

? operator:

Exp1 ? Exp2 : Exp3

↓  
check

(Exp1)  
true

↓  
Exp1

(Exp1)  
false

$y = x > 9 ? 100 : 200 ;$

Same as

If ( $x > 9$ )  $y = 100;$

else  $y = 200;$

## & and \* operators

& → memory address

pointer variable → specifically declared to hold a pointer to an object of specified type.

- Allow functions to modify contents
- support linked lists, binary trees
- dynamic data structures.

& → unary operations (returns memory of operand)

$m = \& Count;$

\* → Complement of &

\* → unary returns the value of the object located at the address. that follows it.

$a = *m;$  → meaning count!

&, \* → higher precedence than arithmetic operators

except unary operators (-) → shares equal precedence.

$\text{Pnt} * a;$  → denote that a will point to a address.  
(hold an address)

char \* ch; → ch is not a character but a pointer to a character.

base type      object of the base type.

mix pointer and non pointer declaration

$\text{int } a, *b, \text{Count};$

{ pointer to integer → [] }

## Sizeof → Compile time unary operator!

- \* gives same result as defined by type of variable & size\_t (loosely = unsigned int).
- \* returns type same as (the type of size\_t (loosely = unsigned int)).

### Comma operator

example \* here is

$x = (y=3, y=4);$

$x=4$

$y=4$

①  $x$  set 3

→ parentheses

②  $x$  set to 4! (final) Answer.

$x=1,2,3,4;$

$x=1 \rightarrow$  not 4

$x=(1,2,3,4);$

→ parentheses

Now  $x=4$

### Dot (.) , arrow (→) operator

- \* Access individual elements of structures & unions (Compound DS under single name).

\* Dot → used with structure / union

\* Arrow → used with pointers to a structure.

Struct employee {

char name [80];

int age;

float wage;

} emp;

→ tag!

emp.wage = 123.23;

P → wage = 123.23;

(meaning dereference P  
to get emp.wage)

Struct employee \*P = &emp;

(\*P).wage = 123.23

### [ ] , ( ) operators

( ) → increase precedence of operations inside it!

[ ] → array indexing!

Highest ( ) → .

! ~ ++ -- - (type) \* & size of

\* / % - & endl → executing & printing

+ -

<< >>

< <= > >=

= = != & ! = : check → executing & wide visibility \*

&

^

!

({} ) braces

&&

||

? :

= . += -= \*= /= etc..

! brace [

: (sc (5 bl) n p) bl = sc

: (sc (T bl)) n p \ bl = sc

wide visibility & scope env? & \*

brace

(brace, bl)

brace

(brace, bl)

$x = f_1(l) + f_2(r);$

can't tell which one

first

### Type conversion

long double, others → also converted to long double

float , " float

exception: unsigned int ← can't be represented by a long.

(Both are converted to unsigned long).

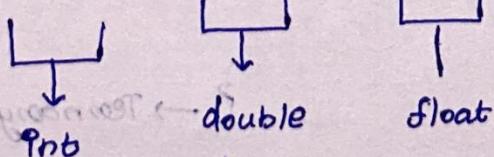
char ch;

int i;

float f;

double d;

result = (ch/i) + (f\*d) - (f+i); q : (callout) return ? sc < p  
"long type" - "suitable" if



PH : q

! suitable because of conversion to int

conversion of int to double

! suitable as int

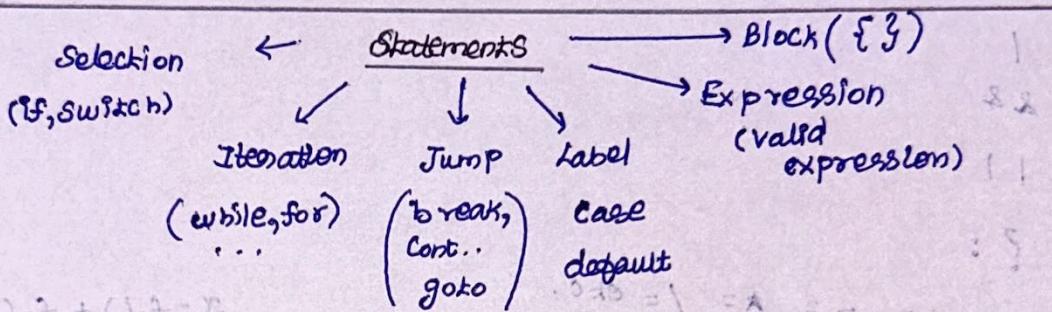
cast  
~~base~~  $x/2$ ;  $\rightarrow$  explicit

