> BIG DATA ECOSYSTEM

Refers to the collection of frameworks, tools, and technologies used to store, process, analyze and visualize massive volumes of data.

Enables organizations to derive insights from structured, semi-structured b unsmichired data:

Key components:

Data Sources - Database, DW; XML & ISON files; Social media, Emails, Video data 2) Data Ingestion - consume the data from various sources (Katka, flume, Sqoop,

3) Data Storage - Distributed file System (HDFS, GFS, Amazon S3)

4) Data Processing - clean data (Apache spark, A stoum, Hadoop Map Reduce)

5) Data Analysis - Take insights 4 analyze the data (Apache Hive, Apache Pig)

Machine learning & Advanced Analytics - (Muib, Apache Mahout)

2) Processed Data Storage - (cassandra, Mongobs, 48ase) storing Nosqu database

(8) Visualization - Tableau, Pawer BI, Quik View

9) Cluster Management - Allocates resources, scheduling Jobs, etc. (YARN)

Processing Visualization Processed netherage basign totale Starage Data | Ingestion > Manager Machine Data Distributed Analytics

> INTRODUCTION TO GOOGLE FILE SYSTEM (GFS).

- GFS is scalable distributed file system for large distributed data-intensive app

Authentication / Authorization

- GFS is designed to handle batch workloads with lots of data.

The system is distributed: multiple machine store copies of every file, and multiple machines try to read/write the same file.

- GFS was originally designed for Google's use case of searching & indexing the web. It was developed around 2003, new version released in 2010 (GF\$2)

characteristics :-

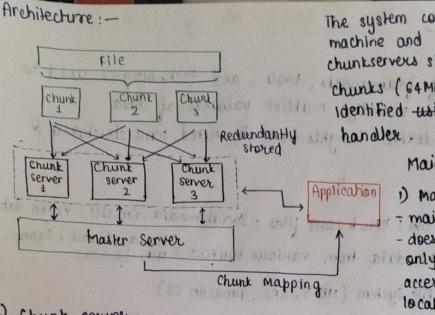
1) fault tolerent - even it disk is corrupted, data in them can be used & restored.

2) Designed for large files - Big data size.

3) Master-Slave Architecture - Uses single master to manage metadala & chunk servers to storce data. Each file is divided into fined size chunks (64 MB)

a) Write Once read many workloads

5) Internal usage - GFS only used within Google Intrastructure; It is not open-source or available for public use.



- Stores actual file data as fixed chunks (64 MB by default).

across different chunkrerver for fault tolerance.

- Serve ruad & write req. from client directly.

- Send & updates status to the master server.

The system consist of a sigle master machine and many chunk servers. These chunk servers store parts of files called chunks (64 MB data blocks), which are identified using by a 64 bit chunk

Main companents :-

maintains all file systems metadata - does not handle actual data;

only metadata (file namespace, access control, file mapping, chunk location) and coordination.

d) client/Application -

- Appl<sup>n</sup> or users that access GFS.

Interacts with the master server to obtain metadata a chunk location.

Communicates directly with chunk-servers for actual data transfer, minimizing the master's load.

working :-

1) Read operation

- Client req. file metadoba from master

- Master returns chunk handle & list of chunkserver.

- Client caches this into 4 reads data directly from chunkserver.

2) Write 1 Append Operation

- client asks master for current chunk b its replicas.

b provides replica locations.

- Client sends data to all replicas; the primary coordinates the write & maintains consistency

write is successful it majority of replicase

\* Heartbeat Mechanism - Mouter regularly checks chunkserver health.

Advantages:-

1) scalable 2) fault Tolerance 3) Cost Effective 4) Efficient for large files

5) Distributed Read/write

Disadvantages: (Masker) 2) Foremal fragmentation

1) single Point failure (master) 2) Internal fragmentation

> HADOOP ARCITECTURE

What is Hadoop? Apache Hadoop is an open source software library that is used to manage data processing & storage in big data applications. Hadoop facilitates analyzing large amount of data parallely & more quickly. Apache Hadop was introduced in 2012, by ASF 10 Apache Software foundation.

(are components: - 1) HDFS 2) YARN 3) MapReduce 4) Hadoop common YARN ( Yet Another Resource Negotiator) - responsible for Resource mangement and jab scheduling in HDFS. Hadoop common - Java libraries & utilities used across modules. > HADOOP STORAGE HDFS (Hadoop Distributed file System) distributed file system for storing a retriving large files HDFS is a with streaming data in record file. Architecture :-- It is a core component of Metadata ops Hadoop ecosystem. Name Node HDFs has Master-Slave architecture. Black aps 1) Name Node (Master) Client - central server that manages the file system namespace V Data Nodes DataNodes Replications and metadata. חם 맘미믐 keeps metadata in memory for fast access. Rack 2 Rack 1 Write Tchent Handles dient request for file operations (open, close - Monitors Data Node via heartbeats & block reports. Ensures data block replication & manages recovery in case of DataNode failure. 3) Data Node (slave) - stores actual data blocks on local disks.

