

## Assignment 1 $\Leftrightarrow$ CS890D0

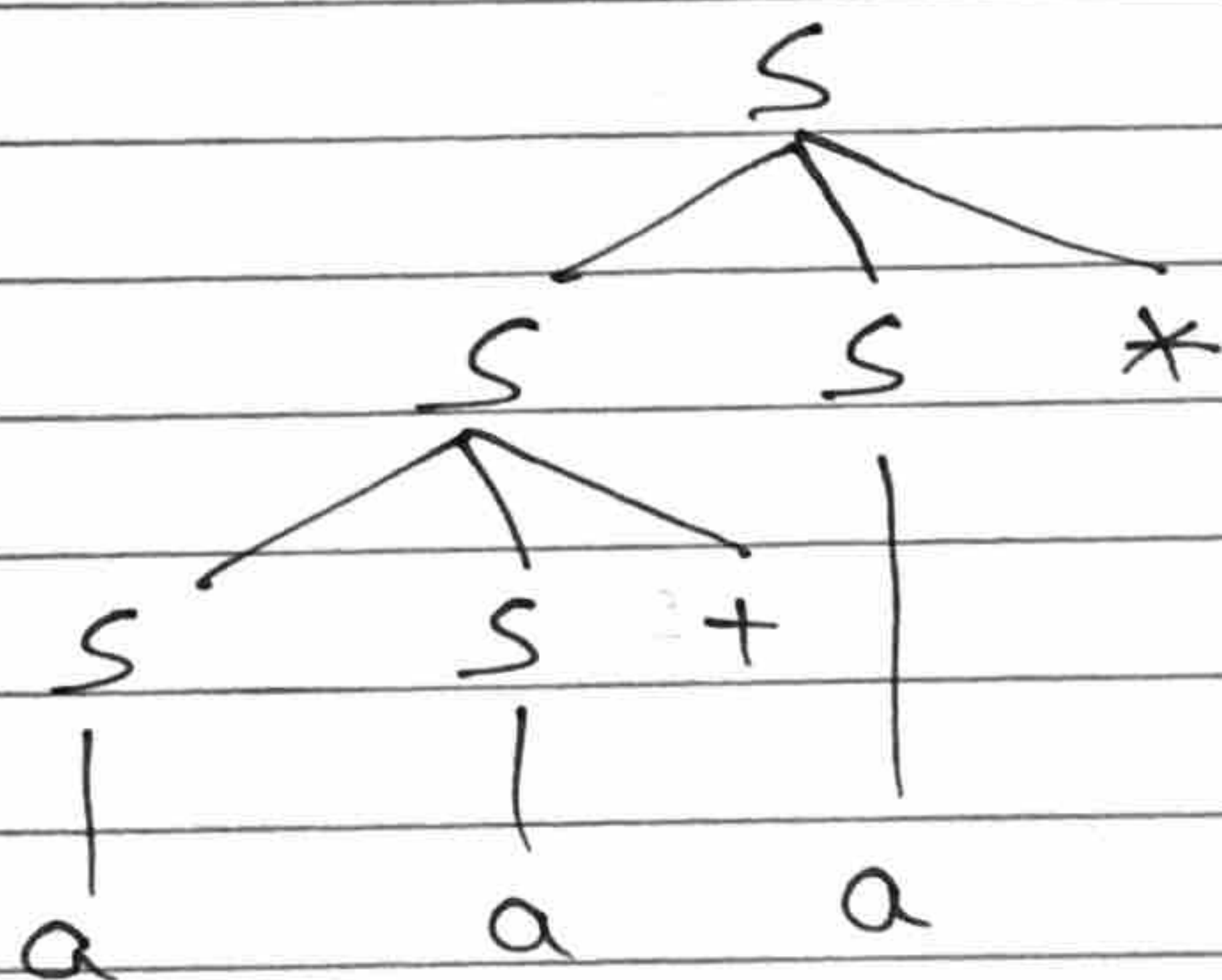
Exercise: 2.2.1 Consider the context free grammar

$$S \rightarrow SS+ \mid SS* \mid a$$

- a] Show how string  $aa+a*$  can be generated by this grammar.

$$\begin{aligned} S &\rightarrow \underline{S}S* \\ &\rightarrow \underline{S}S+S* \\ &\rightarrow a\underline{S}+S* \\ &\rightarrow aa+\underline{S}* \\ &\rightarrow aa+a* \end{aligned}$$

- b] Construct parse tree for this string



- c] What language does this grammar generate? Justify your answer.

$\Rightarrow$  This grammar generates postfix expressions consisting of addition and multiplication. Digits are represented by 'a'.  
 $\Rightarrow$  'a' is only the ~~non~~ terminal symbol in the grammar



So all the 'S' will be replaced by a  
at the end we will be left of a series  
or expression consisting of a's and  
addition and \* multiplication signs, in  
a postfix notation.

### Exercise 2.2.2

What language is generated by the following  
grammars? In each case justify your  
answer.

a)  $S \rightarrow 0S1 \mid 01$

This grammar will have equal no. of 0's  
and 1's

$$L = \{0^n 1^n \mid n \geq 1\}$$

So the smallest string is '01' which has  
one 0 & one 1.

Some strings generated by the grammar

01, 0011, <u>000111</u> , .....		$S \rightarrow 0S1$
↘		$\rightarrow 00S11$
		$\rightarrow 000111$



$$b) \quad S \rightarrow +SS \mid -SS \mid a$$

$L = \{ \text{Prefix expression containing addition and subtraction only with only } a \text{ as a nonterminal} \}$

So if we want to generate  $-+aaa$

$$\begin{aligned} S &\rightarrow -\underline{SS} \\ &\rightarrow -+SSS \\ &\rightarrow -+aSS \\ &\rightarrow -+aaS \\ &\rightarrow -+aaa \end{aligned}$$

So this grammar generates string with a's having addition and subtraction signs in prefix notation.

$$c) \quad S \rightarrow S(S)S \mid \epsilon$$

$L = \{ \text{All properly balanced parentheses which also includes } \epsilon \}$

Some of the valid strings:-

$() , (()) , ()() , (((()))) \dots$

So for every opening parentheses there is a closing parentheses and the string should not start with closing parentheses.



$$d] S \rightarrow aSbS \mid bSaS \mid \epsilon$$

$L = \{ \text{contains equal number of a's and b's} \\ \text{any order and contains epsilon} \}$

Some of the strings produced are

ab, ba, ababba, ...



