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## Sumário

<b>1</b>	<b>Template</b>	<b>2</b>
1.1	Macros	2
<b>2</b>	<b>Numerical algorithms</b>	<b>2</b>
2.1	Triângulo de Pascal	2
2.2	GCD-LCM	2
2.3	Bezout Theorem	2
2.4	Teorema Chinês dos Restos	2
2.5	Crivo de Eratóstenes	3
2.6	Divisores de N	3
2.7	Funções com Números Primos (Crivo, Fatoração, PHI, etc)	3
2.8	Exponenciação Modular Rápida	4
2.9	Exponenciação de Matriz	4
2.10	Brent Cycle Detection	4
2.11	Romberg's method - Calcula Integral (UFS2010)	4
2.12	Pollard's rho algorithm (UFS2010)	5
2.13	Miller-Rabin's algorithm (UFS2010)	5
2.14	Quantidade de dígitos de N! na base B	5
2.15	Quantidade de zeros a direita de N! na base B	5
2.16	Baby Step Giant Step	5
2.17	Primos num intervalo	6
2.18	FFT	6
<b>3</b>	<b>Geometria 2D</b>	<b>6</b>
3.1	Geometria 2D Library	6
<b>4</b>	<b>Polígonos 2D</b>	<b>9</b>
4.1	Polígono 2D Library	9
4.2	Convex Hull	10
4.3	Minimum Enclosing Circle	11
<b>5</b>	<b>Geometria 3D</b>	<b>11</b>
5.1	Geometria 3D Library	11
<b>6</b>	<b>Grafos</b>	<b>14</b>
6.1	Topological Sort	14
6.2	Dijkstra	15
6.3	Floyd-Warshall	15
6.4	Bellman-Ford	16
6.5	Vértices de Articulação e Pontes	16
6.6	Tarjan	16
6.7	Kosaraju	17
6.8	2-Sat	17
6.9	LCA	18
6.10	LCA (Sparse Table)	19
6.11	Maximum Bipartite Matching	19
6.12	Hopcroft Karp - Maximum Bipartite Matching (UNIFEI)	19
6.13	Network Flow (lento)	20
6.14	Network Flow - Dinic	20
6.15	Min Cost Max Flow	21
6.16	Min Cost Max Flow (Stefano)	22
6.17	Tree Isomorphism	23
6.18	Stoer Wagner-Minimum Cut (UNIFEI)	24
6.19	Erdos Gallai (UNIFEI)	24

6.20	Stable Marriage (UNIFEI)	25
6.21	Hungarian Max Bipartite Matching with Cost (UNIFEI)	25
6.22	Blossom	27
<b>7</b>	<b>Estruturas de Dados</b>	<b>28</b>
7.1	BIT	28
7.2	BIT 2D	28
7.3	Sparse Table	29
7.4	RMQ	29
7.5	Persistent Seg Tree	30
7.6	Seg Tree com Lazy	30
7.7	Union-Find	31
7.8	Treap	31
7.9	Seg Tree 2D	32
7.10	Polyce	33
7.11	KD2	34
<b>8</b>	<b>Strings</b>	<b>35</b>
8.1	KMP	35
8.2	Aho Corasick	35
8.3	Suffix Array	37
8.4	Suffix Array (Gugu)	38
8.5	Rolling Hash	38
8.6	Longest Common Prefix with Hash	38
8.7	Minimum Lexicographic Rotation	39
8.8	Longest Palindrome (Manacher algorithm)	39
8.9	Autômato de Sufixos	39
8.10	Suffix Array	40
8.11	Z Algorithm	41
<b>9</b>	<b>PD</b>	<b>41</b>
9.1	Soma acumulada 2D	41
9.2	Knuth Optimization	41
9.3	Convex Hull Trick	42
9.4	Longest Increasing Subsequence	43
9.5	Kadane 1D	43
9.6	Kadane 2D	43
9.7	Knapsack0-1	43
9.8	Edit Distance	43
<b>10</b>	<b>Sorting</b>	<b>44</b>
10.1	Merge Sort com num de Inversoes	44
10.2	Quick Sort	44
<b>11</b>	<b>Miscelânea</b>	<b>44</b>
11.1	Calendário	44

# 1 Template

## 1.1 Macros

```
#include <bits/stdc++.h>
using namespace std;

#define rep(i, a, b) for(int i=(a); i<(b); i++)
#define pb push_back
#define mp make_pair
#define debug(x) cout<<__LINE__<<": "<<#x<<" = "<<x<<endl;
#define debug2(x, y) cout<<__LINE__<<": "<<#x<<" = "<<x<<\
" "<<#y<<" = "<<y<<endl;

#define all(c) (c).begin(), (c).end()
#define F first
#define S second
#define UNIQUE(c) \
    sort(all(c)); \
    (c).resize(unique(all(c))-c.begin());

#define PI 3.1415926535897932384626433832795028841971

typedef long long ll;
typedef pair<int, int> ii;
typedef vector<int> vi;

const int INF = 0x3f3f3f3f;
const double EPS = 1e-9;

inline int cmp(double x, double y = 0, double tol = EPS) {
    return ((x <= y+tol) ? (x+tol < y) ? -1:0:1);
}
```

# 2 Numerical algorithms

## 2.1 Triângulo de Pascal

```
// Calcula os numeros binomiais (N,K) = N!/(K!(N-K)!). (N,K)
// representa o numero de maneiras de criar um subconjunto de tamanho
// K dado um conjunto de tamanho N. A ordem dos elementos nao
// importa.
const int MAXN = 50;
long long C[MAXN][MAXN];
void calc_pascal() {
    memset(C, 0, sizeof(C));
    for (int i = 0; i < MAXN; ++i) {
        C[i][0] = C[i][i] = 1;
        for (int j = 1; j < i; ++j)
            C[i][j] = C[i - 1][j - 1] + C[i - 1][j];
    }
}
/*
Pascal triangle elements:
C(33, 16) = 1.166.803.110 [int limit]
C(34, 17) = 2.333.606.220 [unsigned int limit]
C(66, 33) = 7.219.428.434.016.265.740 [int64_t limit]
C(67, 33) = 14.226.520.737.620.288.370 [uint64_t limit]
```

```
Fatorial
12 != 479.001.600 [(unsigned)int limit]
20 != 2.432.902.008.176.640.000 [(unsigned)int64_t limit]
*/
```

## 2.2 GCD-LCM

```
// Calcula o maior divisor comum entre A e B
ll A, B;
cin >> A >> B;
cout << __gcd(A, B);

// Calcula o menor multiplo comum entre A e B
ll lcm(ll A, ll B) {
    if (A and B)
        return abs(A) / __gcd(A, B) * abs(B);
    else
        return abs(A | B);
}
```

## 2.3 Bezout Theorem

```
// Determina a solucao da equacao a*x+b*y = gcd(a, b), onde a e b sao
// dois numeros naturais. Como chamar: egcd(a, b), Retorna: a tupla
// {gcd(a, b), x, y}. Determina tambem o Inverso Modular.
struct Triple {
    ll d, x, y;
    Triple(ll q, ll w, ll e) : d(q), x(w), y(e) {}
};
Triple egcd(ll a, ll b) {
    if (!b)
        return Triple(a, 1, 0);
    Triple q = egcd(b, a % b);
    return Triple(q.d, q.y, q.x - a / b * q.y);
}
// Retorna o inverso modular de A modulo N
// O inverso modular de um numero A em relacao a N eh um numero X tal
// que (A*X)%N = 1
ll invMod(ll a, ll n) {
    Triple t = egcd(a, n);
    if (t.d > 1)
        return 0;
    return (t.x % n + n) % n;
}
```

## 2.4 Teorema Chinês dos Restos

```
// crt() retorna um X tal que X = a[i] (mod m[i]). Exemplo: Para a[]
// = {1, 2, 3} e m[] = {5, 6, 7} .: X = 206. Requer: Bezout Theorem
// para calcular o inverso modular
#define MAXN 1000
int n;
ll a[MAXN], m[MAXN];
ll crt() {
    ll M = 1, x = 0;
```

```

for (int i = 0; i < n; ++i)
    M *= m[i];
for (int i = 0; i < n; ++i)
    x += a[i] * invMod(M / m[i], m[i]) * (M / m[i]);
return ((x % M) + M) % M;
}

```

## 2.5 Crivo de Eratóstenes

```

bitset<10000005> bs;
vector<int> primos;
void crivo(ll limite = 10000000LL) { // calcula primos ate limite
    primos.clear();
    bs.set();
    bs[0] = bs[1] = 0;
    for (ll i = 2; i <= limite; i++)
        if (bs[i]) {
            for (ll j = i * i; j <= limite; j += i)
                bs[j] = 0;
            primos.push_back(i);
        }
}
bool isPrime(ll N, ll limite) {
    if (N <= limite)
        return bs[N];
    for (int i = 0; i < (int)primos.size(); i++)
        if (N % primos[i] == 0)
            return false;
    return true;
}

```

## 2.6 Divisores de N

```

// Retorna todos os divisores naturais de N em O(sqrt(N)).
vector<ll> divisores(ll N) {
    vector<ll> divisors;
    for (ll div = 1, k; div * div <= N; ++div) {
        if (N % div == 0) {
            divisors.push_back(div);
            k = N / div;
            if (k != div)
                divisors.push_back(k);
        }
    }
    // caso precise ordenado
    sort(divisors.begin(), divisors.end());
    return divisors;
}

```

## 2.7 Funções com Números Primos (Crivo, Fatoração, PHI, etc)

```

// Encontra os fatores primos de N :: N = p1^e1 * ... * pi^ei
// factors armazena em first o fator primo e em segundo seu expoente
map<int, int> factors;

```

```

void primeFactors(ll N) {
    factors.clear();
    while (N % 2 == 0)
        ++factors[2], N >>= 1;
    for (ll PF = 3; PF * PF <= N; PF += 2) {
        while (N % PF == 0)
            N /= PF, factors[PF]++;
    }
    if (N > 1)
        factors[N] = 1;
}

// Funcoess derivadas dos numeros primos
void NumberTheory(ll N) {
    primeFactors(N);
    map<int, int>::iterator f; // iterador
    ll Totient = N; // Totiente ou Euler-Phi de N
    // Totient(N) = qtos naturais x, tal que x < N && gcd(x,N) == 1
    ll numDiv = 1; // Quantidade de divisores de N
    ll sumDiv = 1; // Soma dos divisores de N
    ll sumPF = 0; // Soma dos fatores primos de N (trivial)
    ll numDiffPF = factors.size(); // qtde de fatores distintos

    for (f = factors.begin(); f != factors.end(); f++) {
        ll PF = f->first, power = f->second;
        Totient -= Totient / PF;
        numDiv *= (power + 1);
        sumDiv *= ((ll)pow((double)PF, power + 1.0) - 1) / (PF - 1);
        sumPF += PF;
    }

    printf("Totiente/Euler-Phi de N = %lld\n", Totient);
    printf("qt de divisores de N = %lld\n", numDiv);
    printf("soma dos divisores de N = %lld\n", sumDiv);
    printf("qt de fatores primos distintos = %lld\n", numDiffPF);
    printf("soma dos fatores primos = %lld\n", sumPF);
}

```

```

// Calcula Euler Phi para cada valor do intervalo [1, N]
#define MM 1000010
int phi[MM];
void crivo_euler_phi(int N) {
    for (int i = 1; i <= N; i++)
        phi[i] = i;
    for (int i = 2; i <= N; i++)
        if (phi[i] == i) {
            for (int k = i; k <= N; k += i)
                phi[k] = (phi[k] / i) * (i - 1);
        }
}

// Qtde de fatores primos distintos de cada valor do range [2, MAX_N]
#define MAX_N 1000000
int NDPF[MAX_N]; //
void NumDiffPrimeFactors() {
    memset(NDPF, 0, sizeof NDPF);
    for (int i = 2; i < MAX_N; i++)
        if (NDPF[i] == 0)
            for (int j = i; j < MAX_N; j += i)
                NDPF[j]++;
}

```

```
int main() { return 0; }
```

## 2.8 Exponenciação Modular Rápida

```
/**
 * fastpow() realiza exponenciacao rapida de inteiros
 * #param ll b - base da exponenciacao
 * #param ll expo - expoente
 * #param ll mod - o resultado sera calculado modulo este valor
 * #return - o valor de (b ^ p) % mod
 * #complexidade - O(log(p))
 */

ll fastpow(ll b, ll expo, ll mod) {
    ll ret = 1, pot = b % mod;
    while (expo) {
        if (expo & 1) {
            ret = (ret * pot) % mod;
        }
        pot = (pot * pot) % mod;
        expo >>= 1;
    }
    return ret;
}
```

## 2.9 Exponenciação de Matriz

```
/**
 * fastpow() realiza exponenciacao rapida de matrizes
 * #param matrix_t M - matriz a ser elevada
 * #param ll expo - expoente
 * #param ll mod - o resultado sera calculado modulo este valor
 * #return - o resultado de (M ^ expo) % mod
 * #complexidade - O(size^3 * log(expo))
 */

#define MAX (ChangeMe) // Max size of square matrix

struct matrix_t {
    ll m[MAX][MAX];
    int sz;
    matrix_t(int _sz) {
        memset(m, 0, sizeof(m));
        sz = _sz;
    }
    matrix_t multiply(const matrix_t other, ll mod) {
        matrix_t ret(other.sz);
        rep(i, 0, sz) rep(j, 0, sz) {
            ret.m[i][j] = 0;
            rep(k, 0, sz) ret.m[i][j] =
                (ret.m[i][j] + (m[i][k] * other.m[k][j]) % mod) % mod;
        }
        return ret;
    }
};

matrix_t fastpow(matrix_t M, ll expo, ll mod) {
```

```
matrix_t ret(M.sz);
// Identity matrix
rep(i, 0, ret.sz) rep(j, 0, ret.sz) ret.m[i][j] = (i == j);

while (expo) {
    if (expo & 1)
        ret = ret.multiply(M, mod);
    M = M.multiply(M, mod);
    expo >>= 1;
}
return ret;
}
```

## 2.10 Brent Cycle Detection

```
// Dado uma sequencia formada por uma funcao f(.) e uma semente x0.
// f(x0), f(f(x0)), ..., f(f(...f(x0))), ela pode ser ciclica. Este
// algoritmo retorna o tamanho do ciclo e o valor xi que o inicia.
ii brent_cycle(int x) {
    int p = 1, length = 1, t = x, start = 0;
    int h = f(x);
    while (t != h) {
        if (p == length) {
            t = h;
            p *= 2;
            length = 0;
        }
        h = f(h);
        ++length;
    }
    t = h = x;
    for (int i = length; i != 0; --i)
        h = f(h);
    while (t != h) {
        t = f(t);
        h = f(h);
        ++start;
    }
    return ii(start, length);
}
```

## 2.11 Romberg's method - Calcula Integral (UFS2010)

```
// Calcula a integral de f[a, b]
typedef long double ld;

ld f(double x) {
    // return f(x)
}

ld romberg(ld a, ld b) {
    ld R[16][16], div = (b - a) / 2;
    R[0][0] = div * (f(a) + f(b));
    for (int n = 1; n <= 15; n++, div /= 2) {
        R[n][0] = R[n - 1][0] / 2;
        for (ld sample = a + div; sample < b; sample += 2 * div)
            R[n][0] += div * f(a + sample);
    }
}
```

```

for (int m = 1; m <= 15; m++)
    for (int n = m; n <= 15; n++)
        R[n][m] = R[n][m - 1] +
            1 / (pow(4, m) - 1) * (R[n][m - 1] - R[n - 1][m - 1]);
return R[15][15];
}

```

## 2.12 Pollard's rho algorithm (UFS2010)

```

// Retorna um fator primo de N, util para fatorizacao quando N for
// grande.
ll pollard_rho, pollard_n;
ll f(ll val) { return (val * val + pollard_rho) % pollard_n; }
ll myabs(ll a) { return a >= 0 ? a : -a; }
ll pollard(ll n) {
    srand(unsigned(time(0)));
    pollard_n = n;
    long long d = 1;
    do {
        d = 1;
        pollard_rho = rand() % n;
        long long x = 2, y = 2;
        while (d == 1)
            x = f(x), y = f(f(y)), d = __gcd(myabs(x - y), n);
    } while (d == n);
    return d;
}

```

## 2.13 Miller-Rabin's algorithm (UFS2010)

```

// Teste probabilistico de primalidade
bool miller_rabin(ll n, ll base) {
    if (n <= 1)
        return false;
    if (n % 2 == 0)
        return n == 2;
    ll s = 0, d = n - 1;
    while (d % 2 == 0)
        d /= 2, ++s;
    ll base_d = fastpow(base, d, n);
    if (base_d == 1)
        return true;
    ll base_2r = base_d;
    for (ll i = 0; i < s; ++i) {
        if (base_2r == 1)
            return false;
        if (base_2r == n - 1)
            return true;
        base_2r = base_2r * base_2r % n;
    }
    return false;
}
bool isprime(ll n) {
    if (n == 2 || n == 7 || n == 61)
        return true;
    return miller_rabin(n, 2) && miller_rabin(n, 7) &&
        miller_rabin(n, 61);
}

```

## 2.14 Quantidade de dígitos de N! na base B

```

int NumOfDigitsInFactorial(int N, int B) {
    double logFatN = 0;
    for (int i = 1; i <= N; i++)
        logFatN += log((double)i);
    int nd = floor(logFatN / log((double)B)) + 1;
    return nd;
}

```

## 2.15 Quantiade de zeros a direita de N! na base B

```

// Determina o numero de zeros a direita do fatorial de N na base B
// Ideia: Se a base for B for 10, e fatorarmos N! em fatores primos
// teremos algo como N! = 2^a * 3^b * 5^c ..., como cada par de primos
// 2 e 5 formam 10 que tem um zero, a quantidade seria min(a, c).
int NumOfTrailingZeros(int N, int B) {
    int nfact = fatora(B);
    int zeros = INF;
    // para cada fator de B, aux representa qtas vezes
    // fator[i]^expoente[i] aparece na representacao de N!
    for (int i = 0; i < nfact; i++) {
        int soma = 0;
        int NN = N;
        while (NN) {
            soma += NN / fator[i];
            NN /= fator[i];
        }
        int aux = soma / expoente[i];
        zeros = min(zeros, aux);
    }
    return zeros;
}

```

## 2.16 Baby Step Giant Step

```

// Determinar o menor E tal que B^E = N (mod P), -1 se for impossivel.
// Requer: Bezout Theorem para calcular o inverso modular
ll bsgs(ll b, ll n, ll p) {
    if (n == 1)
        return 0;
    map<ll, int> table;
    ll m = sqrt(p) + 1, pot = 1, pot2 = 1;
    for (int j = 0; j < m; ++j) {
        if (pot == n)
            return j;
        table[(n * invMod(pot, p)) % p] = j;
        pot = (pot * b) % p;
    }
    for (int i = 0; i < m; ++i) {
        if (table.find(pot2) != table.end())
            return i * m + table[pot2];
        pot2 = (pot2 * pot2) % p;
    }
    return -1;
}

```

## 2.17 Primos num intervalo

```
// Encontra os primos no intervalo [n,m]
vector<int> ret;
void primesBetween(int n, int m) {
    ret.clear();
    vector<int> primes(m - n + 1);
    for (int i = 0; i < m - n + 1; ++i)
        primes[i] = 0;
    for (int p = 2; p * p <= m; ++p) {
        int less = (n / p) * p;
        for (int j = less; j <= m; j += p)
            if (j != p && j >= n)
                primes[j - n] = 1;
    }
    for (int i = 0; i < m - n + 1; ++i) {
        if (primes[i] == 0 && n + i != 1) {
            ret.push_back(n + i);
        }
    }
}
```

