

Towards Acceptors and Transducers for Semantics

kevin knight

columbia university, february 2012

Machine Translation

Phrase-based MT

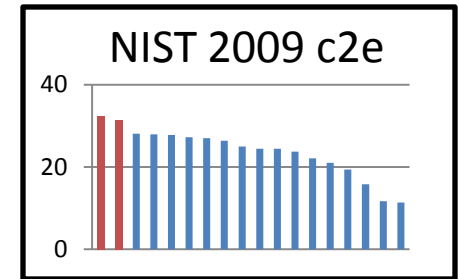
source string → target string

Syntax-based MT


source string → source tree → target tree → target string

Meaning-based MT

source string → source tree → meaning representation → target tree → target string



Meaning-based MT

- What content goes into the meaning representation? 1
 - linguistics, annotation
 - How are meaning representations probabilistically generated, transformed, scored? 2
 - automata theory, efficient algorithms
 - How can a full MT system be built and tested? 3
 - understanding, generation, rule extraction, language modeling, features, training, engineering
- these slides
- 

Automata Frameworks

- How to represent and manipulate linguistic representations?
- Linguistics, NLP, and Automata Theory used to be together (1960s, 70s)
 - Context-free grammars were invented to model human language
 - Tree transducers were invented to model transformational grammar
- They drifted apart
- Renewed connections around MT (this century)
- Role: greatly simplify systems!

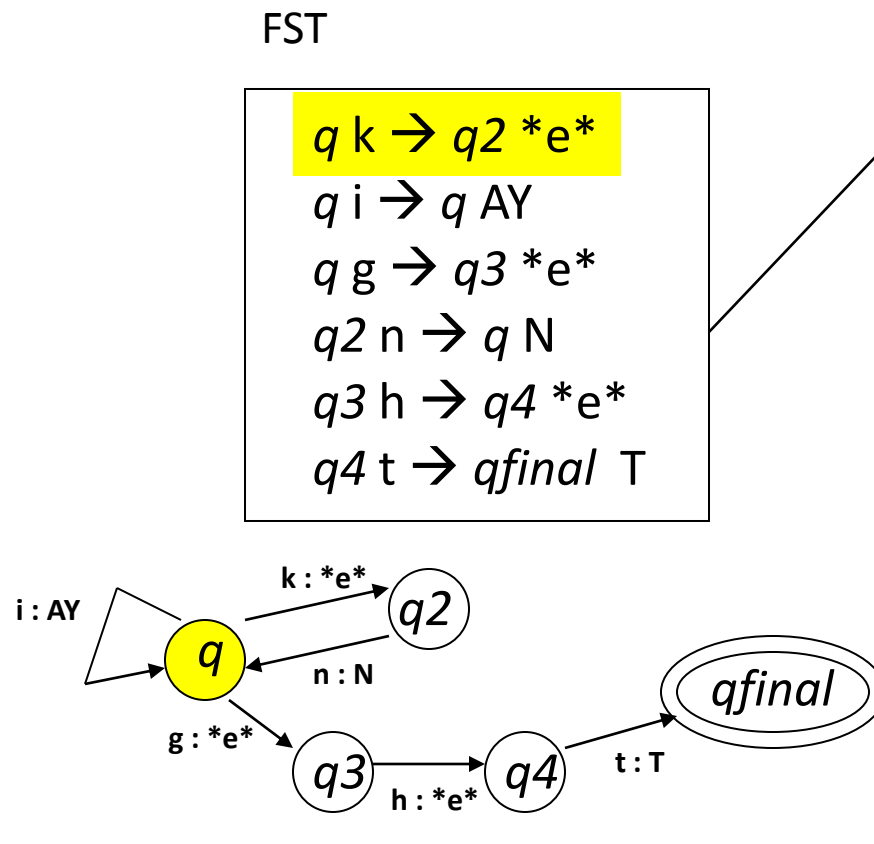
Finite-State Transducer (FST)

Original input:

k
|
n
|
i
|
g
|
h
|
t

Transformation:

q k
|
n
|
i
|
g
|
h
|
t

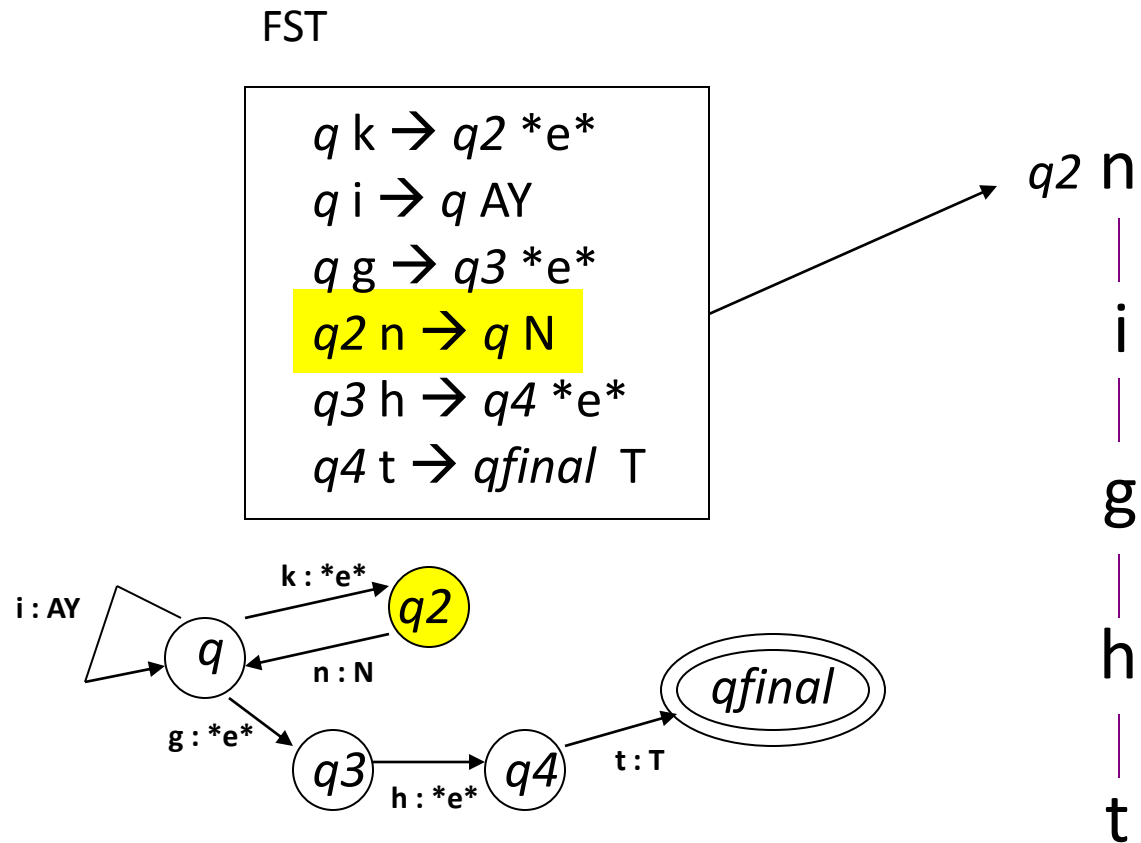


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Transformation:



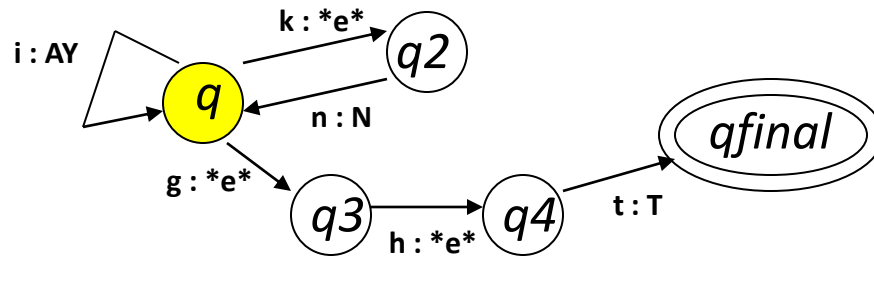
Finite-State Transducer (FST)

Original input:

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i
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g
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t

FST

$q \ k \rightarrow q2 \ *e^*$
 $q \ i \rightarrow q \ AY$
 $q \ g \rightarrow q3 \ *e^*$
 $q2 \ n \rightarrow q \ N$
 $q3 \ h \rightarrow q4 \ *e^*$
 $q4 \ t \rightarrow q_{final} \ T$



Transformation:

N
|
 $q \ i$
|
g
|
h
|
t

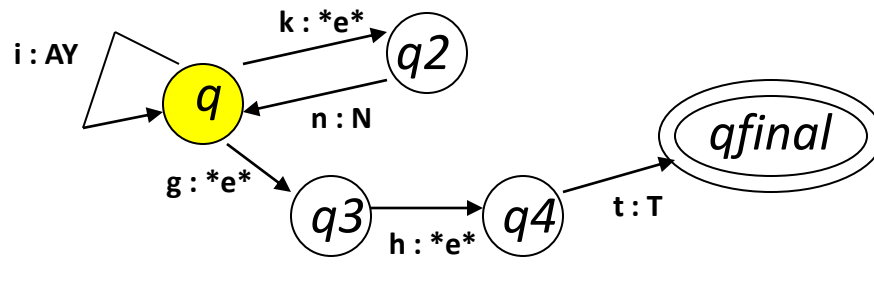
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FST

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 $q \text{ i} \rightarrow q \text{ AY}$
 $q \text{ g} \rightarrow q3 \text{ *e*}$
 $q2 \text{ n} \rightarrow q \text{ N}$
 $q3 \text{ h} \rightarrow q4 \text{ *e*}$
 $q4 \text{ t} \rightarrow q_{final} \text{ T}$



Transformation:

N
|
AY
|
 $q \text{ g}$
|
h
|
t

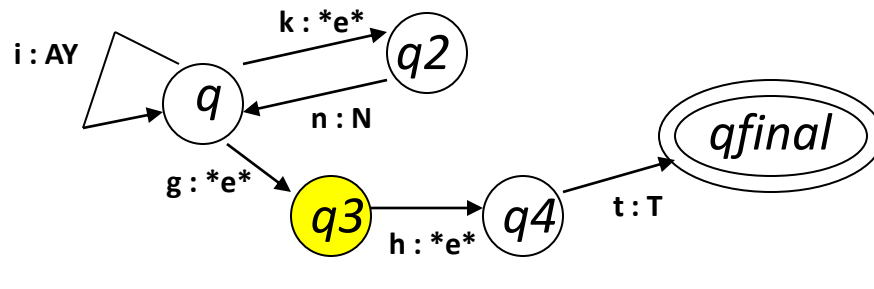
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Transformation:

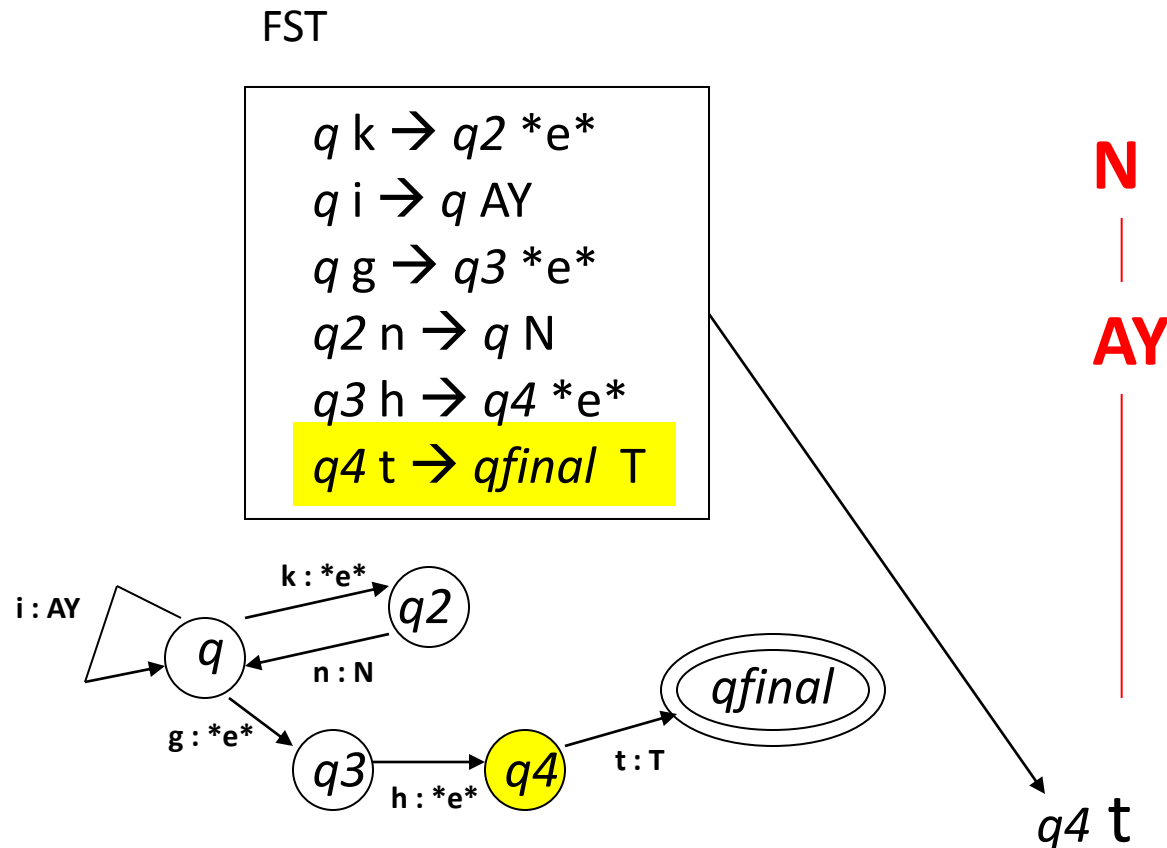
N
|
AY
|
q3 h
|
t

Finite-State Transducer (FST)

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|
t

Transformation:

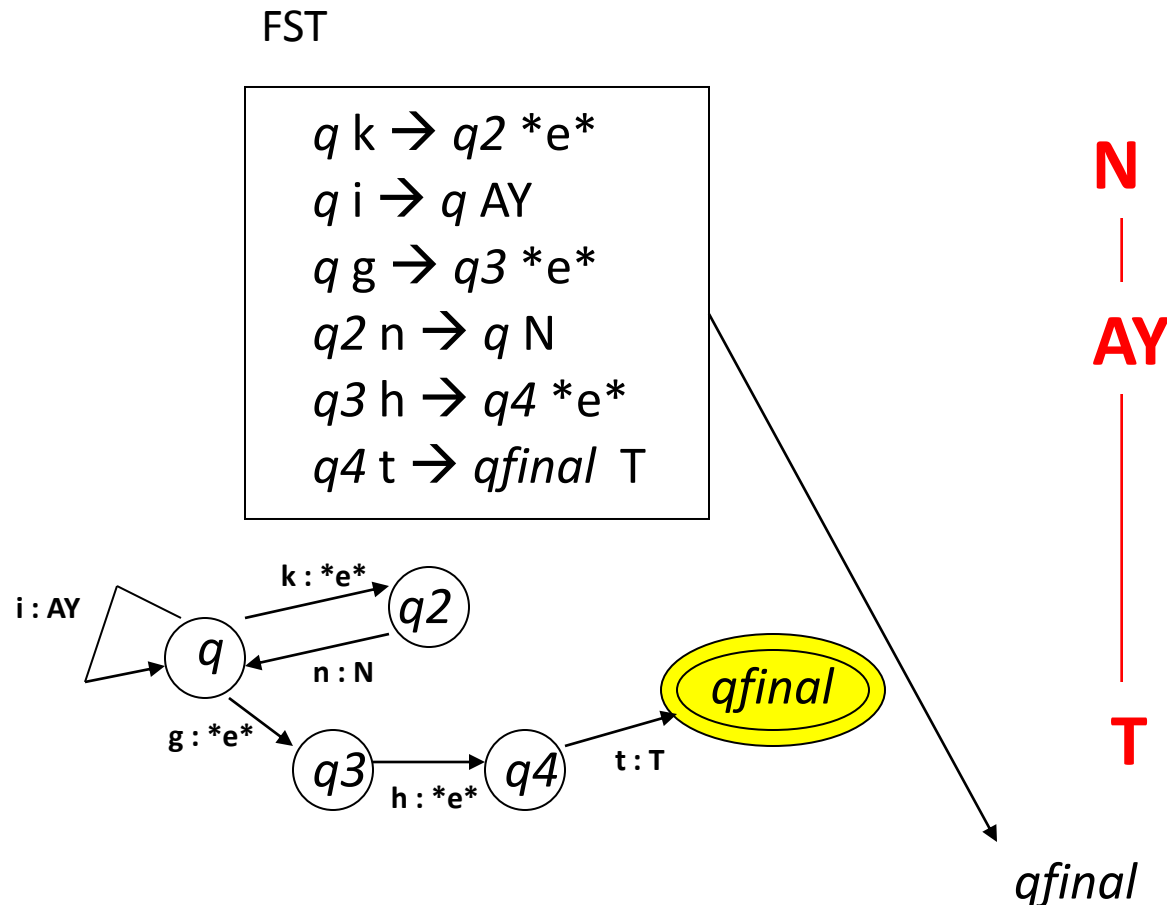


Finite-State Transducer (FST)

Original input:

k
|
n
|
i
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h
|
t

Transformation:



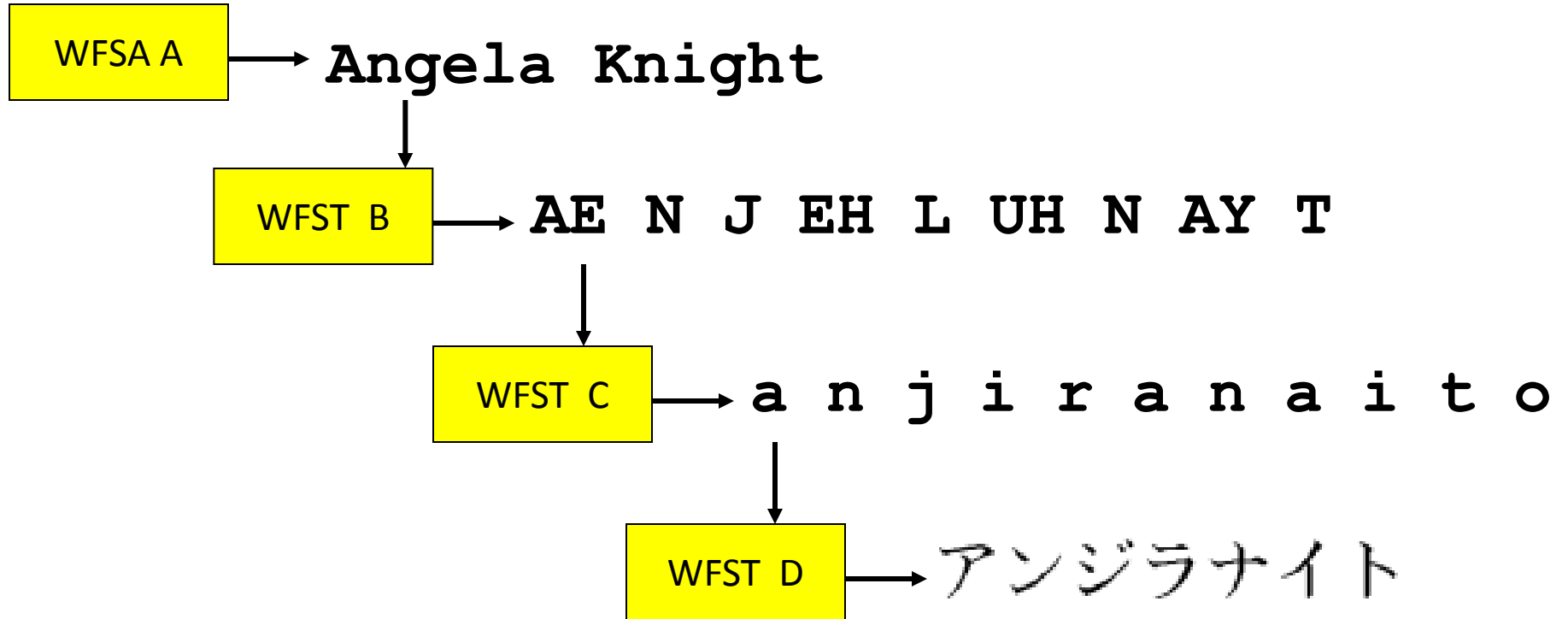
Transliteration



7 input symbols

13 output symbols

Transliteration



Angela Knight

Angela Nite

Andy Law Knight

Angela Nate

+ millions more

WFSA A

Ann Gere Uh

Anne Jill Ahh

Angy Rugh

Ann Zillah

+ millions more

WFST B

AE N J IH R UH N AY T

AH N J IH L UH N AY T OH

+ millions more

WFST C

a n j i r a n a i t o

WFST D

アンジラナイト

DECODE

$$\operatorname{argmax}_E P(E|K) =$$

$$\operatorname{argmax}_E \sum_{es} \sum_{js} P(E) \cdot P(es|E) \cdot \\ P(js|es) \cdot P(K|js)$$

Finite-State String Transducers

- non-deterministic
 - invertible
 - composable
- etc

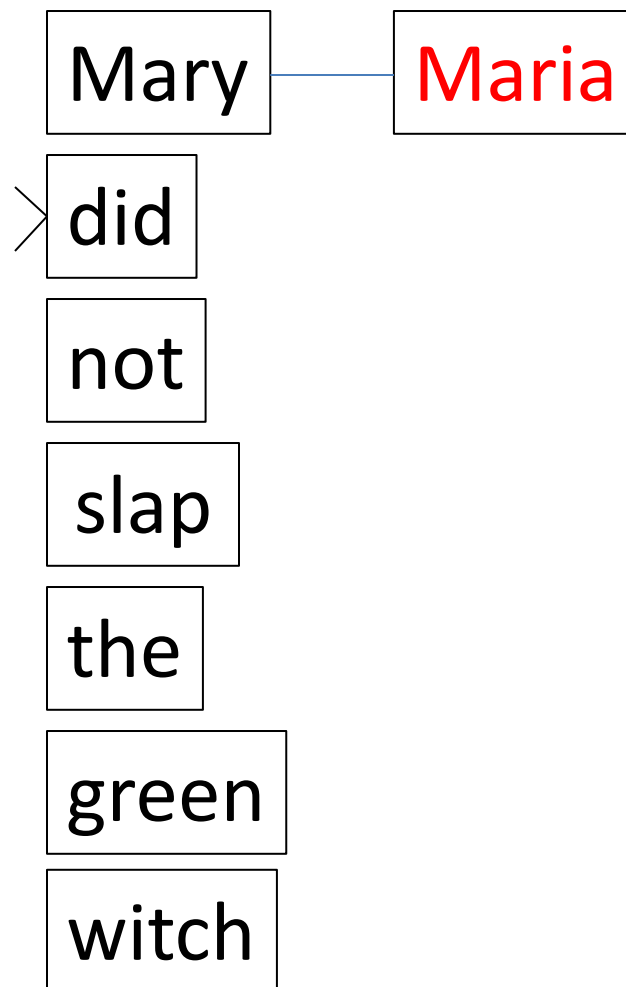
General-Purpose Algorithms for String Automata

N-best paths through an WFSA (Viterbi, 1967; Eppstein, 1998)
EM training	Forward-backward EM (Baum & Welch, 1971; Eisner 2001)
Determinization of weighted string acceptors (Mohri, 1997)
Intersection	WFSA intersection
Application	string \rightarrow WFST \rightarrow WFSA
Transducer composition	WFST composition (Pereira & Riley, 1996)
General-purpose toolkit	FSM (AT&T), Carmel (USC/ISI), OpenFST (Google), ...

