

## Day-12

### Python DSA

#### Leetcode 15 3Sum

<https://leetcode.com/problems/3sum/description/>

#### Bruteforce

```
from typing import List

class Solution:

    def threeSum(self, nums:List[int])>-> List[List[int]]:

        my_set=set()

        n= len(nums)

        for i in range(n):

            for j in range(i+1, n):

                for k in range(j+1,n):

                    if nums[i] + nums[j] + nums[k] ==0:

                        temp= [nums[i], nums[j] , nums[k]]

                        temp.sort()

                        my_set.add(tuple(temp))

        ans = [list(item) for item in my_set]

        return ans

nums = [-1, 0, 1, 2, -1, -4]

sol= Solution()

print(sol.threeSum(nums))
```

#### Time Complexity:

The time complexity is  $O(n^3)$  due to the three nested loops, where  $n$  is the number of elements in the array. This makes the solution inefficient for large input sizes.

### Space Complexity:

The space complexity is  $O(k)$ , where  $k$  is the number of unique triplets found, stored in the set `my_set`.

### Better

```
from typing import List
```

```
class Solution:
```

```
    def threeSum(self, nums: List[int]) -> List[List[int]]:
```

```
        n = len(nums)
```

```
        result = set()
```

```
        for i in range(n):
```

```
            hashset = set()
```

```
            for j in range(i+1, n):
```

```
                third = -(nums[i] + nums[j])
```

```
                if third in hashset:
```

```
                    temp = [nums[i], nums[j], third]
```

```
                    temp.sort()
```

```
                    result.add(tuple(temp))
```

```
            hashset.add(nums[j])
```

```
        ans = list(result)
```

```
        return ans
```

```
nums = [-1, 0, 1, 2, -1, -4]
```

```
sol = Solution()
```

```
print(sol.threeSum(nums))
```

TC-O(N^2)

SC-O(N)

Optimal

```
class Solution:
    def threeSum(self, nums: List[int]) -> List[List[int]]:
        ans = []
        n = len(nums)
        nums.sort()
        for i in range(n):
            if i != 0 and nums[i] == nums[i - 1]:
                continue

            # moving the 2 pointers
            j = i + 1
            k = n - 1
            while j < k:
                total_sum = nums[i] + nums[j] + nums[k]
                if total_sum < 0:
                    j += 1
                elif total_sum > 0:
                    k -= 1
                else:
                    temp = [nums[i], nums[j], nums[k]]
                    ans.append(temp)
                    j += 1
                    k -= 1
                    # skip the duplicates if occurred
                    while j < k and nums[j] == nums[j - 1]:
                        j += 1
                    while j < k and nums[k] == nums[k + 1]:
                        k -= 1

            return ans
```

**Time Complexity:**

The time complexity is  $O(n \log n) + O(n^2)$ . Sorting the array takes  $O(n \log n)$ , and the two-pointer technique within the outer loop runs in  $O(n^2)$ .

**Space Complexity:**

The space complexity is  $O(1)$  for the two-pointer implementation (excluding the space required for the output).