## 2.18 FFT

```
typedef complex<double> comp;
const int MAX_N = 1 << 20;
int rev[MAX_N];
comp roots[MAX_N];

void preCalc(int N, int BASE) {
    for (int i = 1; i < N; ++i)
        rev[i] = (rev[i >> 1] >> 1) + ((i & 1) << (BASE - 1));
    int NN = N >> 1;
    roots[NN] = comp(1, 0);
    roots[NN + 1] = comp(cos(2 * PI / N), sin(2 * PI / N));
    for (int i = 2; i < NN; ++i)
        roots[NN + i] = roots[NN + i - 1] * roots[NN + 1];
    for (int i = NN - 1; i > 0; --i)
        roots[i] = roots[2 * i];
}

void fft(vector<comp> &a, bool invert) {
    int N = a.size();
    if (invert)
        rep(i, 0, N) a[i] = conj(a[i]);
    rep(i, 0, N) if (i < rev[i]) swap(a[i], a[rev[i]]);
    for (int k = 1; k < N; k *= 2) {
        for (int i = 0; i < N; i += 2 * k) {
            rep(j, 0, k) {
                comp B = a[i + j + k] * roots[k + j];
                a[i + j + k] = a[i + j] - B;
                a[i + j] = a[i + j] + B;
            }
        }
    }
    if (invert)
        rep(i, 0, a.size()) a[i] /= N;
}
```

```
vector<comp> multiply_real(vector<comp> a, vector<comp> b,
                          vector<comp> c) {
    int n = a.size();
    int m = b.size();

    int base = 0, N = 1;
    while (N < n + m - 1)
        base++, N <= 1;
    preCalc(N, base);

    a.resize(N, comp(0, 0));
    c.resize(N);

    rep(i, 0, b.size()) a[i] = comp(real(a[i]), real(b[i]));
    fft(a, 0);
    rep(i, 0, N) {
        int j = (N - i) & (N - 1);
        c[i] = (a[i] * a[i] - conj(a[j] * a[j])) * comp(0, -0.25);
    }
    fft(c, 1);
    return c;
}
```

## 3 Geometria 2D

### 3.1 Geometria 2D Library

```
const double EPS = 1e-9;
inline int cmp(double x, double y = 0, double tol = EPS) {
    return ((x <= y + tol) ? (x + tol < y) ? -1 : 0 : 1);
}

struct point {
    double x, y;
    point(double x = 0, double y = 0) : x(x), y(y) {}
    point operator+(point q) const { return point(x + q.x, y + q.y); }
    point operator-(point q) const { return point(x - q.x, y - q.y); }
    point operator*(double t) const { return point(x * t, y * t); }
    point operator/(double t) const { return point(x / t, y / t); }
    int cmp(point q) const {
        if (int t = ::cmp(x, q.x))
            return t;
        return ::cmp(y, q.y);
    }
    bool operator==(point q) const { return cmp(q) == 0; }
    bool operator!=(point q) const { return cmp(q) != 0; }
    bool operator<(point q) const { return cmp(q) < 0; }
};

ostream &operator<<(ostream &os, const point &p) {
    os << "(" << p.x << "," << p.y << ")";
}

#define vec(a, b) (b - a)
typedef vector<point> polygon;

double cross(point a, point b) { return a.x * b.y - a.y * b.x; }
double dot(point a, point b) { return a.x * b.x + a.y * b.y; }
double collinear(point a, point b, point c) {
```

```

    return cmp(cross(b - a, c - a)) == 0;
}
// retorna 1 se R esta a esquerda do vetor P->Q, -1 se estiver a
// direita. 0 se P, Q e R forem colineares
int ccw(point p, point q, point r) { return cmp(cross(q - p, r - p));
}

// Rotaciona um ponto em relacao a origem em 90 graus sentido
// anti-horario
point RotateCCW90(point p) { return point(-p.y, p.x); }
// Rotaciona um ponto em relacao a origem em 90 graus sentido horario
point RotateCW90(point p) { return point(p.y, -p.x); }

// Rotaciona um ponto P em A graus no sentido anti-horario em relacao
// a
// origem; Para rotacionar no sentido horario, basta A ser negativo
point RotateCCW(point p, double a) {
    a = (a / 180.0) * acos(-1.0); // convertendo para radianos
    return point(p.x * cos(a) - p.y * sin(a),
                p.x * sin(a) + p.y * cos(a));
}
// Rotaciona P em A graus em relacao a Q.
point RotateCCW(point p, point q, double a) {
    return RotateCCW(p - q, a) + q;
}
// Tamanho ou norma de um vetor
double abs(point u) { return sqrt(dot(u, u)); }
// Projeta o vetor A sobre a direcao do vetor B
point project(point a, point b) { return b * (dot(a, b) / dot(b, b));
}
// Retorna a projecao do ponto P sobre reta definida por [A,B]
point projectPointLine(point p, point a, point b) {
    return p + project(p - a, b - a);
}
// Retorna o angulo que p faz com +x
double arg(point p) { return atan2(p.y, p.x); }
// Retorna o angulo entre os vetores AB e AC
double arg(point b, point a, point c) {
    point u = b - a, v = c - a;
    return atan2(cross(u, v), dot(u, v));
}

////////////////////////////////////
//////////////////////////////////// Segmentos, Retas //////////////////////////////////////
////////////////////////////////////

// Determina se P esta entre o segmento fechado [A,B], inclusive
bool between(point p, point a, point b) {
    return collinear(p, a, b) && dot(a - p, b - p) <= 0;
}

/* Distancia de ponto P para reta que passa por [A,B]. Armazena em C
 * (por ref) o ponto projecao de P na reta. */
double distancePointLine(point p, point a, point b, point &c) {
    c = projectPointLine(p, a, b);
    return fabs(cross(p - a, b - a)/abs(a - b); // or abs(p-c);
}

/* Distancia de ponto P ao segmento [A,B]. Armazena em C (por ref) o
 * ponto de projecao de P em [A,B]. Se este ponto estiver fora do
 * segmento, eh retornado o mais proximo. */

```

```

double distancePointSeg(point p, point a, point b, point &c) {
    if ((b - a) * (p - a) <= 0) {
        c = a;
        return abs(a - p);
    }
    if ((a - b) * (p - b) <= 0) {
        c = b;
        return abs(b - p);
    }

    c = projectPointLine(p, a, b);
    return fabs(cross(p - a, b - a)/abs(a - b); // or abs(p-c);
}

// Determina se os segmentos [A, B] e [C, D] se tocam
bool seg_intersect(point a, point b, point c, point d) {
    int d1, d2, d3, d4;
    d1 = ccw(c, a, d);
    d2 = ccw(c, b, d);
    d3 = ccw(a, c, b);
    d4 = ccw(a, d, b);
    if (d1 * d2 == -1 && d3 * d4 == -1)
        return true;
    if (d1 == 0 && between(c, a, d))
        return true;
    if (d2 == 0 && between(c, b, d))
        return true;
    if (d3 == 0 && between(a, c, b))
        return true;
    if (d4 == 0 && between(a, d, b))
        return true;
    return false;
}

// Encontra a interseccao das retas (p-q) e (r-s) assumindo que existe
// apenas 1 interseccao. Se for entre segmentos, verificar se
// intersectam primeiro.
point line_intersect(point p, point q, point r, point s) {
    point a = q - p, b = s - r, c = point(cross(p, q), cross(r, s));
    double x = cross(point(a.x, b.x), c);
    double y = cross(point(a.y, b.y), c);
    return point(x, y) / cross(a, b);
}

// determine if lines from a to b and c to d are parallel or collinear
bool LinesParallel(point a, point b, point c, point d) { //!! Nao
                                                         //!! testado
    return fabs(cross(b - a, c - d)) < EPS;
}
bool LinesCollinear(point a, point b, point c,
                    point d) { //!! Nao testado
    return LinesParallel(a, b, c, d) &&
        fabs(cross(a - b, a - c)) < EPS &&
        fabs(cross(c - d, c - a)) < EPS;
}

////////////////////////////////////
//////////////////////////////////// Triangulos //////////////////////////////////////
////////////////////////////////////

bool pointInTriangle(point p, point a, point b, point c) {

```

```

// TODO
}

// Heron's formula - area do triangulo(a,b,c) -1 se nao existe
double area_heron(double a, double b, double c) {
    if (a < b)
        swap(a, b);
    if (a < c)
        swap(a, c);
    if (b < c)
        swap(b, c);
    if (a > b + c)
        return -1;
    return sqrt((a + (b + c)) * (c - (a - b)) * (c + (a - b)) *
                (a + (b - c)) / 16.0);
}

////////////////////////////////////
//////////////////// CIRCULOS //////////////////////
////////////////////////////////////

bool pointInCircle(point p, point c, double radius) {
    // Todo
}

/*Dado dois pontos (A, B) de uma circunferencia e seu raio R, eh
 * possivel obter seus possiveis centros (C1 e C2). Para obter o outro
 * centro, basta inverter os paramentros */
bool circle2PtsRad(point a, point b, double r, point &c) {
    point aux = a - b;
    double d = dot(aux, aux);
    double det = r * r / d - 0.25;
    if (det < 0.0)
        return false;
    double h = sqrt(det);
    c.x = (a.x + b.x) * 0.5 + (a.y - b.y) * h;
    c.y = (a.y + b.y) * 0.5 + (b.x - a.x) * h;
    return true;
}

// Menor distancia entre dois pontos numa esfera de raio r
// lat = [-90,90]; long = [-180,180]
double spherical_distance(double lt1, double lo1, double lt2,
                          double lo2, double r) {
    double pi = acos(-1);
    double a = pi * (lt1 / 180.0), b = pi * (lt2 / 180.0);
    double c = pi * ((lo2 - lo1) / 180.0);
    return r * acos(sin(a) * sin(b) + cos(a) * cos(b) * cos(c));
}

////////////////////////////////////
//////////////////// PLANOS //////////////////////
////////////////////////////////////

// Distancia entre (x,y,z) e plano ax+by+cz=d
double distancePointPlane(double x, double y, double z, double a,
                           double b, double c, double d) {
    return fabs(a * x + b * y + c * z - d) / sqrt(a * a + b * b + c * c);
}

```

```

/**[Inicio] Funcoes que usam numeros complexos para pontos**
typedef complex<double> cxpt;
struct circle {
    cxpt c;
    double r;
    circle(cxpt c, double r) : c(c), r(r) {}
    circle() {}
};
double cross(const cxpt &a, const cxpt &b) {
    return imag(conj(a) * b);
}
double dot(const cxpt &a, const cxpt &b) { return real(conj(a) * b); }

// Area da interseccao de dois circulos
double circ_inter_area(circle &a, circle &b) {
    double d = abs(b.c - a.c);
    if (d <= (b.r - a.r))
        return a.r * a.r * M_PI;
    if (d <= (a.r - b.r))
        return b.r * b.r * M_PI;
    if (d >= a.r + b.r)
        return 0;
    double A = acos((a.r * a.r + d * d - b.r * b.r) / (2 * a.r * d));
    double B = acos((b.r * b.r + d * d - a.r * a.r) / (2 * b.r * d));
    return a.r * a.r * (A - 0.5 * sin(2 * A)) +
        b.r * b.r * (B - 0.5 * sin(2 * B));
}

// Pontos de interseccao de dois circulos
// Intersects two circles and intersection points are in 'inter'
// -1-> outside, 0-> inside, 1-> tangent, 2-> 2 intersections
int circ_circ_inter(circle &a, circle &b, vector<cxpt> &inter) {
    double d2 = norm(b.c - a.c), rS = a.r + b.r, rD = a.r - b.r;
    if (d2 > rS * rS)
        return -1;
    if (d2 < rD * rD)
        return 0;
    double ca = 0.5 * (1 + rS * rD / d2);
    cxpt z = cxpt(ca, sqrt((a.r * a.r / d2) - ca * ca));
    inter.push_back(a.c + (b.c - a.c) * z);
    if (abs(z.imag()) > EPS)
        inter.push_back(a.c + (b.c - a.c) * conj(z));
    return inter.size();
}

// Line-circle intersection
// Intersects (infinite) line a-b with circle c
// Intersection points are in 'inter'
// 0 -> no intersection, 1 -> tangent, 2 -> two intersections
int line_circ_inter(cxpt a, cxpt b, circle c, vector<cxpt> &inter) {
    c.c -= a;
    b -= a;
    cxpt m = b * real(c.c / b);
    double d2 = norm(m - c.c);
    if (d2 > c.r * c.r)
        return 0;
    double l = sqrt((c.r * c.r - d2) / norm(b));
    inter.push_back(a + m + l * b);
    if (abs(l) > EPS)
        inter.push_back(a + m - l * b);
    return inter.size();
}

```



```

}
/**[FIM] Funcoes que usam numeros complexos para pontos**/

```

## 4 Polígonos 2D

### 4.1 Polígono 2D Library

```

/*Poligono eh representado como um array de pontos T[i] sao os
vertices
do poligono. Existe uma aresta que conecta T[i] com T[i+1], e
T[size-1]
com T[0]. Logo assume-se que T[0] != T[size-1]
Poligono simples: Aquele em que as arestas nao se interceptam.
Convexo:
O angulo interno de T[i] com T[i-1] e T[i+1] <= 180. Concavo: Existe
algum i que nao satisfaz a condicao anterior*/

```

```

/* Retorna a area com sinal de um poligono T. Se area > 0, T esta
* listado na ordem CCW */
double signedArea(const polygon &T) {
    double area = 0;
    int n = T.size();
    if (n < 3)
        return 0;
    rep(i, 0, n) area += cross(T[i], T[(i + 1) % n]);
    return (area / 2.0);
}

```

```

/* Retorna a area de um poligono T. (pode ser concavo ou convexo) em
* O(N)*/
double poly_area(const polygon &T) { return fabs(signedArea(T)); }

```

```

/* Retorna a centroide de um poligono T em O(N)*/

```

```

point centroide(const polygon &T) {
    int n = T.size();
    double sgnArea = signedArea(T);

    point c = point(0, 0);
    rep(i, 0, n) {
        int k = (i + 1) % n;
        c = c + (T[i] + T[k]) * cross(T[i], T[k]);
    }
    c = c / (sgnArea * 6.0);
    return c;
}

```

```

/* Retorna o perimetro do poligono T. (pode n funcionar como esperado
* se o poligono for uma linha reta (caso degenerado))*/

```

```

double poly_perimeter(polygon &T) {
    double perimeter = 0;
    int n = T.size();
    if (n < 3)
        return 0;
    rep(i, 0, n) perimeter += abs(T[i] - T[(i + 1) % n]);
    return perimeter;
}

```

```

// tests whether or not a given polygon (in CW or CCW order) is simple

```

```

bool isSimple(const polygon &p) { // nao testado
    for (int i = 0; i < p.size(); i++) {
        for (int k = i + 1; k < p.size(); k++) {
            int j = (i + 1) % p.size();
            int l = (k + 1) % p.size();
            if (i == l || j == k)
                continue;
            if (seg_intersect(p[i], p[j], p[k], p[l]))
                return false;
        }
    }
    return true;
}

```

```

// Retorna True se T for convexo. O(N)

```

```

bool isConvex(polygon &T) {
    int n = T.size();
    if (n < 3)
        return false;

    int giro = 0;
    rep(i, 0, n) { // encontra um giro valido
        int t = ccw(T[i], T[(i + 1) % n], T[(i + 2) % n]);
        if (t != 0)
            giro = t;
    }
    if (giro == 0)
        return false; // todos pontos sao colineares

    rep(i, 0, n) {
        int t = ccw(T[i], T[(i + 1) % n], T[(i + 2) % n]);
        if (t != 0 && t != giro)
            return false;
    }
    return true;
}

```

```

// Determina se P pertence a T, funciona para convexo ou concavo
// -1 borda, 0 fora, 1 dentro. O(N)

```

```

int in_poly(point p, polygon &T) {
    double a = 0;
    int N = T.size();
    rep(i, 0, N) {
        if (between(p, T[i], T[(i + 1) % N]))
            return -1;
        a += arg(T[i], p, T[(i + 1) % N]);
    }
    return cmp(a) != 0;
}

```

```

// determina se P pertence a B, funciona APENAS para convexo

```

```

bool PointInConvexPolygon(point P, const polygon &B) {
    int ini = 1, fim = B.size() - 2, mid, pos = -1;
    int giro = -1; // sentido horario
    while (ini <= fim) {
        mid = (ini + fim) / 2;
        int aux = ccw(B[0], B[mid], P);
        if (aux == giro) {
            pos = mid;
            ini = mid + 1;
        } else {

```

```

    fim = mid - 1;
}
}
if (pos == -1)
    return false;
if (ccw(B[0], B[pos], P) != giro * -1 &&
    ccw(B[0], B[pos + 1], P) != giro &&
    ccw(B[pos], B[pos + 1], P) == giro) // giro || 0 na borda
    return true;
return false;
}

// Determina o poligono interseccao de P e Q
// P e Q devem estar orientados anti-horario.
polygon poly_intersect(polygon &P, polygon &Q) {
    int m = Q.size(), n = P.size();
    int a = 0, b = 0, aa = 0, ba = 0, inflag = 0;
    polygon R;
    while ((aa < n || ba < m) && aa < 2 * n && ba < 2 * m) {
        point p1 = P[a], p2 = P[(a + 1) % n], q1 = Q[b],
            q2 = Q[(b + 1) % m];
        point A = p2 - p1, B = q2 - q1;
        int cross = cmp(cross(A, B)), ha = ccw(p2, q2, p1),
            hb = ccw(q2, p2, q1);
        if (cross == 0 && ccw(p1, q1, p2) == 0 && cmp(dot(A, B)) < 0) {
            if (between(q1, p1, p2))
                R.push_back(q1);
            if (between(q2, p1, p2))
                R.push_back(q2);
            if (between(p1, q1, q2))
                R.push_back(p1);
            if (between(p2, q1, q2))
                R.push_back(p2);
            if (R.size() < 2)
                return polygon();
            inflag = 1;
            break;
        } else if (cross != 0 && seg_intersect(p1, p2, q1, q2)) {
            if (inflag == 0)
                aa = ba = 0;
            R.push_back(line_intersect(p1, p2, q1, q2));
            inflag = (hb > 0) ? 1 : -1;
        }
        if (cross == 0 && hb < 0 && ha < 0)
            return R;
        bool t = cross == 0 && hb == 0 && ha == 0;
        if (t ? (inflag == 1) : (cross >= 0) ? (ha <= 0) : (hb > 0)) {
            if (inflag == -1)
                R.push_back(q2);
            ba++;
            b++;
            b %= m;
        } else {
            if (inflag == 1)
                R.push_back(p2);
            aa++;
            a++;
            a %= n;
        }
    }
    if (inflag == 0) {

```

```

        if (in_poly(P[0], Q))
            return P;
        if (in_poly(Q[0], P))
            return Q;
    }
    R.erase(unique(all(R), R.end()));
    if (R.size() > 1 && R.front() == R.back())
        R.pop_back();
    return R;
}

```

## 4.2 Convex Hull

*/\*Encontra o convex hull de um conjunto de pontos em  $O(N \log N)$   
 pivot: Ponto base para a criacao do convex hull;  
 radial\_lt(): Ordena os pontos em sentido anti-horario (ccw).  
 Input: Conjunto de pontos 2D;  
 Output: Conjunto de pontos do convex hull, no sentido anti-horario;*

*(1) Se for preciso manter pontos colineares na borda do convex hull,  
 essa  
 parte evita que eles sejam removidos;  
 \*/*

```

point pivot;
bool radial_lt(point a, point b) {
    int R = ccw(pivot, a, b);
    if (R == 0) // sao colineares
        return (pivot - a) * (pivot - a) < (pivot - b) * (pivot - b);
    else
        return (R == 1); // 1 se A esta a direita de (pivot->B)
}
vector<point> convexhull(vector<point> &T) {
    // Se for necessario remover pontos duplicados
    sort(T.begin(), T.end()); // ordena por x e por y
    T.resize(unique(T.begin(), T.end()) - T.begin());

    int tam = 0, n = T.size();
    vector<point> U; // convex hull

    int idx = min_element(T.begin(), T.end()) - T.begin();
    // nesse caso, pivot = ponto com menor x, depois menor y
    pivot = T[idx];
    swap(T[0], T[idx]);
    sort(++T.begin(), T.end(), radial_lt);

    /*(1)*/ int k;
    for (k = n - 2; k >= 0 && ccw(T[0], T[n - 1], T[k]) == 0; k--);
    reverse((k + 1) + all(T)); /*(1)*/

    // troque <= por < para manter pontos colineares na borda
    for (int i = 0; i < T.size(); i++) {
        while (tam > 1 && ccw(U[tam - 2], U[tam - 1], T[i]) <= 0)
            U.pop_back(), tam--;
        U.pb(T[i]);
        tam++;
    }
    return U;
}

```

## 4.3 Minimum Enclosing Circle

```
// Finds a circle of the minimum area enclosing a 2D point set.
typedef pair<point, double> circle; // {ponto,raio}
bool in_circle(circle C, point p) { // ponto dentro de circulo?
    return cmp(abs(p - C.first), C.second) <= 0;
}
// menor circulo que engloba o triangulo (P,Q,R)
point circumcenter(point p, point q, point r) {
    point a = p - r, b = q - r, c, ret;
    c = point(dot(a, p + r), dot(b, q + r)) * 0.5;
    ret = point(cross(c, point(a.y, b.y)), cross(point(a.x, b.x), c)) /
        cross(a, b);
    return ret;
}
circle spanning_circle(const vector<point> &T) {
    int n = T.size();
    random_shuffle(all(T));
    circle C(point(), -INF);
    rep(i, 0, n) if (!in_circle(C, T[i])) {
        C = circle(T[i], 0);
        rep(j, 0, i) if (!in_circle(C, T[j])) {
            C = circle((T[i] + T[j]) / 2, abs(T[i] - T[j]) / 2);
            rep(k, 0, j) if (!in_circle(C, T[k])) {
                point O = circumcenter(T[i], T[j], T[k]);
                C = circle(O, abs(O - T[k]));
            }
        }
    }
    return C;
}
```

