

EECE7205 – Project 1

1. Pseudocode

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Max-min-grouping( $A, N, M$ )
   $C[0 \dots N][0 \dots M] \leftarrow -\infty$  // DP table to track max – min sum for each partitioning of elements
   $K[0 \dots N][0 \dots M] \leftarrow -1$  // Partition points table for backtracking the optimal grouping
   $C[0][0] \leftarrow \infty$ 
   $P[0 \dots N] \leftarrow 0$ 
  for  $i \leftarrow 1$  to  $N$  :
     $P[i] \leftarrow P[i - 1] + A[i]$ 
  for  $i \leftarrow 1$  to  $N$  :
     $C[i][1] \leftarrow P[i]$  // DP base case : if only 1 group, use total sum up to  $i$ 
     $K[i][1] \leftarrow 0$  // Group starts from index 0 when there is only one group
  for  $j \leftarrow 2$  to  $M$  : // Loop over number of groups, starting from 2 up to  $M$ 
    for  $i \leftarrow j$  to  $N$  : // Loop over elements to find maximum min – sum for each partition
      for  $k \leftarrow j - 1$  to  $i - 1$  : // Loop to test different partition points for the  $(j - 1)$  – th group
         $\text{sum} \leftarrow P[i] - P[k]$ 
         $\text{min\_sum} \leftarrow \min(C[k][j - 1], \text{sum})$  // Calculate min between previous max – min sum and current sum
        if  $\text{min\_sum} > C[i][j]$  then // If min – sum is greater update DP table
           $K[i][j] \leftarrow k$  // Update partition point in partition table  $k$  for backtracking
   $G[1 \dots M] \leftarrow 0$ 
   $\text{idx} \leftarrow N$ 
  for  $j \leftarrow M$  down to 1 : // Backtrack to find group sizes for each partition
     $k \leftarrow K[\text{idx}][j]$  // Retrieve partition point  $k$  for the  $j$  – th group
     $G[j] \leftarrow \text{idx} - k$ 
     $\text{idx} \leftarrow k$ 
  return  $G[1 \dots M]$ 
```

2. Running time analysis

$$P[i] = \sum_{m=1}^i A[m].$$

As each iteration requires a single addition, this step takes N operations, as the loop runs from $i = 1$ to N . The $\Theta(N)$ notation represents both the upper and lower bounds, since the loop must always execute N times.

Time Complexity: $\Theta(N)$

DP Table Base Case:

$$C[i][1] = P[i] \quad \text{for } i = 1 \text{ to } N.$$

As each assignment is $O(1)$, this step takes N operations, as the loop runs from $i = 1$ to N . The $\Theta(N)$ notation again represents both the upper and lower bounds, as the loop must always execute N times.

Time Complexity: $\Theta(N)$

DP Table Population:

Loop 1 (over groups j) runs from $j = 2$ to M , resulting in M iterations.

Loop 2 (over elements i) runs from $i = j$ to N resulting in N iterations.

For each combination of i and j , loop 3 (over partition points k) runs from $k = j - 1$ to $k = i - 1$, resulting in $O(N)$ iterations.

The three nested loops result in combined time complexity:

$$\sum_{j=2}^M \sum_{i=j}^N \sum_{k=j-1}^{i-1} O(1)$$

Loop 3 Sum (over k): For each fixed i and j , the number of iterations of k ranges from $j - 1$ to $i - 1$, which is $i - j + 1$:

$$\sum_{k=j-1}^{i-1} O(1) = O(i - j + 1)$$

Loop 2 Sum (over i): Summing $O(i - j + 1)$ over i from j to N :

$$\sum_{i=j}^N O(i - j + 1) = O\left(\sum_{i=j}^N (i - j + 1)\right) = O(N^2)$$

Loop 1 Sum (over j): Summing over j from 2 to M :

$$\sum_{j=2}^M O(N^2) = O(M \cdot N^2)$$

Time Complexity: $\Theta(M \cdot N^2)$

Backtracking to Retrieve the Optimal Partition:

The optimal grouping is constructed using the partitions table $K[i][j]$. Since each iteration requires only constant-time operations, this step has a time complexity of $O(M)$. The loop runs from $j = M$ down to 1, resulting in M iterations.

Time Complexity: $\Theta(M)$

Total running time complexity of algorithm:

$$\Theta(M \cdot N^2)$$

3. Input example results

1.

A = {3, 9, 7, 8, 2, 6, 5, 10, 1, 7, 6, 4}

N = 12

M = 3

- **Initialize DP Table for Base case $j = 1$ (Single group):**

$$C[i, 1] = \sum_{m=1}^i A[m]$$

$$\begin{aligned}
C[1, 1] &= \sum_{m=1}^1 A[m] = 3, & C[2, 1] &= \sum_{m=1}^2 A[m] = 12, & C[3, 1] &= \sum_{m=1}^3 A[m] = 19 \\
C[4, 1] &= \sum_{m=1}^4 A[m] = 27, & C[5, 1] &= \sum_{m=1}^5 A[m] = 29, & C[6, 1] &= \sum_{m=1}^6 A[m] = 35 \\
C[7, 1] &= \sum_{m=1}^7 A[m] = 40, & C[8, 1] &= \sum_{m=1}^8 A[m] = 50, & C[9, 1] &= \sum_{m=1}^9 A[m] = 51 \\
C[10, 1] &= \sum_{m=1}^{10} A[m] = 58, & C[11, 1] &= \sum_{m=1}^{11} A[m] = 64, & C[12, 1] &= \sum_{m=1}^{12} A[m] = 68
\end{aligned}$$

• **Compute DP Table for $j = 2$:**

$$C[i, 2] = \max_{1 \leq k \leq i-1} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^i A[m] \right) \right\} \quad \text{for } i = 2 \text{ to } 12$$

$C[1, 2] = \times$ (For $i = 1$ and $j = 2$, cannot partition one element into two groups. So $C[1, 2]$ is not applicable.)

