

CSC 480: Artificial Intelligence I: 2024 Spring, Assignment #2

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Overview:

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Purpose:

To consider:

- Constraint Satisfaction Problems
- Logic
- Logic Programming

I used:

- [Python3 for Problem 1](#)
- [Prolog for Problem 3](#)

Assignment:

0. Free points! (10 Points)

(Because I am such a *nice guy!* :)

1. Constraint Satisfaction Problems (30 Points)

- Write a program to solve the 8-queens problem systematically, with backtracking.

- Now, write another program to solve the 8-queens problem by placing a queen down in her own row, but in a random column. Find a queen with the most conflicts, and randomly assign her a different column.
- Which program did *less* work to solve the 8-queens problem?

Answer

- **The random column placement method usually does less work.**
- Time Complexity:
 - Backtracking: $O(N!)$
 - Random: This complexity is typically polynomial - In the worst case, it might need to run infinite adjustments, but in general, it is usually faster than $O(N!)$. The specific complexity is hard to define strictly but is typically polynomial in nature.
- Space Complexity:
 - Backtracking: $O(N^2)$
 - Random: $O(N)$
- **In conclusion, the random column placement method (`problem1_random.py`) usually does less work to solve the 8-queens problem.**

Code & Output

```

1 # problem1_backtracking.py
2 # python3 ./src/problem1_backtracking.py
3 # Time complexity:  $O(N!)$ 
4 # Speed Complexity:  $O(N^2)$ 
5
6 def print_board(board):
7     for row in board:
8         print(" ".join(str(col) for col in row))
9     print("\n")
10
11 def is_safe(board, row, col):
12     for i in range(row):
13         if board[i][col] == 1:
14             return False
15     for i, j in zip(range(row, -1, -1), range(col, -1, -1)):
16         if board[i][j] == 1:
17             return False
18     for i, j in zip(range(row, -1, -1), range(col, len(board))):

```

```

19         if board[i][j] == 1:
20             return False
21     return True
22
23 def solve_n_queens(board, row):
24     if row >= len(board):
25         return True
26     for col in range(len(board)):
27         if is_safe(board, row, col):
28             board[row][col] = 1
29             if solve_n_queens(board, row + 1):
30                 return True
31             board[row][col] = 0
32     return False
33
34 def main():
35     n = 8
36     board = [[0 for _ in range(n)] for _ in range(n)]
37     if solve_n_queens(board, 0):
38         print_board(board)
39     else:
40         print("No solution found")
41
42 if __name__ == "__main__":
43     main()

```

```

● aydendeng@loop-eduroam-236-181 assignment2 % python3 ./src/problem1_backtracking.py
1 0 0 0 0 0 0 0
0 0 0 0 1 0 0 0
0 0 0 0 0 0 0 1
0 0 0 0 0 1 0 0
0 0 1 0 0 0 0 0
0 0 0 0 0 0 1 0
0 1 0 0 0 0 0 0
0 0 0 1 0 0 0 0

```

Python ▼

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```

1 # problem1_random.py
2 # python3 ./src/problem1_random.py
3 # Time complexity: This complexity is typically polynomial - In the
  worst case, it might need to run infinite adjustments, but in general,
  it is usually faster than  $O(N!)$ . The specific complexity is hard to
  define strictly but is typically polynomial in nature.
4 # Speed Complexity:  $O(N)$ 
5
6 import random

```

```

7
8 def print_board(board):
9     for row in board:
10         print(" ".join(str(col) for col in row))
11     print("\n")
12
13 def calculate_conflicts(board):
14     conflicts = [0] * len(board)
15     for i in range(len(board)):
16         for j in range(i + 1, len(board)):
17             if board[i] == board[j] or abs(board[i] - board[j]) == j -
i:
18                 conflicts[i] += 1
19                 conflicts[j] += 1
20     return conflicts
21
22 def solve_n_queens_random(n):
23     board = [random.randint(0, n - 1) for _ in range(n)]
24     while True:
25         conflicts = calculate_conflicts(board)
26         if max(conflicts) == 0:
27             return board
28         max_conflict = max(conflicts)
29         max_conflict_queens = [i for i, x in enumerate(conflicts) if x
== max_conflict]
30         queen = random.choice(max_conflict_queens)
31         new_position = random.randint(0, n - 1)
32         while new_position == board[queen]:
33             new_position = random.randint(0, n - 1)
34         board[queen] = new_position
35
36 def main():
37     n = 8
38     board = solve_n_queens_random(n)
39     solution = [[0 for _ in range(n)] for _ in range(n)]
40     for i in range(n):
41         solution[i][board[i]] = 1
42     print_board(solution)
43
44 if __name__ == "__main__":
45     main()
46

```