## 5 Geometria 3D

### 5.1 Geometria 3D Library

```
#define LINE 0
#define SEGMENT 1
#define RAY 2
int sgn(double x) { return (x > EPS) - (x < -EPS); }

#define vec(ini, fim) (fim - ini)
struct PT {
    double x, y, z;
    PT() { x = y = z = 0; }
    PT(double x, double y, double z) : x(x), y(y), z(z) {}
    PT operator+(PT q) { return PT(x + q.x, y + q.y, z + q.z); }
    PT operator-(PT q) { return PT(x - q.x, y - q.y, z - q.z); }
    PT operator*(double d) { return PT(x * d, y * d, z * d); }
    PT operator/(double d) { return PT(x / d, y / d, z / d); }
    double dist2() const { return x * x + y * y + z * z; }
    double dist() const { return sqrt(dist2()); }
    bool operator==(const PT &a) const {
        return fabs(x - a.x) < EPS && fabs(y - a.y) < EPS &&
            fabs(z - a.z) < EPS;
    }
};
```

```
double dot(PT A, PT B) { return A.x * B.x + A.y * B.y + A.z * B.z; }
PT cross(PT A, PT B) {
    return PT(A.y * B.z - A.z * B.y, A.z * B.x - A.x * B.z,
        A.x * B.y - A.y * B.x);
}
bool collinear(PT A, PT B, PT C) {
    return sgn(cross(B - A, C - A)) == 0;
}

inline double det(double a, double b, double c, double d) {
    return a * d - b * c;
}
inline double det(double a11, double a12, double a13, double a21,
    double a22, double a23, double a31, double a32,
    double a33) {
    return a11 * det(a22, a23, a32, a33) -
        a12 * det(a21, a23, a31, a33) + a13 * det(a21, a22, a31,
        a32);
}
inline double det(const PT &a, const PT &b, const PT &c) {
    return det(a.x, a.y, a.z, b.x, b.y, b.z, c.x, c.y, c.z);
}

// tamanho do vetor A
double norma(PT A) { return sqrt(dot(A, A)); }

// distancia^2 de (a->b)
double distSq(PT a, PT b) { return dot(a - b, a - b); }

// Projeta vetor A sobre o vetor B
PT project(PT A, PT B) { return B * dot(A, B) / dot(B, B); }

// Verifica se existe interseccao de segmentos
// (assumir que [A,B] e [C,D] sao coplanares)
bool seg_intersect(PT A, PT B, PT C, PT D) {
    return cmp(dot(cross(A - B, C - B), cross(A - B, D - B))) <= 0 &&
        cmp(dot(cross(C - D, A - D), cross(C - D, B - D))) <= 0;
}

// square distance between point and line, ray or segment
double ptLineDistSq(PT s1, PT s2, PT p, int type) {
    double pd2 = distSq(s1, s2);
    PT r;
    if (pd2 == 0)
        r = s1;
    else {
        double u = dot(p - s1, s2 - s1) / pd2;
        r = s1 + (s2 - s1) * u;
        if (type != LINE && u < 0.0)
            r = s1;
        if (type == SEGMENT && u > 1.0)
            r = s2;
    }
    return distSq(r, p);
}

// Distancia de ponto P ao segmento [A,B]
double dist_point_seg(PT P, PT A, PT B) {
    PT PP = A + project(P - A, B - A);
```

```

    if (cmp(norma(A - PP) + norma(PP - B), norma(A - B)) == 0)
        return norma(P - PP); // distance point-line!
    else
        return min(norma(P - A), norma(P - B));
}

// Distance between lines ab and cd. TODO: Test this
double lineLineDistance(PT a, PT b, PT c, PT d) {
    PT v1 = b - a;
    PT v2 = d - c;
    PT cr = cross(v1, v2);
    if (dot(cr, cr) < EPS) {
        PT proj = v1 * (dot(v1, c - a) / dot(v1, v1));
        return sqrt(dot(c - a - proj, c - a - proj));
    } else {
        PT n = cr / sqrt(dot(cr, cr));
        PT p = dot(n, c - a);
        return sqrt(dot(p, p));
    }
}

// Menor distancia do segmento [A,B] ao segmento [C,D] (lento*)
#define dps dist_point_seg
double dist_seg_seg(PT A, PT B, PT C, PT D) {
    PT E = project(A - D, cross(B - A, D - C));
    // distance between lines!
    if (seg_intersect(A, B, C + E, D + E)) {
        return norma(E);
    } else {
        double dA = dps(A, C, D), dB = dps(B, C, D);
        double dC = dps(C, A, B), dD = dps(D, A, B);
        return min(min(dA, dB), min(dC, dD));
    }
}

// Menor distancia do segmento [A,B] ao segmento [C,D] (rapido*)
double dist_seg_seg2(PT A, PT B, PT C, PT D) {
    PT u(B - A), v(D - C), w(A - C);
    double a = dot(u, u), b = dot(u, v);
    double c = dot(v, v), d = dot(u, w), e = dot(v, w);

    double DD = a * c - b * b;
    double sc, sN, sD = DD;
    double tc, tN, tD = DD;
    if (DD < EPS) {
        sN = 0, sD = 1, tN = e, tD = c;
    } else {
        sN = (b * e - c * d);
        tN = (a * e - b * d);
        if (sN < 0) {
            sN = 0, tN = e, tD = c;
        } else if (sN > sD) {
            sN = sD, tN = e + b, tD = c;
        }
    }
    if (tN < 0) {
        tN = 0;
        if (-d < 0)
            sN = 0;
        else if (-d > a)
            sN = sD;
    }
}

```

```

    else {
        sN = -d;
        sD = a;
    }
} else if (tN > tD) {
    tN = tD;
    if ((-d + b) < 0)
        sN = 0;
    else if (-d + b > a)
        sN = sD;
    else {
        sN = -d + b;
        sD = a;
    }
}
sc = fabs(sN) < EPS ? 0 : sN / sD;
tc = fabs(tN) < EPS ? 0 : tN / tD;
PT dP = w + (u * sc) - (v * tc);
return norma(dP);
}

// Distancia de Ponto a Triangulo, dps = dist_point_seg
double dist_point_tri(PT P, PT A, PT B, PT C) {
    PT N = cross(B - A, C - A);
    PT PP = P - project(P - A, N);
    PT R1, R2, R3;
    R1 = cross(B - A, PP - A);
    R2 = cross(C - B, PP - B);
    R3 = cross(A - C, PP - C);

    if (cmp(dot(R1, R2)) >= 0 && cmp(dot(R2, R3)) >= 0 &&
        cmp(dot(R3, R1)) >= 0) {
        return norma(P - PP);
    } else {
        return min(dps(P, A), min(dps(P, B), min(dps(P, C), dps(P, A, C))));
    }
}

// compute a, b, c, d such that all points lie on ax + by + cz = d.
// TODO: test this
void planeFromPts(PT p1, PT p2, PT p3, double &a, double &b, double
    &c,
    double &d) {
    PT normal = cross(p2 - p1, p3 - p1);
    a = normal.x;
    b = normal.y;
    c = normal.z;
    d = -a * p1.x - b * p1.y - c * p1.z;
}

// project point onto plane. TODO: test this
PT ptPlaneProj(PT p, double a, double b, double c, double d) {
    double l =
        (a * p.x + b * p.y + c * p.z + d) / (a * a + b * b + c * c);
    return PT(p.x - a * l, p.y - b * l, p.z - c * l);
}

// distance from point p to plane ax + by + cz + d = 0
double ptPlaneDist(PT p, double a, double b, double c, double d) {
    return fabs(a * p.x + b * p.y + c * p.z + d) /
        sqrt(a * a + b * b + c * c);
}

```

```

}

// distance between parallel planes  $ax + by + cz + d1 = 0$  and
//  $ax + by + cz + d2 = 0$ 
double planePlaneDist(double a, double b, double c, double d1,
                      double d2) {
    return fabs(d1 - d2) / sqrt(a * a + b * b + c * c);
}

// Volume de Tetraedro
double signedTetrahedronVol(PT A, PT B, PT C, PT D) {
    double A11 = A.x - B.x;
    double A12 = A.x - C.x;
    double A13 = A.x - D.x;
    double A21 = A.y - B.y;
    double A22 = A.y - C.y;
    double A23 = A.y - D.y;
    double A31 = A.z - B.z;
    double A32 = A.z - C.z;
    double A33 = A.z - D.z;
    double det = A11 * A22 * A33 + A12 * A23 * A31 + A13 * A21 * A32 -
                 A11 * A23 * A32 - A12 * A21 * A33 - A13 * A22 * A31;
    return det / 6;
}

// Parameter is a vector of vectors of points - each interior vector
// represents the 3 points that make up 1 face, in any order.
// Note: The polyhedron must be convex, with all faces given as
// triangles.
double polyhedronVol(vector<vector<PT>> poly) {
    int i, j;
    PT cent(0, 0, 0);
    for (i = 0; i < poly.size(); i++)
        for (j = 0; j < 3; j++)
            cent = cent + poly[i][j];
    cent = cent * (1.0 / (poly.size() * 3));
    double v = 0;
    for (i = 0; i < poly.size(); i++)
        v += fabs(signedTetrahedronVol(cent, poly[i][0], poly[i][1],
                                       poly[i][2]));
    return v;
}

// Outras implementacoes [Usa struct PT]

struct line { // reta definida por um ponto p e direcao v
    PT p, v;
    line();
    line(const PT &p, const PT &v) : p(p), v(v) { assert(!(v == PT())); }
    bool on(const PT &pt) const { return cross(pt - p, v) == PT(); }
};

struct plane {
    PT n;
    double d;
    plane() : d(0) {}
    plane(const PT &p1, const PT &p2, const PT &p3) {
        n = cross(p2 - p1, p3 - p1);
        d = -dot(n, p1);
        assert(side(p1) == 0);
        assert(side(p2) == 0);
    }

```

```

    assert(side(p3) == 0);
}
int side(const PT &p) const { return sgn(dot(n, p) + d); }
};

```

// intersecao de retas

```

int intersec(const line &l1, const line &l2, PT &res) {
    assert(!(l1.v == PT()));
    assert(!(l2.v == PT()));
    if (cross(l1.v, l2.v) == PT()) {
        if (cross(l1.v, l1.p - l2.p) == PT())
            return 2; // same
        return 0; // parallel
    }
    PT n = cross(l1.v, l2.v);
    PT p = l2.p - l1.p;
    if (sgn(dot(n, p)))
        return 0; // skew
    double t;
    if (sgn(n.x))
        t = (p.y * l2.v.z - p.z * l2.v.y) / n.x;
    else if (sgn(n.y))
        t = (p.z * l2.v.x - p.x * l2.v.z) / n.y;
    else if (sgn(n.z))
        t = (p.x * l2.v.y - p.y * l2.v.x) / n.z;
    else
        assert(false);
    res = l1.p + l1.v * t;
    assert(l1.on(res));
    assert(l2.on(res));
    return 1; // intersects
}

```

// distancia entre 2 retas

```

double dist(const line &l1, const line &l2) {
    PT ret = l1.p - l2.p;
    ret = ret - l1.v * (dot(l1.v, ret) / l1.v.dist2());
    PT tmp = l2.v - l1.v * (dot(l1.v, l2.v) / l1.v.dist2());
    if (sgn(tmp.dist2()))
        ret = ret - tmp * (dot(tmp, ret) / tmp.dist2());
    assert(fabs(dot(ret, l1.v)) < eps);
    assert(fabs(dot(ret, tmp)) < eps);
    assert(fabs(dot(ret, l2.v)) < eps);
    return ret.dist();
}

```

// Retorna os dois pontos mais proximos entre l1 e l2

```

void closest(const line &l1, const line &l2, PT &p1, PT &p2) {
    if (cross(l1.v, l2.v) == PT()) {
        p1 = l1.p;
        p2 = l2.p - l1.v * (dot(l1.v, l2.p - l1.p) / l1.v.dist2());
        return;
    }
    PT p = l2.p - l1.p;
    double t1 =
        (dot(l1.v, p) * l2.v.dist2() - dot(l1.v, l2.v) * dot(l2.v, p)) /
        cross(l1.v, l2.v).dist2();
    double t2 =
        (dot(l2.v, l1.v) * dot(l1.v, p) - dot(l2.v, p) * l1.v.dist2()) /
        cross(l2.v, l1.v).dist2();
    p1 = l1.p + l1.v * t1;

```

```

    p2 = l2.p + l2.v * t2;
    assert(l1.on(p1));
    assert(l2.on(p2));
}

// retorna a intersecao de reta com plano [retorna 1 se intersecao for
// pt]
int cross(const line &l, const plane &pl, PT &res) {
    double d = dot(pl.n, l.v);
    if (sgn(d) == 0) {
        return (pl.side(l.p) == 0) ? 2 : 0;
    }
    double t = (-dot(pl.n, l.p) - pl.d) / d;
    res = l.p + l.v * t;
#ifdef DEBUG
    assert(pl.side(res) == 0);
#endif
    return 1;
}

bool cross(const plane &p1, const plane &p2, const plane &p3,
           PT &res) {
    double d = det(p1.n, p2.n, p3.n);
    if (sgn(d) == 0) {
        return false;
    }
    PT px(p1.n.x, p2.n.x, p3.n.x);
    PT py(p1.n.y, p2.n.y, p3.n.y);
    PT pz(p1.n.z, p2.n.z, p3.n.z);
    PT p(-p1.d, -p2.d, -p3.d);
    res = PT(det(p, py, pz) / d, det(px, p, pz) / d, det(px, py, p) /
            d);
#ifdef DEBUG
    assert(p1.side(res) == 0);
    assert(p2.side(res) == 0);
    assert(p3.side(res) == 0);
#endif
    return true;
}

// retorna reta da intersecao de dois planos
int cross(const plane &p1, const plane &p2, line &res) {
    res.v = cross(p1.n, p2.n);
    if (res.v == PT()) {
        if ((p1.n * (p1.d / p1.n.dist2())) ==
            (p2.n * (p2.d / p2.n.dist2())))
            return 2;
        else
            return 0;
    }
    plane p3;
    p3.n = res.v;
    p3.d = 0;
    bool ret = cross(p1, p2, p3, res.p);
    assert(ret);
    assert(p1.side(res.p) == 0);
    assert(p2.side(res.p) == 0);
    return 1;
}

// testes

```

```

int main() {
    {
        line l;
        l.p = PT(1, 1, 1);
        l.v = PT(1, 0, -1);
        plane p(PT(10, 11, 12), PT(9, 8, 7), PT(1, 3, 2));
        PT res;
        assert(cross(l, p, res) == 1);
    }
    {
        plane p1(PT(1, 2, 3), PT(4, 5, 6), PT(-1, 5, -4));
        plane p2(PT(3, 2, 1), PT(6, 5, 4), PT(239, 17, -42));
        line l;
        assert(cross(p1, p2, l) == 1);
    }
    {
        plane p1(PT(1, 2, 3), PT(4, 5, 6), PT(-1, 5, -4));
        plane p2(PT(1, 2, 3), PT(7, 8, 9), PT(3, -1, 10));
        line l;
        assert(cross(p1, p2, l) == 2);
    }
    {
        plane p1(PT(1, 2, 3), PT(4, 5, 6), PT(-1, 5, -4));
        plane p2(PT(1, 2, 4), PT(4, 5, 7), PT(-1, 5, -3));
        line l;
        assert(cross(p1, p2, l) == 0);
    }

    line l1, l2;
    while (l1.p.load()) {
        l1.v.load();
        l1.v = l1.v - l1.p;
        l2.p.load();
        l2.v.load();
        l2.v = l2.v - l2.p;
        if (l1.v == PT() || l2.v == PT())
            continue;
        PT res;
        int cnt = intersec(l1, l2, res);
        double d = dist(l1, l2);
        if (fabs(d) < eps)
            assert(cnt >= 1);
        else
            assert(cnt == 0);
        PT p1, p2;
        closest(l1, l2, p1, p2);
        assert(fabs((p1 - p2).dist() - d) < eps);
    }
    plane a(PT(1, 0, 0), PT(0, 1, 0), PT(0, 0, 1));
    plane b(PT(-1, 0, 0), PT(0, -1, 0), PT(0, 0, -1));
    line l;
    assert((cross(a, b, l)) == 0);
    return 0;
}

```

## 6 Grafos

### 6.1 Topological Sort

```
// Ordenacao topologia baseado em BFS. Ideia: Processar os vertices
// que nao tem aresta chegando neles. Apos processar, remover as
// arestas dele para seus vizinhos. Os vizinhos que nao tiverem mais
// arestas chegando sao inseridos na fila para serem processados
// depois.
#define MAXV 100001
vector<int> adj[MAXV];
vector<int> ordem;
void topo_sort(int N) {
    queue<int> q;
    // para mudar a ordem que os vertices sao processados pode-se se
    // usar uma priority_queue, outra estrutura para ordenar os vertices
    vector<int> in_degree(N, 0);

    rep(i, 0, N) rep(j, 0, adj[i].size()) in_degree[adj[i][j]]++;

    rep(i, 0, N) if (in_degree[i] == 0) q.push(i);
    while (!q.empty()) {
        int u = q.front();
        q.pop();
        ordem.push_back(u);
        rep(i, 0, adj[u].size()) {
            int v = adj[u][i];
            in_degree[v]--;
            if (in_degree[v] == 0)
                q.push(v);
        }
    }
    if (ordem.size() != N) {
        // grafo contem ciclos, nao eh um DAG
    }
}
int main() { return 0; }
```

## 6.2 Dijkstra

```
/*
 * Encontra o custo do menor caminho uma origem para todos os outros
 * vertices do grafo.
 * So pode ser aplicado em grafos que nao possuem ciclos com peso
 * negativo.
 * Em dist[X] ficara armazenado o custo do menor caminho de src ate X;
 * e em pi[X] ficara o vertice anterior a X neste caminho.
 * Exemplo: src ->, ... , pi[ pi[X] ] -> pi[X] -> X
 * Complexidade  $O((V+E)\log(V))$ 
 */

#define MAXV 100000+10 // quantidade maxima de vertices
typedef long long cost_t; // tipo de variavel para o custo da aresta

int V, E;
vector<pair<int, cost_t> > adj[MAXV];
cost_t dist[MAXV];
int pi[MAXV];

void dijkstra(int src) {
    priority_queue< pair<cost_t, int> > PQ;

    memset(dist, INF, sizeof(dist));
    // nao utilize memset se cost_f for double, use um for
```

```
dist[src] = 0;
PQ.push( make_pair(dist[src], src));

while (!PQ.empty()) {
    pair<cost_t, int> top = PQ.top();
    PQ.pop();

    int u = top.second;
    cost_t d = -top.first;

    if (d != dist[u]) continue;

    rep(i, 0, (int)adj[u].size()) {
        int v = adj[u][i].F;
        cost_t cost_uv = adj[u][i].S;

        if (dist[u] + cost_uv < dist[v]) {
            dist[v] = dist[u] + cost_uv;
            pi[v] = u;
            PQ.push( make_pair(-dist[v], v) );
        }
    }
}
```

## 6.3 Floyd-Warshall

```
#define MAXV 401
int adj[MAXV][MAXV], path[MAXV][MAXV];
int n, m; // #vertices, #arestas
// adj[u][v] = custo de {U->V}
// path[u][v] = k :: K vem logo apos U no caminho ate V
void read_graph() {
    memset(adj, INF, sizeof adj); // para menor caminho
    rep(i, 0, n) adj[i][i] = 0; // para menor caminho
    int u, v, w;
    rep(i, 0, m) {
        cin >> u >> v >> w;
        adj[u][v] = w;
        path[u][v] = v;
    }
}
void floyd() {
    rep(k, 0, n) rep(i, 0, n)
        rep(j, 0, n) if (adj[i][k] + adj[k][j] < adj[i][j]) {
            adj[i][j] = adj[i][k] + adj[k][j];
            path[i][j] = path[i][k];
        }
}
vector<int> findPath(int s, int d) {
    vector<int> Path;
    Path.pb(s);
    while (s != d) {
        s = path[s][d];
        Path.pb(s);
    }
    return Path;
}
```

```

/*Aplicacoes:
1-Encontrar o fecho transitivo (saber se U consegue visitar V)
.: adj[u][v] |= (adj[u][k] & adj[k][v]);
   (inicializar adj com 0)

2-Minimizar a maior aresta do caminho entre U e V
.: adj[u][v] = min(adj[u][v], max(adj[u][k], adj[k][v]));
   (inicializar adj com INF)

3-Maximizar a menor aresta do caminho entre U e V
.: adj[u][v] = max(adj[u][v], min(adj[u][k], adj[k][u]));
   (inicializar adj com -INF)*/

int main() { return 0; }