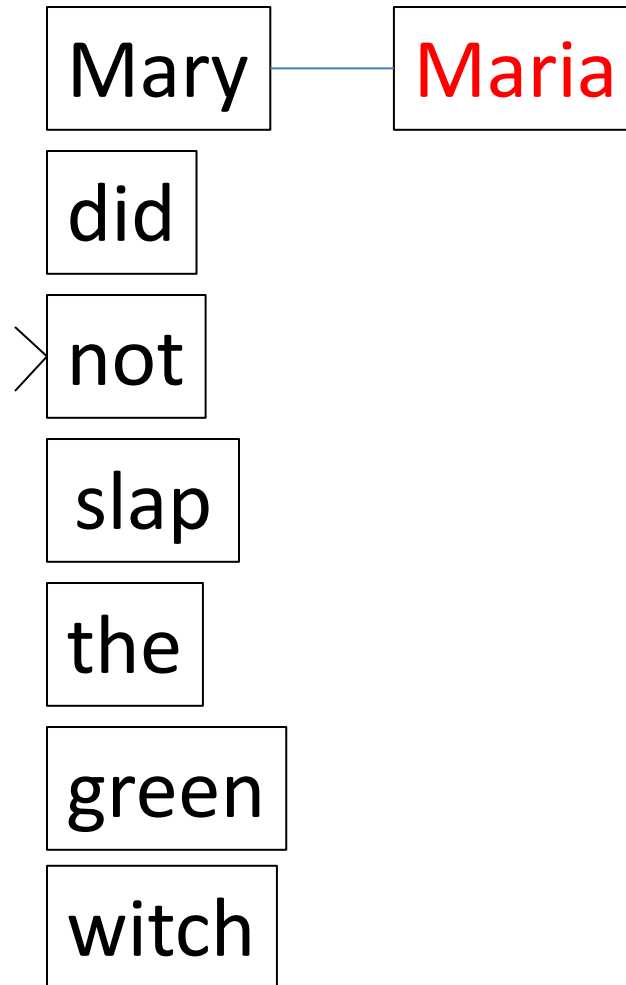
String Transduction for MT

➤ Mary
did
not
slap
the
green
witch

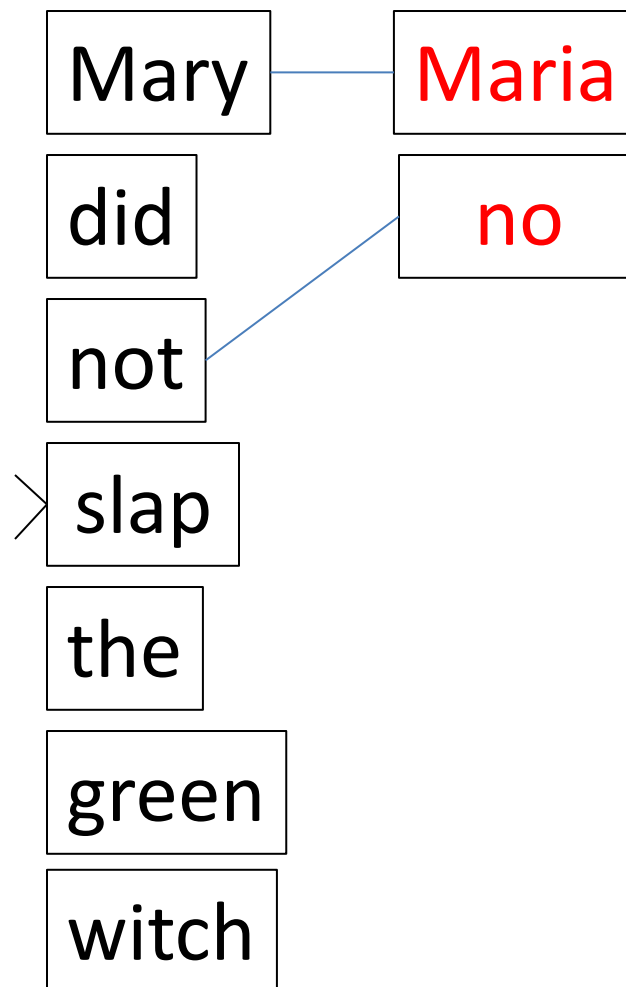
String Transduction for MT



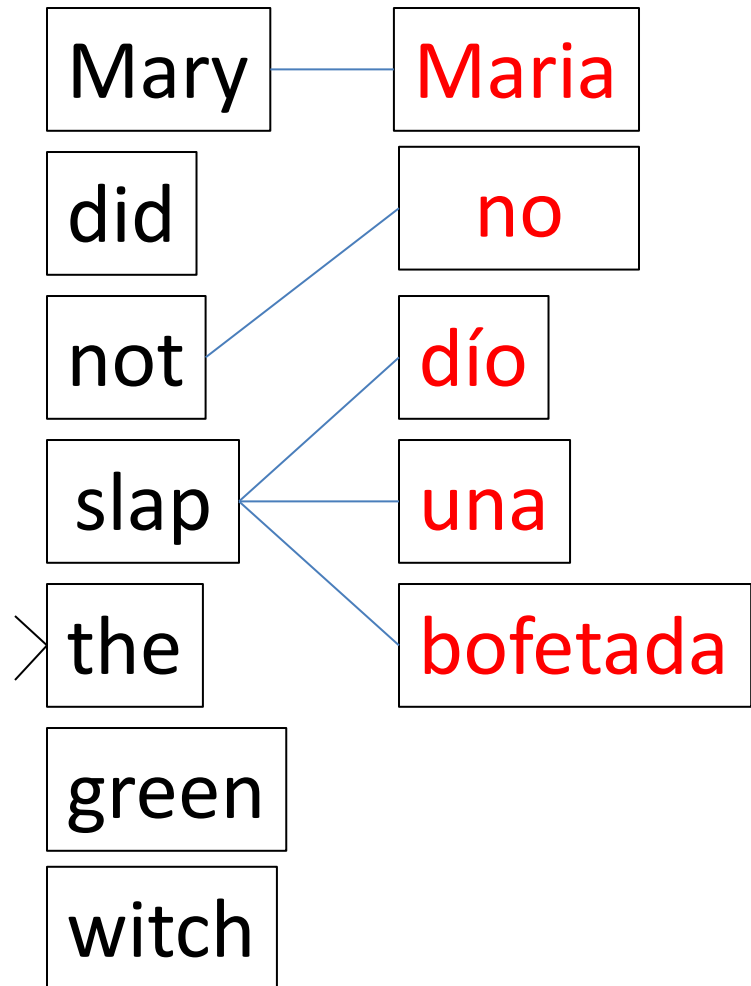
String Transduction for MT



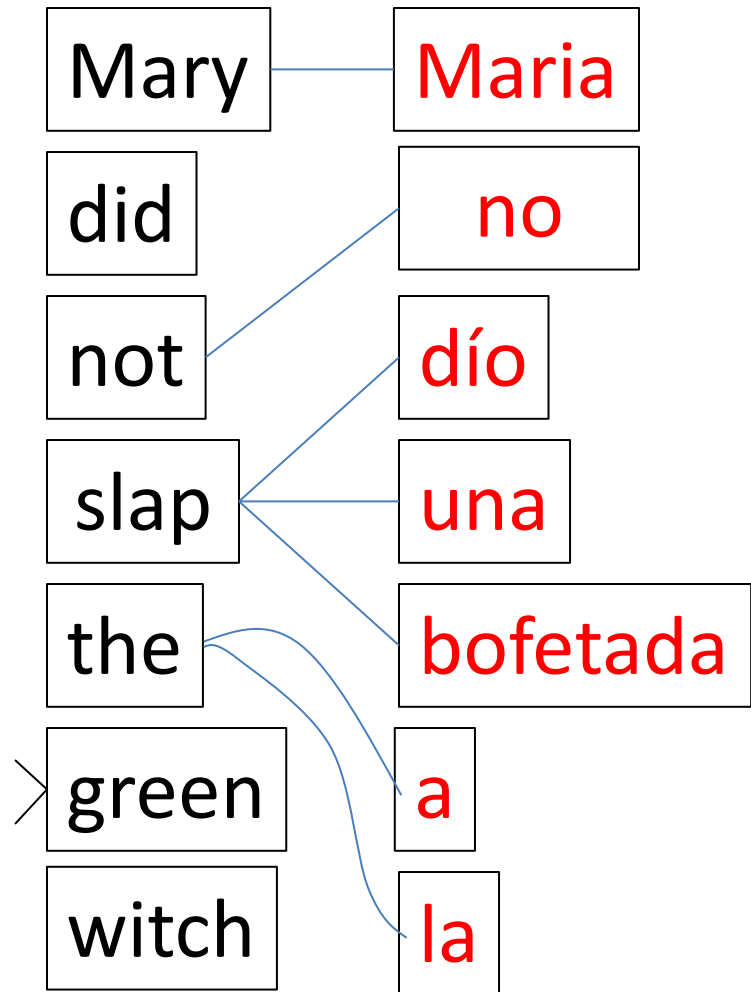
String Transduction for MT



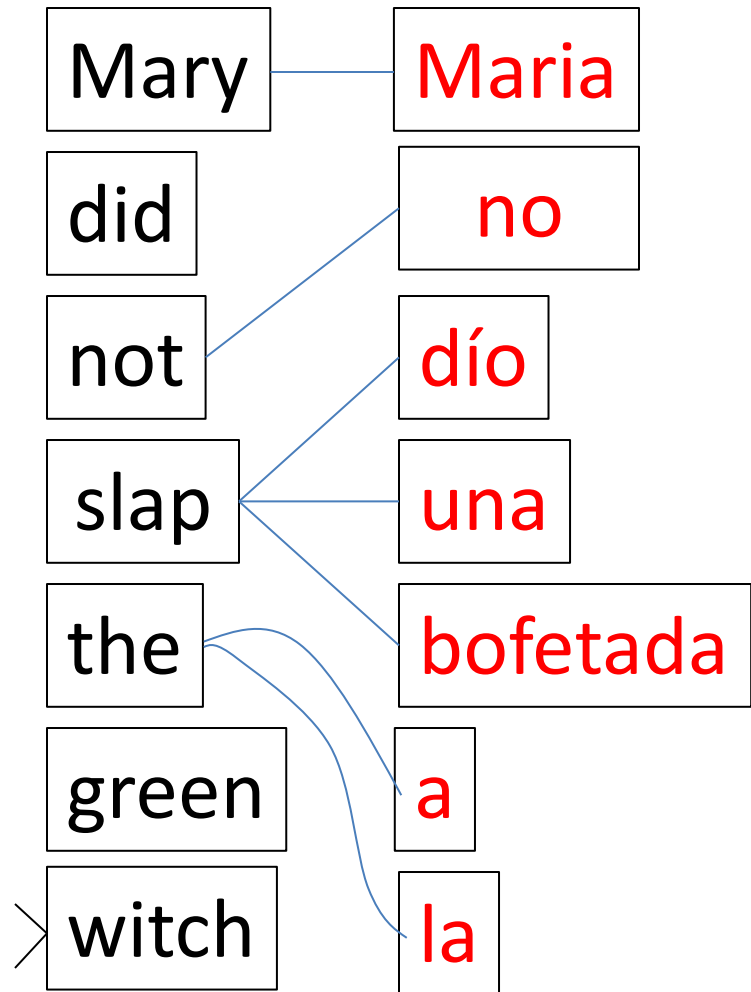
String Transduction for MT



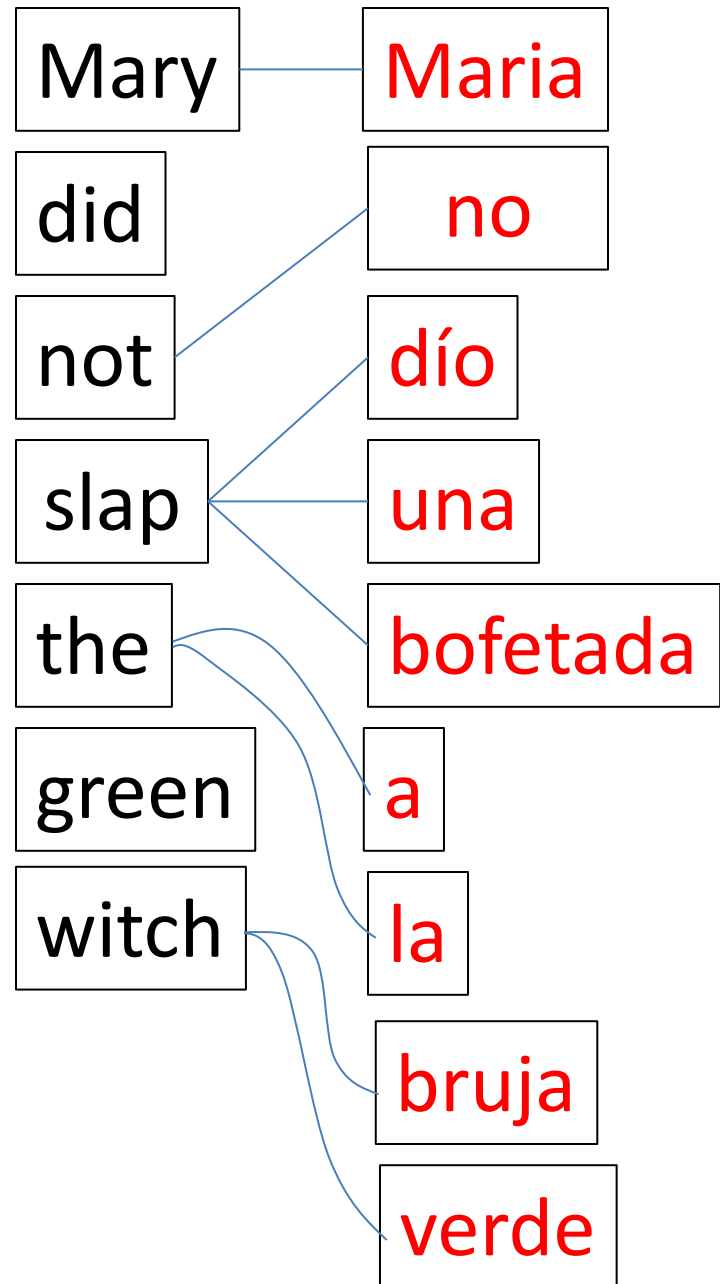
String Transduction for MT



String Transduction for MT



String Transduction for MT



String Transduction for MT

> He

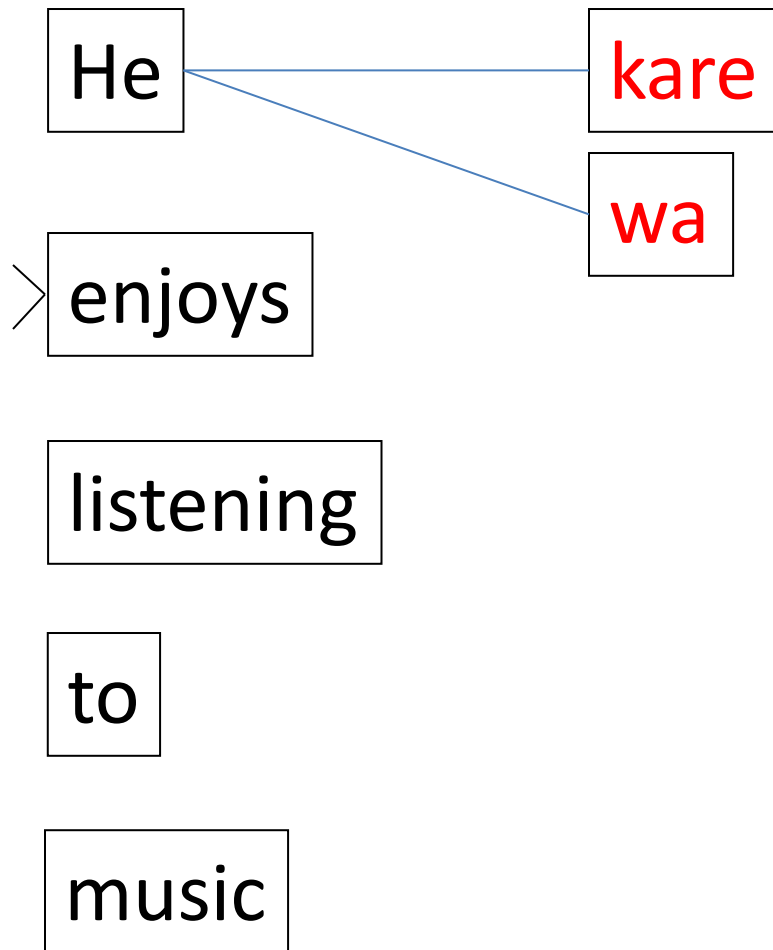
enjoys

listening

to

music

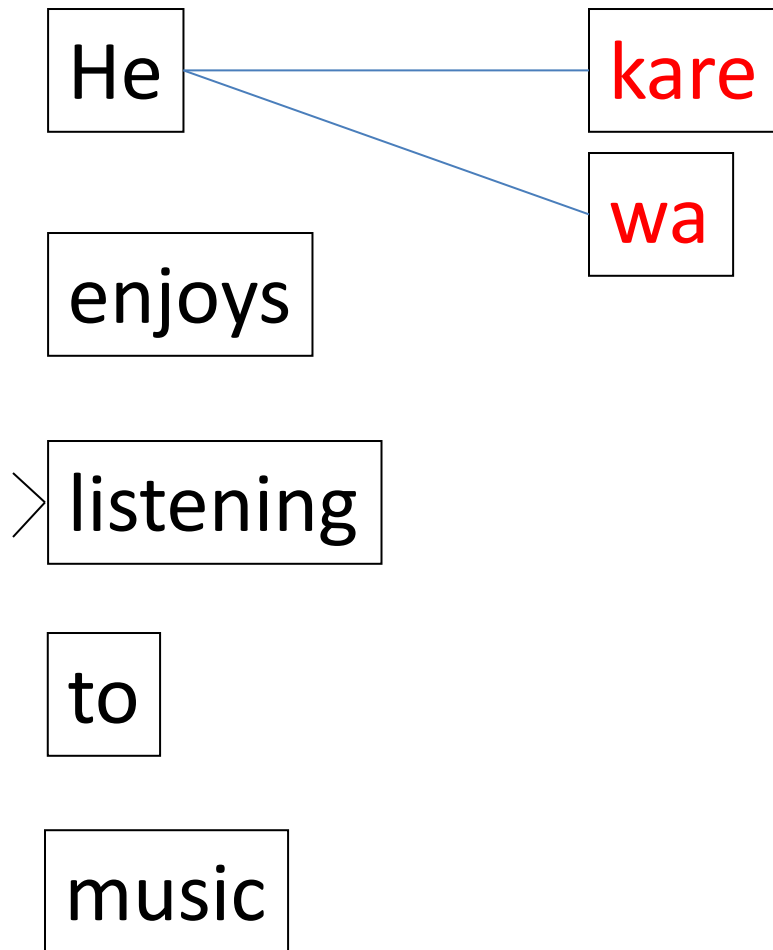
String Transduction for MT



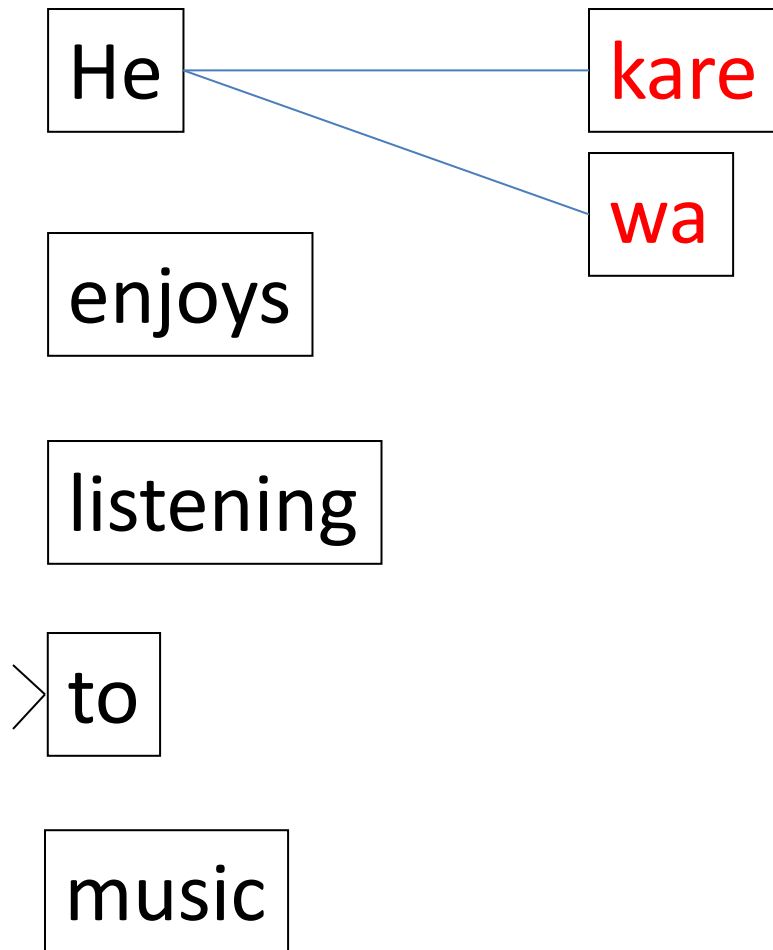
String Transduction for MT



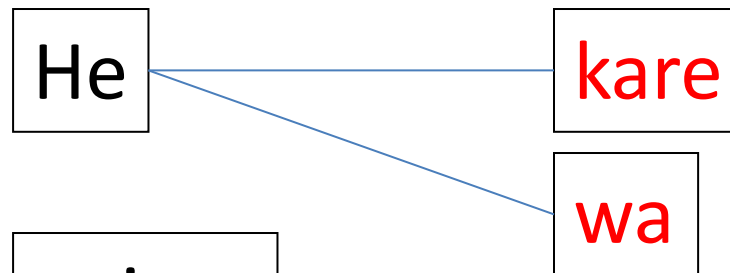
String Transduction for MT



String Transduction for MT



String Transduction for MT



enjoys

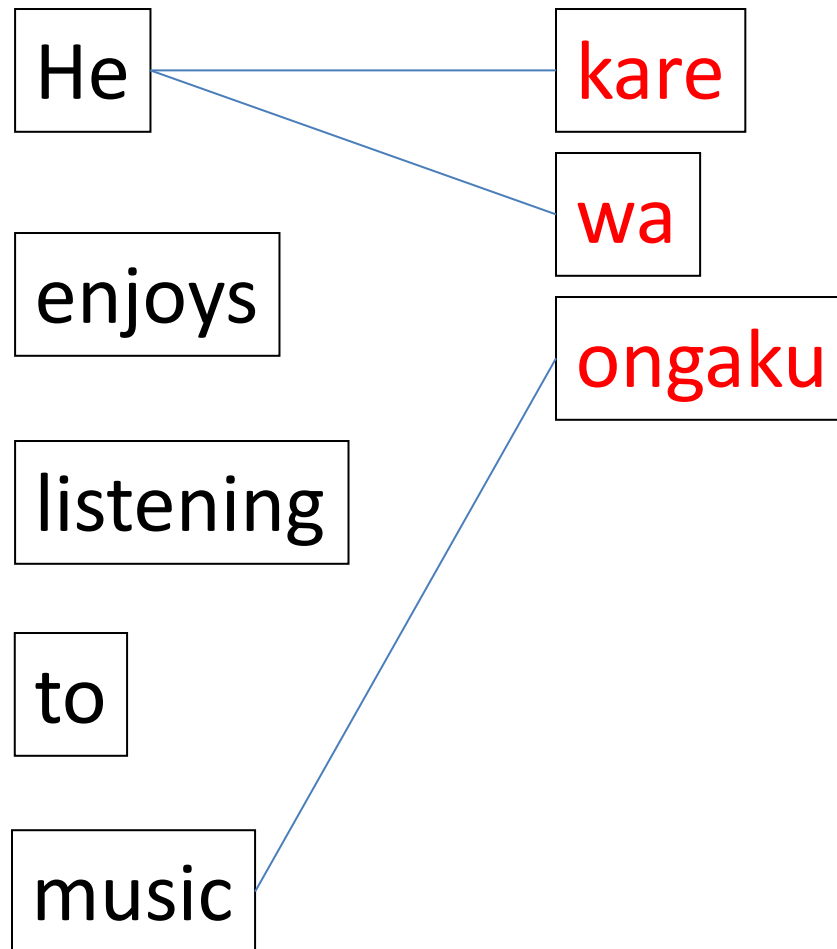
listening

to

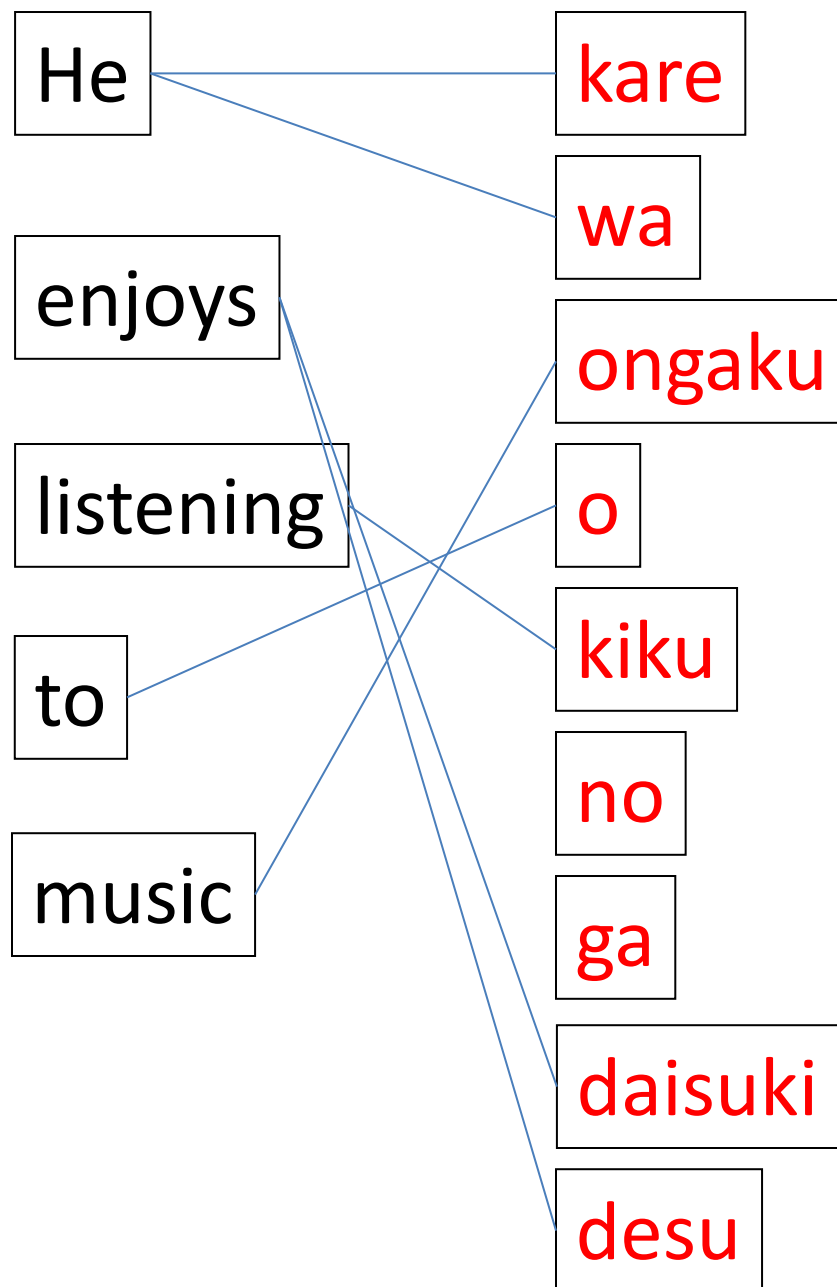
q7382736

> music

String Transduction for MT



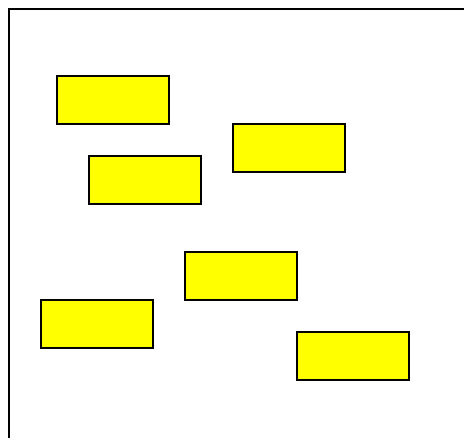
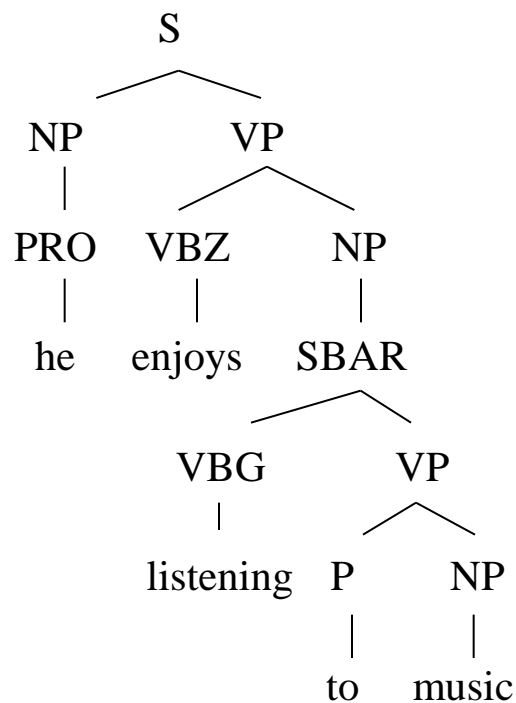
String Transduction for MT



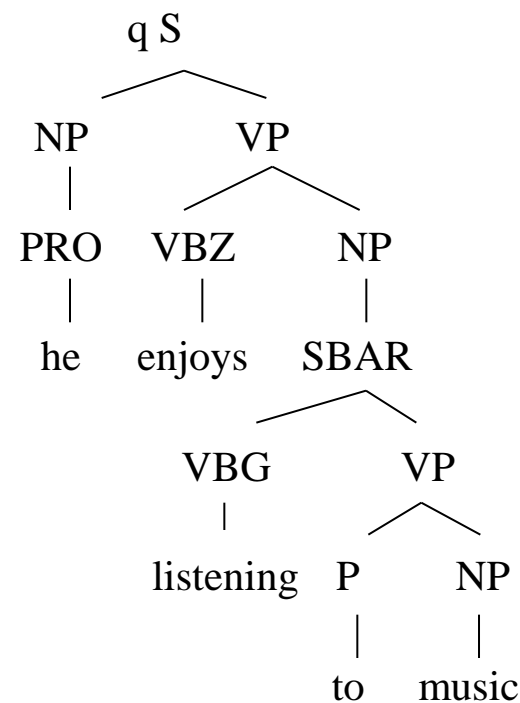
Top-Down Tree Transducer

(W. Rounds 1970; J. Thatcher 1970)

Original input:



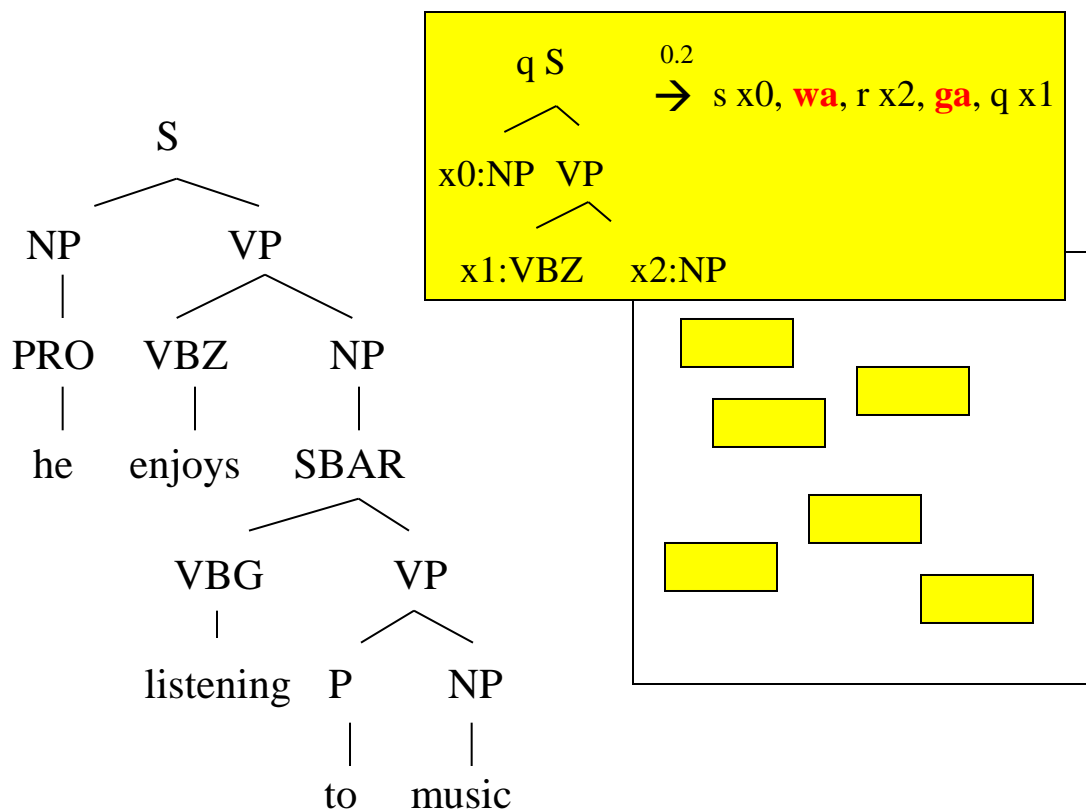
Transformation:



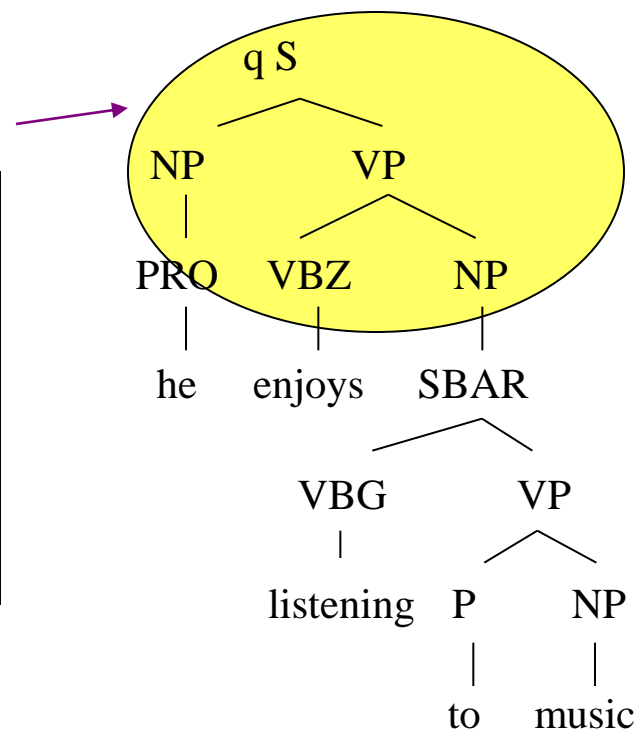
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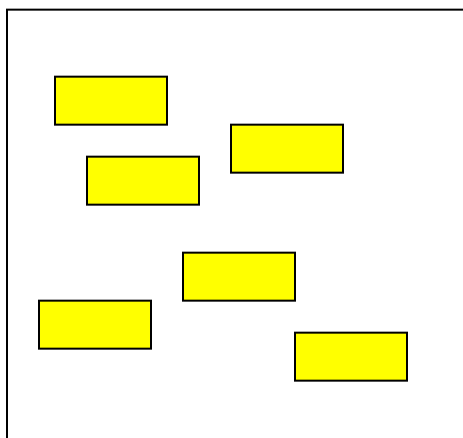
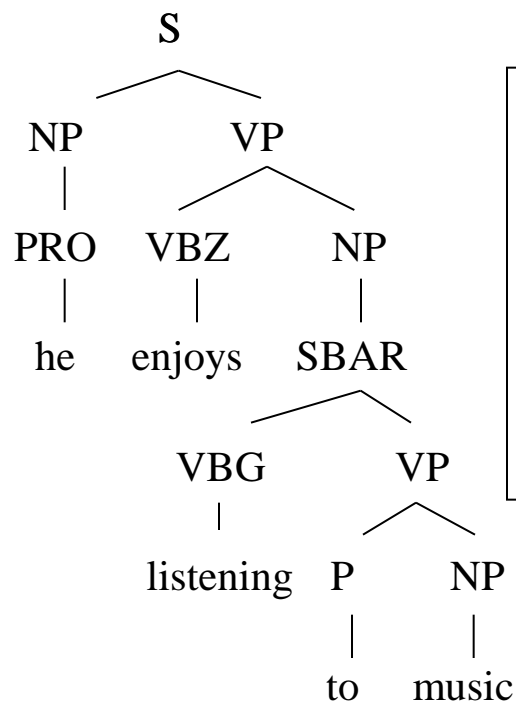
Transformation:



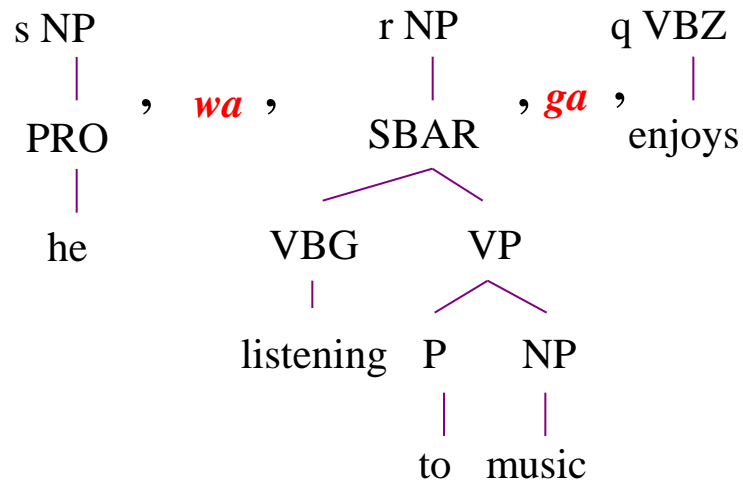
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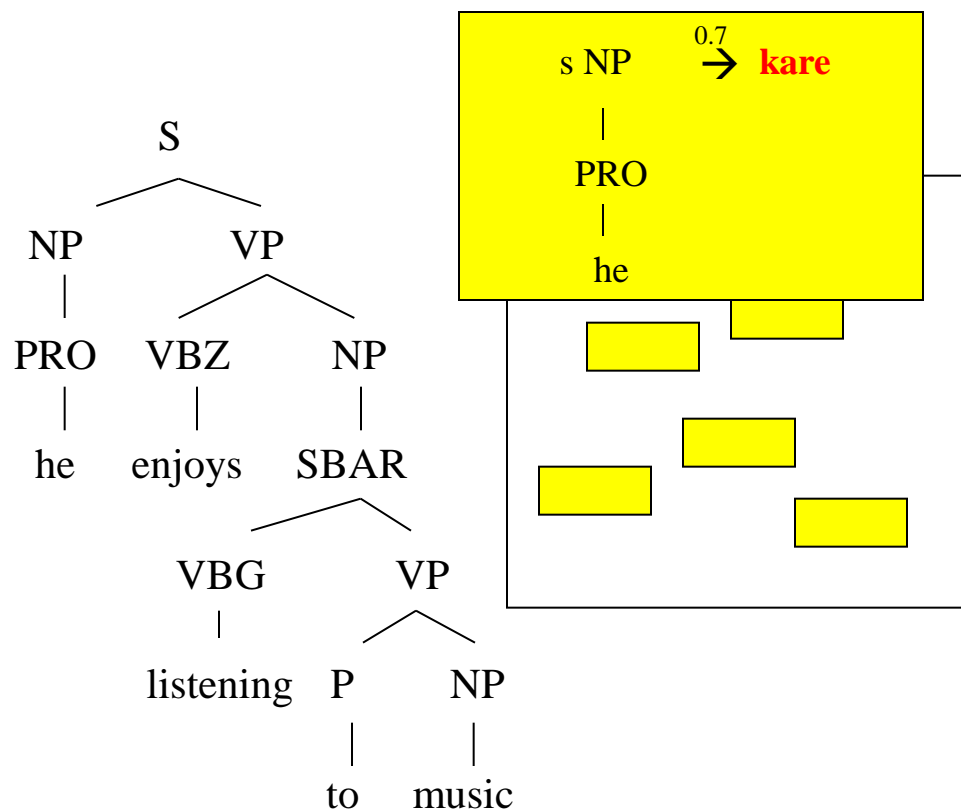
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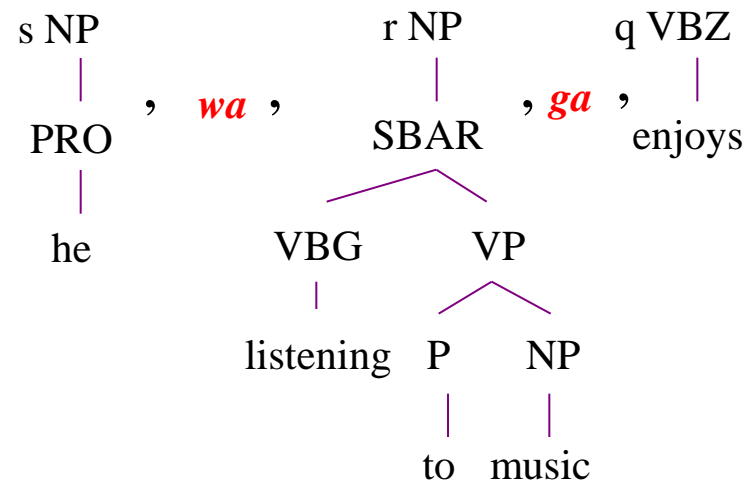
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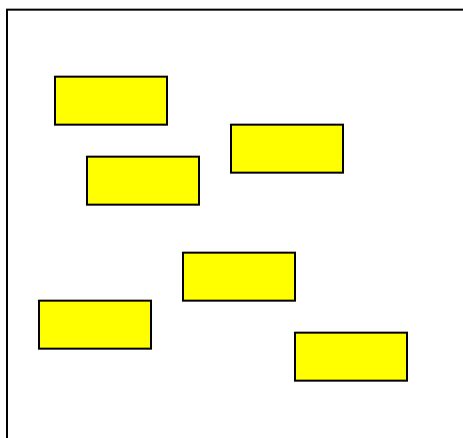
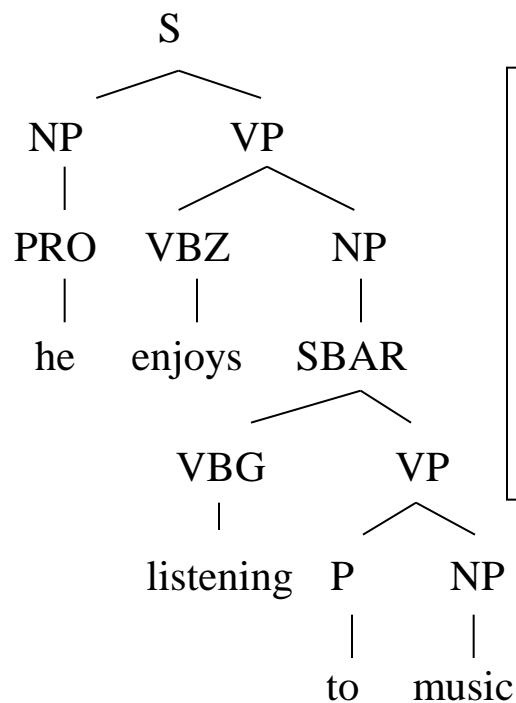
Transformation:



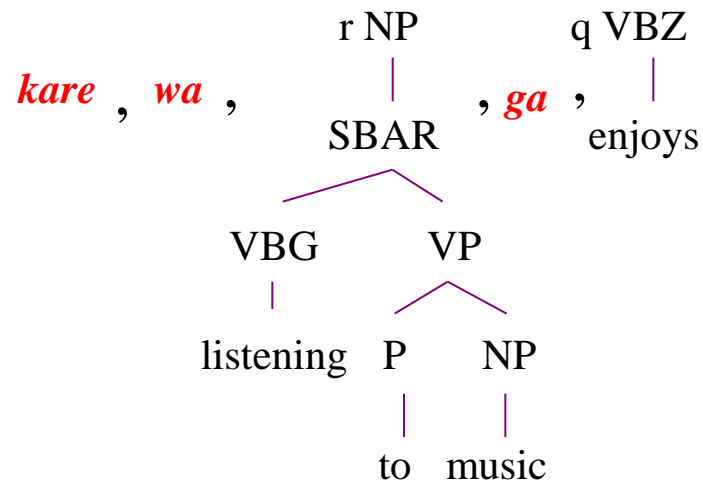
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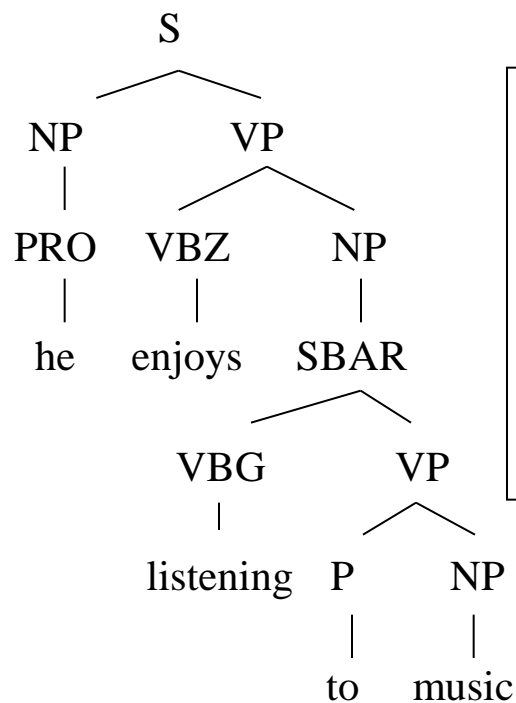
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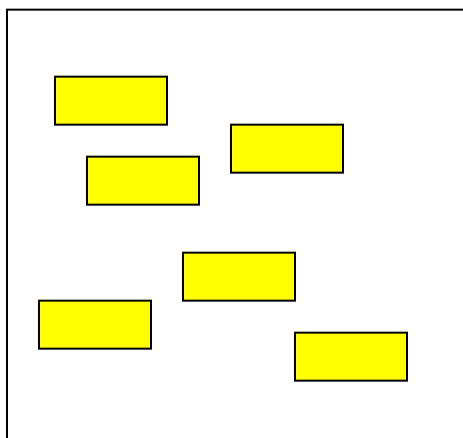
Top-Down Tree Transducer

(W. Rounds 1970; J. Thatcher 1970)

Original input:



Final output:



kare , wa , ongaku , o , kiku , no , ga , daisuki , desu

Top-Down Tree Transducer

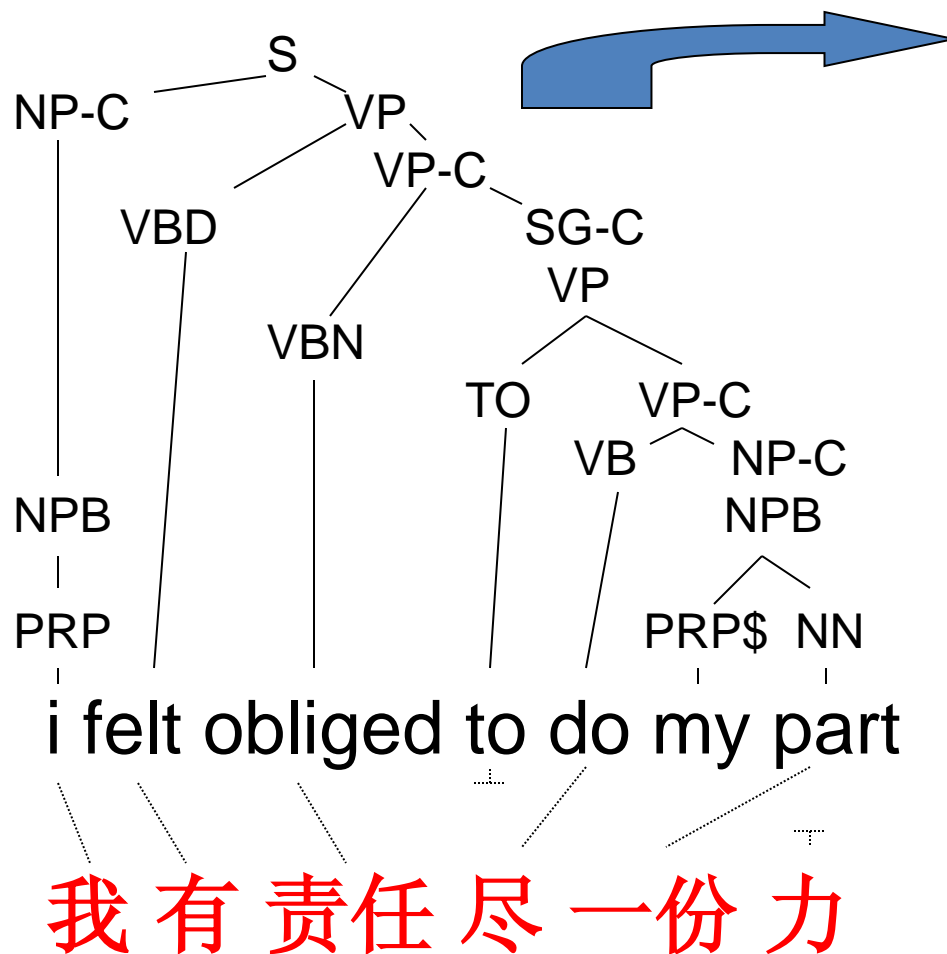
- Introduced by Rounds (1970) & Thatcher (1970)

“Recent developments in the theory of automata have pointed to an extension of the domain of definition of automata from strings to trees ... parts of **mathematical linguistics can be formalized easily** in a tree-automaton setting ...”

(Rounds 1970, “Mappings on Grammars and Trees”, *Math. Systems Theory* 4(3))

- Large theory literature
 - e.g., Gécseg & Steinby (1984), Comon et al (1997)
- Re-connecting with NLP practice
 - e.g., Knight & Graehl (2005), Galley et al (2004, 2006), May & Knight (2006, 2010), Maletti et al (2009)

Tree Transducers Can be Extracted from Bilingual Data (Galley et al, 04)



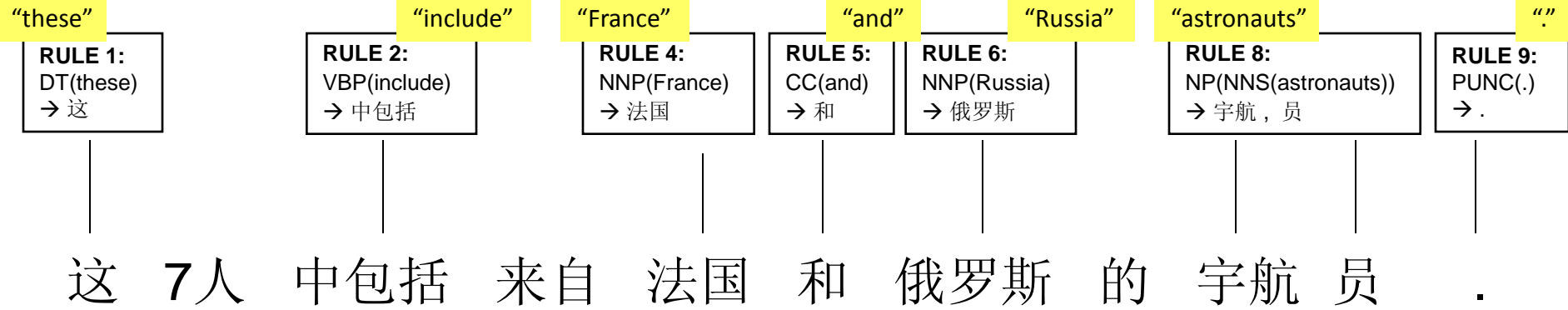
TREE TRANSDUCER RULES:

VBD(felt) → 有
VBN(obliged) → 责任
VB(do) → 尽
NN(part) → 一份
NN(part) → 一份 力
VP-C(x0:VBN x1:SG-C) → x0 x1
VP(TO(to) x0:VP-C) → x0
...
S(x0:NP-C x1:VP) → x0 x1

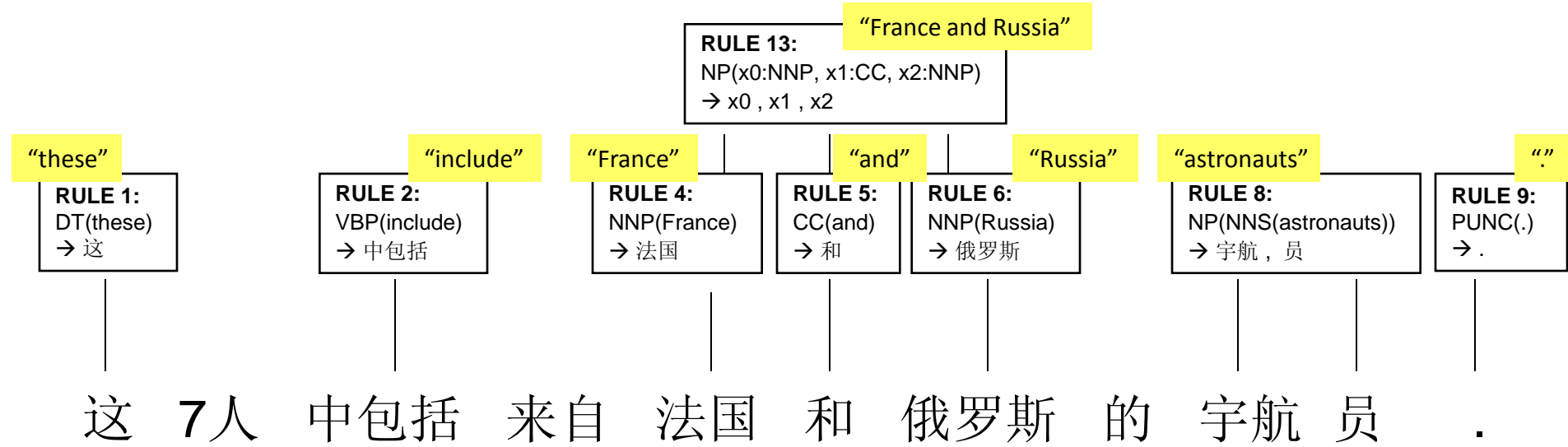
Syntax-Based Decoding

这 7人 中包括 来自 法国 和 俄罗斯 的 宇航 员 .

