

Assignment 1 Report

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1 Program Design

This train simulation is implemented in OpenMPI. Primary design considerations are:

- Each edge is simulated by one process as required in the specifications.
- There is a master process responsible for the following:
 - Distribution of information (map, lines) to slave processes.
 - Keeping track of train states (travelling/stationary).
 - Keeping track of station door open/close times.
 - Synchronising time across all slaves.
- Slave processes hold queues of trains and send them to each other, and reports to master whenever train states change.

Assumptions

- Only one train can open its doors at each station at any one time, regardless of direction.
- Train stations have infinite capacity for waiting trains.
- Time units are discrete and can have no subdivisions
 - **Implication:** It is sufficient to store all time units as integers instead of floating point numbers
- Trains must open their doors for at least 1 unit of time.
 - **Implication:** We round every randomly generated door open time up to the nearest integer

2 Points to Note / Implementation Details

- Current simulation time needs to be shared across all processes. In addition, time can only be advanced after all processes have completed the actions to be done in the current tick.
 - **Implication:** Explicit synchronisation messages need to be sent between master and slaves before and after the advancement of time.

- Each slave process maintains two queues of trains. The first queue contains the trains that are waiting to enter the edge (*entry queue*). The second queue contains the trains that are waiting (arrived and/or door open) at the next station but have yet to close their doors, and the time at which they would close their doors (*exit queue*).
 - For convenience, the same queue implementation is used for both queues. The queue stores pairs where first element is the train and second element is the time to dequeue the train. *Entry queue*'s elements simply have garbage values for their second element.
- The logic for each slave process at each tick is as follows:
 1. If a current train is occupying the edge and it can leave at the current time, the process will query master to find out when this train will finish closing its doors. The process will then add this train and the time received into its *exit queue*. The edge will then be marked as available.
 2. If there are trains waiting to access the edge, the process will dequeue the first one and set it as the current train.
 3. If the slave process should spawn trains for any of the lines its edge falls on, it will do so, query master for their door closing time, and enqueue them (see step 1 above).
 4. If the train at the head of *exit queue* can close its doors at the **next** tick, the train is sent to the process containing the next edge that it should traverse.
 - Since adjacent edges would only receive these trains at step 6, trains sent at this step would prematurely begin traversing the next edge.
 - This ensures that before step 2, every slave process already holds all trains that can begin traversing the edge (if the edge is unoccupied).
 5. The slave process sends “no more trains” messages to all adjacent edges.
 6. The slave process waits for messages from all adjacent edges. If a train is sent, it is enqueued into the *entry queue*. When “no more trains” messages are received from all adjacent edges, this step is complete.
 7. The slave process informs master that it has completed all actions for the current tick.
 8. The slave process waits for an OpenMPI broadcast message that would advance time.
- The logic for the master process at each tick is as follows:
 1. The master process waits for messages from slave processes. If a request for next door close time is received, it computes the appropriate value, updates station statistics and sends it back. If it receives a message that all actions for the current tick have been completed, it increments a counter. When all slave processes have reported completion, this step is complete.
 2. The master process prints the per-tick output, and broadcasts the next tick time.
 3. If the simulation has reached completion, instead of broadcasting the next tick time, master will broadcast the shutdown signal.
- The generic `MPI_Send` is used for slave-slave communication. Since messages are small, they are buffered and do not cause any deadlock. Deadlock does not occur when running all test cases in the lab. However, in the event that `MPI_Send` is not buffered and deadlock occurs, this can be resolved in the following ways:

- Non-blocking sends can be used for slave steps 4-5 and `MPI_Wait` called after slave step 6. This assumes that the non-blocking send will make process in the background even before `MPI_Wait` is called.
- Alternatively, a window can be created via `MPI_Win_create`. All information to be communicated in slave steps 4-5 will be written to its own window. When all slave processes have finished updating their own window (synchronised with `MPI_Barrier`), `MPI_Get` calls can be used in slave step 6 to obtain the relevant information.

3 Execution Time

3.1 Testcase Used

The Ruby script used to generate testcases for the previous OpenMP thread-based implementation is adapted to generate the various graphs to be used as input for the OpenMPI process-based implementation. To ensure fairness, we use the same number of threads as number of processes. In order to do this, the graph size thus changes for different number of processes since one process represents one edge. Since the graphs are generated by generating random Minimum Spanning Trees (MST), $e = v - 1$ where e is the number of edges and v is the number of vertices. Thus, we need to generate a graph with e edges, the graph must consist of $e + 1$ vertices. Note that each edge is unidirectional.

