An Euler tour of a strongly connected, directed graph G=(V,E) is a cycle that traverses each edge of G exactly once, although it may visit a vertex more than once. Show that G has an Euler tour if and only if indegree(v)=outdegree(v) for each vertex $v \in V$.

Proof:

Sufficiency Let G_n be a strongly connected directed graph with n vertices, for which indegree(v)=outdegree(v) for each vertex $v \in V$.

Base case: G_2 has a Euler tour. Let the vertices be u and v As we have a graph of two vertices with indegree(u) = outdegree(u) and indegree(v) = outdegree(v), we can travel from u to v and then back to u k times, resulting in an Euler tour

Hypothesis: Lets assume that there exists an Euler tour in any arbitrary graph G_n .

Inductive step: We have to prove that G_{n+1} has an Euler tour. Every vertex in G_n+1 has equal indegree and outdegree. So, for any vertex V, we can arbitrarily pair up all of V's in edges and out edges. Now we can remove vertex V from G_n+1 and for every pair of edges (u,V) and (V,W) we add an edge (u,W) in G_n . In this way, in-degree and out-degree of U and U remains the same. Thus, there is an Euler tour in G_n

(Inductive Hypothesis). Let the Eulerian cycle be $C=x_0,x_1,...,u,w,...,x_0$

Now we add the vertex V back to Gn and traverse edge (u, V) and then (v, w) in the place of (u, w) in C. The resulting path $P=x_0, x_1, ..., u, v, w, ..., x_0$ is also an Eulerian cycle in Gn+1. Thus, Gn+1 has an Euler tour. Thus G has an Euler tour if indegree(v)=outdegree(v) for each vertex $v \in V$

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