

expectation \rightarrow

$$\langle k \rangle = \int_{k_{\min}}^{k_{\max}} k p(k) dk = c \int_{k_{\min}}^{k_{\max}} k k^{-\alpha} dk$$

$$= c \int_{k_{\min}}^{k_{\max}} k^{1-\alpha} dk$$

$$= c \left[\frac{k^{2-\alpha}}{2-\alpha} \right]_{k_{\min}}^{k_{\max}}$$

now let's say $\alpha > 2$. (diverges otherwise)

to calculate var \rightarrow

$$\langle k^2 \rangle = c \int_{k_{\min}}^{k_{\max}} k^2 p(k) dk = c \left[\frac{k^{3-\alpha}}{3-\alpha} \right]_{k_{\min}}^{k_{\max}}$$

now α needs to be greater than 3.

However,
we usually find that α is between 2 & 3.
 $\alpha \in [2, 3]$

$$\langle k \rangle \pm \frac{\sigma_k}{\sqrt{n}}$$

$n \rightarrow$ no. of nodes.
 $\sigma \rightarrow$ s.d.

\downarrow
confidence interval

to calculate σ , we need $\alpha > 3$. Hence we can't calculate σ when $\alpha \in [2, 3]$.

\therefore we have $\langle k \rangle$, but we cannot calculate the confidence.

\therefore we call such distribution to be a scale free distribution.

$$\langle k^t \rangle \approx c \left[\frac{k^{t+1-\alpha}}{t+1-\alpha} \right]_{k_{\min}}^{k_{\max}}$$

scale free distributions are usually self similar structures.

→ PARETO'S Power Law distribution
 80:20 → wealth distribution

FRAC TALS

$$p(k) = a' k^{-\alpha}$$

↓
 multiplying by a to
 change the scale.

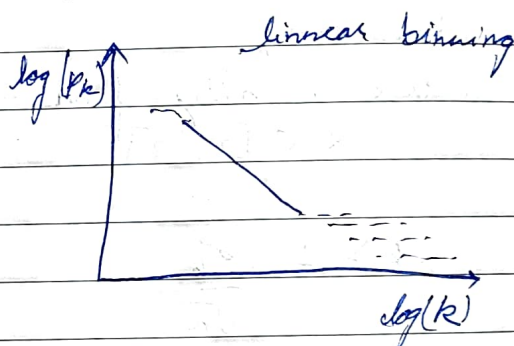
→ the distribution remains the
 same. Just multiplied by a
 constant.

80:20 → 20% of the nodes has 80% of the edges
 & 80% of the nodes has 20% of the edges.

Draw the distribution →

N_k
 ↪ no. of
 nodes with degree k
 $k = 1, \dots, k_{\max}$

$$p_k = \frac{N_k}{N} \rightarrow \text{total no. of nodes.}$$



then make a log log plot & observe if it's linear or
 not.

logarithmic binning →
 $b_0 = 1$
 $b_1 = 2, 3$
 $b_2 = 4, 5, 6, 7$
 degrees

typically we don't do logarithmic binning till degree
 10.

CDF → Rank - frequency method.

arrange degrees from highest to lowest.

calculate for exponential distribution.

k_i/n

10	1	1/7
9	2	2/7
8	3	3/7
7	4	4/7
6	5	
3	6	
2	7	

→ this is not just for power law

$$\int_{n_{1/2}}^{\infty} p(n) dn$$

$$= \frac{1}{2}$$

$n_{1/2}$ → divides the distribution into 2 equal parts.

$$\int_{n_{\min}}^{\infty} p(n) dn$$

$$\int_{n_{1/2}}^{\infty} n p(n) dn$$

$$= \frac{-\left(\frac{\alpha-2}{\alpha-1}\right)}{2}$$

→ for a power law distribution

$$\int_{n_{\min}}^{\infty} n p(n) dn$$

$k_{\max} \sim f(N)$ → how does k_{\max} depend on the no. of nodes.