```

aydendeng@loop-eduroam-236-181 assignment2 % python3 ./src/problem1_random.py
0 1 0 0 0 0 0 0
0 0 0 0 0 0 1 0
0 0 0 0 1 0 0 0
0 0 0 0 0 0 0 1
1 0 0 0 0 0 0 0
0 0 0 1 0 0 0 0
0 0 0 0 0 1 0 0
0 0 1 0 0 0 0 0

```

2. Logic representation and reasoning (30 Points)

(NOTE: Below we are using mathematical notation. Variables are `lowercase`. Constants are `Uppercase`.)

Questions:

- a. Translate the sentences above into First Order Logic. This will be your knowledge base.
 - Every runner is health-conscious
 - Someone is health conscious if and only if that person sees the doctor regularly
 - Anyone who sees the doctor regularly and who has a health problem has that problem found early
 - Anyone who has a health problem and who has that problem found early gets that problem successfully treated
 - `mike` is a person and is a runner.
 - `mike` has the health problem `diabetes`.
- b. Translate your sentences to disjunctive form for resolution.
- c. Use resolution to show that `diabetes` problem of `mike` will be successfully treated.

Answer:

- a. Translate the sentences above into First Order Logic. This will be your knowledge base.
 - Every runner is health-conscious:
 - $\forall x(Runner(x) \rightarrow HealthConscious(x))$
 - Someone is health conscious if and only if that person sees the doctor regularly
 - $\forall x(HealthConscious(x) \leftrightarrow SeesDoctorRegularly(x))$
 - Anyone who sees the doctor regularly and who has a health problem has that problem found early

- $\forall x \forall y (\text{SeesDoctorRegularly}(x) \wedge \text{HasHealthProblem}(x,y) \rightarrow \text{ProblemFoundEarly}(x,y))$
- Anyone who has a health problem and who has that problem found early gets that problem successfully treated
 - $\forall x \forall y (\text{HasHealthProblem}(x,y) \wedge \text{ProblemFoundEarly}(x,y) \rightarrow \text{SuccessfullyTreated}(x,y))$
- `mike` is a person and is a runner.
 - $\text{Person}(\text{mike}) \wedge \text{Runner}(\text{mike})$
- `mike` has the health problem `diabetes` .
 - $\text{HasHealthProblem}(\text{mike}, \text{diabetes})$

b. Translate your sentences to disjunctive form for resolution.

- Every runner is health-conscious:
 - $\forall x (\text{Runner}(x) \rightarrow \text{HealthConscious}(x))$
 - Since
 - $P \rightarrow Q = \neg P \vee Q$
 - Then
 - $\forall x (\text{Runner}(x) \rightarrow \text{HealthConscious}(x)) = \forall x (\neg \text{Runner}(x) \vee \text{HealthConscious}(x))$
- Someone is health conscious if and only if that person sees the doctor regularly
 - $\forall x (\text{HealthConscious}(x) \leftrightarrow \text{SeesDoctorRegularly}(x))$
 - Since
 - $P \leftrightarrow Q = (P \rightarrow Q) \wedge (Q \rightarrow P), P \rightarrow Q = \neg P \vee Q$
