



DEPARTMENT OF PHILOSOPHY,
LINGUISTICS AND THEORY OF SCIENCE

ON THE GRAMMAR OF PROOF

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Abstract

Brief summary of research question, background, method, results...

Preface

Acknowledgements, etc.

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1935- Add section numbers to sections On the Grammar of Proof

Introduction

The central concern of this thesis is the syntax of mathematics, programming languages, and their respective mutual influence, as conceived and practiced by mathematicians and computer scientists. From one vantage point, the role of syntax in mathematics may be regarded as a 2nd order concern, a topic for discussion during a Fika, an artifact of ad hoc development by the working mathematician whose real goals are producing genuine mathematical knowledge. For the programmers and computer scientists, syntax may be regarded as a matter of taste, with friendly debates recurring regarding the use of semicolons, brackets, and white space. Yet, when viewed through the lens of the propositions-as-types paradigm, these discussions intersect in new and interesting ways. When one introduces a third paradigm through which to analyze the use of syntax in mathematics and programming, namely linguistics, I propose what some may regard as superficial detail, indeed becomes a central paradigm raising many interesting and important questions.

Beyond Computational Trinitarianism

The doctrine of computational trinitarianism holds that computation manifests itself in three forms: proofs of propositions, programs of a type, and mappings between structures. These three aspects give rise to three sects of worship: Logic, which gives primacy to proofs and propositions; Languages, which gives primacy to programs and types; Categories, which gives primacy to mappings and structures.[1]

We begin this discussion of the three relationships between three respective fields, mathematics, and logic. The trinity Figure 1 as they are aptly named, are related via both formal and informal methods. The propositions as types paradigm, for example, is a heuristic and yet it also has many examples of successful ideas translating between the domains. Alternatively, the interpretation of a Type Theory(TT) into a category theory is incredibly *formal*.

We hope this thesis will help clarify another possible dimension in this diagram, that of Linguistics. And while the different categories may resemble religions in their own right, with communities convinced that they have a canonical perspective. Questioning the holy trinity is an act of a heresy, and it is the goal of this thesis to be a bit heretical by including a much less well understood perspective, name Linguistics, which may provide additional challenges and insights into the trinity.

We begin by showing how the trinity give rise to formal semantic interpretations of natural language. Semantics is just linguistic phenomenon worth investigating in these domains, and could be replaced by other linguistic paradigms. This work is alternatively concerned with syntax.

Additionally, we can ask how do the trinity embed into natural language. These are the most *informal* arrows of the “holy tetrahedron”, or at least one reading of

