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Hands-On 6.

$(n) O + n \log n = (n) T$

Q3) Mathematically derive the average runtime complexity of the non-random pivot version of quicksort.

$(n) O + n \log n = (n) T$

Let) For array of size n ,

$(n) O + n \log n = (n) T$

recurrence relation for $T(n)$

$$1) T(n) = T(k) + T(n-k-1) + O(n)$$

k - Number of elements in left sub-array

$O(n)$ - Time partition of array.

$k-1$

$$2) T(n) = \frac{1}{n} \sum_{k=0}^{n-1} T(k) + T(n-k-1) + O(n)$$

Now, Since left & right part follow symmetry

$$3) T(n) = \frac{2}{n} \sum_{k=0}^{n-1} T(k) + O(n)$$

Let's assume $S(n) = \sum_{k=0}^{n-1} T(k)$

$$\begin{aligned} \therefore S(n) &= \int_0^n n \cdot \log n \cdot dn = \frac{n^2 \log n}{2} - \frac{n^2}{4} \\ &= \frac{n^2 \log n}{2} \end{aligned}$$

$$\therefore T(n) = \frac{2}{n} \cdot n^2 \log n + O(n)$$

$$\approx n \log n + O(n)$$

$$\Rightarrow T(n) = O(n \log n)$$

$$(n)O + (1-x-x)T + (n)T = (n)T \quad |$$

$$(n)O + (1-x-x)T + (n)T \geq \frac{1}{1-x} = (n)T \quad |$$

$$(n)O + (n)T \geq \frac{1}{1-x} = (n)T \quad |$$

$$(n)T \geq \frac{1}{1-x} = (n)T \quad |$$

$$-n - n \log n = -n \log n \quad |$$

$$\frac{1}{1-x}$$