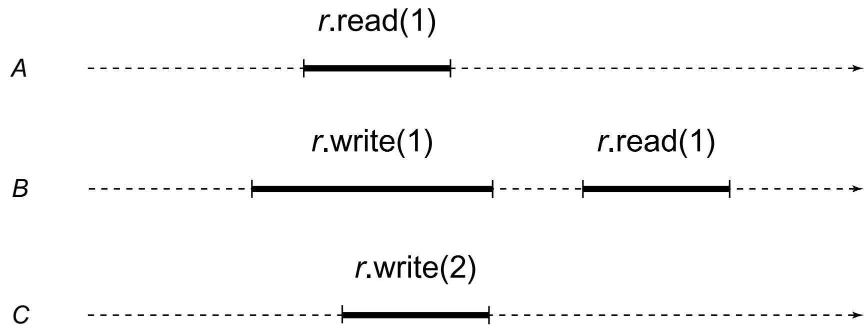

Figure 2: Second history for Problem 3.

The pop() method reads the value of tail and then traverses the array in descending order from the tail to slot zero. For each slot, it swaps None with the current contents, returning the first non-None item it finds. If all slots are None, the method returns None, indicating an empty stack.

- Give an example execution showing that the linearization point for push() cannot occur at the line

  ```
  val j = tail.getAndIncrement()
  ```

  Hint: give an execution where two push() calls are *not linearized* (*i.e.*, take effect) in the order they execute this line.
- Give another example execution showing that the linearization point for push() cannot occur at the line

  ```
  items.set(j, Some(x))
  ```

- Since these are the only two memory accesses in push(), we must conclude that push() has no single fixed linearization point. Does this mean push() is not linearizable? Can you then show an execution history of AGMStack that is not linearizable or argue that *all* histories produced by AGMStack are linearizable somehow (for the latter option you might want to split the set of alls histories into a finite set of "classes" of histories and argue for linearizability of each such class)?