Spacing & parentheses - "Easy to read"

$x = 10/y \sim (127/x);$  ] same!  
 $x = 10 / y \sim ((127)/x);$

\* Additional space / parentheses  $\rightarrow$  don't slow down.

\* Gives more classification.



see slides No 1 true/false  $\rightarrow$  Boolean expressions

if (scalar - int / char / pointer / float / bool)

rand()  $\rightarrow$  <standard>

if (?)  $\rightarrow$  good if  
 if (?)  $\rightarrow$  good because it is followed by else  
 if (?)  $\rightarrow$  These two are associated!

?  $\rightarrow$  Alternative

1 > 2 ? printf("Hello") : printf("Hi");

O/P : Hi

?  $\rightarrow$  Ternary operator!

\* It is not restricted to variable alone/assignment

we can do anything!

## Conditional expression!

\* Any valid expression can control if (or) ? operator! (not just restricted to relational & logical operators as in PASCAL/BASIC)

True → non zero
False → zero

## Division:

```
switch (expression) {
```

```
    Case constant 1 :
```

```
        :
        break;
```

```
    Case constant 2 :
```

```
        :
        break;
```

```
    default
```

```
        Statement
```

Int  
Char  
Constants

## Switch

\* float not allowed!

\* case % executed (matched) → until break

\* No match → default (optional)

c99: At least 1 or 3 statements (case)

'low : high efficiency' (no. of cases)

\* switch → only test for equality

\* No two case - identical expression constant!

\* char - Automatically converted to int.

```
switch (n) {
```

```
    case 1
```

```
        printf("E");
```

```
    case 2
```

```
        printf("A");
```

(only true:

execute  
until break

or  
end of  
switch)

2

ARIPRAKASH

1

HARIPIRAKASH

3

IPIRAKASH

→ No break!

1

HARIPIRAKASH Hello

```
    case 5
```

```
        printf("PRAKASH");
```

```
        break;
```

```
    default
```

```
        printf("Hello");
```

## Nested Switch

```
switch (n) {
```

```
    case 1
```

```
        switch (n)
```

```
        {
```

```
        }
```

outer correct doesn't mean inner also correct.

Inner - evaluate & follow the  
same!

## Iteration

`for (initialize; check; increment/decrement)`

**Check: Condition**

Variants:

`for (x=0, y=0; x+y < 10; ++x)`

break cont. → break

break → break

\*\*\*\*\*

(79) Page 37 \*

} (multiple) values

29

14 - 14

`for ( : x < 10; ) { }`  
↓  
(no init, no incr/decr)

`for ( ; ; )`

`for ( ; ; ) → oo loop;`

\* A statement may be empty!

\* Body of for loop - also may be empty!

"Show" ≠ "Show"

↳ Remove space?

`for ( ; *str == ' ' ; str++);`

Advance leading spaces!

Note: arr[20] = "Hello";

\* we get ordinary pointer! that points somewhere

\* Incrementing that pointer leads to invalid memory

∴ It is a pointer to an array!

array: non modifiable left values (values) → No operators modify the array it self.

problem: Create explicit pointers

eg: char \*pt = arr;

`for ( ; *start == ' ' ; start++) ;` do copy, because  $start[0] = "Hello"$ ;  
↳ value error!