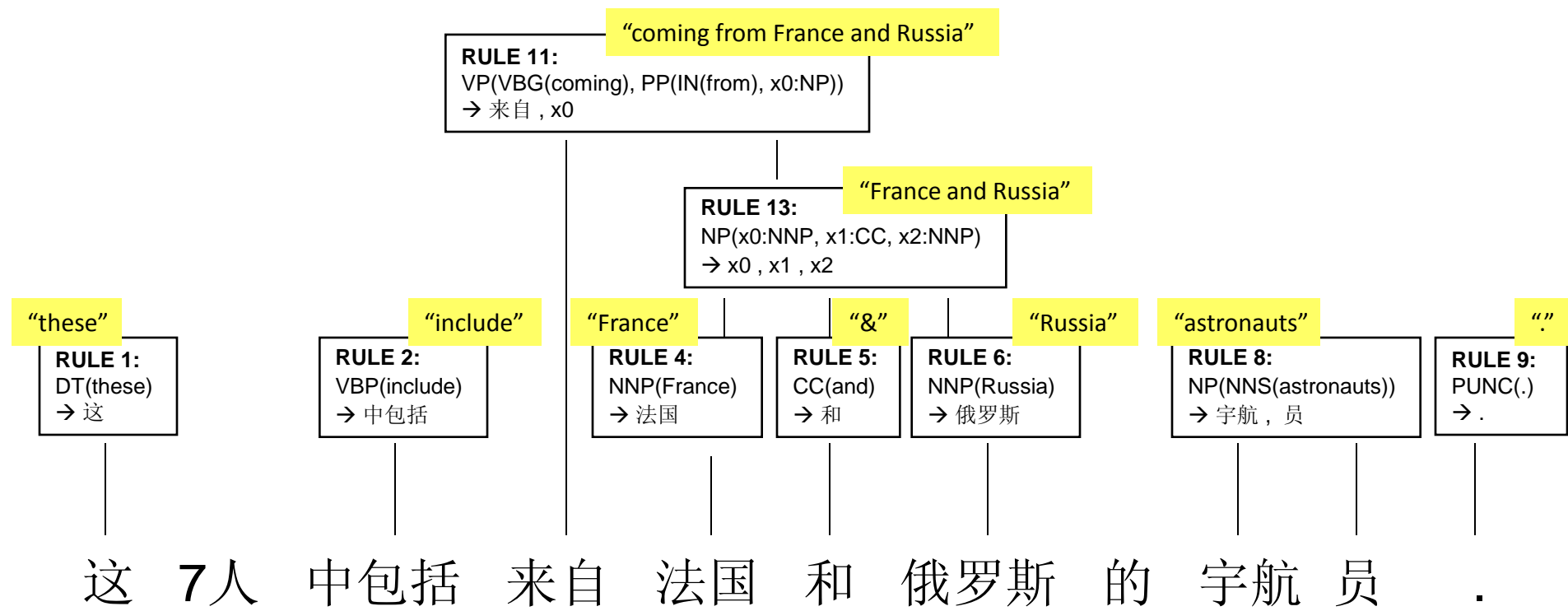
Syntax-Based Decoding



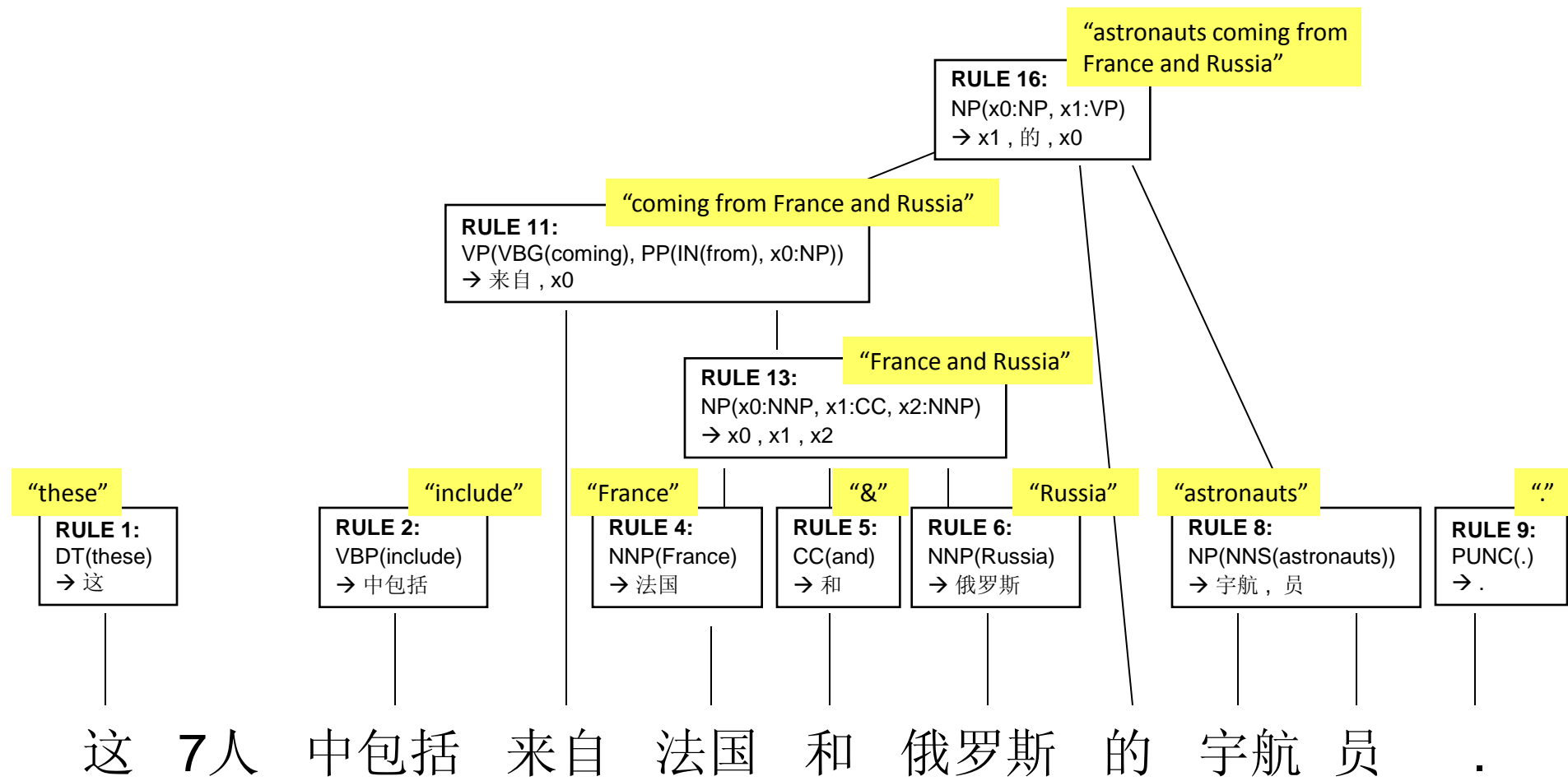
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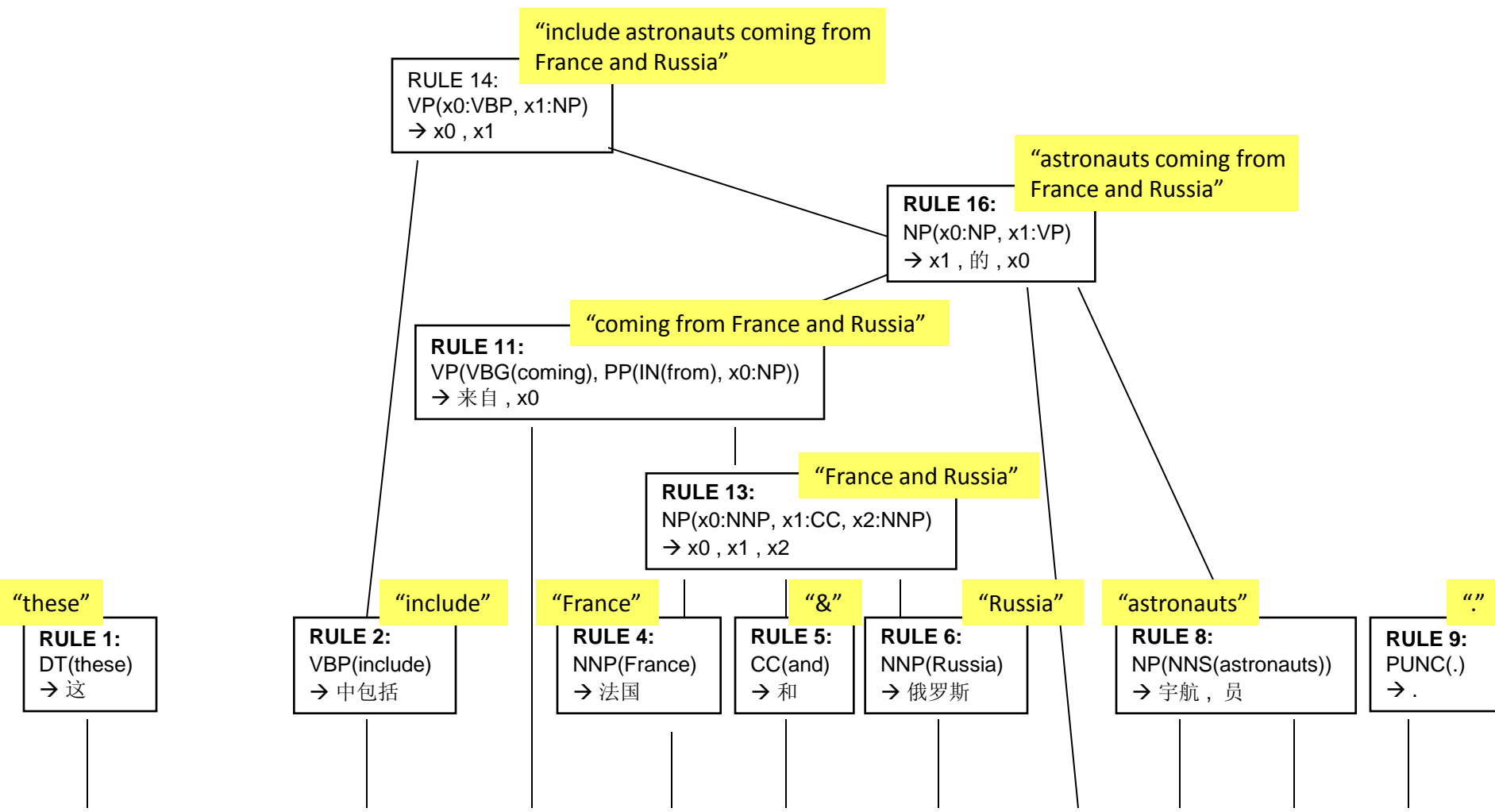


Syntax-Based Decoding



Syntax-Based Decoding





这 7人中包括来自法国和俄罗斯的宇航员 .

RULE 15:
S(x0:NP, x1:VP, x2:PUNC)
→ x0 , x1 , x2

"These 7 people include astronauts coming from France and Russia"

RULE 14:
VP(x0:VBP, x1:NP)
→ x0 , x1

"include astronauts coming from France and Russia"

"astronauts coming from France and Russia"

RULE 16:
NP(x0:NP, x1:VP)
→ x1 , 的 , x0

"coming from France and Russia"

RULE 11:
VP(VBG(coming), PP(IN(from), x0:NP))
→ 来自 , x0

"France and Russia"

RULE 13:
NP(x0:NNP, x1:CC, x2:NNP)
→ x0 , x1 , x2

RULE 10:
NP(x0:DT, CD(7), NNS(people))
→ x0 , 7人

"these 7 people"

"these"

RULE 1:
DT(these)
→ 这

"include"

RULE 2:
VBP(include)
→ 中包括

"France"

RULE 4:
NNP(France)
→ 法国

"&"

RULE 5:
CC(and)
→ 和

"Russia"

RULE 6:
NNP(Russia)
→ 俄罗斯

"astronauts"

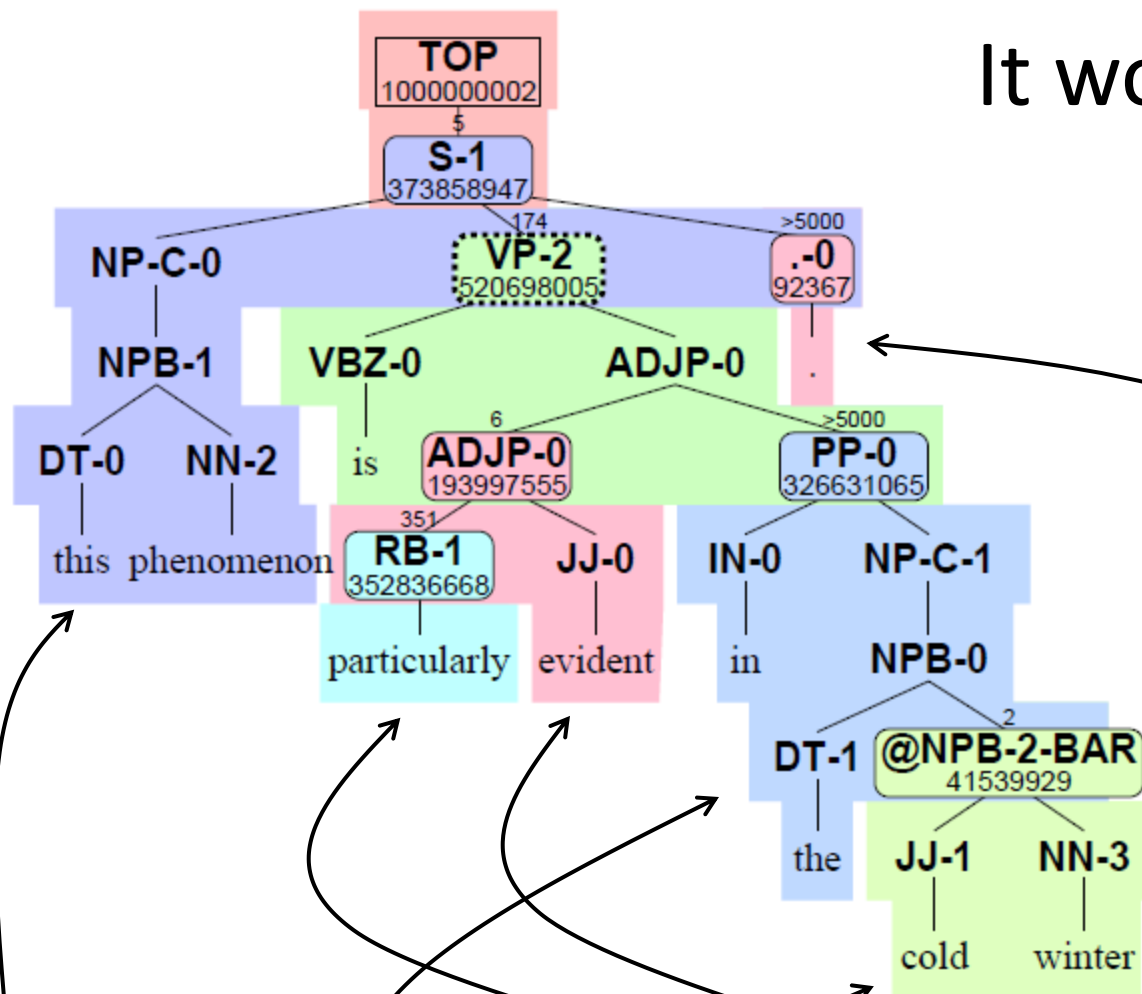
RULE 8:
NP(NNS(astronauts))
→ 宇航 , 员

"."

RULE 9:
PUNC(.)
→ .

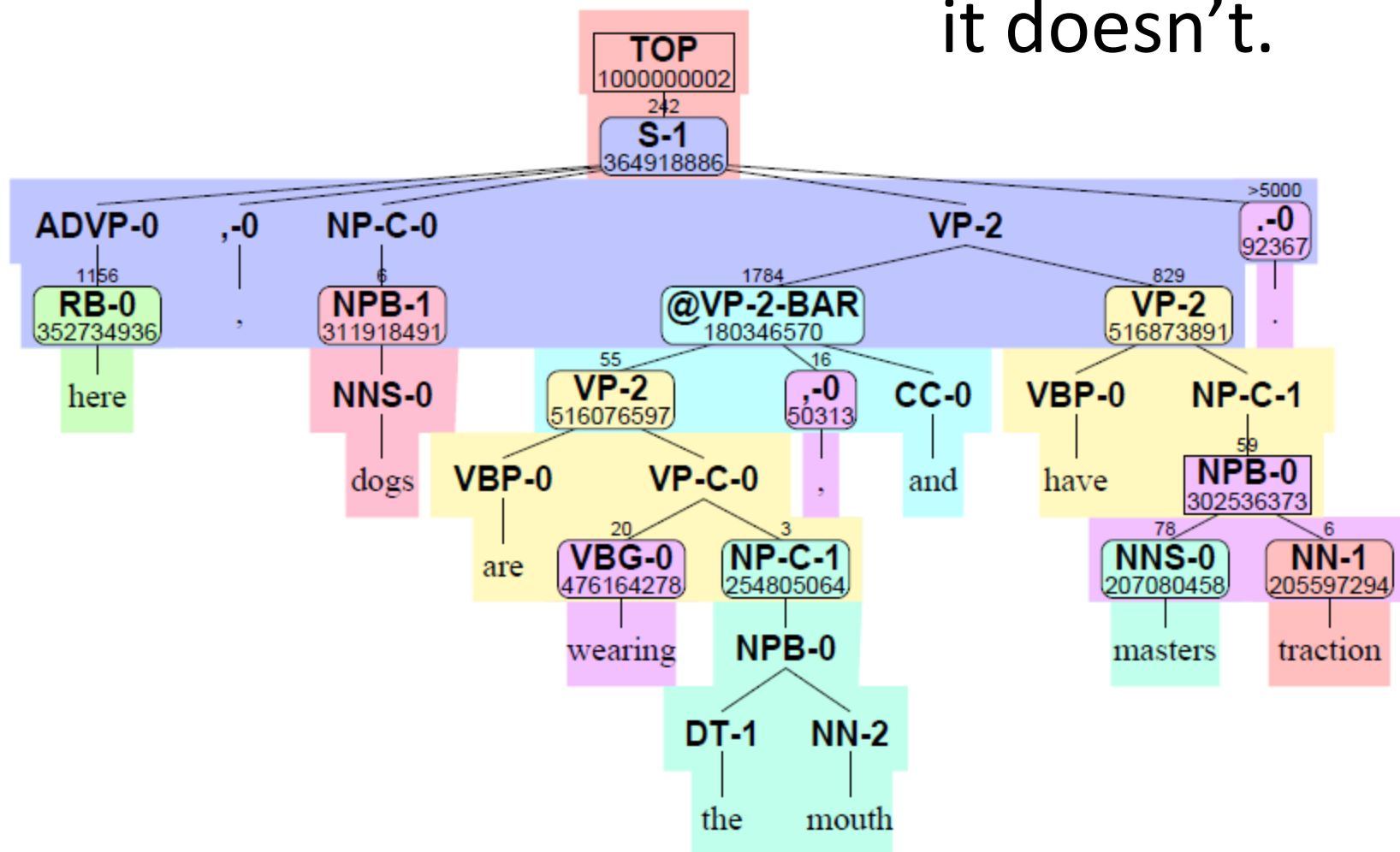
这 7人 中包括 来自 法国 和 俄罗斯 的 宇航 员 .

It works...



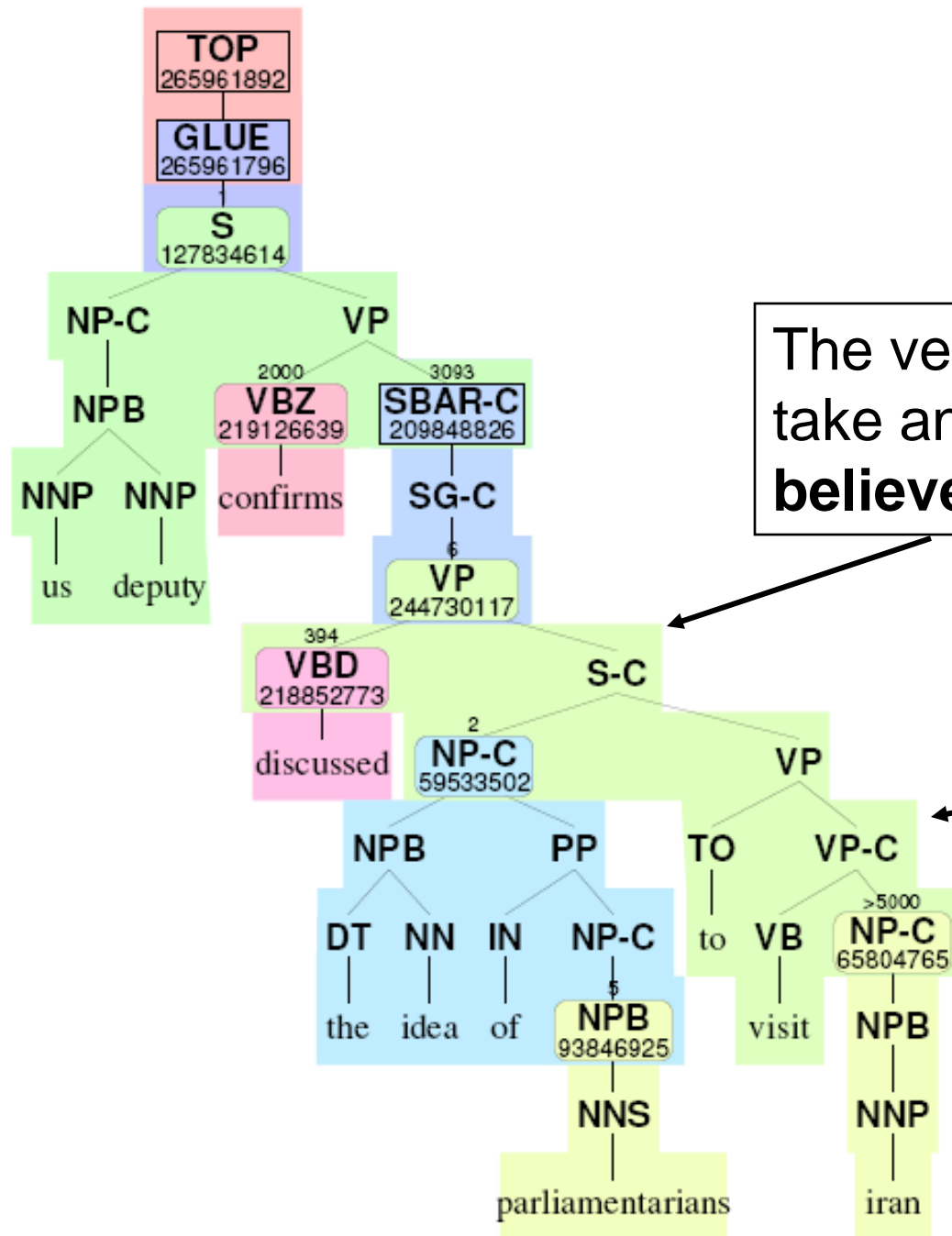
这种现象在寒冷的冬季尤其明显。

...except when
it doesn't.



在这里，狗都配戴嘴套，并有主人牵引。

...except when
it doesn't.



The verb **discuss** doesn't
take an S argument (like
believe and **realize** do)

An **idea** can't
visit a **place**

General-Purpose Algorithms for Tree Automata

	String Automata Algorithms	Tree Automata Algorithms
N-best paths through an WFSA (Viterbi, 1967; Eppstein, 1998)	... trees in a weighted forest (Jiménez & Marzal, 2000; Huang & Chiang, 2005)
EM training	Forward-backward EM (Baum/Welch, 1971; Eisner 2003)	Tree transducer EM training (Graehl & Knight, 2004)
Determinization of weighted string acceptors (Mohri, 1997)	... of weighted tree acceptors (Borchardt & Vogler, 2003; May & Knight, 2005)
Intersection	WFSA intersection	Tree acceptor intersection
Applying transducers	string \rightarrow WFST \rightarrow WFSA	tree \rightarrow TT \rightarrow weighted tree acceptor
Transducer composition	WFST composition (Pereira & Riley, 1996)	Many tree transducers not closed under composition (Maletti et al 09)
General-purpose tools	FSM, Carmel, OpenFST	Tiburon (May & Knight 10)

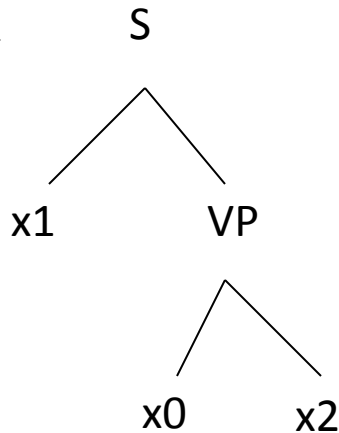
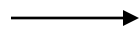
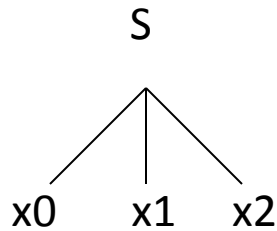
Tree Transducers

every rule has this form

one-level LHS

multilevel RHS

T – top-down
L – linear (non-copying)
N – non-deleting



arabic verb

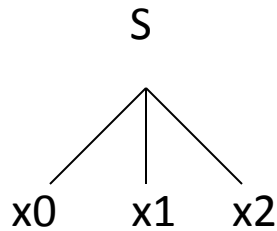
arabic subject

arabic object

LNT

Tree Transducers

one-level LHS

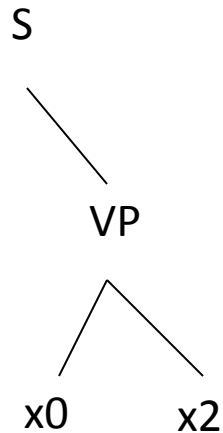


arabic verb

arabic subject

arabic object

multilevel RHS



T – top-down

L – linear (non-copying)

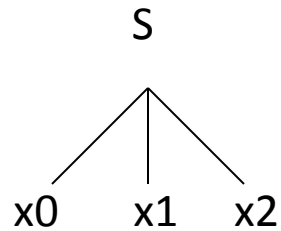
N – non-deleting

LT

can delete subtrees

Tree Transducers

one-level LHS

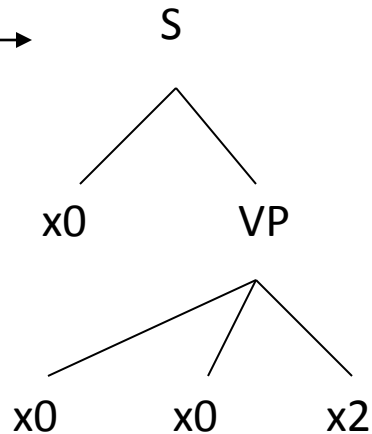


arabic verb

arabic subject

arabic object

multilevel RHS



T – top-down

L – linear (non-copying)

N – non-deleting

T

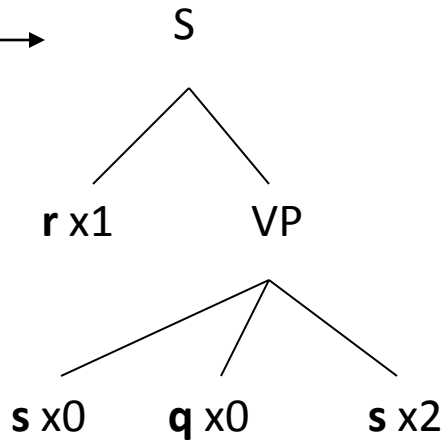
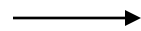
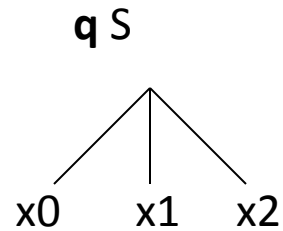
can copy & delete subtrees

Tree Transducers

one-level LHS

multilevel RHS

T – top-down
L – linear (non-copying)
N – non-deleting



arabic verb

arabic subject

arabic object

T

LT

LNT

all employ **states**

Choices

- There are many different tree transducers
- What is good for NLP?
- Investigate formal properties related to:
 - expressiveness
 - modularity
 - inclusiveness
 - learnability

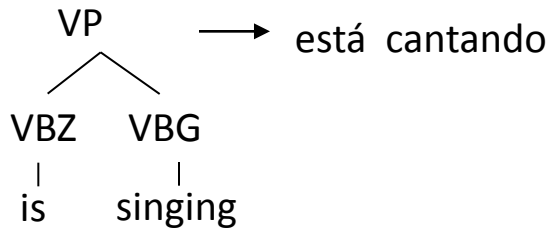
Choices

- There are many different tree transducers
- What is good for NLP?
- Investigate formal properties related to:
 - expressiveness (suggestions coming)
 - modularity (closed under composition)
 - inclusiveness (does everything FST does)
 - learnability (polynomial time EM)