```

## 6.4 Bellman-Ford

```

// Menor custo de uma origem s para todos vertices em O(V^3).
// bellman() retorna FALSE se o grafo tem ciclo com custo negativo.
// dist[v] contem o menor custo de s ate v.
#define MAXV 400

// Vertices indexados em 0.
int V, E; // #vertices, #arestas
vector<ii> adj[MAXV];
ll dist[MAXV];

bool bellman(int s) {
    rep(i, 0, V) dist[i] = INF;
    dist[s] = 0;
    rep(i, 0, V - 1) rep(u, 0, V) {
        rep(j, 0, adj[u].size()) {
            int v = adj[u][j].F, duv = adj[u][j].S;
            dist[v] = min(dist[v], dist[u] + duv);
        }
    }
    // verifica se tem ciclo com custo negativo
    rep(u, 0, V) rep(j, 0, adj[u].size()) {
        int v = adj[u][j].F, duv = adj[u][j].S;
        if (dist[v] > dist[u] + duv)
            return false;
    }
    return true;
}

int main() { return 0; }

```

## 6.5 Vértices de Articulação e Pontes

```

#define MAXV 100001
vector<int> adj[MAXV];
int dfs_num[MAXV], dfs_low[MAXV], dfs_parent[MAXV];
int dfscounter, V, dfsRoot, rootChildren, ans;
int articulation[MAXV], articulations;
vector<ii> bridges;

void articulationPointAndBridge(int u) {
    dfs_low[u] = dfs_num[u] = dfscounter++;

```

```

    rep(i, 0, adj[u].size()) {
        int v = adj[u][i];
        if (dfs_num[v] == -1) {
            dfs_parent[v] = u;
            if (u == dfsRoot)
                rootChildren++;

            articulationPointAndBridge(v);
            if (dfs_low[v] >= dfs_num[u])
                articulation[u] = true;
            if (dfs_low[v] > dfs_num[u])
                bridges.pb(mp(u, v));

            dfs_low[u] = min(dfs_low[u], dfs_low[v]);
        } else if (v != dfs_parent[u])
            dfs_low[u] = min(dfs_low[u], dfs_num[v]);
    }
}

int main() {
    // read graph
    dfscounter = 0;
    rep(i, 0, V) {
        dfs_low[i] = dfs_parent[i] = articulation[i] = 0;
        dfs_num[i] = -1;
    }
    articulations = 0;
    bridges.clear();
    rep(i, 0, V) if (dfs_num[i] == -1) {
        dfsRoot = i;
        rootChildren = 0;
        articulationPointAndBridge(i);
        articulation[dfsRoot] = (rootChildren > 1);
    }
    printf("#articulations = %d\n", articulations);
    rep(i, 0, V) if (articulation[i]) printf("Vertex %d\n", i);
    printf("#bridges = %d\n", bridges.size());
    rep(i, 0, bridges.size())
        printf("Bridge %d<->%d\n", bridges[i].F, bridges[i].S);
    return 0;
}

```

## 6.6 Tarjan

```

#define MAXV 100010
vector<int> adj[MAXV];
int V;
int dfs_num[MAXV], dfs_low[MAXV], vis[MAXV], SCC[MAXV];
int dfsCounter, numSCC;
vector<int> S; // global variables

void tarjanSCC(int u) {
    dfs_low[u] = dfs_num[u] = dfsCounter++; // dfs_low[u] <= dfs_num[u]
    S.push_back(u); // stores u in a vector based on order of
                    // visitation
    vis[u] = 1;
    rep(i, 0, adj[u].size()) {
        int v = adj[u][i];
        if (dfs_num[v] == -1)
            tarjanSCC(v);
    }
}

```



```

    if (vis[v]) // condition for update
        dfs_low[u] = min(dfs_low[u], dfs_low[v]);
}
if (dfs_low[u] ==
    dfs_num[u]) { // if this is a root (start) of an SCC
    while (true) {
        int v = S.back();
        S.pop_back();
        vis[v] = 0;
        SCC[v] = numSCC; // wich SCC this vertex belong
        if (u == v)
            break;
    }
    numSCC++;
}
}

int main() {
    // read graph
    rep(i, 0, V) {
        dfs_num[i] = -1;
        dfs_low[i] = vis[i] = 0;
        SCC[i] = -i;
    }
    dfsCounter = numSCC = 0;
    rep(i, 0, V) if (dfs_num[i] == -1) tarjanSCC(i);
    rep(i, 0, V) printf("vertice %d, componente %d\n", i, SCC[i]);
    return 0;
}

```

## 6.7 Kosaraju

```

// Encotra componentes conexos. Mesmo que Tarjan
#define MAXV 100000
#define DFS_WHITE 0
vector<int> adj[2][MAXV]; // adj[0][] original, adj[1][] transposto
vector<int> S, dfs_num;
int N, numSCC, SCC[MAXV];

void Kosaraju(int u, int t, int comp) {
    dfs_num[u] = 1;
    if (t == 1)
        SCC[u] = comp;
    for (int j = 0; j < (int)adj[t][u].size(); j++) {
        int v = adj[t][u][j];
        if (dfs_num[v] == DFS_WHITE)
            Kosaraju(v, t, comp);
    }
    S.push_back(u);
}

void doit() { // chamar na main
    S.clear();
    dfs_num.assign(N, DFS_WHITE);
    for (int i = 0; i < N; i++)
        if (dfs_num[i] == DFS_WHITE)
            Kosaraju(i, 0, -1);
    numSCC = 0;
    dfs_num.assign(N, DFS_WHITE);
    for (int i = N - 1; i >= 0; i--)
        if (dfs_num[S[i]] == DFS_WHITE) {

```

```

            Kosaraju(S[i], 1, numSCC);
            numSCC++;
        }
    }
    printf("There are %d SCCs\n", numSCC);
}

int main() { return 0; }

```

## 6.8 2-Sat

```

#define MAXV 100001
// 2-sat -Codigo do problema X-Mart
// vertices indexado em 1
vector<int> adj[2 * MAXV];
vector<int> radj[2 * MAXV];
int seen[2 * MAXV], comp[2 * MAXV], order[2 * MAXV], ncomp, norder;
int N; // #variaveis
int n; // #vertices

#define NOT(x) ((x <= N) ? (x + N) : (x - N))
#define quero 1
void add_edge(int a, int b, int opcao) {
    if (a > b)
        swap(a, b);
    if (b == 0)
        return;
    if (a == 0) {
        if (opcao == quero)
            adj[NOT(b)].pb(b);
        else
            adj[b].pb(NOT(b));
    } else { // normal...
        if (opcao == quero) {
            adj[NOT(a)].pb(b);
            adj[NOT(b)].pb(a);
        } else {
            a = NOT(a);
            b = NOT(b);
            adj[NOT(a)].pb(b);
            adj[NOT(b)].pb(a);
        }
    }
}

void init() {
    rep(i, 0, n + 1) {
        adj[i].clear();
        radj[i].clear();
    }
}

void dfs1(int u) {
    seen[u] = 1;
    rep(i, 0, adj[u].size()) if (!seen[adj[u][i]]) dfs1(adj[u][i]);
    order[norder++] = u;
}

void dfs2(int u) {
    seen[u] = 1;
    rep(i, 0, radj[u].size()) if (!seen[radj[u][i]]) dfs2(radj[u][i]);
    comp[u] = ncomp;
}

void strongly_connected_components() {

```

```

rep(v, 1, n + 1) rep(i, 0, (int)adj[v].size())
    radj[adj[v][i]].pb(v);

norder = 0;
memset(seen, 0, sizeof seen);
rep(v, 1, n + 1) if (!seen[v]) dfs1(v);

ncomp = 0;
memset(seen, 0, sizeof seen);
for (int i = n - 1, u = order[n - 1]; i >= 0; u = order[--i])
    if (!seen[u]) {
        dfs2(u);
        ncomp++;
    }
}

bool sat2() {
    strongly_connected_components();
    rep(i, 1, n + 1) if (comp[i] == comp[NOT(i)]) return false;
    return true;
}

int main() {
    int Clientes;
    while (cin >> Clientes >> N) {
        if (Clientes == 0 && N == 0)
            break;
        n = 2 * N;
        init();
        int u, v;
        rep(i, 0, Clientes) {
            scanf("%d %d", &u, &v);
            add_edge(u, v, quero);
            scanf("%d %d", &u, &v);
            add_edge(u, v, !quero);
        }
        sat2() ? printf("yes\n") : printf("no\n");
    }
    return 0;
}

```

## 6.9 LCA

*/\*Lowest Common Ancestor (LCA) entre dois vertices A, B de uma arvore. LCA(A,B) = ancestral mais proximo de A adj B. O codigo abaixo tambem calcula a menor aresta do caminho entre A adj B. Para saber quantas arestas tem entre A adj B basta fazer: level[A]+level[B]-2\*level[lca(A,B)] Pode-se modificar para retorna a distancia entre A adj B. Como usar: (1) ler a arvore em adj[] adj W[], chamar doit(raiz), passando a raiz da arvore. Indexar em 0 os vertices (2) A funcao retorna o LCA adj a menor aresta entre A adj B. \*/*

```

#define MAXV 101000
const int maxl = 20; // profundidade maxima 2^(maxl) > MAXV
int pai[MAXV][maxl + 1]; // pai[v][i] = pai de v subindo 2^i arestas
int dist[MAXV][maxl + 1]; // dist[v][i] = menor aresta de v subindo
                          // 2^i arestas
int level[MAXV]; // level[v] = #arestas de v ate a raiz

int N, M; // numero de vertices adj arestas

```

```

vector<pair<int, int>> adj[MAXV]; // {v,custo}

void dfs(int v, int p, int peso) {
    level[v] = level[p] + 1;
    pai[v][0] = p;
    dist[v][0] = peso; // aresta de v--pai[v]
    for (int i = 1; i <= maxl; i++) {
        pai[v][i] = pai[pai[v][i - 1]][i - 1]; // subindo 2^i arestas
        dist[v][i] = min(dist[v][i - 1], dist[pai[v][i - 1]][i - 1]);
    }
    rep(i, 0, adj[v].size()) {
        int viz = adj[v][i].F;
        int cost = adj[v][i].S;
        if (viz == p)
            continue;
        dfs(viz, v, cost);
    }
}

void doit(int root) {
    level[root] = 0;
    for (int i = 0; i <= maxl; i++)
        pai[root][i] = root, dist[root][i] = INF;
    rep(i, 0, adj[root].size()) {
        int viz = adj[root][i].F;
        int cost = adj[root][i].S;
        dfs(viz, root, cost);
    }
}

pair<int, int> lca(int a, int b) {
    int menor = INF; // valor da menor aresta do caminho a->b
    if (level[a] < level[b])
        swap(a, b);

    for (int i = maxl; i >= 0; i--) {
        if (level[pai[a][i]] >= level[b]) {
            menor = min(menor, dist[a][i]);
            a = pai[a][i];
        }
    }
    if (a != b) {
        for (int i = maxl; i >= 0; i--) {
            if (pai[a][i] != pai[b][i]) {
                menor = min(menor, min(dist[a][i], dist[b][i]));
                a = pai[a][i];
                b = pai[b][i];
            }
        }
        // ultimo salto
        menor = min(menor, min(dist[a][0], dist[b][0]));
        a = pai[a][0];
        b = pai[b][0];
    }
    return make_pair(a, menor);
}

int main() { return 0; }

```

## 6.10 LCA (Sparse Table)

```

/*
Encontra o lca usando sparse table.
O(NlogN) de pre-processamento
O(1) em cada consulta

Como usar:
    main: level[root] = 0
        dfs(root, root);
    consulta:
        lca(u,v)
*/

typedef pair<int, int> ii;

#define MAXN (1e5 + 1);
#define LOGN (23)

vector<int> adj[MAXN];
int level[MAXN];

vector<int> num;
int f[MAXN];
ii st[4 * MAXN][LOGN];

void dfs(int u, int p) {
    level[u] = level[p] + 1;

    f[u] = num.size();
    num.pb(u);

    rep(i, 0, (int)adj[u].size()) {
        if (adj[u][i] == p)
            continue;
        dfs(adj[u][i], u);
        num.pb(u);
    }
}

ii comb(ii left, ii right) { return min(left, right); }

void SparseTable() {
    rep(i, 0, (int)num.size()) st[i][0] =
        make_pair(level[num[i]], num[i]);

    rep(k, 1, LOGN) for (int i = 0; (i + (1 << k) - 1) <
        (int)num.size(); i++) st[i][k] =
        comb(st[i][k - 1], st[i + (1 << (k - 1))][k - 1]);
}

int lca(int u, int v) {
    int l = f[u];
    int r = f[v];

    int k = log2(r - l + 1);
    return comb(st[l][k], st[r - (1 << k) + 1][k]).second;
}

```

## 6.11 Maximum Bipartite Matching

```

// Encontra o casamento bipartido maximo. Set de vertices X e Y.
// x = [0,X-1], y = [0,Y-1]. match[y] = x - contem quem esta casado
// com y. Teorema de Konig - Num grafo bipartido, o matching eh igual
// ao minimum vertex cover. Complexidade O(nm)
#define MAXV 1000
vector<int> adj[MAXV];
int match[MAXV], V, X, Y;
bool vis[MAXV];

int aug(int v) {
    if (vis[v])
        return 0;
    vis[v] = true;
    rep(i, 0, adj[v].size()) {
        int r = adj[v][i];
        if (match[r] == -1 || aug(match[r])) {
            match[r] = v; // augmenting path
            return 1;
        }
    }
    return 0;
}

int matching(int X, int Y) {
    int V = X + Y;
    rep(i, 0, V) match[i] = -1;
    int mcbm = 0;
    rep(i, 0, X) {
        rep(j, 0, X) vis[j] = false;
        mcbm += aug(i);
    }
    return mcbm;
}

int main() { return 0; }

```

## 6.12 Hopcroft Karp - Maximum Bipartite Matching (UNI-FEI)

```

/*Encontra o casamento bipartido maximo em O(sqrt(V)*E)
1) Chamar init(L,R) #vertices da esquerda, #vertices da direita
2) Usar addEdge(Li,Ri) para adicionar a aresta Li -> Ri
3) maxMatching() retorna o casamento maximo.
   matching[Rj] -> armazena Li */

#define MAXN1 3010
#define MAXN2 3010
#define MAXM 6020
int n1, n2, edges, last[MAXN1], pre[MAXM], head[MAXM];
int matching[MAXN2], dist[MAXN1], Q[MAXN1];
bool used[MAXN1], vis[MAXN1];

void init(int L, int R) {
    n1 = L, n2 = R;
    edges = 0;
    fill(last, last + n1, -1);
}

```

```

}
void addEdge(int u, int v) {
    head[edges] = v;
    pre[edges] = last[u];
    last[u] = edges++;
}
void bfs() {
    fill(dist, dist + n1, -1);
    int sizeQ = 0;
    for (int u = 0; u < n1; ++u) {
        if (!used[u]) {
            Q[sizeQ++] = u;
            dist[u] = 0;
        }
    }
    for (int i = 0; i < sizeQ; i++) {
        int u1 = Q[i];
        for (int e = last[u1]; e >= 0; e = pre[e]) {
            int u2 = matching[head[e]];
            if (u2 >= 0 && dist[u2] < 0) {
                dist[u2] = dist[u1] + 1;
                Q[sizeQ++] = u2;
            }
        }
    }
}
bool dfs(int u1) {
    vis[u1] = true;
    for (int e = last[u1]; e >= 0; e = pre[e]) {
        int v = head[e];
        int u2 = matching[v];
        if (u2 < 0 || !vis[u2] && dist[u2] == dist[u1] + 1 && dfs(u2)) {
            matching[v] = u1;
            used[u1] = true;
            return true;
        }
    }
    return false;
}
int maxMatching() {
    fill(used, used + n1, false);
    fill(matching, matching + n2, -1);
    for (int res = 0;;) {
        bfs();
        fill(vis, vis + n1, false);
        int f = 0;
        for (int u = 0; u < n1; ++u)
            if (!used[u] && dfs(u))
                ++f;
        if (!f)
            return res;
        res += f;
    }
}

int main() { return 0; }

```

## 6.13 Network Flow (lento)

// Ford-Fulkerson para fluxo maximo

```

#define MAXV 250
vector<int> edge[MAXV];
int cap[MAXV][MAXV];
bool vis[MAXV];

void init() {
    rep(i, 0, MAXV) edge[i].clear();
    memset(cap, 0, sizeof cap);
}

void add(int a, int b, int cap_ab, int cap_ba) {
    edge[a].pb(b), edge[b].pb(a);
    cap[a][b] += cap_ab, cap[b][a] += cap_ba;
}

int dfs(int src, int snk, int fl) {
    if (vis[src])
        return 0;
    if (snk == src)
        return fl;
    vis[src] = 1;
    rep(i, 0, edge[src].size()) {
        int v = edge[src][i];
        int x = min(fl, cap[src][v]);
        if (x > 0) {
            x = dfs(v, snk, x);
            if (!x)
                continue;
            cap[src][v] -= x;
            cap[v][src] += x;
            return x;
        }
    }
    return 0;
}

int flow(int src, int snk) {
    int ret = 0;
    while (42) {
        memset(vis, 0, sizeof vis);
        int delta = dfs(src, snk, 1 << 30);
        if (!delta)
            break;
        ret += delta;
    }
    return ret;
}

int main() { return 0; }

```

## 6.14 Network Flow - Dinic

```

// Dinic para fluxo maximo
// Grafo indexado em 1
// Inicializar maxN, maxE.
// Chamar init() com #nos, source e sink. Montar o grafo chamando
// add(a,b,c1,c2), sendo c1 cap. de a->b e c2 cap. de b->a
#define FOR(i, a, b) for (int i = a; i <= b; i++)
#define SET(c, v) memset(c, v, sizeof c)
const int maxN = 5000;
const int maxE = 70000;

```

```

const int inf = 1000000005;

int nnode, nedge, src, snk;
int Q[maxN], pro[maxN], fin[maxN], dist[maxN];
int flow[maxE], cap[maxE], to[maxE], prox[maxE];

void init(int _nnode, int _src, int _snk) {
    nnode = _nnode, nedge = 0, src = _src, snk = _snk;
    FOR(i, 1, nnode) fin[i] = -1;
}

void add(int a, int b, int c1, int c2) {
    to[nedge] = b, cap[nedge] = c1, flow[nedge] = 0,
    prox[nedge] = fin[a], fin[a] = nedge++;
    to[nedge] = a, cap[nedge] = c2, flow[nedge] = 0,
    prox[nedge] = fin[b], fin[b] = nedge++;
}

bool bfs() {
    SET(dist, -1);
    dist[src] = 0;
    int st = 0, en = 0;
    Q[en++] = src;
    while (st < en) {
        int u = Q[st++];
        for (int e = fin[u]; e >= 0; e = prox[e]) {
            int v = to[e];
            if (flow[e] < cap[e] && dist[v] == -1) {
                dist[v] = dist[u] + 1;
                Q[en++] = v;
            }
        }
    }
    return dist[snk] != -1;
}

int dfs(int u, int fl) {
    if (u == snk)
        return fl;
    for (int &e = pro[u]; e >= 0; e = prox[e]) {
        int v = to[e];
        if (flow[e] < cap[e] && dist[v] == dist[u] + 1) {
            int x = dfs(v, min(cap[e] - flow[e], fl));
            if (x > 0) {
                flow[e] += x, flow[e ^ 1] -= x;
                return x;
            }
        }
    }
    return 0;
}

ll dinic() {
    ll ret = 0;
    while (bfs()) {
        FOR(i, 1, nnode) pro[i] = fin[i];
        while (true) {
            int delta = dfs(src, inf);
            if (!delta)
                break;
            ret += delta;
        }
    }
}