`char arr[15] = "Hello";`

`char *a = arr;`

~~return to the original state~~ `for ( ; *a == ' ' ; a++) ;`

`printf("%s", a);`

Ans: Hello

'No whitespace-leading'

`char *arr = " -- Hello";`

`for ( ; *arr == ' ' ; arr++) ;`

`printf("%s", arr);`

→ equal

! whitespace will print this

! whitespace ← local

Time delay

\* `for (t=0; t < Some_Value; t++)` → setup time

embedded programming!

declare within for loop

`for (int i=0; i<10; i++)`

! important distinction of code is declaration

! (if you want memory leak, if you're not careful)

! goal is straightforward a  
while ( $i < \text{lenString}$ )

! keeping do Padding → fill with spaces, finally \0

. always present after do

do while

space

\0

do {

:

do at least once, then check!

? while(. . .) doesn't work like do while

Continue

Jump statements

→ Return

→ goto

↓ → break

↓ → exit()

Return:

\* Return from a function (Jump back to where took egg!)

\* Return - may or may not have associated value

\* Return - only return non void function!

\* In void → Preturn without value is used.

Return function → must return a value!

\* Function stops at 1st return!

\* void function → No return value.

return expression;

goto

\* occasionally has its uses - narrow situations - jumping out of set of nested loops.

\* Can't jump b/w functions!

\* label → identifier!

goto label;

switch control ← (++) Jumps back to label

label;

switch control

Break

\* Terminate a case in switch statement

\* Terminate a loop. (only inside loop - where break is)!

break;      Kbdhit() → returns 0 → if key is not pressed!

else non zero value.

'control'

(Break causes exit from innerloop alone?)

exit() → not a jump statement

\* Function in std library

\* Termination of the entire program section

to operating system.

\* general form:

void exit (int return\_code);

return code → exit status

\* Usually 0000 is used as return code!

\* Also: common: EXIT\_SUCCESS, EXIT\_FAILURE

Macros for return code.

exit code 0

<stdlib.h> → exit()

(use for debugging!)

exit() → quit the program!

Continue!

- \* Doesn't terminate → Just skips to next iteration.
- \* cause increments; But after Continue statement won't be executed in that particular iteration alone!

eg: if space: don't point - Just skip!

Expression statements: valid expression followed by semicolon

eg: func(); // function call

statement → iteration

a = a + b; // assignments

b + f(); // valid - no use (as never stored/used)

{} ; // empty statement!

(++int i=5; i<8)

Block Statement

{

i = 120;

→ Free Standing

printf("%d", i); → 120

(i = 120) block statement!

g

printf("%d", g); → 120

not like func - creating local values?

\*\*\*\*\* i=0  
\*  
\* j=end.  
\*  
\*  
\*\*\*\*\* i=end

'Patterns'

hollow square with diagonal

i == j

→ left diagonal

$$\begin{aligned}(0, 4) &= 5 \\ (2, 3) &= 5 \\ (3, 2) &= 5 \\ (4, 1) &= 5.\end{aligned}$$

*				*
	*			
		*		
			*	

Right  
diagonal  
 $i+j=n$

! check condition so know if next iteration

INITIAL STATE

*	*	*	*	*	*	*	*	*	*
*	*	*	*	*	*	*	*	*	*
*	*	*	*	*	*	*	*	*	*
*	*	*	*	*	*	*	*	*	*
*	*	*	*	*	*	*	*	*	*

From the loop it's

obvious!

check condition

! corresponding flow ← (for)

for ( $i=0$ ;  $i < n$ ;  $i++$ )

{ for ( $k=0$ ;  $k < n-1-i$ ;  $k++$ )

    printf ("\* ");

for ( $j=0$ ;  $j < n$ ;  $j++$ )

    printf ("\* ");

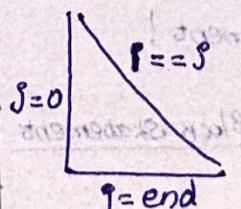
    printf ("\n");

}

relationship of iterations following below

### mirrored - Rhombus -

0	1	2	3	4
*				
*	*			
*		*		
*			*	
*	*	*	*	*



for ( $i=0$ ;  $i < n$ ;  $i++$ )

    for ( $j=0$ ;  $j < i$ ;  $j++$ )

        printf ("\* ");

    printf ("\n");

for ( $i=n-1$ ;  $i >= 0$ ;  $i--$ )

    for ( $j=0$ ;  $j < i$ ;  $j++$ )

        printf ("\* ");

    printf ("\n");

( $i=0$ ;  $i < n$ ;  $i++$ )

( $i=n-1$ ) || ( $j=0$ ) || ( $i=j$ )

else ← ; ( $i=0$ ) ?