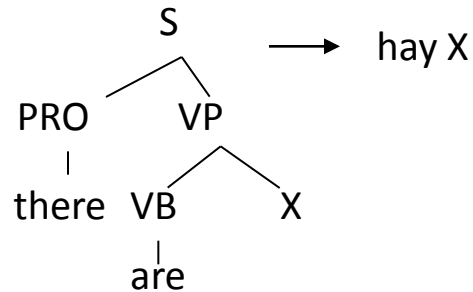
Expressiveness

some necessary things for machine translation

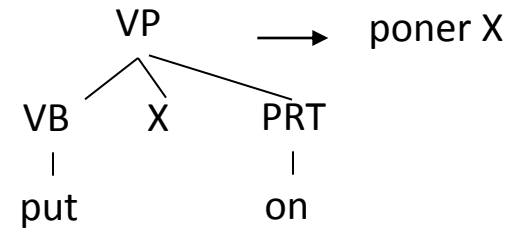
Phrasal Translation



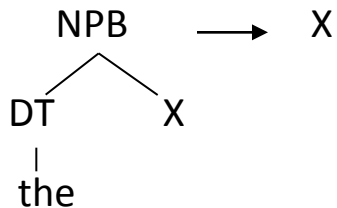
Non-constituent Phrases



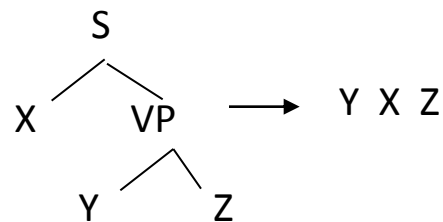
Non-contiguous Phrases



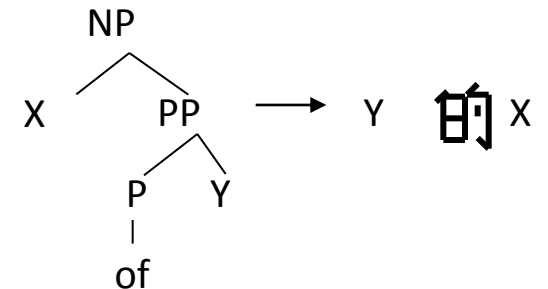
Context-Sensitive Word Insertion



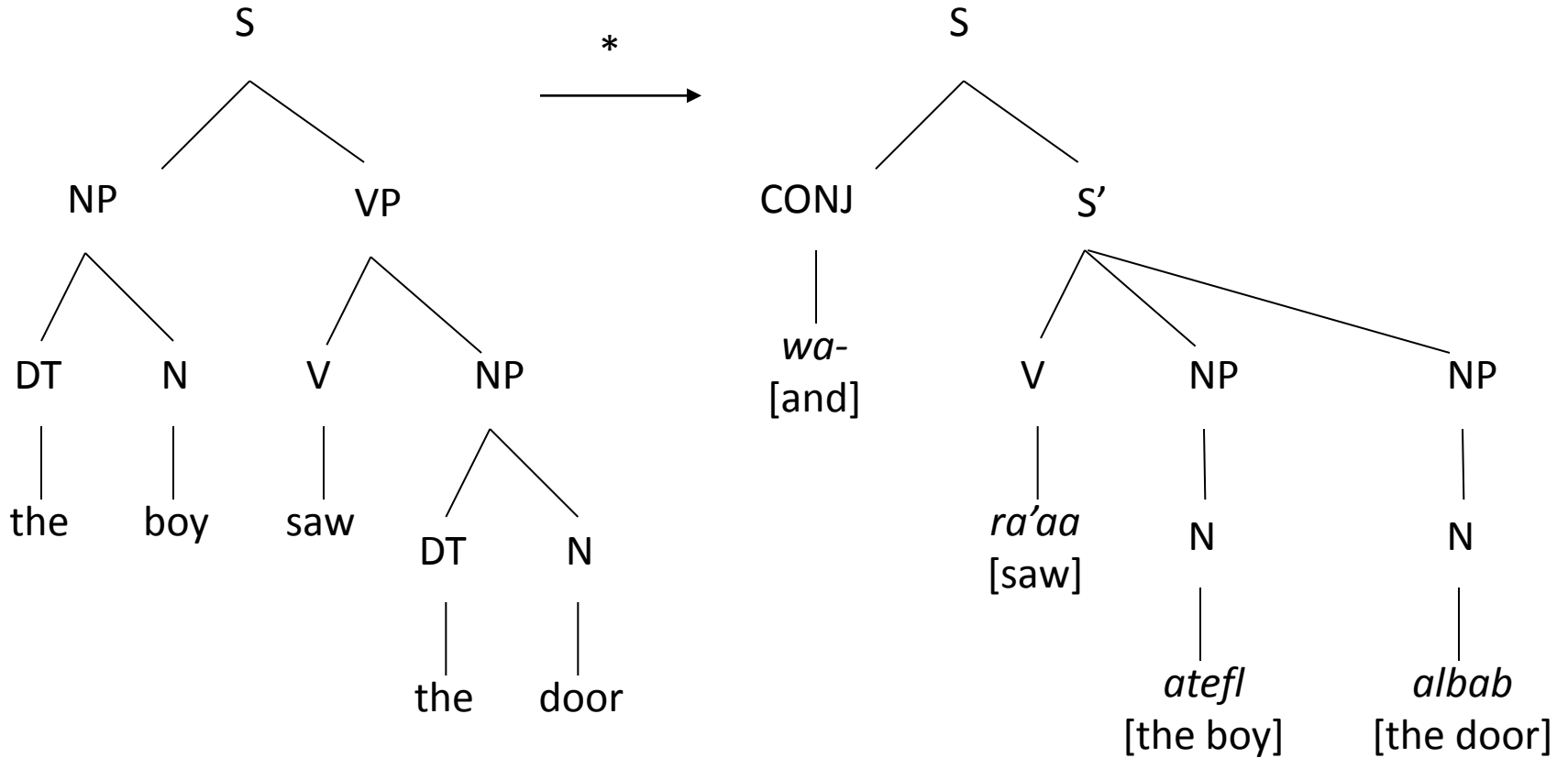
Re-Ordering



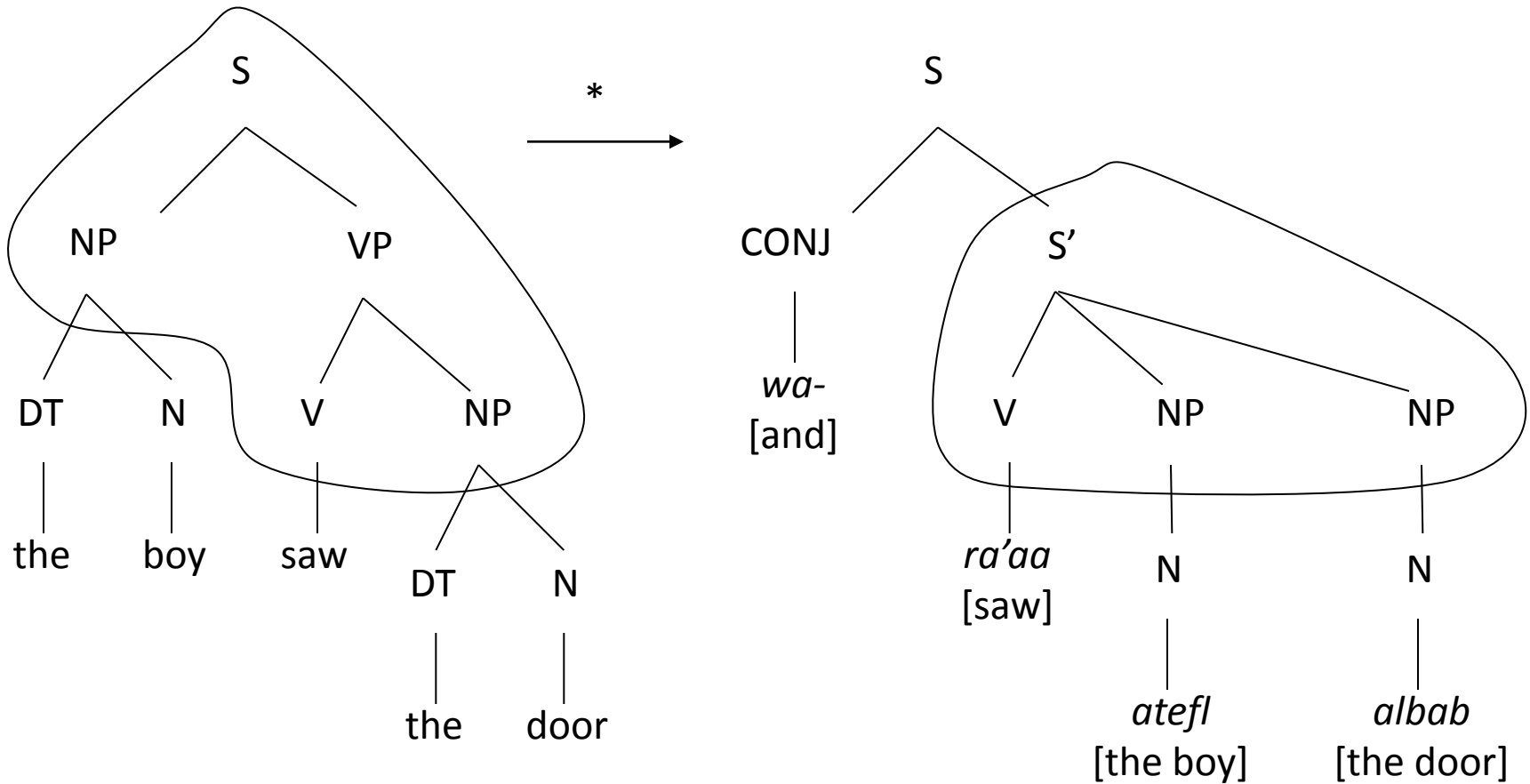
Lexicalized Re-Ordering



Expressiveness

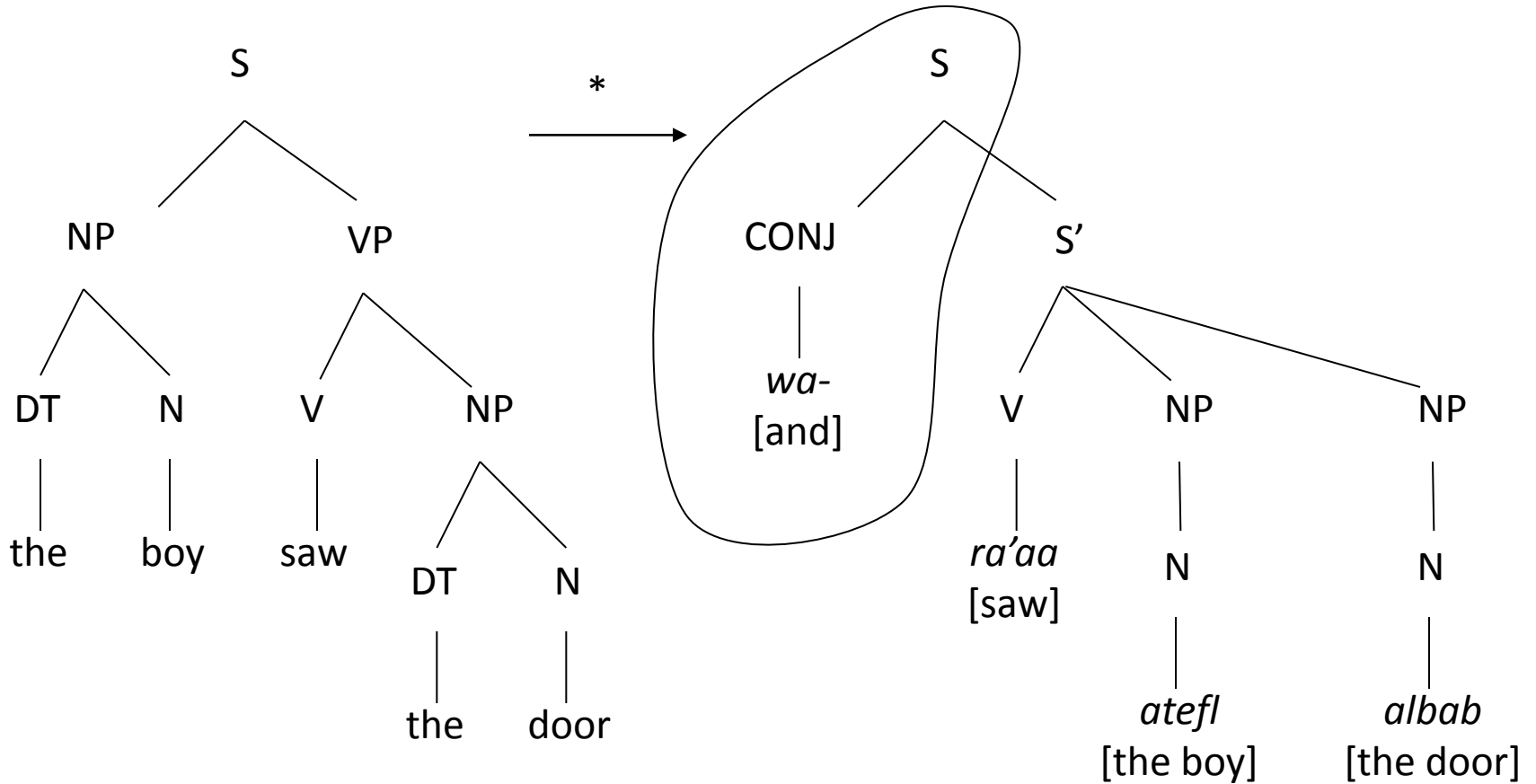


Expressiveness



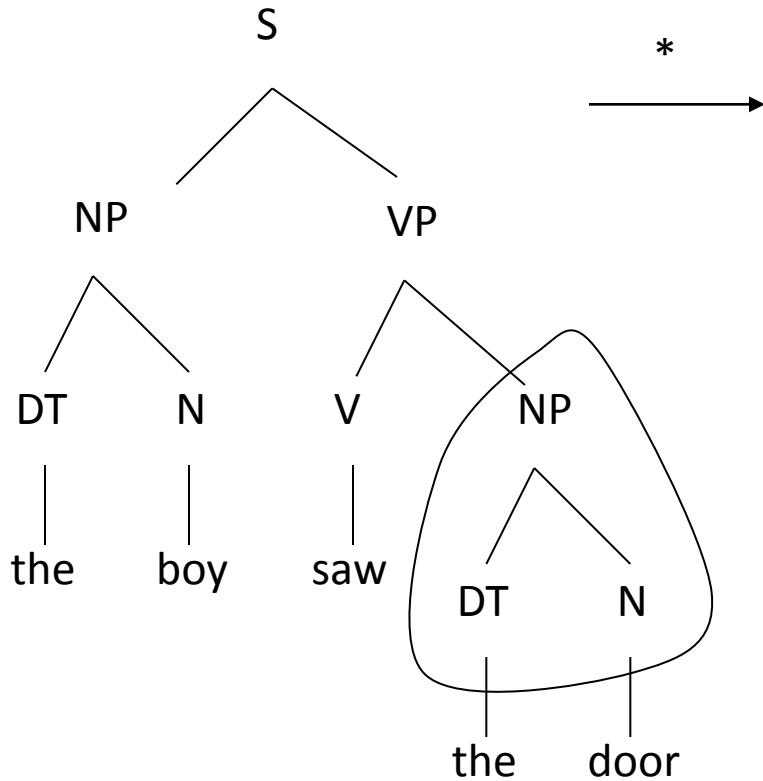
Local rotation

Expressiveness



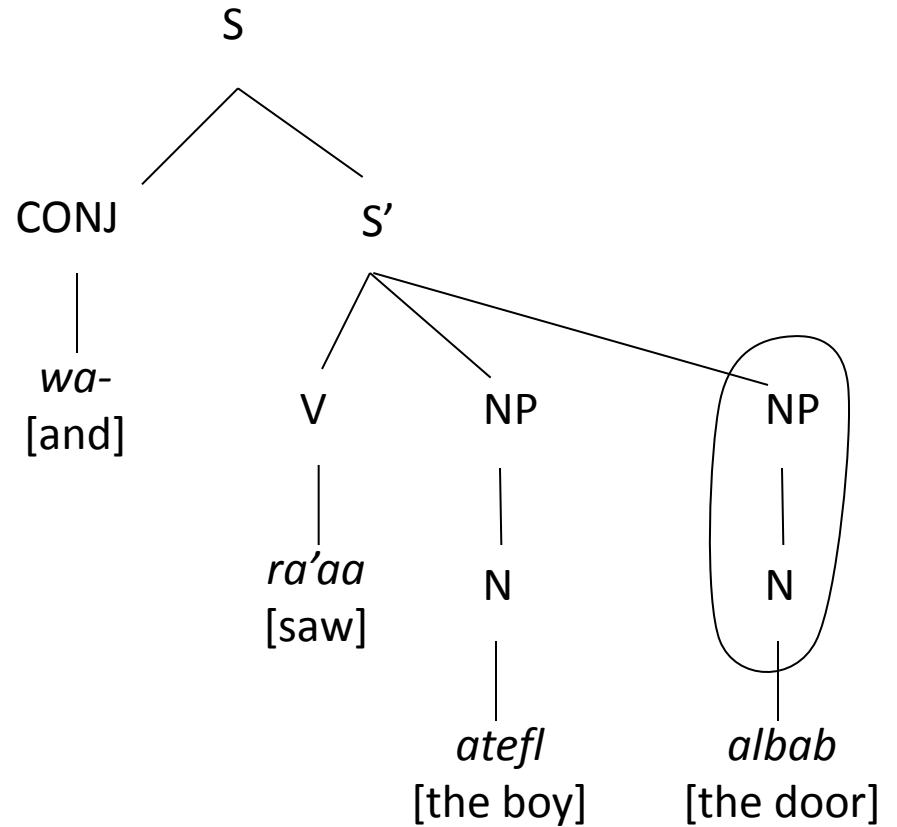
Produce output without
consuming input

Expressiveness



*

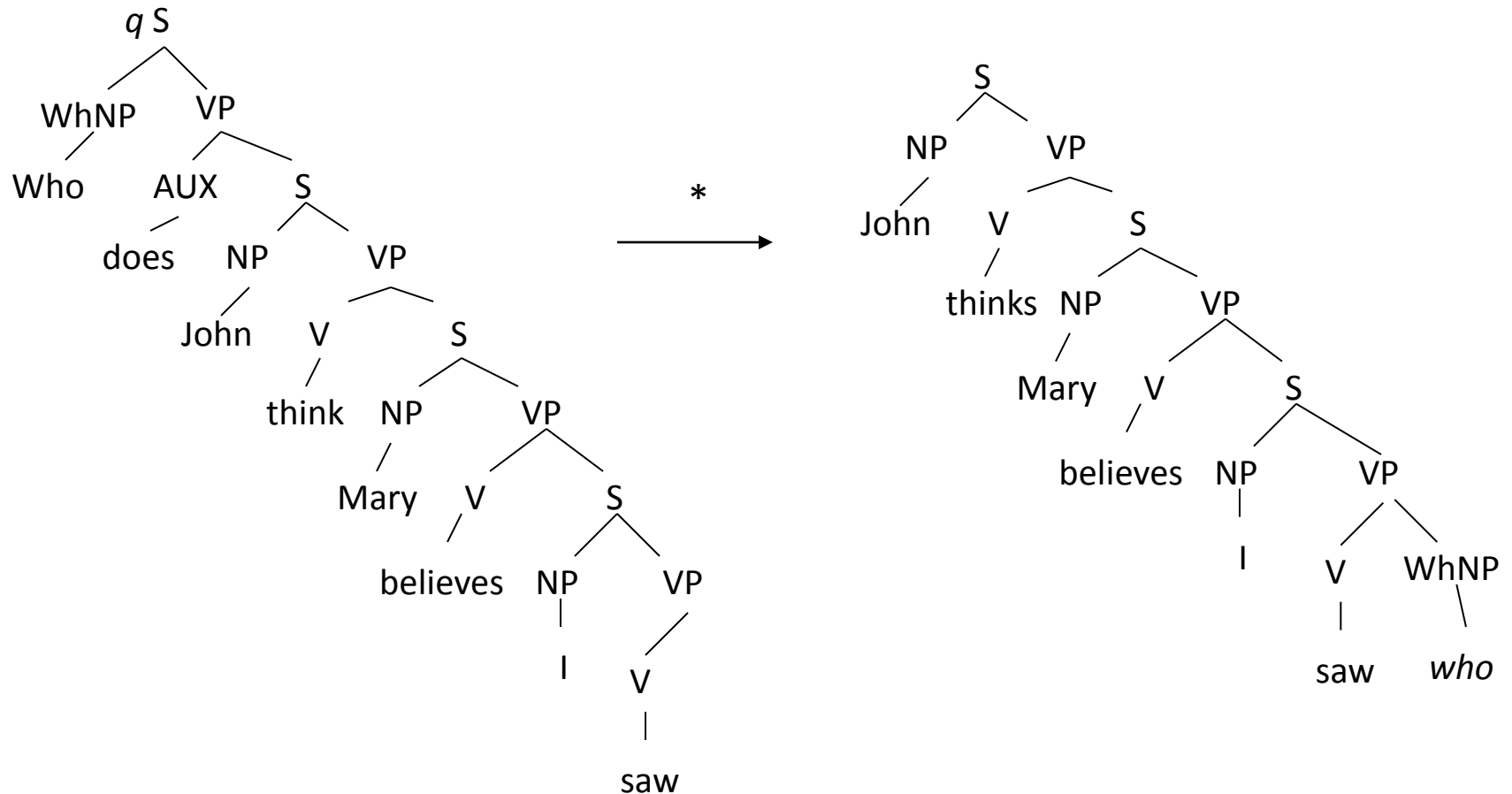
→



Lookahead

Expressiveness

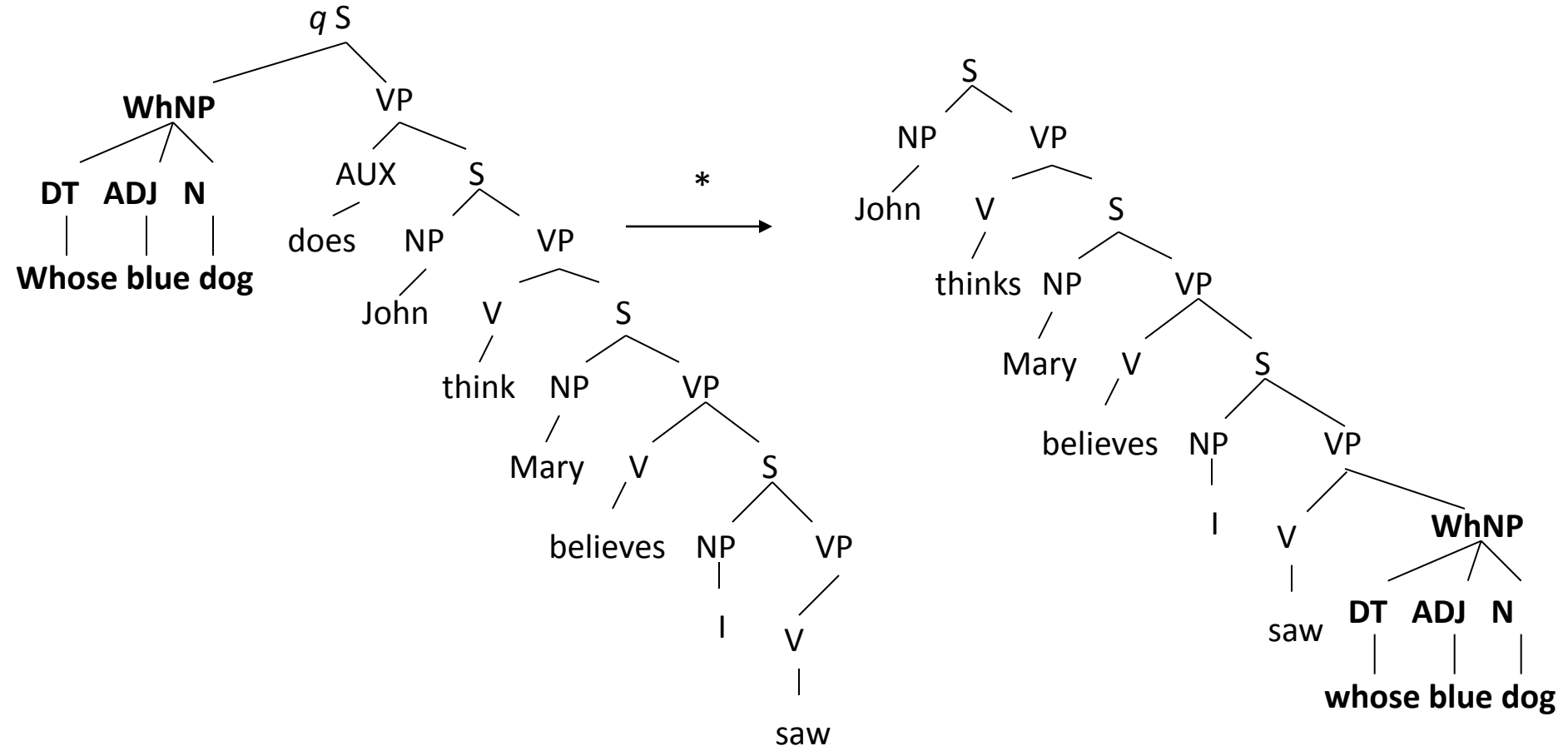
Who does John think Mary believes I saw? → John thinks Mary believes I saw **who**?



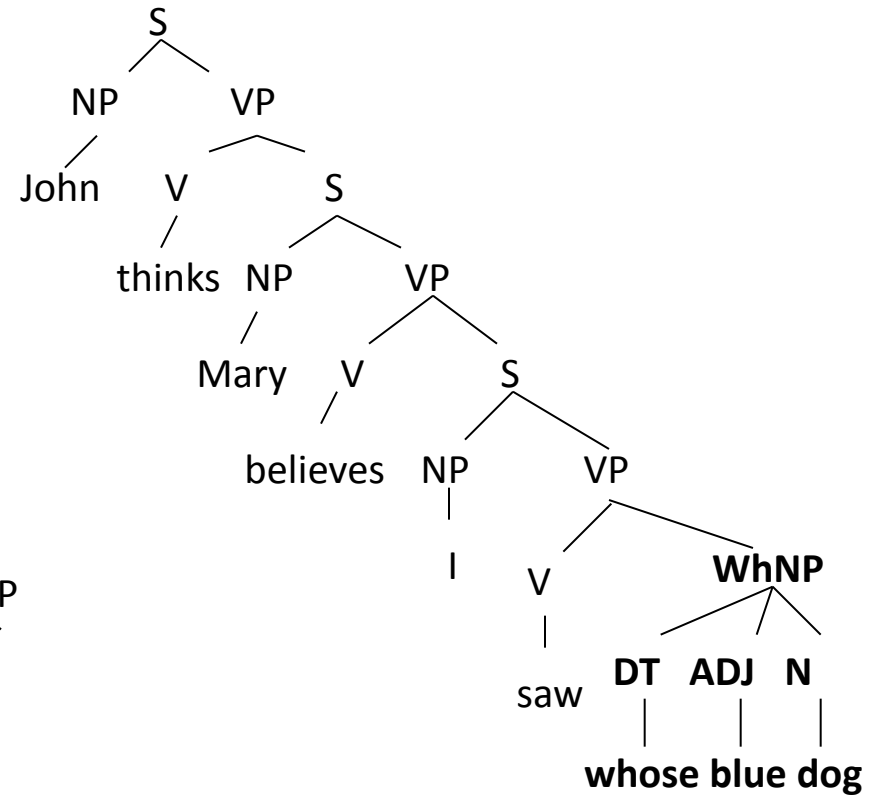
Long distance movement
of bounded material

Expressiveness

Whose blue dog does John think Mary believes I saw? → John thinks Mary believes I saw **whose blue dog**?

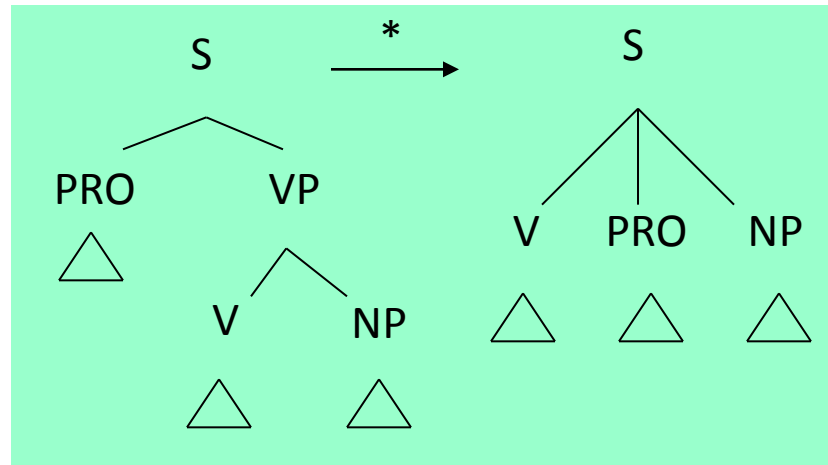


Long distance movement
of unbounded material



Expressiveness

Let's mainly focus on local rotation:



T – top-down
L – linear (non-copying)
N – non-deleting



T – top-down
 L – linear (non-copying)
 N – non-deleting

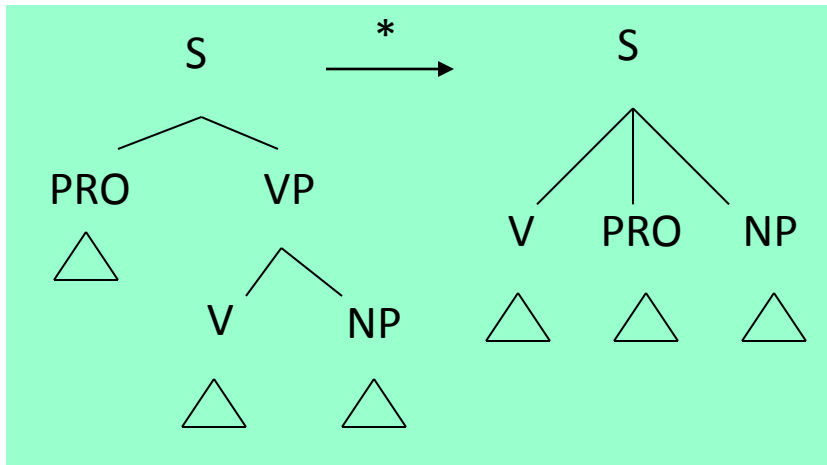
copying

non-copying

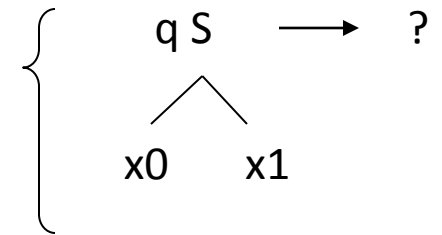
deleting

non-deleting

Expressiveness:



T
 ↑
 LT
 ↑
 LNT



T – top-down
 L – linear (non-copying)
 N – non-deleting

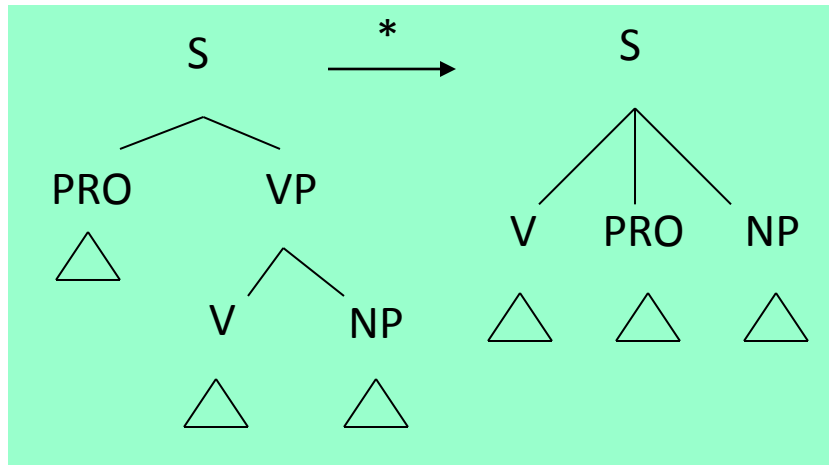
copying

non-copying

deleting

non-deleting

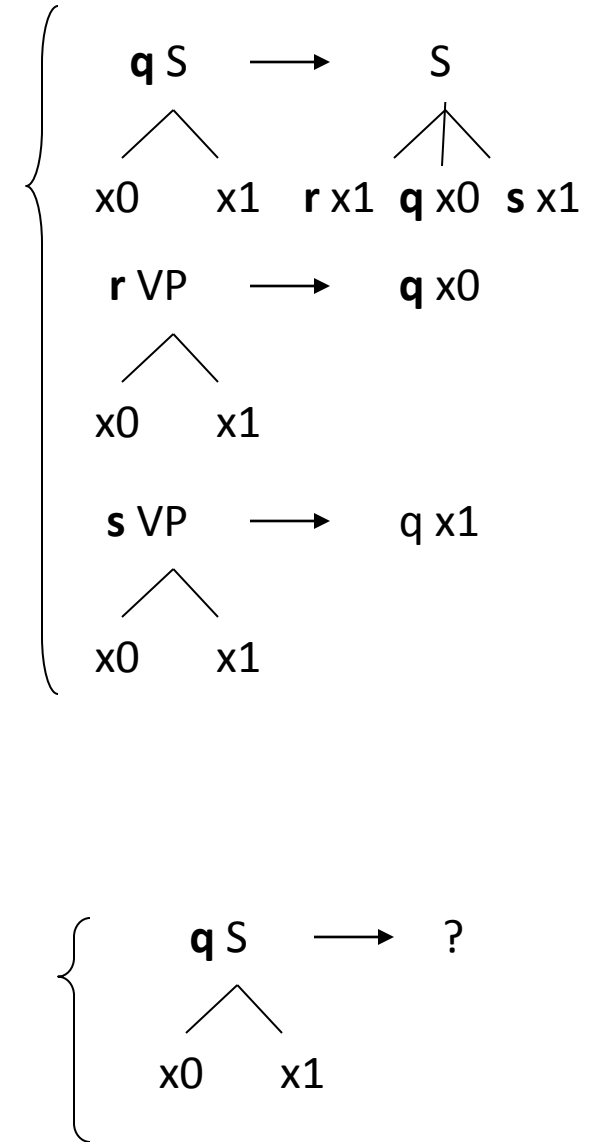
Expressiveness:



T

LT

LNT

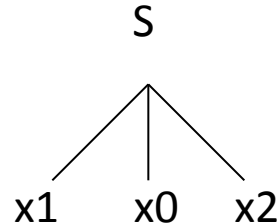
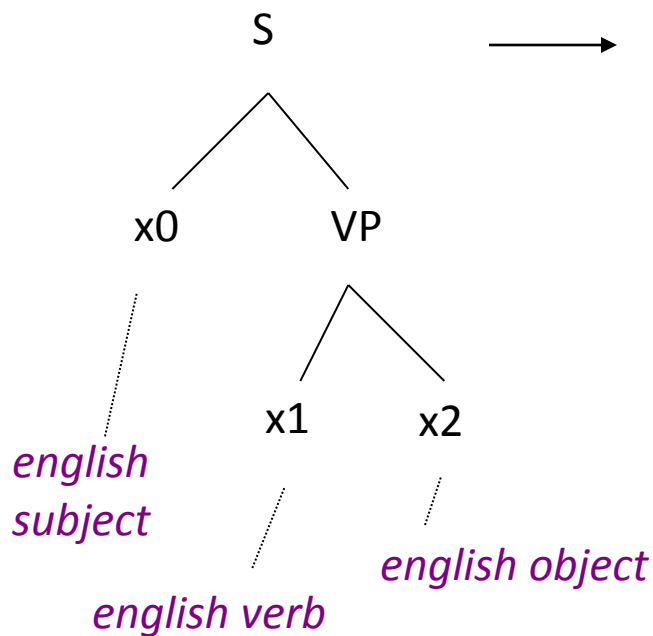


Extended (x-) Transducers

multilevel LHS

multilevel RHS

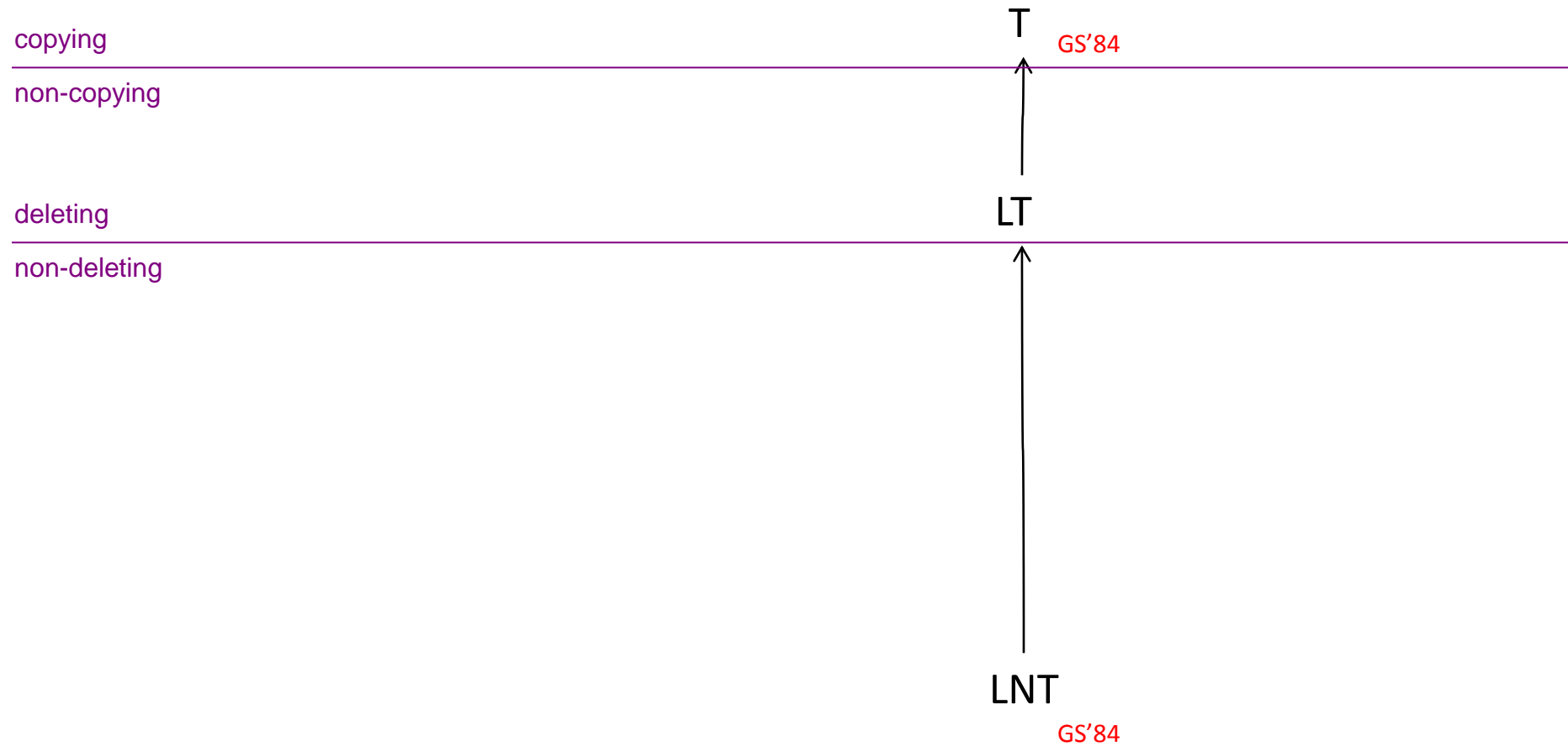
T – top-down
L – linear (non-copying)
N – non-deleting
x – extended LHS

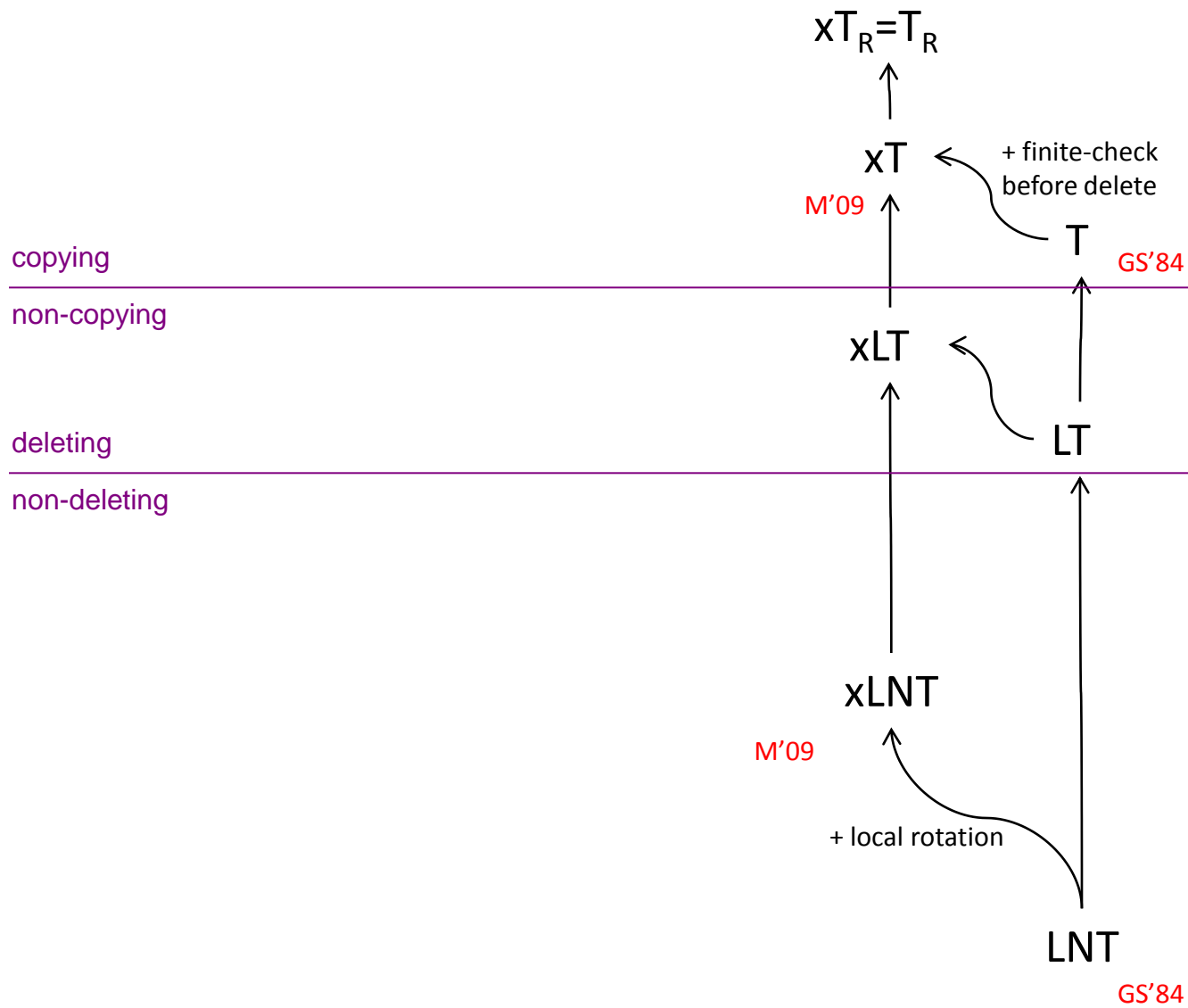


xLNT

can grab more structure

- possibility mentioned in [Rounds 70]
- variant defined in [Dauchet 76]
- used for practical MT by [Galley et al 04, 06]
- studied formally by [Maletti et al 09]





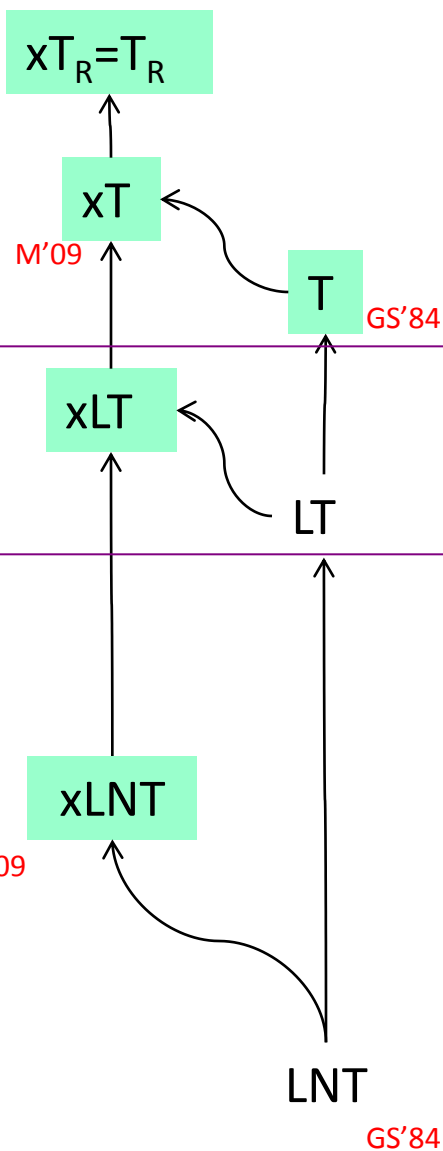
Expressive enough for local rotation

copying

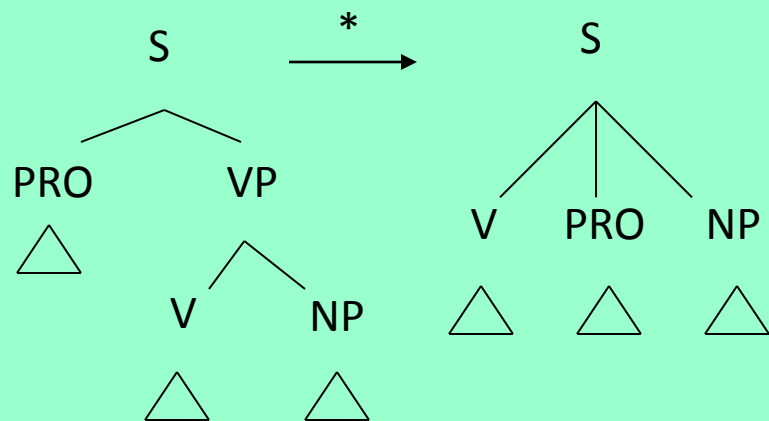
non-copying

deleting

non-deleting



Expressiveness:



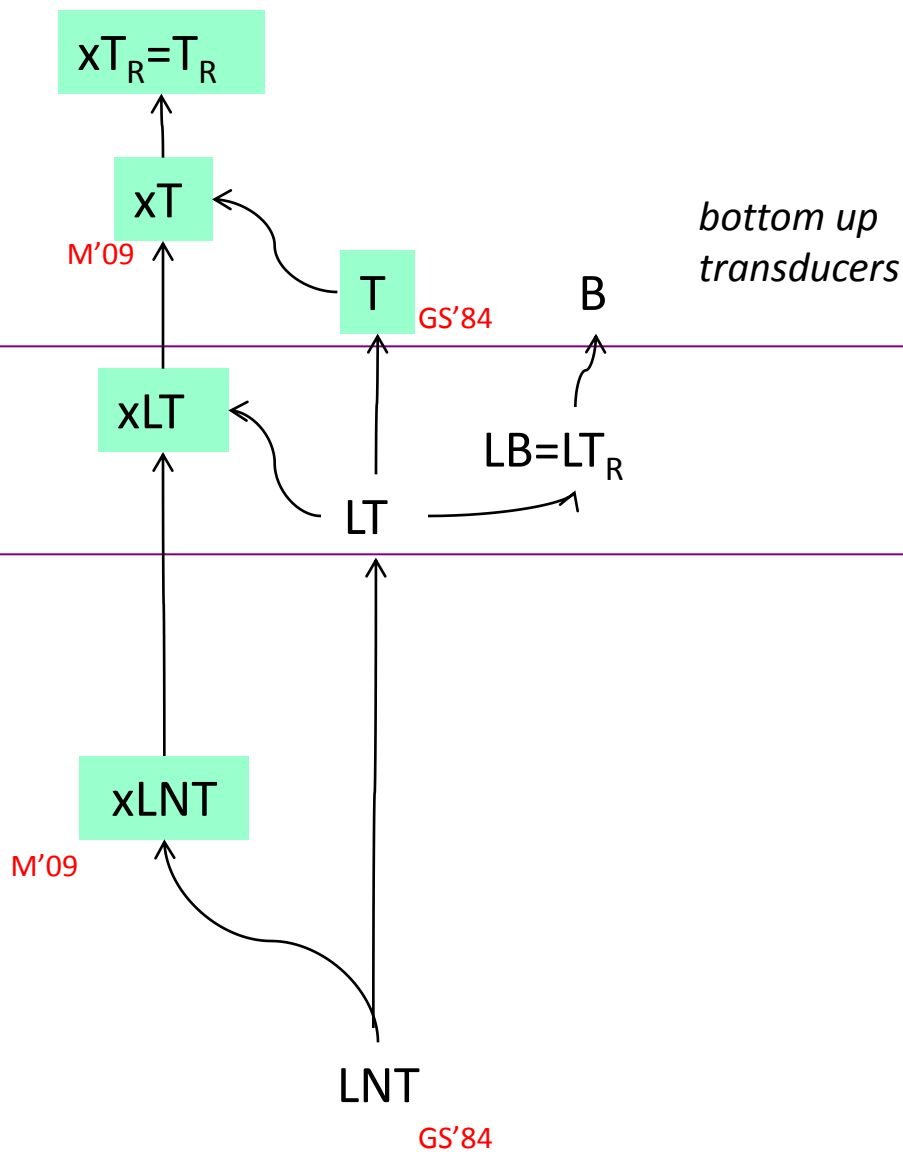
Expressive enough for local rotation

copying

non-copying

deleting

non-deleting



Expressive enough for local rotation

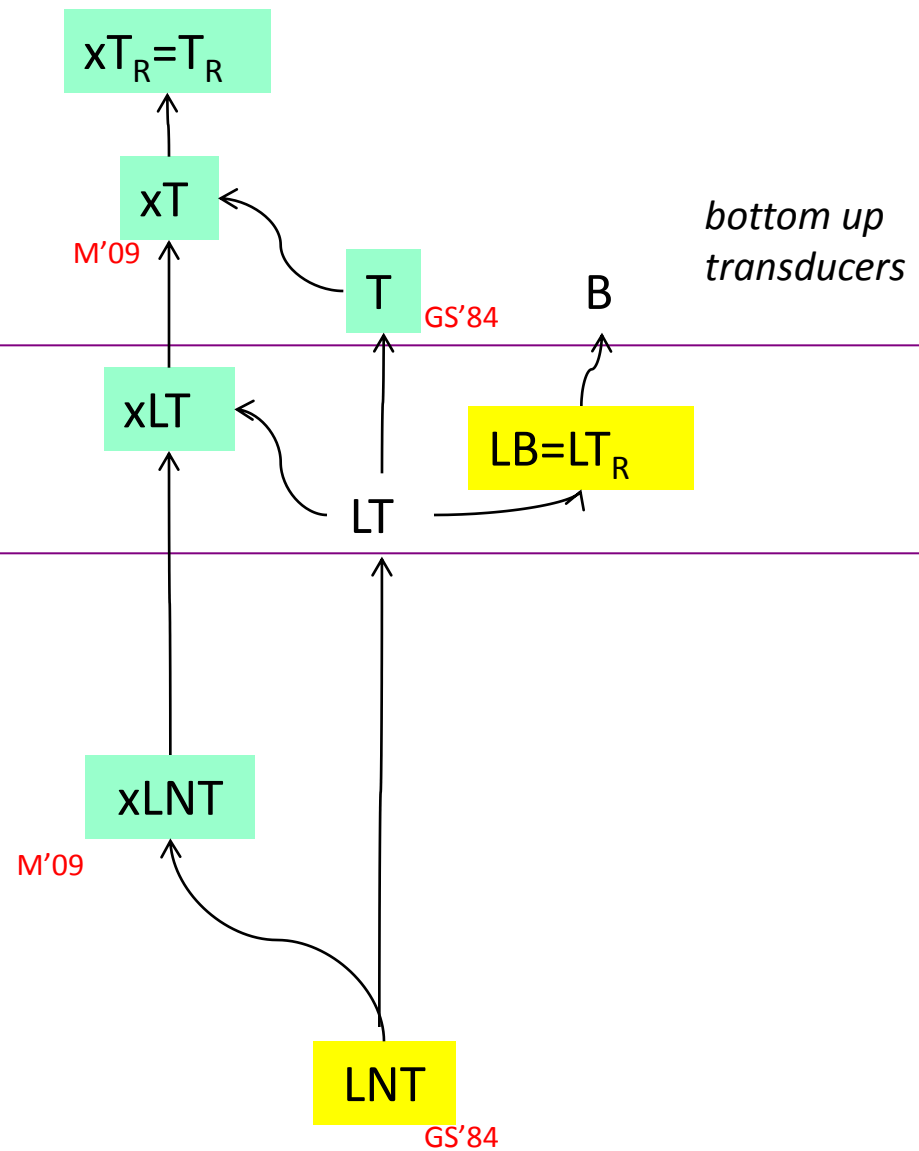
Closed under composition

copying

non-copying

deleting

non-deleting



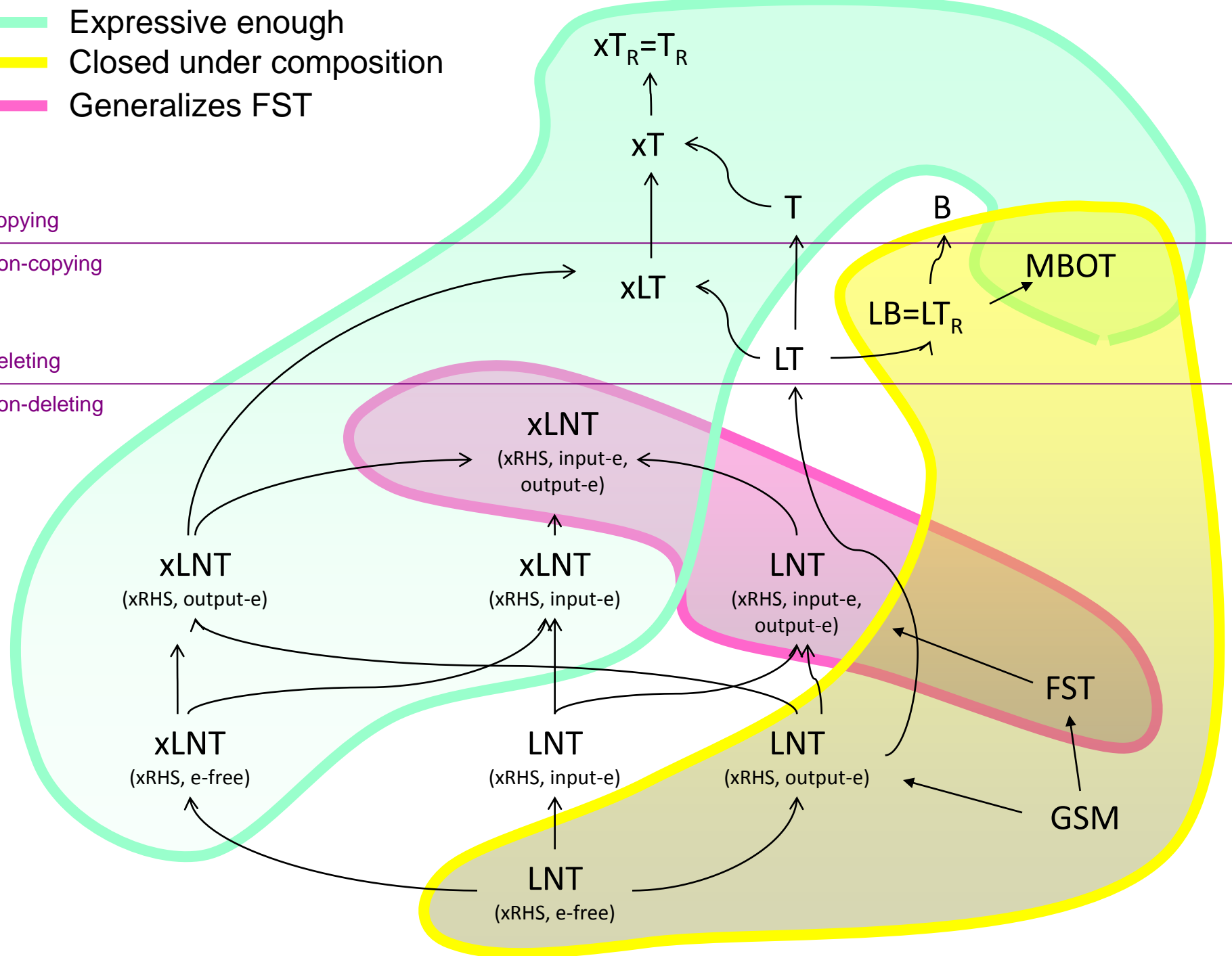
- Expressive enough
- Closed under composition
- Generalizes FST

copying

non-copying

deleting

non-deleting



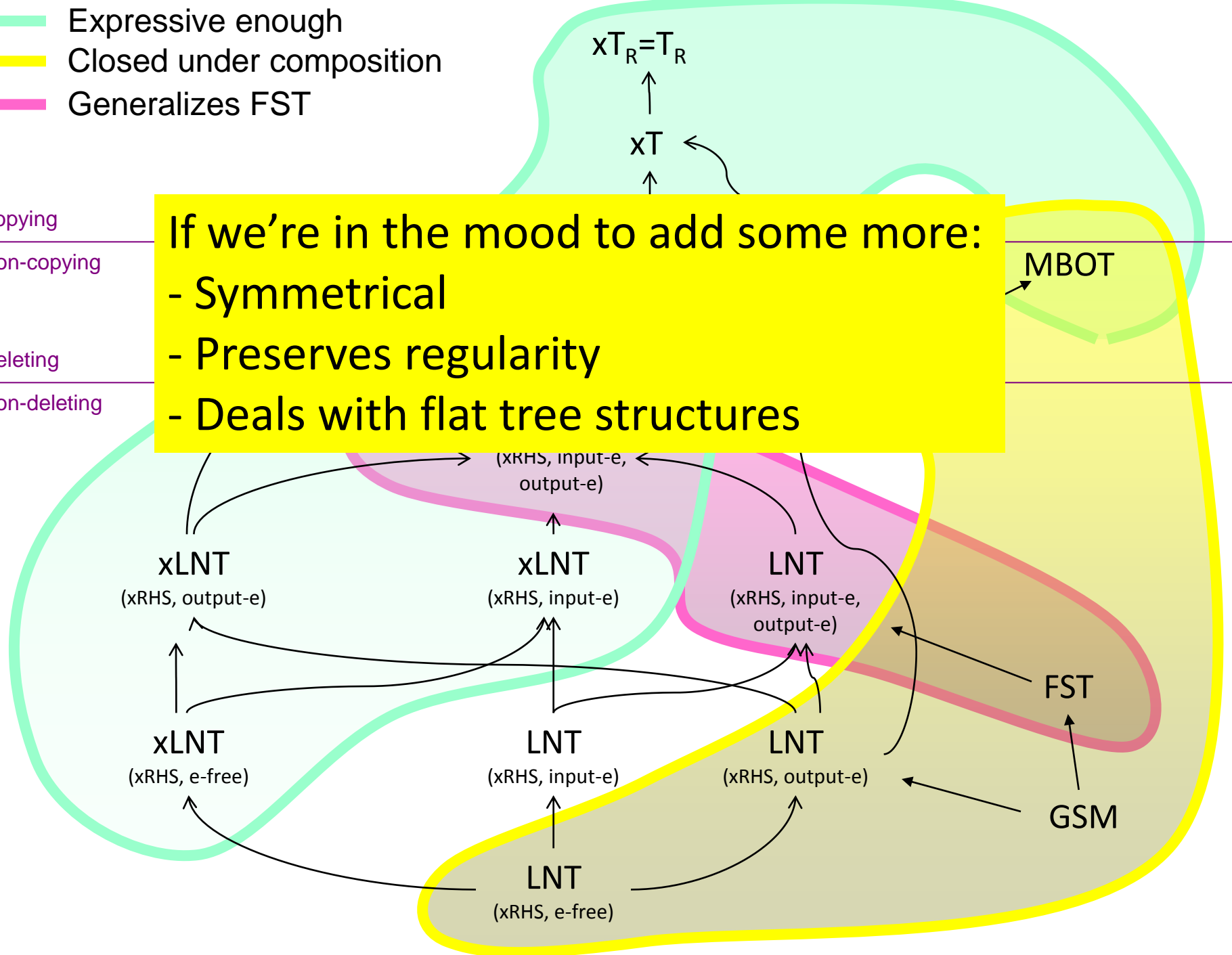
- Expressive enough
- Closed under composition
- Generalizes FST

If we're in the mood to add some more:

- Symmetrical
- Preserves regularity
- Deals with flat tree structures

copying
non-copying

deleting
non-deleting



General-Purpose Algorithms for Tree Automata

	String Automata Algorithms	Tree Automata Algorithms
N-best paths through an WFSA (Viterbi, 1967; Eppstein, 1998)	... trees in a weighted forest (Jiménez & Marzal, 2000; Huang & Chiang, 2005)
EM training	Forward-backward EM (Baum/Welch, 1971; Eisner 2003)	Tree transducer EM training (Graehl & Knight, 2004)
Determinization of weighted string acceptors (Mohri, 1997)	... of weighted tree acceptors (Borchardt & Vogler, 2003; May & Knight, 2005)
Intersection	WFSA intersection	Tree acceptor intersection
Applying transducers	string \rightarrow WFST \rightarrow WFSA	tree \rightarrow TT \rightarrow weighted tree acceptor
Transducer composition	WFST composition (Pereira & Riley, 1996)	Many tree transducers not closed under composition (Maletti et al 09)
General-purpose tools	FSM, Carmel, OpenFST	Tiburon (May & Knight 10)

Machine Translation

Phrase-based MT



Syntax-based MT



Meaning-based MT



Equivalent Semantic Representations

"The boy wants to go."

LOGICAL FORM

$\exists w, b, g : \text{instance}(w, \text{WANT}) \wedge$
 $\text{instance}(g, \text{GO}) \wedge$
 $\text{instance}(b, \text{BOY}) \wedge$
 $\text{agent}(w, b) \wedge$
 $\text{patient}(w, g) \wedge$
 $\text{agent}(g, b)$

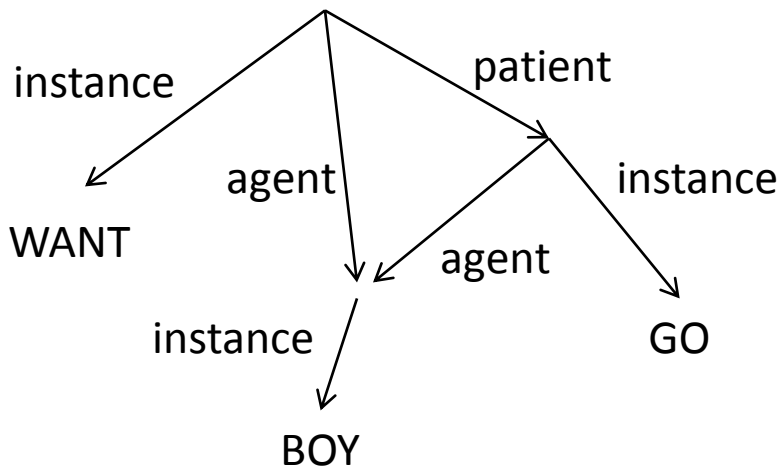
PENMAN

(w / WANT
:agent (b / BOY)
:patient (g / GO
:agent b)))

PATH EQUATIONS

((x0 instance) = WANT
((x1 instance) = BOY
((x2 instance) = GO
((x0 agent) = x1
((x0 patient) = x2
((x2 agent) = x1

DIRECTED ACYCLIC GRAPH



FEATURE STRUCTURE

instance: WANT
agent: [1 [instance: BOY]]
patient: [instance: GO
agent: [1]]

Equivalent Semantic Representations

“The boy wants to go.”

LOGICAL FORM

$\exists w, b, g : \text{instance}(w, \text{WANT}) \wedge$
 $\text{instance}(g, \text{GO}) \wedge$
 $\text{instance}(b, \text{BOY}) \wedge$
 $\text{agent}(w, b) \wedge$
 $\text{patient}(w, g) \wedge$
 $\text{agent}(g, b)$

PENMAN

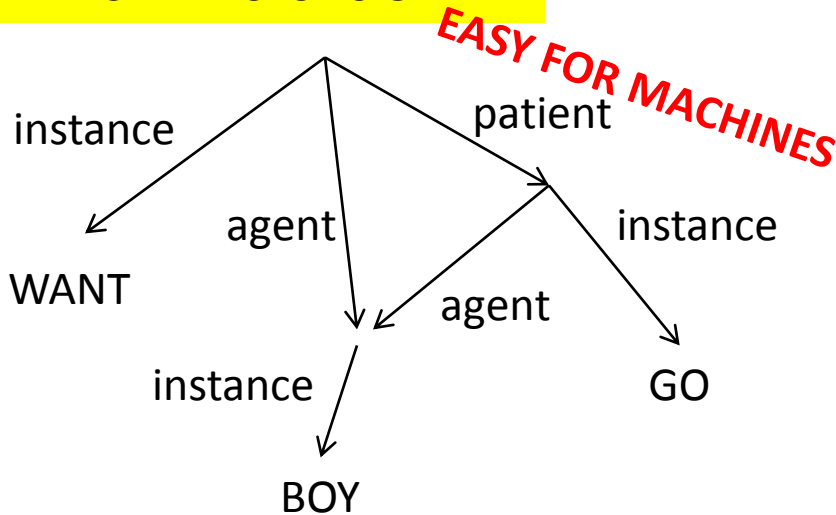
(w / WANT
:agent (b / BOY)
:patient (g / GO
:agent b)))

EASY FOR PEOPLE

PATH EQUATIONS

((x0 instance) = WANT
((x1 instance) = BOY
((x2 instance) = GO
((x0 agent) = x1
((x0 patient) = x2
((x2 agent) = x1

DIRECTED ACYCLIC GRAPH



FEATURE STRUCTURE

instance: WANT
agent: [1 [instance: BOY]]
patient: [instance: GO
agent: [1]]

Example

“Government forces closed on rebel outposts on Thursday, showering the western mountain city of Zintan with missiles and attacking insurgents holed up near the Tunisian border, according to rebel sources.”

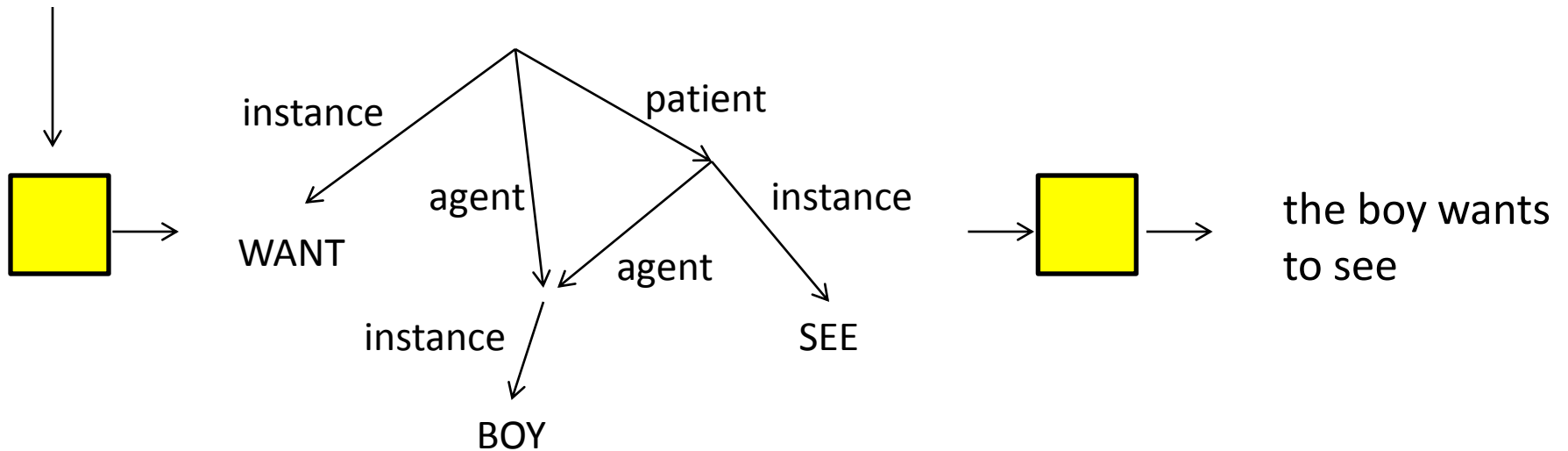
```
(s / say
  :agent (s2 / source :mod (r / rebel))
  :patient (a / and
    :op1 (c / close-on
      :agent (f / force :mod (g / government))
      :patient (o / outpost :mod (r2 / rebel))
      :temporal-locating (t / thursday))
    :op2 (s / shower
      :agent f
      :patient (c2 / city
        :mod (m / mountain)
        :mod (w / west)
        :name "Zintan")
      :instrument (m2 / missile))
    :op3 (a2 / attack
      :agent f
      :patient (i / insurgent
        :agent-of (h / hole-up
          :pp-near (b / border :poss (c3 / country
            :name "Tunisia"]
```

“more logical
than a parse tree”

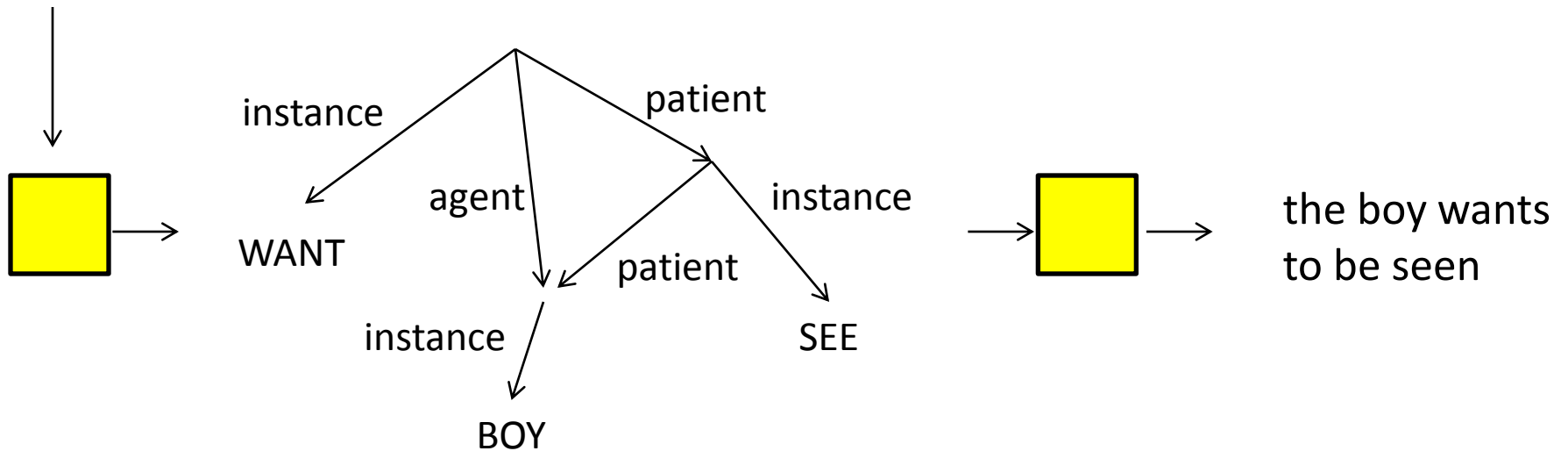
General-Purpose Algorithms for Feature Structures (Graphs)

	String Automata Algorithms	Tree Automata Algorithms	Graph Automata Algorithms?
N-best paths through an WFSA (Viterbi, 1967; Eppstein, 1998)	... trees in a weighted forest (Jiménez & Marzal, 2000; Huang & Chiang, 2005)	
EM training	Forward-backward EM (Baum/Welch, 1971; Eisner 2003)	Tree transducer EM training (Graehl & Knight, 2004)	
Determinization...	... of weighted string acceptors (Mohri, 1997)	... of weighted tree acceptors (Borchardt & Vogler, 2003; May & Knight, 2005)	
Intersection	WFSA intersection	Tree acceptor intersection	
Applying transducers	string \rightarrow WFST \rightarrow WFSA	tree \rightarrow TT \rightarrow weighted tree acceptor	
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General tools	FSM, Carmel, OpenFST	Tiburón (May & Knight 10)	

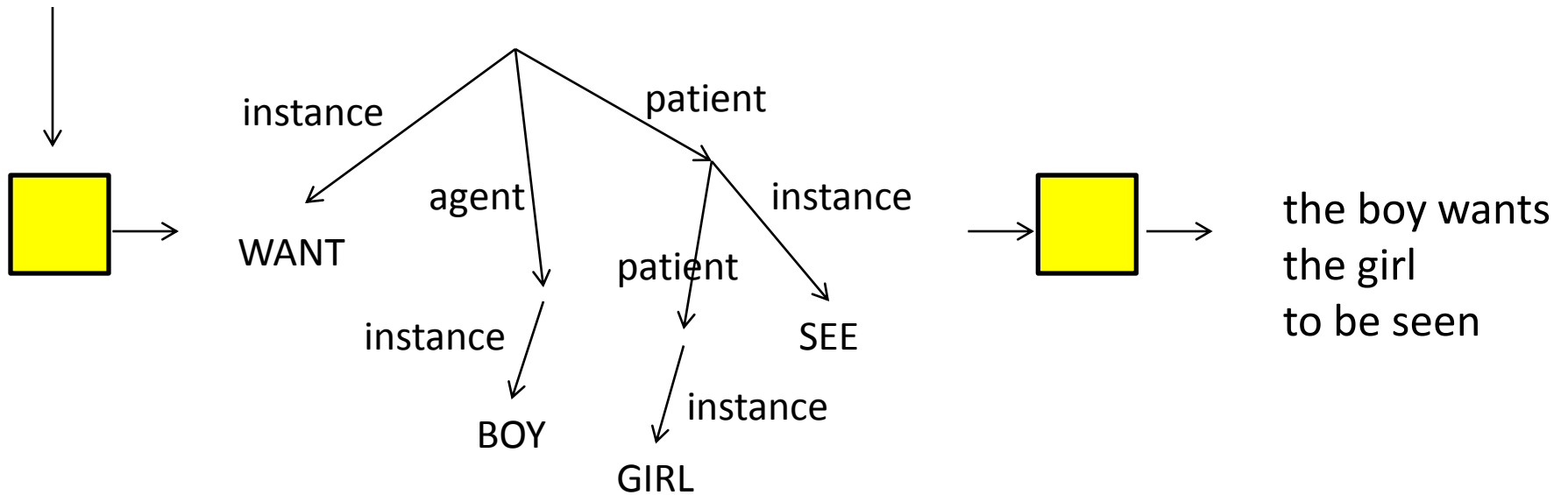
Mapping Between Meaning and Text



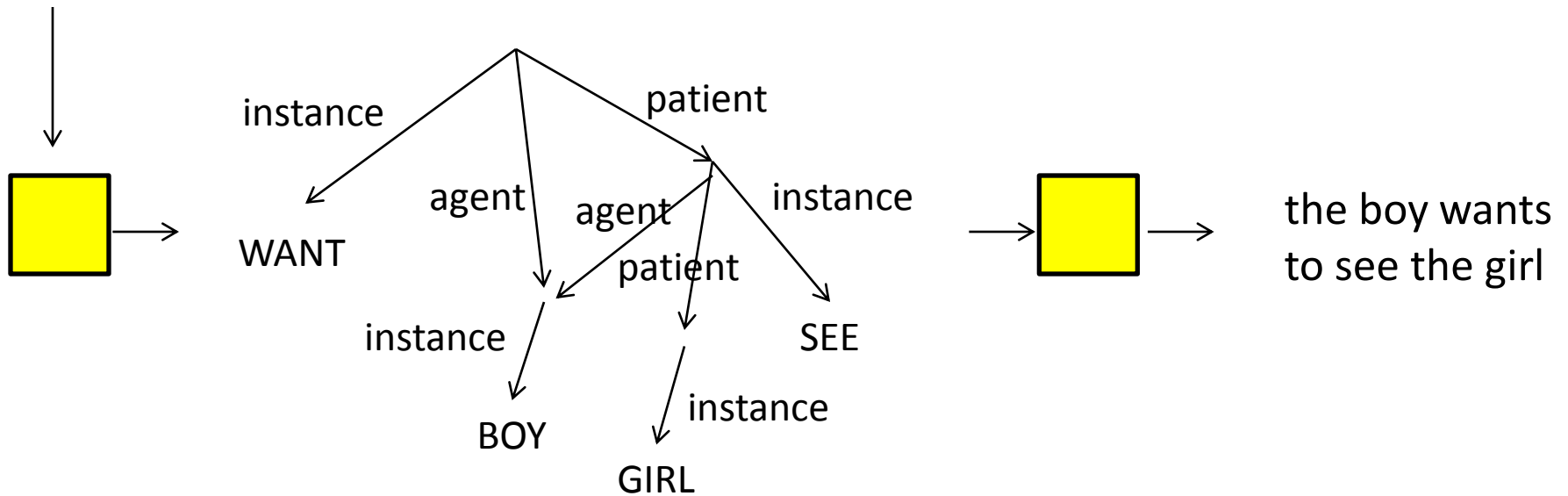
Mapping Between Meaning and Text



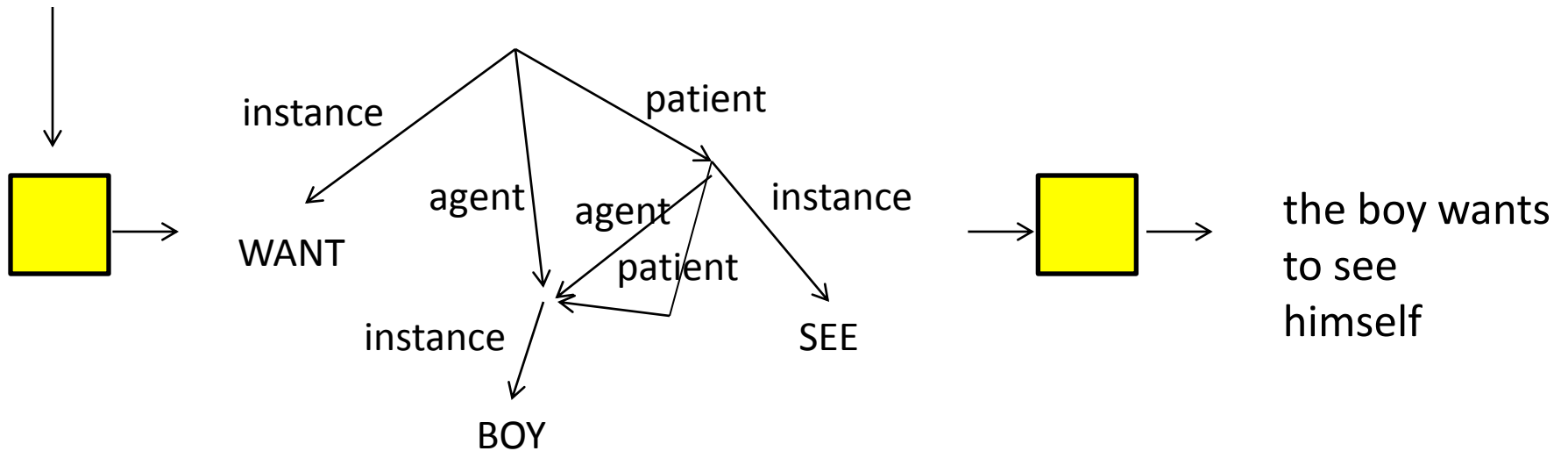
Mapping Between Meaning and Text



Mapping Between Meaning and Text



Mapping Between Meaning and Text



Automata Frameworks

- Unification grammar: string-to-semantics
(Moore 89)
- Hyperedge-replacement graph grammars
(Drewes et al 97)
- DAG acceptors (Hart 75)
- DAG-to-tree transducers (Kamimura & Slutski 82)

Automata Frameworks

- Unification grammar (Moore 89)
- Hyperedge-replacement (Drewes et al 97)
- DAG acceptors (Hart)
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Math. Systems Theory 15, 225–249 (1982)

Mathematical
Systems Theory

Transductions of Dags and Trees

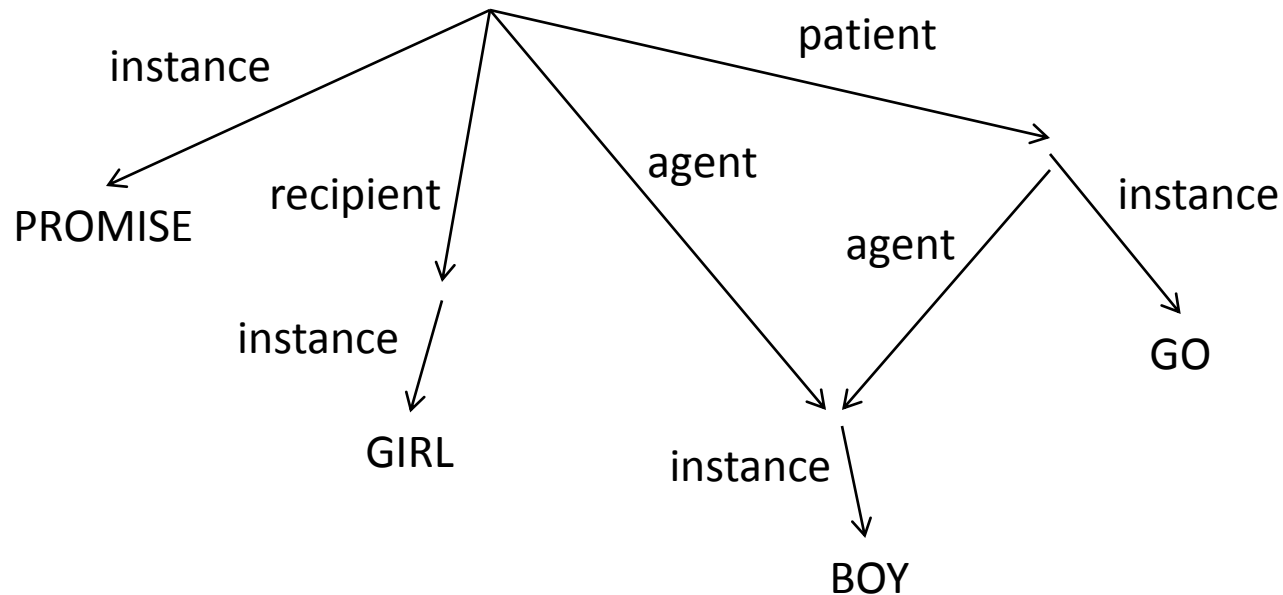
Tsutomu Kamimura* and Giora Slutzki

Department of Computer Science, University of Kansas, Lawrence, Kansas

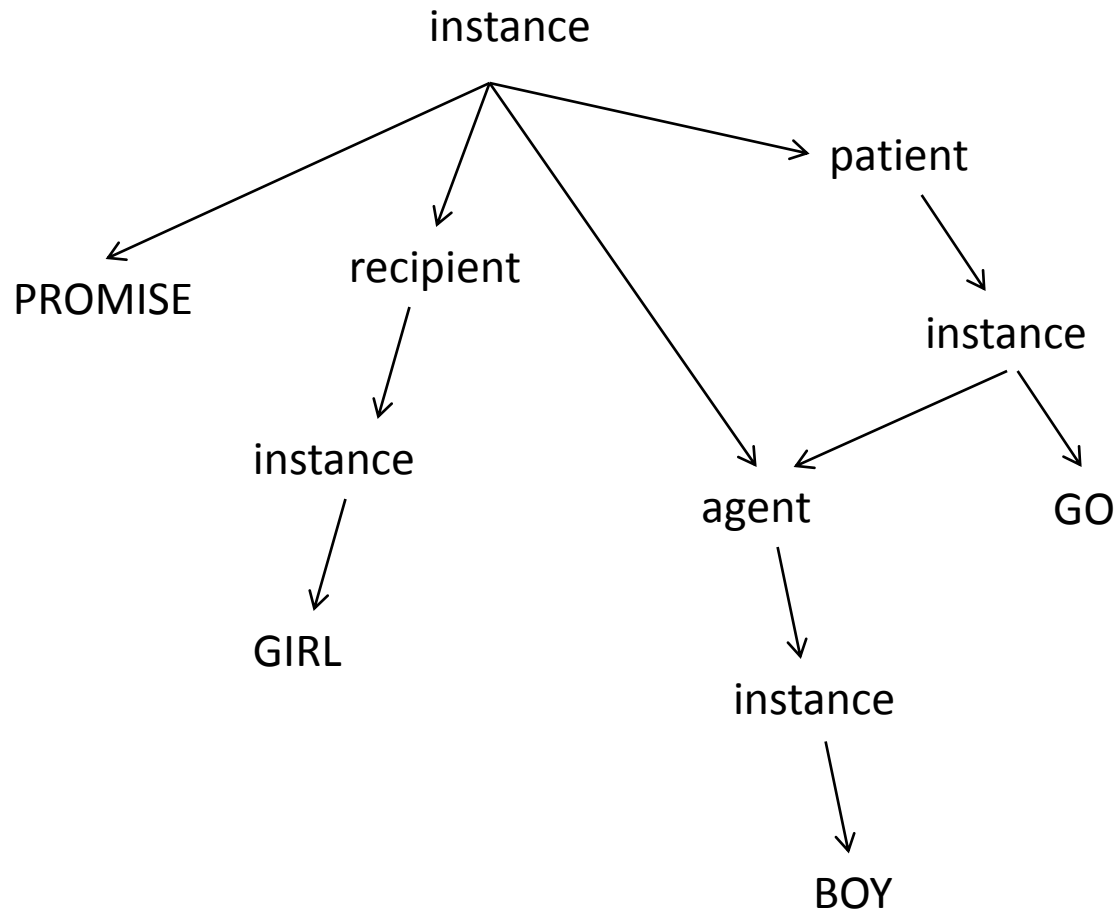
Abstract. Directed acyclic graphs (dags) model derivations of phrase-structure grammars analogously to the way that trees model derivations of context-free grammars.