$S \rightarrow aSbS$ $\rightarrow abSaSbS$ $\rightarrow ababSbS$ $\rightarrow ababS$ $\rightarrow ababba$	$S \rightarrow aSbS$ $\rightarrow abSaSbS$ $\rightarrow ababSbS$ $\rightarrow ababS$ $\rightarrow ababba$
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As we can see we have equal no. of a's and b's in every string produced by grammar.

$$e] S \rightarrow a \mid S+S \mid SS \mid S^* \mid (S)$$

$L = \{ \text{Regular expression with terminal } a \}$

$SS$ : concatenation

$S+S$ : One or more occurrence

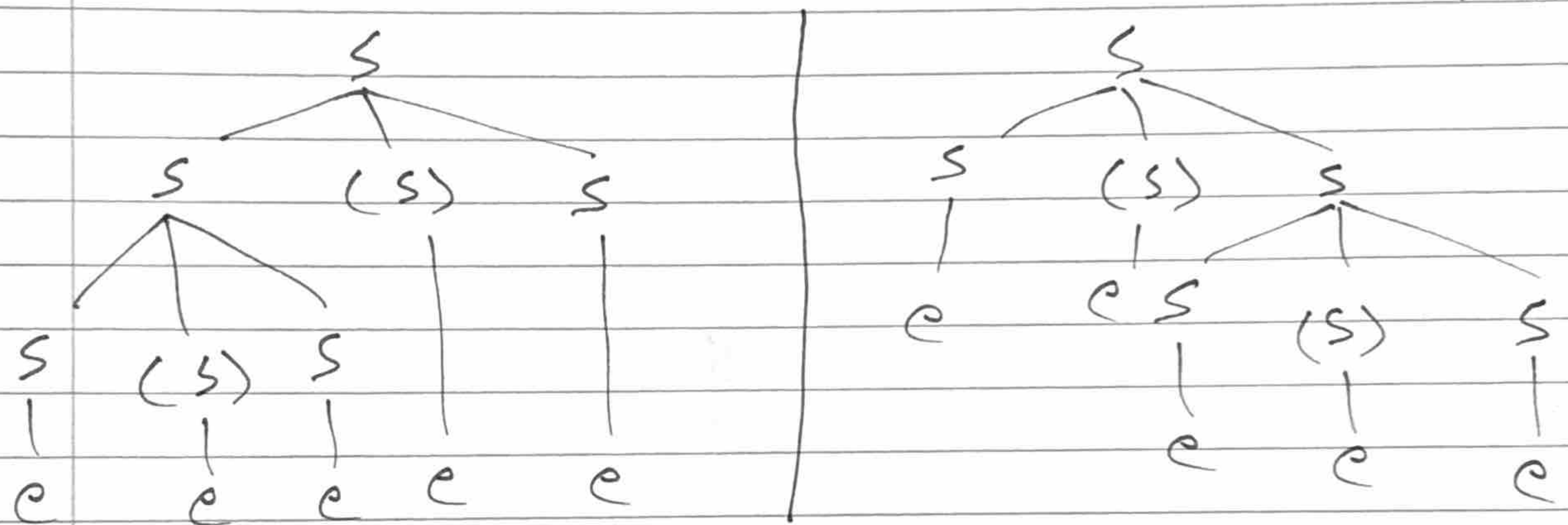
$S^*$ : zero or more occurrence

### Exercise 2.2.3

Which of the grammars in exercise 2.2.2 are ambiguous?

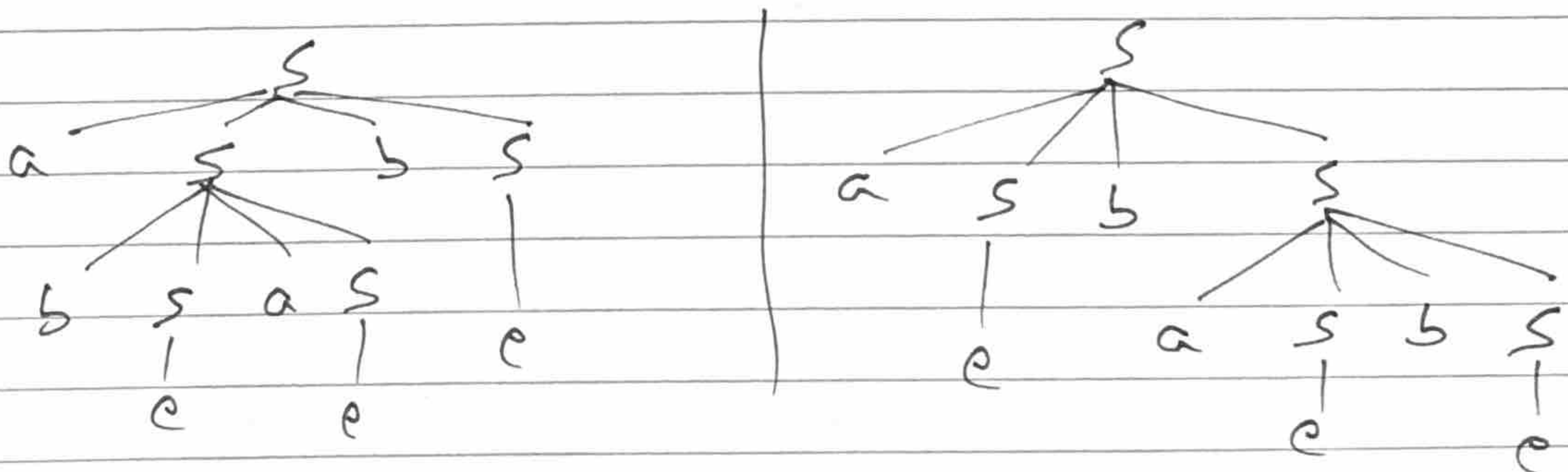
c]  $S \rightarrow S(S)S \mid \epsilon$  Ambiguous

