

Program Understanding

Carnegie Mellon University Pittsburgh, PA 15213-3890

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Program Understanding

Outline

This lecture will cover:

- Static Analysis
- Program Slicing
- Program Plans
- [Reverse Engineering]

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Analysis

Extracts and derives information not explicitly available from data gathering

Traditional metrics

Query mechanisms for pattern matching

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Presentation

Spatial and visual data

Descriptive and depictive information

Sense integration: look, feel, etc

Some current issues:

- Integration of various visual representations
- High flexibility
- Context based visualization
- Integration of various visualization techniques

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Program Understanding – Supporting Techniques

- Static Analysis
 - Control Flow
 - Data Flow
- Program Slicing
- Program Plans

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Static Analysis - Introduction

Static analysis of code a program is the analysis of the code without regard to its execution or input.

What analysis is useful for understanding:

- Control flow analysis; what pieces of the code would be executed and in what sequence
- Data flow analysis; how does information flow within a program and across programs

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Control Flow – Introduction

Control Flow

- Used to identify the possible paths through the program
- The flow is represented as a directed graph with splits and joins
- Identify loops

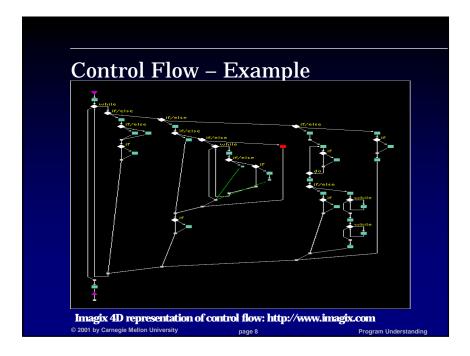
Control Flow represented as a graph of Basic Blocks

- Sequence of operations with 1-entry and 1-exit (usually a sequence of statements)
- Unique start point where program begins
- Edge between basic blocks shows the flow

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Control Flow - Code View

Another example of visualizing the control flow of a program is using a Control Structure Diagram (CSD). CSD is a algorithmic level graphical representation for software.

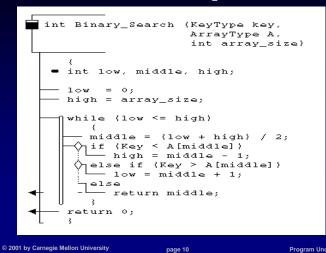
The following notations are used:

- Sequential flow straight line
- If/The/Else/Switch statements diamonds
- For/While elongated loop
- Loop exit arrow
- Function open-ended box

The GRASP project at Auburn University http://www.eng.auburn.edu/department/cse/research/grasp/

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Control Flow - Example - CSD



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Data Flow - Introduction

Data Flow is used to analyze the flow of data throughout a program and between program

Local Data Flow Analysis

- Analyze the effects of each statement
 - variable(s) defined
 - set of variable(s) referenced
- Compose the effects to derive information from beginning of each basic block to the statement

Data Flow Analysis

 Propagate basic block information over entire Control Flow Graph

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Data Flow - Statement

Suppose we have the statement S1:

S1: a = b + c

Then:

• defined(S1) = {a}

• referenced(S1) = {b, c}

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Data Flow – Types of Analysis - 1

Gen

 The set of statements where variable definitions are created in the basic block

Kill

 Set of statements that contain variables that are redefined in the basic block.

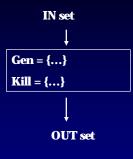
1: a := b + c 2: d := e + f 3: f := a + b Gen: {1, 2, 3} Kill: {}

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Data Flow – Types of Analysis - 1



 $\mathbf{OUT} = \mathbf{Gen} \ \cup \ \mathbf{(IN-Kill)}$

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Data Flow - Types of Analysis - 2

Def-Use (DU) Chain

 Connects a definition of a variable to all of the possible uses of the variable

Use-Def (UD) Chain

Connects a use of a variable with all possible definitions of the variable

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Data Flow - Interprocedural

Parameter Passing Mechanisms

- Call-by-value
- Call-by-reference (call-by address, call-by-location)
- Copy restore (copy-in copy-out, value result)
- Call-by-name

Procedures and Functions as parameters

Identifier scoping problems with nested procedures

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Data Flow – Difficulties

Difficulties in understanding data flow:

- Variable Aliasing (different ways to reference the same variable)
- Various parameter passing mechanisms
- Pointers

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Program Slicing – Introduction - 1

Program Slice definition:

A slice is taken with respect to a slicing criterion $\langle s, v \rangle$, which specifies a location (statement s) and a variable (v).

For statement s and variable v, the slice of program P with respect to the slicing criterion $\langle s, v \rangle$ includes only those statements of P needed to capture the behavior of v at s.