$$C[2, 2] = \max_{1 \leq k \leq 1} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^2 A[m] \right) \right\} = \min(3, 9) = 3 \quad (\text{Selected } k = 1)$$

$$C[3, 2] = \max_{1 \leq k \leq 2} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^3 A[m] \right) \right\} = \max(3, 7) = 7 \quad (\text{Selected } k = 2)$$

$$C[4, 2] = \max_{1 \leq k \leq 3} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^4 A[m] \right) \right\} = \max(3, 12, 8) = 12 \quad (\text{Selected } k = 2)$$

$$C[5, 2] = \max_{1 \leq k \leq 4} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^5 A[m] \right) \right\} = \max(3, 12, 10, 2) = 12 \quad (\text{Selected } k = 2)$$

$$C[6, 2] = \max_{1 \leq k \leq 5} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^6 A[m] \right) \right\} = \max(3, 12, 16, 8, 6) = 16 \quad (\text{Selected } k = 3)$$

$$C[7, 2] = \max_{1 \leq k \leq 6} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^7 A[m] \right) \right\} = \max(3, 12, 19, 13, 11, 5) = 19 \quad (\text{Selected } k = 3)$$

$$C[8, 2] = \max_{1 \leq k \leq 7} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^8 A[m] \right) \right\} = \max(3, 12, 19, 23, 21, 15, 10) = 23 \quad (\text{Selected } k = 4)$$

$$C[9, 2] = \max_{1 \leq k \leq 8} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^9 A[m] \right) \right\} = \max(3, 12, 19, 24, 22, 16, 11, 1) = 24 \quad (\text{Selected } k = 4)$$

$$C[10, 2] = \max_{1 \leq k \leq 9} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^{10} A[m] \right) \right\} = \max(3, 12, 19, 27, 29, 23, 18, 8, 7) = 29 \quad (\text{Selected } k = 5)$$

$$C[11, 2] = \max_{1 \leq k \leq 10} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^{11} A[m] \right) \right\} = \max(3, 12, 19, 27, 29, 29, 24, 14, 13, 6) = 29 \quad (\text{Selected } k = 5)$$

$$C[12, 2] = \max_{1 \leq k \leq 11} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^{12} A[m] \right) \right\} = \max(3, 12, 19, 27, 29, 33, 28, 18, 17, 10, 4) = 33 \quad (\text{Selected } k = 6)$$

- Compute DP Table for $j = 3$:

$$C[i, 3] = \max_{2 \leq k \leq i-1} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^i A[m] \right) \right\} \quad \text{for } i = 3 \text{ to } 12$$

$C[1, 3] = \times$ (For $i = 1$ and $j = 3$, cannot partition one element into three groups. So $C[1, 3]$ is not applicable.)

$C[2, 3] = \times$ (For $i = 2$ and $j = 3$, cannot partition two elements into three groups. So $C[2, 3]$ is not applicable.)

$$C[3, 3] = \max_{2 \leq k \leq 2} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^3 A[m] \right) \right\} = \min(3, 7) = 3 \quad (\text{Selected } k = 2)$$

$$C[4, 3] = \max_{2 \leq k \leq 3} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^4 A[m] \right) \right\} = \max(3, 7) = 7 \quad (\text{Selected } k = 3)$$

$$C[5, 3] = \max_{2 \leq k \leq 4} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^5 A[m] \right) \right\} = \max(3, 7, 2) = 7 \quad (\text{Selected } k = 3)$$

$$C[6, 3] = \max_{2 \leq k \leq 5} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^6 A[m] \right) \right\} = \max(3, 7, 8, 6) = 8 \quad (\text{Selected } k = 4)$$

$$C[7, 3] = \max_{2 \leq k \leq 6} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^7 A[m] \right) \right\} = \max(3, 7, 12, 11, 5) = 12 \quad (\text{Selected } k = 4)$$

$$C[8, 3] = \max_{2 \leq k \leq 7} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^8 A[m] \right) \right\} = \max(3, 7, 12, 12, 15, 10) = 15 \quad (\text{Selected } k = 6)$$

$$C[9, 3] = \max_{2 \leq k \leq 8} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^9 A[m] \right) \right\} = \max(3, 7, 12, 12, 16, 11, 1) = 16 \quad (\text{Selected } k = 6)$$

$$C[10, 3] = \max_{2 \leq k \leq 9} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^{10} A[m] \right) \right\} = \max(3, 7, 12, 12, 16, 18, 8, 7) = 18 \quad (\text{Selected } k = 7)$$

$$C[11, 3] = \max_{2 \leq k \leq 10} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^{11} A[m] \right) \right\} = \max(3, 7, 12, 12, 16, 19, 14, 13, 6) = 19 \quad (\text{Selected } k = 7)$$

$$C[12, 3] = \max_{2 \leq k \leq 11} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^{12} A[m] \right) \right\} = \max(3, 7, 12, 12, 16, 19, 18, 17, 10, 4) = 19 \quad (\text{Selected } k = 7)$$

Table 1: DP Table $C[i][j]$

$j \setminus i$	1	2	3	4	5	6	7	8	9	10	11	12
1	3	12	19	27	29	35	40	50	51	58	64	68
2	–	3	7	12	12	16	19	23	24	29	29	33
3	–	–	3	7	7	8	12	15	16	18	19	19

Table 2: Partition Points $K[i][j]$

$j \setminus i$	1	2	3	4	5	6	7	8	9	10	11	12
1	0	0	0	0	0	0	0	0	0	0	0	0
2	–	1	2	2	2	3	3	4	4	5	5	6
3	–	–	2	3	3	4	4	6	6	7	7	7

- Maximum minimal sum achievable by partitioning $A[1...12]$ into $M = 3$ groups is $C[12, 3] = 19$.

