

# Covalence Design and Implementation

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## 1. Core Calculus

We begin by giving the core calculus for the covalence kernel, as formalized in Lean 4. This is a variant of extensional Martin-Löf Type Theory (MLTT) based on a single judgement,  $\Gamma \vdash a \equiv b : A$ , which reads “ $a$  is equal to  $b$  at type  $A$  in context  $\Gamma$ .” The rules and grammar given below correspond to those in the formalization, which is in the locally-nameless style [1].

### 1.1. Grammar

Terms  $a, A, \varphi$  consist of:

**Variables**  $x$

**Universes**  $\mathcal{U}_\ell$  parametrized by a *level*  $\ell \in \mathbb{N}$

**Unit and empty types**  $1_\ell, *_\ell, 0_\ell$

**Equations**  $a =_A b$

**Dependent functions**  $\Pi_\ell A.B, \lambda_\ell A.(b : B)$

**Applications**  $\text{app}_{A.B}(f, a)$

**Dependent pairs**  $\Sigma_\ell A.B, \text{mk}_{\ell A.B}(a, b)$

**Pair projections**  $\pi_{lA.B}(e), \pi_{rA.B}(e)$

**Dependent if-then-else**  $\text{ite}(\varphi, A, a, b)$

**Propositional truncations**  $\|A\|$

**Choice terms**  $\varepsilon A.\varphi$

**Natural numbers**  $\mathbb{N}, 0, s$

**Iteration**  $\text{rec}_{\mathbb{N}}(C, n, z, s)$

We introduce syntax sugar:

**Propositions**  $2 = \mathcal{U}_0$

**Booleans**  $\top := 1_0, \perp := 0_0$

**Functions**  $A \rightarrow B := \Pi A.B^\uparrow$

**Negations**  $\neg A := A \rightarrow \perp$

**Existential quantifiers**  $\exists A.\varphi := \|\Sigma A.\varphi\|$

### 1.2. Judgements

Judgement	Reading	Definition
$\Gamma \vdash a \equiv b : B$	“In $\Gamma$ , $a$ is equal to $b$ at type $A$ ”	Primitive
$\Gamma \text{ ok}$	“ $\Gamma$ is a well-formed context”	$\Gamma \vdash 0 \equiv 0 : \mathbb{N}$
$\Gamma \vdash a : A$	“In $\Gamma$ , $a$ has type $A$ ”	$\Gamma \vdash a \equiv a : A$

Judgement	Reading	Definition
$\Gamma \vdash A$	“In $\Gamma$ , $A$ is inhabited”	$\exists a. \Gamma \vdash a : A$

### 1.3. Rules

#### 1.3.1. Context well-formedness

$$\begin{array}{c}
\frac{}{\cdot \vdash 0 \equiv 0 : \mathbb{N}} \text{nil-ok} \quad \frac{\Gamma \vdash 0 \equiv 0 : \mathbb{N} \quad x \notin \Gamma \quad \Gamma \vdash A : \mathcal{U}_\ell}{\Gamma, x : A \vdash 0 \equiv 0 : \mathbb{N}} \text{cons-ok} \\
\\
\frac{\Gamma \text{ ok}}{\Gamma \vdash \mathcal{U}_\ell : \mathcal{U}_{\ell+1}} \mathcal{U} \quad \frac{\Gamma \text{ ok} \quad \Gamma(x) = A}{\Gamma \vdash x : A} \text{var}
\end{array}$$