In this paper we introduce translations of such dags which naturally extend the bottom-up tree translations. Composition results of these dag-to-tree transformations are studied. It is shown that every “recursively enumerable tree language” can be obtained from a recognizable dag language by such a transduction. Tree languages obtained from some subsets of recognizable dag languages by these transductions are investigated.

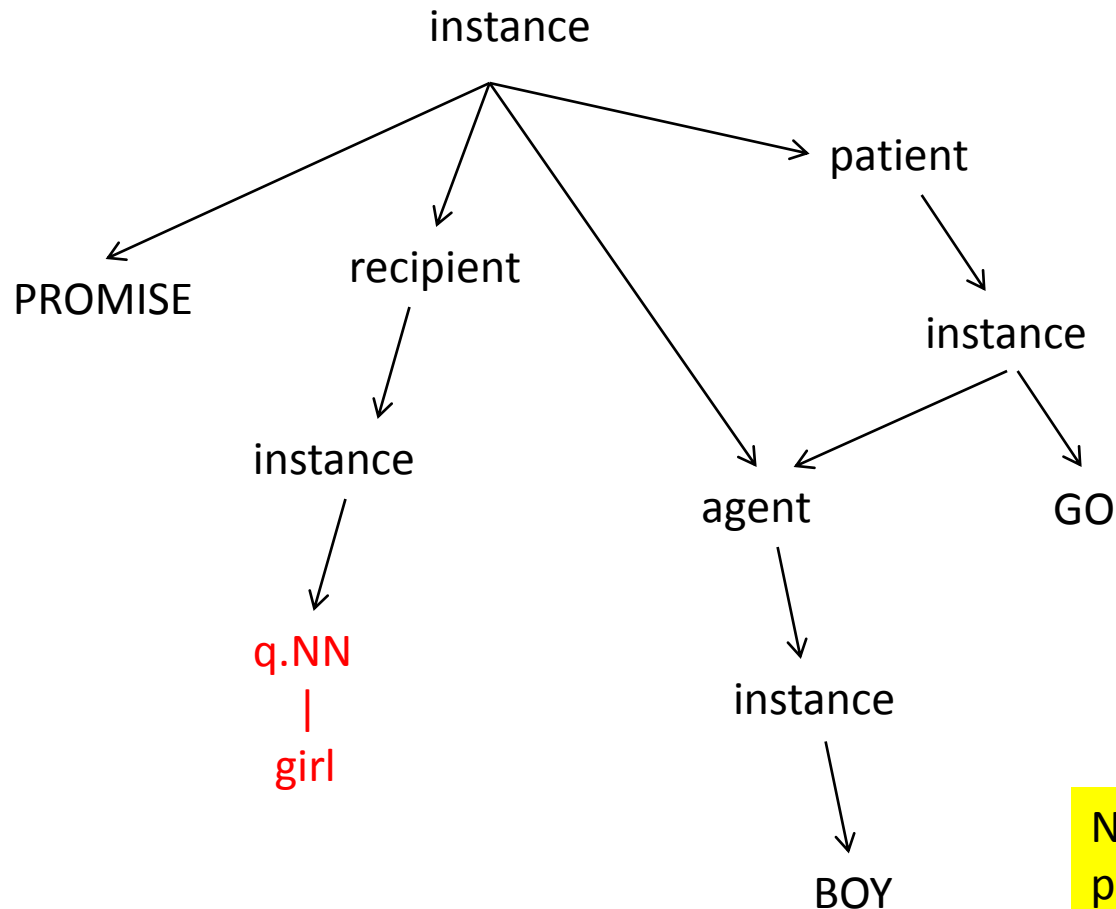
Bottom-Up DAG-to-Tree Transduction (Kamimura & Slutski 82)



Bottom-Up DAG-to-Tree Transduction (Kamimura & Slutski 82)

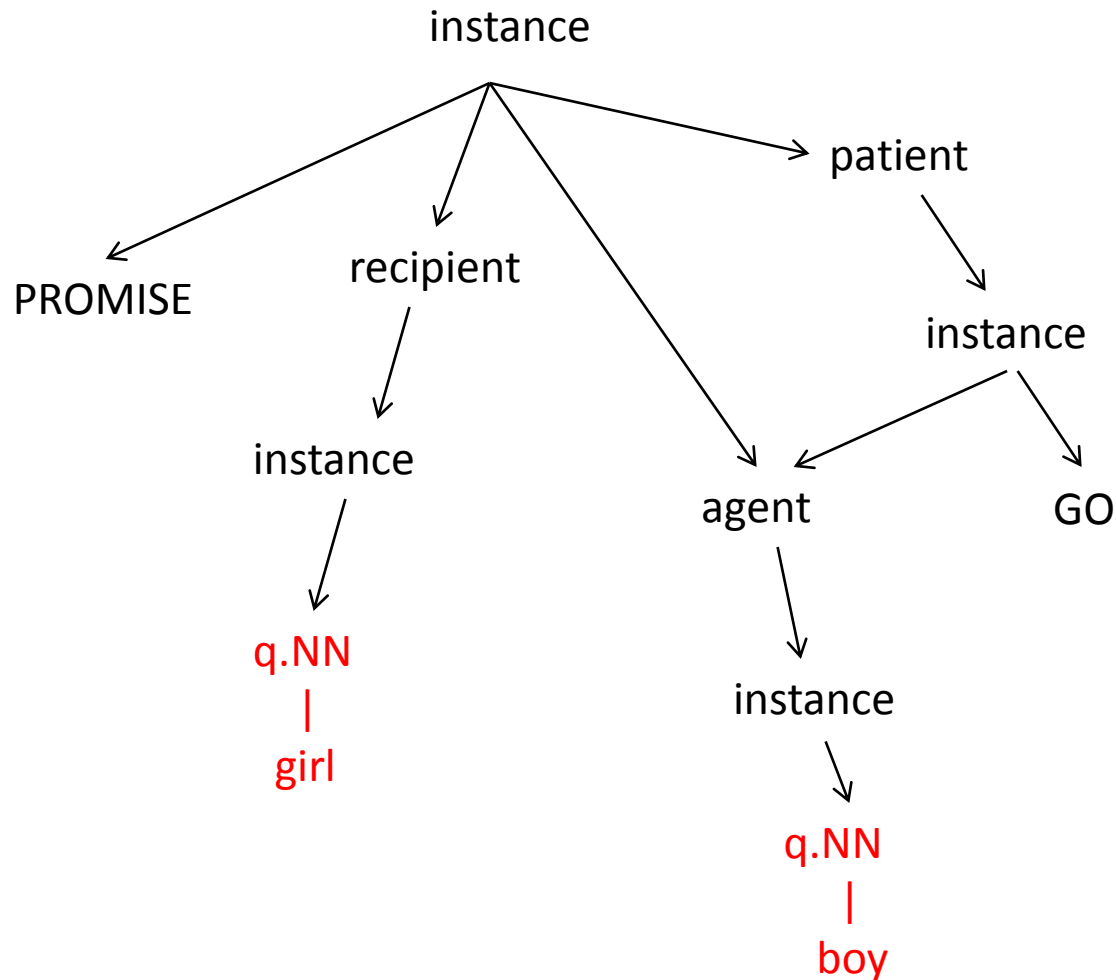


Bottom-Up DAG-to-Tree Transduction (Kamimura & Slutski 82)

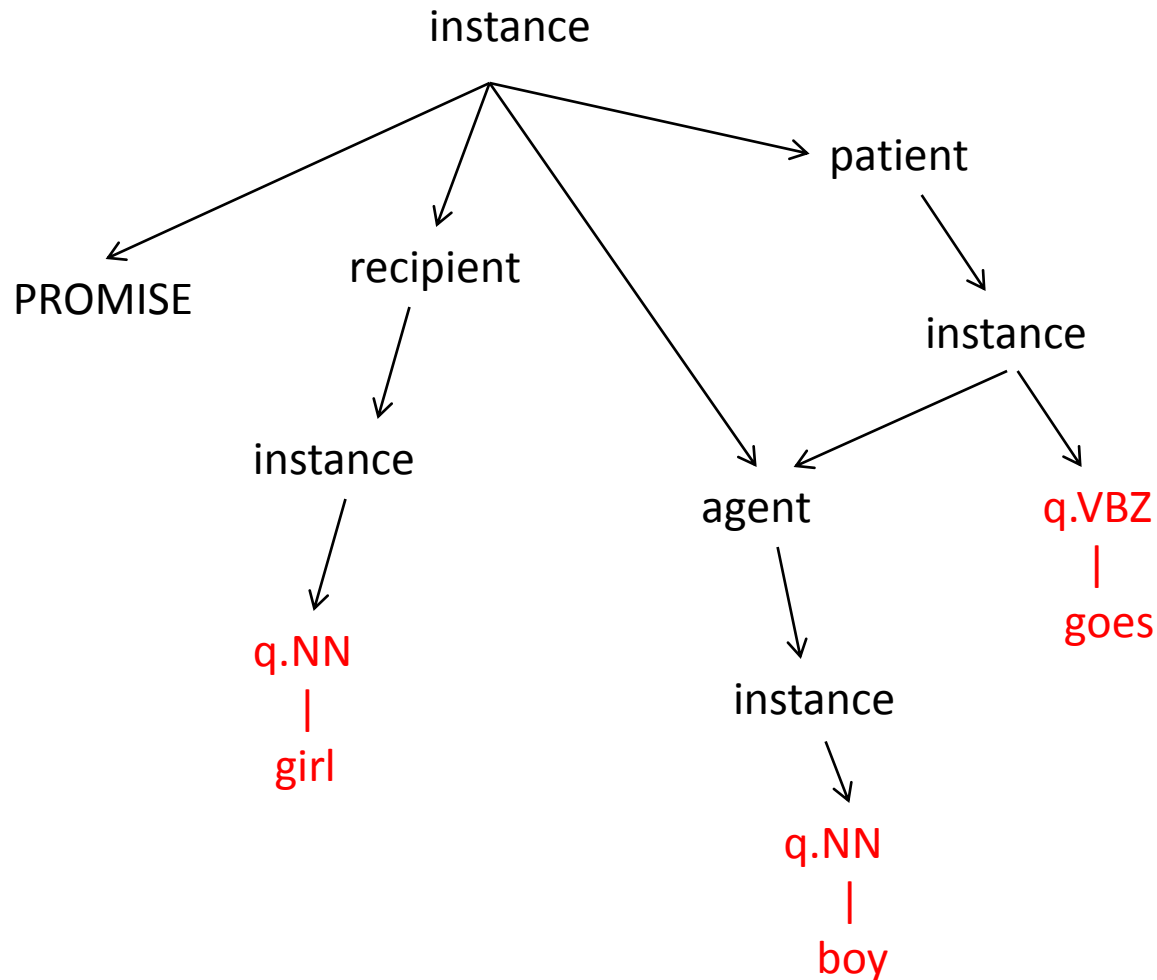


NOTE: States not
precise! Rules not
shown!

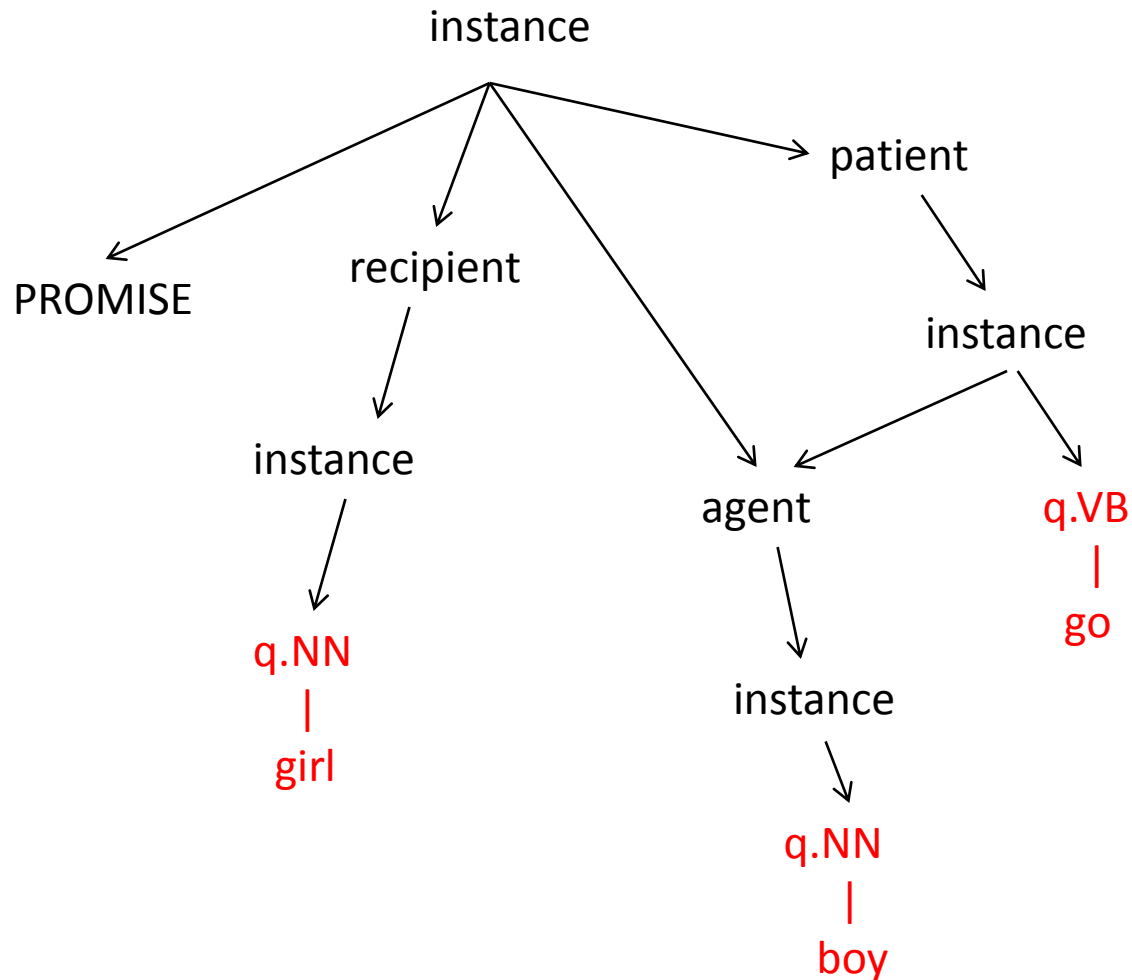
Bottom-Up DAG-to-Tree Transduction (Kamimura & Slutski 82)



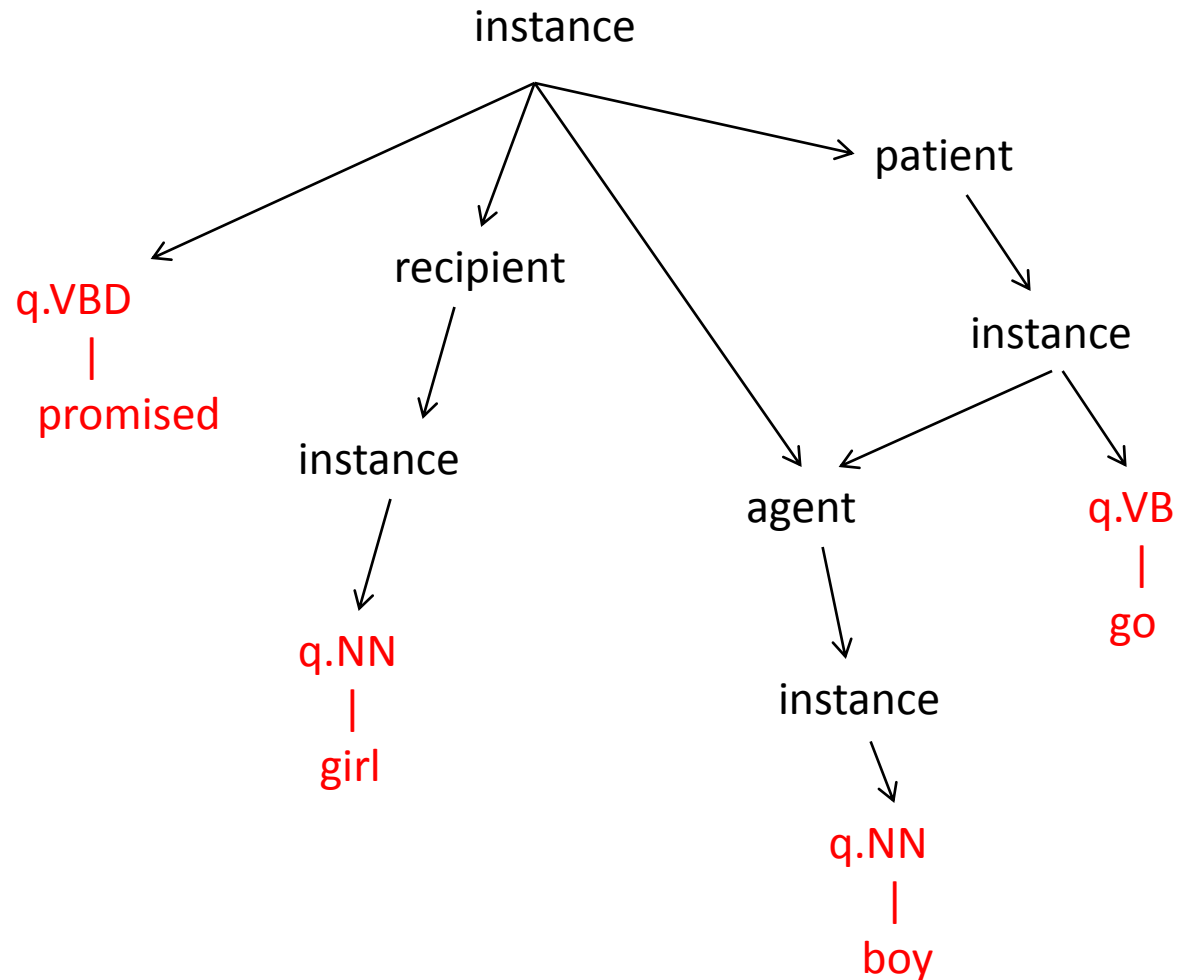
Bottom-Up DAG-to-Tree Transduction (Kamimura & Slutski 82)



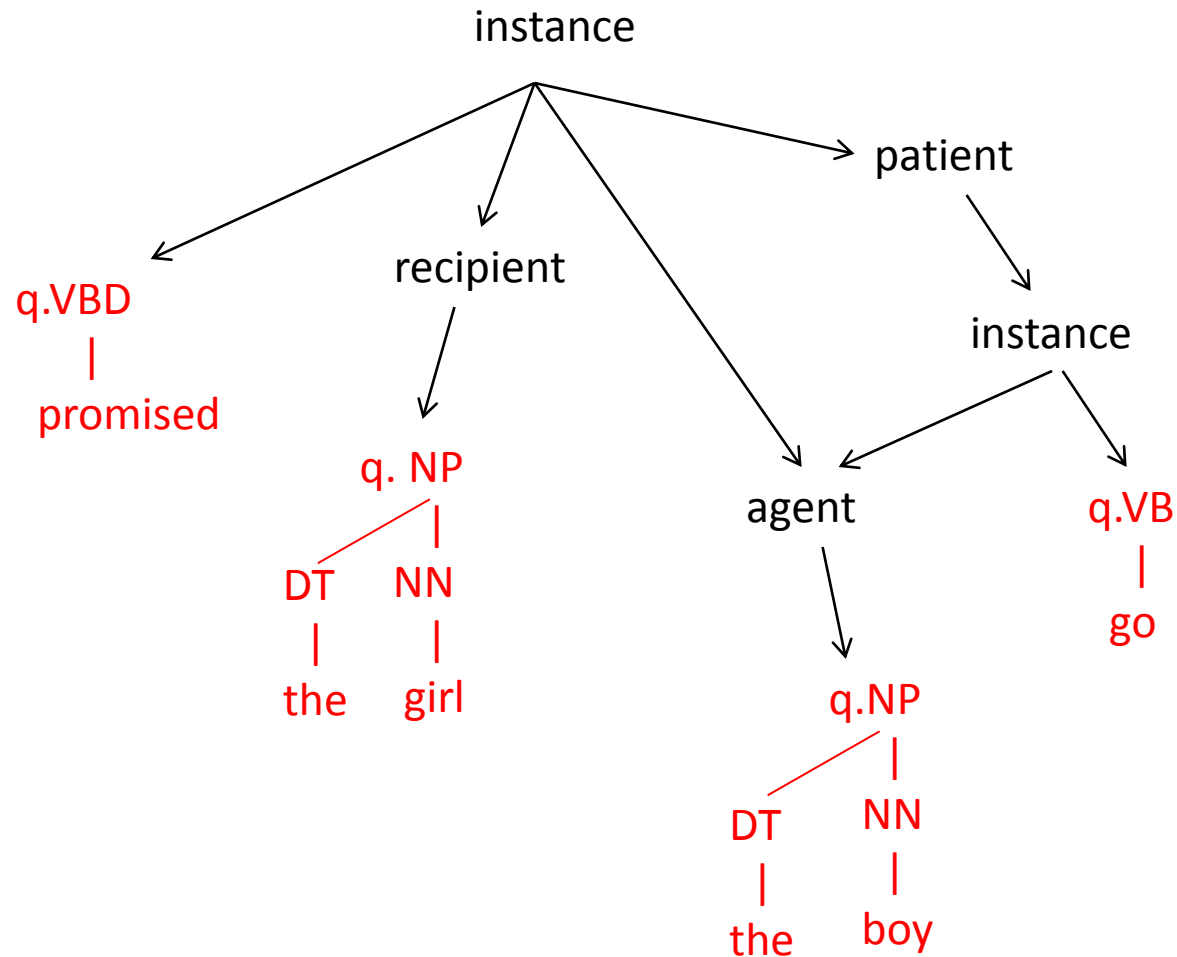
Bottom-Up DAG-to-Tree Transduction (Kamimura & Slutski 82)



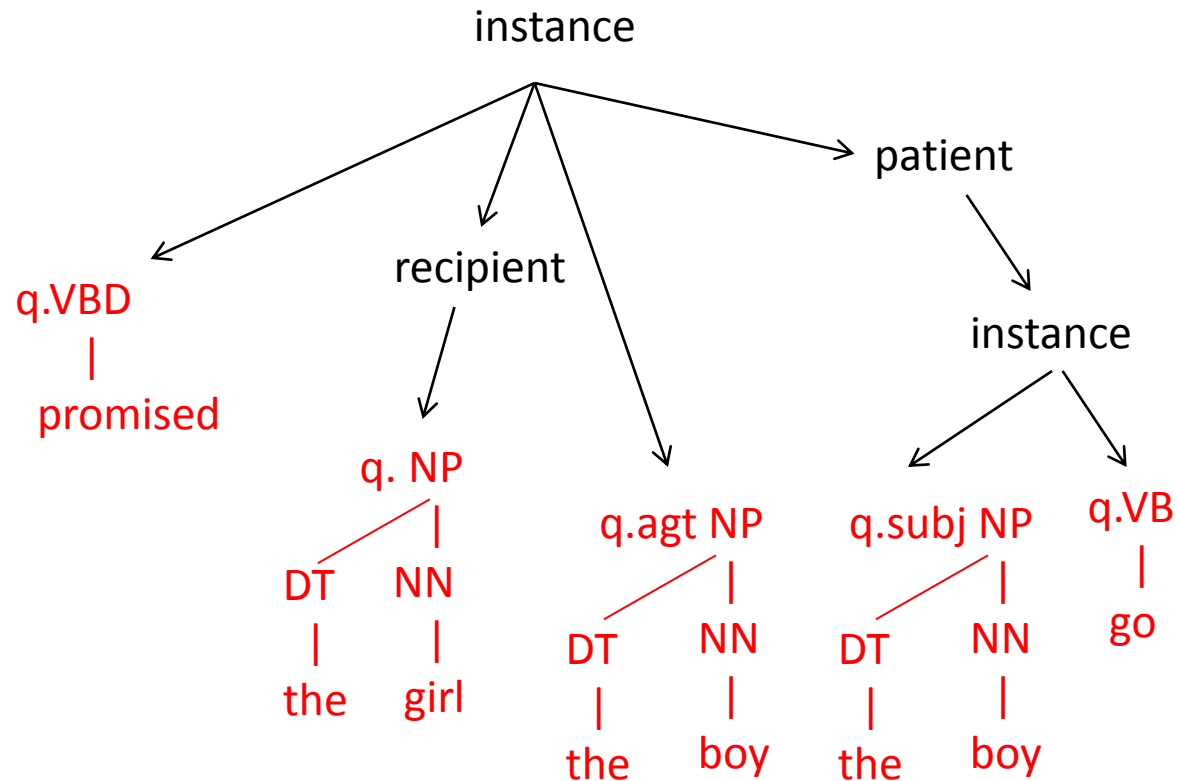
Bottom-Up DAG-to-Tree Transduction (Kamimura & Slutski 82)



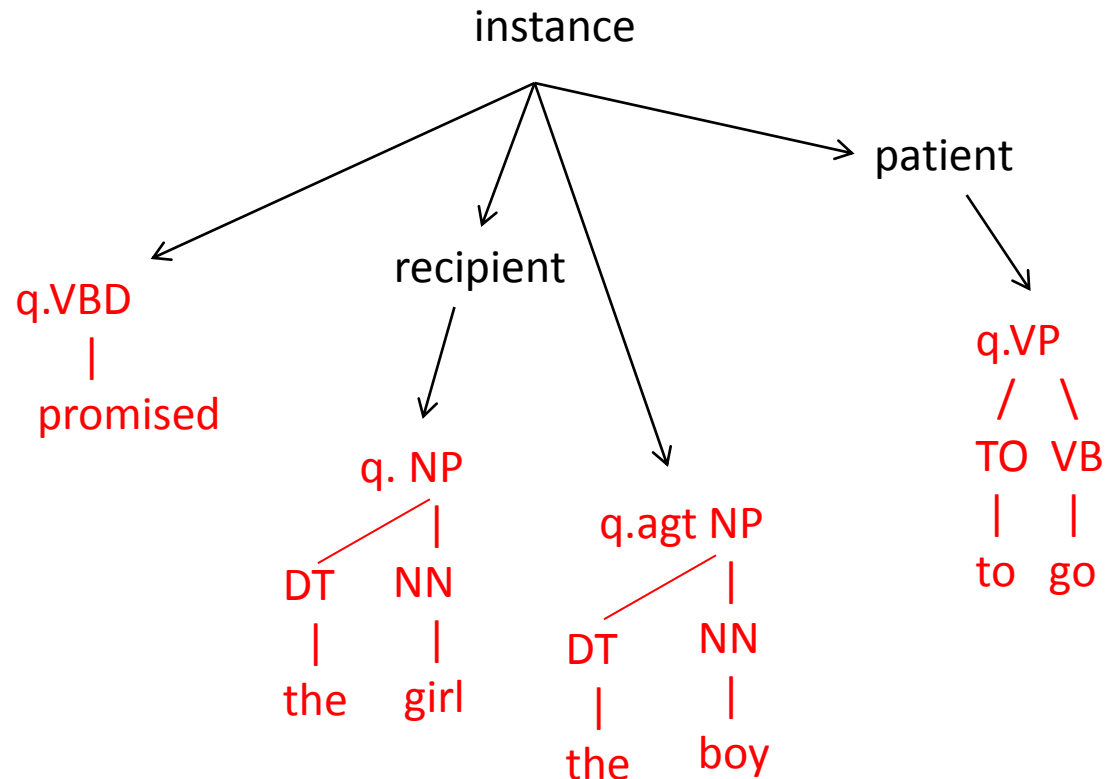
Bottom-Up DAG-to-Tree Transduction (Kamimura & Slutski 82)



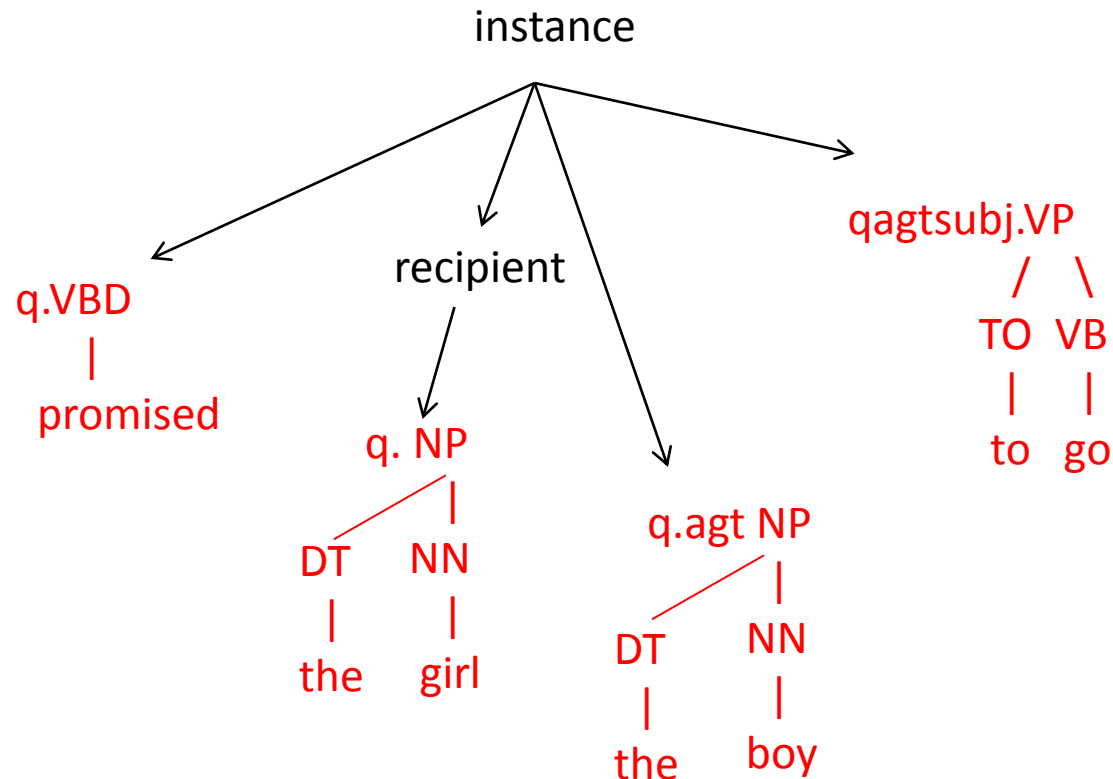
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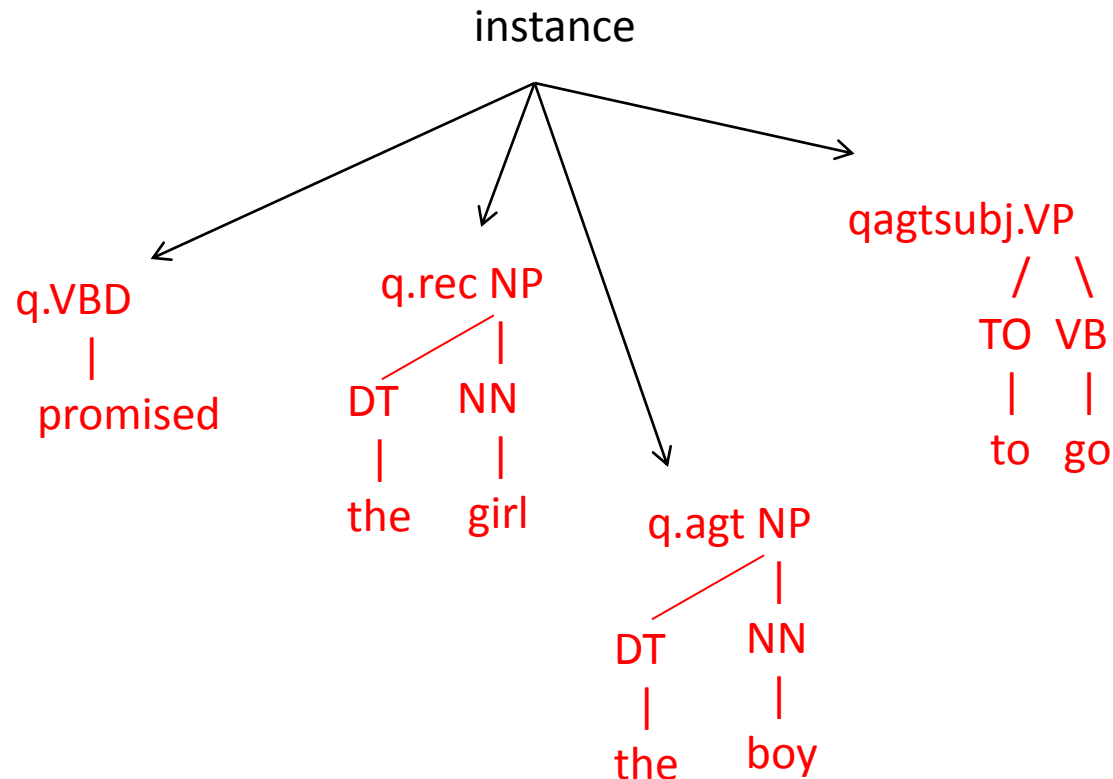
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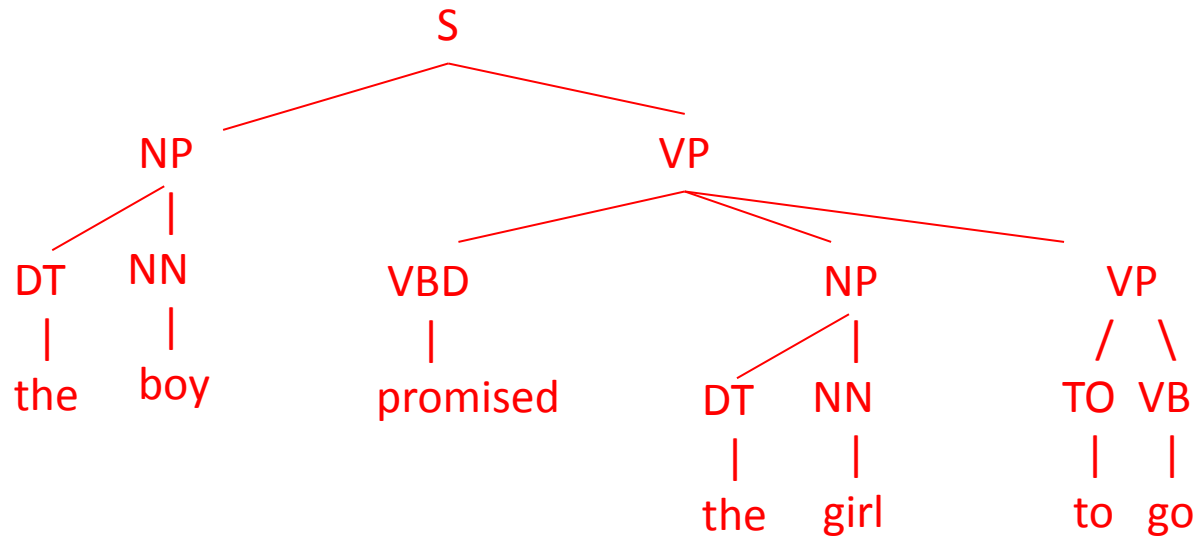
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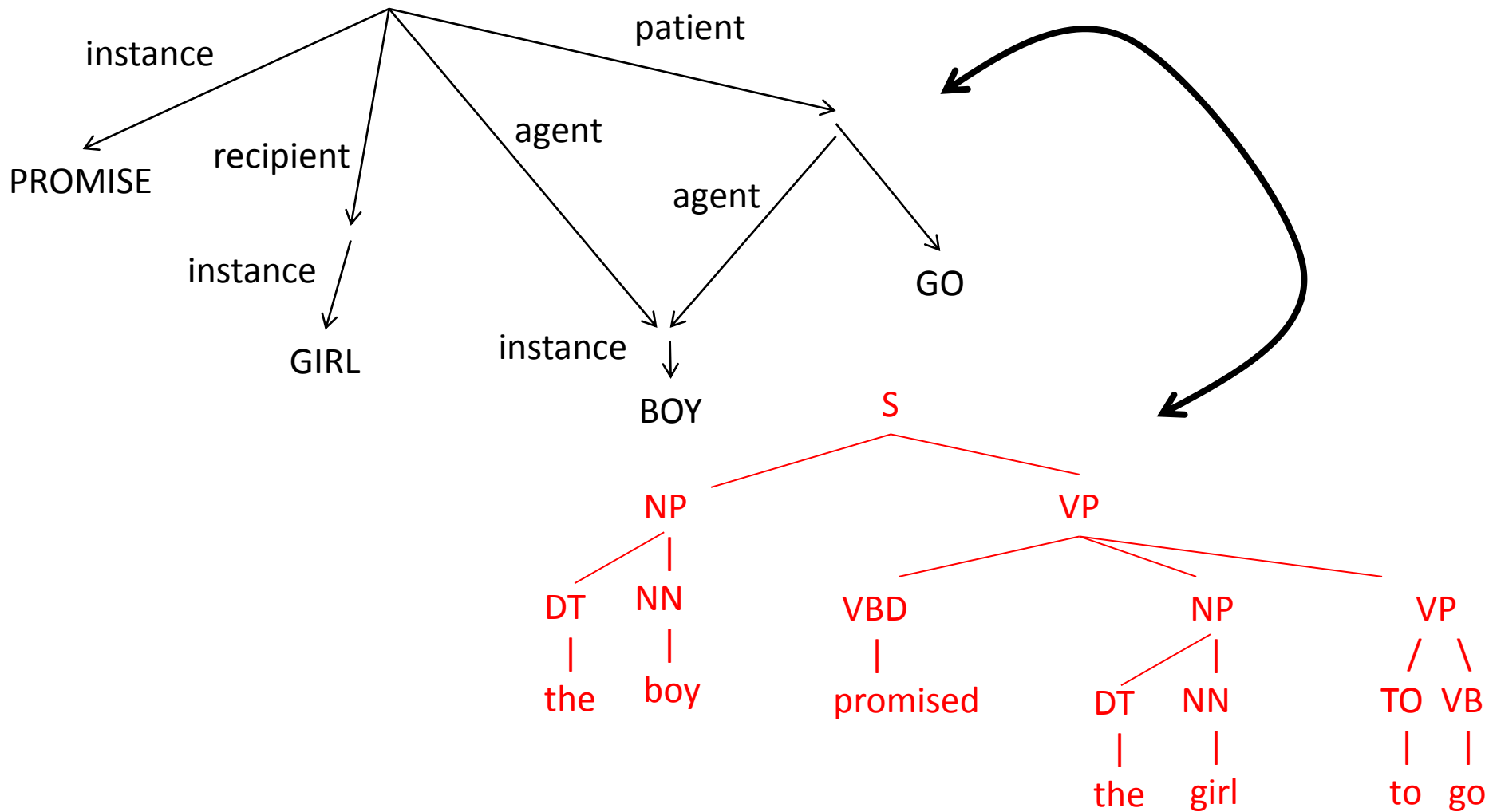
Bottom-Up DAG-to-Tree Transduction (Kamimura & Slutski 82)