- Backtracking through the DP table to determine the optimal grouping $G[1 \dots M]$ by finding partition points k that result in $C[12, 3] = 19$.

Partitioning $A[1..12]$ into optimal groups based on $k=7$ and $k=3$:

- Group 1: $A[1..3] = [3, 9, 7]$
- Group 2: $A[4..7] = [8, 2, 6, 5]$
- Group 3: $A[8..12] = [10, 1, 7, 6, 4]$

Grouping array $G[1 \dots 3]$: $G = [3, 4, 5]$

- Determine summation of the elements in the j -th group of array A .

- Group 1 ($A[1 \dots 3]$): $B[1] = 3 + 9 + 7 = 19$
- Group 2 ($A[4 \dots 7]$): $B[2] = 8 + 2 + 6 + 5 = 21$
- Group 3 ($A[8 \dots 12]$): $B[3] = 10 + 1 + 7 + 6 + 4 = 28$

Sum of elements in each group $B[1 \dots 3]$: $B = [19, 21, 28]$

- The minimal group sum:

$$\min_{1 \leq j \leq 3} B[j] = \min(19, 21, 28) = 19$$

The minimal group sum matches $C[12, 3] = 19$, confirming that the minimal group sum is maximized.

Optimal Grouping Sizes ($G[1..M]$):

Group 1: 3 elements

Group 2: 4 elements

Group 3: 5 elements

Group Details:

Group 1: $A[1..3] = [3, 9, 7]$, Sum = 19

Group 2: $A[4..7] = [8, 2, 6, 5]$, Sum = 21

Group 3: $A[8..12] = [10, 1, 7, 6, 4]$, Sum = 28

Optimal Grouping G : 3 4 5

.....

2.

$A = \{4, 2, 5, 1, 6, 7, 3, 8, 2, 4\}$

$N = 10$

$M = 3$

- **Initialize DP Table for Base case $j = 1$ (Single group):**

$$C[i, 1] = \sum_{m=1}^i A[m]$$

$$C[1, 1] = \sum_{m=1}^1 A[m] = 4, \quad C[2, 1] = \sum_{m=1}^2 A[m] = 6, \quad C[3, 1] = \sum_{m=1}^3 A[m] = 11,$$

$$C[4, 1] = \sum_{m=1}^4 A[m] = 12, \quad C[5, 1] = \sum_{m=1}^5 A[m] = 18, \quad C[6, 1] = \sum_{m=1}^6 A[m] = 25,$$

$$C[7, 1] = \sum_{m=1}^7 A[m] = 28, \quad C[8, 1] = \sum_{m=1}^8 A[m] = 36, \quad C[9, 1] = \sum_{m=1}^9 A[m] = 38,$$

$$C[10, 1] = \sum_{m=1}^{10} A[m] = 42.$$

• **Compute DP Table for $j = 2$:**

$$C[i, 2] = \max_{1 \leq k \leq i-1} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^i A[m] \right) \right\}, \quad \text{for } i = 2 \text{ to } 10$$

$C[1, 2] = \times$ (For $i = 1$ and $j = 2$, cannot partition one element into two groups. So $C[1, 2]$ is not applicable.)

$$C[2, 2] = \max_{1 \leq k \leq 1} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^2 A[m] \right) \right\} = \min(2, 5) = 2 \quad (\text{Selected } k = 1)$$

$$C[3, 2] = \max_{1 \leq k \leq 2} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^3 A[m] \right) \right\} = \max(4, 5) = 5 \quad (\text{Selected } k = 2)$$

$$C[4, 2] = \max_{1 \leq k \leq 3} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^4 A[m] \right) \right\} = \max(4, 6, 1) = 6 \quad (\text{Selected } k = 2)$$

$$C[5, 2] = \max_{1 \leq k \leq 4} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^5 A[m] \right) \right\} = \max(4, 6, 7, 6) = 7 \quad (\text{Selected } k = 3)$$

$$C[6, 2] = \max_{1 \leq k \leq 5} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^6 A[m] \right) \right\} = \max(4, 6, 11, 12, 7) = 12 \quad (\text{Selected } k = 4)$$

$$C[7, 2] = \max_{1 \leq k \leq 6} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^7 A[m] \right) \right\} = \max(4, 6, 11, 12, 10, 3) = 12 \quad (\text{Selected } k = 4)$$

$$C[8, 2] = \max_{1 \leq k \leq 7} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^8 A[m] \right) \right\} = \max(4, 6, 11, 12, 18, 11, 8) = 18 \quad (\text{Selected } k = 5)$$

$$C[9, 2] = \max_{1 \leq k \leq 8} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^9 A[m] \right) \right\} = \max(4, 6, 11, 12, 18, 13, 10, 2) = 18 \quad (\text{Selected } k = 5)$$

$$C[10, 2] = \max_{1 \leq k \leq 9} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^{10} A[m] \right) \right\} = \max(4, 6, 11, 12, 18, 17, 14, 6, 4) = 18 \quad (\text{Selected } k = 5)$$

• **Compute DP Table for $j = 3$:**

$$C[i, 3] = \max_{k=2}^{i-1} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^i A[m] \right) \right\}, \quad \text{for } i = 3 \text{ to } 10$$

$C[1, 3] = \times$ (For $i = 1$ and $j = 3$, cannot partition one element into three groups. So $C[1, 3]$ is not applicable.)
