

Bubble Blast

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Abstract

The article is about game playing algorithm for the game known as Bubble Blast. The article tries to explore different strategies for maximizing the score obtained in the game. Many algorithms are presented and it is shown that the algorithm with strategy of minimum color block first performs significantly better than all its counterparts. Experimental results are presented that give insight into the size of blocks with respect to the number of colors and size of grid as well as analysis of time taken to solve the game.

I. INTRODUCTION

Bubble Blast is a well known one player puzzle game. It was initially released as "Chain shot" in 1985. There has been other versions of the game with slightly different or same rules namely Clickomania, HMaki, Samegame, Jawbreaker, Bubbles, Bubble breaker etc. It has also been ported to numerous computer and mobile platforms.

II. RULES

The typical version of the Bubble burst game consists of a gameboard consisting of a screen of differently colored bubbles in a matrix. The gameboard is a typical $N \times N$ matrix with each position in the matrix having a corresponding color. In the bubble burst version of the game, different positions on the matrix are placed with different colored bubbles. The player's move is to click on a particular bubble and eliminate the block of that particular block if one exists.

I. What is a block ?

If two bubbles are of the same color and they are connected, then they are part of the same block. This relation is transitive in nature, for example, if a bubble A and B are part of the same block and bubbles B and C, then bubbles A and C are also part of the same block.

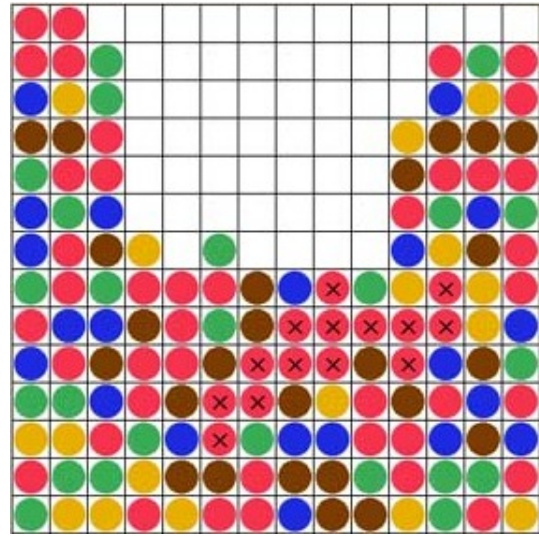


Figure 1: A typical setup of Bubble burst game showing a grid of 14×14 with 5 color of ball used. The bubbles with cross on them show the bubbles of the selected block.

In figure 1 we can see a block shown by the cross mark on the member bubbles of the block.

II. What is connectivity ?

There are two types of connectivity that can be considered :

- **4 Neighbour Connectivity** : In this connectivity only 4 bubbles are considered connected to a particular bubble. The bubbles considered connected are vertically upwards, vertically downward, adjacent left and adjacent right bubble to the given

bubble.

- **8 Neighbour Connectivity** : In this connectivity along with those considered in 4 neighbour, four extra bubbles are considered connected. These four are diagonally upward-left, diagonally upward-right, diagonally downward-left and diagonally downward-right.

For the purpose of our study, we will consider neighbour connectivity only.

III. What happens after elimination of block ?

The player when clicks on any bubble, if the bubble is part of any block, then all the members of that block are eliminated. After the elimination of a bubble, to fill the void created by the elimination all the bubbles upwards of the void are pulled down until no void is left in that vertical column. This ensures that there is no void in the matrix or in other words there is no position on gameboard which is empty but any position vertically above has bubble present. Another situation is what happens when the whole column becomes empty. Here we will shift all the leftwards column of an empty column by one spot to the right to fill the empty column. This way the empty column is shifted to leftmost corner.

IV. Scoring

There are many versions of this game each having its own scoring some having n^2 , $n(n-1)$ or $n(n-2)$, where n is the size of the block eliminated. One thing common among all these scoring is that in all the versions the scoring depends on the square of block size.

For the purpose of our study, we will consider the score to be n^2 .

V. Game Ending

A standard game ends when the player has no moves left or in other words there are no other

blocks present on the gameboard.

III. ALGORITHM

The strategy for the game is defined as any rule that gives the priority order by which block has to be eliminated so as to maximize the score. There are different strategies which can be applied to this game so as to maximize the score.

Our Algorithm uses the "**Minimum Color Block First**" strategy. In this particular strategy, first of all the colors are sorted in the order of their number of occurrence. The sorted order of colors gives us the priority order of the blocks which has to be followed for the entire game. For example, if the order by number of occurrence initially is red, blue, green, yellow then during the game whenever a block of red color is present it will be the first one to be eliminated followed by blue, green and yellow. This given priority of color is followed throughout the game.

