

### 3.4 - Hashing

Sunday, July 13, 2025 11:29 AM

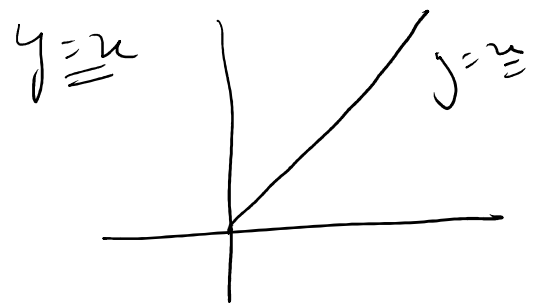
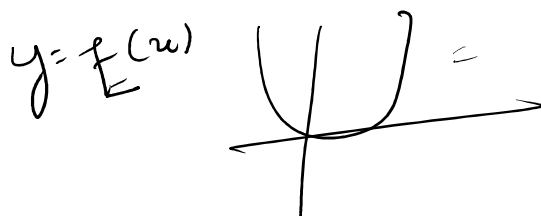
- Linear  $- O(n)$
- Binary  $- O(\log n)$

→ Technique which helps you perform search in an optimized manner ( $O(1)$ )

# Functions

$y = f(x)$   $\rightarrow$  independent variable  
 $y$   $\rightarrow$  dependent variable

$y = x^2$   $\rightarrow$  parabolic function  
 $x=1, y=1$   
 $x=2, y=4$



$$y = x^{0.10}$$

$$\Rightarrow x=13, y=3.1$$

$$x=21, y=1$$

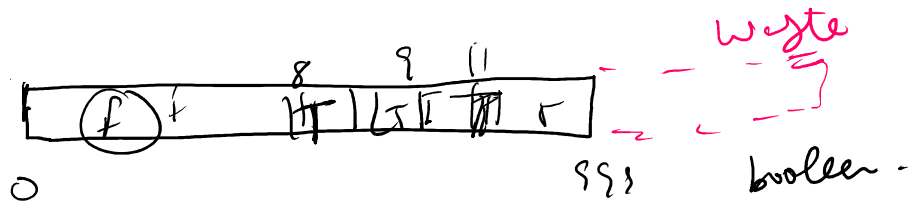
$$y = x^{0.10 \pm \epsilon} \Rightarrow [0, \epsilon-1]$$

$$\frac{6/0.10}{0.09}$$

# Concept of Hashing

(2)

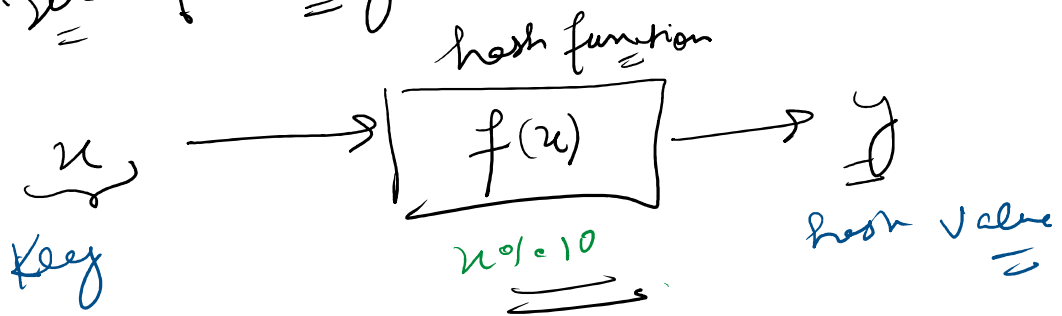
$\Rightarrow \{ \textcircled{8}, 9, 11, 10, 5 \}$