#### 1.3.2. Congruence Rules

$$\begin{array}{c}
\frac{\Gamma \text{ ok}}{\Gamma \vdash \mathbf{1}_\ell : \mathcal{U}_\ell} \mathbf{1} \quad \frac{\Gamma \text{ ok}}{\Gamma \vdash *_\ell : \mathbf{1}_\ell} * \quad \frac{\Gamma \text{ ok}}{\Gamma \vdash \mathbf{0}_\ell : \mathcal{U}_\ell} \mathbf{0} \quad \frac{\Gamma \text{ ok}}{\Gamma \vdash \mathbb{N} : \mathcal{U}_1} \mathbb{N} \quad \frac{\Gamma \text{ ok}}{\Gamma \vdash \mathbf{s} : \mathbb{N} \rightarrow \mathbb{N}} \mathbf{s} \\
\\
\frac{\Gamma \vdash A \equiv A' : \mathcal{U}_\ell \quad \Gamma \vdash a \equiv a' : A \quad \Gamma \vdash b \equiv b' : A}{\Gamma \vdash (a =_A b) \equiv (a' =_{A'} b') : \mathbf{2}} \text{eqn} \\
\\
\frac{\Gamma \vdash A \equiv A' : \mathcal{U}_m \quad \Gamma, x : A \vdash B^x \equiv B'^x : \mathcal{U}_n}{\Gamma \vdash \Pi_{\text{imax}(m,n)} A.B \equiv \Pi_{\text{imax}(m,n)} A'.B' : \mathcal{U}_{\text{imax}(m,n)}} \Pi \\
\\
\frac{\Gamma \vdash A \equiv A' : \mathcal{U}_m \quad \Gamma, x : A \vdash B^x \equiv B'^x : \mathcal{U}_n \quad \Gamma \vdash f \equiv f' : \Pi A.B \quad \Gamma \vdash a \equiv a' : A}{\Gamma \vdash \text{app}_{A.B}(f, a) \equiv \text{app}_{A'.B'}(f', a') : B^a} \text{app} \\
\\
\frac{\Gamma \vdash A \equiv A' : \mathcal{U}_m \quad \Gamma, x : A \vdash B^x \equiv B'^x : \mathcal{U}_n \quad \Gamma, x : A \vdash b^x \equiv b'^x : B^x}{\Gamma \vdash \lambda_{\text{imax}(m,n)} A.(b : B) \equiv \lambda_\ell A'.(b' : B') : \Pi_{\text{imax}(m,n)} A.B} \lambda \\
\\
\frac{\Gamma \vdash A \equiv A' : \mathcal{U}_m \quad \Gamma, x : A \vdash B^x \equiv B'^x : \mathcal{U}_n}{\Gamma \vdash \Sigma_{m \sqcup n} A.B \equiv \Sigma_{m \sqcup n} A'.B' : \mathcal{U}_{m \sqcup n}} \Sigma \\
\\
\frac{\Gamma \vdash A \equiv A' : \mathcal{U}_m \quad \Gamma, x : A \vdash B^x \equiv B'^x : \mathcal{U}_n \quad \Gamma \vdash a \equiv a' : A \quad \Gamma \vdash b \equiv b' : B^a}{\Gamma \vdash \text{mk}_{\ell A.B}(a, b) \equiv \text{mk}_{\ell A'.B'}(a', b') : \Sigma_{m \sqcup n} A.B} \text{mk} \\
\\
\frac{\Gamma \vdash A \equiv A' : \mathcal{U}_m \quad \Gamma, x : A \vdash B^x \equiv B'^x : \mathcal{U}_n \quad \Gamma \vdash e \equiv e' : \Sigma_{m \sqcup n} A.B}{\Gamma \vdash \pi_{l A.B}(e) \equiv \pi_{l A'.B'}(e') : A} \pi_l \\
\\
\frac{\Gamma \vdash A \equiv A' : \mathcal{U}_m \quad \Gamma, x : A \vdash B^x \equiv B'^x : \mathcal{U}_n \quad \Gamma \vdash e \equiv e' : \Sigma_{m \sqcup n} A.B}{\Gamma \vdash \pi_{r A.B}(e) \equiv \pi_{r A'.B'}(e') : B^{\pi_{l A.B}(e)}} \pi_r \\
\\
\frac{\Gamma \vdash \varphi \equiv \varphi' : \mathbf{2} \quad \Gamma \vdash A \equiv A' : \mathcal{U}_\ell \quad \Gamma, x : \varphi \vdash a \equiv a' : A \quad x \notin \text{fv}(a) \quad \Gamma, y : \neg \varphi \vdash b \equiv b' : A \quad y \notin \text{fv}(b)}{\Gamma \vdash \text{ite}(\varphi, A, a, b) \equiv \text{ite}(\varphi', A', a', b') : A} \text{dite} \\
\\
\frac{\Gamma \vdash A \equiv A' : \mathcal{U}_\ell}{\Gamma \vdash \|A\| \equiv \|A'\| : \mathbf{2}} \text{tr} \quad \frac{\Gamma \vdash A \equiv A' : \mathcal{U}_\ell \quad \Gamma \vdash \|A\| \equiv \top : \mathbf{2} \quad \Gamma, x : A \vdash \varphi^x \equiv \varphi'^x : \mathbf{2}}{\Gamma \vdash \varepsilon A.\varphi \equiv \varepsilon A'.\varphi' : A} \varepsilon \\
\\
\frac{\Gamma, x : \mathbb{N} \vdash C^x \equiv C'^x : \mathcal{U}_\ell \quad \Gamma \vdash n \equiv n' : \mathbb{N} \quad \Gamma \vdash z \equiv z' : C^0 \quad \Gamma, x : \mathbb{N} \vdash s^x \equiv s'^x : C^x \rightarrow C^{\text{app}_{\mathbb{N}, \mathbb{N}}(s, x)}}{\Gamma \vdash \text{rec}_{\mathbb{N}}(C, n, z, s) \equiv \text{rec}_{\mathbb{N}}(C', n', z', s') : C^n} \text{rec}_{\mathbb{N}}
\end{array}$$