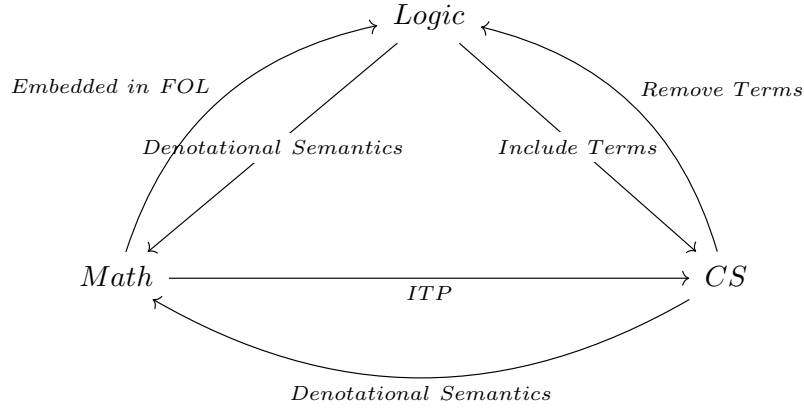


Figure 1: The Holy Trinity

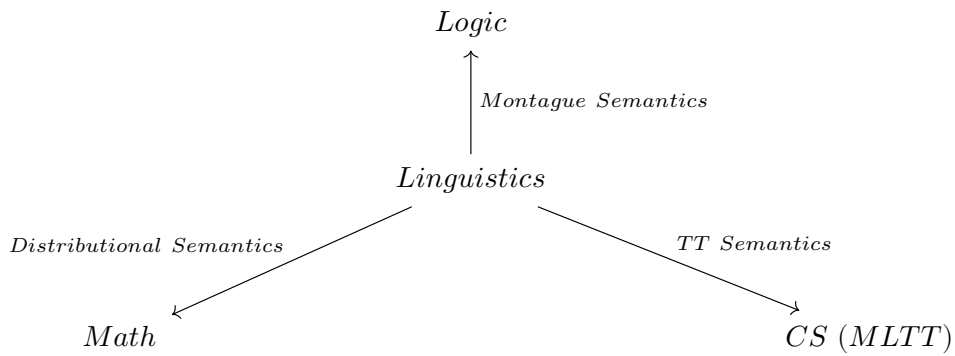


Figure 2: Formal Semantics

it. One can analyze mathematics using linguistic methods, or try to give a natural language justification of Intuitionistic Type Theory using Martin-Löf's meaning explanations.

In this work, we will see that there are multiple GF grammars which model some subset of each member of the trinity. Studying these grammars, and asking how they can be used in applications for mathematicians, logicians, and computer scientists, is an important practical and philosophical question.

Therefore, we hope this attempt at giving the language of mathematics, in particular how propositions and proofs are expressed and thought about in that language, a stronger foundation.

Example

To get a feel for this syntactic paradigm, let us look at a basic mathematical example: that of a group homomorphism, as expressed in a variety of sources.

Definition 1 *In mathematics, given two groups, $(G, *)$ and (H, \cdot) , a group homomorphism from $(G, *)$ to (H, \cdot) is a function $h : G \rightarrow H$ such that for all u and v in G it holds that*

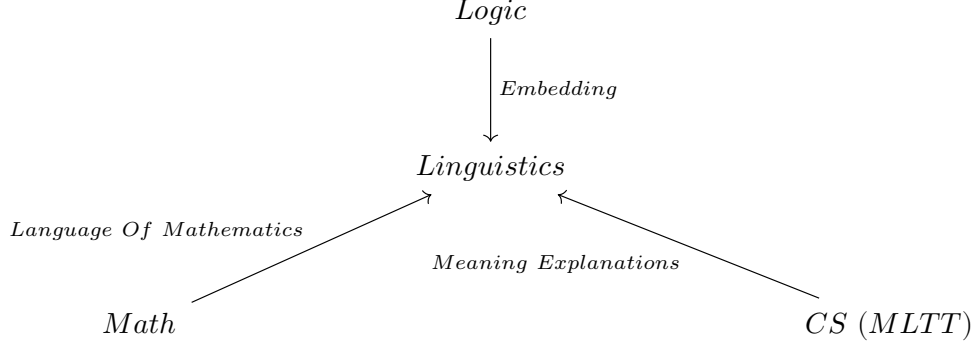


Figure 3: Interpretations of Natural Language

$$h(u * v) = h(u) \cdot h(v)$$

Definition 2 Let $G = (G, \cdot)$ and $G' = (G', *)$ be groups, and let $\phi : G \rightarrow G'$ be a map between them. We call ϕ a **homomorphism** if for every pair of elements $g, h \in G$, we have

$$\phi(g * h) = \phi(g) \cdot \phi(h)$$

Definition 3 Let G, H , be groups. A map $\phi : G \rightarrow H$ is called a group homomorphism if

$$\phi(xy) = \phi(x)\phi(y) \text{ for all } x, y \in G$$

(Note that xy on the left is formed using the group operation in G , whilst the product $\phi(x)\phi(y)$ is formed using the group operation H .)

Definition 4 Classically, a group is a monoid in which every element has an inverse (necessarily unique).

We inquire the reader to pay attention to nuance and difference in presentation that is normally ignored or taken for granted by the fluent mathematician.

If one want to distill the meaning of each of these presentations, there is a significant amount of subliminal interpretation happening very much analogous to our innate linguistic usage. The inverse and identity are discarded, even though they are necessary data when defining a group. The order of presentation of information is incostent, as well as the choice to use symbolic or natural language information. In (3), the group operation is used implicitly, and its clarification a side remark.

Details aside, these all mean the same thing-don't they? This thesis seeks to provide an abstract framework to determine whether two linguistically nuanced presentations mean the same thing via their syntactic transformations.

These syntactic transformations come in two flavors : parsing and linearization, and are natively handled by a Logical Framework (LF) for specifying grammars : Grammatical Framework (GF).