For the purpose of comparison, the algorithm is compared with other algorithms which use different strategies, namely :

- **Random Blocks**: In this particular strategy any block is selected at random and that particular block is eliminated. As this strategy is random the score obtained is generally unpredictable and has high variance in the score.
- **Maximum Size Block First**: In this particular strategy the block with the minimum number of member bubbles in it is given highest priority.
- **Minimum Size Block**: In this particular strategy the block with the maximum number of member bubbles in it is given highest priority.

IV. RESULTS

There are many results of the game we analyze with respect to changing parameters and game playing strategies. First of which is :

I. Average Block Size

The average block size is the average of the block size of all the blocks present on the gameboard initially before making any moves. It depends on both the numbers of colors used and the size of the gameboard. Here we have found the average block size by taking the average of 100 random gameboard configuration for a given parameter.

We vary both number of color used and size of the matrix and then note the resulting change in the average block size.

Given below is the average block size for different color on a 25×25 board and :

Table 1: Average block size for given number of color

Number of color used	Average block Size
2	10.677
3	4.367
4	3.32
5	2.915
5	2.478
10	2.344

The relation between number of colors and average value of block size is plotted on a graph in figure 2.

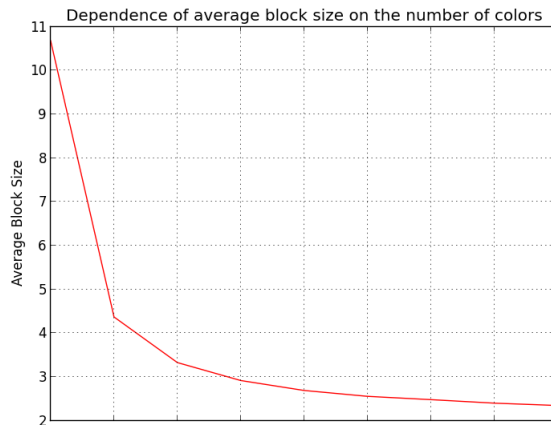


Figure 2: The relation between average block size and number of color.

It is evident from Figure 2 that as the number of color increases the average block size decreases. The average block size decreases steeply initially at low values for number of

color but as the number of colors used the change becomes very small.

Given below is the average block size for different size of the board using 5 colors.

Table 2: Average block size for given size of the board

Size of Grid	Average block Size
1×1	0.0
2×2	1.236
3×3	2.371
4×4	2.692
5×5	2.716
10×10	2.820
15×15	2.876
20×20	2.908
25×25	2.912

The relation between number of colors and average value of block size is plotted on a graph in figure 2.

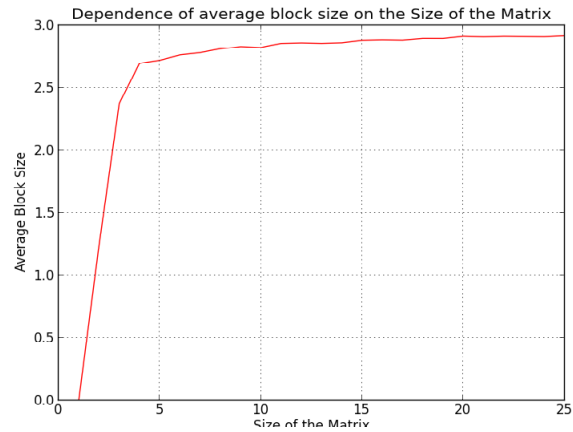


Figure 3: The relation between average block size and number of color.

It can be observed from Figure 3 that as the size of the gameboard increases there is a steep increase in the average block size initially but after at larger gameboard average block size start approaching a constant value.

II. Time Complexity

It is important to know the time complexity of an algorithm for its analysis. We can infer

about the time complexity by the average time taken by it for different cases. It has already been proved in [1] that finding the maximum score sequence of step is NP-complete. In the table we have the average time for 50 cases for a given $N \times N$ matrix for each N .

Table 3: Average time for a $N \times N$ matrix with varying N

Size of the Matrix	Average time taken in ms
3×3	0.599
4×4	1.75
5×5	3.24
6×6	7.499
10×10	46.650
15×15	193.876
20×25	489.399

As the average time taken increases steeply with increase in size N , the graph between log of time taken and size N is plotted.