### 1.3.3. Equations

$$\begin{array}{c}
\frac{\Gamma \vdash A : \mathbf{2} \quad \Gamma \vdash a : A}{\Gamma \vdash a \equiv *_0 : \mathbf{2}} \text{!}_* \quad \frac{\Gamma \vdash a : \mathbf{0}_\ell}{\Gamma \vdash \top \equiv \perp : \mathbf{2}} \perp_0 \quad \frac{\Gamma \vdash a \equiv b : A}{\Gamma \vdash (a =_A b) \equiv \top : \mathbf{2}} \top_= \\
\\
\frac{\Gamma \vdash A : \mathcal{U}_m \quad \Gamma, x : A \vdash B^x : \mathcal{U}_n \quad \Gamma \vdash a : A \quad \Gamma, x : A \vdash b^x : B^x}{\Gamma \vdash \text{app}_{A.B}(\lambda_{\text{imax}(m,n)} A.(b : B), a) \equiv b^a : B^a} \beta \\
\\
\frac{\Gamma \vdash A : \mathcal{U}_m \quad \Gamma, x : A \vdash B^x : \mathcal{U}_n \quad \Gamma \vdash a : A \quad \Gamma \vdash b : B^a}{\Gamma \vdash \pi_{lA.B}^l(\text{mk}_{m \sqcup n A.B}(a, b)) \equiv a : A} \beta_{\pi_l} \\
\\
\frac{\Gamma \vdash A : \mathcal{U}_m \quad \Gamma, x : A \vdash B^x : \mathcal{U}_n \quad \Gamma \vdash a : A \quad \Gamma \vdash b : B^a}{\Gamma \vdash \pi_{rA.B}^r(\text{mk}_{m \sqcup n A.B}(a, b)) \equiv b : B^a} \beta_{\pi_r} \\
\\
\frac{\Gamma \vdash \varphi \equiv \top : \mathbf{2} \quad \Gamma \vdash A : \mathcal{U}_\ell \quad \Gamma, x : \varphi \vdash a : A \quad x \notin \text{fv}(a) \quad \Gamma, y : \neg \varphi \vdash b : A \quad y \notin \text{fv}(b)}{\Gamma \vdash \text{ite}(\varphi, A, a, b) \equiv a : A} \beta_\top \\
\\
\frac{\Gamma \vdash \varphi \equiv \perp : \mathbf{2} \quad \Gamma \vdash A : \mathcal{U}_\ell \quad \Gamma, x : \varphi \vdash a : A \quad x \notin \text{fv}(a) \quad \Gamma, y : \neg \varphi \vdash b : A \quad y \notin \text{fv}(b)}{\Gamma \vdash \text{ite}(\varphi, A, a, b) \equiv b : A} \beta_\perp \\
\\
\frac{\Gamma, x : \mathbb{N} \vdash C^x : \mathcal{U}_\ell \quad \Gamma \vdash z : C^0 \quad \Gamma, x : \mathbb{N} \vdash s^x : C^x \rightarrow C^{\text{app}_{\mathbb{N}.N}(s, x)}}{\Gamma \vdash \text{rec}_{\mathbb{N}}(C, 0, z, s) \equiv z : C^0} \text{rec}_{\mathbb{N}-0} \\
\\
\frac{\Gamma, x : \mathbb{N} \vdash C^x : \mathcal{U}_\ell \quad \Gamma \vdash n : \mathbb{N} \quad \Gamma \vdash z : C^0 \quad \Gamma, x : \mathbb{N} \vdash s^x : C^x \rightarrow C^{\text{app}_{\mathbb{N}.N}(s, x)}}{\Gamma \vdash \text{rec}_{\mathbb{N}}(C, \text{app}_{\mathbb{N}.N}(s, n), z, s) \equiv \text{app}_{C^n.C^{\text{app}_{\mathbb{N}.N}(s, n)}}(s, \text{rec}_{\mathbb{N}}(C, n, z, s)) : C^{\text{app}_{\mathbb{N}.N}(s, n)}} \text{rec}_{\mathbb{N}-s}
\end{array}$$

