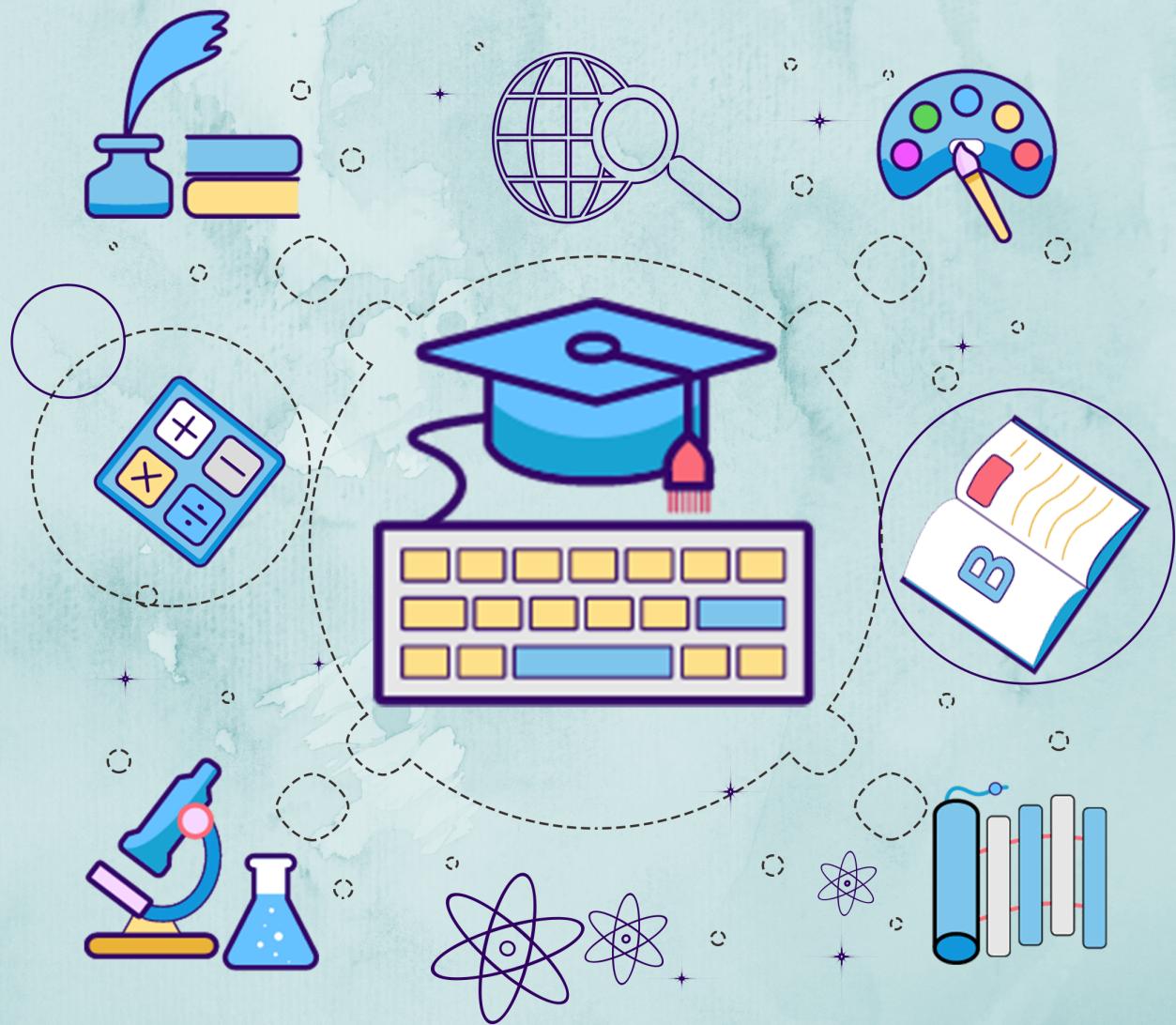


Kerala Notes



SYLLABUS | STUDY MATERIALS | TEXTBOOK

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KTU STUDY MATERIALS

LOGIC SYSTEM DESIGN

CST 203

Module 2

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SOLVED QUESTION PAPER

Module : 2

- Logic Gate
- Boolean Algebra
- Kmap
- Tabulation method

Logic Gate

1. Basic Gate : (a) NOT

(b) AND

(c) OR

2. Universal Gate: (a) NAND

(b) NOR

3. Arithmetic Gate : (a) XOR

(b) XNOR

$$A + A = A$$

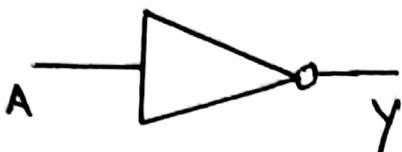
A	B	XOR
0	0	0
1	0	1
0	1	1
1	1	0



I. BASIC Gate

(a) NOT Gate (Inverted)