To generate:  $()()$  we have two parse trees



d]  $S \rightarrow aSbS \mid bSaS \mid \epsilon$  Ambiguous

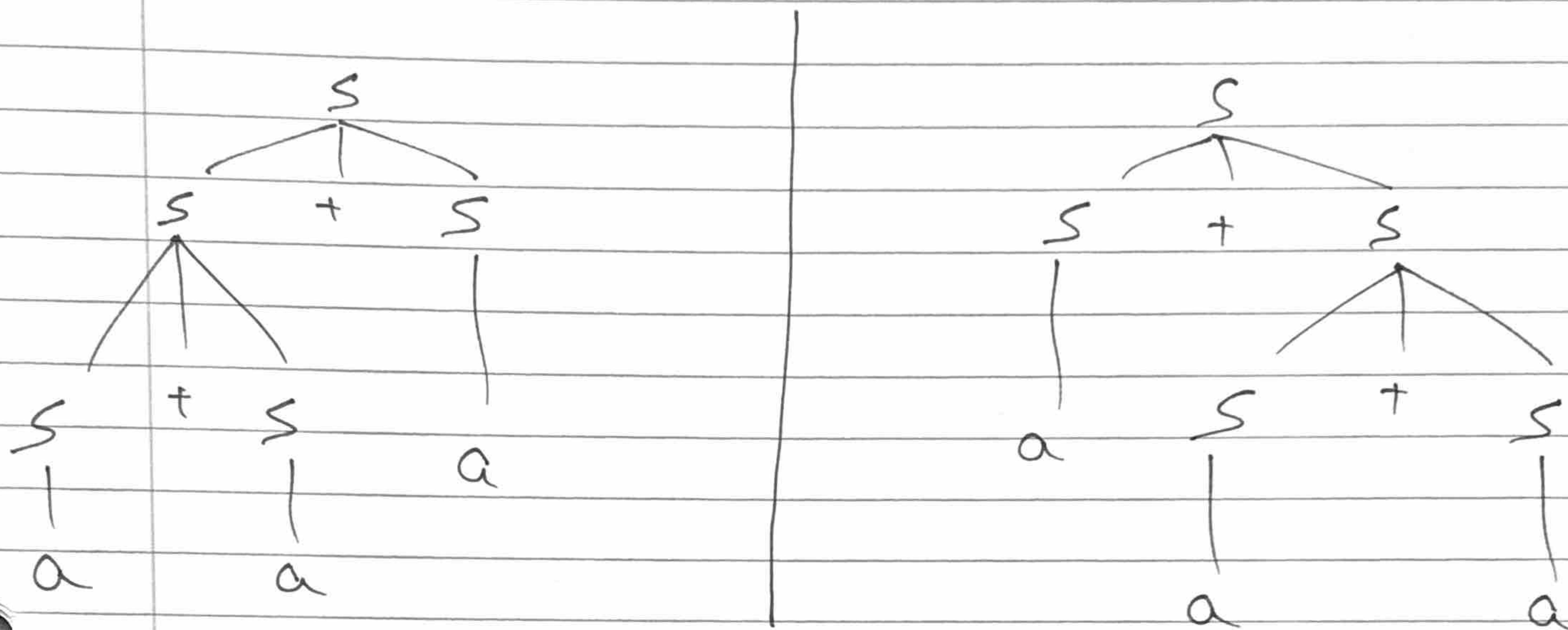
To generate:  $abab$  we have two parse trees





e]  $S \rightarrow a \mid S+S \mid SS \mid S^* \mid (S) \rightarrow \text{Ambiguous}$

To generate  $a+a+a$  we have two parse trees.



Exercise:- 2.2.4

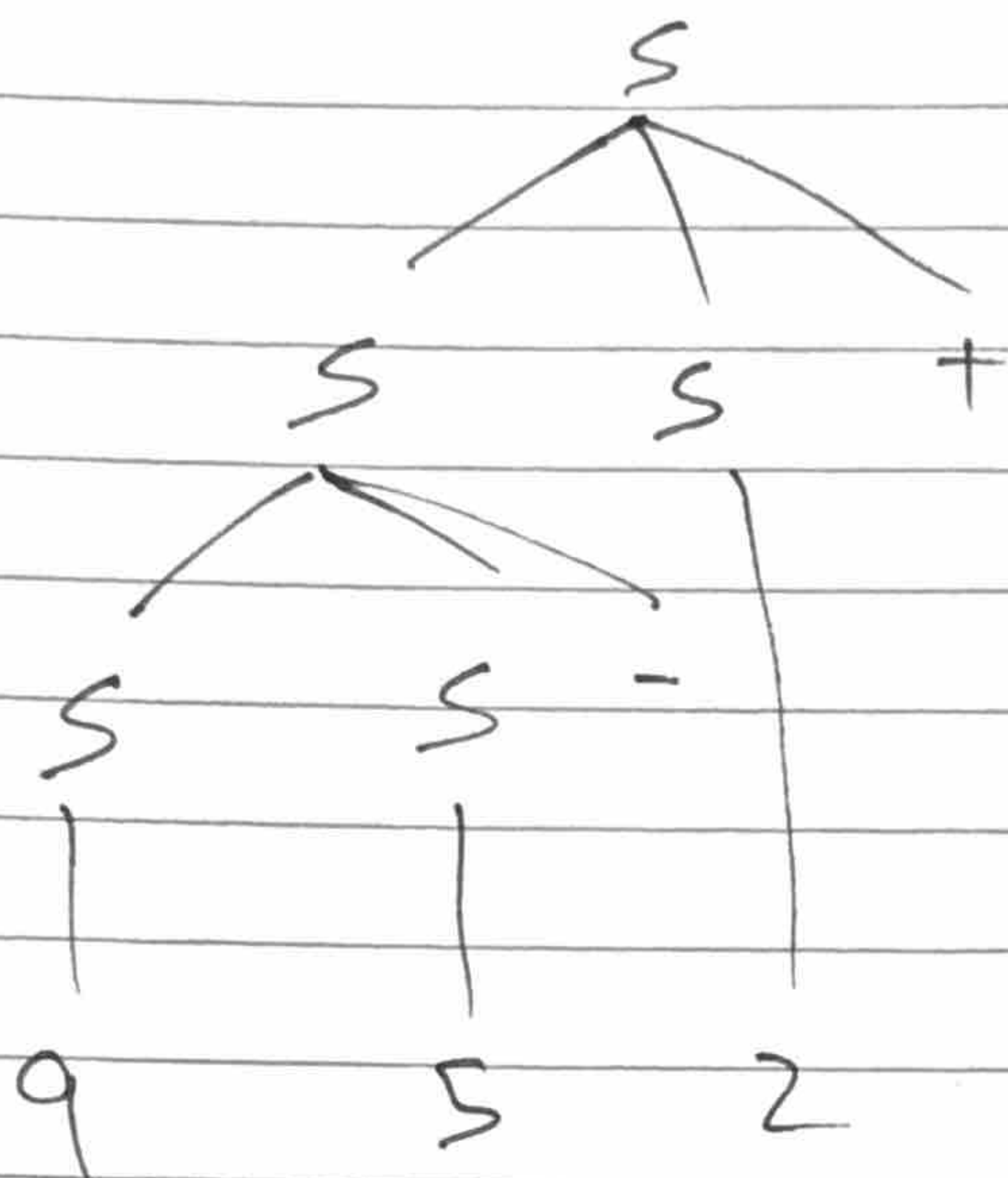
Construct unambiguous context-free grammar for each of following languages. In each case show that your grammar is correct.

a] Arithmetic expressions in postfix notation.