This time round, since we are generating smaller maps too, we reconfigured the thresholds of the graph generation to require 3 termini (vertices with degree 1) and that each line must have at least 2 stations. Maximum distance (maximum edge weight) between stations is 9.

The testcases all specified 10,000 time ticks. We ran testcases for 8, 16, 32, and 64 processes/threads. For the OpenMPI programme, we ran it on 8, 16 and 32 cores using a rankfile across 4 Xeon nodes in the lab. We ran each test twice, once with the per-tick status output enabled, and another time with it disabled.

Below you can find a sample input and visualisation of the adjacency matrix and the train lines for 8 edges/trains. All testcase input files and testcase generator script can all be found in the Appendix.

```

1      5
2      Sengkang,Bukit Panjang,Mattar,Damai,Botanic Gardens
3      0 0 0 0 6
4      0 0 0 8 9
5      0 0 0 0 5
6      0 8 0 0 0
7      6 9 5 0 0
8      0.8,0.6,0.5,0.8,1.0
9      Sengkang,Botanic Gardens,Bukit Panjang,Damai
10     Mattar,Botanic Gardens,Bukit Panjang,Damai
11     Sengkang,Botanic Gardens,Mattar
12     10000
13     3,3,2

```

Figure 1: Testcase for 8 edges/trains

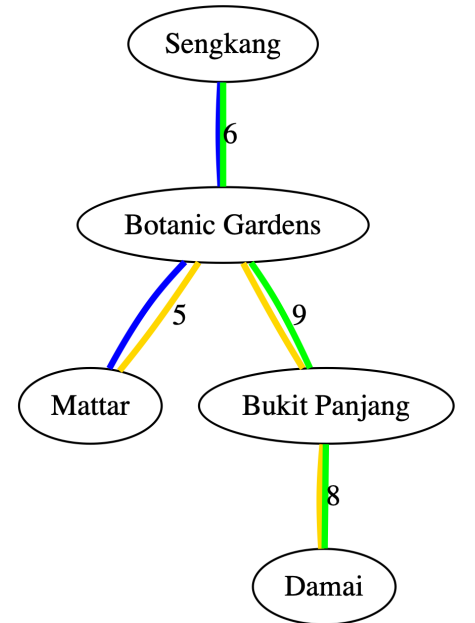


Figure 2: Map of the train lines for 8 edges/trains

3.2 Raw Data Collected

Legend:

- num_edge = Number of edges
- num_trains = Number of trains
- num_cores = Number of cores
- short = TRUE means no per-tick statistics is printed

num_edge / num_trains	num_cores	short	time
8	8	TRUE	1.152
8	8	FALSE	1.229
8	thread	TRUE	0.067
8	thread	FALSE	0.142
16	8	TRUE	1.466
16	8	FALSE	1.942
16	16	TRUE	10.216
16	16	FALSE	10.270
16	thread	TRUE	0.076
16	thread	FALSE	0.289
32	8	TRUE	2.004
32	8	FALSE	2.459
32	16	TRUE	12.244
32	16	FALSE	12.293
32	32	TRUE	15.611
32	32	FALSE	15.843
32	thread	TRUE	1.197
32	thread	FALSE	3.621
64	8	TRUE	3.640
64	8	FALSE	5.306
64	16	TRUE	382.949
64	16	FALSE	470.589
64	32	TRUE	39.227
64	32	FALSE	46.471
64	thread	TRUE	2.772
64	thread	FALSE	8.345

Table 1: Data collected

3.3 Comparison Summary

Number of cores	8 edges/trains	16 edges/trains	32 edges/trains	64 edges/trains
thread	0.067	0.076	1.197	2.772
8	1.152	1.466	2.004	3.64
16		10.216	12.244	382.949
32			15.611	39.227

