# The Fixed Point Characterization of CTL

Formal Methods – Unit V

#### Lecture Objectives

- Understand fixpoint theory in the context of temporal logic
- Explore least and greatest fixpoints
- Learn how CTL operators can be expressed using fixpoints
- Apply fixpoint characterization to model checking

### Introduction to Fixpoints

- A fixpoint of a function f is a value x such that f(x) = x
- Important in defining recursive properties in logic and computation

#### Types of Fixpoints

- Least Fixpoint (LFP):
  The smallest solution x such that f(x) = x
- Greatest Fixpoint (GFP):
  The largest solution x such that f(x) = x
- Used respectively in 'until' and 'globally' CTL operators

## Why Fixpoints in CTL?

- CTL operators like EU (exists until) and EG (exists globally) are inherently recursive
- Fixpoints provide a clean, formal way to define their semantics

#### Fixed Point Characterization – EU

- $E[p U q] = least fixpoint of: \lambda Z. q V (p \wedge EX Z)$
- Start with states satisfying q and grow backward using p and EX

#### Fixed Point Characterization – EG

- EG p = greatest fixpoint of:  $\lambda Z$ . p  $\wedge$  EX Z
- Start with all states and eliminate those not satisfying the condition

## Example: E[p U q]

- 1. Initialize  $Z0 = \{s \mid q \text{ holds in } s\}$
- 2. Repeat  $Z(i+1) = Z(i) \cup \{s \mid p \text{ holds in } s \text{ and } EX Z(i)\}$
- 3. Stop when Z(i+1) = Z(i) (convergence reached)

### Example: EG p

- 1. Initialize  $Z0 = \{s \mid p \text{ holds in } s\}$
- 2. Repeat  $Z(i+1) = Z(i) \cap \{s \mid \exists \text{ successor in } Z(i)\}$
- 3. Continue until Z(i+1) = Z(i)

#### Visualizing Fixpoint Computation

- Iterative marking of states in the Kripke structure
- Convergence indicates property satisfaction in initial state

### Relation to Model Checking

- Labeling algorithm is built on fixpoint evaluation
- Efficient and forms basis of many tools (e.g., NuSMV, SPIN)

### Fixpoints and Modal Mu Calculus

- CTL is a fragment of modal μ-calculus
- μ-calculus uses fixpoint operators μ (least) and ν (greatest) explicitly

### Strengths of Fixpoint View

- Mathematically precise and elegant
- Supports compositional reasoning
- Used in advanced verification tools

#### Summary

- Fixpoints define semantics of recursive temporal operators
- EU  $\rightarrow$  least fixpoint, EG  $\rightarrow$  greatest fixpoint
- Efficient fixpoint algorithms enable model checking
- Foundation for advanced logic like μ-calculus

#### Questions

- How does the fixpoint process guarantee termination?
- Can all CTL properties be expressed using fixpoints?
- Why are some fixpoints least and others greatest?