- serve read and write req. from client as directed by the NameNade.

Report back to Name Node with block information & health status.

files and Blocks - the file is the doubt which we want to store, when we store a file into HDFS it's broken to blocks, the default size of each

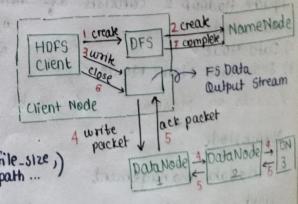
128/256 MB or 64 MB. Each Block is replicated for providing fault tolerance & availability, the

default number is 3.

ANATOMY OF FILE WRITE AND READ

1) HDFs client calls create() fon DFS (Distributed file system) to create a file.

2) DFS interacts with Namenade through RPC (Remote Procedure Call) to create a new file in filesystem namespace (with file-size,) associated with it.



3) The Namenode makes sure that the file doesn't already exist & client has the night permissions to creak file. If all these checks pass, the Namenode makes a record of the new file. The DFS returns an FSData output Stream 4) for client to start writing data to Datanode.

As client starts writing data, the DFS Octipul Stream (wrapped in FSDOS)

splits the client's data into packets & writes it to an internal queue called data queue. This dataqueue is used by DataStreamer, which tells

Name Node to allocate new blocks. The Name Node returns a list of suitable

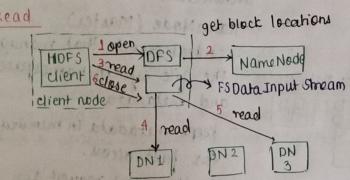
Data Moder for replication.

5) The DFS Output Stream also maintains another queue of packets, called ack queue, which is waiting for the acknowledgement from Data Nodes.

6) The HDFs client calls the close () method on the stream when it functions

writing data.

7) The FS Data Output Stream then sends an acknowledgement to Name Node.



1) HDFS dient calls open () on file system object.

2) DFS then calls NN using RPC to get the location of the blocks of a file. For each block NN returns the address of the close DN from the client that countain 0

2

N

a copy of that block. Then DFS returns an FSDOS to the client from where

the client can read the data. 3) Then the dient calls the read () method on FSData Input stream abject.

4) The DFSInpul stream (wrapped in FSDIS) which contains the addresses of the blocks of the file, connects to the closest DN to read the first block 5) Upon reaching the end of the file, DFSIS closses the connection with that

DataNade and finds the best suited DN for the next black.

() when the client has finished reading the data, it calls close () on the

FSDIS.

NameNode, Lecondary NameNode, and DataNode.

Apache Hadoop 2 consists of foll Daemons Daemons = process

4) Resource Manager / 1) Name Node

5) Node Manager V 2) Data Node V

3) Secondary Name Node V I works on Master I works on blave machine

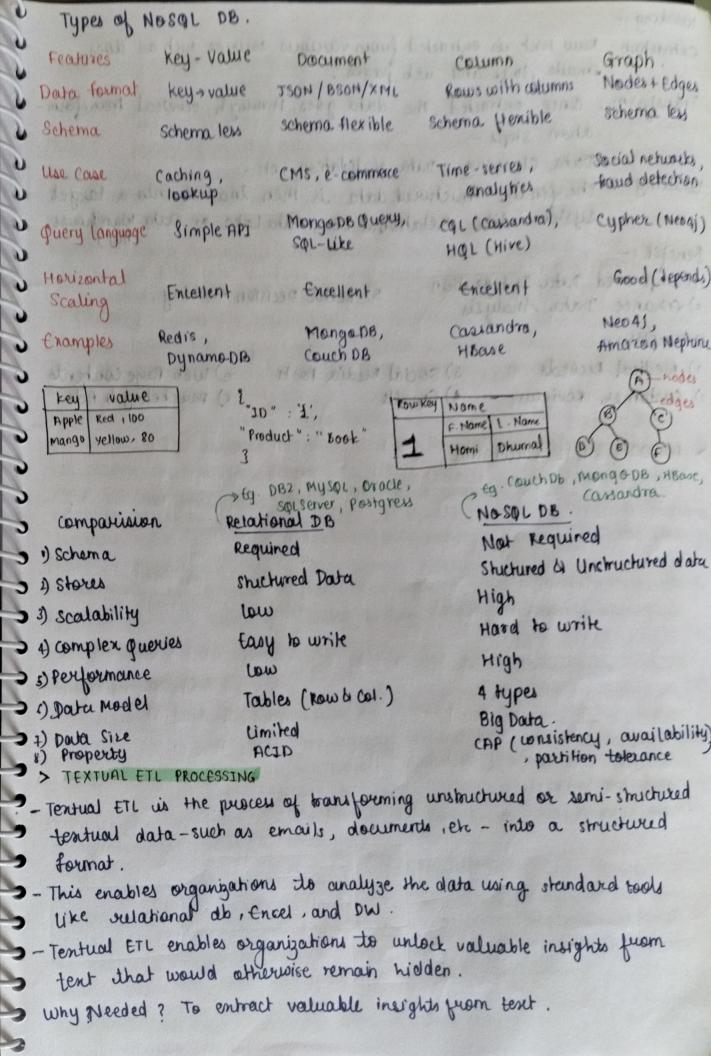
## NameNode

Metadala management

hadoop-daemon. It start namenocle start name node hadoop-daemon. Sh stop namenocle stop name node