 - Then
 - $\forall x (\text{HealthConscious}(x) \leftrightarrow \text{SeesDoctorRegularly}(x))$
 - $= \forall x ((\text{HealthConscious}(x) \rightarrow \text{SeesDoctorRegularly}(x)) \wedge (\text{SeesDoctorRegularly}(x) \rightarrow \text{HealthConscious}(x)))$
 - $= \forall x ((\neg \text{HealthConscious}(x) \vee \text{SeesDoctorRegularly}(x)) \wedge (\neg \text{SeesDoctorRegularly}(x) \vee \text{HealthConscious}(x)))$
 - $= (\neg \text{HealthConscious}(x) \vee \text{SeesDoctorRegularly}(x)) \wedge (\neg \text{SeesDoctorRegularly}(x) \vee \text{HealthConscious}(x))$
 - Hence
 - $\forall x (\text{HealthConscious}(x) \leftrightarrow \text{SeesDoctorRegularly}(x)) = (\neg \text{HealthConscious}(x) \vee \text{SeesDoctorRegularly}(x)) \wedge (\neg \text{SeesDoctorRegularly}(x) \vee \text{HealthConscious}(x))$

$\forall \text{HealthConscious}(x)$

- Anyone who sees the doctor regularly and who has a health problem has that problem found early
 - $\forall x \forall y (\text{SeesDoctorRegularly}(x) \wedge \text{HasHealthProblem}(x,y) \rightarrow \text{ProblemFoundEarly}(x,y))$
 - Since
 - $P \rightarrow Q = \neg P \vee Q$
 - De Morgan's Laws: $\neg(A \wedge B) \equiv \neg A \vee \neg B$
 - Then
 - $\neg(\text{SeesDoctorRegularly}(x) \wedge \text{HasHealthProblem}(x,y)) \equiv \neg \text{SeesDoctorRegularly}(x) \vee \neg \text{HasHealthProblem}(x,y)$
 - $\forall x \forall y (\neg \text{SeesDoctorRegularly}(x) \vee \neg \text{HasHealthProblem}(x,y) \vee \text{ProblemFoundEarly}(x,y))$
 - Hence
 - $\neg \text{SeesDoctorRegularly}(x) \vee \neg \text{HasHealthProblem}(x,y) \vee \text{ProblemFoundEarly}(x,y)$
 - Anyone who has a health problem and who has that problem found early gets that problem successfully treated
 - $\forall x \forall y (\text{HasHealthProblem}(x,y) \wedge \text{ProblemFoundEarly}(x,y) \rightarrow \text{SuccessfullyTreated}(x,y))$
 - Since
 - $P \rightarrow Q = \neg P \vee Q$
 - De Morgan's Laws: $\neg(A \wedge B) \equiv \neg A \vee \neg B$
 - Then
 - $\neg(\text{HasHealthProblem}(x,y) \wedge \text{ProblemFoundEarly}(x,y)) \equiv \neg \text{HasHealthProblem}(x,y) \vee \neg \text{ProblemFoundEarly}(x,y)$
 - $\forall x \forall y (\neg \text{HasHealthProblem}(x,y) \vee \neg \text{ProblemFoundEarly}(x,y) \vee \text{SuccessfullyTreated}(x,y))$
 - Hence
 - $\neg \text{HasHealthProblem}(x,y) \vee \neg \text{ProblemFoundEarly}(x,y) \vee \text{SuccessfullyTreated}(x,y)$
 - mike is a person and is a runner.
 - $\text{Person}(\text{mike}) \wedge \text{Runner}(\text{mike})$
 - $\text{Person}(\text{mike})$
 - $\text{Runner}(\text{mike})$

- `mike` has the health problem `diabetes` .

- $HasHealthProblem(mike, diabetes)$