Table 2: Summarised Comparison when per-tick statistics are not printed

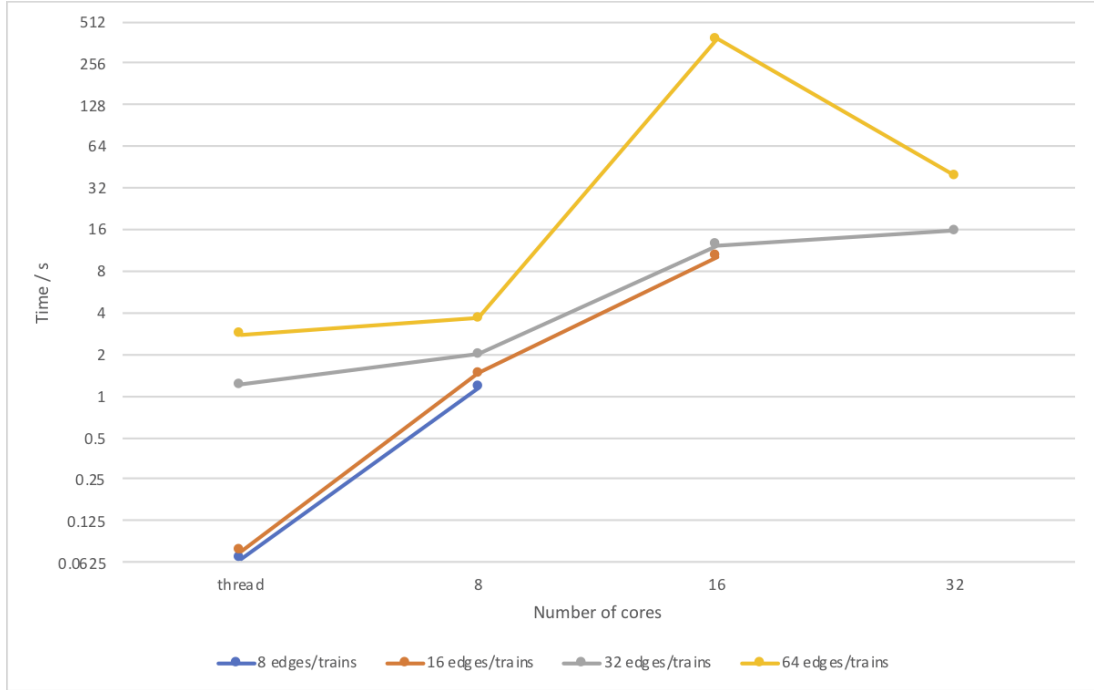


Figure 3: The comparison summary plotted

3.4 Additional results

We also tried running with 1000 trains for each line using the the map containing 64 edges. For fair comparison, we also disabled per-tick output. You can find the results below.

Configuration	time
65 processes locally	2.833
65 proceses over 32 cores in 4 nodes	42.253
Multithreaded	127.469

Table 3: Results of running map for 64 edges with 1000 trains

4 Discussion

It can be observed that the wall clock time taken for the simulation to complete increases as number of threads increases. This is because as number of threads increases, there is more contention of resources – in our case contention for each train (thread) in using links (edges) between stations (vertices) and contention for each train in opening door at each station. We are managing this using an implicit queue through implementing a timekeeper to track the next allowed event to occur, protected by marking that section as critical. As such, for each link (edge) or station (vertex), only one train (thread) is able to register itself to the timekeeper at any given time – others will have to wait.

However, we have an interesting observation as well. Notice that the execution time behaves very differently when the number of threads is beyond the number of logical cores (20 cores). For small input size (100 ticks), when the number of threads is beyond the number of logical cores, the execution time actually falls. However, as input size gets larger (1,000 and 10,000 ticks), execution time increases. This can be explained that when the number of threads are below the number of logical cores, all the threads are running concurrently, resulting in more lock contention (not to be confused with resource contention). However, when number of threads is above the number of logical cores, the threads take turns to wake up and do work, resulting in less lock contention. For smaller

input size, the lock contention time actually outweighs the execution time, resulting in the fall in execution time. However, for larger input size, there is a large overhead in context-switching which outweighs the effect of lock contention time. This is supported by the data we collected on number of context-switches for 100 ticks and 1,000 ticks.

We also observe another trend – that the variance in execution time falls when the number of threads exceed the number of logical cores. We currently have no explanation on this, but we suspect that the compiler does an optimisation when the number of threads exceed the number of logical cores.

5 Bonus

We chose to find the minimum number of trains per line required to maximise throughput in the network. We measured throughput by parsing the simulation output and counting the number of edge traversals by the trains. This optimisation was done on the example input map.