Symbol



Truth table

input A output Y

0	1
1	0

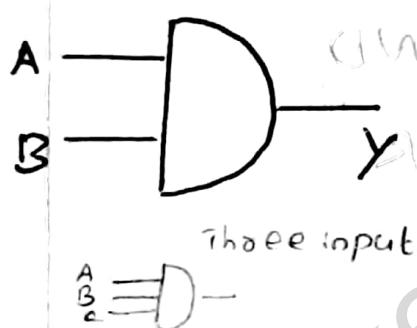
Boolean

Expression

$$Y = \bar{A}$$

(b) AND Gate (multiplexer)

Symbol



Truth table

A	B	Y
0	0	0
0	1	0
1	0	0
1	1	1

Boolean expression

$$Y = A \cdot B$$

(c) OR Gate

Symbol



Truth table

A	B	Y
0	0	0
0	1	1
1	0	1
1	1	1

Boolean Expression

$$Y = A + B$$

2. Universal Gate

Universal because we can make any other gate by using NAND or NOR.

(a) NAND Gate

Symbol (ANAN+NOT)



Truth table

A	B	Y
0	0	1
0	1	1
1	0	1
1	1	0

Boolean expression

$$Y = \overline{AB}$$

(b) NOR Gate

Symbol (OR+NOT)



Truth table

A	B	Y
0	0	1
0	1	0
1	0	0
1	1	0

Boolean expression

$$Y = \overline{A+B}$$

3. Arithmetic Gate

of base 2 (000 to 111)

addition with increment/decrement

2) parallel addition using register

file controller

(a) XOR (Exclusive OR)

Symbol

Truth Table

A	B	Y
0	0	0
0	1	1
1	0	1
1	1	0

Boolean Expression

$$Y = A \oplus B$$

$$Y = \bar{A}B + A\bar{B}$$

Output is high (1) if inputs are different

$$\text{Eg: } A=0 / B=0$$

$$1 \cdot 0 + 0 \cdot 1 = 0$$

$$A=0 / B=1$$

$$1 \cdot 1 + 0 \cdot 0 = 1 + 0 = 1$$

(b) XNOR (Exclusive XNOR)

Symbol

XOR + NOT

Truth Table

A	B	Y
0	0	1
0	1	0
1	1	0
1	0	1

Boolean Expression

$$Y = A \cdot \bar{B}$$

$$\bar{A}B + A\bar{B}$$

output is high (1) if inputs are same

Boolean Algebra

It is a set of rules used to simplify/minimize the Boolean expression without changing its functionality

1. Complement Rule : $\bar{\bar{A}} = A$

Eg: $A=0, \bar{A}=1, \bar{\bar{A}}=0$

2. AND Rule :

$$A \cdot A = A$$

$$A \cdot 0 = 0$$

$$A \cdot 1 = A$$

$$A \cdot \bar{A} = 0$$

3. OR Rule :

$$A + A = A$$

$$A + 0 = A$$

$$* A + 1 = 1$$

$$A + \bar{A} = 1$$

$$\text{eg: } \underbrace{\bar{A}Bc}_{=0} + \underbrace{A\bar{B}c}_{=0} +$$

$$= 1$$

$$\underbrace{AB+AC+1}_{=1}$$

4. Distributive Law:

$$A \cdot (B+C) = AB + AC$$

$$A + (B \cdot C) = (A+B)(A+C)$$

$$\text{Eg: } A + \bar{A}B \rightarrow (A + \bar{A}) \cdot (A + B)$$

$$(A + \bar{A}) \cdot (A + B) = 1 \cdot (A + B) = A + B$$

$$\text{Eg: } B + \bar{B}C \rightarrow (B + \bar{B}) \cdot (B + C)$$

$$= 1 \cdot (B + C) = \underline{\underline{B + C}}$$

5. Commutative law:

$$A + B = A \cdot B$$

$$A \cdot B = B \cdot A$$

6. Associative law:

$$(A \cdot B) \cdot C = A \cdot (B \cdot C)$$

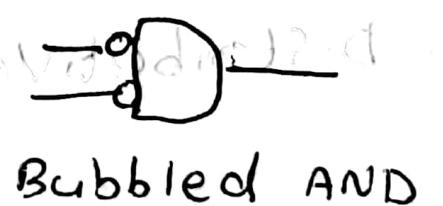
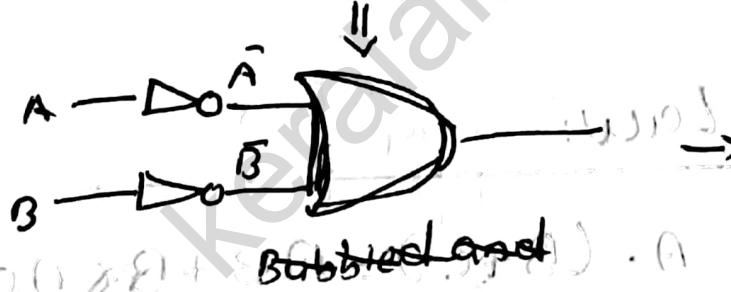
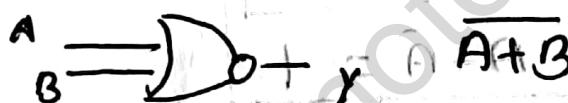
7. DeMorgan's Law:

$$\overline{A+B} = \bar{A} \cdot \bar{B}$$

$$\overline{AB} = \bar{A} + \bar{B}$$

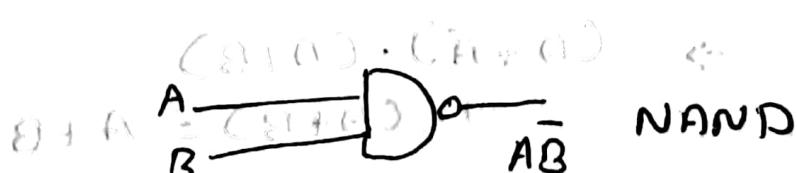
$$1. \overline{A+B} = \bar{A} \cdot \bar{B}$$

\downarrow
NOR gate



Bubbled AND

$$2. \overline{AB} = \bar{A} + \bar{B}$$



$$(S+D) \cdot (S+D) \leftarrow S(S+D) + D(S+D)$$



bubbled OR

NOR
~~AND~~
= Bubbled AND

AND
~~OR~~
= Bubbled OR

Questions:

1. $A \cdot A = A$

2. $A + (B \cdot C) = (A + B) (A + C) = A + [B + C]$

3. $A + 1 = 1$

4. $A + \bar{A} = 1$

Sum of Product (SOP)

$$Q = A \cdot B \cdot \bar{C} + A \cdot \bar{B} \cdot C + \bar{A} \cdot B \cdot C$$

Product of sum (POS)

$$Q = (A + B) \cdot (\bar{B} + C) \cdot (A + \bar{C})$$

$$SOP = \underbrace{A \cdot B \cdot \bar{C}}_{0\sigma} + \underbrace{A \bar{B} C}_{0\sigma} + \underbrace{\bar{A} B C}_{0\sigma}$$

$$(POS) \underbrace{(A + B + C)}_{(and)} \cdot \underbrace{(A + B + \bar{C})}_{(and)} \cdot \underbrace{(A + \bar{B} + \bar{C})}_{(and)}$$

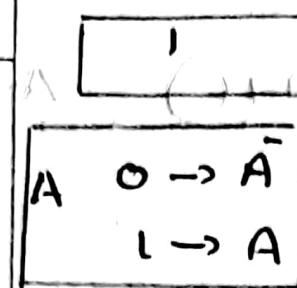
and

(and and
C, and)

Sum of product (SOP)

A	B	C	O/P
0	0	0	0
0	0	1	0
0	1	0	0
1	1	1	1
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	1

considered parameters



min + term (cm)

$$= (\bar{A} \cdot B \cdot C) + (A \cdot \bar{B} \cdot \bar{C}) +$$

$$(A \cdot B \cdot \bar{C}) + (\bar{A} \cdot \bar{B} \cdot \bar{C})$$

$$= m_3 + m_4$$

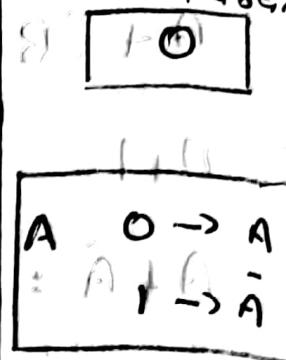
$$\Sigma_m (3, 4, 6, 7)$$

(sum of min team)

Product of sum (POS)

A	B	C	O/P
0	0	0	0
0	0	1	1
0	1	0	1
1	1	0	0
1	0	1	0
1	0	0	1
1	1	1	1

considered parameters



Max + term (cm)

$$= (A + B + C) \cdot (A + \bar{B} + \bar{C})$$

$$(A + \bar{B} + C) \cdot (A + \bar{B} + \bar{C})$$

$$+ 3 \cdot 8 \cdot A = 96$$

$$= M_0 \cdot M_3 \cdot M_6 \cdot M_7$$

$$= \Sigma M (0, 3, 6, 7)$$

$$\prod M (0, 3, 6, 7)$$

(product of max team)