$S \rightarrow SS+ \mid SS- \mid SS* \mid SS/ \mid id$   
 $id \rightarrow 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$

If we want to show  $95-2+$  then we will have only one parse tree



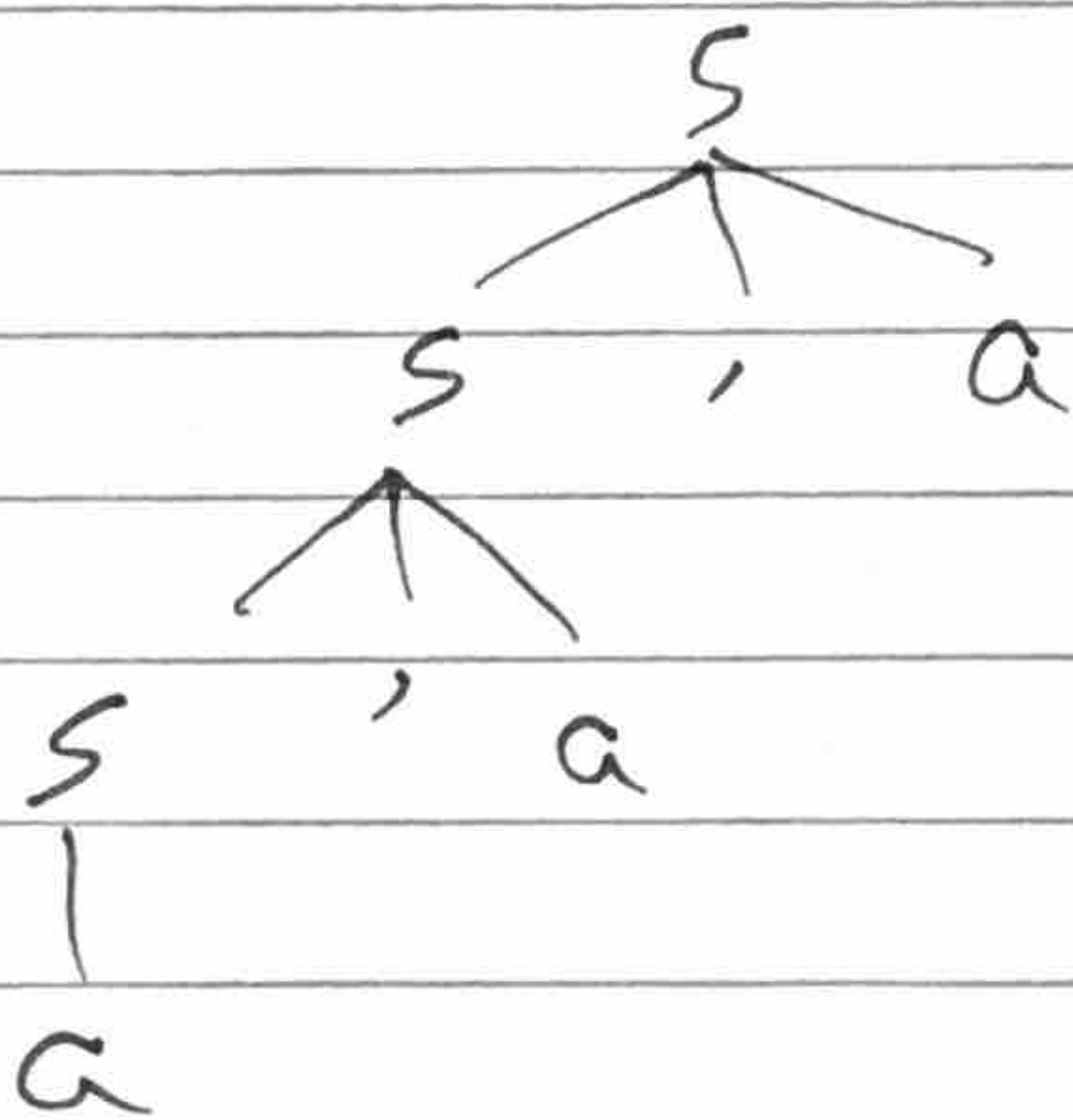


The above parse tree shows that the grammar is correct and unambiguous.

- 5] Left associative ~~if~~ list of identifiers separated by commas.

$$S \rightarrow S, a \mid a$$

To generate  $a, a, a$  we have one parse tree



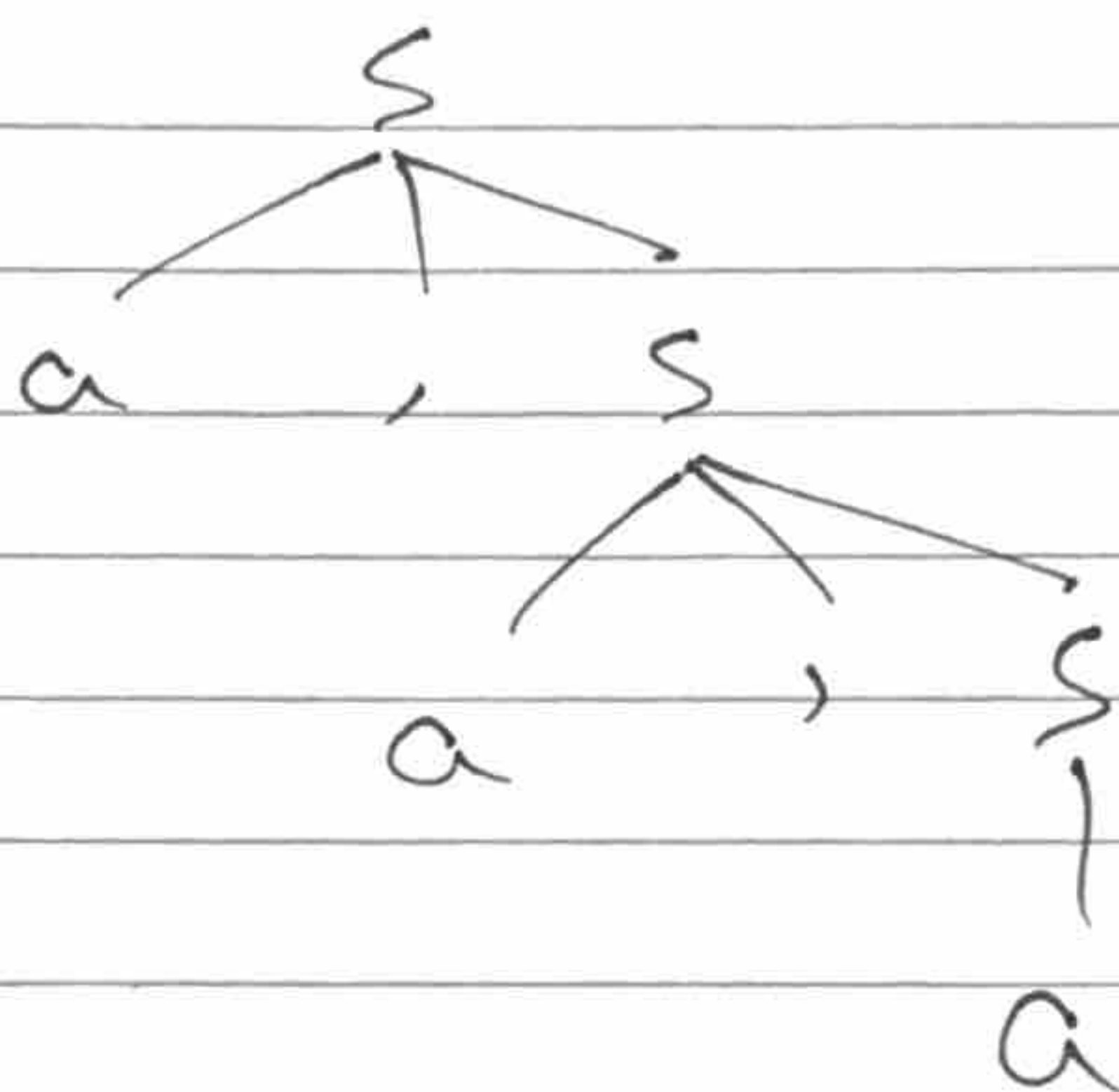
This shows the grammar is correct and unambiguous.