				*
			*	*
		*	*	
	*	*	*	
*	*	*	*	*

$j=0$

$i=0$  \* \* \* \*

\*

$i=1$  \*

\*

$i=2$

*				
	*			
		*		
			*	

$i=(0, 1)$

$i=(1, 2)$

$i=(2, 3)$

$i=(3, 4)$

★	★	★	★	★
★	★	★	★	
★	★	★	★	
★	★			
★				

	0	1	2	3	4
0	*	*	*	*	*
1	*			*	
2	*		*		
3	*	*			
4	*				

$$\text{diagonal: } 9+3 = (n-1)$$

1 = 0

$$J = 0$$

J

$\left( \begin{array}{l} \text{for } i = 0; i < n; i++ \\ \quad \text{for } j = 0; j < n - i; j++ \end{array} \right)$

★	★	★	★	★
	★			★
		★		★
			★	★
	9	A	dictor	55235
		al	A	dictor
				9
				9=0

A grid pattern consisting of stars and asterisks arranged in a grid. The grid has 10 columns and 10 rows. The first column contains 10 stars. The second column contains 9 stars and 1 asterisk. The third column contains 8 stars and 2 asterisks. This pattern continues, with each subsequent column containing one less star and one more asterisk than the previous one. The last column contains 1 asterisk.

( $\theta = 0$ ;  $\theta < \omega$ ;  $\theta + +$ )

$K=0; K < n-1; K++$

## Space

$j=0; j < (2 * 9) + 1; j++$

☆

$$\text{hollow: } g=0, \theta=n-1, \quad S=(2\pi)^n$$

invested!

hollow

9 = 0



Number pattern 14

5	5	5	5	5
5	4	4	4	4
5	4	3	3	3
5	4	3	2	2
5	4	3	2	1

Numbers pattern 15

1	2	3	4	5
2	3	4	5	5
3	4	5	5	5
4	5	5	5	5
5	5	5	5	5

Pattern 16

1	2	3	4	5
2	3	4	5	1
3	4	5	2	1
4	5	3	2	1
5	4	3	2	1

Pattern 17

1	2	3	4	5
2	1	2	3	4
3	2	1	2	3
4	3	2	1	2
5	4	3	2	1

left  
right

5	5	5	5	5	5
4	4	4	4	4	4
3	3	3	3	3	4
3	2	2	2	3	4
3	2	1	2	3	4
3	2	1	2	3	4
5	5	5	5	5	5

5	5	5	5	5	5
4	4	4	4	4	4
3	3	3	3	3	3
2	2	2	2	2	1
1					

Pattern 18

5	5	5	5	5	5	5	5
4	4	4	4	4	4	4	5
3	3	3	3	3	3	4	5
2	2	2	2	2	3	4	5
1	2	3	4	5	6	7	8
5	4	3	2	1	2	3	4
5	4	3	3	3	3	3	4
5	4	4	4	4	4	4	5
5	5	5	5	5	5	5	5

3	4	5
4	5	
.	5	

match again

details (A)

miss (B)

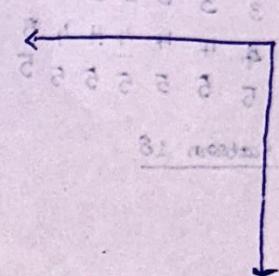
goal (C)

out (D)