A	B	C	D	Min Term SOP (m _i)	SOP $0 \rightarrow \bar{A}$ $I \rightarrow A$	CMS
0	0	0	0	$\bar{A} \bar{B} \bar{C} \bar{D}$ (m ₀)	$A + B + C + D$ (M ₀)	
1	0	0	0	$\bar{A} \bar{B} \bar{C} D$ (m ₁)	$A + B + C + \bar{D}$ (M ₁)	
2	0	0	1	$\bar{A} \bar{B} C \bar{D}$ (m ₂)	$A + B + \bar{C} + D$ (M ₂)	
3	0	0	1	$\bar{A} \bar{B} C D$ (m ₃)	$A + B + \bar{C} + \bar{D}$ (M ₃)	
4	0	1	0	$\bar{A} B \bar{C} \bar{D}$ (m ₄)	$A + \bar{B} + C + D$ (M ₄)	
5	0	1	0	$\bar{A} B \bar{C} D$ (m ₅)	$A + \bar{B} + C + \bar{D}$ (M ₅)	
6	0	1	1	$\bar{A} \cdot B C \bar{D}$ (m ₆)	$A + \bar{B} + \bar{C} + D$ (M ₆)	
7	0	1	1	$\bar{A} B C D$ (m ₇)	$A + \bar{B} + \bar{C} + \bar{D}$ (M ₇)	
8	1	0	0	$A \bar{B} \bar{C} \bar{D}$ (m ₈)	$\bar{A} + B + C + D$ (M ₈)	
9	1	0	0	$A \bar{B} \bar{C} D$ (m ₉)	$\bar{A} + B + C + \bar{D}$ (M ₉)	
10	1	0	1	$A \bar{B} C \bar{D}$ (m ₁₀)	$\bar{A} + B + \bar{C} + D$ (M ₁₀)	
11	1	0	1	$A \bar{B} C D$ (m ₁₁)	$\bar{A} + B + \bar{C} + \bar{D}$ (M ₁₁)	
12	1	1	0	$A B \bar{C} \bar{D}$ (m ₁₂)	$\bar{A} + \bar{B} + C + D$ (M ₁₂)	
13	1	1	0	$A B \bar{C} D$ (m ₁₃)	$\bar{A} + \bar{B} + C + \bar{D}$ (M ₁₃)	
14	1	1	1	$A B C \bar{D}$ (m ₁₄)	$A + \bar{B} + \bar{C} + \bar{D}$ (M ₁₄)	
15	1	1	1	$A B C D$ (m ₁₅)	$\bar{A} + \bar{B} + \bar{C} + \bar{D}$ (M ₁₅)	

Problem:

A, B, C, D, E find m_{18} , M_{18}

$$\begin{array}{r} 2 \mid 18(0) \\ 2 \mid 9(1) \\ 2 \mid 4(0) \\ 2 \mid 2(0) \\ \quad . \end{array}$$

$$18 \rightarrow 10010$$

$$m_{18} = A \bar{B} \bar{C} \bar{D} \bar{E}$$

$$m_{18} = \bar{A} + \bar{B} + \bar{C} + \bar{D} + \bar{E}$$

[Gated then
 $\gamma = 4$ bit]

Here largest no
 is 9 it is a
 Four bit no

$$m_0 = 0000 \quad \bar{A} \bar{B} \bar{C} \bar{D}$$

$$m_3 = 0011 \quad \bar{A} \bar{B} \bar{C} D$$

$$m_5 = 0101 \quad \bar{A} \bar{B} \bar{C} D$$

$$m_7 = 0111 \quad \bar{A} \bar{B} \bar{C} D$$

$$m_9 = 1001 \quad A \bar{B} \bar{C} D$$

$$= (\bar{A} \bar{B} \bar{C} \bar{D}) + (\bar{A} \bar{B} C D) +$$

$$(\bar{A} B \bar{C} \bar{D}) + (\bar{A} B C D) +$$

$$(A \bar{B} \bar{C} D)$$

0	0	0	0	1	0
0	0	0	1	1	1
0	0	1	0	1	2
0	0	1	1	1	3
0	1	0	0	1	4
0	1	0	1	1	5
0	1	1	0	1	6
0	1	1	1	1	7
1	0	0	0	1	8
1	0	0	1	1	9
1	0	1	0	1	10
1	0	1	1	1	11
1	1	0	0	1	12
1	1	0	1	1	13
1	1	1	0	1	14
1	1	1	1	1	15

M_0, M_5, M_7

Here log KeralaNotes
it is a 3 bit NO

000 101 111

$$M_0 = 000 = A + B + C$$

$$M_5 = 101 = \bar{A} + B + \bar{C}$$

$$M_7 = 111 = \bar{A} + \bar{B} + \bar{C}$$

$$\Rightarrow (A + B + C) \cdot (\bar{A} + B + \bar{C}) \cdot (\bar{A} + \bar{B} + \bar{C})$$