5.1 Notes

- Every edge traversal is counted as 1 regardless of the time taken.
- For convenience, partial traversals that occur because the simulation ends before they can complete are counted. This has an insignificant impact on the result since the maximum number of partial traversals is 16 (number of edges in the graph) while total traversals is in the order of $\sim 10,000$.
- Each line has the same number of trains regardless of line metrics (number of stations, distance between stations, station popularity)
- Simulations were run for 10000 ticks.
- All lines have the same number of trains.

5.2 Results

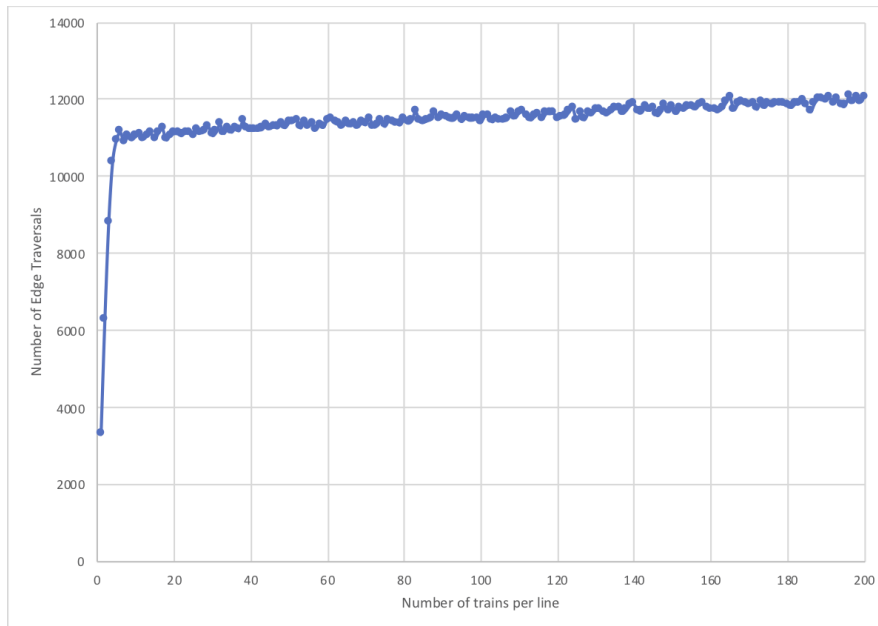


Figure 4: Scatter plot of number of edge traversals against number of trains per line

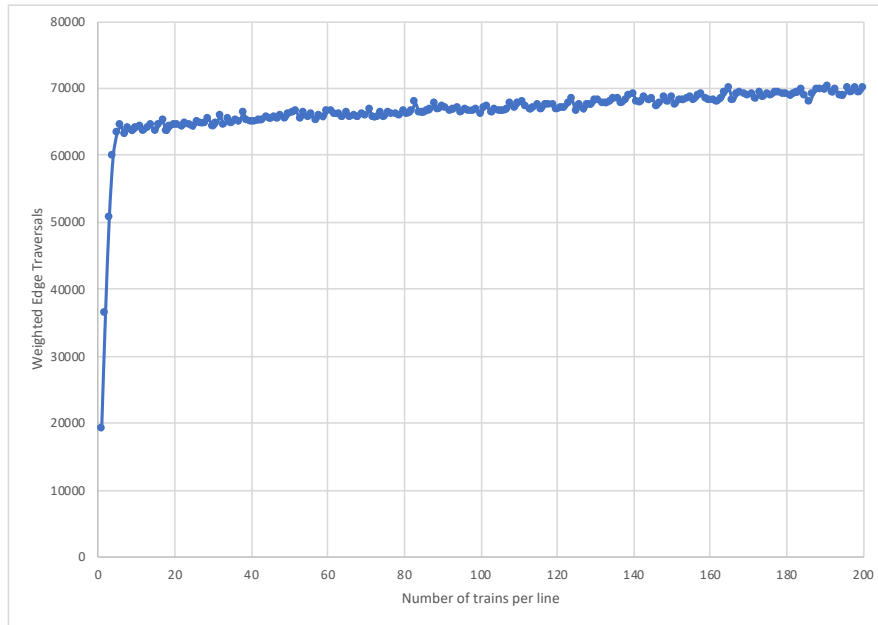


Figure 5: Scatter plot of weighted edge traversals against number of trains per line