9	8	7	6	5	4	3	2
9	8	7	6	5	4		
9	8	7	6	5			
9	8	7	6				
9	8	7					
9	8						
9							

1	2	3	4	5	6	7	8	9	10
36	37	38	39	40	41	42	43	44	11
35	64	65	66	67	68	69	70	45	12
34	63	84	85	86	87	88	71	46	13
33	62	83	96	97	98	89	72	47	14
32	61	82	95	100	99	90	73	48	15
31	60	81	94	93	92	91	74	49	16
30	59	80	79	78	77	76	75	50	17
29	58	57	56	55	54	53	52	51	18
28	27	26	25	24	23	22	21	20	19

0	1	2	3	4	5	6	7	8	9
0	1	2	3	4	5	6	7	8	9
1	36	37	38	39	40	41	42	43	44
2	35	64	65	66	67	68	69	70	45
3	34	63	84	85	86	87	88	71	46
4	33	62	83	96	97	98	89	72	47
5	32	61	82	95	100	99	90	73	48
6	31	60	81	94	93	92	91	74	49
7	30	59	80	79	78	77	76	75	50
8	29	58	57	56	55	54	53	52	51
9	28	27	26	25	24	23	22	21	20



①  $[0][0]$  to  $[0][9]$       |     $[1][9]$  to  $[9,9]$       |     $[9,8]$  to  $[9,0]$       |     $[9,0],[0,0]$   
     "left to right"      |    "Right to left"      |

3	8	4	3	0	5	8	K
6	4	3	8	P			
H	3	0	T	3	P		
3	0	T	8	P			
3	0	T	8	P			

more detail

1) Go right

2) Down

3) Go left

4) Go up!

$L \rightarrow R$	$down$	$right \rightarrow left$	$up$
0,0 to 0,9	1,9 to 9,9	9,9 to 0,0	8,0 to 0,0
1,1 to 1,8	2,8 to 8,8	8,7 to 8,1	7,1 to 2,1
2,2 to 2,7	3,7 to 7,7	7,6 to 7,2	6,2 to 3,2
3,3 to 3,6	4,6 to 6,6	6,5 to 6,3	5,3 to 4,3
4,4 to 4,5	5,5	5,4	
9,9 to 9=0st map		left	
9,9 to 9 goes to 9	area [9] [9]	, 9++	
9 goes to 9	area [9]	1 to 9	left = 0 right =

left right  
from 0 to 9

① area [left] [9]  
area [9] [down]

1	2	3	4	5
16	17	18	19	6
15	24	25. 20	7	
14	23	22	21	8
13	12	11	10	9

left = 0  
right = 4  
top = 0  
bottom = 4

\* left to right

① area [left] [9]

$[9 = left \text{ to } right]$   
 $(0) \text{ to } (4)$

\* Top to down

area [9] [right]

$[9 \text{ from top} + 1 \text{ to bottom}]$

\* Right to left.

area [bottom] [9]

$[9 \text{ from right} + 1 \text{ to left}]$

\* down to top

area [9] [left] .  $[9 \text{ from bottom-1 to top+1}]$

Now

left = 1, right = 3

top = 1, bottom = 3.

0,0 at 0,8 It seems like we have enough 0,1

1,0 at 1,7

2,0 at 2,6

3,0 at 3,5

'AS Squeeze' - Top = left

Right = bottom.

1	2	3	4	5
16	17	18	19	6
15	24	25	20	7
14	23	22	21	8
19	12	11	10.	9

①  $\text{left} = 0, \text{right} = 4$

$\text{area}[\text{left}][\text{right}] \rightarrow \text{from left to right}$

$\text{area}[\text{left}][\text{right}] \rightarrow \text{from left+1 to right}$

$\text{area}[\text{right}][\text{right}] \rightarrow \text{from right-1 to left}$

$\text{area}[\text{left}][\text{right}] \rightarrow \text{from right-1 to left+1}$

②  $\text{left} = 1, \text{right} = 3$

$\text{area}[\text{left}][\text{right}] \rightarrow \text{from left to right}$

$\text{area}[\text{left}][\text{right}] \rightarrow \text{from left+1 to right}$

$\text{area}[\text{right}][\text{right}] \rightarrow \text{from right-1 to left}$