Questions

1. Min term of 20

$$\begin{array}{r}
 2 | 20(0) \\
 2 | 10(0) \\
 2 | 5(1) \\
 2 | 2(0) \\
 \hline
 1
 \end{array}
 32A + 16A + 8A = 902$$

$$(A + B + C)(\bar{A} + B + \bar{C})(\bar{A} + \bar{B} + C) = 20$$

20 \rightarrow 10100

Min term of 20 = $A \cdot \bar{B} \cdot C \cdot \bar{D} \cdot E$

2. Max term of 56

$$\begin{array}{r}
 2 | 56(0) \\
 2 | 28(0) \\
 2 | 14(0) \\
 2 | 7(1) \\
 2 | 3(1) \\
 \hline
 1
 \end{array}$$

$$\begin{aligned}
 &= 111000 \\
 &= \cancel{ABCDEF} \\
 &= \bar{A} \cdot \bar{B} \cdot \bar{C} \cdot \bar{D} \cdot \bar{E} \cdot \bar{F}
 \end{aligned}$$

3.

Summation of minterms ($2, 3, 5, 7$)

$$m_2 = 0010 = \bar{A}B\bar{C} \quad \text{m } 000$$

$$m_3 = 0011 = \bar{A}BC$$

$$m_5 = 0101 = A\bar{B}\bar{C}$$

$$m_7, 0111 = A\bar{B}C$$

$$= (\bar{A}\bar{B}\bar{C}) + (\bar{A}\bar{B}C) + (A\bar{B}\bar{C}) + (A\bar{B}C)$$

Standard / canonical SOP

standard
↑

$$SOP = \underbrace{ABC}_{OR} + \underbrace{\bar{A}BC} + \underbrace{AB\bar{C}} + \underbrace{A\bar{B}\bar{C}}$$

$$SPOS = \underbrace{(A+B+C)}_{(A+B+\bar{C})} \cdot \underbrace{(A+\bar{B}+\bar{C})}_{(A+\bar{B}+C)} \cdot \underbrace{(\bar{A}+B+C)}_{(A+\bar{B}+\bar{C})}$$

non standard
Eg: for SOP:

- $\underbrace{A}_{\text{missing}} + \underbrace{Bc}_{\text{missing}} + \underbrace{\bar{A}\bar{B}\bar{C}}_{\text{missing}}$
- $\underbrace{\bar{B}\bar{C}}_{\text{missing}} + \underbrace{A}_{\text{missing}} \rightarrow \text{missing}$
- $\underbrace{\bar{B}c}_{\text{missing}} + \underbrace{\bar{A}}_{\text{missing}} \rightarrow \text{missing}$
- $\underbrace{Bc}_{\text{missing}} + \underbrace{B\bar{C}}_{\text{missing}}$

- $(A+B+C) \cdot \underbrace{(\bar{A}+\bar{B})}_{\substack{\text{C} \\ \bar{C}}} \cdot \underbrace{(\bar{B}+\bar{C})}_{\substack{\text{A} \\ \bar{A}}} \rightarrow \text{missing}$
- $\underbrace{(A)}_{\substack{\text{B} \\ \bar{C}}} + \underbrace{(\bar{A}+\bar{B}+\bar{C})}_{\substack{\text{A} \\ \bar{B} \\ \bar{C}}} \rightarrow \text{missing}$

$\rightarrow \bar{B}\bar{C} + (\bar{A}+\bar{B}+\bar{C})(\bar{B}\bar{A} + \bar{B}\bar{A})$

$$= \bar{B}\bar{A} + \bar{B}\bar{A} + \bar{B}\bar{A} + \bar{B}\bar{A} + \bar{B}\bar{A} + \bar{B}\bar{A}$$

Step 1: Write down all the terms

Step 2: If one or more variable are missing, expand value by multiplying with sum of each one of the missing variable and its complement

Step 3: Drop out Redundant terms

Examples :

$$\begin{array}{ll} A & \bar{B} \\ \bar{A} & \bar{B} \\ \bar{B} & B \\ B & \bar{B} \end{array}$$

$$(\bar{A}+A)\bar{B} + (\bar{B}+B)(\bar{A}+\bar{B})A$$

$$+ (\bar{B}+\bar{B}+\bar{B}+\bar{B})A \quad \cancel{\bar{B}A} + \cancel{B}A$$

Q convert the sop into standard form

$$1. \bar{A}BC + \bar{A}\bar{B}C + \bar{B}C + \bar{A}\bar{B}\bar{C}$$

Step 1:

$$\bar{A}BC + \bar{A}\bar{B} + \bar{B}C + \bar{A}\bar{B}\bar{C}$$

\bar{C} $A\bar{A}$

$$2. \Rightarrow \bar{A}BC + \bar{A}\bar{B}(C + \bar{C}) + \bar{B}C(A + \bar{A}) + \bar{A}\bar{B}C$$

$$\Rightarrow \underline{\bar{A}BC} + \underline{\bar{A}\bar{B}C} + \underline{\bar{A}\bar{B}\bar{C}} + A\bar{B}C + \cancel{\bar{A}\bar{B}C} + \cancel{\bar{A}\bar{B}\bar{C}}$$

