

تقسیم و غلبه - بخش سوم



Divide and Conquer III

پارادایم تقسیم و غلبه



Divide the problem into a number of subproblems that are smaller instances of the same problem.

Conquer the subproblems by solving them recursively. If the subproblem sizes are small enough, however, just solve the subproblems in a straightforward manner.

Combine the solutions to the subproblems into the solution for the original problem.

- Tournament Scheduling
- Convex Hull

برنامه ریزی تورنمنت ←
پوشش محدب ←

زمان بندی تورنمنت



- In a tournament with N teams,
 - هرف - each team must play against other teams through $(N-1)$ matches,
 - فید - one team can play at most one match per day.
- 1 • How many days a tournament of N teams can be fulfilled?
- 2 • What is the minimum days an N -tournament can be fulfilled?
- 3 • How an N -tournament is scheduled in a optimal way?

زمان بندی تورنمنت



- In a *tournament* with N teams,
 - each team must play against other teams through $(N-1)$ matches,
 - one team can play at most one match per day.
- How many days a tournament of N teams can be fulfilled?
- What is the minimum days an N -tournament can be fulfilled?
- How an N -tournament is scheduled in a optimal way?
- *First of all, how to represent a scheduling timetable?*

اینکه شکل هارو
جواب را

چگونه
قرار داد کنیم

$N=4$

Day 1	
1	2
3	4

Day 2	
1	3
2	4

Day 3	
1	4
2	3

team	Day 1	Day 2	Day 3
1	2	3	4
2	1	4	3
3	4	1	2
4	3	2	1

زمان بندی تورنمنت



- What is the minimum days an N -tournament can be fulfilled?
 - $(N-1)$ days, if N is even
 - N days, if N is odd

team	Day 1	Day 2	Day 3
1	2	3	—
2	1	—	3
3	—	1	2

$N=3$

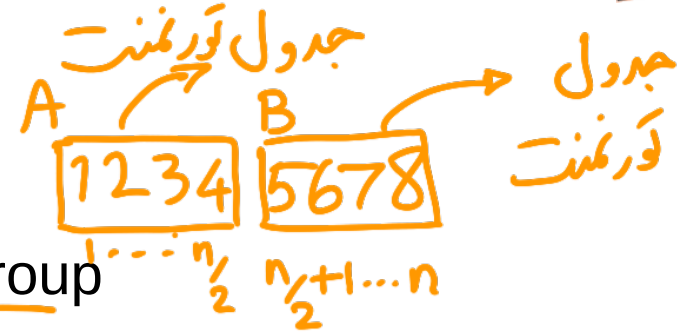
team	Day 1	Day 2	Day 3
1	2	3	4
2	1	4	3
3	4	1	2
4	3	2	1

$N=4$



یک الگوریتم تقسیم و غلبه برای زمان بندی تورنمنت

- Let $N=2^k$, (این فرض را بعداً کناری می‌گذاریم)
- ✓ Divide teams into two groups of $N/2$
- ✓ Schedule a $(N/2)$ -tournament for each group
 - Only one recursive call for problem of size $(N/2)$
 - Duplicate the timetable for $(N/2)$ and relabel matching teams, i.e add $N/2$ to the team number
- ✓ Schedule $(N/2)$ matches between pairs of two groups
 - Use the ``soccer hand-shaking`` fashion



$N=2$

team	Day 1
1	2
2	1

$N=4$

	team	Day 1	Day 2	Day 3
Group A	1	2		
	2	1		
Group B	3	2		
	4	1		



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- Schedule $(N/2)$ matches between pairs of two groups
 - Use the ``soccer hand-shaking`` fashion

$N=2$

team	Day 1
1	2
2	1

$N=4$

	team	Day 1	Day 2	Day 3
<u>Group A</u>	1	2		
	2	1		
<u>Group B</u>	3	4		
	4	3		

انجام
تورنمنت های درون گروهی



یک الگوریتم تقسیم و غلبه برای زمان بندی تورنمنت

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- Schedule a $(N/2)$ -tournament for each group
 - Only one recursive call for problem of size $(N/2)$
 - Duplicate the timetable for $(N/2)$ and relabeled matching teams, i.e add $N/2$ to the team number
- Schedule $(N/2)$ matches between pairs of two groups
 - Use the ``soccer hand-shaking`` fashion