```

```

    }
    return ret;
}

int main() { return 0; }

```

## 6.15 Min Cost Max Flow

*// Criar o grafo chamando MCMF g(V), onde g eh o grafo e V a qtde de  
// vertices (indexado em 0). Chamar g.add(u,v,cap,cost) para add a  
// aresta u->v, se for bidirecional, chamar tbm g.add(v,u,cap,cost)*

```

struct MCMF {
    typedef int ctype;
    enum { MAXN = 550, INF = INT_MAX };
    struct Edge {
        int x, y;
        ctype cap, cost;
    };
    vector<Edge> E;
    vector<int> adj[MAXN];
    int N, prev[MAXN];
    ctype dist[MAXN], phi[MAXN];

    MCMF(int NN) : N(NN) {}

    void add(int x, int y, ctype cap, ctype cost) { // cost >= 0
        Edge e1 = {x, y, cap, cost}, e2 = {y, x, 0, -cost};
        adj[e1.x].push_back(E.size());
        E.push_back(e1);
        adj[e2.x].push_back(E.size());
        E.push_back(e2);
    }

    void mcmf(int s, int t, ctype &flowVal, ctype &flowCost) {
        int x;
        flowVal = flowCost = 0;
        memset(phi, 0, sizeof(phi));
        while (true) {
            for (x = 0; x < N; x++)
                prev[x] = -1;
            for (x = 0; x < N; x++)
                dist[x] = INF;
            dist[s] = prev[s] = 0;

            set<pair<ctype, int>> Q;
            Q.insert(make_pair(dist[s], s));
            while (!Q.empty()) {
                x = Q.begin()->second;
                Q.erase(Q.begin());
                for (vector<int>::iterator it = adj[x].begin();
                    it != adj[x].end(); it++) {
                    const Edge &e = E[*it];
                    if (e.cap <= 0)
                        continue;
                    ctype cc = e.cost + phi[x] - phi[e.y];
                    if (dist[x] + cc < dist[e.y]) {
                        Q.erase(make_pair(dist[e.y], e.y));
                        dist[e.y] = dist[x] + cc;
                        prev[e.y] = *it;
                    }
                }
            }
        }
    }
}

```

```

        Q.insert(make_pair(dist[e.y], e.y));
    }
}
if (prev[t] == -1)
    break;

ctype z = INF;
for (x = t; x != s; x = E[prev[x]].x)
    z = min(z, E[prev[x]].cap);
for (x = t; x != s; x = E[prev[x]].x) {
    E[prev[x]].cap -= z;
    E[prev[x] ^ 1].cap += z;
}
flowVal += z;
flowCost += z * (dist[t] - phi[s] + phi[t]);
for (x = 0; x < N; x++)
    if (prev[x] != -1)
        phi[x] += dist[x];
}
}
};

int main() { return 0; }

```

## 6.16 Min Cost Max Flow (Stefano)

```

#define MAX_V 2003
#define MAX_E 2 * 3003
// Inicializar MAX_V e MAX_E corretamente. Chamar init(_V) com a qtde
// de vertices (indexado em 0) mesmo que seja bidirecional. Adicionar
// as arestas duas vezes no main(). Complexiade (rapido)

typedef int cap_type;
typedef long long cost_type;
const cost_type inf = LLONG_MAX;

int V, E, pre[MAX_V], last[MAX_V], to[MAX_E], nex[MAX_E];
bool visited[MAX_V];
cap_type flowVal, cap[MAX_E];
cost_type flowCost, cost[MAX_E], dist[MAX_V], pot[MAX_V];

void init(int _V) {
    memset(last, -1, sizeof(last));
    V = _V;
    E = 0;
}

void add_edge(int u, int v, cap_type _cap, cost_type _cost) {
    to[E] = v, cap[E] = _cap;
    cost[E] = _cost, nex[E] = last[u];
    last[u] = E++;
    to[E] = u, cap[E] = 0;
    cost[E] = -_cost, nex[E] = last[v];
    last[v] = E++;
}

// only if there is initial negative cycle
void BellmanFord(int s, int t) {
    bool stop = false;

```

```

    for (int i = 0; i < V; ++i)
        dist[i] = inf;
    dist[s] = 0;

    for (int i = 1; i <= V && !stop; ++i) {
        stop = true;

        for (int j = 0; j < E; ++j) {
            int u = to[j ^ 1], v = to[j];

            if (cap[j] > 0 && dist[u] != inf &&
                dist[u] + cost[j] < dist[v]) {
                stop = false;
                dist[v] = dist[u] + cost[j];
            }
        }
    }

    for (int i = 0; i < V; ++i)
        if (dist[i] != inf)
            pot[i] = dist[i];
}

void mcmf(int s, int t) {
    flowVal = flowCost = 0;
    memset(pot, 0, sizeof(pot));

    BellmanFord(s, t);

    while (true) {
        memset(pre, -1, sizeof(pre));
        memset(visited, false, sizeof(visited));
        for (int i = 0; i < V; ++i)
            dist[i] = inf;

        priority_queue<pair<cost_type, int>>> Q;
        Q.push(make_pair(0, s));
        dist[s] = pre[s] = 0;

        while (!Q.empty()) {
            int aux = Q.top().second;
            Q.pop();

            if (visited[aux])
                continue;
            visited[aux] = true;

            for (int e = last[aux]; e != -1; e = nex[e]) {
                if (cap[e] <= 0)
                    continue;
                cost_type new_dist =
                    dist[aux] + cost[e] + pot[aux] - pot[to[e]];
                if (new_dist < dist[to[e]]) {
                    dist[to[e]] = new_dist;
                    pre[to[e]] = e;
                    Q.push(make_pair(-new_dist, to[e]));
                }
            }
        }

        if (pre[t] == -1)

```

```

        break;

    cap_type f = cap[pre[t]];
    for (int i = t; i != s; i = to[pre[i] ^ 1])
        f = min(f, cap[pre[i]]);
    for (int i = t; i != s; i = to[pre[i] ^ 1]) {
        cap[pre[i]] -= f;
        cap[pre[i] ^ 1] += f;
    }

    flowVal += f;
    flowCost += f * (dist[t] - pot[s] + pot[t]);

    for (int i = 0; i < V; ++i)
        if (pre[i] != -1)
            pot[i] += dist[i];
    }

    int main() { return 0; }

```

## 6.17 Tree Isomorphism

*// Verifica se dado duas arvores, desconsiderando o rotulo dos  
// vertices, elas tem a mesma forma.*

```

typedef vector<int> vi;
#define sz(a) (int)a.size()
#define fst first
#define snd second

struct tree {
    int n;
    vector<vi> adj;
    tree(int n) : n(n), adj(n) {}
    void add_edge(int src, int dst) {
        adj[src].pb(dst);
        adj[dst].pb(src);
    }
    vi centers() {
        vi prev;
        int u = 0;
        for (int k = 0; k < 2; ++k) {
            queue<int> q;
            prev.assign(n, -1);
            q.push(prev[u] = u);
            while (!q.empty()) {
                u = q.front();
                q.pop();
                for (auto i : adj[u]) {
                    if (prev[i] >= 0)
                        continue;
                    q.push(i);
                    prev[i] = u;
                }
            }
        }
        vi path = {u};
        while (u != prev[u])
            path.pb(u = prev[u]);
        int m = sz(path);

```

```

        if (m % 2 == 0)
            return {path[m / 2 - 1], path[m / 2]};
        else
            return {path[m / 2]};
    }
    vector<vi> layer;
    vi prev;
    int levelize(int r) {
        prev.assign(n, -1);
        prev[r] = n;
        layer = {{r}};
        while (true) {
            vi next;
            for (auto u : layer.back()) {
                for (int v : adj[u]) {
                    if (prev[v] >= 0)
                        continue;
                    prev[v] = u;
                    next.pb(v);
                }
            }
            if (next.empty())
                break;
            layer.pb(next);
        }
        return sz(layer);
    }

    bool isomorphic(tree S, int s, tree T, int t) {
        if (S.n != T.n)
            return false;
        if (S.levelize(s) != T.levelize(t))
            return false;
        vector<vi> longcodeS(S.n + 1), longcodeT(T.n + 1);
        vi codeS(S.n), codeT(T.n);
        for (int h = S.layer.size() - 1; h >= 0; h--) {
            map<vi, int> bucket;
            for (int u : S.layer[h]) {
                sort(all(longcodeS[u]));
                bucket[longcodeS[u]] = 0;
            }
            for (int u : T.layer[h]) {
                sort(all(longcodeT[u]));
                bucket[longcodeT[u]] = 0;
            }
            int id = 0;
            for (auto &p : bucket)
                p.snd = id++;
            for (int u : S.layer[h]) {
                codeS[u] = bucket[longcodeS[u]];
                longcodeS[S.prev[u]].pb(codeS[u]);
            }
            for (int u : T.layer[h]) {
                codeT[u] = bucket[longcodeT[u]];
                longcodeT[T.prev[u]].pb(codeT[u]);
            }
        }
        return codeS[s] == codeT[t];
    }
}

```

```

bool isomorphic(tree S, tree T) {
    auto x = S.centers(), y = T.centers();
    if (sz(x) != sz(y))
        return false;
    if (isomorphic(S, x[0], T, y[0]))
        return true;
    return sz(x) > 1 and isomorphic(S, x[1], T, y[0]);
}

int main() {
    int N, u, v;
    cin >> N;
    tree A(N + 2), B(N + 2);
    rep(i, 0, N - 1) {
        scanf("%d %d", &u, &v);
        u--, v--;
        A.add_edge(u, v);
    }
    rep(i, 1, N) {
        scanf("%d %d", &u, &v);
        u--, v--;
        B.add_edge(u, v);
    }
    puts(isomorphic(A, B) ? "S" : "N");
}

```

## 6.18 Stoer Wagner-Minimum Cut (UNIFEI)

```

/*
Retorna o corte minimo do grafo
(Conjunto de arestas que caso seja removido, desconecta o grafo)
Input: n = #vertices, g[i][j] = custo da aresta (i->j)
Output: Retorna o corte minimo
Complexidade: O(N^3)
*/

// Maximum number of vertices in the graph
#define NN 101
// Maximum edge weight (MAXW * NN * NN must fit into an int)
#define MAXW 110

// Adjacency matrix and some internal arrays
int g[NN][NN], v[NN], w[NN], na[NN], n;
bool a[NN];
int stoer_wagner() {
    // init the remaining vertex set
    for (int i = 0; i < n; i++)
        v[i] = i;
    // run Stoer-Wagner
    int best = MAXW * n * n;
    while (n > 1) {
        // initialize the set A and vertex weights
        a[v[0]] = true;
        for (int i = 1; i < n; i++) {
            a[v[i]] = false;
            na[i - 1] = i;
            w[i] = g[v[0]][v[i]];
        }
        // add the other vertices
        int prev = v[0];

```

```

        for (int i = 1; i < n; i++) {
            // find the most tightly connected non-A vertex
            int zj = -1;
            for (int j = 1; j < n; j++)
                if (!a[v[j]] && (zj < 0 || w[j] > w[zj]))
                    zj = j;
            // add it to A
            a[v[zj]] = true;
            // last vertex?
            if (i == n - 1) {
                // remember the cut weight
                best = min(best, w[zj]);

                // merge prev and v[zj]
                for (int j = 0; j < n; j++)
                    g[v[j]][prev] = g[prev][v[j]] += g[v[zj]][v[j]];
                v[zj] = v[--n];
                break;
            }
            prev = v[zj];
            // update the weights of its neighbours
            for (int j = 1; j < n; j++)
                if (!a[v[j]])
                    w[j] += g[v[zj]][v[j]];
        }
    }
    return best;
}

int main() { return 0; }

```

## 6.19 Erdos Gallai (UNIFEI)

```

// Determina se existe um grafo tal que b[i] eh o grau do i-esimo
// vertice. Vertices indexado em 1. Apenas armazenar em b[1...N] e
// chamar EGL()
long long b[100005], n;
long long dmax, dmin, dsum, num_degs[100005];

bool basic_graphical_tests() { // Sort and perform some simple tests
    // on the sequence

    int p = n;
    memset(num_degs, 0, (n + 1) * sizeof(long long));

    dmax = dsum = n = 0;
    dmin = p;
    for (int d = 1; d <= p; d++) {
        if (b[d] < 0 || b[d] >= p)
            return false;
        else if (b[d] > 0) {
            if (dmax < b[d])
                dmax = b[d];
            if (dmin > b[d])
                dmin = b[d];
            dsum = dsum + b[d];
            n++;
            num_degs[b[d]]++;
        }
    }
    if (dsum % 2 || dsum > n * (n - 1))

```



```

    return false;
    return true;
}
bool EGL() {
    long long k, sum_deg, sum_nj, sum_jnj, run_size;

    if (!basic_graphical_tests())
        return false;
    if (n == 0 || 4 * dmin * n >= (dmax + dmin + 1) * (dmax + dmin + 1))
        return true;

    k = sum_deg = sum_nj = sum_jnj = 0;
    for (int dk = dmax; dk >= dmin; dk--) {
        if (dk < k + 1)
            return true;

        if (num_degs[dk] > 0) {
            run_size = num_degs[dk];
            if (dk < k + run_size)
                run_size = dk - k;
            sum_deg += run_size * dk;

            for (int v = 0; v < run_size; v++) {
                sum_nj += num_degs[k + v];
                sum_jnj += (k + v) * num_degs[k + v];
            }
            k += run_size;
            if (sum_deg > k * (n - 1) - k * sum_nj + sum_jnj)
                return false;
        }
    }
    return true;
}
int main() { return 0; }

```

## 6.20 Stable Marriage (UNIFEI)

*/\*Seja um conjunto de m homens e n mulheres, onde cada pessoa tem uma preferencia por outra de sexo oposto. O algoritmo produz o casamento estavel de cada homem com uma mulher. Estavel:*

- Cada homem se casara com uma mulher diferente ( $n \geq m$ )
- Dois casais H1M1 e H2M2 nao serao instaveis.

*Dois casais H1M1 e H2M2 sao instaveis se:*

- H1 prefere M2 ao inves de M1, e
- M1 prefere H2 ao inves de H1.

*Entrada*

(1)  $m = \text{\#homens}$ ,  $n = \text{\#mulheres}$

(2)  $R[x][y] = i$ ,  $i$ : eh a ordem de preferencia do homem  $y$  pela mulher  $x$

*Obs.: Quanto maior o valor de  $i$  menor eh a preferencia do homem  $y$  pela mulher  $x$*

(3)  $L[x][i] = y$ : A mulher  $y$  eh a  $i$ -esima preferencia do homem  $x$

*Obs.:  $0 \leq i \leq n-1$ , quanto menor o valor de  $i$  maior eh a preferencia do homem  $x$  pela mulher  $y$*

*Saida*

$L2R[i]$ : a mulher do homem  $i$  (sempre entre 0 e  $n-1$ )

$R2L[j]$ : o homem da mulher  $j$  (-1 se a mulher for solteira)

*Complexidade  $O(m^2)$*

*\*/*

```

#define MAXM 1000
#define MAXW 1000
int L[MAXM][MAXW];
int R[MAXW][MAXM];
int L2R[MAXM], R2L[MAXW];
int m, n;
int p[MAXM];

void stableMarriage() {
    static int p[MAXM];
    memset(R2L, -1, sizeof(R2L));
    memset(p, 0, sizeof(p));
    for (int i = 0; i < m; ++i) {
        int man = i;
        while (man >= 0) {
            int wom;
            while (42) {
                wom = L[man][p[man]++];
                if (R2L[wom] < 0 || R[wom][man] > R[wom][R2L[wom]])
                    break;
            }
            int hubby = R2L[wom];
            R2L[L2R[man] = wom] = man;
            man = hubby;
        }
    }
}

int main() { return 0; }

```

## 6.21 Hungarian Max Bipartite Matching with Cost (UNIFEI)

*/\*Encontra o casamento bipartido maximo/minimo com peso nas arestas*

*Criar o grafo:*

*Hungarian G(L, R, ehMaximo)*

$L = \text{\#vertices a esquerda}$

$R = \text{\#vertices a direita}$

*ehMaximo = variavel booleana que indica se eh casamento maximo ou minimo*

*Adicionar arestas:*

$G.add\_edge(x, y, \text{peso})$

$x = \text{vertice da esquerda no intervalo } [0, L-1]$

$y = \text{vertice da direita no intervalo } [0, R-1]$

*peso = custo da aresta*

*obs: tomar cuidado com multiplas arestas.*

*Resultado:*

$match\_value = \text{soma dos pesos dos casamentos}$

$pairs = \text{quantidade de pares (x-y) casados}$

$xy[x] = \text{vertice y casado com x}$

$yx[y] = \text{vertice x casado com y}$

*Complexidade do algoritmo:  $O(V^3)$*

*Problemas resolvidos: SCITIES (SPOJ)*

*\*/*

```

struct Hungarian {

```

```

enum { MAXN = 150, INF = 0x3f3f3f3f };
int cost[MAXN][MAXN];
int xy[MAXN], yx[MAXN];
bool S[MAXN], T[MAXN];
int lx[MAXN], ly[MAXN], slack[MAXN], slackx[MAXN], prev[MAXN];
int match_value, pairs;
bool ehMaximo;
int n;

```

```

Hungarian(int L, int R, bool _ehMaximo = true) {
    n = max(L, R);
    ehMaximo = _ehMaximo;
    if (ehMaximo)
        memset(cost, 0, sizeof cost);
    else
        memset(cost, INF, sizeof cost);
}

```

```

void add_edge(int x, int y, int peso) {
    if (!ehMaximo)
        peso *= (-1);
    cost[x][y] = peso;
}

```

```

int solve() {
    match_value = 0;
    pairs = 0;
    memset(xy, -1, sizeof(xy));
    memset(yx, -1, sizeof(yx));
    init_labels();
    augment();
    for (int x = 0; x < n; ++x)
        match_value += cost[x][xy[x]];
    return match_value;
}

```

```

void init_labels() {
    memset(lx, 0, sizeof(lx));
    memset(ly, 0, sizeof(ly));
    for (int x = 0; x < n; ++x)
        for (int y = 0; y < n; ++y)
            lx[x] = max(lx[x], cost[x][y]);
}

```

```

void augment() {
    if (pairs == n)
        return;
    int x, y, root;
    int q[MAXN], wr = 0, rd = 0;
    memset(S, false, sizeof(S));
    memset(T, false, sizeof(T));
    memset(prev, -1, sizeof(prev));
    for (x = 0; x < n; ++x)
        if (xy[x] == -1) {
            q[wr++] = root = x;
            prev[x] = -2;
            S[x] = true;
            break;
        }
    for (y = 0; y < n; ++y) {
        slack[y] = lx[root] + ly[y] - cost[root][y];

```

```

        slackx[y] = root;
    }
    while (true) {
        while (rd < wr) {
            x = q[rd++];
            for (y = 0; y < n; ++y)
                if (cost[x][y] == lx[x] + ly[y] && !T[y]) {
                    if (yx[y] == -1)
                        break;
                    T[y] = true;
                    q[wr++] = yx[y];
                    add(yx[y], x);
                }
            if (y < n)
                break;
        }
        if (y < n)
            break;
        update_labels();
        wr = rd = 0;
        for (y = 0; y < n; ++y)
            if (!T[y] && slack[y] == 0) {
                if (yx[y] == -1) {
                    x = slackx[y];
                    break;
                } else {
                    T[y] = true;
                    if (!S[yx[y]]) {
                        q[wr++] = yx[y];
                        add(yx[y], slackx[y]);
                    }
                }
            }
        if (y < n)
            break;
    }
    if (y < n) {
        ++pairs;
        for (int cx = x, cy = y, ty; cx != -2; cx = prev[cx], cy = ty) {
            ty = xy[cx];
            yx[cy] = cx;
            xy[cx] = cy;
        }
        augment();
    }
}

```

```

void add(int x, int prevx) {
    S[x] = true;
    prev[x] = prevx;
    for (int y = 0; y < n; ++y)
        if (lx[x] + ly[y] - cost[x][y] < slack[y]) {
            slack[y] = lx[x] + ly[y] - cost[x][y];
            slackx[y] = x;
        }
}

```

```

void update_labels() {
    int x, y, delta = INF;
    for (y = 0; y < n; ++y)
        if (!T[y])

```

```

    delta = min(delta, slack[y]);
    for (x = 0; x < n; ++x)
        if (S[x])
            lx[x] -= delta;
    for (y = 0; y < n; ++y)
        if (T[y])
            ly[y] += delta;
    for (y = 0; y < n; ++y)
        if (!T[y])
            slack[y] -= delta;
}

int casouComX(int x) { return xy[x]; }

int casouComY(int y) { return yx[y]; }
};

// O código abaixo resolve o problema scities (Spoj)
int main() {
    int casos;
    cin >> casos;
    while (casos--) {
        int L, R;
        cin >> L >> R;
        Hungarian G(L, R, true);

        int x, y, w, aux[L][R];
        memset(aux, 0, sizeof aux);
        while (scanf("%d %d %d", &x, &y, &w) != EOF) {
            if (x == 0 && y == 0 && w == 0)
                break;
            aux[x - 1][y - 1] += w;
        }
        for (int x = 0; x < L; x++) {
            for (int y = 0; y < R; y++) {
                if (aux[x][y] != 0) {
                    G.add_edge(x, y, aux[x][y]);
                }
            }
        }
        printf("%d\n", G.solve());
    }
    return 0;
}