3. check repeat

$$\Rightarrow \underline{\bar{A}BC} + \underline{\bar{A}\bar{B}C} + \underline{\bar{A}\bar{B}\bar{C}} + A\bar{B}C$$

$$2. Y = A + \bar{B}C \text{ for } A, B, C$$

BC

$$\frac{A}{\bar{B}\bar{C}} + \frac{1}{A}$$

$\bar{B}C$ A

$\bar{B}C$ \bar{A}

$B\bar{C}$

BC

$$A(B + \bar{B})(C + \bar{C}) + \bar{B}C(A + \bar{A})$$

$$\cancel{AB} + \cancel{A\bar{B}} \quad A(BC + BC + \bar{B}C + \bar{B}\bar{C}) +$$

$$\bar{B}CA + \bar{B}C\bar{A}$$

$A \bar{B} C$

$$\Rightarrow A B C + A B \bar{C} + A \bar{B} C + A \bar{B} \bar{C} +$$

minimise $\bar{A} B C$
 $\bar{A} B \bar{C}$

$$A \bar{B} C + \bar{A} C \bar{B}$$

$$\Rightarrow A \bar{B} \bar{C} + A B \bar{C} + A \bar{B} C + A \bar{B} \bar{C} + \bar{A} C \bar{B}$$

3. convert the equation into standard pos

$$(A+B+C) (\bar{A}+\bar{C}) (\bar{A}+\bar{B})$$

1	0	0
1	1	0
1	0	1

$$= (A+B+C) (\bar{A}+\bar{C}+B) (\bar{A}+C+\bar{B})$$

$$(\bar{A}+\bar{B}+C) (\bar{A}+\bar{B}+\bar{C})$$

arrange =

$$= (A+B+C) (\bar{A}+\bar{B}+C) (\bar{A}+\bar{B}+\bar{C})$$
 ~~$(\bar{A}+\bar{B}+C) \cdot (\bar{A}+\bar{B}+\bar{C})$~~

$$\rightarrow (A+B+\bar{C}) (\bar{A}+B+C) (\bar{A}+\bar{B}+C)$$

$$\bar{B} \bar{A} = X \quad (\bar{A}+\bar{B}+\bar{C})$$

$$[100] x = 8 \quad x = A$$

$$\bar{X} \bar{X} = X$$

Content

NAND \rightarrow • NOT
 • AND
 • OR
 • NOR

NOR \rightarrow • NOT
 • AND

• OR
 • NAND

\Rightarrow NAND (NOT + AND)

=Do-

Boolean expression

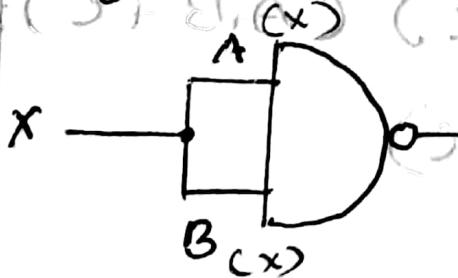
$$Y = \overline{AB}$$

Truth table		
A	B	y
0	0	1
0	1	1
1	0	1
1	1	0

(i) NOT Function

NAND \rightarrow NOT

An inverter can be made from NAND gate by connecting all the input together and creating, in effect a single common input



Proof

$$Y = \overline{AB}$$

Putting $A=x$ $B=x$ [equal]

$$Y = \overline{xx}$$

By using DeMorgan's Law [Page No. 112]

$$Y = \bar{X} + \bar{\bar{X}} \quad [\text{Reference: } A + A = A]$$

$$Y = \bar{\bar{X}}$$

It means it is without any output complement of

So convert to digital logic with input X and output Y.

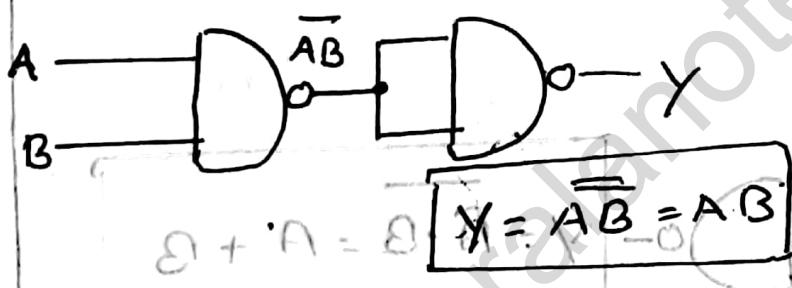
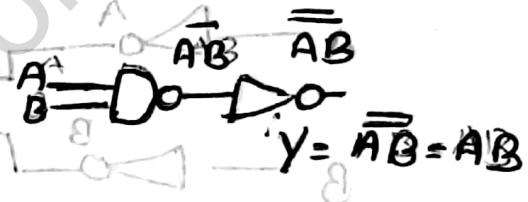
(ii) **AND Function** [Relation: Non-complement $\rightarrow A + A \text{ AND }$]