### 1.3.4. Axioms

$$\begin{array}{c}
\frac{\Gamma \vdash a =_A b}{\Gamma \vdash a \equiv b : A} \text{ext} \quad \frac{\Gamma \vdash A : \mathbf{2} \quad \Gamma \vdash B : \mathbf{2} \quad \Gamma \vdash A \rightarrow B \quad \Gamma \vdash B \rightarrow A}{\Gamma \vdash A \equiv B : \mathbf{2}} \text{propext} \\
\\
\frac{\Gamma \vdash A : \mathcal{U}_m \quad \Gamma, x : A \vdash B^x : \mathcal{U}_n \quad \Gamma \vdash e : \Sigma_{m \sqcup n} A.B}{\Gamma \vdash \text{mk}_{m \sqcup n A.B}(\pi_{m \sqcup n A.B}^l(e), \pi_{m \sqcup n A.B}^r(e)) \equiv e : \Sigma_{m \sqcup n} A.B} \text{ext}_\Sigma
\end{array}$$

### 1.3.5. Closure

$$\begin{array}{c}
\frac{\Gamma \vdash a \equiv b : A \quad \Gamma \vdash b \equiv c : A}{\Gamma \vdash a \equiv c : A} \text{trans} \quad \frac{\Gamma \vdash a \equiv b : A}{\Gamma \vdash b \equiv a : A} \text{symm} \\
\\
\frac{\Gamma \vdash a \equiv b : A \quad \Gamma \vdash A \equiv B : \mathcal{U}_\ell}{\Gamma \vdash a \equiv b : B} \text{cast}
\end{array}$$

## Bibliography

- [1] A. Charguéraud, “The Locally Nameless Representation,” *J. Autom. Reason.*, vol. 49, no. 3, pp. 363–408, 2012, doi: [10.1007/S10817-011-9225-2](https://doi.org/10.1007/S10817-011-9225-2).