$O(1)$  if (all  $[u] == \text{true}$ ) present  
else not present

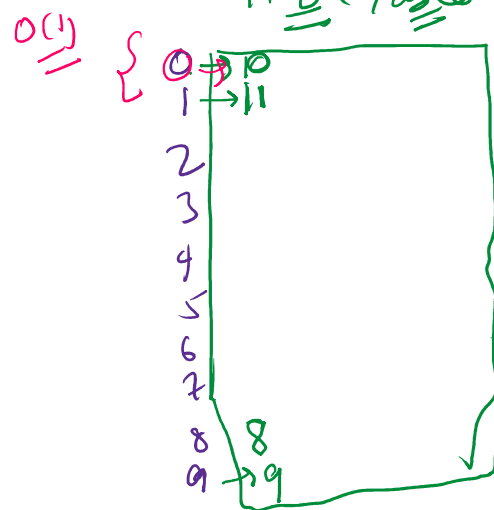
$\Rightarrow$  Technique of Hashing

$\textcircled{0}$   $\textcircled{18} \rightarrow$



$\{ 8, \textcircled{11}, \textcircled{10}, \textcircled{9} \} \rightarrow \{ 8, 1, 0, 9 \}$   
keySet

$\textcircled{0} \textcircled{10} \rightarrow [u] = 0$



$\Rightarrow$  size of Hash Table depends on hash function

$$n = n \cdot \epsilon \Rightarrow \lceil n \cdot \epsilon \rceil$$

$$f(x) = \frac{u \cdot x \cdot \varepsilon}{\varepsilon} \Rightarrow [0 - \varepsilon - 1] \rightarrow \textcircled{\varepsilon}$$

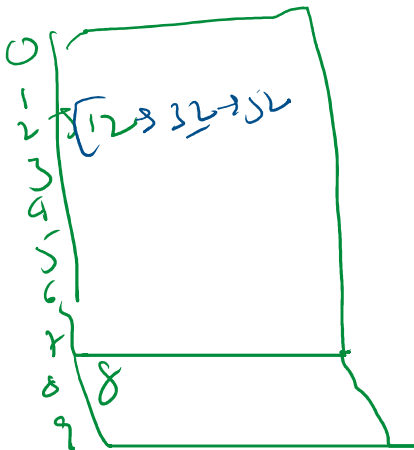
⇒ Collision

→ same values for diff multiple keys

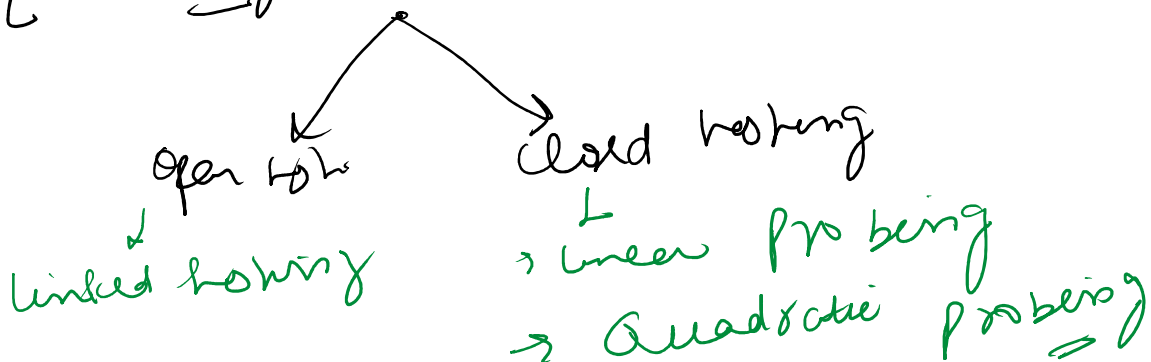
$$\{8, 12, \underline{32}, 9, 11\} \rightarrow \boxed{u=110}$$

$$\downarrow$$

$$\{8, \underline{2}, \textcircled{2}, 9, 11\}$$



# [Handling collision]



• → 0 → 0 → 0 → 0

# If the no. of collisions are high, then  
the complexity in open hashing goes to  
 $O(n)$ .

# Closed hashing

→ Linear probing ⇒

$$h(u) + i \Rightarrow i$$

$$f(u) = [h(u) + i]$$

$$h(u) = u \% 10$$

$$f(u) = [u \% 10 + i] \Rightarrow i = 0, 1, 2, \dots$$

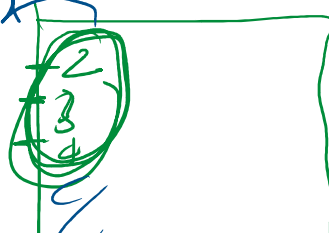
Ex { 8, 12, 7, 22, 14, 21 } 10, 4, 22, 4