c] Right associative list of identifiers separated by commas.

$$S \rightarrow a, S \mid a$$

to generate  $a, a, a$  we have one parse tree



This shows the above grammar as correct and unambiguous.

d] Arithmetic expression of integers and identifiers with the four binary operators  $+$ ,  $-$ ,  $*$ ,  $/$ .

$$S \rightarrow S + N \mid S - N \mid N$$

$$N \rightarrow N * T \mid N / T \mid T$$

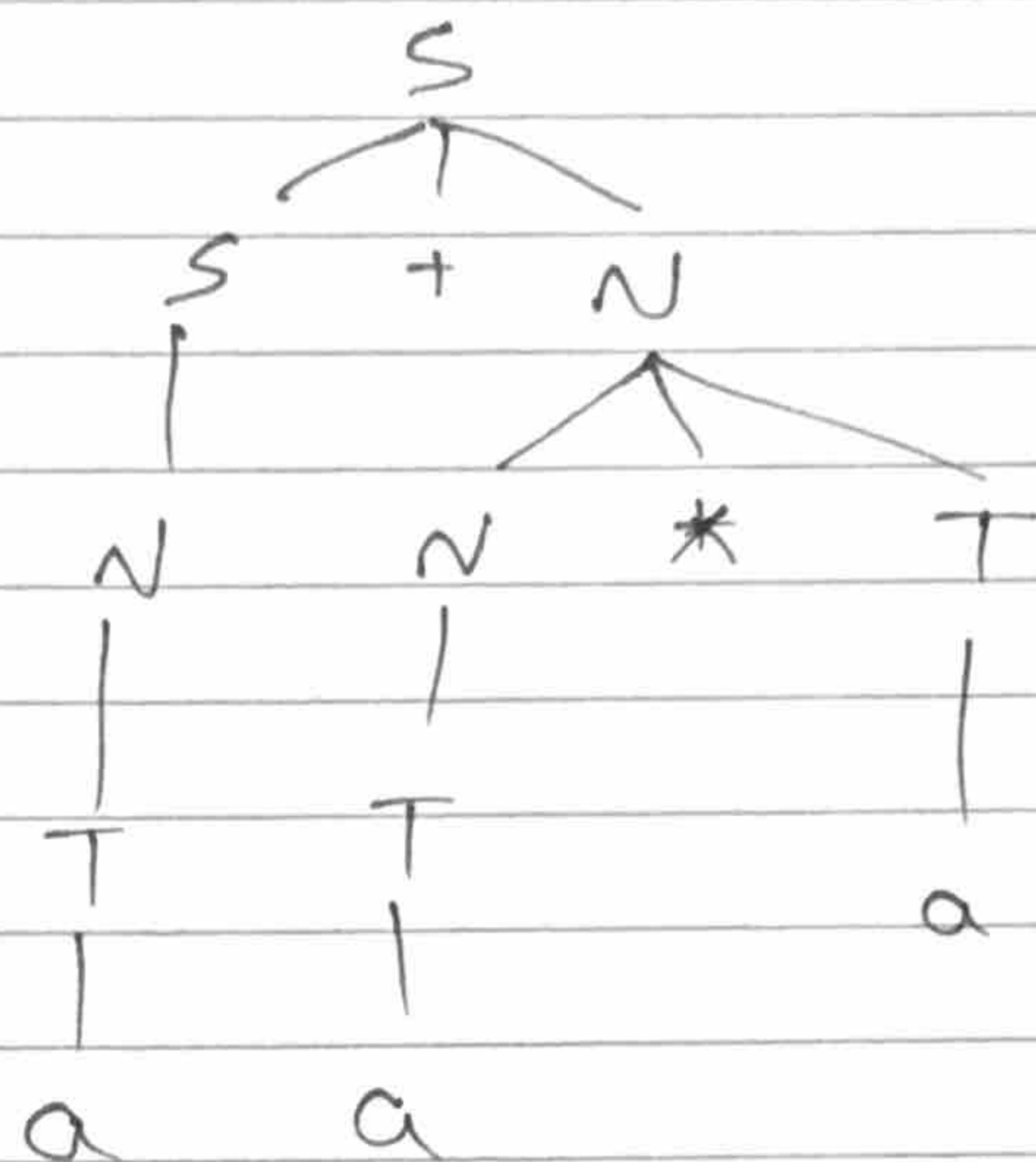
$$T \rightarrow \text{num} \mid a \mid (S)$$

To generate  $a + a * a$ , so the multiplication is to be performed first.

In parse tree if any operation ~~an~~ has higher precedence then it should be lower in the tree.



$a + a * a$



as  $*$  is lower in the parse tree it will be executed first.

Also there will be only one parse tree generated for above expression.