Appendix A: Testcases and configuration files used

8 edges

```
1 5
2 Sengkang,Bukit Panjang,Mattar,Damai,Botanic Gardens
3 0 0 0 0 6
4 0 0 0 8 9
5 0 0 0 0 5
6 0 8 0 0 0
7 6 9 5 0 0
8 0.8,0.6,0.5,0.8,1.0
9 Sengkang,Botanic Gardens,Bukit Panjang,Damai
10 Mattar,Botanic Gardens,Bukit Panjang,Damai
11 Sengkang,Botanic Gardens,Mattar
12 10000
13 3,3,2
```

16 edges

```
1 9
2 Cashew,City Hall,Sixth Avenue,Fernvale,Kovan,Holland Village,Tuas
   ↪ Crescent,Sengkang,Orchard
3 0 0 9 0 0 0 0 0 0
4 0 0 0 0 0 0 0 0 9
5 9 0 0 7 0 1 7 0 0
6 0 0 7 0 0 0 0 6 0
7 0 0 0 0 0 3 0 0 0
8 0 0 1 0 3 0 0 0 1
9 0 0 7 0 0 0 0 0 0
10 0 0 0 6 0 0 0 0 0
11 0 9 0 0 0 1 0 0 0
12 0.6,0.4,0.3,0.5,0.9,1.0,0.7,1.0,0.7
13 City Hall,Orchard,Holland Village,Sixth Avenue,Tuas Crescent
14 City Hall,Orchard,Holland Village,Sixth Avenue,Fernvale,Sengkang
15 Cashew,Sixth Avenue,Fernvale,Sengkang
16 10000
17 5,5,6
```

32 edges

```
1 17
2 Jalan Besar,Buona Vista,Sixth Avenue,Punggol,Little India,Kovan,Farrer
   ↪ Park,Newton,Kupang,Phoenix,Kranji,Keat Hong,Bras
   ↪ Basah,MacPherson,Compassvale,Mattar,Samudera
3 0 0 0 0 0 1 0 0 0 0 0 0 0 0 0 0
4 0 0 0 0 0 0 0 0 0 2 0 0 0 0 0 0
5 0 0 0 0 1 0 0 0 2 0 0 0 0 0 0 6
6 0 0 0 0 0 9 0 0 0 0 0 0 0 0 0 0
7 0 0 1 0 0 0 0 0 0 5 5 0 0 4 8 0
8 1 0 0 9 0 0 0 0 0 6 0 6 3 0 0 0
9 0 0 0 0 0 0 5 0 0 0 0 0 0 0 0 0
10 0 0 0 0 0 0 5 0 0 0 4 0 0 0 0 0
```



```

11 0 0 2 0 0 0 0 0 0 0 0 0 0 0 0 0 0
12 0 2 0 0 0 0 0 0 0 0 0 9 0 0 0 0 0
13 0 0 0 0 5 6 0 4 0 9 0 0 0 0 0 0 0
14 0 0 0 0 5 0 0 0 0 0 0 0 0 0 0 0 0
15 0 0 0 0 0 6 0 0 0 0 0 0 0 0 0 0 0
16 0 0 0 0 0 3 0 0 0 0 0 0 0 0 0 0 0
17 0 0 0 0 4 0 0 0 0 0 0 0 0 0 0 0 0
18 0 0 0 0 8 0 0 0 0 0 0 0 0 0 0 0 0
19 0 0 6 0 0 0 0 0 0 0 0 0 0 0 0 0 0
20 0.7,0.1,1.0,0.7,0.1,1.0,0.3,0.9,0.8,1.0,1.0,0.6,0.9,0.2,0.2,0.4,0.8
21 Bras Basah,Kovan,MacPherson
22 Jalan Besar,Kovan,Punggol
23 Punggol,Kovan,Bras Basah
24 10000
25 11,11,10