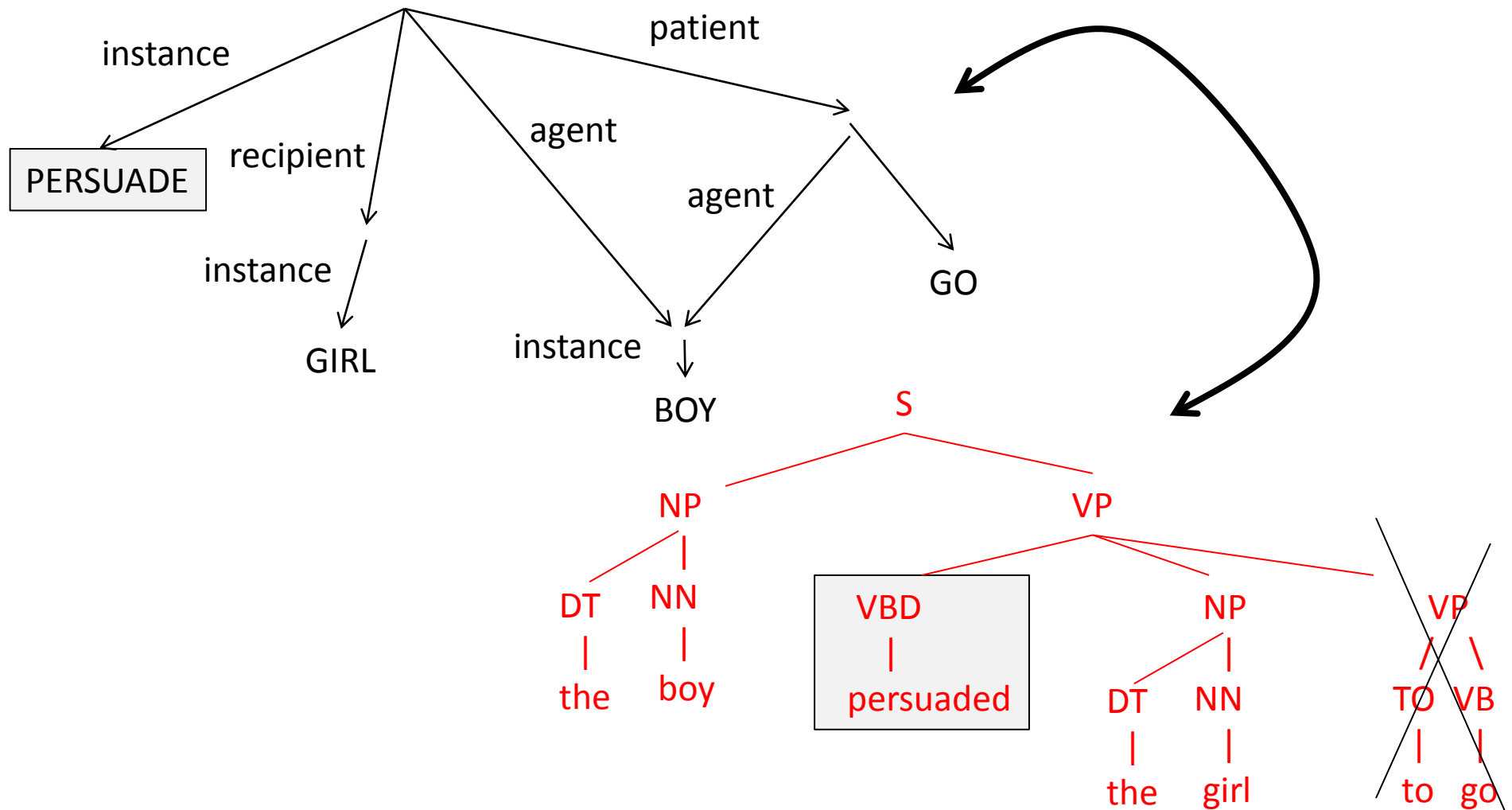
Bottom-Up DAG-to-Tree Transduction (Kamimura & Slutski 82)



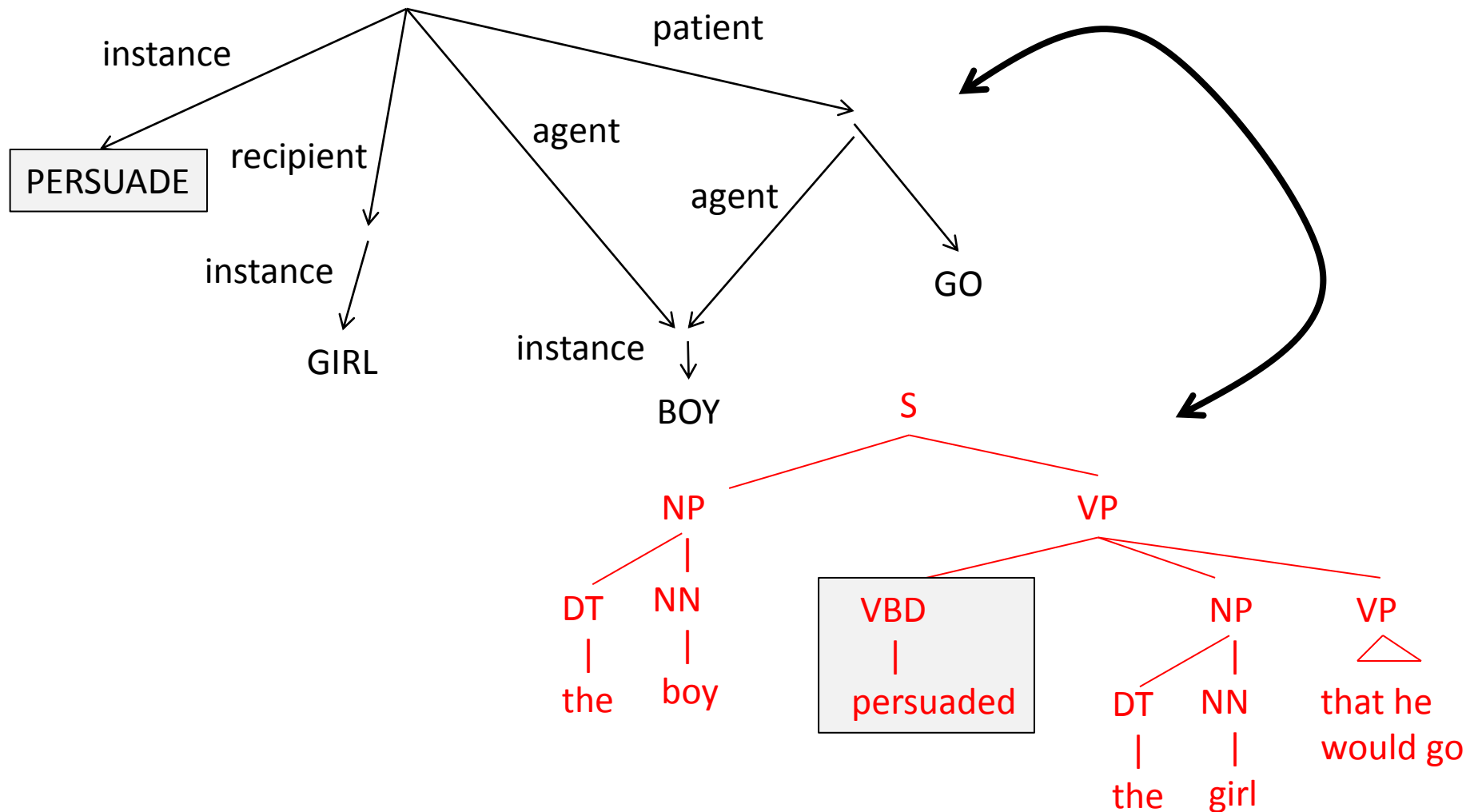
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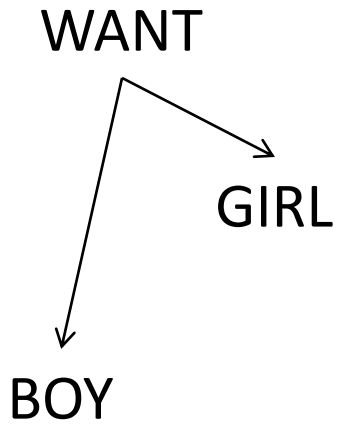


A Semantic Microworld

Node Labels: {WANT, BELIEVE, BOY, GIRL, 0}

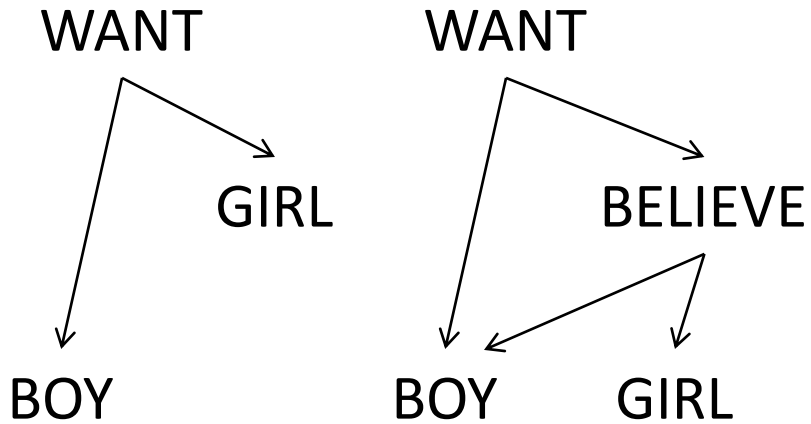
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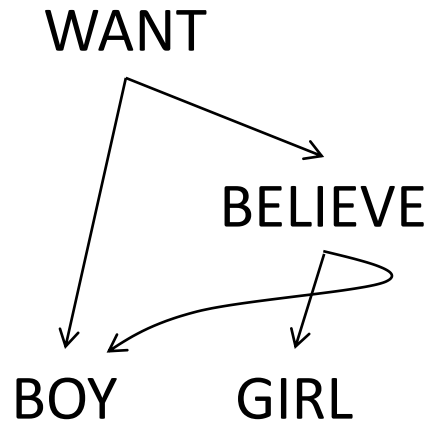
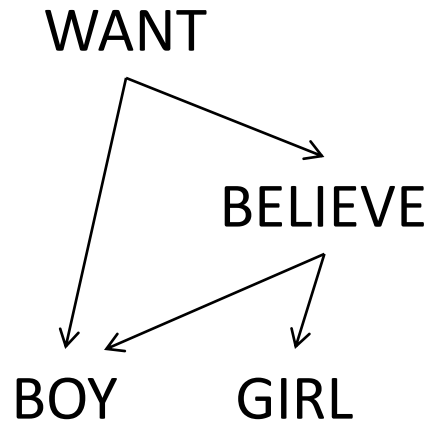
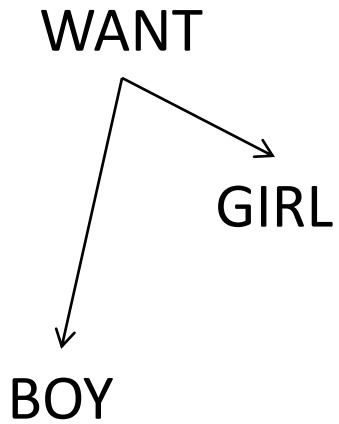
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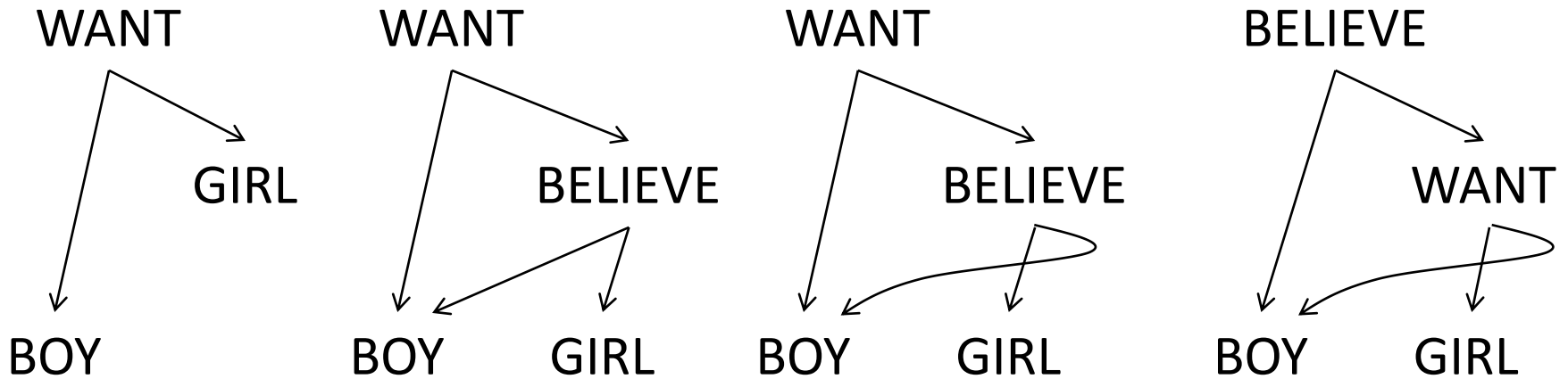
A Semantic Microworld

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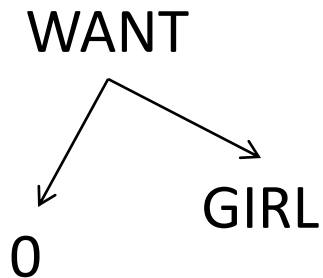
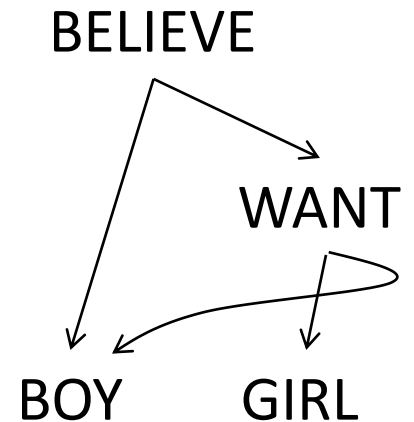
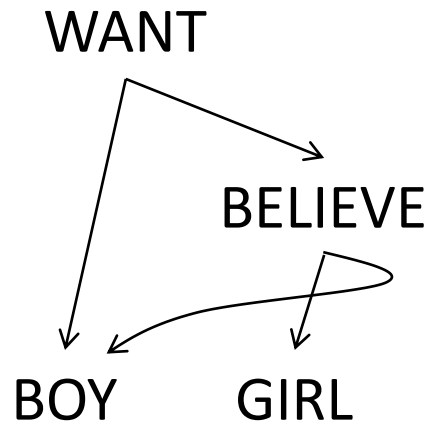
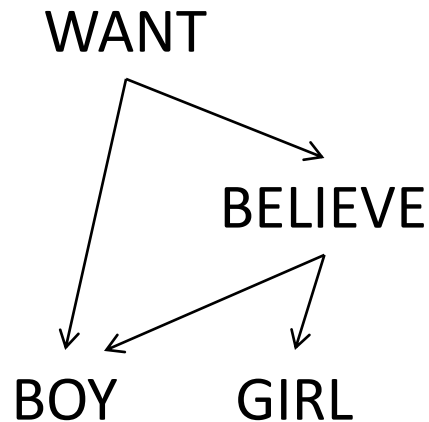
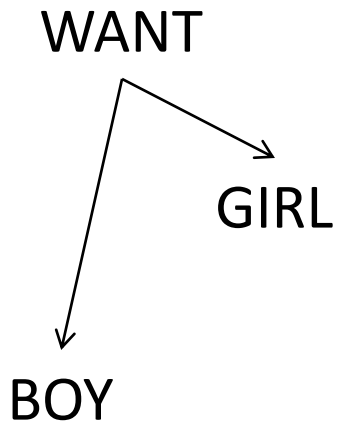
A Semantic Microworld

Node Labels: {WANT, BELIEVE, BOY, GIRL, 0}



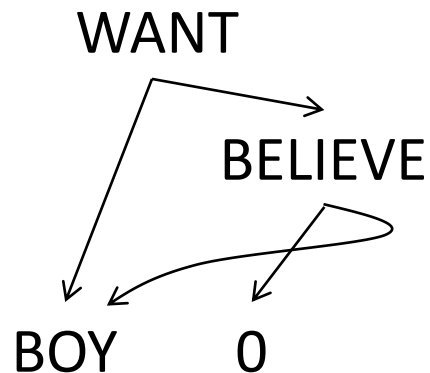
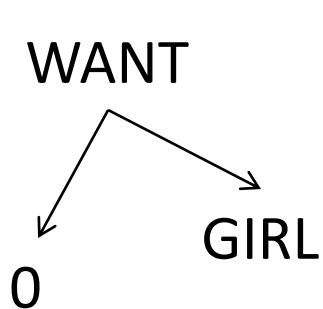
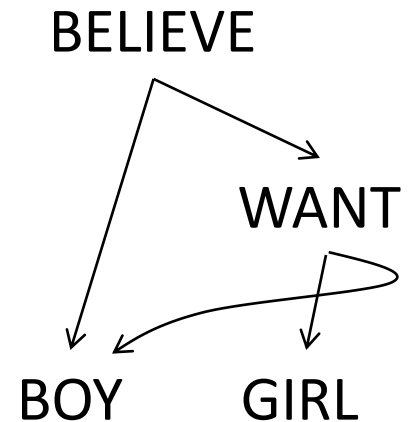
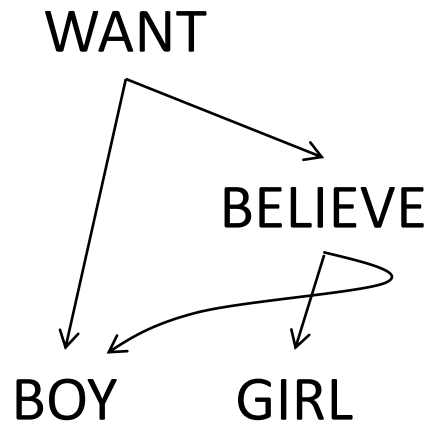
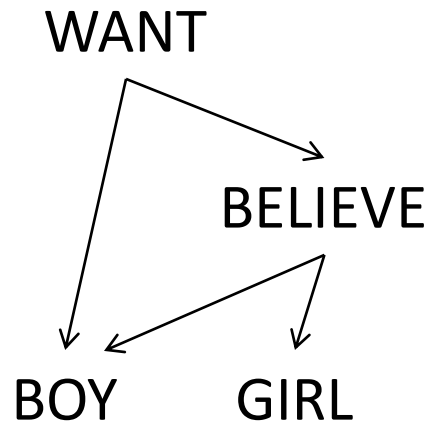
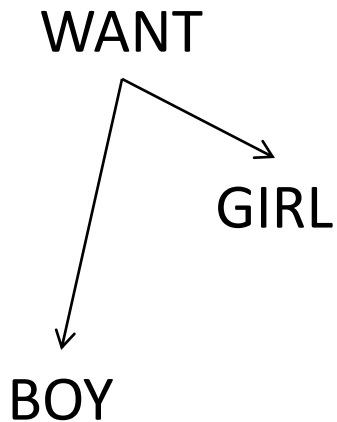
A Semantic Microworld

Node Labels: {WANT, BELIEVE, BOY, GIRL, 0}



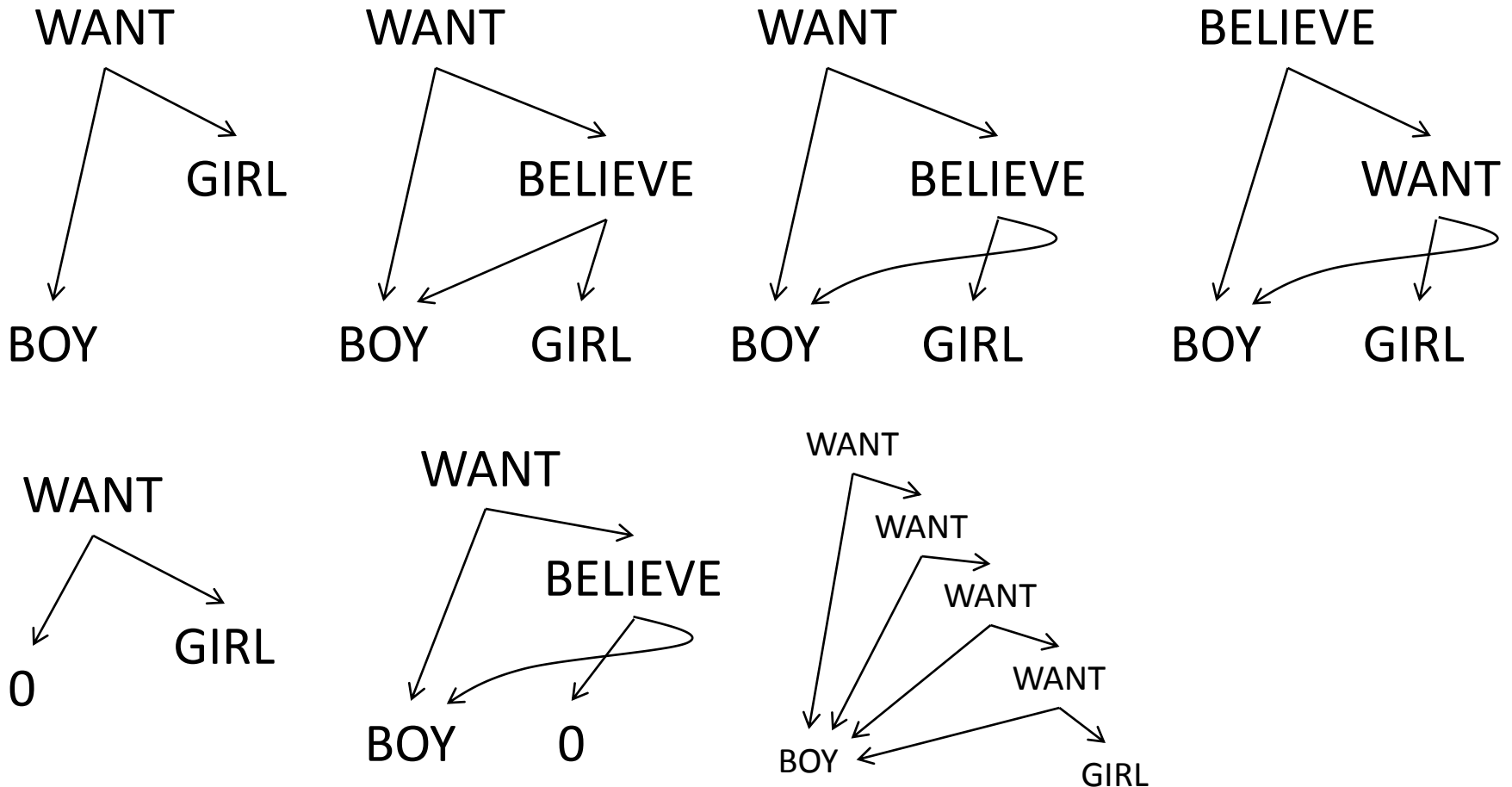
A Semantic Microworld

Node Labels: {WANT, BELIEVE, BOY, GIRL, 0}



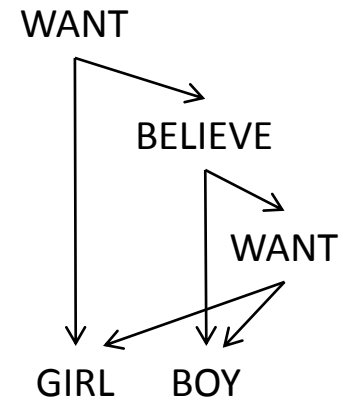
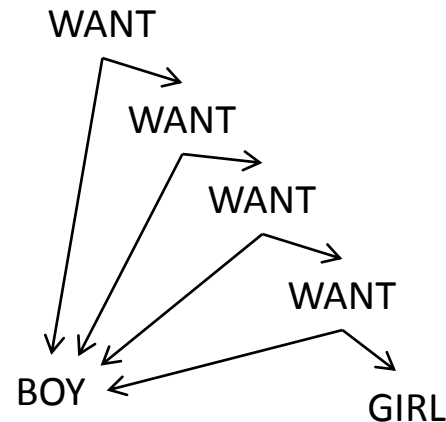
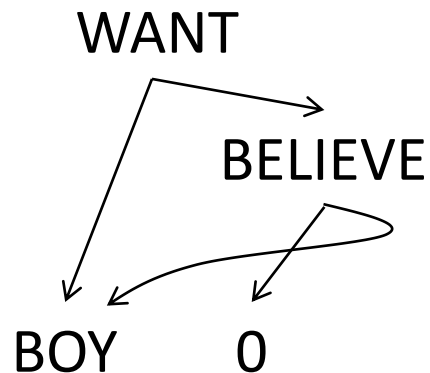
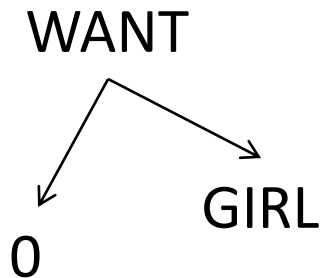
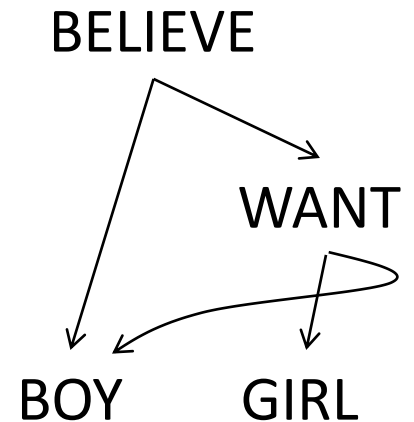
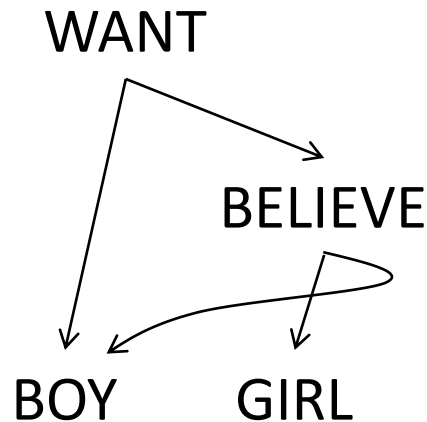
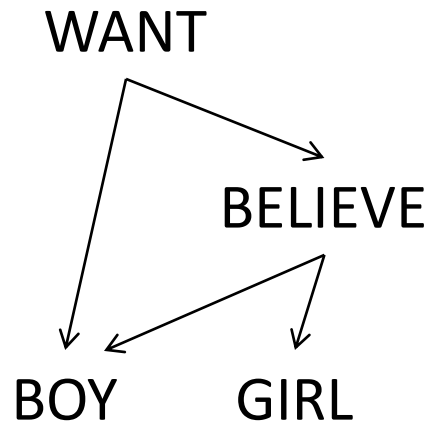
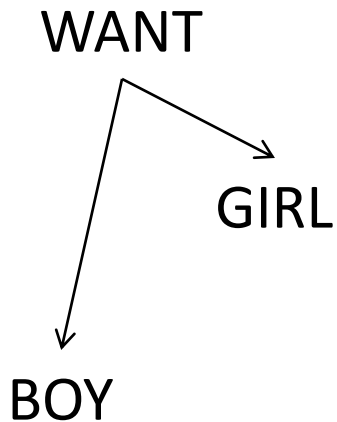
A Semantic Microworld

Node Labels: {WANT, BELIEVE, BOY, GIRL, 0}



A Semantic Microworld

Node Labels: {WANT, BELIEVE, BOY, GIRL, 0}



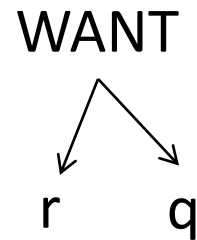
DAG Acceptor

q
q \rightarrow WANT(r q)
q \rightarrow BELIEVE(r q)
q \rightarrow r | 0
r \rightarrow BOY | GIRL | 0
[r r] \rightarrow r
[r r r] \rightarrow r

q

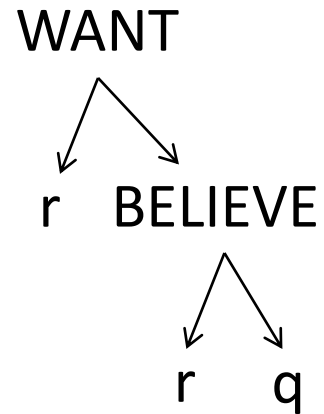
DAG Acceptor

q
q \rightarrow WANT(r q)
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q \rightarrow r | 0
r \rightarrow BOY | GIRL | 0
[r r] \rightarrow r
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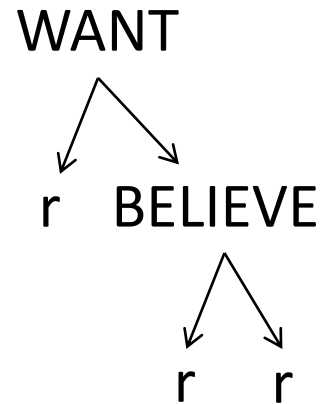
DAG Acceptor

q
q \rightarrow WANT(r q)
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q \rightarrow r | 0
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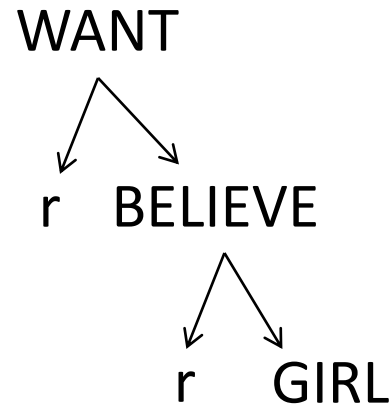
DAG Acceptor

q
q \rightarrow WANT(r q)
q \rightarrow BELIEVE(r q)
q \rightarrow r | 0
r \rightarrow BOY | GIRL | 0
[r r] \rightarrow r
[r r r] \rightarrow r



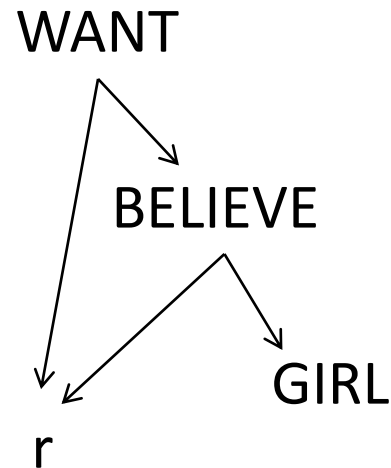
DAG Acceptor

q
q \rightarrow WANT(r q)
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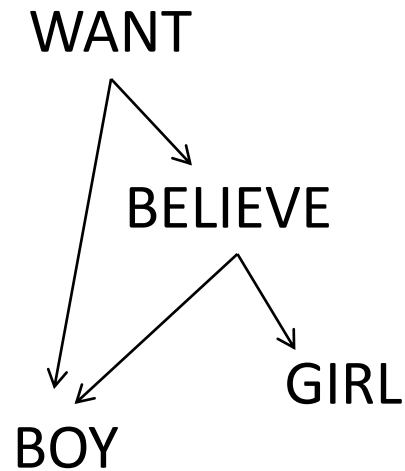
DAG Acceptor

q
q \rightarrow WANT(r q)
q \rightarrow BELIEVE(r q)
q \rightarrow r | 0
r \rightarrow BOY | GIRL | 0
[r r] \rightarrow r
[r r r] \rightarrow r



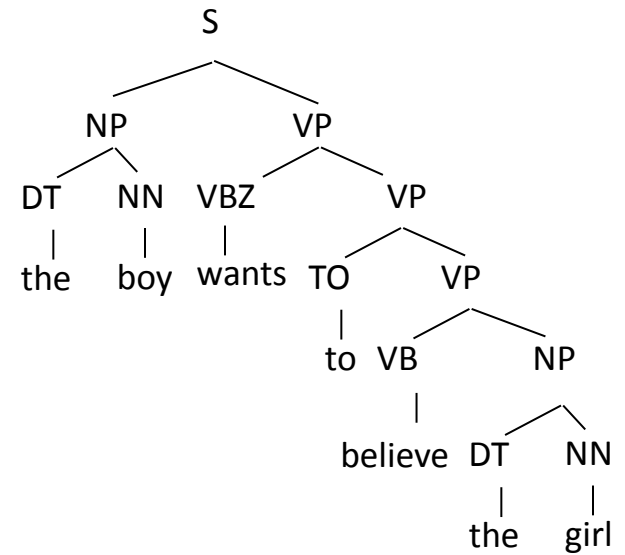
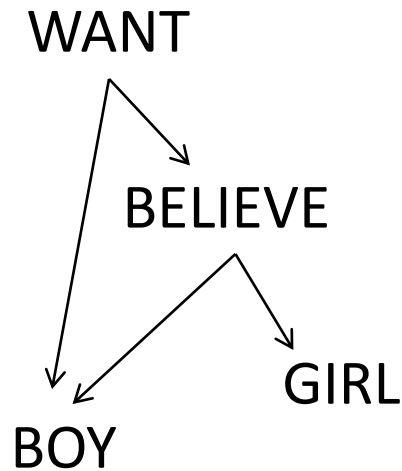
DAG Acceptor

q
q \rightarrow WANT(r q)
q \rightarrow BELIEVE(r q)
q \rightarrow r | 0
r \rightarrow BOY | GIRL | 0
[r r] \rightarrow r
[r r r] \rightarrow r



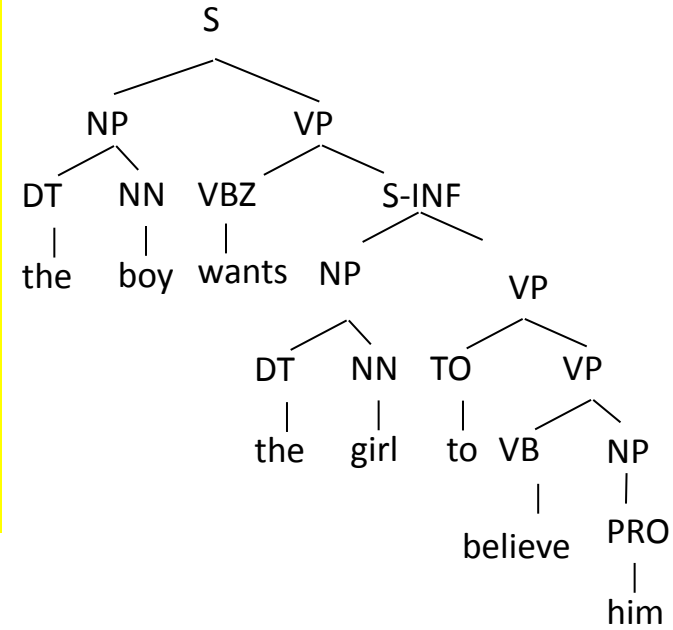
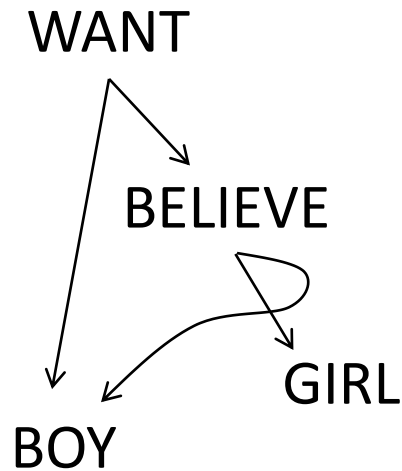
DAG Transducer

q
q → WANT(r q)
q → BELIEVE(r q)
q → r | 0
r → BOY | GIRL | 0
[r r] → r
[r r r] → r