$\text{area}[\text{left}][\text{right}] \rightarrow \text{from right-1 to left+1}$

$\text{left} = 2, \text{right} = 2.$

\*  $\text{area}[\text{left}][\text{left}] = \text{left}$

[ $\text{left}$  to  $\text{left}$ ]

\*  $\text{area}[\cdot \cdot] = \text{no!}$

[ $\text{left}$  to  $\text{left}$ ]

\*  $\text{area}[\cdot]$

[ $\text{left}$  to  $\text{left}$ ] no! decrements!

\*  $\text{area}[\cdot \cdot] = \text{no!}$

[ $\text{left}$  to  $\text{left}$ ]

$\frac{N}{2}$  times

Facts:  $10/2 = 5$  times. 0 to 4 (5 times)  $\Rightarrow \frac{N}{2}$

$$\frac{5}{2} = 2$$

① 0,0 to 0,4  
1,4 to 4,4  
4,3 to 4,0  
3,0 to 0,0

much at right \*

got at right \*

② 1,1 to 1,3  
2,3 to 2,3  
3,2 to 3,1  
2,1 to 2,2

No. of steps:  $\frac{N}{2}$ .

3,2 to 3,1 = right \* = total move

2,1 to 2,2 = moved 1 unit



end  $\alpha * i - i$

$$\alpha - 1 = 1$$

$$4 - \alpha = 2$$

$$6 - 3 = 3$$

$$8 - 4 = 4$$

	0	1	2	3	4
0	1				
1	2	6			
2	3	7	10		
3	4	8	11	13	
4	5	9	12	14	15

1	2	3	4	5
---	---	---	---	---

6	7	8	9	10	11	12	13	14	15
---	---	---	---	----	----	----	----	----	----

for ( $i=0; i < n; i++$ )

for ( $j=0; j \leq i; j++$ )

array[i][j]

$E = 1 - \frac{1}{4} = 1 - (0 - 1) * 5 = 6$

$E = 1 - \frac{1}{4} = 1 - (0 - 2) * 5 = 7$

① array[0][0] → 1

② array[0][1] → 2

[0][1] → 6

array[0][2] → 3

[1][2] → 7

[2][2] → 10

Note: need to leave - empty spaces (in arrays)

### Efficient method - Inplace execution

1	0			
2	6			
3	7	10		
4	8	11	13	
5	9	12	14	15

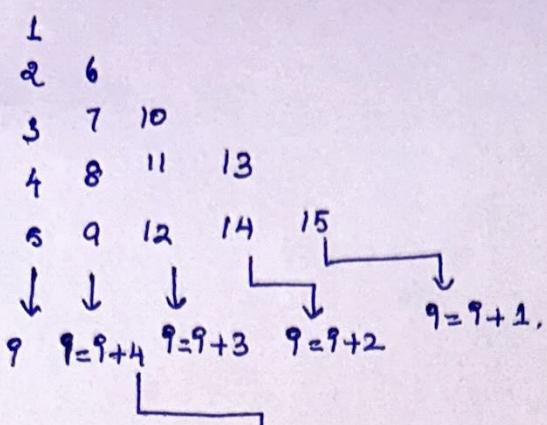
$\downarrow$      $\downarrow$      $\downarrow$      $\downarrow$      $\downarrow$   
 $i+4$      $i+3$      $i+2$      $i+1$

1				
2	7			
3	8	12		
4	9	13	16	
5	10	14	17	19

$\downarrow$      $\downarrow$      $\downarrow$      $\downarrow$      $\downarrow$   
 $i+5$      $i+4$      $i+3$      $i+2$      $i+1$

$(i-1) + 5 = i$

$(i-1) * 5 + 1 = i$



printf("%d", r)

for (j=0, k=n-1, l=j+k; j <= r; j++, k--, l=k+l)

printf("%d", l);

\* Initialize incrementing value by  $n-1$

\* Then each time Value = Previous value + incrementing value  
↳ decreased by 1 each time.