```

64 edges

```

1 33
2 Punggol Point,Fernvale,Orchard,Joo Koon,Ten Mile
   ↳ Junction,Bartley,Kallang,Fajar,Bayfront,Nibong,Sengkang,Toa
   ↳ Payoh,MacPherson,Esplanade,Newton,Boon Lay,Sam Kee,Sengkang,Kangkar,Little India,Tai
   ↳ Seng,Tiong Bahru,Soo Teck,Geylang Bahru,Buona Vista,Choa Chu Kang,Cheng
   ↳ Lim,Bishan,Botanic Gardens,Dhoby Ghaut,Kent Ridge,Changi Airport,Tampines
3 0 1 0 0 0 0 9 0 0 0 0 0 0 0 0 2 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
4 1 0 0 7 0 0 0 0 0 0 0 0 0 0 0 6 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
5 0 0 0 0 0 0 0 0 0 8 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
6 0 7 0 0 0 8 0 0 0 4 0 0 0 0 0 0 6 9 0 0 9 0 8 0 0 0 0 0 0 0 0 7
7 0 0 0 0 0 0 3 0 0 0 0 0 0 0 0 0 0 0 0 8 0 0 0 0 0 0 0 0 0 0 0 0
8 0 0 0 8 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 4 0 4 0 0
9 9 0 0 0 0 0 0 0 0 3 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 6 0 0 0 0 0
10 0 0 0 0 3 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
11 0 0 0 0 0 0 0 0 0 0 0 0 0 0 2 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
12 0 0 8 4 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
13 0 0 0 0 0 0 3 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 1 0 0 0 0
14 0 0 0 0 0 0 0 0 0 0 0 0 4 4 7 0 0 0 0 0 4 0 0 0 0 0 0 0 0 0 0 0
15 0 0 0 0 0 0 0 0 0 0 4 0 0 0 0 3 0 0 0 0 0 0 0 0 0 0 0 0 0 6 0 0
16 0 0 0 0 0 0 0 0 0 0 4 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
17 2 0 0 0 0 0 0 0 2 0 0 7 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
18 0 6 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
19 0 0 0 0 0 0 0 0 0 0 0 0 3 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
20 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 6 0 0 0 0 0 0 0 0 0 0 0 0 0 0
21 0 0 0 6 0 0 0 0 0 0 0 0 0 0 0 0 6 0 0 0 0 0 0 0 0 0 0 0 0 0 0
22 0 0 0 9 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 6 0 0 0 0 0 0 0 0 0 0
23 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 5 0 0 0 0 3 0 0 0 0 0 0 0
24 0 0 0 0 8 0 0 0 0 0 4 0 0 0 0 0 0 5 0 0 0 0 0 0 0 0 0 0 0 0 0
25 0 0 0 9 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 5 0 0 0 0 0 0 0 0 0 0
26 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 6 0 0 0 0 0 0 0 0 0 0 0 0 0 0
27 0 0 0 8 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
28 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 5 0 0 0 0 0 0 0 0 0 0 0 0 0 0
29 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 3 0 0 0 0 0 0 0 0 0 0 0 0 0 0
30 0 0 0 0 0 0 6 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
31 0 0 0 0 0 4 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0

```

```

32 0 0 0 0 0 0 0 0 0 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
33 0 0 0 0 0 4 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
34 0 0 0 0 0 0 0 0 0 0 0 0 6 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
35 0 0 0 7 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
36 0.3,0.6,0.6,1.0,0.4,0.5,0.3,0.4,1.0,0.4,0.2,1.0,0.2,0.2,0.3,0.8,0.8,0.4,0.8,0.8,0.9,0.6,
   ↪ 0.4,0.9,0.7,0.7,0.4,0.8,0.1,0.5,0.3,0.7,0.5
37 Fajar,Ten Mile Junction,Tiong Bahru,Toa Payoh,Newton,Punggol Point,Fernvale,Joo
   ↪ Koon,Little India,Geylang Bahru
38 Sam Kee,MacPherson,Toa Payoh,Newton,Punggol Point,Fernvale,Joo Koon,Little India,Geylang
   ↪ Bahru
39 Kent Ridge,Bartley,Joo Koon,Fernvale,Punggol Point,Newton,Toa Payoh,MacPherson,Changi
   ↪ Airport
40 10000
41 21,21,22

```

Appendix B: Ruby script used to generate test cases

Below is the code listing of the ruby script used to generate test cases. Essentially, it does:

1. Create a random adjacency matrix with diagonal = 0
2. Find the MST of the random graph created
3. Ensure that there are enough vertices with degree = 1, else go back to step 1
4. Enumerate the 2-combinations of the vertices with degree = 1, and pick 3 randomly.
5. For each of the three 2-combinations, assign them to be the termini of each line.
6. Using breadth-first-search, find the path between the two vertices for each pair of termini.
7. Ensure that the path is long enough, else go back to step 4.

```

1 require 'set'
2
3 MIN_NUM_TERMINI = 3
4 MIN_NUM_STATIONS_LINE = 2