 $C[2, 3] = \times$ (For $i = 2$ and $j = 3$, cannot partition two elements into three groups. So $C[2, 3]$ is not applicable.)

$$C[3, 3] = \max_{2 \leq k \leq 2} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^3 A[m] \right) \right\} = \min(2, 5) = 2 \quad (\text{Selected } k = 2)$$

$$C[4, 3] = \max_{2 \leq k \leq 3} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^4 A[m] \right) \right\} = \max \{2, 1\} = 2 \quad (\text{Selected } k = 2)$$

$$C[5, 3] = \max_{2 \leq k \leq 4} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^5 A[m] \right) \right\} = \max \{2, 5, 6\} = 6 \quad (\text{Selected } k = 4)$$

$$C[6, 3] = \max_{2 \leq k \leq 5} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^6 A[m] \right) \right\} = \max \{2, 5, 6, 7\} = 7 \quad (\text{Selected } k = 5)$$

$$C[7, 3] = \max_{2 \leq k \leq 6} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^7 A[m] \right) \right\} = \max \{2, 5, 6, 7, 3\} = 7 \quad (\text{Selected } k = 5)$$

$$C[8, 3] = \max_{2 \leq k \leq 7} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^8 A[m] \right) \right\} = \max \{2, 5, 6, 7, 11, 8\} = 11 \quad (\text{Selected } k = 6)$$

$$C[9, 3] = \max_{2 \leq k \leq 8} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^9 A[m] \right) \right\} = \max \{2, 5, 6, 7, 12, 10, 2\} = 12 \quad (\text{Selected } k = 6)$$

$$C[10, 3] = \max_{2 \leq k \leq 9} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^{10} A[m] \right) \right\} = \max \{2, 5, 6, 7, 12, 12, 6, 4\} = 12 \quad (\text{Selected } k = 6)$$

Table 3: DP Table $C[i][j]$

$j \backslash i$	1	2	3	4	5	6	7	8	9	10
1	4	6	11	12	18	25	28	36	38	42
2	–	2	5	6	7	12	12	18	18	18
3	–	–	2	2	6	7	7	11	12	12

Table 4: Partition Points $K[i][j]$

$j \backslash i$	1	2	3	4	5	6	7	8	9	10
1	0	0	0	0	0	0	0	0	0	0
2	–	1	2	2	3	4	4	5	5	5
3	–	–	2	2	4	5	5	6	6	6

- Maximum minimal sum achievable by partitioning $A[1 \dots 10]$ ($i=10$) into $M = 3$ ($j=3$) groups is $C[10, 3] = 12$.
- Backtracking through the DP table to determine optimal grouping $G[1 \dots M]$ by finding partition points k that leads to $C[10, 3] = 12$.

Partitioning $A[1 \dots 10]$ into optimal groups based on $k=4$ and $k=6$:

- Group 1: $A[1 \dots 4] = [4, 2, 5, 1]$
- Group 2: $A[5 \dots 6] = [6, 7]$
- Group 3: $A[7 \dots 10] = [3, 8, 2, 4]$

Grouping array $G[1 \dots 3]$: $G = [4, 2, 4]$

- Determine summation of the elements in the j -th group of array A .
 - Group 1 ($A[1 \dots 4]$): $B[1] = 4 + 2 + 5 + 1 = 12$
 - Group 2 ($A[4 \dots 7]$): $B[2] = 6 + 7 = 13$
 - Group 3 ($A[8 \dots 12]$): $B[3] = 3 + 8 + 2 + 4 = 17$

Sum of elements in each group $B[1 \dots 3]$: $B = [12, 13, 17]$

- The minimal group sum:

$$\min_{1 \leq j \leq 3} B[j] = \min(12, 13, 17) = 12$$

The minimal group sum matches $C[10, 3] = 12$, confirming that the minimal group sum is maximized.

Optimal Grouping Sizes ($G[1..M]$):

Group 1: 4 elements

Group 2: 2 elements

Group 3: 4 elements

Group Details:

Group 1: $A[1 \dots 4] = [4, 2, 5, 1]$, Sum = 12

Group 2: $A[5 \dots 6] = [6, 7]$, Sum = 13

Group 3: $A[7 \dots 10] = [3, 8, 2, 4]$, Sum = 17

Optimal Grouping G : 4 2 4

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3.