c. Use resolution to show that `diabetes` problem of `mike` will be successfully treated.

i. $Goal: SuccessfullyTreated(mike, diabetes)$

1. Since

a. $\forall x (Runner(x) \rightarrow HealthConscious(x))$

b. Then

i. $Runner(mike) \rightarrow HealthConscious(mike)$

ii. $\neg Runner(mike) \vee HealthConscious(mike)$

2. Since

a. $\forall x (HealthConscious(x) \leftrightarrow SeesDoctorRegularly(x))$

b. Then

i. $(HealthConscious(mike) \rightarrow SeesDoctorRegularly(mike)) \wedge (SeesDoctorRegularly(mike) \rightarrow HealthConscious(mike))$

ii. $(\neg HealthConscious(mike) \vee SeesDoctorRegularly(mike)) \wedge (\neg SeesDoctorRegularly(mike) \vee HealthConscious(mike))$

3. Since

a. $\forall x \forall y (SeesDoctorRegularly(x) \wedge HasHealthProblem(x, y) \rightarrow ProblemFoundEarly(x, y))$

b. Then

i. $SeesDoctorRegularly(mike) \wedge HasHealthProblem(mike, diabetes) \rightarrow ProblemFoundEarly(mike, diabetes)$

ii. $\neg SeesDoctorRegularly(mike) \vee \neg HasHealthProblem(mike, diabetes) \vee ProblemFoundEarly(mike, diabetes)$

4. Since

a. $\forall x \forall y (HasHealthProblem(x, y) \wedge ProblemFoundEarly(x, y) \rightarrow SuccessfullyTreated(x, y))$

b. Then

i. $HasHealthProblem(mike, diabetes) \wedge ProblemFoundEarly(mike, diabetes) \rightarrow SuccessfullyTreated(mike, diabetes)$

ii. $\neg HasHealthProblem(mike, diabetes) \vee \neg ProblemFoundEarly(mike, diabetes) \vee SuccessfullyTreated(mike, diabetes)$

5. We knew

- a. Runner(mike)
 - b. HasHealthProblem(mike,diabetes)
 - ii. Let's proof
 - 1. Since
 - a. Runner(mike) & \neg Runner(mike) \vee HealthConscious(mike)
 - b. Then
 - i. HealthConscious(mike)
 - 2. Since
 - a. HealthConscious(mike)
 - b. \neg HealthConscious(mike) \vee SeesDoctorRegularly(mike)
 - c. Then
 - i. SeesDoctorRegularly(mike)
 - 3. Since
 - a. SeesDoctorRegularly(mike)
 - b. HasHealthProblem(mike,diabetes)
 - c. \neg SeesDoctorRegularly(mike) \vee \neg HasHealthProblem(mike,diabetes) \vee ProblemFoundEarly(mike,diabetes)
 - d. Then
 - i. ProblemFoundEarly(mike,diabetes)
 - 4. Since
 - a. HasHealthProblem(mike,diabetes)
 - b. ProblemFoundEarly(mike,diabetes)
 - c. \neg HasHealthProblem(mike,diabetes) \vee \neg ProblemFoundEarly(mike,diabetes) \vee SuccessfullyTreated(mike,diabetes)
 - i. Then
 - 1. SuccessfullyTreated(mike,diabetes)
 - 5. Hence, the diabetes problem of mike will be successfully treated.
-

3. Prolog (30 Points)

(**NOTE:** Below we are using Prolog notation. Variables are `Uppercase` . Constants are `lowercase` .)

- In class I showed how 2 sorted-lists can be merged in a sorted fashion. That means we are so close to implementing `merge-sort` ! We will implement `merge-sort` below.

Example output:

```
1  mergeSort([],L).
2  L = []
3  mergeSort([2],L).
4  L = [2]
5  mergeSort([4,2,3,1],L).
6  L = [1, 2, 3, 4]
7  mergeSort([5,4,6,2,7,3,1,8],L).
8  L = [1, 2, 3, 4, 5, 6, 7, 8]
```

Our program will be implemented with 5 predicates:

- `mergeSort/2` : which in comes an unsorted list, out goes the sorted list.
- `merge/3` : which merges the
- `split2ways/3` :
- `insertIn1/5` and `insertIn2/5` :

So let us go!

Steps:

1. Finish `mergeSort/2` :

```
1 mergeSort([],[]) :- !.           # Base case when list is empty
2 mergeSort([X],[]) :- !.         # Base case when list has only 1
3
4 # YOUR CODE TO HANDLE GENERAL CASE
5 # (1) Split the in coming list into 2 sublists
6 # (2) Recursively sort both lists
7 # (3) Merge both lists into the output list
```

2. Here is all the code for `merge/3`.

(This is all of it, there is nothing else to add.)

```
1 merge([],L,L).
2 merge(L,[],L).
3 merge([F0|R0],[F1|R1],[F0|L]) :-
4     (F0 < F1),
5     merge(R0,[F1|R1],L).
6 merge([F0|R0],[F1|R1],[F1|L]) :-
7     (F0 >= F1),
8     merge([F0|R0],R1,L).
```

3. Here is all the code for `split2ways/3`.

(This is all of it, there is nothing else to add.)

```
1 split2ways(Original,R1,R2) :-
2     insertIn1(Original,[],[],R1,R2).
```

4. You must finish the predicates

- `insertIn1()` (`insertIn1/5`),
- `insertIn2()` (`insertIn2/5`),

Both have 5 arguments:

- The incoming list
- 2 lists which hold temporary lists that build the answers
- 2 lists which will hold the answers

Both predicates will have the same base case (`_facts_`) corresponding to when the list to split is empty. It is:

```
1 insertIn1([],R1,R2,R1,R2).
```

It means "When there are no items to split, the temporary lists `R1` and `R2` are also the answer lists."

Your job is to write the 2 rules that build the temporary lists:

```
1 # YOUR insertIn1 RULE HERE
2
3 insertIn1([],R1,R2,R1,R2).      # This is the base case
4
5 # YOUR insertIn2 RULE HERE
6
7 insertIn2([],R1,R2,R1,R2).      # This is the base case
```

What are the remaining 2 rules?

- Hints:

- Both should separate the list to split into a head and a tail (`[H|T]`).
- Both should build the a list a their respective position
- `insertIn1/5` should call `insertIn2/5`
- `insertIn2/5` should call `insertIn1/5`

Answer:

Code

```
1 mergeSort([], []) :- !.
2 mergeSort([X], [X]) :- !.
3 mergeSort(List, Sorted) :-
4     split2ways(List, Left, Right),
5     mergeSort(Left, SortedLeft),
6     mergeSort(Right, SortedRight),
7     merge(SortedLeft, SortedRight, Sorted).
8
9 merge([], L, L).
10 merge(L, [], L).
11 merge([F0|R0], [F1|R1], [F0|L]) :-
12     F0 < F1,
13     merge(R0, [F1|R1], L).
14 merge([F0|R0], [F1|R1], [F1|L]) :-
```

```

15     F0 >= F1,
16     merge([F0|R0], R1, L).
17
18 split2ways(Original, R1, R2) :-
19     insertIn1(Original, [], [], R1, R2).
20
21 insertIn1([], R1, R2, R1, R2).
22 insertIn1([H|T], R1, R2, FinalR1, FinalR2) :-
23     insertIn2(T, [H|R1], R2, FinalR1, FinalR2).
24
25 insertIn2([], R1, R2, R1, R2).
26 insertIn2([H|T], R1, R2, FinalR1, FinalR2) :-
27     insertIn1(T, R1, [H|R2], FinalR1, FinalR2).
28

```

Output

```

1  ?- [src/problem3_merge_sort].
2  true.
3
4  ?- mergeSort([], L).
5  L = [].
6
7  ?- mergeSort([2], L).
8  L = [2].
9
10 ?- mergeSort([4, 2, 3, 1], L).
11 L = [1, 2, 3, 4] .
12
13 ?- mergeSort([5, 4, 6, 2, 7, 3,
14             1, 8], L).
15 L = [1, 2, 3, 4, 5, 6, 7, 8]

```

```

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```

```

For online help and background, visit https://www.swi-prolog.org
For built-in help, use ?- help(Topic). or ?- apropos(Word).

```

```

?- [src/problem3_merge_sort].
true.

```

```

?- mergeSort([], L).
L = [].

```

```

?- mergeSort([2], L).
L = [2].

```

```

?- mergeSort([4, 2, 3, 1], L).
L = [1, 2, 3, 4] .

```

```

?- mergeSort([5, 4, 6, 2, 7, 3, 1, 8], L).
L = [1, 2, 3, 4, 5, 6, 7, 8] █

```