```

```

5 MRT_STATION_NAMES = ["Jurong East", "Bukit Batok", "Bukit Gombak", "Choa Chu Kang", "Yew
↳ Tee", "Kranji", "Marsiling", "Woodlands", "Admiralty", "Sembawang", "Yishun",
↳ "Khatib", "Yio Chu Kang", "Ang Mo Kio", "Bishan", "Braddell", "Toa Payoh", "Novena",
↳ "Newton", "Orchard", "Somerset", "Dhoby Ghaut", "City Hall", "Raffles Place", "Marina
↳ Bay", "Marina South Pier", "Pasir Ris", "Tampines", "Simei", "Tanah Merah", "Bedok",
↳ "Kembangan", "Eunos", "Paya Lebar", "Aljunied", "Kallang", "Lavender", "Bugis", "City
↳ Hall", "Raffles Place", "Tanjong Pagar", "Outram Park", "Tiong Bahru", "Redhill",
↳ "Queenstown", "Commonwealth", "Buona Vista", "Dover", "Clementi", "Jurong East",
↳ "Chinese Garden", "Lakeside", "Boon Lay", "Pioneer", "Joo Koon", "Gul Circle", "Tuas
↳ Crescent", "Tuas West Road", "Tuas Link", "Expo", "Changi Airport", "HarbourFront",
↳ "Outram Park", "Chinatown", "Clarke Quay", "Dhoby Ghaut", "Little India", "Farrer
↳ Park", "Boon Keng", "Potong Pasir", "Woodleigh", "Serangoon", "Kovan", "Hougang",
↳ "Buangkok", "Sengkang", "Punggol", "Dhoby Ghaut", "Bras Basah", "Esplanade",
↳ "Promenade", "Nicoll Highway", "Stadium", "Mountbatten", "Dakota", "Paya Lebar",
↳ "MacPherson", "Tai Seng", "Bartley", "Serangoon", "Lorong Chuan", "Bishan",
↳ "Marymount", "Caldecott", "Botanic Gardens", "Farrer Road", "Holland Village", "Buona
↳ Vista", "one-north", "Kent Ridge", "Haw Par Villa", "Pasir Panjang", "Labrador Park",
↳ "Telok Blangah", "HarbourFront", "Bayfront", "Marina Bay", "Bukit Panjang", "Cashew",
↳ "Hillview", "Beauty World", "King Albert Park", "Sixth Avenue", "Tan Kah Kee",
↳ "Botanic Gardens", "Stevens", "Newton", "Little India", "Rochor", "Bugis",
↳ "Promenade", "Bayfront", "Downtown", "Telok Ayer", "Chinatown", "Fort Canning",
↳ "Bencoolen", "Jalan Besar", "Bendemeer", "Geylang Bahru", "Mattar", "MacPherson",
↳ "Ubi", "Kaki Bukit", "Bedok North", "Bedok Reservoir", "Tampines West", "Tampines",
↳ "Tampines East", "Upper Changi", "Expo", "Choa Chu Kang", "South View", "Keat Hong",
↳ "Teck Whye", "Phoenix", "Bukit Panjang", "Petir", "Pending", "Bangkit", "Fajar",
↳ "Segar", "Jelapang", "Senja", "Ten Mile Junction", "Sengkang", "Compassvale",
↳ "Rumbia", "Bakau", "Kangkar", "Ranggung", "Cheng Lim", "Farmway", "Kupang",
↳ "Thanggam", "Fernvale", "Layar", "Tongkang", "Renjong", "Punggol", "Cove",
↳ "Meridian", "Coral Edge", "Riviera", "Kadaloor", "Oasis", "Damai", "Sam Kee", "Teck
↳ Lee", "Punggol Point", "Samudera", "Nibong", "Sumang", "Soo Teck"]

6
7 def generate_random_graph(s, max_weight)
8   Array.new(s) { |i| Array.new(s) { |j| i == j ? 0 : rand(1..max_weight) } }
9 end
10
11 def print_graph(matrix)
12   matrix.map { |row| row.join(' ') }.join("\n")
13 end
14
15 def prim(matrix)
16   cost = Array.new(matrix.length, Float::INFINITY)
17   parent = Array.new(matrix.length, nil)
18   visited = Array.new(matrix.length, false)
19
20   # start from the first vertex
21   cost[0] = 0
22   parent[0] = -1
23
24   matrix.length.times do
25     u = nil
26     min_weight = Float::INFINITY