$$f(u) = [u \% 10 + i]$$

$$f'(u) = [u \% 10 + 1]$$

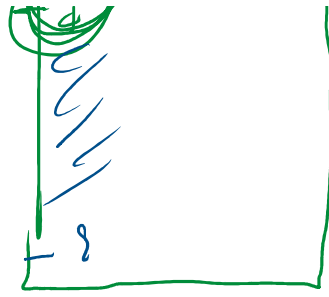
$$f'(u) = [u \% 10 + 1]$$

{ 22, 32, 42, 2 } clustering



0	
1	21
2	→ 12
3	→ 22
4	→ 14
5	→ 104
6	
7	→ 7
8	→ 8
9	

Problems of linear  
probing



(Clustering)

# Load Factor

⇒ No. of entries / size of the hash table

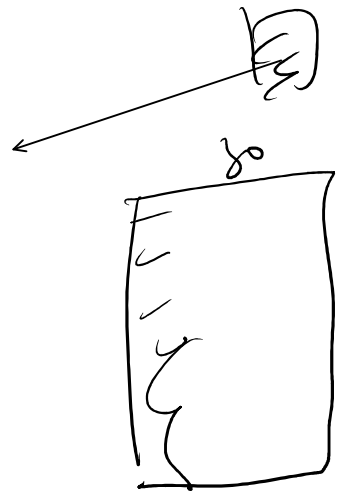
⇒ a good hash func<sup>n</sup> keeps the load factor below 0.75.

⇒ when load factor becomes large, we do rehashing

$n = 15$

→

$n = 1.80$

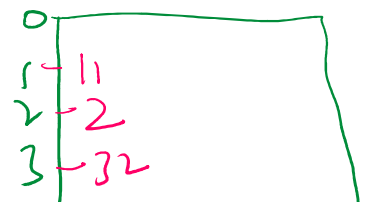


⇒ Quadratic Probing

$$f(u) = \left[ \underbrace{h(u)}_{n \cdot 10} + \underbrace{c^2}_{c=0,1,2,\dots} \right] \% 10$$

{ 8, 12, 11, 32, 14, 102 }

$$9 + 16 = 25 \pmod{10} = 5$$



$\{8, 12, 11, 32, 14, 102\}$

$2 + 0^2$   
 $2 + 1^2$   
 $2 + 2^2$

$2 + 1^2 = 3$

$J+1$   
 $J+4$   
 $J+5$   
 $J+16$

2  
 32  
 14  
 102  
 8

$\{22, 72, 42, 52\}$   
 9

2  
 3  
 6

# Data Structures in Java based on the concept of Hashing

$\Rightarrow$  HashSet  
 $\Rightarrow$  HashMap

Based on the concept of Hashing

↓  
Based on the concept of Hashing  
HashTable

## # Hash Set

⇒ DS which only stores unique values.

{ 2, 3, 5, 2, 4, 3 }

↓

{ 2, 3, 5, 4 }

⇒ Hash set stores values in random order

TC

Insertion  $\rightarrow O(1)$

Search  $\rightarrow O(1)$

Delete  $\rightarrow O(1)$

## # Hash Map

↳ collection of key-value pairs.

{ Rohit  $\rightarrow$  40  
Loki  $\rightarrow$  60 } } key-value

{ Rohit  $\rightarrow$  40  
Kohli  $\rightarrow$  60  
KL Rahul  $\rightarrow$  0

} key-value

$\Rightarrow$  keys are always unique  
but values can be duplicate

T.C  
insertion  $\rightarrow O(1)$   
read  $\rightarrow O(1)$   
deletion  $\rightarrow O(1)$

$\Rightarrow$  How to iterate on Hash Map

$\Rightarrow$  mp.keySet()  
 $\Rightarrow$  mp.values()

$\rightarrow$  for each loop  
because there is no  
concept of indexing.

map.put (key, value)  
map.get (key)  $\rightarrow$  <sup>return</sup> value  
null

⊕ while doing a get operation on a mp,  
always first check whether the key is



⊕) When ... always first check whether the key is present or not.

int u = map.get (key)  
↓  
need pointer-  
exception