Q] Add unary plus and minus to the arithmetic operators of (d)

$$S \rightarrow S + N \mid S - N \mid N$$

$$N \rightarrow N * T \mid N / T \mid T$$

$$T \rightarrow +F \mid -F \mid F$$

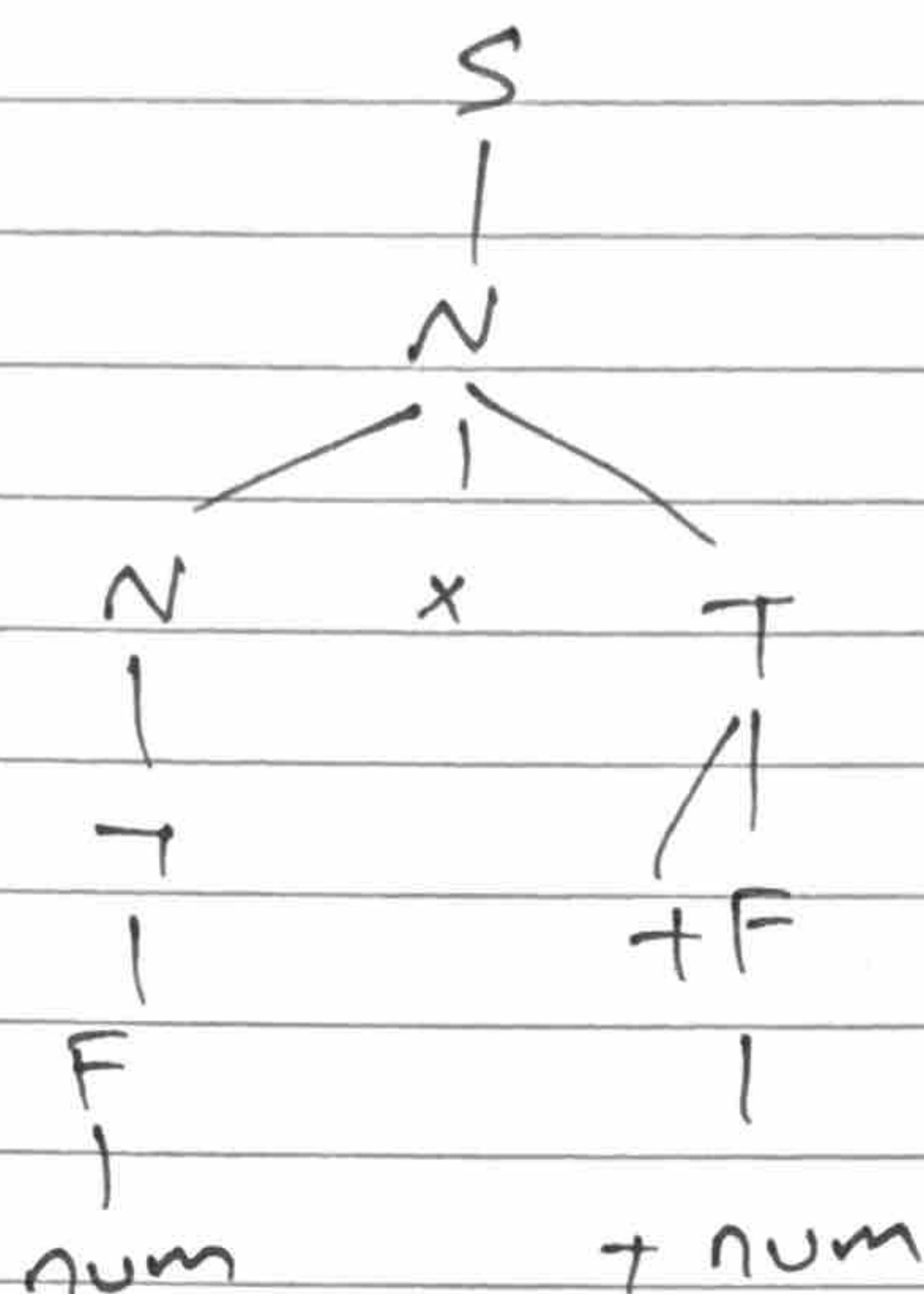
$$F \rightarrow id \mid num \mid (S)$$

Example:-  $+num \ num * \ num$

The unary operator has high precedence as compared to the multiplication. So, it will come at lower level in parse tree than multiplication.



+num num x num



### Exercise 2.2.6

Construct a context-free grammar for roman numerals.

⇒ In decimal the roman value is represented by the following table.