NAND \rightarrow AND

[$A = \bar{A}$] [With respect to]

An AND function is generated by simply inverting the output of NAND gate

NAND + NOT \rightarrow AND



		AND	
A	B	AB	AB-bar
0	0	0	1
0	1	0	1
1	0	0	1
1	1	1	0

A	B	AB	AB-bar
0	0	0	1
0	1	0	1
1	0	0	1
1	1	1	0

$$\overline{\overline{AB}} = AB$$

we know how to connect NAND using here. $= \overline{D_o - D_o} = D_o - \overline{D_o}$

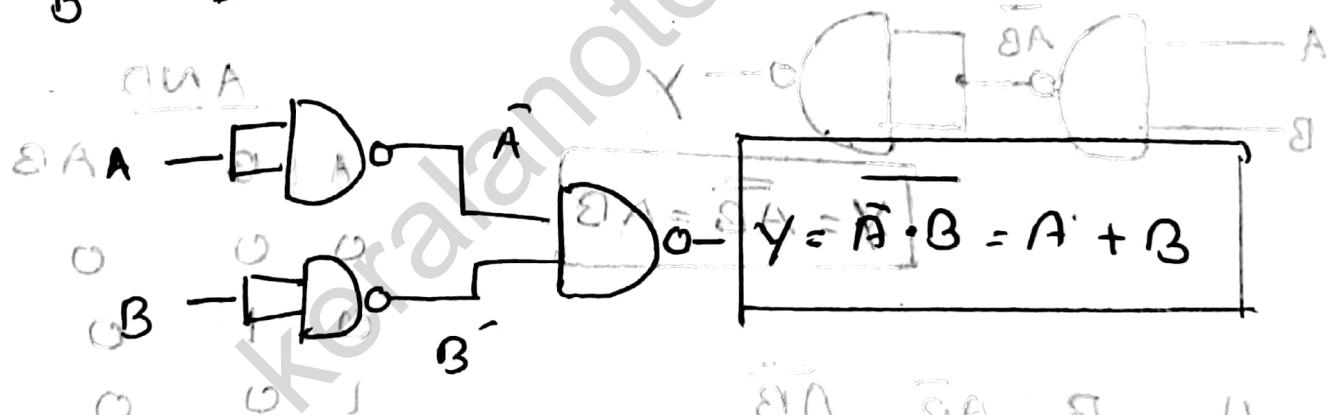
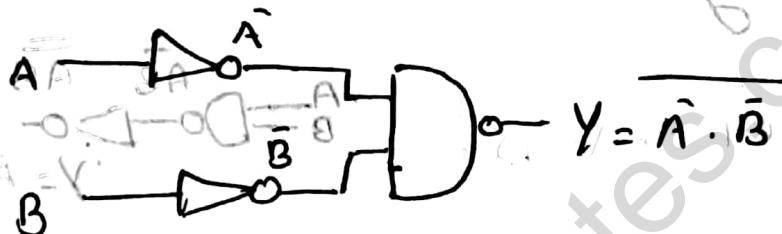
(iii) OR Function

~~NAND \rightarrow OR~~ $X + \bar{X} = Y$

An OR function is generated by simply inverting both the inputs of NAND gate.

$$Y = A + B$$

$$= \bar{\bar{A}} + \bar{\bar{B}} \quad [\text{DeMorgan's law}] \quad [\bar{\bar{A}} = A]$$