$$\int_{k_{\min}}^{\infty} p(k) dk = \frac{1}{N}$$

$$C = (\alpha-1) k_{\min}^{\alpha-1}$$

$$C \int_{k_{\min}}^{\infty} k^{-\alpha} dk = \frac{C k_{\min}^{1-\alpha}}{\alpha-1} = \frac{1}{N}$$

$$k_{\max} = k_{\min} N^{-(1/(1-\alpha))}$$

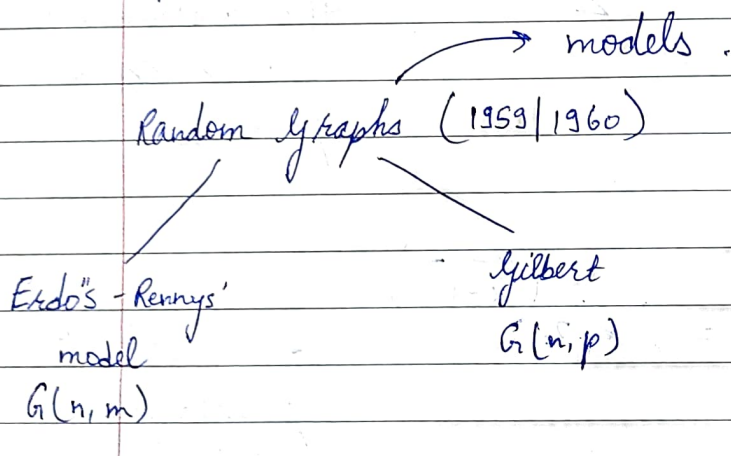
$$k_{\max} = k_{\min} N^{(1/(\alpha-1))}$$

$\alpha > 1$

robust properties \rightarrow properties that do not change on changing the network distribution size.

\rightarrow these properties give the insights on the distribution.

- hubs \rightarrow power-law degree distribution
 - small world \rightarrow short avg. path length
 - high clustering
 - sparse networks.
- Rule .



$n \rightarrow$ no. of nodes
 $m \rightarrow$ no. of links.
 $p \rightarrow$ probability of connecting 2 nodes

in random graphs, one thinks of averages.

~~to~~ calculating those averages is a little difficult for a $G(n, m)$ model as compared to a $G(n, p)$ model.

for a $G(n, p)$ model, the no. of links aren't going to be the same for different configurations of the network.