Roman Symbol	I	V	X	L	C	D	M
Decimal value	1	5	10	50	100	500	1000

A number can be generated with the combination of above roman numerals such as

I → 1  
 II → 2  
 III → 3  
 IV → 4  
 ⋮  
 VII → 7  
 IX → 9



RN  $\rightarrow$  ~~too~~ thousand hundred ten unit

thousand  $\rightarrow$  M | M | MMM |  $\epsilon$

hundred  $\rightarrow$  sh | CD | Dsh | CM

sh  $\rightarrow$  c | cc | ccc |  $\epsilon$

ten  $\rightarrow$  st | XL | Lst | XC

st  $\rightarrow$  x | xx | xxx |  $\epsilon$

unit  $\rightarrow$  su | IV | Vsu | IX

su  $\rightarrow$  I | II | III |  $\epsilon$

Example: 49

RN  $\rightarrow$  thousand hundred ten unit

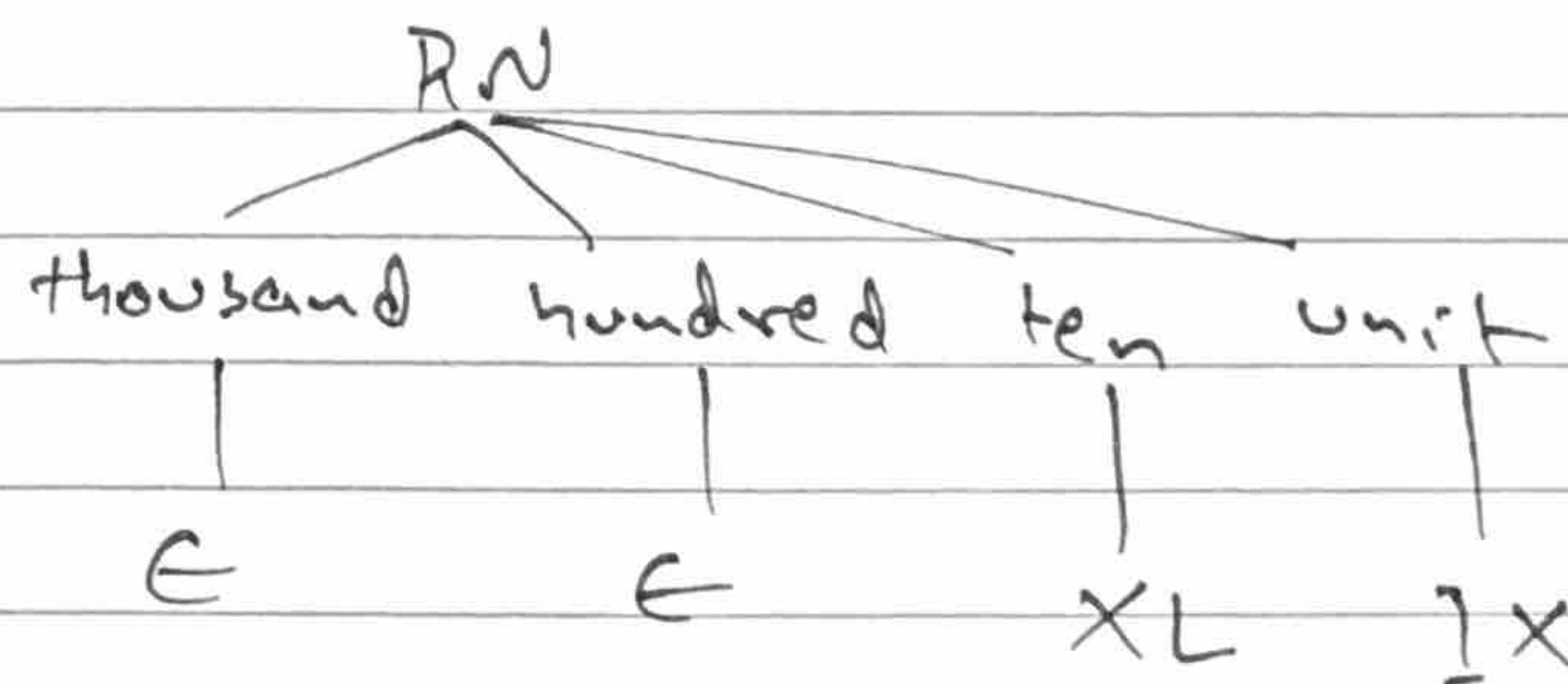
$\rightarrow$   $\epsilon$  hundred ten unit

$\rightarrow$   $\epsilon$  ten unit

$\rightarrow$  XL unit

$\rightarrow$  XLIX

Parse Tree





### Exercise 2.3.2

Construct a syntax-directed translation scheme that translates arithmetic expression from postfix notation into infix notation. Give annotated parse trees for inputs  $95-2*$  and  $952*-$ .

Following ~~notat~~ grammar is for postfix notation.

$$S \rightarrow SF+ | SF- | F$$

$$F \rightarrow FT* | FT/ | T$$

$$T \rightarrow \text{num} | (S)$$

$$\text{num} \rightarrow [0 \dots 9]$$

Annotated ~~parse tree~~ Translation schemes

$$S \rightarrow S \{ \text{print}(' + ') \} F + | S \{ \text{print}(' - ') \} F - | F$$

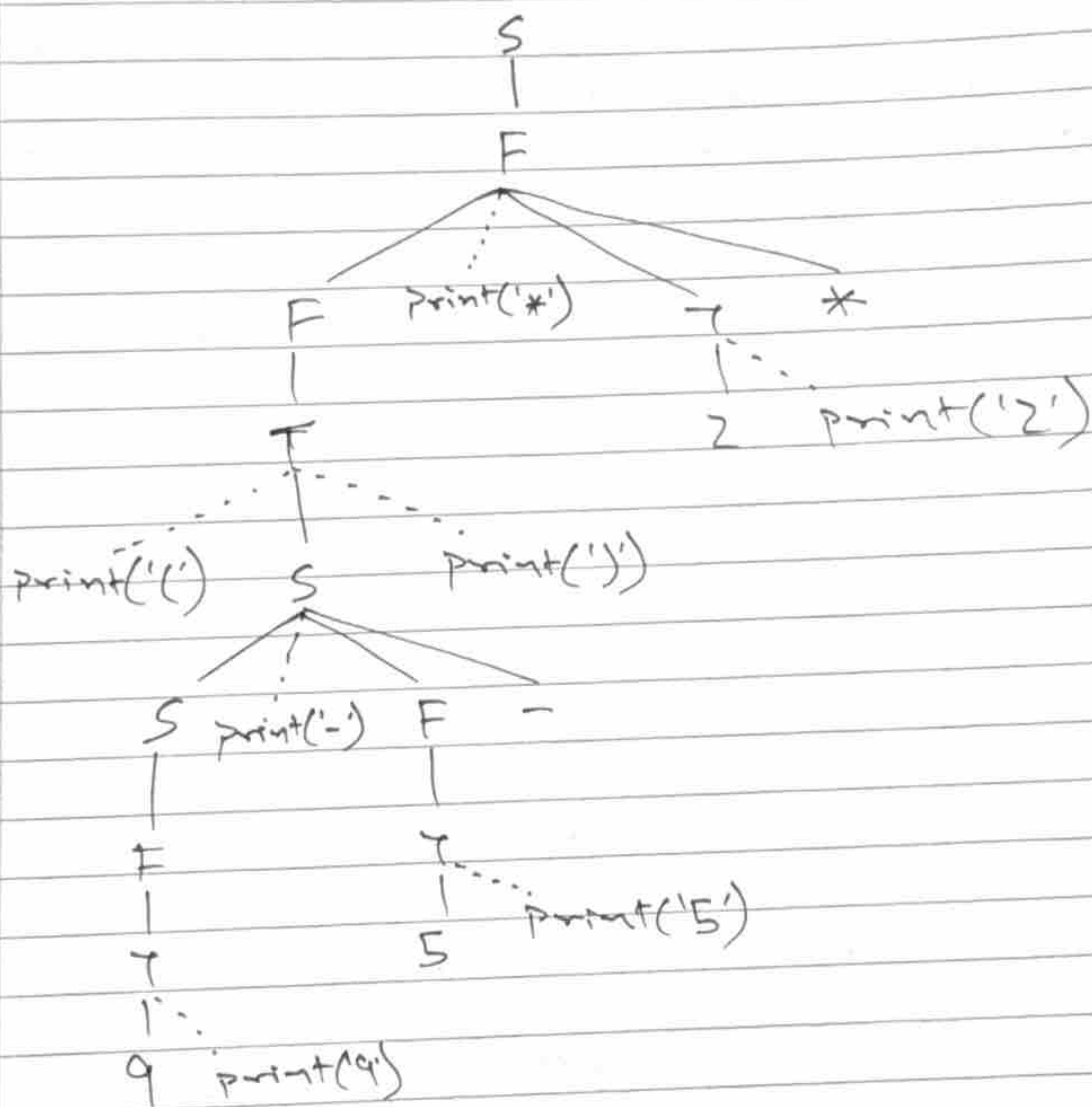
$$F \rightarrow F \{ \text{print}(' * ') \} T * | F \{ \text{print}(' / ') \} T / | T$$