$A = \{5, 2, 4, 7, 1, 3, 6, 8, 2, 4, 9\}$

$N = 11$

$M = 4$

- **Initialize DP Table for Base case $j = 1$ (Single group):**

$$C[i, 1] = \sum_{m=1}^i A[m]$$

$$C[1, 1] = \sum_{m=1}^1 A[m] = 5, \quad C[2, 1] = \sum_{m=1}^2 A[m] = 7, \quad C[3, 1] = \sum_{m=1}^3 A[m] = 11,$$

$$C[4, 1] = \sum_{m=1}^4 A[m] = 18, \quad C[5, 1] = \sum_{m=1}^5 A[m] = 19, \quad C[6, 1] = \sum_{m=1}^6 A[m] = 22,$$

$$C[7, 1] = \sum_{m=1}^7 A[m] = 28, \quad C[8, 1] = \sum_{m=1}^8 A[m] = 36, \quad C[9, 1] = \sum_{m=1}^9 A[m] = 38,$$

$$C[10, 1] = \sum_{m=1}^{10} A[m] = 42, \quad C[11, 1] = \sum_{m=1}^{11} A[m] = 51$$

- **Compute DP Table for $j = 2$:**

$$C[i, 2] = \max_{1 \leq k \leq i-1} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^i A[m] \right) \right\}, \quad \text{for } i = 2 \text{ to } 11$$

$C[1, 2] = \times$ (For $i = 1$ and $j = 2$, cannot partition 1 element into 2 groups. So $C[1, 2]$ is not applicable.)

$$C[2, 2] = \max_{1 \leq k \leq 1} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^2 A[m] \right) \right\} = \min(5, 2) = 2 \quad (\text{Selected } k = 1)$$

$$C[3, 2] = \max_{1 \leq k \leq 2} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^3 A[m] \right) \right\} = \max(5, 4) = 5 \quad (\text{Selected } k = 1)$$

$$C[4, 2] = \max_{1 \leq k \leq 3} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^4 A[m] \right) \right\} = \max(5, 7, 7) = 7 \quad (\text{Selected } k = 2)$$

$$C[5, 2] = \max_{1 \leq k \leq 4} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^5 A[m] \right) \right\} = \max(5, 7, 8, 1) = 8 \quad (\text{Selected } k = 3)$$

$$C[6, 2] = \max_{1 \leq k \leq 5} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^6 A[m] \right) \right\} = \max(5, 7, 11, 4, 3) = 11 \quad (\text{Selected } k = 3)$$

$$C[7, 2] = \max_{1 \leq k \leq 6} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^7 A[m] \right) \right\} = \max(5, 7, 11, 10, 9, 6) = 11 \quad (\text{Selected } k = 3)$$

$$C[8, 2] = \max_{1 \leq k \leq 7} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^8 A[m] \right) \right\} = \max(5, 7, 11, 18, 17, 14, 8) = 18 \quad (\text{Selected } k = 4)$$

$$C[9, 2] = \max_{1 \leq k \leq 8} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^9 A[m] \right) \right\} = \max(5, 7, 11, 18, 19, 16, 10, 2) = 19 \quad (\text{Selected } k = 5)$$

$$C[10, 2] = \max_{1 \leq k \leq 9} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^{10} A[m] \right) \right\} = \max(5, 7, 11, 18, 19, 20, 14, 6, 4) = 20 \quad (\text{Selected } k = 6)$$

$$C[11, 2] = \max_{1 \leq k \leq 10} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^{11} A[m] \right) \right\} = \max(5, 7, 11, 18, 19, 22, 23, 15, 13, 9) = 23 \quad (\text{Selected } k = 7)$$

• **Compute DP Table for $j = 3$:**

$$C[i, 3] = \max_{2 \leq k \leq i-1} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^i A[m] \right) \right\}, \quad \text{for } i = 3 \text{ to } 11$$

$C[1, 3] = \times$ (For $i = 1$ and $j = 3$, cannot partition one element into three groups. So $C[1, 3]$ is not applicable.)
 $C[2, 3] = \times$ (For $i = 2$ and $j = 3$, cannot partition two elements into three groups. So $C[2, 3]$ is not applicable.)

$$C[3, 3] = \max_{2 \leq k \leq 2} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^3 A[m] \right) \right\} = \min(2, 4) = 2 \quad (\text{Selected } k = 2)$$

$$C[4, 3] = \max_{2 \leq k \leq 3} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^4 A[m] \right) \right\} = \max(2, 5) = 5 \quad (\text{Selected } k = 3)$$

$$C[5, 3] = \max_{2 \leq k \leq 4} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^5 A[m] \right) \right\} = \max(2, 5, 1) = 5 \quad (\text{Selected } k = 3)$$

$$C[6, 3] = \max_{2 \leq k \leq 5} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^6 A[m] \right) \right\} = \max(2, 5, 4, 3) = 5 \quad (\text{Selected } k = 3)$$

$$C[7, 3] = \max_{2 \leq k \leq 6} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^7 A[m] \right) \right\} = \max(2, 5, 7, 8, 6) = 8 \quad (\text{Selected } k = 5)$$

$$C[8, 3] = \max_{2 \leq k \leq 7} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^8 A[m] \right) \right\} = \max(2, 5, 7, 8, 11, 8) = 11 \quad (\text{Selected } k = 6)$$

$$C[9, 3] = \max_{2 \leq k \leq 8} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^9 A[m] \right) \right\} = \max(2, 5, 7, 8, 11, 10, 2) = 11 \quad (\text{Selected } k = 6)$$

$$C[10, 3] = \max_{2 \leq k \leq 9} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^{10} A[m] \right) \right\} = \max(2, 5, 7, 8, 11, 11, 6, 4) = 11 \quad (\text{Selected } k = 6)$$

$$C[11, 3] = \max_{2 \leq k \leq 10} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^{11} A[m] \right) \right\} = \max(2, 5, 7, 8, 11, 11, 15, 13, 9) = 15 \quad (\text{Selected } k = 8)$$

• **Compute DP Table for $j = 4$:**

$$C[i, 4] = \max_{3 \leq k \leq i-1} \left\{ \min \left(C[k, 3], \sum_{m=k+1}^i A[m] \right) \right\}, \quad \text{for } i = 4 \text{ to } 11$$

$C[1, 4] = \times$ (For $i = 1$ and $j = 4$, cannot partition 1 element into 2 groups. So $C[1, 3]$ is not applicable.)
