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Hands-On 4

Problem : Fibonacci:

Call Stack:

$\text{fib}(5) \rightarrow \text{fib}(4) \rightarrow \text{fib}(3) \rightarrow \text{fib}(2) \rightarrow \text{fib}(1) \rightarrow \text{fib}(0)$   
 $\rightarrow \text{fib}(1)$   
 $\rightarrow \text{fib}(2) \rightarrow \text{fib}(1) \rightarrow \text{fib}(0)$   
 $\rightarrow \text{fib}(3) \rightarrow \text{fib}(2) \rightarrow \text{fib}(1) \rightarrow \text{fib}(0)$   
 $\rightarrow \text{fib}(1)$

→ Problem 1:

Q2.) Time Complexity

Time Complexity for number of levels is

$$O(\log M) \text{ of } (N)$$

For the merging of array is  $O(M)$

$$\begin{aligned} \text{Total time complexity} &= O(\log M) \cdot O(M) \\ &= O(M \log M) \end{aligned}$$

Consider all rows and columns

$$T(K, N) = O((K \cdot N) \cdot \log(K \cdot N))$$



Q3) We can improve the algorithm by implementing min-heap with optimized priority queues and libraries like `heapq`. This allows to store only necessary data and minimize memory overhead. Multi-threading can also be used as one of the alternatives.

→ Problem 2:

Q2.) Time complexity:

If array is sorted:

$O(N) \rightarrow$  for input

$O(N) \rightarrow$  for removing duplicates

$O(N)$  combined:  $O(N)$

If array is unsorted:

$O(N) \rightarrow$  for input

$O(N \log N) \rightarrow$  for sorting



$O(N) \rightarrow$  for removing duplicates

$\therefore$  Combined :  $O(N \log N)$

Q3-) Approach is already efficient. If the array is sorted, it takes time complexity of  $O(N)$ .