```

## 6.22 Blossom

```

// Encontra o emparelhamento maximo em um grafo nao direcionado.
// Armazenar em n a quantidade de vertice e em mat[][] as adjacencias.
// edmond(n) retorna o emparelhamento maximo.
typedef vector<int> VI;
typedef vector<vector<int>> VVI;

int mat[205][205], n;

int lf[205];
VVI adj;
VI vis, inactive, match;
int N;

```

```

bool dfs(int x, VI &blossom) {
    if (inactive[x])
        return false;
    int i, y;
    vis[x] = 0;
    for (i = adj[x].size() - 1; i >= 0; i--) {
        y = adj[x][i];
        if (inactive[y])
            continue;
        if (vis[y] == -1) {
            vis[y] = 1;
            if (match[y] == -1 || dfs(match[y], blossom)) {
                match[y] = x;
                match[x] = y;
                return true;
            }
        }
        if (vis[y] == 0 || blossom.size()) {
            blossom.push_back(y);
            blossom.push_back(x);
            if (blossom[0] == x) {
                match[x] = -1;
                return true;
            }
            return false;
        }
    }
    return false;
}

bool augment() {
    VI blossom, mark;
    int i, j, k, s, x;
    for (i = 0; i < N; i++) {
        if (match[i] != -1)
            continue;
        blossom.clear();
        vis = VI(N + 1, -1);
        if (!dfs(i, blossom))
            continue;
        s = blossom.size();
        if (s == 0)
            return true;

        mark = VI(N + 1, -1);
        for (j = 0; j < s - 1; j++) {
            for (k = adj[blossom[j]].size() - 1; k >= 0; k--)
                mark[adj[blossom[j]][k]] = j;
        }

        for (j = 0; j < s - 1; j++) {
            mark[blossom[j]] = -1;
            inactive[blossom[j]] = 1;
        }

        adj[N].clear();
        for (j = 0; j < N; j++) {
            if (mark[j] != -1)
                adj[N].pb(j), adj[j].pb(N);
        }
    }
}

```

```

match[N] = -1;
N++;
if (!augment())
    return false;
N--;

for (j = 0; j < N; j++) {
    if (mark[j] != -1)
        adj[j].pop_back();
}
for (j = 0; j < s - 1; j++) {
    inactive[blossom[j]] = 0;
}

x = match[N];
if (x != -1) {
    if (mark[x] != -1) {
        j = mark[x];
        match[blossom[j]] = x;
        match[x] = blossom[j];
        if (j & 1)
            for (k = j + 1; k < s; k += 2) {
                match[blossom[k]] = blossom[k + 1];
                match[blossom[k + 1]] = blossom[k];
            }
        else
            for (k = 0; k < j; k += 2) {
                match[blossom[k]] = blossom[k + 1];
                match[blossom[k + 1]] = blossom[k];
            }
    }
    return true;
}
return false;
}

int edmond(int n) {
    int i, j, ret = 0;
    N = n;
    adj = VVI(2 * N + 1);
    for (i = 0; i < n; i++) {
        for (j = i + 1; j < n; j++) {
            if (mat[i][j]) {
                adj[i].pb(j);
                adj[j].pb(i);
            }
        }
    }
    match = VI(2 * N + 1, -1);
    inactive = VI(2 * N + 1);
    while (augment())
        ret++;
    return ret;
}

```

## 7 Estruturas de Dados

### 7.1 BIT

```

// Permite realizar operacoes de query e update em um vetor em O(logN)
// Obs: A[] deve ser indexado em 1, nao em 0.
#define MAXN 100001
ll ft[MAXN];
ll A[MAXN];
int N;

// ATUALIZA UM INDICE i, CONSULTA UM INTERVALO (i,j)
// update(i, valor) faz A[i] += valor em log(N)
void update(int i, ll valor) {
    for (; i <= N; i += i & -i)
        ft[i] += valor;
}

// query(i) retorna a soma A[1] + ... + A[i] em log(N)
ll query(int i) {
    ll sum = 0;
    for (; i > 0; i -= i & -i)
        sum += ft[i];
    return sum;
}

// query(i,j) retorna a soma A[i] + A[i+1] + ... + A[j] em log(N)
ll query(int i, int j) { return query(j) - query(i - 1); }

// ATUALIZA UM INTERVALO (i,j), CONSULTA UM ELEMENTO i

// range_update(i,j,valor) faz A[k] += valor, para i <= k <= j em
// log(N) query(i): retorna o valor de A[i] em log(N)
void range_update(int i, int j, ll valor) {
    update(i, valor);
    update(j + 1, -valor);
}

// Just a wrapper function... Returns the value at A[i] after
// range_update( ) calls
ll point_query(int i) {
    return query(i);
}

int main() { return 0; }

```

### 7.2 BIT 2D

```

#define MAXL 3001
#define MAXC 3001
ll ft[MAXL][MAXC];
int L, C;
// update(x,y,v) incrementa v na posicao (x,y)
// .: M[x][y] += v em O(log(N))
void update(int x, int y, int v) {
    for (; x <= L; x += x & -x)

```

```

    for (int yy = y; yy <= C; yy += yy & -yy)
        ft[x][yy] += v;
}

// query(x,y) retorna o somatorio da submatriz definida por
// (1,1)-(x,y) .: sum += M[i][j] para todo 1 <= i <= x e 1 <= j <= y,
// em O(log(N))
ll query(int x, int y) {
    if (x <= 0 || y <= 0)
        return 0;
    ll sum = 0;
    for (; x > 0; x -= x & -x)
        for (int yy = y; yy > 0; yy -= yy & -yy)
            sum += ft[x][yy];
    return sum;
}

// query(x1,y1,x2,y2) retorna o somatorio da submatriz definida por
// (x1,x1) -- (x2,y2) .: sum += M[i][j] para todo x1 <= i <= x2 e
// y1 <= j <= y2, em O(log(N))
ll query(int x1, int y1, int x2, int y2) {
    return query(x2, y2) - query(x2, y1 - 1) - query(x1 - 1, y2) +
        query(x1 - 1, y1 - 1);
}

// A ideia de atualizar um intervalo (submatriz) e consultar um
// elemento (i,j) tambem sao validos

```

## 7.3 Sparse Table

```

/*
Resolve problemas de consulta a intervalos (RSQ, RMQ etc) de um vetor
estatico, ou seja, os valores nao sofrem update.
Alterar a funcao comb() de acordo (min, max, soma etc)
Pre-processamento O(NlogN) e consulta em O(1).
N = tamanho do vetor a[]
a[] deve ser indexado em 0
*/
const int MAXN = (1e6 + 1);
#define LOGN (21)
int st[MAXN][LOGN];
int N, a[MAXN];

int comb(int left, int right) { return min(left, right); }

void SparseTable() {
    rep(k, 0, LOGN) for (int i = 0; (i + (1 << k) - 1) < N; i++)
        st[i][k] =
            k ? comb(st[i][k - 1], st[i + (1 << (k - 1))][k - 1]) :
                a[i];
}

int query(int l, int r) {
    int k = log2(r - l + 1);
    return comb(st[l][k], st[r - (1 << k) + 1][k]);
}

```

## 7.4 RMQ

```

// Range Minimum Query: idx do menor elemento num intervalo de um
// array. Permite consultas e updates no array em O(logN). ATENCAO:
// Array A[] deve ser indexado em 0;
#define MAXN 500000
int A[MAXN], T[4 * MAXN];
int N; // #number of elements in A[]
int neutro = -1;

// combina o resultado de dois segmentos
int combine(int p1, int p2) {
    if (p1 == -1)
        return p2;
    if (p2 == -1)
        return p1;
    if (A[p1] <= A[p2])
        return p1;
    else
        return p2;
}

// chamar build() apos preencher o vetor A[]. O(N)
void build(int no = 1, int a = 0, int b = N - 1) {
    if (a == b) {
        T[no] = a;
    } else {
        int m = (a + b) / 2;
        int esq = 2 * no;
        int dir = esq + 1;
        build(esq, a, m);
        build(dir, m + 1, b);
        T[no] = combine(T[esq], T[dir]);
    }
}

// Modifica A[i] em O(logN), neste caso A[i] = v
void update(int i, int v, int no = 1, int a = 0, int b = N - 1) {
    if (a > i || b < i)
        return;
    if (a == i && b == i) {
        A[i] = v;
        T[no] = i; // desnecessario ;p
        return;
    }
    int m = (a + b) / 2;
    int esq = 2 * no;
    int dir = esq + 1;
    update(i, v, esq, a, m);
    update(i, v, dir, m + 1, b);
    T[no] = combine(T[esq], T[dir]);
}

// Retorna o idx k do menor valor A[k] no intervalo [i,j] em O(logN)
int query(int i, int j, int no = 1, int a = 0, int b = N - 1) {
    if (a > j || b < i)
        return neutro;
    if (a >= i && b <= j)
        return T[no];
    int m = (a + b) / 2;

```

```

int esq = 2 * no;
int dir = esq + 1;

int p1 = query(i, j, esq, a, m);
int p2 = query(i, j, dir, m + 1, b);
return combine(p1, p2);
}

int main() { return 0; }

```

## 7.5 Persistent Seg Tree

```

const int MAX = 1e5 + 5;
const int MAXT = 80 * MAX;
int a[MAX], tree[MAXT];
int L[MAXT], R[MAXT];
vector<int> root;
int cnt;

void build(int no, int i, int j) {
    if (i == j) {
        tree[no] = a[i];
        return;
    }
    L[no] = ++cnt;
    R[no] = ++cnt;
    build(L[no], i, (i + j) / 2);
    build(R[no], (i + j) / 2 + 1, j);
    tree[no] = tree[L[no]] + tree[R[no]];
}

int update(int no, int i, int j, int p, int v) {
    int NO = ++cnt;
    if (i == j) {
        tree[NO] = v;
        return NO;
    }

    L[NO] = L[no];
    R[NO] = R[no];

    if (p <= (i + j) / 2) L[NO] = update(L[NO], i, (i + j) / 2, p, v);
    else R[NO] = update(R[NO], (i + j) / 2 + 1, j, p, v);

    tree[NO] = tree[L[NO]] + tree[R[NO]];
    return NO;
}

int query(int no, int i, int j, int l, int r) {
    if (i > r || j < l) return 0;
    if (i >= l && j <= r) return tree[no];

    int left = query(L[no], i, (i + j) / 2, l, r);
    int right = query(R[no], (i + j) / 2 + 1, j, l, r);
    return left + right;
}

void init() {
    memset(tree, 0, sizeof tree);
    root.clear();
}

```

```

root.pb(0);
cnt = 0;
}

int main(int argc, char **argv) {
    query(root[k], 1, N, l, r)); // consulta na versao k
    root.pb(update(root[sz], 1, N, l, v)); // update
    return 0;
}

```

## 7.6 Seg Tree com Lazy

```

// RSG agora com queries e updates em intervalos. Precisa de Lazy
// Propagation. Array A[] deve ser indexado em 0. Nem sempre o array
// que sera modificado armazena apenas um valor, nesse caso usamos
// struct para representar cada no.
#define MAXN 500000
ll A[MAXN], tree[4 * MAXN], lazy[4 * MAXN];
int N;
int neutro = 0;

// funcao que realiza o merge de um intervalo, pode ser *, -, min,
// max, etc...
int combine(int segEsq, int segDir) { return segEsq + segDir; }

void build(int no = 1, int a = 0, int b = N - 1) {
    if (a == b) {
        tree[no] = A[a];
        return;
    }
    int m = (a + b) / 2;
    int esq = 2 * no;
    int dir = esq + 1;
    build(esq, a, m);
    build(dir, m + 1, b);
    tree[no] = combine(tree[esq], tree[dir]);
}

void propagate(int no, int a, int b) {
    if (lazy[no] != 0) {
        // esta parte depende do problema, neste caso queremos adicionar o
        // valor lazy[no] no intervalo [a,b], mas estamos atualizando
        // apenas o noh que representa este intervalo
        tree[no] += (b - a + 1) * lazy[no];
        if (a != b) {
            lazy[2 * no] += lazy[no];
            lazy[2 * no + 1] += lazy[no];
        }
        lazy[no] = 0;
    }
}

// update(i,j,v) faz A[k] += v, para i <= k <= j, em log(N)
void update(int i, int j, ll v, int no = 1, int a = 0, int b = N - 1) {
    if (lazy[no])
        propagate(no, a, b);
    if (a > j || b < i)
        return;
    if (a >= i && b <= j) {

```

```

    lazy[no] += v; // atualiza apenas a flag da raiz da subarvore
    propagate(no, a, b);
    return;
}
int m = (a + b) / 2;
int esq = 2 * no;
int dir = esq + 1;
update(i, j, v, esq, a, m);
update(i, j, v, dir, m + 1, b);
tree[no] = combine(tree[esq], tree[dir]);
}

// query(i,j) retorna o somatorio A[i] + A[i+1] + ... + A[j]
ll query(int i, int j, int no = 1, int a = 0, int b = N - 1) {
    if (lazy[no])
        propagate(no, a, b);
    if (a > j || b < i)
        return neutro;
    if (a >= i && b <= j)
        return tree[no];
    int m = (a + b) / 2;
    int esq = 2 * no;
    int dir = esq + 1;
    ll q1 = query(i, j, esq, a, m);
    ll q2 = query(i, j, dir, m + 1, b);
    return combine(q1, q2);
}

int main() { return 0; }

```

## 7.7 Union-Find

```

// Conjuntos Disjuntos. Inicialmente cada elemento eh lider de seu
// proprio conjunto. Operacoes de join(u,v) fazem com que os conjuntos
// que u e v pertencem se unam. find(u) retorna o lider do conjunto
// que u esta contido.
#define MAXV 100000
int V, pai[MAXV], rnk[MAXV], size[MAXV];

void init() { rep(i, 0, V) pai[i] = i, rnk[i] = 0, size[i] = 1; }

int find(int v) {
    if (v != pai[v])
        pai[v] = find(pai[v]);
    return pai[v];
}

void join(int u, int v) {
    u = find(u);
    v = find(v);
    if (u == v)
        return;

    if (rnk[u] < rnk[v])
        swap(u, v);
    pai[v] = u; // add v no conjunto de u
    size[u] += size[v];
    if (rnk[u] == rnk[v])
        rnk[u]++;
}

```

```

bool same_set(int u, int v) { return find(u) == find(v); }

int main() { return 0; }

```

## 7.8 Treap

```

typedef struct node {
    int prior, size;
    int val; // value stored in the array
    int sum; // whatever info you want to maintain in segtree for each
             // node
    int lazy; // whatever lazy update you want to do
    int rev;
    struct node *l, *r;
} node;
typedef node *pnode;
int sz(pnode t) { return t ? t->size : 0; }
void upd_sz(pnode t) {
    if (t)
        t->size = sz(t->l) + 1 + sz(t->r);
}

void lazy(pnode t) {
    if (!t || t->lazy == -1)
        return;
    t->val = t->lazy; // operation of lazy
    t->sum = t->lazy * sz(t);
    if (t->l)
        t->l->lazy = t->lazy; // propagate lazy
    if (t->r)
        t->r->lazy = t->lazy;
    t->lazy = -1;
}

void reset(pnode t) {
    if (t)
        t->sum = t->val; // no need to reset lazy coz when we call this
                        // lazy would itself be propagated
}

// combining two ranges of segtree
void combine(pnode &t, pnode l, pnode r) {
    if (!l || !r)
        return void(t = l ? l : r);
    t->sum = l->sum + r->sum;
}

void operation(pnode t) { // operation of segtree
    if (!t)
        return;
    reset(t); // reset the value of current node assuming it now
              // represents a single element of the array
    lazy(t->l);
    lazy(t->r); // imp:propagate lazy before combining t->l,t->r;
    combine(t, t->l, t);
    combine(t, t, t->r);
}

void push(pnode t) {
    if (!t || !t->rev)
        return;
    t->rev = false;
    swap(t->l, t->r);
}

```

```

    if (t->l)
        t->l->rev ^= true;
    if (t->r)
        t->r->rev ^= true;
}

void split(pnode t, pnode &l, pnode &r, int pos, int add = 0) {
    if (!t)
        return void(l = r = NULL);
    push(t);
    lazy(t);
    int curr_pos = add + sz(t->l);
    if (curr_pos <= pos) // element at pos goes to left subtree(l)
        split(t->r, t->r, r, pos, curr_pos + 1), l = t;
    else
        split(t->l, l, t->l, pos, add), r = t;
    upd_sz(t);
    operation(t);
}

// l->leftarray, r->rightarray, t->resulting array
void merge(pnode &t, pnode l, pnode r) {
    push(l);
    push(r);
    lazy(l);
    lazy(r);
    if (!l || !r)
        t = l ? l : r;
    else if (l->prior > r->prior)
        merge(l->r, l->r, r), t = l;
    else
        merge(r->l, l, r->l), t = r;
    upd_sz(t);
    operation(t);
}

pnode init(int val) {
    pnode ret = new node;
    ret->prior = rand();
    ret->size = 1;
    ret->val = val;
    ret->sum = val;
    ret->lazy = -1;
    ret->rev = 0;
    ret->l = NULL, ret->r = NULL;
    return ret;
}

int range_query(pnode t, int l, int r) { //[l,r]
    pnode L, mid, R;
    split(t, L, mid, l - 1);
    split(mid, t, R, r - 1); // note: r-1!!
    int ans = t->sum;
    merge(mid, L, t);
    merge(t, mid, R);
    return ans;
}

void range_update(pnode t, int l, int r, int val) { //[l,r]
    pnode L, mid, R;
    split(t, L, mid, l - 1);
    split(mid, t, R, r - 1); // note: r-1!!
    t->lazy = val; // lazy_update
    merge(mid, L, t);
    merge(t, mid, R);
}

```

```

}

void reverse(pnode t, int l, int r) {
    pnode L, mid, R;
    split(t, L, mid, l - 1);
    split(mid, mid, R, r - 1);
    mid->rev ^= true;
    merge(t, L, mid);
    merge(t, t, R);
}

void output(pnode t) {
    if (!t)
        return;
    push(t);
    lazy(t);
    output(t->l);
    printf("%d ", t->val);
    output(t->r);
}

int valor(int val) { return val & 1 ? 0 : 1; }

int main() {
    int P, Q;
    while (scanf("%d %d", &P, &Q) != EOF) {
        pnode tree = NULL, T1 = NULL, T2 = NULL, T3 = NULL;
        int val;
        rep(i, 0, P) {
            scanf("%d", &val);
            split(tree, T1, T2, i);
            merge(T1, T1, init(valor(val)));
            merge(tree, T1, T2);
        }
        while (Q--) {
        }
    }
}

```

## 7.9 Seg Tree 2D

```

struct node {
    int qt;
    int f1, f2, f3, f4;
};

node new_node() {
    node ret;
    ret.qt = ret.f1 = ret.f2 = ret.f3 = ret.f4 = 0;
    return ret;
}

vector<node> tree;
int cnt = 0;

bool inRange(int x1, int x2, int y1, int y2, int a1, int a2, int b1,
             int b2) {
    if (x2 < x1 || y2 < y1)
        return false;
    if (x2 < a1 || x1 > a2)
        return false;
    if (y2 < b1 || y1 > b2)
        return false;
}

```



```

    return false;
return true;
}

void update(int no, int x1, int x2, int y1, int y2, int a1, int a2,
           int b1, int b2, int val) {
    if (no == cnt)
        tree[cnt++] = new_node();

    if (x1 >= a1 && x2 <= a2 && y1 >= b1 && y2 <= b2) {
        tree[no].qt = val;
        return;
    }

    int f1 = 0, f2 = 0, f3 = 0, f4 = 0;
    if (inRange(x1, (x1 + x2) / 2, y1, (y1 + y2) / 2, a1, a2, b1, b2)) {
        if (!tree[no].f1)
            tree[no].f1 = cnt;
        update(tree[no].f1, x1, (x1 + x2) / 2, y1, (y1 + y2) / 2, a1, a2,
              b1, b2, val);
    }
    if (inRange(x1, (x1 + x2) / 2, (y1 + y2) / 2 + 1, y2, a1, a2, b1,
              b2)) {
        if (!tree[no].f2)
            tree[no].f2 = cnt;
        update(tree[no].f2, x1, (x1 + x2) / 2, (y1 + y2) / 2 + 1, y2, a1,
              a2, b1, b2, val);
    }
    if (inRange((x1 + x2) / 2 + 1, x2, y1, (y1 + y2) / 2, a1, a2, b1,
              b2)) {
        if (!tree[no].f3)
            tree[no].f3 = cnt;
        update(tree[no].f3, (x1 + x2) / 2 + 1, x2, y1, (y1 + y2) / 2, a1,
              a2, b1, b2, val);
    }
    if (inRange((x1 + x2) / 2 + 1, x2, (y1 + y2) / 2 + 1, y2, a1, a2,
              b1,
              b2)) {
        if (!tree[no].f4)
            tree[no].f4 = cnt;
        update(tree[no].f4, (x1 + x2) / 2 + 1, x2, (y1 + y2) / 2 + 1, y2,
              a1, a2, b1, b2, val);
    }

    if (tree[no].f1)
        f1 = tree[tree[no].f1].qt;
    if (tree[no].f2)
        f2 = tree[tree[no].f2].qt;
    if (tree[no].f3)
        f3 = tree[tree[no].f3].qt;
    if (tree[no].f4)
        f4 = tree[tree[no].f4].qt;

    tree[no].qt = f1 + f2 + f3 + f4;
}

int query(int no, int x1, int x2, int y1, int y2, int a1, int a2,
          int b1, int b2) {
    if (!inRange(x1, x2, y1, y2, a1, a2, b1, b2) || no >= cnt ||
        tree[no].qt == 0)
        return 0;
}