$$T \rightarrow \text{num} \{ \text{print}(' \text{num.value} ') \}$$

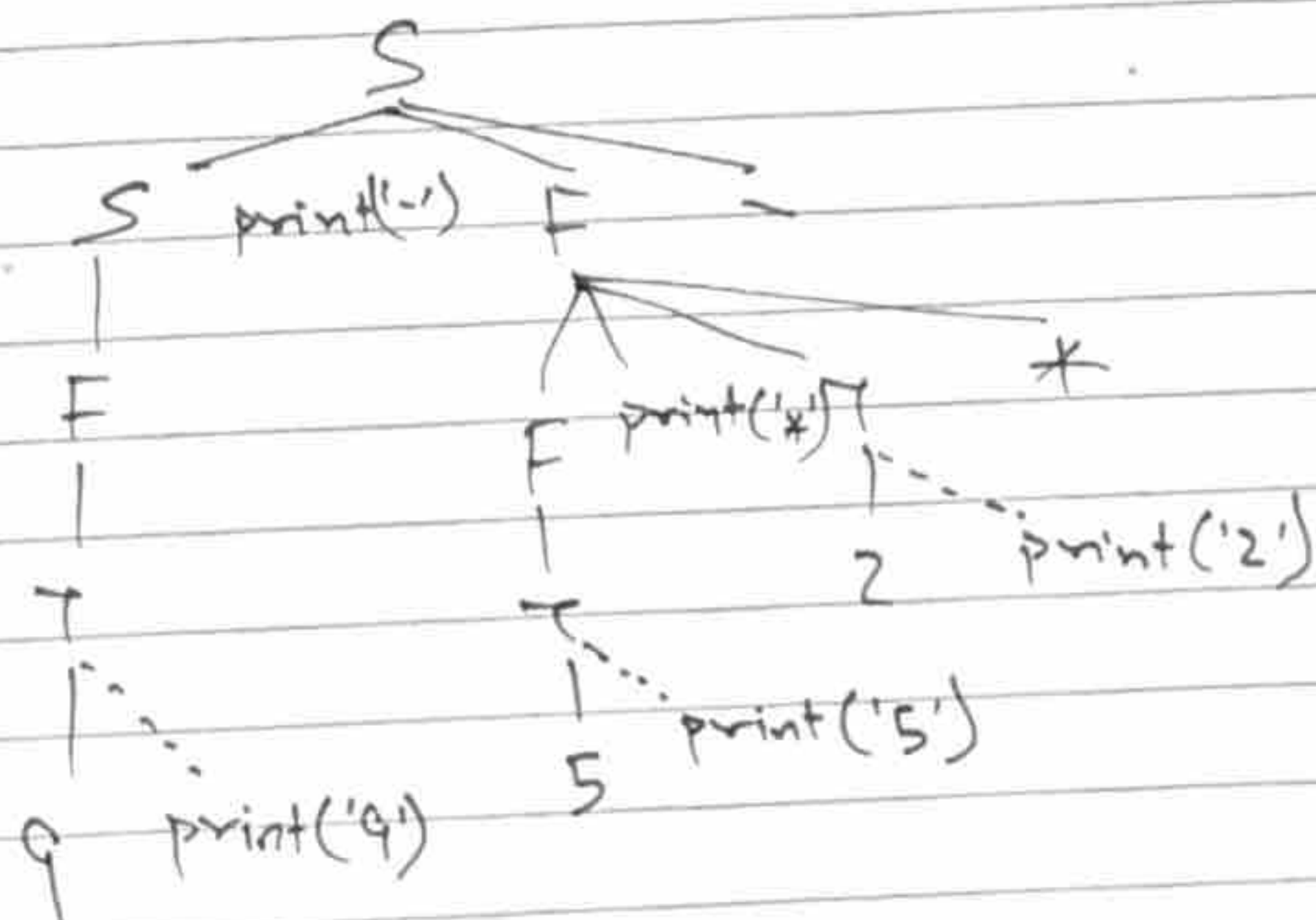
$$| \{ \text{print}(' ( ') \} S \{ \text{print}(' ) ') \}$$



⇒ Annotated parse tree for 95-2\*



⇒ Annotated parse tree for 952\*-





### Exercise 2.3.3

Construct a syntax directed translation scheme that translates integers into roman numerals.

num  $\rightarrow$  thousand hundred ten digit  
 $\{ \text{num.roman} = \text{thousand.roman} \parallel \text{hundred.roman}$   
 $\{ \parallel \text{ten.roman} \parallel \text{digit.roman};$   
 $\text{print(num.roman)} \}$

thousand  $\rightarrow$  low  $\{ \text{thousand.roman} = \text{repeat}('M', \text{low.v}) \}$

hundred  $\rightarrow$  low  $\{ \text{hundred.roman} = \text{repeat}('C', \text{low.v}) \}$

| 4  $\{ \text{hundred.roman} = 'CD' \}$   
| high  $\{ \text{hundred.roman} = 'D' \parallel \text{repeat}('X', \text{high.v} - 5) \}$   
| 9  $\{ \text{hundred.roman} = 'CM' \}$

ten  $\rightarrow$  low  $\{ \text{ten.roman} = \text{repeat}('X', \text{low.v}) \}$   
| 4  $\{ \text{ten.roman} = 'XL' \}$   
| high  $\{ \text{ten.roman} = 'L' \parallel \text{repeat}('X', \text{high.v} - 5) \}$   
| 9  $\{ \text{ten.roman} = 'XC' \}$

digit  $\rightarrow$  low  $\{ \text{digit.roman} = \text{repeat}('I', \text{low.v}) \}$

| 4  $\{ \text{digit.roman} = 'IV' \}$

| high  $\{ \text{digit.roman} = 'V' \parallel \text{repeat}('I', \text{high.v} - 5) \}$

| 9  $\{ \text{digit.roman} = 'IX' \}$



low  $\rightarrow 0 \{low.v = 0\}$

11  $\{low.v = 1\}$

12  $\{low.v = 2\}$

13  $\{low.v = 3\}$

high  $\rightarrow$  ~~4~~ 5  $\{high.v = 5\}$

16  $\{high.v = 6\}$

17  $\{high.v = 7\}$

18  $\{high.v = 8\}$

x ————— x ————— x ————— x —