Data Node hadoop-daemon. sh start datanode start ! hadoop-daemon sh stop datanode shop: secondary Name Node - used for taking the hourly backup of the data. In case, the Hadoop cluster fails, or crashes, the secondary Name Node will take howely backup of data & store this data with a file name fsimage. query for Edit logs in regular Secondary Name Node Updated FsImage Editlogs otek statement on dream copy the updated ssimage back to FSImage Name Node Mapkeduce is a programming model that allows us to perform parallel e distributed processing on huge datasets. - Mapkeduce convist of 2 tasks-Map () 1) Map Reduce() 2) Reduce. Map () Reduce phase takes place Map () after the mapper phase has been completed. Reduce map Tousky Tasks mapper is responsible for splitting is mapping the data, while responsible for shuffling and reducing the data. 1) Splitting phase: Input data is split into smaller chunks. 2) Mapping phase: Churks are converted into < key, value > pairs with their occurrence frequency. Here if a word occurs multiple times, it will be stored sorting phase: system performs sout & sends the output to the as single value Reducer. sorting is done based on keys, not value. 4) Reducing phase: Combines the result of outputs, which is stored in HDFS. Result Splitting List ( k2, V2) Enample:-(K1, V)) > (Bear (111)) -> Bear (2) Deer,1 pear, 2 Bear,1 7 (Deer Bear River) River,1 3 car (1,1,1) - (car (3)) car, 3 Deer Bear River car, 1 car car River Deer, 2 car, 1 River, 1 (Deer (1,1)) -(Deer (2)) River, 2 Bear) >Deer car Deer, 1 5 (River (1,1)) - River (2) cor, 1 Bear . 1

Advantages of Mapreduce: 1) Parallel Processing: - Job is divided among multiple nodes W mouster 2) Data locality: Instead of moving data to the processing unit, we are moving the processing unit to the data. 1 slave MAP-REDUCE TASKS Jab - The entire mapkeduce appin, submitted by user. Task - Jobs are split into smaller units called tasks. Map Tast - nun on input splits. Reduce Task - work on intermediate data. JOB TRACKER AND TASK TRACKER 1) Job Tracker - Acts as master node for Mapkeduce job enecution. Responsibilities - Receive MapReduce jobs from client. Communicates with NN to determine tocation of data Blocks. Monitors the perogress & status of all tooks through negular updates (heartheats ) from Took Trackers - manages resource allocation & overall job scheduling in the cluster 2) Task Tracker - Runs on ON (slave node) and enecutes tasks assigned by the Job Trackers. Responsibilities - Receives map & reduce tasks from Job Trackers and launches them as separate perocesses Monitors & manages encution of these tasks, capturing output and enit codes. - Periodically unds heartbeat signals & status updates to Job tracker to confirm it is alive and report task progresses Tast Tracker Data node 1 Task Tracker Data node 2 Fask Tracker Data node 3 Name Node > INTRODUCTION TO NOSQU All db that does not require a fixed schema for storing data & can store both structured and unstructured data. Reasons for choosing Nosqu db:a) High availability 1) Herrogenous data a) High scalability 5) High functionality 3) High performance



working:-Entraction: Row tent is entracted from various sources such as does, emails, logs and web pages. Transform: Once entracted, the data goes through several transform. ation steps -Data Cleaning, Data Validation, etc. loading: The final, smuchwed data is then loaded into a target system like Relational DB, Data Lake or Data Warehouse Benefit: -1) Improved Data Integration 3) Better Decision Making. 2) Enhance Data Analysis Encumples: -1) medical Records 5) Web Page contents
() marketing & Ads 3) Social Media Posts 2) Customer Keviews 4) Emails HADOOP SHELL COMMANDS. and water los word holy s in the best about the best of the control of the noted haverages Othlet, of C 19 set relance Dinempolusq1 1. How steps to come that the power days proportion to warmed out in 173 feeters! because it the such as encuer, declimate, the last the testing Larkach and the old opposite at continuent toldies in the lac boils of Morganise on strict stratter dealer at well-gine see island sta but of , asblety manor sales allo blazar bodt. Tool that many at great about the door of a babasay plant