```

```

    if (x1 >= a1 && x2 <= a2 && y1 >= b1 && y2 <= b2)
        return tree[no].qt;

    int f1 = 0, f2 = 0, f3 = 0, f4 = 0;
    if (tree[no].f1)
        f1 = query(tree[no].f1, x1, (x1 + x2) / 2, y1, (y1 + y2) / 2, a1,
              a2, b1, b2);
    if (tree[no].f2)
        f2 = query(tree[no].f2, x1, (x1 + x2) / 2, (y1 + y2) / 2 + 1, y2,
              a1, a2, b1, b2);
    if (tree[no].f3)
        f3 = query(tree[no].f3, (x1 + x2) / 2 + 1, x2, y1, (y1 + y2) / 2,
              a1, a2, b1, b2);
    if (tree[no].f4)
        f4 = query(tree[no].f4, (x1 + x2) / 2 + 1, x2, (y1 + y2) / 2 + 1,
              y2, a1, a2, b1, b2);

    return f1 + f2 + f3 + f4;
}

void erase() {
    tree.clear();
    vector<node> xua;
    swap(tree, xua);
    tree.resize(1000010);
    cnt = 0;
}

int main() { return 0; }

```

## 7.10 Polyce

// <https://codeforces.com/blog/entry/11080>

```

#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>

using namespace std;
using namespace __gnu_pbds;

typedef tree<int, // tipo da variavel
            null_type,
            less<int>, // funcao de comparacao(greater, less_equal,
            rb_tree_tag, tree_order_statistics_node_update>
            ordered_set;

void newSet() {
    // fuciona como um set normal, mas ha 2 funcoes especiais: log(n)
    ordered_set T;
    ordered_set::iterator it;
    int k = *T.find_by_order(0); // retorna o K-esimo elemento segundo
    // a funcao de comparacao
    int kk = T.order_of_key(0); // retorna a posicao que um elemento
    // encaixaria segundo a funcao de
    // comparacao
}

#include <ext/rope>

```

```
using namespace __gnu_cxx;

void newVector() {
    // funciona como um vector, mas consegue algo a mais: (log(n))
    rope<int> v;
    rope<int>::iterator it;
    int l, r; // segmento
    rope<int> cur =
        v.substr(l, r - l + 1); // copia um segmento do vector
    v.erase(l, r - l + 1); // apaga um segmento
    v.insert(v.mutable_begin(), cur); // insere um segmento
    for (it = cur.mutable_begin(); it != cur.mutable_end(); it++)
        cout << *it << " "; // percorre ele
}

int main() { return 0; }
```

## 7.11 KD2

```
struct point {
    int x, y, z;
    point(int x = 0, int y = 0, int z = 0) : x(x), y(y), z(z) {}
    point operator-(point q) { return point(x - q.x, y - q.y, z - q.z); }
    int operator*(point q) { return x * q.x + y * q.y + z * q.z; }
};

typedef vector<point> polygon;
priority_queue<double> vans;
int NN, CC, KK, DD;

struct KDTreeNode {
    point p;
    int level;
    KDTreeNode *below, *above;
    KDTreeNode(const point &q, int lev) {
        p = q;
        level = lev;
        below = above = 0;
    }
    ~KDTreeNode() { delete below, above; }
    int diff(const point &pt) {
        switch (level) {
            case 0:
                return pt.x - p.x;
            case 1:
                return pt.y - p.y;
            case 2:
                return pt.z - p.z;
        }
        return 0;
    }
    ll distSq(point &q) { return (p - q) * (p - q); }
    int rangeCount(point &pt, ll K) {
        int count = (distSq(pt) <= K * K) ? 1 : 0;
        if (count)
            vans.push(-sqrt(distSq(pt)));

        int d = diff(pt);
        if (~d <= K && above != 0)
```

```
        count += above->rangeCount(pt, K);
        if (d <= K && below != 0)
            count += below->rangeCount(pt, K);
        return count;
    }
};

class KDTree {
public:
    polygon P;
    KDTreeNode *root;
    int dimension;
    KDTree() {}
    KDTree(polygon &poly, int D) {
        P = poly;
        dimension = D;
        root = 0;
        build();
    }
    ~KDTree() { delete root; }
    // count the number of pairs that has a distance less than K
    ll countPairs(ll K) {
        ll count = 0;
        rep(i, 0, P.size()) count += root->rangeCount(P[i], K) - 1;
        return count;
    }

protected:
    void build() {
        // random_shuffle(all(P));
        rep(i, 0, P.size()) { root = insert(root, P[i], -1); }
    }
    KDTreeNode *insert(KDTreeNode *t, const point &pt, int parentLevel)
    {
        if (t == 0) {
            t = new KDTreeNode(pt, (parentLevel + 1) % dimension);
            return t;
        } else {
            int d = t->diff(pt);
            if (d <= 0)
                t->below = insert(t->below, pt, t->level);
            else
                t->above = insert(t->above, pt, t->level);
        }
        return t;
    }
};

int main() {
    point e;
    e.z = 0;
    polygon p;
    set<ii> st;

    while (scanf("%d %d %d %d", &NN, &CC, &KK, &DD) != EOF) {
        p.clear();
        KK = min(NN, KK);
        st.clear();

        rep(i, 0, NN) {
            scanf("%d %d", &e.x, &e.y);
```

```

    st.insert(mp(e.x, e.y));
    p.pb(e);
}

KDTree tree(p, 2);
int ans = 0;
rep(i, 0, CC) {
    scanf("%d %d", &e.x, &e.y);
    if (st.count(mp(e.x, e.y)))
        continue;

    ll at = 0;
    rep(i, 0, 30) {
        at = ll(1) << i;
        while (!vans.empty())
            vans.pop();
        int aux = tree.root->rangeCount(e, at);
        if (aux >= KK)
            break;
    }
    double sum = 0.0;
    rep(i, 0, KK) {
        sum += -vans.top();
        vans.pop();
    }
    if (sum >= DD)
        ans++;
}
printf("%d\n", ans);
}
return 0;
}

```

## 8 Strings

### 8.1 KMP

*// obs: A funcao strstr (char\* text, char\* pattern) da biblioteca  
 // <cstring> implementa KMP (C-ANSI). A funcao retorna a primeira  
 // ocorrencia do padrao no texto, KMP retorna todas. nres -> O numero  
 // de ocorrencias do padrao no texto res[] -> posicoes das nres  
 // ocorrencias do padrao no texto Complexidade do algoritmo: O(n+m)\**

```

#define MAXN 100001
int pi[MAXN], res[MAXN], nres;
void kmp(string text, string pattern) {
    nres = 0;
    pi[0] = -1;
    rep(i, 1, pattern.size()) {
        pi[i] = pi[i - 1];
        while (pi[i] >= 0 && pattern[pi[i] + 1] != pattern[i])
            pi[i] = pi[pi[i]];
        if (pattern[pi[i] + 1] == pattern[i])
            ++pi[i];
    }
    int k = -1; // k+1 eh o tamanho do match atual
    rep(i, 0, text.size()) {
        while (k >= 0 && pattern[k + 1] != text[i])

```

```

        k = pi[k];
        if (pattern[k + 1] == text[i])
            ++k;
        if (k + 1 == pattern.size()) {
            res[nres++] = i - k;
            k = pi[k];
        }
    }
}

```

### 8.2 Aho Corasick

```

const int cc = 26;
const int MAX = 100;

int cnt;
int sig[MAX][cc];
int term[MAX];
int T[MAX];
int v[MAX];

inline int C(char c) { return c - '0'; }

void add(string s, int id) {
    int x = 0;
    rep(i, 0, s.size()) {
        int c = C(s[i]);
        if (!sig[x][c]) {
            term[cnt] = 0;
            sig[x][c] = cnt++;
        }
        x = sig[x][c];
    }
    term[x] = 1;
    v[id] = x;
}

void aho() {
    queue<int> q;
    rep(i, 0, cc) {
        int x = sig[0][i];
        if (!x)
            continue;
        q.push(x);
        T[x] = 0;
    }
    while (!q.empty()) {
        int u = q.front();
        q.pop();
        rep(i, 0, cc) {
            int x = sig[u][i];
            if (!x)
                continue;
            int v = T[u];
            while (v && !sig[v][i])
                v = T[v];
            v = sig[v][i];
            T[x] = v;
            term[x] += term[v];
            q.push(x);
        }
    }
}

```

```

    }
}

// Conta a quantidade de palavras de exatamente l caracteres que se
// pode formar com um determinado alfabeto, dado que algumas palavras
// sao "proibidas"

int mod = 1e9 + 7;
ll pd[100][MAX];

ll solve(int pos, int no) {
    if (pos == 0)
        return 1;
    if (pd[pos][no] != -1)
        return pd[pos][no];
    ll ans = 0;
    rep(i, 0, cc) {
        int v = no;
        while (v && !sig[v][i])
            v = T[v];
        v = sig[v][i];
        if (term[v])
            continue;
        ans = (ans + solve(pos - 1, v)) % mod;
    }
    return pd[pos][no] = ans;
}

void Qttd_de_Palavras() {
    while (1) {
        memset(sig, 0, sizeof sig);
        memset(pd, -1, sizeof pd);
        cnt = 1;
        int l = readInt();
        if (!l)
            break;
        int n = readInt();
        string pattern;
        rep(i, 0, n) {
            cin >> pattern;
            add(pattern, i);
        }
        aho();
        ll ans = 0;
        rep(i, 1, l + 1) ans = (ans + solve(i, 0)) % mod;
        printf("%d\n", ans);
    }
}

// Verifica quais padroes ocorreram em um texto

int alc[MAX];

void busca(string s) {
    int x = 0;
    rep(i, 0, s.size()) {
        int c = C[s[i]];
        while (x && !sig[x][c])
            x = T[x];
        x = sig[x][c];
    }
}

```

```

        alc[x] = 1;
    }
}

void Ql_Ocorreu() {
    string pattern, text;
    while (getline(cin, text)) {
        if (text == "*")
            break;
        memset(sig, 0, sizeof sig);
        memset(alc, 0, sizeof alc);
        cnt = 1;
        int n;
        cin >> n;
        rep(i, 0, n) {
            cin >> pattern;
            add(pattern, i);
        }
        aho();
        busca(text);
        for (int i = cnt - 1; i >= 0; i--) {
            if (alc[i])
                alc[T[i]] = 1;
        }
        rep(i, 0, n) {
            int u = v[i];
            if (alc[u])
                printf("Ocorreu\n");
            else
                printf("Nao ocorreu\n");
        }
    }
}

// Total de ocorrencias de cada padrao em uma string, mesmo com
// sufixos iguais
ll busca2(string s) {
    ll x = 0, cont = 0;
    rep(i, 0, s.size()) {
        int c = C[s[i]];
        while (x && !sig[x][c])
            x = T[x];
        x = sig[x][c];
        cont += term[x];
    }
    return cont;
}

void Qnts_vezes_Ocorreu() {
    string text, pattern;
    while (cin >> text) {
        if (text == "*")
            break;
        memset(sig, 0, sizeof sig);
        cnt = 1;
        int n = readInt();
        rep(i, 0, n) {
            cin >> pattern;
            add(pattern, i);
        }
        aho();
    }
}

```

```

    rep(i, 1, 10) debug(T[i]) cout << busca2(text) << endl;
}
}

// Encontra a primeira ocorrencia de cada padrao em uma string
void busca3(string s) {
    int x = 0;
    rep(i, 0, s.size()) {
        int c = C(s[i]);
        while (x && !sig[x][c])
            x = T[x];
        x = sig[x][c];
        if (!alc[x])
            alc[x] = i + 1;
    }
}

void Onde_Ocorreu() {
    string pattern, text;
    int tam[1000];
    while (cin >> text) {
        if (text == "*")
            break;
        memset(sig, 0, sizeof sig);
        memset(alc, 0, sizeof alc);
        cnt = 1;
        int n;
        cin >> n;
        rep(i, 0, n) {
            cin >> pattern;
            tam[i] = pattern.size();
            add(pattern, i);
        }

        aho();
        busca3(text);
        for (int i = cnt - 1; i >= 0; i--) {
            alc[T[i]] = min(alc[i], alc[T[i]]);
        }
        rep(i, 0, n) {
            int u = v[i];
            if (alc[u] != INF) {
                int k = alc[u] - tam[i] + 1;
                printf("De %d a %d\n", k, alc[u]);
            } else
                printf("Nao ocorreu\n");
        }
    }
}

```

### 8.3 Suffix Array

```

#define MAX 100010
#define MAX_N 100010
char T[MAX_N];
ll n;
int RA[MAX_N], tempRA[MAX_N];
int SA[MAX_N], tempSA[MAX_N];
int c[MAX_N];
int Phi[MAX_N], PLCP[MAX_N], LCP[MAX_N];

```

```

void countingSort(int k) {
    int i, sum, maxi = max((ll)300, n);
    memset(c, 0, sizeof c);
    for (i = 0; i < n; i++)
        c[i + k < n ? RA[i + k] : 0]++;
    for (i = sum = 0; i < maxi; i++) {
        int t = c[i];
        c[i] = sum;
        sum += t;
    }
    for (i = 0; i < n; i++)
        tempSA[c[SA[i] + k < n ? RA[SA[i] + k] : 0]++] = SA[i];
    for (i = 0; i < n; i++)
        SA[i] = tempSA[i];
}

void constructSA() {
    int i, k, r;
    for (i = 0; i < n; i++)
        RA[i] = T[i];
    for (i = 0; i < n; i++)
        SA[i] = i;

    for (k = 1; k < n; k <= 1) {
        countingSort(k);
        countingSort(0);
        tempRA[SA[0]] = r = 0;
        for (i = 1; i < n; i++)
            tempRA[SA[i]] = (RA[SA[i]] == RA[SA[i - 1]] &&
                             RA[SA[i] + k] == RA[SA[i - 1] + k])
                             ? r
                             : ++r;

        for (i = 0; i < n; i++)
            RA[i] = tempRA[i];
        if (RA[SA[n - 1]] == n - 1)
            break;
    }
}

void computeLCP() {
    int i, L;
    Phi[SA[0]] = -1;
    for (i = 1; i < n; i++)
        Phi[SA[i]] = SA[i - 1];
    for (i = L = 0; i < n; i++) {
        if (Phi[i] == -1) {
            PLCP[i] = 0;
            continue;
        }
        while (T[i + L] == T[Phi[i] + L])
            L++;
        PLCP[i] = L;
        L = max(L - 1, 0);
    }
    for (i = 0; i < n; i++) {
        LCP[i] = PLCP[SA[i]];
    }
}

int main() {
    // concatenar $ no final

```

## 8.4 Suffix Array (Gugu)

```

}

const int MAX = 100010;
int gap, tam, sa[MAX], pos[MAX], lcp[MAX], tmp[MAX];

bool sufixCmp(int i, int j) {
    if (pos[i] != pos[j])
        return pos[i] < pos[j];
    i += gap, j += gap;
    return (i < tam && j < tam) ? pos[i] < pos[j] : i > j;
}

void buildSA(char s[]) {
    tam = strlen(s);
    for (int i = 0; i < tam; i++)
        sa[i] = i, pos[i] = s[i], tmp[i] = 0;

    for (gap = 1; gap <= tam; gap *= 2) {
        sort(sa, sa + tam, sufixCmp);
        tmp[0] = 0;
        for (int i = 0; i < tam - 1; i++)
            tmp[i + 1] = tmp[i] + sufixCmp(sa[i], sa[i + 1]);
        for (int i = 0; i < tam; i++)
            pos[sa[i]] = tmp[i];
        if (tmp[tam - 1] == tam - 1)
            break;
    }
}

ll buildLCP(char s[]) {
    ll sum = 0;
    for (int i = 0, k = 0; i < tam; i++) {
        if (pos[i] == tam - 1)
            continue;
        for (int j = sa[pos[i] + 1]; s[i + k] == s[j + k];)
            k++;
        lcp[pos[i] + 1] = k;
        sum += k;
        if (k > 0)
            k--;
    }
    return sum;
}

void PrintAll(char s[]) {
    printf("SA\ttam\tLCP\tSuffix\n");
    rep(i, 0, tam) printf("%2d\t%2d\t%2d\t%s\n", sa[i], tam - sa[i],
        lcp[i], s + sa[i]);
}

ll num_subs(ll m) { return (ll)tam * (tam + 1) / 2 - m; }
ll num_subsrn() {
    ll ret = 0;
    rep(i, 1, tam) if (lcp[i] > lcp[i - 1]) ret += lcp[i] - lcp[i - 1];
    return ret;
}

void printans(char s[], int n) {
    int maior = 0, id = -1;
    rep(i, 0, tam) if (lcp[i] > n && lcp[i] > maior) maior = lcp[i],
        id = i;

    if (id == -1)
        printf("*");
}

```

```

else
    rep(i, sa[id], sa[id] + maior) printf("%c", s[i]);
printf("\n");
}

char s[MAX];
int main() {
    while (1) {
        scanf("%s", s);
        if (s[0] == '*')
            break;

        buildSA(s);
        ll m = buildLCP(s);

        PrintAll(s); // printa sa, lcp, suffixs
        // printf("%lld\n", num_subs(m)); //numero de substrings nao
        // repetidas printf("%lld\n", num_subsrn()); //numero de
        // substrings
        // que se repete printans(s, 2); //maior substring de tamanho
        // maior
        // ou igual a n que se repete
    }
}

```

## 8.5 Rolling Hash

```

// Permite encontrar um hash de uma substring de S. precompute O(n),
// my_hash O(1)
#define NN 1000006
const ll mod = 1e9 + 7; // modulo do hash
const ll x = 33; // num. primo > que o maior caracter de S.
ll H[NN], X[NN];
ll V(char c) { return c - 'A'; }
ll my_hash(int i, int j) {
    ll ret = H[j];
    if (!i)
        return ret;
    return ((ret - (H[i - 1] * X[j - i + 1]) % mod) + mod) % mod;
}

void precompute(string s) {
    X[0] = 1;
    rep(i, 1, NN) X[i] = (X[i - 1] * x) % mod;
    H[0] = V(s[0]);
    rep(i, 1, s.size()) H[i] = ((H[i - 1] * x) % mod + V(s[i])) % mod;
}

```

## 8.6 Longest Common Prefix with Hash

```

// Longest Common Prefix between S[i..] and S[j..]
int lcp(int i, int j, int tam) {
    int lo = 0, hi = tam, ans;
    while (lo <= hi) {
        int mid = (lo + hi) / 2;
        if (my_hash(i, i + mid - 1) == my_hash(j, j + mid - 1)) {
            ans = mid;
            lo = mid + 1;
        } else

```

```

    hi = mid - 1;
}
return ans;
}

```

## 8.7 Minimum Lexicographic Rotation

```

// Retorna a menor string lexicografica de s. Necessario my_hash() e
// lcp()
string min_lex_rot(string s) {
    int t = s.size();
    precompute(s); // hashing
    s += s;
    int idx = 0;
    for (int i = 1; i < t; i++) {
        // tam do prefix comum
        int len = lcp(i, idx, t);
        if (s[i + len] < s[idx + len])
            idx = i;
    }
    return s.substr(idx, t);
}