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Program Slicing – Introduction - 2

Applications of program slicing:

- understanding
- debugging
- testing
- parallelization
- integration
- software quality
- software maintenance
- software metrics

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Program Slicing – Introduction - 3

Program Slicing was first introduced by Weiser. He introduced the concept of an executable backwards static slice.

- executable slice is required to be an executable program
- backwards because of the direction the edges are traversed when computing the slice using a dependence graph
- static because they are computed as the solution to a static analysis problem (without considering the program's input)

Many applications of program slicing (such as debugging) do not require executable slices.

M. Weiser, Program Slicing, Proceedings of ICSE 1981, 439-449.

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Program Slicing - Introduction - 4

Forward slicing (introduced by Horwitz et al.)

- "What statements are affected by the value of *v* at statement s?".

S. Horwitz and T. Reps and D. Binkley, *Interprocedural Slicing using dependence graphs*, Proceedings of the ACM SIGPLAN 88 Conference on Programming Language Design and Implementation, 1988.

Dynamic Slicing (introduced by Korel and Laski)

- A slice is computed for a particular fixed input.

B. Korel and J. Laski, *Dynamic Program Slicing*, Information Processing Letters, 29(3), Oct 1988, 155-163

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Program Slicing - Flow Graphs - 1

Slicing of a flow graph is a two-step process:

- 1. Compute the data flow information
- 2. Use this information to extract a slice

To obtain the data flow information for statement *n* we first obtain:

- REF(n) the set of variables that are referenced in n
- DEF(n) the set of variables defined (given a value) in n

The data flow information is the set of relevant variables at each node *n*.

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Program Slicing - Flow Graphs - 2

For the slice with respect to <s, v> the relevant set for each node contains the variables whose values transitively affect the computation of v at s.

A statement n is in the slice if it assigns a value to a variable relevant at n and the slice taken with respect to any predicate node that directly controls n's execution.

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Program Slicing - Flow Graphs - 3

Relevant sets for the slice taken with respect to <*n*,*v*> are computed as follows:

- 1. Initialize all relevant sets to the empty set.
- 2. Insert *v* into *relevant(n)*.
- 3. For *m*, *n*'s immediate predecessor, assign relevant(m) the value (relevant(n) DEF(m)) ∪ (REF(m) if relevant(n) ∩ DEF(m) ≠ {})
- (REF(m) if relevant(n) ∩ DEF(m) ≠ {})
 4. Working backwards, repeat step 3 for m's predecessors until n_{initial} is reached

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Program Slicing - Flow Graph - 4

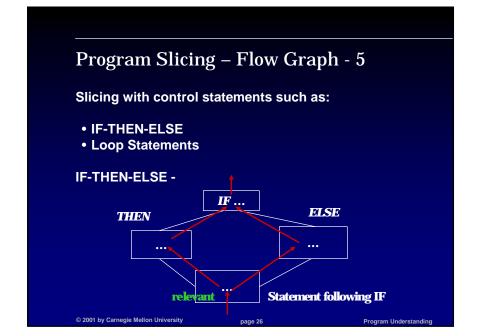
n	statement	refs(n)	defs(n)	relevant(n)
1	b = 1		b	
2	c = 2		С	b
3	d = 3		d	b,c
4	a = d	d	а	b,c
5	d = b + d	b,d	d	b,c
6	b = b + 1	b	b	b,c
7	a = b + c	b,c	а	b,c
8	print a	а		а

Slice on <8,a>: {7, 6, 2, 1}

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Program Slicing – Flow Graph - 6

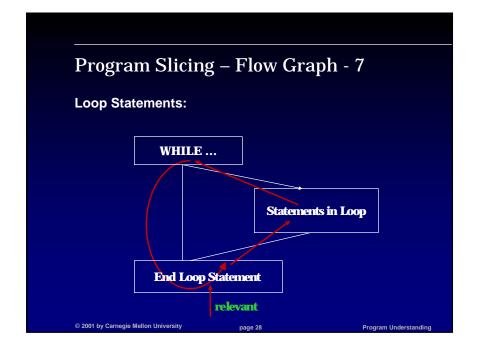
n	statement	refs(n)	defs(n)	control(n)	relevant(n)
1	b = 1	·	b		
2	c = 2		С		b
3	d = 3		d		b,c
4	a = d	d	а		b,c,d
5	if (a) then	а			b,c,d
6	d = b + d	b,d	d	5	b,d
7	c = b + d	b,d	С	5	b,d
8	else			5	b,c
9	b = b + 1	b	b	8	b,c
10	d = b + 1	b	d	8	b,c
11	endif				b,c
12	a = b + c	b,c	а		b,c
13	print a	а			а

Slice on <13,a>: {12, 11, 9, 8, 7, 6, 5, 4, 3, 2, 1}

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Program Slicing – Flow Graph - 8