We now show yet another definition of a group homomorphism formalized in the Agda programming language:

[TODO: replace monoidhom with grouphom]

theres supposed to be agda here

```

monoidHom : {ℓ : Level}
            → ((monoid' a _ _ _ _ _ ) (monoid' a' _ _ _ _ _ ) : Monoid' {ℓ} )
            → (a → a') → Type ℓ
monoidHom
  (monoid' A ε _ • _ left-unit right-unit assoc carrier-set)
  (monoid' A1 ε1 _ •1 _ left-unit1 right-unit1 assoc1 carrier-set1)
  f
  = (m1 m2 : A) → f (m1 • m2) ≡ (f m1) •1 (f m2)

```

While the first three definitions above are should be linguistically comprehensible to a non-mathematician, this last definition is most certainly not. While may carry degree of comprehension to a programmer or mathematician not exposed to Agda, it is certainly comprehensible to a computer : that is, it typechecks on a computer where Cubical Agda is installed. While GF is designed for multilingual syntactic transformations and is targeted for natural language translation, it's underlying theory is largely based on ideas from the compiler communities. A cousin of the BNF Converter (BNFC), GF is fully capable of parsing programming languages like Agda! And while the above definition is just another concrete syntactic presentation of a group homomorphism, it is distinct from the natural language presentations above in that the colors indicate it has indeed type checked.

While this example may not exemplify the power of Agda's type checker, it is of considerable interest to many. The typechecker has merely assured us that monoidHom, is a well-formed type. The type-checker is much more useful than is immediately evident: it delegates the work of verifying that a proof is correct, that is, the work of judging whether a term has a type, to the computer. While it's of practical concern is immediate to any exploited grad student grading papers late on a Sunday night, its theoretical concern has led to many recent developments in modern mathematics. Thomas Hales solution to the Kepler Conjecture was seen as unverifiable by those reviewing it. This led to Hales outsourcing the verification to Interactive Theorem provers HOL Light and Isabelle, during which led to many minor corrections in the original proof which were never spotted due to human oversight.

Fields Medalist Vladimir Voevodsky, had the experience of being told one day his proof of the Milnor conjecture was fatally flawed. Although the leak in the proof was patched, this experience of temporarily believing much of his life's work invalid led him to investigate proof assistants as a tool for future thought. Indeed,

this proof verification error was a key event that led to the Univalent Foundations Project [2].

While Agda and other programming languages are capable of encoding definitions, theorems, and proofs, they have so far seen little adoption, and in some cases treated with suspicion and scorn by many mathematicians. This isn't entirely unfounded : it's a lot of work to learn how to use Agda or Coq, software updates may cause proofs to break, and the inevitable errors we humans are instilled in these Theorem Provers. And that's not to mention that Martin-Löf Type Theory, the constructive foundational project which underlies these proof assistants, is rejected by those who dogmatically accept the law of the excluded middle and ZFC as the word of God.

It should be noted, the constructivist rejects neither the law of the excluded middle nor ZFC. She merely observes them, and admits their handiness in certain situations. Excluded middle is indeed a helpful tool, as many mathematicians may attest. The contention is that it should be avoided whenever possible - proofs which don't rely on it, or it's corollary of proof by contradiction, are much more amenable to formalization in systems with decidable type checking. And ZFC, while serving the mathematicians of the early 20th century, is lacking when it comes to the higher dimensional structure of n-categories and infinity groupoids.

What these theorem provers give the mathematician is confidence that her work is correct, and even more importantly, that the work which she takes for granted and references in her work is also correct. The task before us is then one of religious conversion. And one doesn't undertake a conversion by simply by preaching. Foundational details aside, this thesis is meant to provide a blueprint for the syntactic reformation that must take place.

It doesn't ask the mathematician to relinquish the beautiful language she has come to love in expressing her ideas. Rather, it asks her to make a compromise for the time being, and use a Controlled Natural Language (CNL) to develop her work. In exchange she'll get the confidence that Agda provides. Not only that, she'll be able to search through a library, to see who else has possibly already postulated and proved her conjecture. A version of this grandiose vision is explored in The Formal Abstracts Project.

References

- [1] Harper, R. (2011). The holy trinity.
- [2] The Univalent foundations program & Institute for advanced study (Princeton, N. (2013). *Homotopy Type Theory: Univalent Foundations of Mathematics*.

Appendices