 $C[2, 4] = \times$ (For $i = 2$ and $j = 4$, cannot partition 2 elements into 3 groups. So $C[2, 3]$ is not applicable.)
 $C[3, 4] = \times$ (For $i = 3$ and $j = 4$, cannot partition 3 elements into 4 groups. So $C[3, 3]$ is not applicable.)

$$\begin{aligned}
C[4, 4] &= \max_{3 \leq k \leq 3} \left\{ \min \left(C[k, 3], \sum_{m=k+1}^4 A[m] \right) \right\} = 2 \quad (\text{Selected } k = 3) \\
C[5, 4] &= \max_{3 \leq k \leq 4} \left\{ \min \left(C[k, 3], \sum_{m=k+1}^5 A[m] \right) \right\} = \max(2, 1) = 2 \quad (\text{Selected } k = 3) \\
C[6, 4] &= \max_{3 \leq k \leq 5} \left\{ \min \left(C[k, 3], \sum_{m=k+1}^6 A[m] \right) \right\} = \max(2, 4, 3) = 4 \quad (\text{Selected } k = 4) \\
C[7, 4] &= \max_{3 \leq k \leq 6} \left\{ \min \left(C[k, 3], \sum_{m=k+1}^7 A[m] \right) \right\} = \max(2, 5, 5, 5) = 5 \quad (\text{Selected } k = 4) \\
C[8, 4] &= \max_{3 \leq k \leq 7} \left\{ \min \left(C[k, 3], \sum_{m=k+1}^8 A[m] \right) \right\} = \max(2, 5, 5, 5, 8) = 8 \quad (\text{Selected } k = 7) \\
C[9, 4] &= \max_{3 \leq k \leq 8} \left\{ \min \left(C[k, 3], \sum_{m=k+1}^9 A[m] \right) \right\} = \max(2, 5, 5, 5, 8, 2) = 8 \quad (\text{Selected } k = 7) \\
C[10, 4] &= \max_{3 \leq k \leq 9} \left\{ \min \left(C[k, 3], \sum_{m=k+1}^{10} A[m] \right) \right\} = \max(2, 5, 5, 5, 8, 6, 4) = 8 \quad (\text{Selected } k = 7) \\
C[11, 4] &= \max_{3 \leq k \leq 10} \left\{ \min \left(C[k, 3], \sum_{m=k+1}^{11} A[m] \right) \right\} = \max(2, 5, 5, 5, 8, 11, 11, 9) = 11 \quad (\text{Selected } k = 8)
\end{aligned}$$

Table 5: DP Table $C[i][j]$

$j \setminus i$	1	2	3	4	5	6	7	8	9	10	11
1	5	7	11	18	19	22	28	36	38	42	51
2	–	2	5	7	8	11	11	18	19	20	23
3	–	–	2	5	5	5	8	11	11	11	15
4	–	–	–	2	2	4	5	8	8	8	11

Table 6: Partition Points $K[i][j]$

$j \setminus i$	1	2	3	4	5	6	7	8	9	10	11
1	0	0	0	0	0	0	0	0	0	0	0
2	–	1	1	2	3	3	3	4	5	6	7
3	–	–	2	3	3	3	5	6	6	6	8
4	–	–	–	3	3	4	4	7	7	7	8

- Maximum minimal sum achievable by partitioning $A[1 \dots 11]$ into $M = 4$ groups is $C[11, 4] = 11$.
- Backtracking through the DP table to determine optimal grouping $G[1 \dots M]$ by finding partition points k that leads to $C[11, 4] = 11$.

Partitioning $A[1 \dots 11]$ into optimal groups based on $k=3$, $k=6$ and $k=8$:

- Group 1: $A[1 \dots 3] = [5, 2, 4]$
- Group 2: $A[4 \dots 6] = [7, 1, 3]$

- Group 3: $A[7 \dots 8] = [6, 8]$
- Group 4: $A[9 \dots 11] = [2, 4, 9]$

Grouping array $G[1 \dots 3]$: $G = [3, 3, 2, 3]$

- Determine summation of the elements in the j -th group of array A .
 - Group 1 ($A[1 \dots 3]$): $B[1] = 5 + 2 + 4 = 11$
 - Group 2 ($A[4 \dots 6]$): $B[2] = 7 + 1 + 3 = 11$
 - Group 3 ($A[7 \dots 8]$): $B[3] = 6 + 8 = 14$
 - Group 3 ($A[9 \dots 11]$): $B[4] = 2 + 4 + 9 = 15$

Sum of elements in each group $B[1 \dots 4]$: $B = [11, 11, 14, 15]$

- The minimal group sum:

$$\min_{1 \leq j \leq 4} B[j] = \min(11, 11, 14, 15) = 11$$

The minimal group sum matches $C[11, 4] = 11$, confirming that the minimal group sum is maximized.

Optimal Grouping Sizes ($G[1..M]$):

Group 1: 3 elements

Group 2: 3 elements

Group 3: 2 elements

Group 4: 3 elements

Group Details:

Group 1: $A[1 \dots 3] = [5, 2, 4]$, Sum = 11

Group 2: $A[4 \dots 6] = [7, 1, 3]$, Sum = 11

Group 3: $A[7 \dots 8] = [6, 8]$, Sum = 14

Group 4: $A[9 \dots 11] = [2, 4, 9]$, Sum = 15

Optimal Grouping G : 3 3 2 3

■

4.