Loop Example:

n	statement	refs(n)	defs(n)	control(n)	relevant(n)	relevant(n)
					Iter 1	Iter 2
1	b = 1		b			
2	c = 2		С			b
3	d = 5		d			b,c
4	a = 3		а			b,c
5	While (a < 10)	а			a,b,c	a,b,c
6	b = b + c	b,c	b	5	b,c	b,c
7	c = c + 1	С	С	5	b	b,c
8	a = b	b	а	5	b	b,c
9	EndWhile			5	а	
10	print a	а			а	

Slice on <10,a>: {9, 8, 7, 6, 5, 4, 2, 1}

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Program Slicing – Flow Graph - 9

0 While (a) 1 $x_n = x_{n-1}$ 2 $x_{n-1} = x_{n-2}$... n $x_1 = x_0$ n + 1 EndWhile

If $\mathbf{x}_{\mathbf{n}}$ in slicing criteria - need n passes through loop

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Program Slicing - Dynamic

- Dynamic Program Slicing
 - Only the dependencies that occur in a specific execution of the program are taken into account.
 - A dynamic slicing criterion specifies the input, occurrence of a statement, and a variable
 - dynamic slicing assumes a fixed input for a program whereas a static slice does not make assumptions about the input.
- Hybrid approaches that use both static and dynamic slicing also exist.

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Program Slicing - Dynamic e.g. -1

```
statement
   read(n)
   i := 1
3
   while (i <= n) do
4
   begin
5
     if (i mod 2 = 0) then
6
        x := 17
      else
8
        x := 18;
9
     i := i + 1
10
   end;
   write(x);
11
```

What is dynamic slice with criterion (n = 2, 11, x)?

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Program Slicing – Dynamic e.g. -2

n	statement
1	read(n)
2	i := 1
2 3 4 5	while (i <= n) do
4	begin
5	if (i mod $2 = 0$) then
6	x := 17
7	else
8	
9	i := i + 1
10	end;
11	write(x);

Dynamic slice with criterion (n=2, 11, x) is entire program without line 8.

Static slice (11, x) is the entire program.

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Interlude

See John Field's slides.

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Program Plans – Introduction

The goal is to recognize clichés using plans.

A cliché is a pattern that appears frequently in programs (e.g., algorithms, data structures, domain-specific patterns).

A plan is an abstract representation of a cliché.

Representation is at the semantic level rather than at the syntactic level.

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Clichés - Examples

Data structure clichés: lists, trees, tables, vectors, matrices

Algorithmic clichés: list, tree, graph traversals; iterators, applicators, manipulators; linear, binary, hash searches; event handler; exception handler

ADT clichés: dictionary, priority queue, heaps

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Plan Recognition

Approaches:

- Top-down: Start with set of goals to be achieved; determine what plans can achieve these goals; connect these plans to source code patterns.
 - Problem: Requires detailed advance knowledge otherwise connection to code is unrealistic
- Bottom-up: Start with source code; identify plans that match source code; infer higher-level goals from these plans.
 - Problem: Combinatorial explosion of alternatives
- Hybrid: top-down and bottom-up

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Plan Recognition - Method

Typical method of plan recognition

- An effective (language-independent) program representation
- A translator to transform source text into this program representation
- A library of programming plans representing clichés at various levels of abstraction
- A plan recognizer which parses the program to recognize plans stored in the library
- The result is a tree or lattice with program components at the leaves, programming plans, and the goals of the program at the root
- Bottom-up program understanding

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Plan Recognition - Issues

Syntactic variations: Recognizer works on the basis of structure information only; syntactic variations lead to the same paraphrase, modulo identifiers Non-contiguousness: Recognizer works with graph structures, can accommodate equivalent sequences of statements

Implementation variations: Similar programs are matched against the same plans, lead to the same paraphrases

Recognition algorithm depends polynomially on size of the program and plan library; graph grammars and graph recognition algorithms deployed

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Reverse Engineering Activities

The three main Reverse Engineering activities:

- Data Gathering
- Knowledge Organization
- Information Exploration

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Data Gathering

Raw data is used to identify a system's artifacts and relationships

Techniques include:

- Static source code analysis (parsing)
- Dynamic Analysis (profiling)
- Informal extraction (interviewing)

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Knowledge Organization

Goals of Knowledge Organization are:

- Efficient storage of knowledge
- Permit automated analysis
- Reflect user's perspective

Classical data models

- Hierarchical
- Network
- Relational

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Abstraction Mechanisms

Abstraction: Selective emphasis on detail

Common Mechanisms:

- Classification
- Aggregation
- Generalization

Conceptual Modeling

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Information Exploration

Probably the most important activity:

- Data gathering: necessary to begin
- Knowledge organization: structure model
- Information Exploration: understanding

Composite Activities:

- Navigation
- Analysis
- Presentation

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Navigation

Traverse non-linear information structures

Link relationships:

- Component hierarchies
- Inheritances
- Control and data flow

Hypotheses postulation \Rightarrow exploration

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