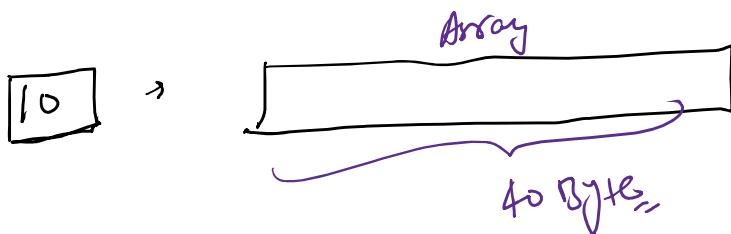


} → Data Structure

#



int arr[] = new int[10]

(4)

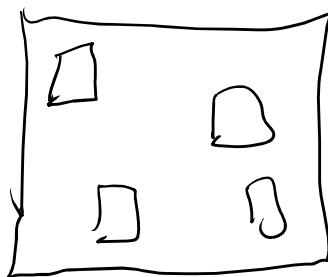
ArrayList →



dynamic size

continuous array:

(10) → 40 Bytes

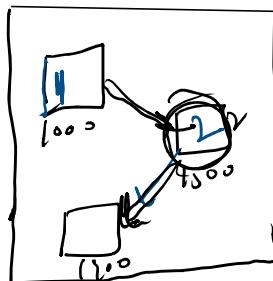


[40 Bytes]

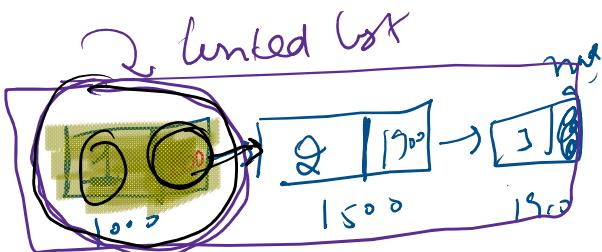
→ Linked list solves the problem of fragmented memory:

$$3 \times 4 = 12$$

J

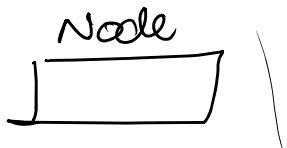


1, 2, 3

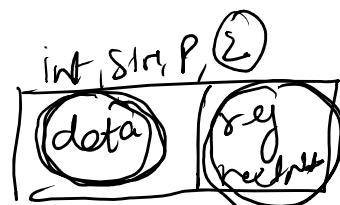


Node

⇒ Node: Basic Building block of a linked list

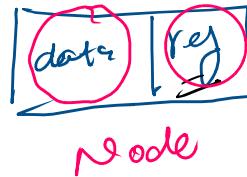


class of  
int data



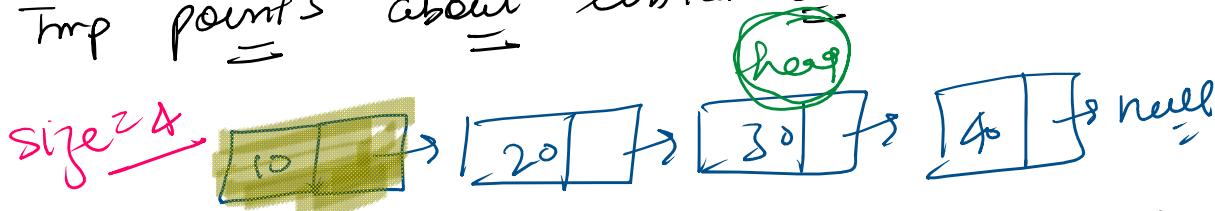
Person P is not Person  
try 6) { }  
int data;  
Node next;

```
class Node {  
    int data;  
    string str;  
    Node next;  
}
```



Person  $p$  = new Person()

# Temp points about = listed list

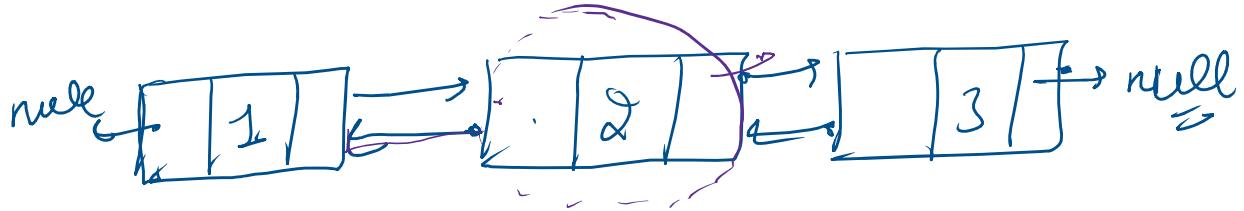


\* size of linked list = no. of Nodes in the LL.

→ first Node is called head of the L-List

→ We can only access a linked list via the head node

→ A LL in which we can only move in a single direction is called a singly linked list.



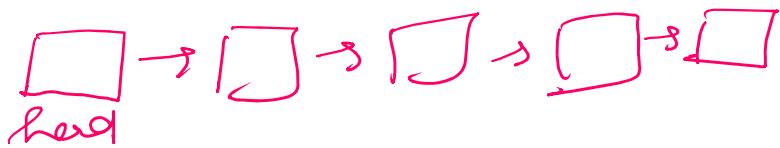
→ A linked list in which we can move in both the directions is called a doubly LL

```
class Node {
    data
    {Node next;
     & Node prev;
}
```

}

# Strategy to solve linked list questions:

1) Always draw a general LL diagram of 4-5 nodes.

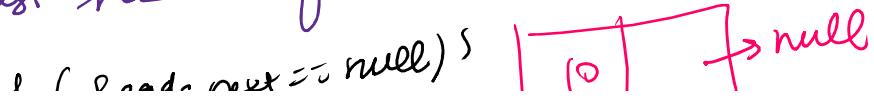


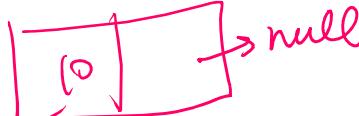
2) Dry run the movement of the head pointer.

3) Think of two edge cases

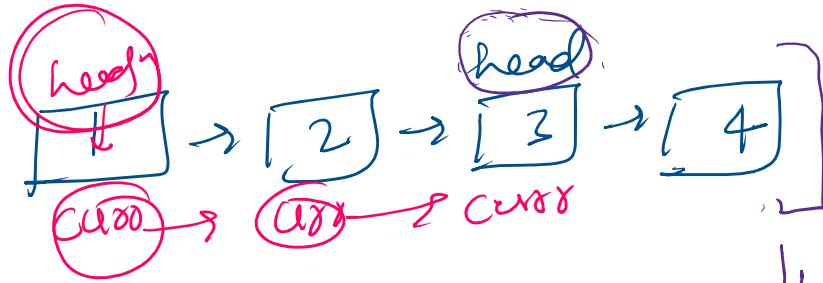
→ if list is empty  
if (head == null);

→ if list has only one node  
if (node.next == null);



$\rightarrow$  if  $exn \neq \text{null}$   
 if ( $\text{head} \cdot \text{next} = \text{null}$ )  
 }  


④ Vo Imp



→ generally we don't prefer to move head because  
 it is the only way to access our  
 LL.

→ we make a copy head ('curr') and move it to