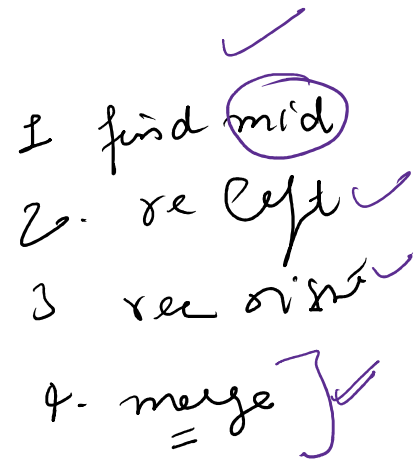


Saturday, July 5, 2025 10:41 AM



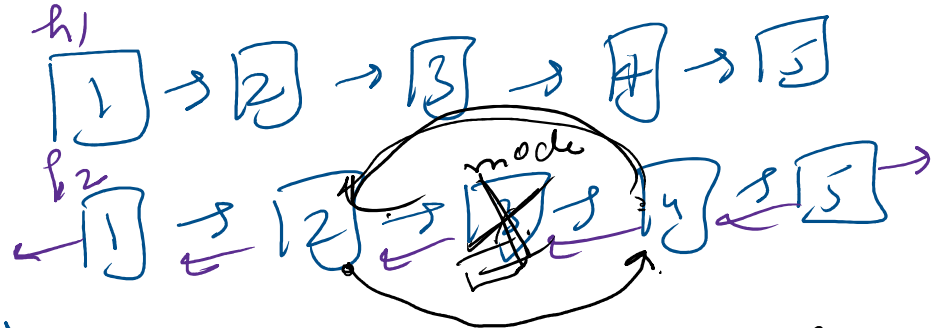
$\text{null} \leftarrow [1] \rightarrow [2] \rightarrow [3] \rightarrow [4] \rightarrow [5] \rightarrow \text{null}$

node of
in data
nearest.
Node prev;
1

4. ~~next~~

Time complexity analysis of LL

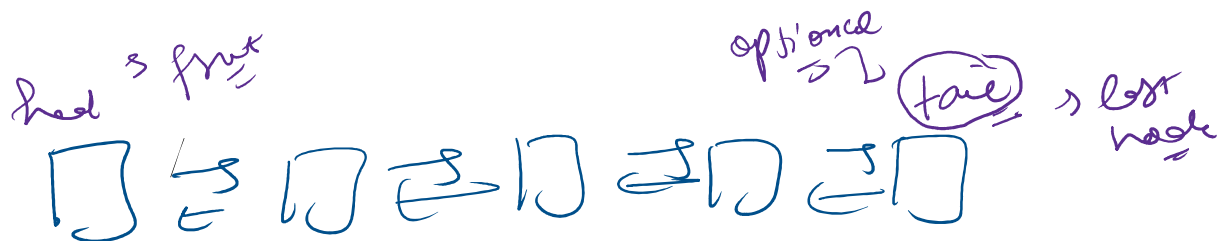
③



$1 \rightarrow 2 \rightarrow \text{next} \rightarrow 3 \rightarrow 4 \rightarrow 5$
 $\text{node} \cdot \text{prev} \cdot \text{next} = \text{node} \cdot \text{next}$
 $\text{node} \cdot \text{next} \cdot \text{prev} = \text{node} \cdot \text{prev}$

$\rightarrow O(1)$

\rightarrow Deletion in doubly LL is $O(1)$, and more efficient than array, array list and singly LL. But it comes with an assumption that input node is given to us.