1	
2	4
7	11 16
22	29 37 46
56	67 79 92 106

1+1 = 2	16+6 = 22
2+2 = 4	22+7 = 29
4+3 = 7	37 = 29+8
7+4 = 11	87+9 = 46
11+5 = 16	

1	→ 1,1
3 2	→ 2,2
4 5 6	→ 3,3
10 9 8 7	→ 4,4
11 12 13 14 15	→ 5,5.

∴  $\text{fmax} = \cancel{9}$  → Index Atg

eg:

$q=2 \rightarrow 3 \quad 2$

$q=4 \rightarrow 10 \quad 9 \quad 8 \quad 7$

skip:  $q+1$

Not working

$K=2+9-1$   
 $3 \text{ to } 2$   
 $K=7+9-1 = 10$

10 to 7

(around  $A \leftarrow A + 1$ )  $B=N$  part 20  
if ( $q \geq 2$  ! = 0)  $\rightarrow$  Not working

$q+j \rightarrow$  points  $01 \rightarrow$  See code!

else {

`for (g=1; g<=n; g++)`

{  
    {  $i+j=9$ ,  $i-g=H$ ,  $(g=2)$  not  
 $g_f = (g \% d \neq 0)$   
 $i+j+g = 9 \geq 6$

(else)  
`for (g=1; g<=9; g++, k++)`

: (2 to 6) 2nd row  
points `("e%d", k);`

g  
g

else temp=k;

{  
    `for (g=k+1; g>temp; g++, k++)`

1st & 2nd row values  
{  
    points

initials `printf ("e% d", g)`

3

$$S_0 = i + j$$

$$P_S = i + g_0$$

$$i + P_S = r_0$$

$$i + H = P + r_0$$

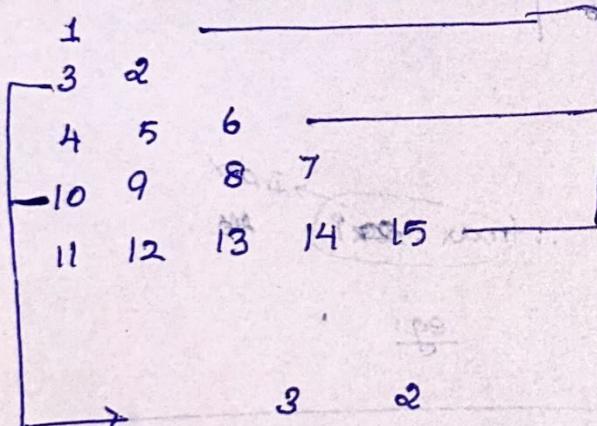
$$S_0 = i + j$$

$$H = S + S_0$$

$$r_0 = S + A$$

$$H = P + r_0$$

Logic



$$i = g + 1$$

$$K=1$$

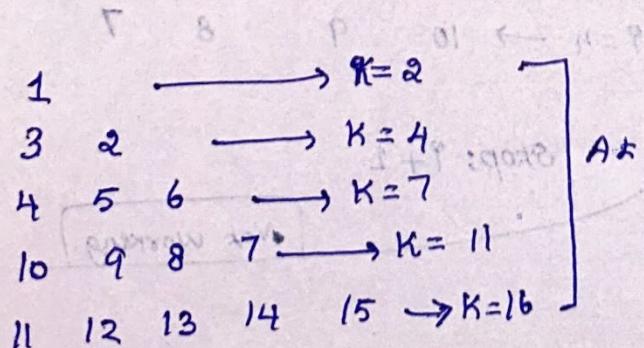
$$K=2$$

$$K=3$$

$$\vdots$$

$$S = S + A$$

) continuation of previous row values.  
(let take K)



At the end.

using K=2 (Row 1 to 2 ele, 4 to 4 elements)

$$g = K + g - 1 \rightarrow g = K$$

$$3 \rightarrow 2$$

$$10 \rightarrow 7$$

See pattern 51.