```

## 8.8 Longest Palindrome (Manacher algorithm)

```

string preprocess(string s) {
    int n = s.length();
    if (n == 0)
        return "^$";

    string ret = "^";
    for (int i = 0; i < n; i++)
        ret += "#" + s.substr(i, 1);
    ret += "$";
    return ret;
}

string longestPalindrome(string s) {
    L = C = s.size();
    string T = preprocess(s);
    int n = T.length();
    int *P = new int[n];
    int C = 0, R = 0;

    for (int i = 1; i < n - 1; i++) {
        int i_mirror = 2 * C - i;
        P[i] = (R > i) ? min(R - i, P[i_mirror]) : 0;
        while (T[i + 1 + P[i]] == T[i - 1 - P[i]])
            P[i]++;
        if (i + P[i] > R) {
            C = i;
            R = i + P[i];
        }
    }

    int maxLen = 0;
    int centerIndex = 0;
}

```

```

for (int i = 1; i < n - 1; i++) {
    if (!P[i])
        continue;
    if (P[i] > maxLen) {
        maxLen = P[i];
        centerIndex = i;
    }
}
delete[] P;
return s.substr((centerIndex - 1 - maxLen) / 2, maxLen);
}

```

## 8.9 Autômato de Sufixos

```

struct state {
    int len, link;
    int next[26];
};

const int MAXN = 200020;
state st[2 * MAXN]; // vetor que armazena os estados
int sz; // contador do numero de estados
int last; // numero do estado que corresponde ao texto todo

void sa_init() {
    sz = 1;
    last = 0;
    st[0].len = 0;
    st[0].link = -1;
    rep(i, 0, 26) st[0].next[i] = 0;
    // limpa o mapeamento de transicoes
}

void sa_extend(int c, ll &ans) {
    int cur = sz++; // novo estado a ser criado
    st[cur].len = st[last].len + 1;
    rep(i, 0, 26) st[cur].next[i] = 0;
    int p; // variavel que itera sobre os estados terminais
    for (p = last; p != -1 && !st[p].next[c]; p = st[p].link) {
        st[p].next[c] = cur;
    }
    if (p == -1) { // nao ocorreu transicao c nos estados terminais
        st[cur].link = 0;
        ans += st[cur].len;
    } else { // ocorreu transicao c no estado p
        int q = st[p].next[c];
        if (st[p].len + 1 == st[q].len) {
            st[cur].link = q;
        } else {
            int clone = sz++; // criacao do vertice clone de q
            st[clone].len = st[p].len + 1;
            rep(i, 0, 26) st[clone].next[i] = st[q].next[i];
            st[clone].link = st[q].link;
            for (; p != -1 && st[p].next[c] == q; p = st[p].link) { // atualizacao das transicoes c
                st[p].next[c] = clone;
            }
            st[q].link = st[cur].link = clone;
        }
        ans += st[cur].len - st[st[cur].link].len;
    }
}

```

```

    }
    // atualizacao do estado que corresponde ao texto
    last = cur;
}

bool busca_automato(int m, string p) {
    int i, pos = 0;
    for (i = 0; i < m; i++) {
        if (st[pos].next[p[i]] == 0) {
            return false;
        } else {
            pos = st[pos].next[p[i]];
        }
    }
    return true;
}

int maior_tamanho_em_comum(string s, string t) {
    ll nothing = 0;
    // Constroi o automato com o primeiro texto
    sa_init();
    for (int i = 0; i < (int)s.size(); i++)
        sa_extend(s[i] - 'a', nothing);
    int estado = 0, tamanho = 0, maior = 0;
    // Passando pelos caracteres do segundo texto
    for (int i = 0; i < (int)t.size(); ++i) {
        while (estado && !st[estado].next[t[i] - 'a']) {
            estado = st[estado].link;
            tamanho = st[estado].len;
        }
        if (st[estado].next[t[i] - 'a']) {
            estado = st[estado].next[t[i] - 'a'];
            tamanho++;
        }
        if (tamanho > maior) {
            maior = tamanho;
        }
    }
    return maior;
}

int main() {
    char s[MAXN];
    char p[MAXN];
    while (gets(s)) {
        sa_init();
        int tam = strlen(s);
        ll ans = 0;
        rep(i, 0, tam) { sa_extend(s[i] - 'a', ans); }
        gets(p);
        printf("%d\n", maior_tamanho_em_comum(s, p));
    }
    return 0;
}

```

## 8.10 Suffix Array

```

#define MAX 100010
#define MAX_N 100010
char T[MAX_N];

```

```

ll n;
int RA[MAX_N], tempRA[MAX_N];
int SA[MAX_N], tempSA[MAX_N];
int c[MAX_N];
int Phi[MAX_N], PLCP[MAX_N], LCP[MAX_N];

void countingSort(int k) {
    int i, sum, maxi = max((ll)300, n);
    memset(c, 0, sizeof c);
    for (i = 0; i < n; i++)
        c[i + k < n ? RA[i + k] : 0]++;
    for (i = sum = 0; i < maxi; i++) {
        int t = c[i];
        c[i] = sum;
        sum += t;
    }
    for (i = 0; i < n; i++)
        tempSA[c[SA[i] + k < n ? RA[SA[i] + k] : 0]++] = SA[i];
    for (i = 0; i < n; i++)
        SA[i] = tempSA[i];
}

void constructSA() {
    int i, k, r;
    for (i = 0; i < n; i++)
        RA[i] = T[i];
    for (i = 0; i < n; i++)
        SA[i] = i;

    for (k = 1; k < n; k <= 1) {
        countingSort(k);
        countingSort(0);
        tempRA[SA[0]] = r = 0;
        for (i = 1; i < n; i++)
            tempRA[SA[i]] = (RA[SA[i]] == RA[SA[i - 1]] &&
                             RA[SA[i] + k] == RA[SA[i - 1] + k])
                             ? r
                             : ++r;
        for (i = 0; i < n; i++)
            RA[i] = tempRA[i];
        if (RA[SA[n - 1]] == n - 1)
            break;
    }
}

void computeLCP() {
    int i, L;
    Phi[SA[0]] = -1;
    for (i = 1; i < n; i++)
        Phi[SA[i]] = SA[i - 1];
    for (i = L = 0; i < n; i++) {
        if (Phi[i] == -1) {
            PLCP[i] = 0;
            continue;
        }
        while (T[i + L] == T[Phi[i] + L])
            L++;
        PLCP[i] = L;
        L = max(L - 1, 0);
    }
    for (i = 0; i < n; i++) {

```



```

    LCP[i] = PLCP[SA[i]];
}
}
int main() {
    // concatenar $ no final
}

```

## 8.11 Z Algorithm

```

// Algorithm produces an array Z where Z[i] is the length of the
// longest substring starting from S[i] which is also a prefix of S.
string s;
vector<int> z;

void Z() {
    int n = s.size(), L = 0, R = 0;
    z.assign(n, 0);

    for (int i = 1; i < n; i++) {
        if (i > R) {
            L = R = i;
            while (R < n && s[R - L] == s[R])
                R++;
            z[i] = R - L;
            R--;
        } else {
            int k = i - L;
            if (z[k] < R - i + 1)
                z[i] = z[k];
            else {
                L = i;
                while (R < n && s[R - L] == s[R])
                    R++;
                z[i] = R - L;
                R--;
            }
        }
    }
}

```

## 9 PD

### 9.1 Soma acumulada 2D

```

/*Retorna o somatorio dos elementos de uma submatriz em O(1).
 * Submatriz definida por canto superior esquerdo (x1,y1) e canto
 * inferior direito (x2,y2) .: x1 <= x2 && y1 <= y2 */
#define MAXN 3000
int N, M; // linhas colunas
long long V[MAXN + 2][MAXN + 2]; // matriz da entrada
long long S[MAXN + 2][MAXN + 2]; // matriz com as somas acumuladas

// precomputa as somas em O(N*M)
void precal() {
    rep(x, 0, N) rep(y, 0, M) {
        S[x][y] = V[x][y];
        if (x > 0)

```

```

        S[x][y] += S[x - 1][y];
        if (y > 0)
            S[x][y] += S[x][y - 1];
        if (x > 0 && y > 0)
            S[x][y] -= S[x - 1][y - 1];
    }
}
// retorna a soma da submatriz em O(1)
long long sum(int x1, int y1, int x2, int y2) {
    long long soma = S[x2][y2];
    if (x1 > 0)
        soma -= S[x1 - 1][y2];
    if (y1 > 0)
        soma -= S[x2][y1 - 1];
    if (x1 > 0 && y1 > 0)
        soma += S[x1 - 1][y1 - 1];
    return soma;
}

```

## 9.2 Knuth Optimization

```

int N, B, C, yep, save[MAXN][MAXN], sav[MAXN];
ll n[MAXN], mc[MAXN][MAXN], se[MAXN], sd[MAXN], pd[MAXN][MAXN];

ll solve(int i, int k) {
    if (i == N)
        return 0;
    if (k == 1)
        return pd[i][k] = mc[i][N - 1];
    if (pd[i][k] != -1)
        return pd[i][k];

    ll ret = LINF;
    int ini = i, fim = N - k + 1, best = -1;
    if (i && save[i - 1][k])
        ini = save[i - 1][k];
    if (save[i][k - 1])
        fim = save[i][k - 1] + 1;

    rep(l, ini, fim) {
        ll aux = solve(l + 1, k - 1) + mc[i][l];
        if (ret > aux) {
            best = l;
            ret = aux;
        }
    }
    save[i][k] = best;
    return pd[i][k] = ret;
}

int main() {
    rep(i, 0, N) scanf("%lld", &n[i]);

    se[0] = n[0];
    rep(i, 1, N) se[i] = se[i - 1] + n[i];

    sd[N - 1] = n[N - 1];
    for (int i = N - 2; i >= 0; i--)
        sd[i] = sd[i + 1] + n[i];
}

```

```

rep(i, 1, N) pd[0][i] = pd[0][i - 1] + se[i - 1];
for (int i = N - 2; i >= 0; i--)
    pd[N - 1][i] = pd[N - 1][i + 1] + sd[i + 1];

rep(i, 1, N) {
    rep(j, i + 1, N) pd[i][j] = pd[i - 1][j] - n[i - 1] * (j - i + 1);
}
for (int i = N - 2; i >= 0; i--) {
    for (int j = i - 1; j >= 0; j--)
        pd[i][j] = pd[i + 1][j] - n[i + 1] * (i - j + 1);
}

rep(i, 0, N) {
    if (pd[i][i + 1] < pd[i + 1][i])
        mc[i][i + 1] = pd[i][i + 1], save[i][i + 1] = i + 1;
    else
        mc[i][i + 1] = pd[i + 1][i], save[i][i + 1] = i;
    rep(j, i + 2, N) {
        int ini = save[i][j - 1];
        mc[i][j] = pd[i][ini] + pd[j][ini], save[i][j] = ini;
        rep(k, ini + 1, j + 1) {
            ll a = pd[i][k] + pd[j][k];
            if (mc[i][j] <= a)
                break;
            mc[i][j] = a;
            save[i][j] = k;
        }
    }
    rep(j, 0, N + 1) { pd[i][j] = -1, save[i][j] = 0; }
}

rep(j, 0, N + 1) pd[N][j] = -1, save[N][j] = 0;

solve();

return 0;
}

```

### 9.3 Convex Hull Trick

```

bool bad(int l1, int l2, int l3) {
    return (B[l3] - B[l1]) * (M[l1] - M[l2]) <
           (B[l2] - B[l1]) * (M[l1] - M[l3]);
}

void add(long long m, long long b) {
    M.push_back(m);
    B.push_back(b);
    while (M.size() >= 3 &&
           bad(M.size() - 3, M.size() - 2, M.size() - 1)) {
        M.erase(M.end() - 2);
        B.erase(B.end() - 2);
    }
}

long long query(long long x) {
    if (pointer >= M.size())
        pointer = M.size() - 1;
    while (pointer < M.size() - 1 &&
           M[pointer + 1] * x + B[pointer + 1] <
           M[pointer] * x + B[pointer])
        pointer++;
}

```

```

return M[pointer] * x + B[pointer];
}

struct hux {
    int a, b, id;
};

bool my_sort(hux a, hux b) {
    return a.b != b.b ? a.b > b.b : a.a > b.a;
}

const ll LINF = 1LL << 52;
const double EPS = 1e-9;
const int MAXV = 100010;

double intersept(hux a, hux b) {
    return double(b.b - a.b) / (a.a - b.a);
}

vector<pair<double, double>> convex_hux(const vector<hux> &v) {
    int p = 0, n = v.size(), bestai = v[0].a;
    double cross = 0.0;
    pair<double, int> aux;

    priority_queue<pair<double, int>> pq;
    vector<pair<double, double>> ret(n + 1, mp(-1, -1));

    pq.push(mp(cross, p));
    ret[v[p].id].F = cross, ret[v[p].id].S = LINF;

    rep(i, 1, n) {
        aux = pq.top();
        cross = aux.F, p = aux.S;

        if (v[i].a <= bestai)
            continue;
        bestai = v[i].a;

        double new_cross = intersept(v[i], v[p]);
        while (new_cross <= cross + EPS) {
            pq.pop();
            ret[v[p].id] = mp(-1.0, -1.0);

            aux = pq.top();
            cross = aux.F, p = aux.S;

            new_cross = intersept(v[i], v[p]);
        }

        pq.push(mp(new_cross, i));
        ret[v[p].id].S = new_cross;
        ret[v[i].id].F = new_cross;
        ret[v[i].id].S = LINF;
    }

    // rep(i, 0, n) cout << ret[i].F << " " << ret[i].S << "\n";

    return ret;
}

```

## 9.4 Longest Increasing Subsequence

```
// Maior subsequencia crescente
#define MAX_N 100
int vet[MAX_N], P[MAX_N], N;
void reconstruct_print(int end) {
    int x = end;
    stack<int> s;
    while (P[x] >= 0) {
        s.push(vet[x]);
        x = P[x];
    }
    printf("%d", vet[x]);
    while (!s.empty()) {
        printf(" ", %d", s.top());
        s.pop();
    }
}
int lis() {
    int L[MAX_N], L_id[MAX_N];
    int li = 0, lf = 0; // lis ini, lis end
    rep(i, 0, N) {
        int pos = lower_bound(L, L + li, vet[i]) - L;
        L[pos] = vet[i];
        L_id[pos] = i;
        P[i] = pos ? L_id[pos - 1] : -1;
        if (pos + 1 > li) {
            li = pos + 1;
            lf = i;
        }
    }
    reconstruct_print(lf);
    return li;
}
```

## 9.5 Kadane 1D

```
// Encontra maior soma contigua positiva num vetor em O(N). {s,f}
// contem o intervalo de maior soma.
int Kadane1D(int vet[], int N, int &s, int &f) {
    int ret = -INF, sum, saux;
    sum = s = f = saux = 0;
    rep(i, 0, N) {
        sum += vet[i];
        if (sum > ret) {
            ret = sum;
            s = saux;
            f = i;
        }
        if (sum < 0) {
            sum = 0;
            saux = i + 1;
        }
    }
    return ret;
}
```

## 9.6 Kadane 2D

```
/*Maior soma de uma sub-matriz a partir de valores positivos.
 * [x1,y1]=upper-left, [x2,y2]=bottom-right*/
int L, C, pd[MAX_L], mat[MAX_L][MAX_C];
int x1, y1, x2, y2;
int Kadane2D() {
    int ret = 0, aux;
    rep(left, 0, C) {
        rep(i, 0, L) pd[i] = 0;
        rep(right, left, C) {
            rep(i, 0, L) pd[i] += mat[i][right];
            int sum = aux = 0;
            rep(i, 0, L) { // Kadane1D
                sum += pd[i];
                if (sum > ret)
                    ret = sum, x1 = aux, y1 = left, x2 = i, y2 = right;
                if (sum < 0)
                    sum = 0, aux = i + 1;
            }
        }
    }
    return ret;
}
```

## 9.7 Knapsack0-1

```
//[IME] 0-1 Knapsack, v-valores, w-pesos, Cap-capacidade
int mochila01(vector<int> v, vector<int> w, int Cap) {
    int n = v.size();
    int dp[n + 1][Cap + 1];
    for (int i = 0; i <= n; i++)
        dp[i][0] = 0;
    for (int j = 0; j <= Cap; j++)
        dp[0][j] = 0;
    for (int i = 1; i <= n; i++)
        for (int j = 1; j <= Cap; j++) {
            if (w[i - 1] > j)
                dp[i][j] = dp[i - 1][j];
            else
                dp[i][j] =
                    max(dp[i - 1][j], v[i - 1] + dp[i - 1][j - w[i - 1]]);
        }
    return dp[n][Cap];
}
```

## 9.8 Edit Distance

```
//[IME] menor custo para transformar a em b, dado as operacoes de
// inserir, remover e substituir caracteres de a
int editDistance(string a, string b) {
    int cost, insertCost = 1, deletCost = 1, substCost = 1;
    int m = a.size();
    int n = b.size();
    int d[m + 1][n + 1];
    for (int i = 0; i <= m; i++)
```

```

    d[i][0] = i * deletCost;
    for (int j = 0; j <= n; j++)
        d[0][j] = j * insertCost;
    for (int i = 1; i <= m; i++)
        for (int j = 1; j <= n; j++) {
            if (a[i - 1] == b[j - 1])
                cost = 0;
            else
                cost = substCost;
            d[i][j] =
                min(d[i - 1][j] + deletCost,
                    min(d[i][j - 1] + insertCost, d[i - 1][j - 1] + cost));
        }
    return d[m][n];
}

```

## 10 Sorting

### 10.1 Merge Sort com num de Inversoes

```

// Ordena arr aplicando mergesort e conta o numero de inversoes
void merge(int *arr, int size1, int size2, ll &inversions) {
    int temp[size1 + size2 + 2];
    int ptr1 = 0, ptr2 = 0;

    while (ptr1 + ptr2 < size1 + size2) {
        if (ptr1 < size1 && arr[ptr1] <= arr[size1 + ptr2] ||
            ptr1 < size1 && ptr2 >= size2)
            temp[ptr1 + ptr2] = arr[ptr1++];

        if (ptr2 < size2 && arr[size1 + ptr2] < arr[ptr1] ||
            ptr2 < size2 && ptr1 >= size1) {
            temp[ptr1 + ptr2] = arr[size1 + ptr2++];
            inversions += size1 - ptr1;
        }
    }

    for (int i = 0; i < size1 + size2; i++)
        arr[i] = temp[i];
}

void mergeSort(int *arr, int size, ll &inversions) {
    if (size == 1)
        return;

    int size1 = size / 2, size2 = size - size1;
    mergeSort(arr, size1, inversions);
    mergeSort(arr + size1, size2, inversions);
    merge(arr, size1, size2, inversions);
}

```

### 10.2 Quick Sort

```

// No main, chamar quicksort(array, 0, tam-1);
int partition(int s[], int l, int h) {
    int i, p, firsthigh;
    p = h;

```

```

    firsthigh = l;
    for (i = l; i < h; i++)
        if (s[i] < s[p]) {
            swap(s[i], s[firsthigh]);
            firsthigh++;
        }
    swap(s[i], s[firsthigh]);
    return firsthigh;
}

void quicksort(int s[], int l, int h) {
    int p;
    if ((h - l) > 0) {
        p = partition(s, l, h);
        quicksort(s, l, p - 1);
        quicksort(s, p + 1, h);
    }
}

```

## 11 Miscelânea

### 11.1 Calendário

```

// Routines for performing computations on dates. In these routines,
// months are expressed as integers from 1 to 12, days are expressed
// as integers from 1 to 31, and years are expressed as 4-digit
// integers.
string dayOfWeek[] = {"Mon", "Tue", "Wed", "Thu", "Fri", "Sat",
    "Sun"};

// converts Gregorian date to integer (Julian day number)
int dateToInt(int m, int d, int y) {
    return 1461 * (y + 4800 + (m - 14) / 12) / 4 +
        367 * (m - 2 - (m - 14) / 12 * 12) / 12 -
        3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 + d - 32075;
}

// converts integer (Julian day number) to Gregorian date:
// month/day/year
void intToDate(int jd, int &m, int &d, int &y) {
    int x, n, i, j;
    x = jd + 68569;
    n = 4 * x / 146097;
    x -= (146097 * n + 3) / 4;
    i = (4000 * (x + 1)) / 1461001;
    x -= 1461 * i / 4 - 31;
    j = 80 * x / 2447;
    d = x - 2447 * j / 80;
    x = j / 11;
    m = j + 2 - 12 * x;
    y = 100 * (n - 49) + i + x;
}

// converts integer (Julian day number) to day of week
string intToDay(int jd) { return dayOfWeek[jd % 7]; }

int main() {
    int jd = dateToInt(3, 24, 2004);
    int m, d, y;
    intToDate(jd, m, d, y);
    string day = intToDay(jd);
    // expected output:

```

```
// 2453089
// 3/24/2004
// Wed
cout << jd << endl
      << m << "/" << d << "/" << y << endl
```

```
      << day << endl;
}
```

---