

8.8 Securing Wireless LANs

Security is a particularly important concern in wireless networks, where radio waves carrying frames can propagate far beyond the building containing the wireless base station and hosts. In this section we present a brief introduction to wireless security. For a more in-depth treatment, see the highly readable book by Edney and Arbaugh [Edney 2003].

The issue of security in 802.11 has attracted considerable attention in both technical circles and in the media. While there has been considerable discussion, there has been little debate—there seems to be universal agreement that the original 802.11 specification contains a number of serious security flaws. Indeed, public domain software can now be downloaded that exploits these holes, making those who use the vanilla 802.11 security mechanisms as open to security attacks as users who use no security features at all.

In the following section, we discuss the security mechanisms initially standardized in the 802.11 specification, known collectively as **Wired Equivalent Privacy (WEP)**. As the name suggests, WEP is meant to provide a level of security similar to that found in wired networks. We'll then discuss a few of the security holes in WEP and discuss the 802.11i standard, a fundamentally more secure version of 802.11 adopted in 2004.

8.8.1 Wired Equivalent Privacy (WEP)

The IEEE 802.11 WEP protocol was designed in 1999 to provide authentication and data encryption between a host and a wireless access point (that is, base station) using a symmetric shared key approach. WEP does not specify a key management algorithm, so it is assumed that the host and wireless access point have somehow agreed on the key via an out-of-band method. Authentication is carried out as follows:

1. A wireless host requests authentication by an access point.
2. The access point responds to the authentication request with a 128-byte nonce value.
3. The wireless host encrypts the nonce using the symmetric key that it shares with the access point.
4. The access point decrypts the host-encrypted nonce.

If the decrypted nonce matches the nonce value originally sent to the host, then the host is authenticated by the access point.

The WEP data encryption algorithm is illustrated in Figure 8.30. A secret 40-bit symmetric key, K_S , is assumed to be known by both a host and the access point. In addition, a 24-bit Initialization Vector (IV) is appended to the 40-bit key to create a 64-bit key that will be used to encrypt a single frame. The IV will change from one

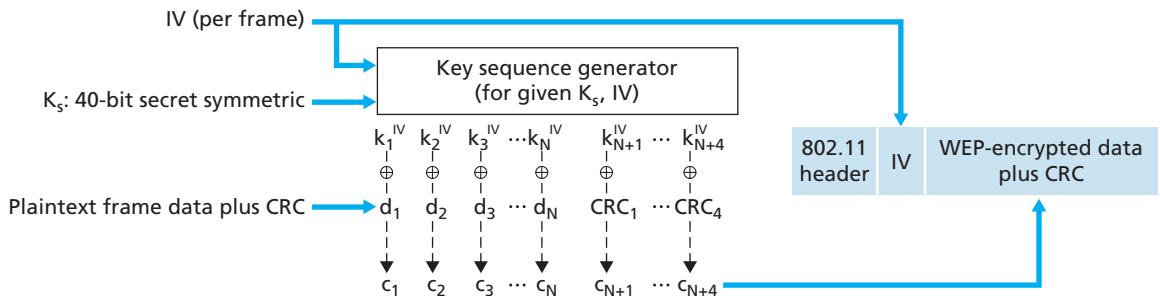


Figure 8.30 ♦ 802.11 WEP protocol

frame to another, and hence each frame will be encrypted with a different 64-bit key. Encryption is performed as follows. First a 4-byte CRC value (see Section 5.2) is computed for the data payload. The payload and the four CRC bytes are then encrypted using the RC4 stream cipher. We will not cover the details of RC4 here (see [Schneier 1995] and [Edney 2003] for details). For our purposes, it is enough to know that when presented with a key value (in this case, the 64-bit (K_s , IV) key), the RC4 algorithm produces a stream of key values, k_1^{IV} , k_2^{IV} , k_3^{IV} , ... that are used to encrypt the data and CRC value in a frame. For practical purposes, we can think of these operations being performed a byte at a time. Encryption is performed by XOR-ing the i th byte of data, d_i , with the i th key, k_i^{IV} , in the stream of key values generated by the (K_s , IV) pair to produce the i th byte of ciphertext, c_i :

$$c_i = d_i \oplus k_i^{IV}$$

The IV value changes from one frame to the next and is included in *plaintext* in the header of each WEP-encrypted 802.11 frame, as shown in Figure 8.30. The receiver takes the secret 40-bit symmetric key that it shares with the sender, appends the IV, and uses the resulting 64-bit key (which is identical to the key used by the sender to perform encryption) to decrypt the frame:

$$d_i = c_i \oplus k_i^{IV}$$

Proper use of the RC4 algorithm requires that the same 64-bit key value *never* be used more than once. Recall that the WEP key changes on a frame-by-frame basis. For a given K_s (which changes rarely, if ever), this means that there are only 2^{24} unique keys. If these keys are chosen randomly, we can show [Walker 2000; Edney 2003] that the probability of having chosen the same IV value (and hence used the same 64-bit key) is more than 99 percent after only 12,000 frames. With 1 Kbyte frame sizes and a data transmission rate of 11 Mbps, only a few seconds are