$A = \{3, 1, 4, 1, 5, 9, 2, 6, 5, 3\}$

$N = 10$

$M = 3$

- **Initialize DP Table for Base case $j = 1$ (Single group):**

$$C[i, 1] = \sum_{m=1}^i A[m]$$

$$C[1, 1] = \sum_{m=1}^1 A[m] = 3, \quad C[2, 1] = \sum_{m=1}^2 A[m] = 4, \quad C[3, 1] = \sum_{m=1}^3 A[m] = 8,$$

$$C[4, 1] = \sum_{m=1}^4 A[m] = 9, \quad C[5, 1] = \sum_{m=1}^5 A[m] = 14, \quad C[6, 1] = \sum_{m=1}^6 A[m] = 23,$$

$$C[7, 1] = \sum_{m=1}^7 A[m] = 25, \quad C[8, 1] = \sum_{m=1}^8 A[m] = 31, \quad C[9, 1] = \sum_{m=1}^9 A[m] = 36,$$

$$C[10, 1] = \sum_{m=1}^{10} A[m] = 39$$

• **Compute DP Table for $j = 2$:**

$$C[i, 2] = \max_{1 \leq k \leq i-1} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^i A[m] \right) \right\}, \quad \text{for } i = 2 \text{ to } 10$$

$C[1, 2] = \times$ (For $i = 1$ and $j = 2$, cannot partition 1 element into 2 groups. So $C[1, 2]$ is not applicable.)

$$C[2, 2] = \max_{1 \leq k \leq 1} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^2 A[m] \right) \right\} = \min(3, 1) = 1 \quad (\text{Selected } k = 1)$$

$$C[3, 2] = \max_{1 \leq k \leq 2} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^3 A[m] \right) \right\} = \max(3, 4) = 4 \quad (\text{Selected } k = 2)$$

$$C[4, 2] = \max_{1 \leq k \leq 3} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^4 A[m] \right) \right\} = \max(3, 4, 1) = 4 \quad (\text{Selected } k = 2)$$

$$C[5, 2] = \max_{1 \leq k \leq 4} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^5 A[m] \right) \right\} = \max(3, 4, 6, 5) = 6 \quad (\text{Selected } k = 3)$$

$$C[6, 2] = \max_{1 \leq k \leq 5} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^6 A[m] \right) \right\} = \max(3, 4, 8, 9, 9) = 9 \quad (\text{Selected } k = 4)$$

$$C[7, 2] = \max_{1 \leq k \leq 6} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^7 A[m] \right) \right\} = \max(3, 4, 8, 9, 11, 2) = 11 \quad (\text{Selected } k = 5)$$

$$C[8, 2] = \max_{1 \leq k \leq 7} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^8 A[m] \right) \right\} = \max(3, 4, 8, 9, 14, 8, 6) = 14 \quad (\text{Selected } k = 5)$$

$$C[9, 2] = \max_{1 \leq k \leq 8} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^9 A[m] \right) \right\} = \max(3, 4, 8, 9, 14, 13, 11, 5) = 14 \quad (\text{Selected } k = 5)$$

$$C[10, 2] = \max_{1 \leq k \leq 9} \left\{ \min \left(C[k, 1], \sum_{m=k+1}^{10} A[m] \right) \right\} = \max(3, 4, 8, 9, 14, 16, 14, 8, 3) = 16 \quad (\text{Selected } k = 6)$$

• **Compute DP Table for $j = 3$:**

$$C[i, 3] = \max_{2 \leq k \leq i-1} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^i A[m] \right) \right\}, \quad \text{for } i = 3 \text{ to } 10$$

$C[1, 3] = \times$ (For $i = 1$ and $j = 3$, cannot partition one element into three groups. So $C[1, 3]$ is not applicable.)
 $C[2, 3] = \times$ (For $i = 2$ and $j = 3$, cannot partition two elements into three groups. So $C[2, 3]$ is not applicable.)

$$C[3, 3] = \max_{2 \leq k \leq 2} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^3 A[m] \right) \right\} = \min(1, 4) = 1 \quad (\text{Selected } k = 2)$$

$$C[4, 3] = \max_{2 \leq k \leq 3} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^4 A[m] \right) \right\} = \max(1, 1) = 1 \quad (\text{Selected } k = 2)$$

$$C[5, 3] = \max_{2 \leq k \leq 4} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^5 A[m] \right) \right\} = \max(1, 4, 4) = 4 \quad (\text{Selected } k = 3)$$

$$C[6, 3] = \max_{2 \leq k \leq 5} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^6 A[m] \right) \right\} = \max(1, 4, 4, 6) = 6 \quad (\text{Selected } k = 5)$$

$$C[7, 3] = \max_{2 \leq k \leq 6} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^7 A[m] \right) \right\} = \max(1, 4, 4, 6, 2) = 6 \quad (\text{Selected } k = 5)$$

$$C[8, 3] = \max_{2 \leq k \leq 7} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^8 A[m] \right) \right\} = \max(1, 4, 4, 6, 8, 6) = 8 \quad (\text{Selected } k = 6)$$

$$C[9, 3] = \max_{2 \leq k \leq 8} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^9 A[m] \right) \right\} = \max(1, 4, 4, 6, 9, 11, 5) = 11 \quad (\text{Selected } k = 7)$$

$$C[10, 3] = \max_{2 \leq k \leq 9} \left\{ \min \left(C[k, 2], \sum_{m=k+1}^{10} A[m] \right) \right\} = \max(1, 4, 4, 6, 9, 11, 8, 3) = 11 \quad (\text{Selected } k = 7)$$

Table 7: DP Table $C[i][j]$ (Maximum Minimal Sums)

$j \setminus i$	1	2	3	4	5	6	7	8	9	10
1	3	4	8	9	14	23	25	31	36	39
2	–	1	4	4	6	9	11	14	14	16
3	–	–	1	1	4	6	6	8	11	11

Table 8: Partition Points $K[i][j]$

$j \setminus i$	1	2	3	4	5	6	7	8	9	10
1	0	0	0	0	0	0	0	0	0	0
2	–	1	2	2	3	4	5	5	5	6
3	–	–	2	2	3	5	5	6	7	7

- Maximum minimal sum achievable by partitioning $A[1 \dots 10]$ into $M = 3$ groups is $C[10, 3] = 11$.