*Hand
shaking*

Group A	<u>1</u>	<u>2</u>
Group B	<u>3</u>	<u>4</u>

←

	team	Day 1	Day 2	Day 3
Group A	<u>1</u>	2	<u>3</u>	
	<u>2</u>	1	<u>4</u>	
Group B	<u>3</u>	4	<u>1</u>	
	<u>4</u>	3	<u>2</u>	



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- Let $N=2^k$,
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- Schedule a $(N/2)$ -tournament for each group
 - Only one recursive call for problem of size $(N/2)$,
 - Duplicate the timetable for $(N/2)$ and relabeled matching teams, i.e add $N/2$ to the team number.
- Schedule $(N/2)$ matches between pairs of two groups
 - Use the ``soccer hand-shaking`` fashion.

Hand
shaking

Group A	1	2	1	2
Group B	3	4	<u>4</u>	<u>3</u>



	team	Day 1	Day 2	Day 3
Group A	1	2	3	<u>4</u>
	2	1	4	<u>3</u>
Group B	3	4	1	<u>2</u>
	4	3	2	<u>1</u>

تورنمنت بین گروهی



یک الگوریتم تقسیم و غلبه برای زمان بندی تورنمنت

مثال $N=7$ تیم نخودی

- If N is odd,

- Solve the problem for $(N+1)$ -tournament,
- Remove the team with label N and replace its matches by ``rest``.

به تیم اضافه کنید $N=8$

- If N is even, but $N/2$ is odd,

مثال $N=6$

- Use the divide-and-conquer as for $N=2^k$,
- But schedule a match for two ``rest`` teams on a day.

$N=3$

team	Day 1	Day 2	Day 3
1	2	3	--
2	1	--	3
3	--	1	2

$N=6$

Group A

Group B

team	Day 1	Day 2	Day 3	Day 4	Day 5
1	2	3	--		
2	1	--	3		
3	--	1	2		
4	5	6	--		
5	4	--	6		
6	--	4	5		

+3



یک الگوریتم تقسیم و غلبه برای زمان بندی تورنمنت

- If N is odd,
 - Solve the problem for $(N+1)$ -tournament,
 - Remove the team with label N and replace its matches by ``rest``.
- If N is even, but $N/2$ is odd,
 - Use the divide-and-conquer as for $N=2^k$,
 - But schedule a match for two ``rest`` teams on a day.

✓ تورنمنت درون گروهی به علاوه یک مسابقه بین گروهی برای هر تیم

$N=3$

team	Day 1	Day 2	Day 3
1	2	3	--
2	1	--	3
3	--	1	2

$N=6$

Group A

Group B

team	Day 1	Day 2	Day 3	Day 4	Day 5
1	2	3	<u>4</u>		
2	1	<u>5</u>	3		
3	<u>6</u>	1	2		
4	5	6	<u>1</u>		
5	4	<u>2</u>	6		
6	<u>3</u>	4	5		



یک الگوریتم تقسیم و غلبه برای زمان بندی تورنمنت

- If N is odd,
 - Solve the problem for $(N+1)$ -tournament,
 - Remove the team with label N and replace its matches by ``rest``.
- If N is even, but $N/2$ is odd,
 - Use the divide-and-conquer as for $N=2^k$,
 - But schedule a match for two ``rest`` teams on a day.

Hand shaking

Group A	1	2	3
Group B	4	5	6

← ignore

Group A	1	2	3
Group B	5	6	4

✓

Group A	1	2	3
Group B	6	4	5

✓

$N=6$

team	Day 1	Day 2	Day 3	Day 4	Day 5
1	2	3	4	5	6
2	1	5	3	6	4
3	6	1	2	4	5
4	5	6	1	3	2
5	4	2	6	1	3
6	3	4	5	2	1

↓ ↓

آنالیز الگوریتم تقسیم و غلبه برای حل مسئله زمانبندی تورنمنت

فرض کنید $T(n)$ نشان دهنده هزینه زمان اجرای الگوریتم تقسیم و غلبه

برای حل تورنمنت n بازی است.

$$T(n) = \begin{cases} T(n/2) + \Theta(n^2) & \text{if } n \text{ is even} \\ T(n+1) + \Theta(n) & \text{if } n \text{ is odd} \end{cases}$$

$$\Rightarrow T(n) = T\left(\left\lfloor \frac{n+1}{2} \right\rfloor\right) + \Theta(n^2)$$

$$\Rightarrow T(n) = \Theta(n^2) \quad \checkmark$$

حالت سوم قضیه
اساسی