```

```

29     # Find unvisited vertex with minimum cost
30     cost.zip(visited).each_with_index do |zipped, i|
31         c, v = zipped
32         if c < min_weight and !v
33             min_weight = c
34             u = i
35         end
36     end
37     visited[u] = true
38
39     matrix[u].zip(cost, visited).each_with_index do |zipped, i|
40         m, c, v = zipped
41         if m > 0 && !v && c > m
42             cost[i] = m
43             parent[i] = u
44         end
45     end
46 end
47
48 result = Array.new(matrix.length) { Array.new(matrix.length, 0) }
49
50 (1..matrix.length).each do |i|
51     result[i][parent[i]] = result[parent[i]][i] = matrix[i][parent[i]]
52 end
53
54 result
55 end
56
57 def bfs(matrix, termini)
58     from, to = termini
59     open_set = []
60     closed_set = Set[]
61     meta = {}
62
63     root = from
64     meta[root] = nil
65     open_set.unshift(root)
66
67     while !open_set.empty? do
68         subtree_root = open_set.shift
69         if subtree_root == to
70             return construct_path(subtree_root, meta)
71         end
72         matrix[subtree_root].each_with_index.select { |w, i| w > 0 }.map { |x| x.last }.each
73             ↪ do |child|
74                 next if closed_set.include?(child)
75                 if !open_set.include?(child)
76                     meta[child] = subtree_root
77                     open_set.unshift(child)
78                 end
79             end
80         closed_set.add(subtree_root)
81     end

```

```

81 end
82
83 def construct_path(state, meta)
84   result = [state]
85   while !meta[state].nil? do
86     state = meta[state]
87     result.append(state)
88   end
89   result.reverse
90 end
91
92 def permutate_sum(n)
93   (0..n).to_a.flat_map { |i| (0..(n - i)).to_a.map { |j| [i, j, n - i - j] } }
94 end
95
96 def usage_message
97   puts "Invalid args"
98   puts "Usage: ruby test_case_generator.rb <max_weight> <num_tick>"
99   exit 1
100 end
101
102 # Start of main
103 if ARGV.length != 2
104   usage_message
105 end
106
107 max_weight, tick = ARGV.map { |a| a.to_i }
108
109 if max_weight <= 0
110   usage_message
111 end
112
113
114 dir_name = "test-#{Time.now.strftime("%Y%m%d-%H%M")}"
115 Dir.mkdir(dir_name)
116
117 [[3, 3, 2], [5, 5, 6], [11, 11, 10], [21, 21, 22]].each do |trains|
118   n = trains.inject(&:+)
119   s = n / 2 + 1
120
121   puts n, s
122
123   primmed = nil
124   termini = []
125
126   while termini.length < MIN_NUM_TERMINI do
127     graph = generate_random_graph(s, max_weight)
128     primmed = prim(graph)
129     termini = primmed
130       .each_with_index.select do |row, i|
131         row.reduce(0) { |acc, weight| acc += weight > 0 ? 1 : 0 } == 1
132       end
133     .map { |pair| pair.last }

```

```

134 end
135
136 stations = MRT_STATION_NAMES.sample(s)
137 popularities = Array.new(s) { rand(1..10) / 10.0 }
138
139 green_line = []
140 yellow_line = []
141 blue_line = []
142 while green_line.length < MIN_NUM_STATIONS_LINE || yellow_line.length <
  ↪ MIN_NUM_STATIONS_LINE || blue_line.length < MIN_NUM_STATIONS_LINE do
143   green_termini, yellow_termini, blue_termini = termini.combination(2).to_a.sample(3)
144
145   green_line = bfs(primmed, green_termini)
146   yellow_line = bfs(primmed, yellow_termini)
147   blue_line = bfs(primmed, blue_termini)
148 end
149 File.open("#{dir_name}/#{n}_#{trains.join("-")}", "w") do |f|
150   f.puts s
151   f.puts stations.join(",")
152   f.puts print_graph(primmed)
153   f.puts popularities.join(",")
154   f.puts green_line.map { |s| stations[s] }.join(",")
155   f.puts yellow_line.map { |s| stations[s] }.join(",")
156   f.puts blue_line.map { |s| stations[s] }.join(",")
157   f.puts tick
158   f.puts trains.join(",")
159 end
160 end

```