DAG Transducer

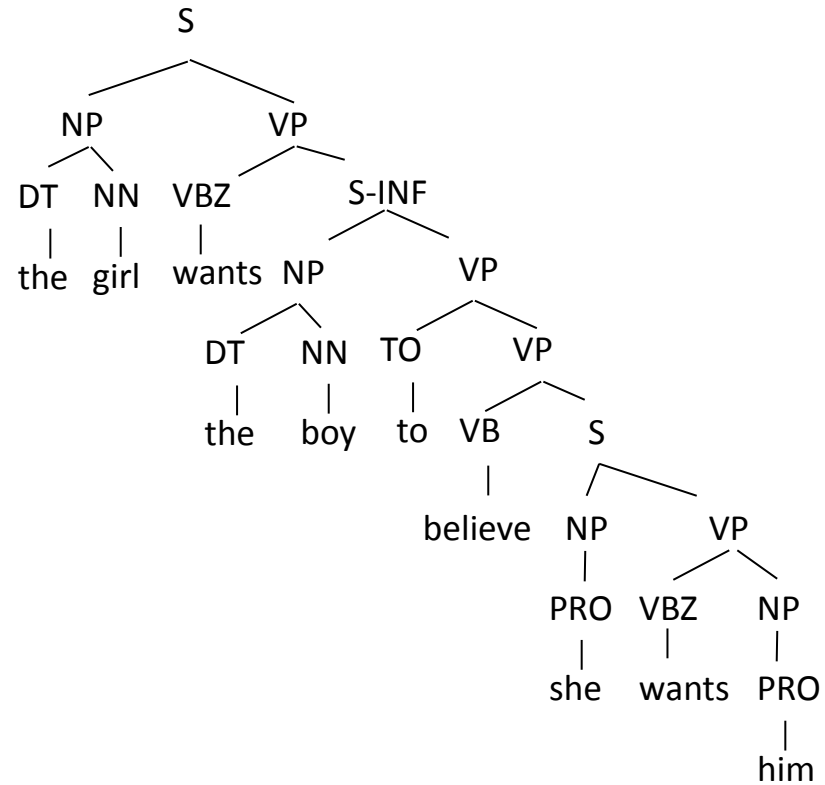
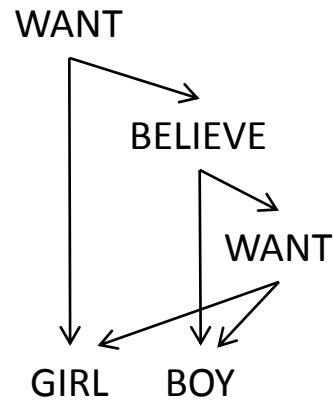
q
q → WANT(r q)
q → BELIEVE(r q)
q → r | 0
r → BOY | GIRL | 0
[r r] → r
[r r r] → r



or, "the boy wants to be
believed by the girl"

DAG Transducer

q
q → WANT(r q)
q → BELIEVE(r q)
q → r | 0
r → BOY | GIRL | 0
[r r] → r
[r r r] → r



DAG Transducer

$q_s.WANT(x, y) \rightarrow S(q_{nomb}.x, wants, q_{accg}.y)$
 $q_s.WANT(x, y) \rightarrow S(q_{nomb}.x, wants, q_{reflb}.y)$
 $q_s.WANT(x, y) \rightarrow S(q_{nomg}.x, wants, q_{acccb}.y)$
 $q_s.WANT(x, y) \rightarrow S(q_{nomg}.x, wants, q_{reflg}.y)$
 $q_s.WANT(x, y) \rightarrow S(q_{nomb}.x, wants, q_{infb}.y)$
 $q_s.WANT(x, y) \rightarrow S(q_{nomg}.x, wants, q_{infg}.y)$
 $q_{infb}.WANT(x, y) \rightarrow INF(q_{zerob}.x, to\ want, q_{accg}.y)$
 $q_{infb}.WANT(x, y) \rightarrow INF(q_{zerob}.x, to\ want, q_{reflb}.y)$
 $q_{infb}.WANT(x, y) \rightarrow INF(q_{accg}.x, to\ want, q_{acccb}.y)$
 $q_{infb}.WANT(x, y) \rightarrow INF(q_{accg}.x, to\ want, q_{reflg}.y)$
 $q_{infb}.WANT(x, y) \rightarrow INF(q_{zerob}.x, to\ want, q_{infb}.y)$
 $q_{infb}.WANT(x, y) \rightarrow INF(q_{accg}.x, to\ want, q_{infg}.y)$
 $q_{infg}.WANT(x, y) \rightarrow INF(q_{zerog}.x, to\ want, q_{acccb}.y)$
 $q_{infg}.WANT(x, y) \rightarrow INF(q_{zerog}.x, to\ want, q_{reflg}.y)$
 $q_{infg}.WANT(x, y) \rightarrow INF(q_{acccb}.x, to\ want, q_{accg}.y)$
 $q_{infg}.WANT(x, y) \rightarrow INF(q_{acccb}.x, to\ want, q_{reflg}.y)$
 $q_{infg}.WANT(x, y) \rightarrow INF(q_{zerog}.x, to\ want, q_{infg}.y)$
 $q_{infg}.WANT(x, y) \rightarrow INF(q_{acccb}.x, to\ want, q_{infb}.y)$
 $q_s.BELIEVE(x, y) \rightarrow S(q_{nomb}.x, believes, q_{accg}.y)$
 $q_s.BELIEVE(x, y) \rightarrow S(q_{nomb}.x, believes, q_{reflb}.y)$
 $q_s.BELIEVE(x, y) \rightarrow S(q_{nomg}.x, believes, q_{acccb}.y)$
 $q_s.BELIEVE(x, y) \rightarrow S(q_{nomg}.x, believes, q_{reflg}.y)$
 $q_s.BELIEVE(x, y) \rightarrow S(q_{nomb}.x, believes\ that, q_s.y)$

$q_s.BELIEVE(x, y) \rightarrow S(q_{nomg}.x, believes, q_s.y)$
 $q_{infb}.BELIEVE(x, y) \rightarrow INF(q_{zerob}.x, to\ believe, q_{accg}.y)$
 $q_{infb}.BELIEVE(x, y) \rightarrow INF(q_{zerob}.x, to\ believe, q_{reflb}.y)$
 $q_{infb}.BELIEVE(x, y) \rightarrow INF(q_{accg}.x, to\ believe, q_{acccb}.y)$
 $q_{infb}.BELIEVE(x, y) \rightarrow INF(q_{accg}.x, to\ believe, q_{reflg}.y)$
 $q_{infb}.BELIEVE(x, y) \rightarrow INF(q_{zerob}.x, to\ believe, thatq_s.y)$
 $q_{infb}.BELIEVE(x, y) \rightarrow INF(q_{accg}.x, to\ believe, thatq_s.y)$
 $q_{infg}.BELIEVE(x, y) \rightarrow INF(q_{zerog}.x, to\ believe, q_{acccb}.y)$
 $q_{infg}.BELIEVE(x, y) \rightarrow INF(q_{zerog}.x, to\ believe, q_{reflg}.y)$
 $q_{infg}.BELIEVE(x, y) \rightarrow INF(q_{acccb}.x, to\ believe, q_{accg}.y)$
 $q_{infg}.BELIEVE(x, y) \rightarrow INF(q_{acccb}.x, to\ believe, q_{reflb}.y)$
 $q_{infg}.BELIEVE(x, y) \rightarrow INF(q_{zerog}.x, to\ believe\ that, q_s.y)$
 $q_{infg}.BELIEVE(x, y) \rightarrow INF(q_{acccb}.x, to\ believe\ that, q_s.y)$
 $q_{nomb}.BOY \rightarrow NP(the\ boy)$
 $q_{acccb}.BOY \rightarrow NP(the\ boy)$
 $[q_{nomb}, q_{nomb}].BOY \rightarrow NP(the\ boy), NP(he)$
 $[q_{nomb}, q_{acccb}].BOY \rightarrow NP(the\ boy), NP(him)$
 $[q_{nomb}, q_{reflb}].BOY \rightarrow NP(the\ boy), NP(himself)$
 $[q_{nomb}, q_{zerob}].BOY \rightarrow NP(the\ boy), 0$
 $q_{nomg}.GIRL \rightarrow NP(the\ girl)$
 $q_{accg}.GIRL \rightarrow NP(the\ girl)$
 $[q_{nomg}, q_{nomg}].GIRL \rightarrow NP(the\ girl), NP(she)$
 $[q_{nomg}, q_{accg}].GIRL \rightarrow NP(the\ girl), NP(her)$
 $[q_{nomg}, q_{reflg}].GIRL \rightarrow NP(the\ girl), NP(herself)$
 $[q_{nomg}, q_{zerog}].GIRL \rightarrow NP(the\ girl), 0$

Two DAG Transducers

foreign
text

→ -1

```
(s / say
:agent (s2 / source :mod (r / rebel))
:patient (a / and
:op1 (c / close-on
:agent (f / force :mod (g / government))
:patient (o / outpost :mod (r2 / rebel))
:temporal-locating (t / thursday))
:op2 (s / shower
:agent f
:patient (c2 / city
:mod (m / mountain)
:mod (w / west)
:name "Zintan")
:instrument (m2 / missile))
:op3 (a2 / attack
:agent f
:patient (i / insurgent
:agent-of (h / hole-up
:pp-near (b / border
:poss (c3 / country
:name "Tunisia"]
```

→

Government forces closed on rebel outposts on Thursday, showering the western mountain city of Zintan with missiles and attacking insurgents holed up near the Tunisian border, according to rebel sources.

DAG Transducer
applies in reverse

General-Purpose Algorithms

	String Automata Algorithms	Tree Automata Algorithms	Graph Automata Algorithms?
N-best paths through an WFSA (Viterbi, 1967; Eppstein, 1998)	... trees in a weighted forest (Jiménez & Marzal, 2000; Huang & Chiang, 2005)	
EM training	Forward-backward EM (Baum/Welch, 1971; Eisner 2003)	Tree transducer EM training (Graehl & Knight, 2004)	
Determinization...	... of weighted string acceptors (Mohri, 1997)	... of weighted tree acceptors (Borchardt & Vogler, 2003; May & Knight, 2005)	
Intersection	WFSA intersection	Tree acceptor intersection	
Applying transducers	string \rightarrow WFST \rightarrow WFSA	tree \rightarrow TT \rightarrow weighted tree acceptor	
Transducer composition	WFST composition (Pereira & Riley, 1996)	Many tree transducers not closed under composition (Maletti et al 09)	
General tools	FSM, Carmel, OpenFST	Tiburon (May & Knight 10)	

Final Thought

Used
in SMT

2020

2010

2000

1990

Finite-state Transducer

Top-down tree transducer

Graph automata

1960

1970

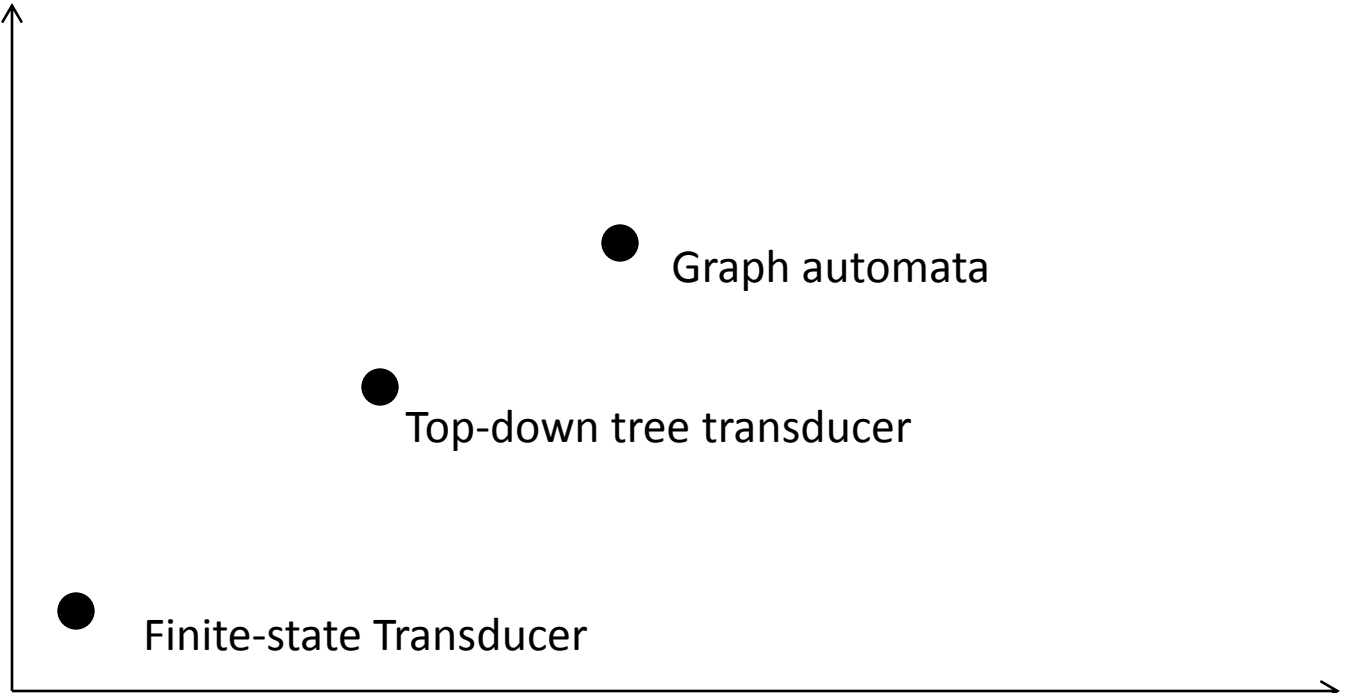
1980

1990

2000

2010

Automata Framework



Final Thought

Used
in SMT

2020

2010

2000

1990

Finite-state Transducer

Top-down tree transducer

Graph automata

???

1960

1970

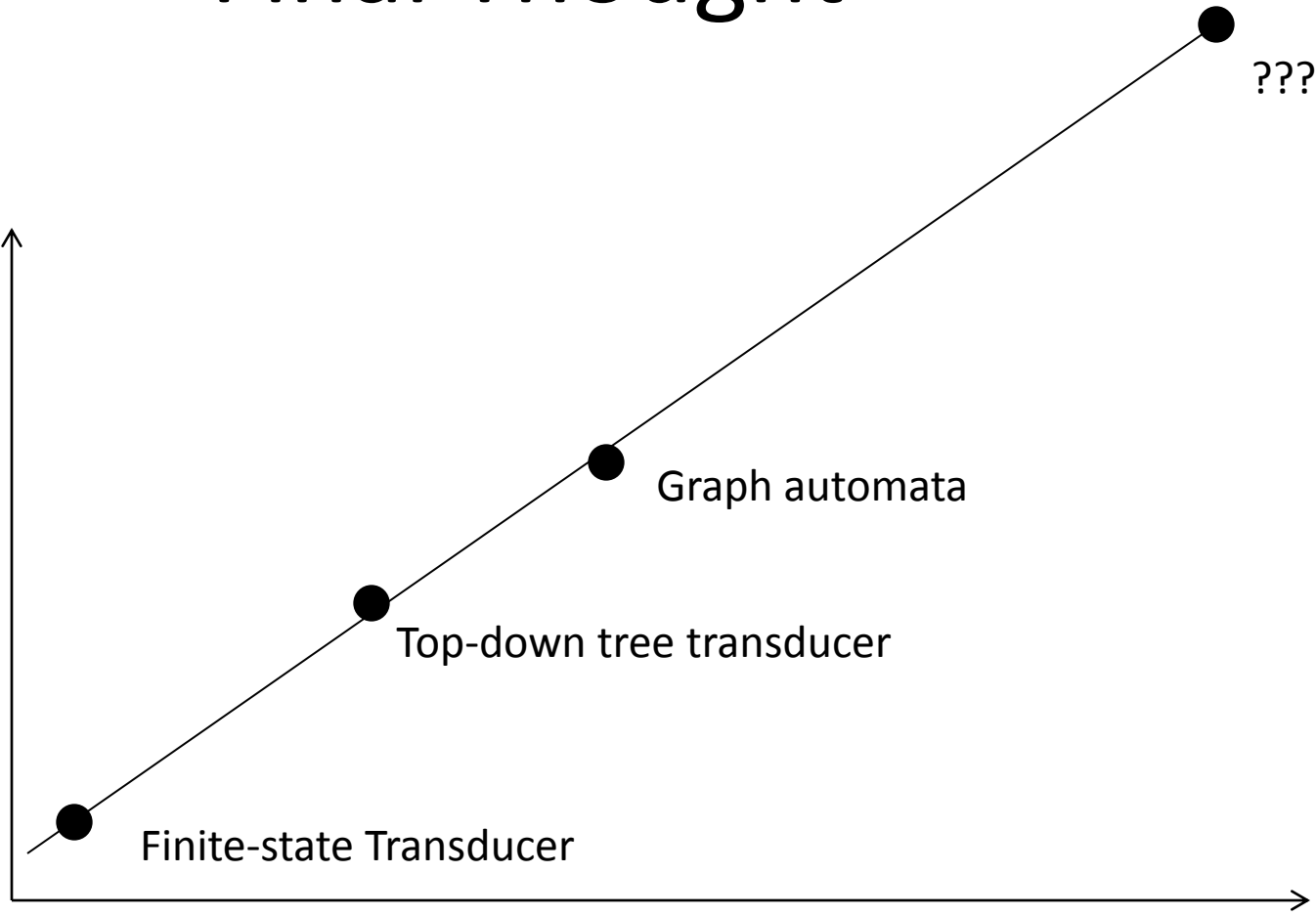
1980

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2010

Automata Framework

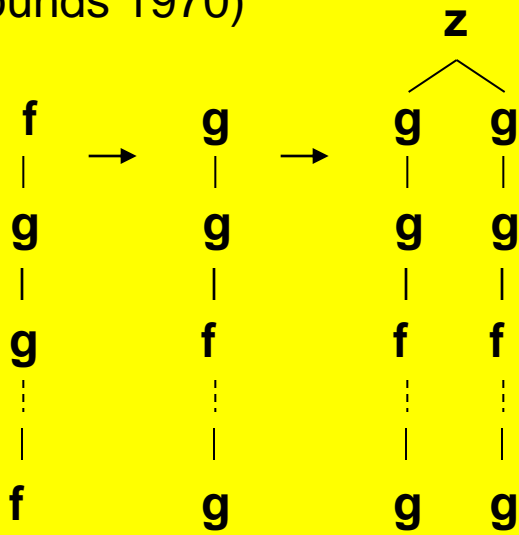


end

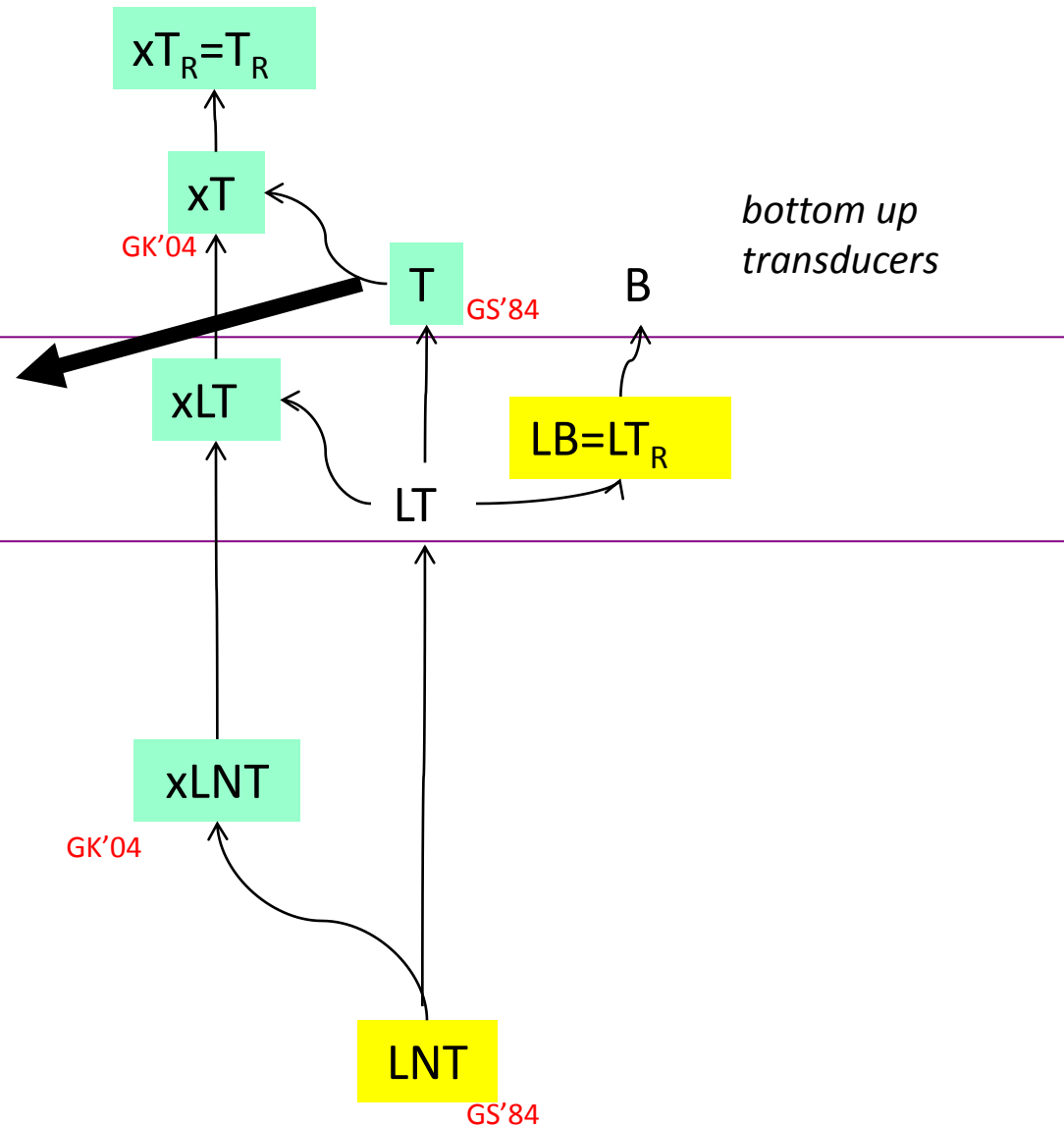
Expressive enough for local rotation

Closed under composition

T is not closed under composition
(Rounds 1970)



T cannot *Nprocess-then-copy*



Expressive enough for local rotation

Closed under composition

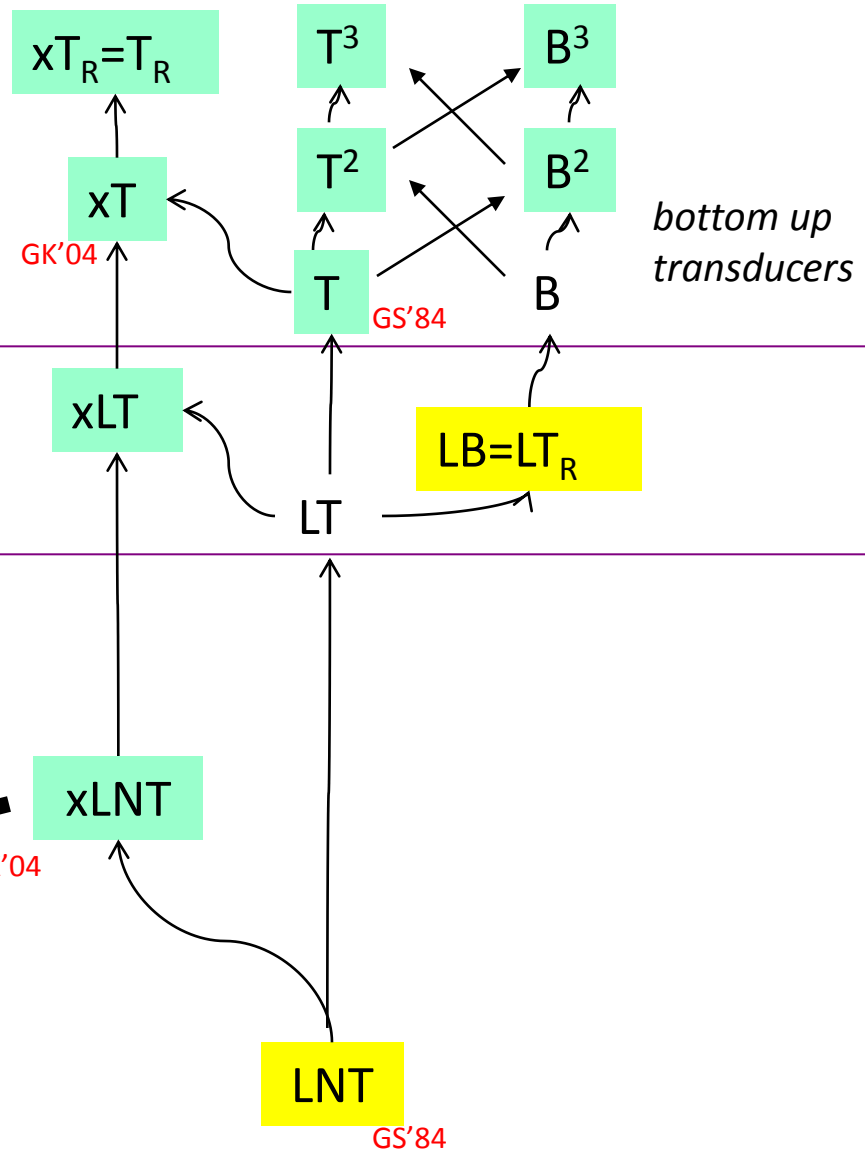
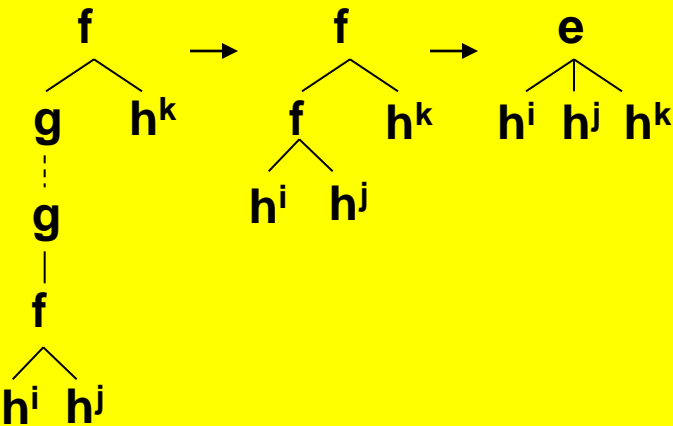
copying

non-copying

deleting

non-deleting

xLNT is **not closed** under composition



Expressive enough for local rotation

Closed under composition

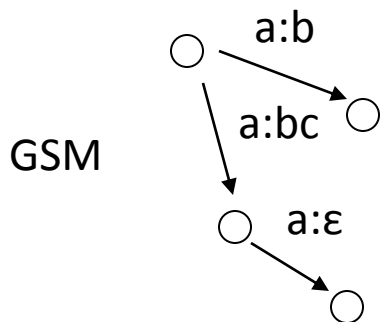
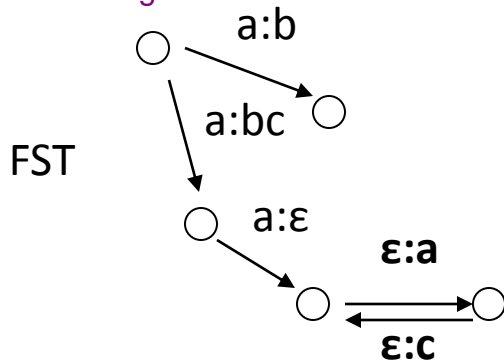
Generalizes FST

copying

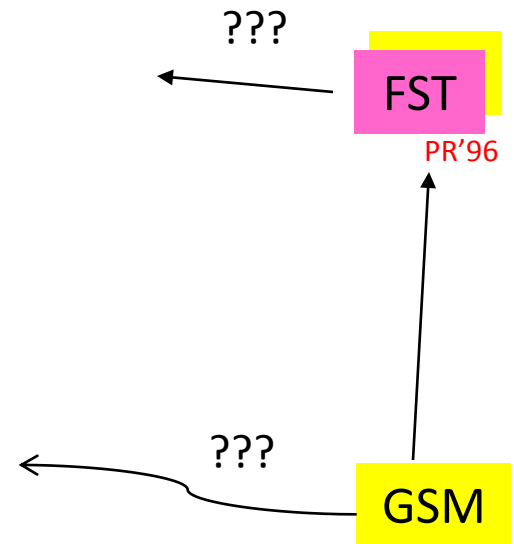
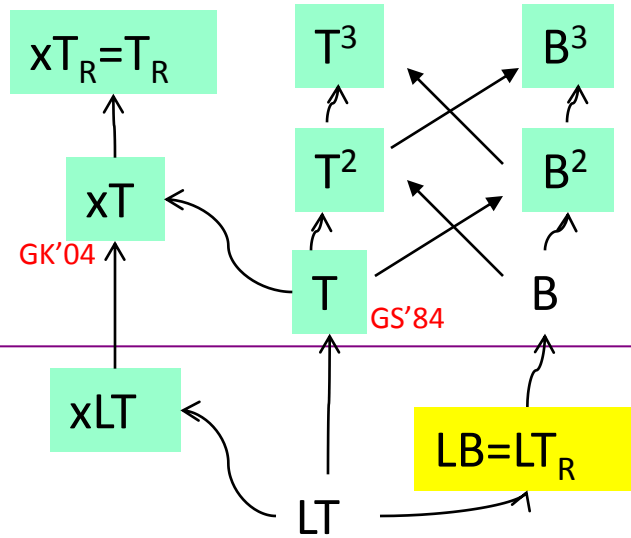
non-copying

deleting

non-deleting



every transition must eat up one input symbol.



Expressive enough for local rotation

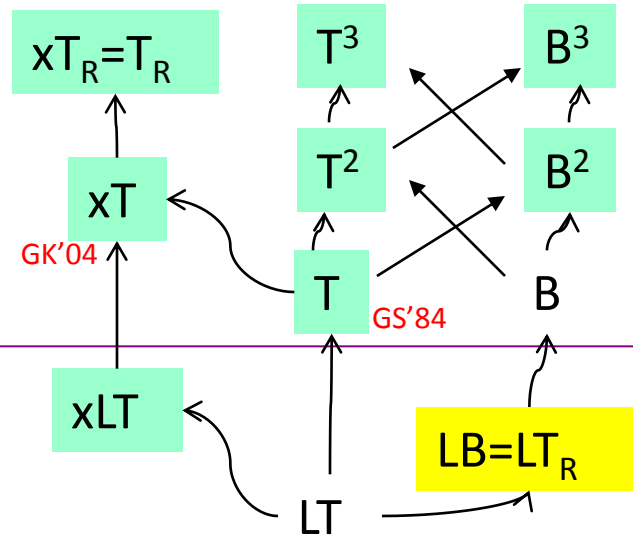
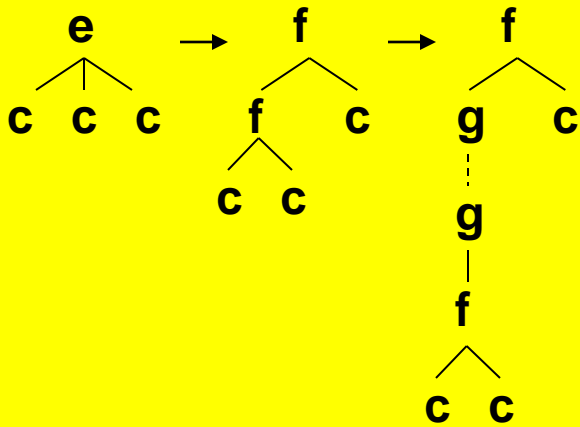
Closed under composition

Generalizes FST

copying

non-copying

LNT w/ input-epsilon is **not closed** under composition



GK'04

xLNT
(input-e,
output-e)

xLNT
(input-e)

LNT
(input-e,
output-e)

FST
PR'96

xLNT
(e-free)
GHKM'04

LNT
(input-e)

LNT
(output-e)
GS'84

LNT
(e-free)

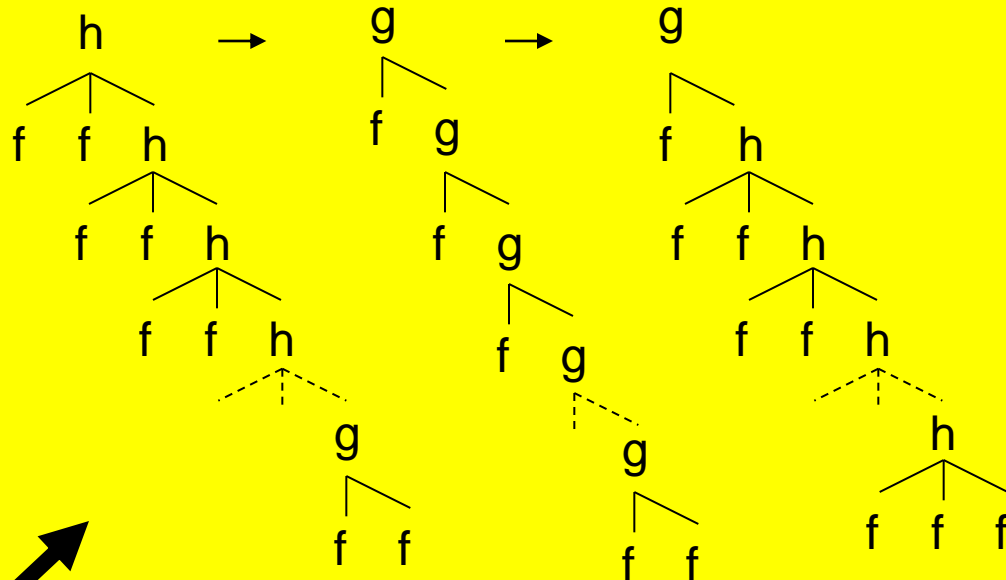
GSM

Expressive enough for local rotation

Closed under composition

Generalizes FST

xLNT (e-free) is **not closed** under composition



xLNT
(output-e)

xLNT
(e-free)

GHKM'04

LNT
(input-e)

LNT
(e-free)

output-e)

LNT
(output-e)

GS'84

FST

PR'96

GSM

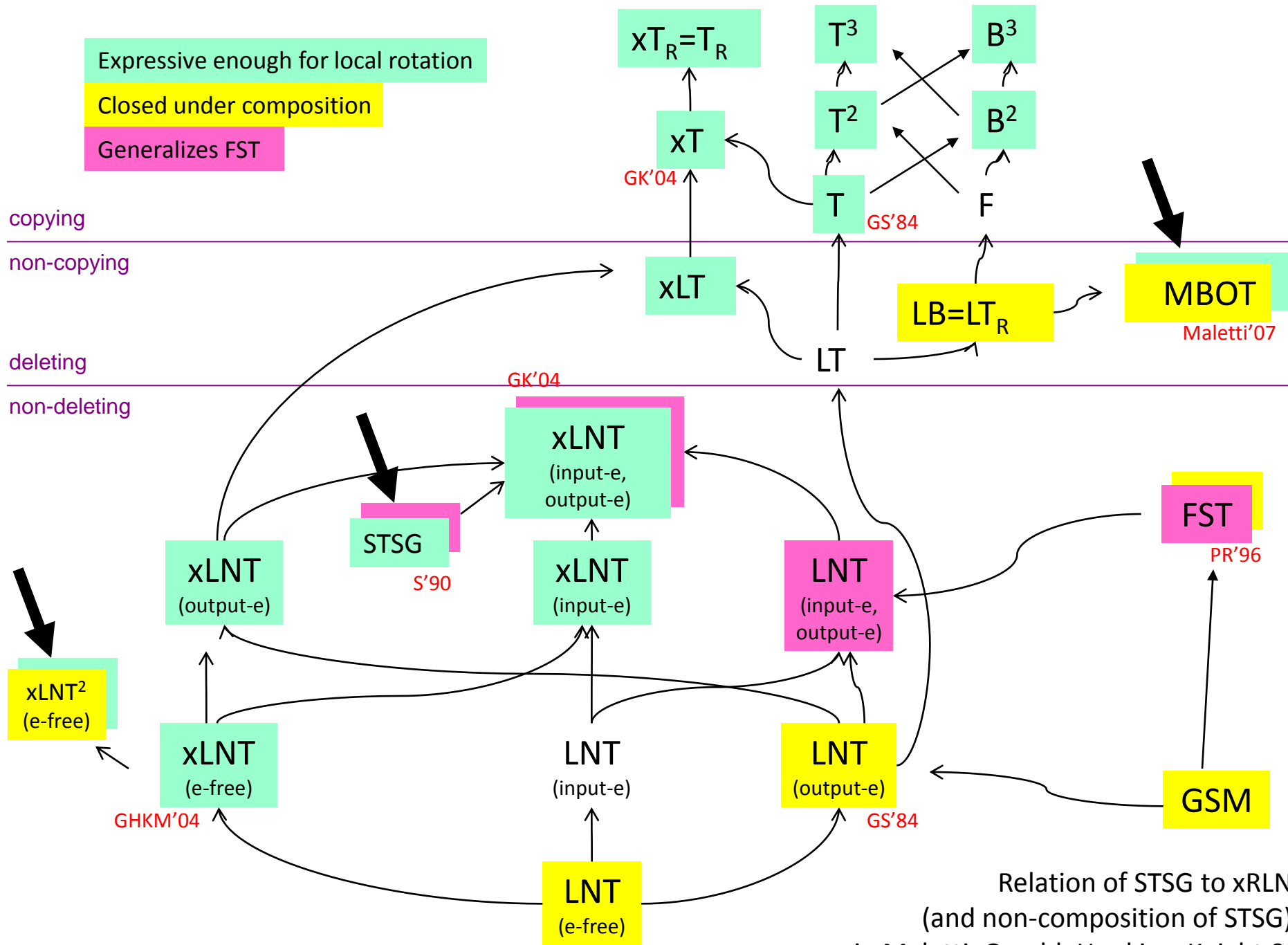
Composition theorems for e-free in Arnold & Dauchet 82

copying

non-copying

deleting

non-deleting



Relation of STSG to xRLN (and non-composition of STSG) in Maletti, Graehl, Hopkins, Knight 09