- Backtracking through the DP table to determine optimal grouping $G[1 \dots M]$ by finding partition points k that leads to $C[10, 3] = 11$.

Partitioning $A[1 \dots 10]$ into optimal groups based on $k=5$, and $k=7$:

– Group 1: $A[1 \dots 5] = [3, 1, 4, 1, 5]$

– Group 2: $A[6 \dots 7] = [9, 2]$

– Group 3: $A[8 \dots 10] = [6, 5, 3]$

Grouping array $G[1 \dots 3]$: $G = [5, 2, 3]$

- Determine summation of the elements in the j -th group of array A .

- Group 1 ($A[1 \dots 5]$): $B[1] = 3 + 1 + 4 + 1 + 5 = 14$
- Group 2 ($A[6 \dots 7]$): $B[2] = 9 + 2 = 11$
- Group 3 ($A[8 \dots 10]$): $B[3] = 6 + 5 + 3 = 14$

Sum of elements in each group $B[1 \dots 3]$: $B = [14, 11, 14]$

- The minimal group sum:

$$\min_{1 \leq j \leq 3} B[j] = \min(14, 11, 14) = 11$$

The minimal group sum matches $C[10, 3] = 11$, confirming that the minimal group sum is maximized.

Optimal Grouping Sizes ($G[1..M]$):

Group 1: 5 elements

Group 2: 2 elements

Group 3: 3 elements

Group Details:

Group 1: $A[1 \dots 5] = [3, 1, 4, 1, 5]$, Sum = 14

Group 2: $A[6 \dots 7] = [9, 2]$, Sum = 11

Group 3: $A[8 \dots 10] = [6, 5, 3]$, Sum = 14

Optimal Grouping G : 5 2 3

—

4. Implementation

```

1 #include <iostream>
2 #include <vector>
3 #include <algorithm>
4 #include <climits>
5 #include <numeric>
6 using namespace std;
7
8 vector<int> max_min_grouping(const vector<int>& A, int N, int M) {
9     // Initialize DP table C[i][j] which stores the maximum
10    // minimal sum that can be achieved by partitioning the i elements into j groups
11    vector<vector<int>>> C(N + 1, vector<int>(M + 1, INT_MIN));
12    // K table to store the partition points for reconstructing the optimal groups
13    vector<vector<int>>> K(N + 1, vector<int>(M + 1, -1));
14    // Base case: For 0 elements and 0 groups, the value max infinity as maximizing the minimum values
15    C[0][0] = INT_MAX;
16    // Sum array P to quickly calculate sums over any subarray of A
17    vector<int> P(N + 1, 0);
18    for (int i = 1; i <= N; ++i) {
19        P[i] = P[i - 1] + A[i - 1];
20    }
21    // DP base case: Only 1 group (j=1)
22    for (int i = 1; i <= N; ++i) {
23        C[i][1] = P[i];
24        K[i][1] = 0;
25    }
26
27    // DP table populate using the recurrence relation for each possible number of groups j (from 2 to M) and
28    // elements i (from j to N)
29    // Try multiple partition points from k until (j-1) to partition the first i elements into j groups by
30    // considering different
31    for (int j = 2; j <= M; ++j) {
32        for (int i = j; i <= N; ++i) {
33            for (int k = j - 1; k < i; ++k) {
34                // Sum of elements from A[k+1] to A[i]
35                int sum = P[i] - P[k];

```

```

34         // Minimum value from partitioning the first k elements into (j-1) groups and including the
new sum
35         int min_sum = min(C[k][j - 1], sum);
36         // Update the DP table if minimum value is larger minimum sum
37         if (min_sum > C[i][j]) {
38             C[i][j] = min_sum;
39             K[i][j] = k; // Store the partition point for backtracking
40         }
41     }
42 }
43 }
44
45 // Backtrack to reconstruct the optimal grouping G[1..M] by using the partition points stored in K
46 vector<int> G(M, 0);
47 int idx = N;
48 for (int j = M; j >= 1; --j) {
49     int k = K[idx][j]; // Get the partition point for group
50     G[j - 1] = idx - k; // Store the size of the j-th group
51     idx = k; // Move to the previous partition point
52 }
53
54 // Return the grouping sizes G[1..M], where each G[j] represents the number of elements
55 return G;
56 }
57
58 int main() {
59     vector<int> A = {3,1,4,1,5,9,2,6,5,3};
60     int N = A.size();
61     int M = 3;
62
63     vector<int> G = max_min_grouping(A, N, M);
64
65     cout << "Optimal Grouping G: ";
66     for (int g : G) {
67         cout << g << " ";
68     }
69     cout << endl;
70
71     return 0;
72 }

```