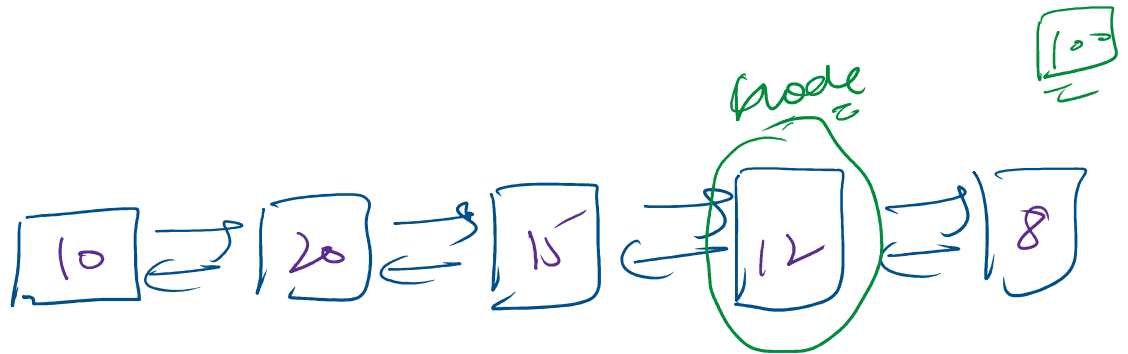


\rightarrow Edge cases.

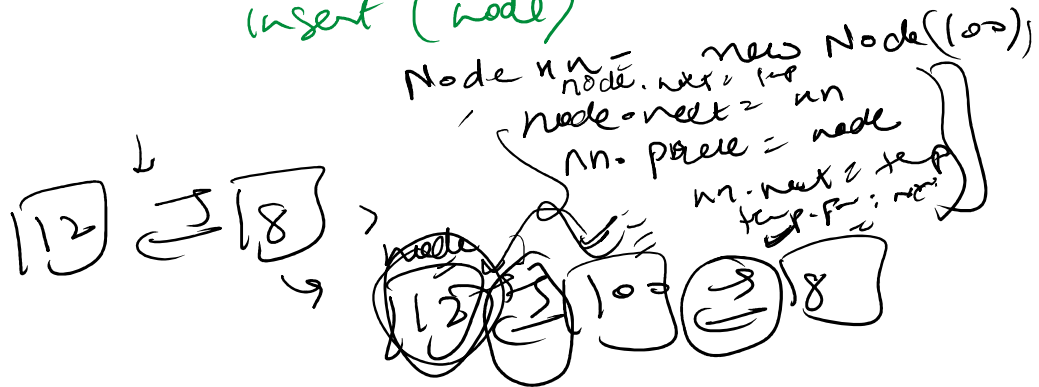
⇒ Edge Cases:

- 1) head == null (empty list)
- 2) head == tail (single list)

- 3) head (delete first node)
- 4) tail (deletion of last node)

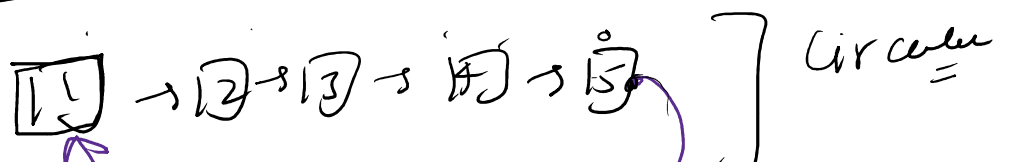


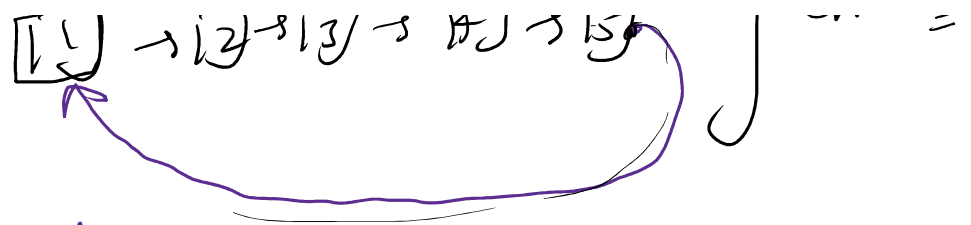
insert (node)



Circular LL

1 → 2 → 3 → 4 → 5 → 5





① Last ^{the next} pointer always points to head

SLL while (curr != null)

CLL while (curr.next == head)