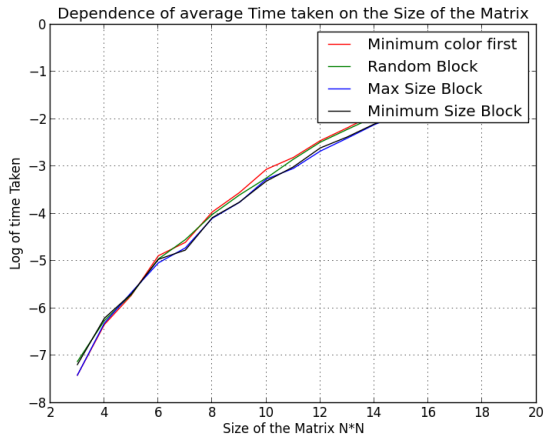


Figure 4: The relation between average time taken and size N in a $N \times N$ gameboard.

It can be easily observed from figure 4 that the log of average time taken is linear to the size N for all the algorithms tried. Hence it can be said that average time taken is exponential to N for a $N \times N$ gameboard.

So the time complexity of each algorithm compared is exponention to N for a $N \times N$ gameboard.

III. Score

Our main aim is to maximise the score obtained, therefore the score obtained for different algorithms is tested on different parameters.

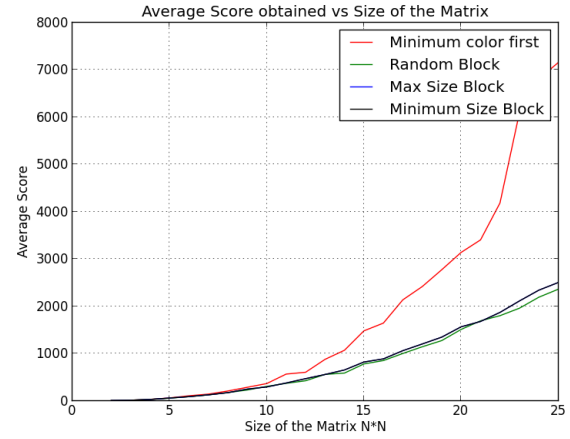


Figure 5: The relation between average score and size N in a $N \times N$ gameboard.

It can be observed from the figure 5 that our algorithm score significantly better than all other algorithms on all varying size of gameboard.

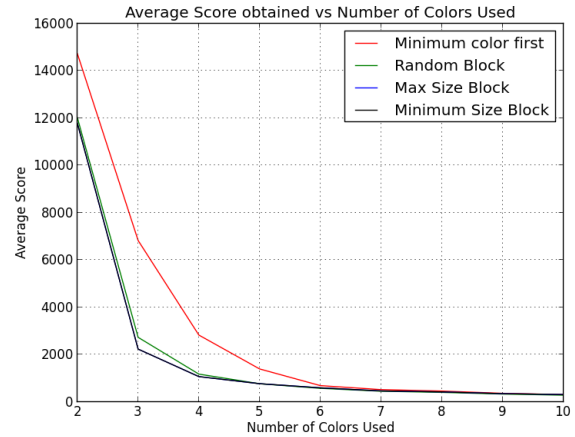


Figure 6: The relation between average score and size N in a $N \times N$ gameboard.

It can be observed from the figure 6 that our algorithm score significantly better than all other algorithms on all varying size of gameboard. Finally the algorithms are tested on 100 different configuration of a 25×25 gameboard and the average of all scores is compared.

Table 4: Results for a given algorithm on 25×25 matrix with 5 colors

Algorithm	Score	S.Deviation	Time
MinColor	6839.58	2348.061	1.621
Random	2407.8	228.582	1.548
Max Size	2508.14	162.268	1.161
Min Size	2576.14	180.268	1.184

From the table given we can confirm that the algorithm with strategy of "**Minimum Color Block First**" perform better than algorithm with other strategies in general.

V. CONCLUSION

I. Time Complexity

It is evident that the time complexity is exponential to N for a $N \times N$ gameboard for all

of the algorithms tried. The time complexity could not be improved beyond this.

II. Algorithms

Among all the algorithms tried it can surely said that our algorithm has the best score on varying size of the gameboard and number of color used.

REFERENCES

- [1] M.P.D. Schadd, M.H.M. Winands, H.J. van den Herik and H. Aldewereld. Addressing NP-complete puzzles with Monte-Carlo methods. In *Proceedings of the AISB 2008 Symposium on Logic and the Simulation of Interaction and Reasoning, Volume 9*, pages 55 to 61, The Society for the Study of Artificial Intelligence and Simulation of Behaviour, 2008.