$G(n, p)$ → probability of connecting node pairs

average no. of links → $\langle m \rangle$

$$\langle m \rangle = {}^n C_2 \cdot p = \binom{n}{2} p$$

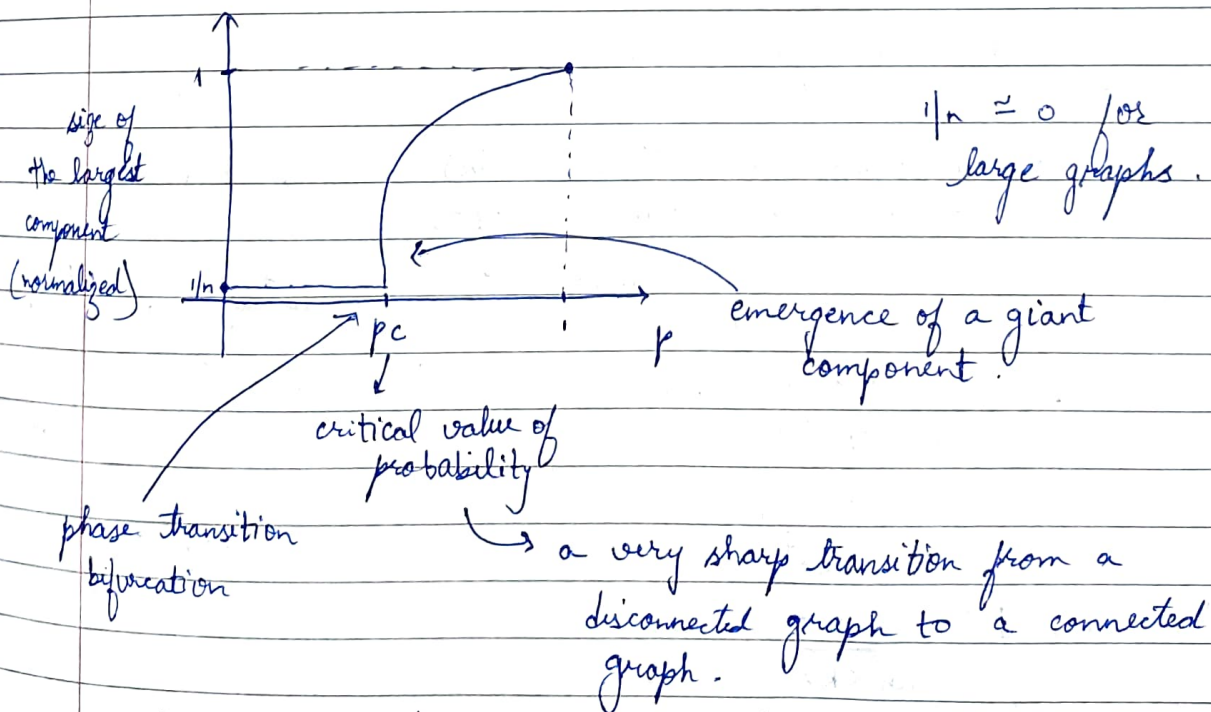
average degree → $\langle k \rangle$

$$\langle k \rangle = \frac{2m}{n} = \frac{2 \langle m \rangle}{n} = (n-1)p \approx np$$

↓
for large graphs.

Density → f

$$f = p = \frac{\langle m \rangle}{\binom{n}{2}}$$



$$\langle k \rangle = 1 \text{ at } p_c$$

if $\langle k \rangle < 1$ → disconnected graph (mostly)
if $\langle k \rangle > 1$ → connected graph (mostly)

sparse connected graphs \rightarrow usually trees
 \hookrightarrow avg degree close to 1.

$$\langle k \rangle \approx 1$$

$$\langle k \rangle \approx np$$

$$p = \frac{\langle k \rangle}{n}$$

$p \ll 1 \rightarrow$ tree kind of a structure
 \downarrow

clustering coefficient close to 0.

for a random graph \rightarrow clustering coefficient $\approx p$
 i.e. for $p \ll 1$, such graphs have a low
 value of clustering coefficient.

$$[C = p]$$

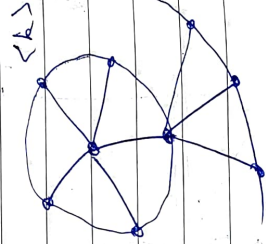
Average path length \rightarrow

$$\text{average degree} = \langle k \rangle$$

no. of neighbours on 1st level = $\langle k \rangle$

on 2nd level = $\langle k \rangle^2$

3rd level = $\langle k \rangle^3$



$$\therefore n = 1 + \langle k \rangle + \langle k \rangle^2 + \dots + \langle k \rangle^d$$

$$= \frac{\langle k \rangle^{d+1} - 1}{\langle k \rangle - 1}$$

$$\approx \langle k \rangle^d$$

$$d \approx \frac{\ln n}{\ln \langle k \rangle}$$

$$d = 3$$

Dunbar's no. = 75 \rightarrow avg. no. of people a person knows in their lifetime.
 let's work with 100.

$$\therefore n = 100^d$$

Degree distribution \rightarrow

$$P(k) = p_k = \binom{n-1}{k} p^k (1-p)^{n-1-k}$$

\downarrow
 k neighbours out of $n-1$ possible neighbours
 (Binomial distribution)

$$P(k) p_k \lim_{n \rightarrow \infty} \rightarrow \langle k \rangle = np$$

$$p_k \approx \frac{\langle k \rangle^k e^{-\langle k \rangle}}{k!} \rightarrow \text{poisson distribution}$$

$$\text{mean} = \langle k \rangle$$

$$\text{variance} = \langle k \rangle$$

the short avg. path length is obtained because of the randomness in a graph & not because of hubs.

clustering, degree distribution & power law behaviour isn't just because of randomness. It depicts the kind of graph we're working with.

if we try to make things regular, let's say increase the clustering coefficient, we might end up increasing the avg. path length.

∴ if we combine the idea of regular & random graphs, we might end up with a network having a high clustering coeff. & a short avg. path length.

small world phenomenon

collective phenomenon

science

waltz & Strogatz

→ make a graph with a lot of triangles, high clustering.

form a bridge

remove rest of the edges at random.