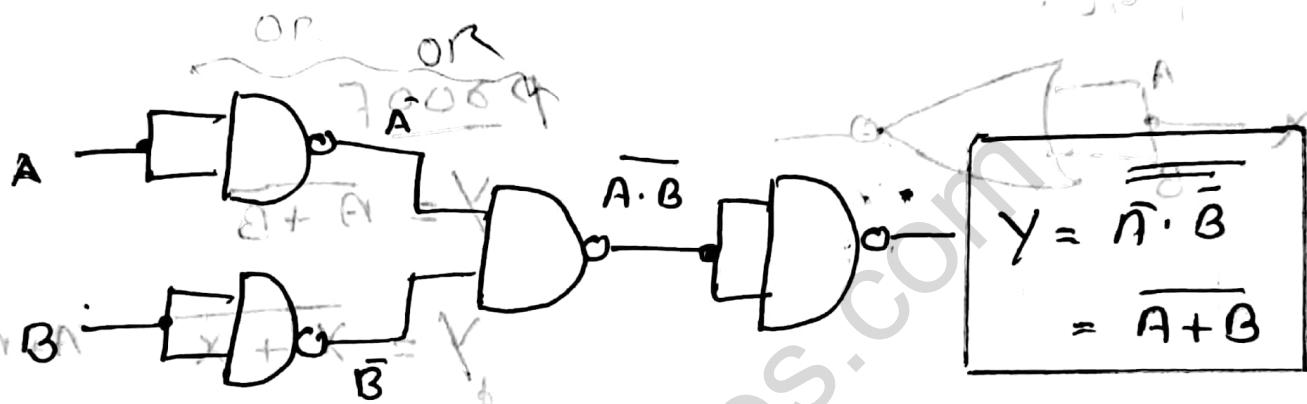
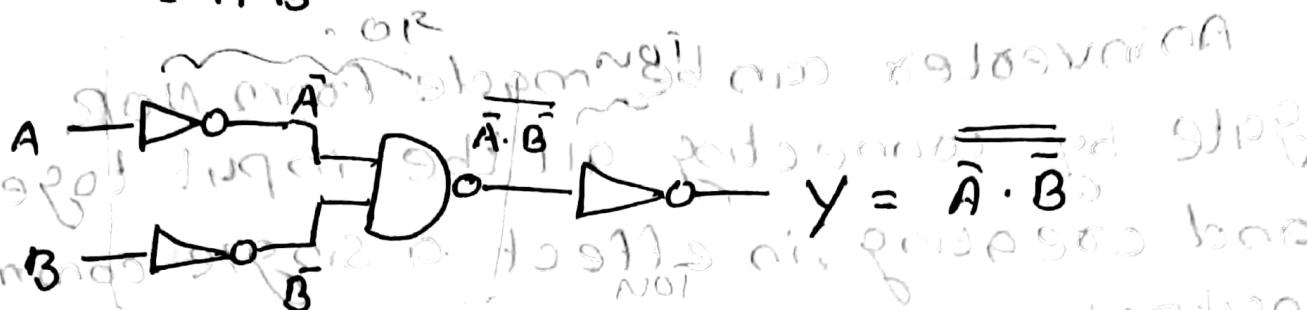
(iv) NOR Function [OR+NOT]

If we connect a NOT at OR output we get NOR Function

A NOR function is generated by simply inverting both the inputs of NAND gate and invert the output.

$$Y = \overline{A+B}$$

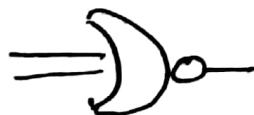
$$= \bar{A} \cdot \bar{B} \quad [\text{Demorgan's law}] \quad [\bar{\bar{A}} = A]$$



\Rightarrow NOR Gate ($\text{NOT} + \text{OR}$)

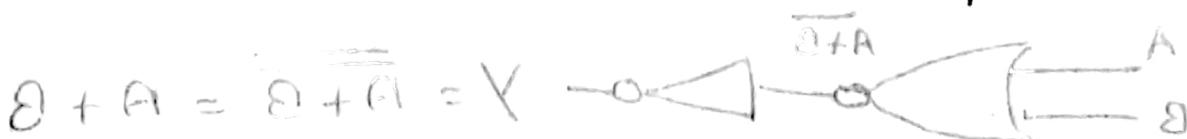
$\rightarrow \text{NOT}$

Boolean Expression

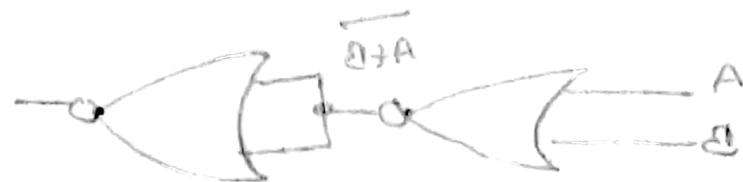


$$Y = \overline{A+B}$$

A	B	Y
0	0	1
1	0	0
1	1	0



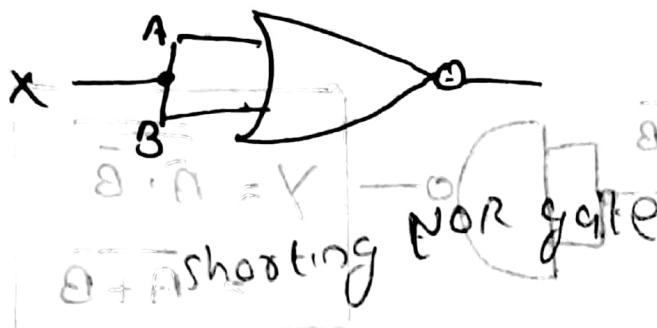
$$B+A = \overline{\overline{B}+\overline{A}} = X$$



(i) NOT function

$\text{NOR} \rightarrow \text{NOT}$

An inverter can be made from NOR gate by connecting all the input together and creating, in effect a single common output.



Proof

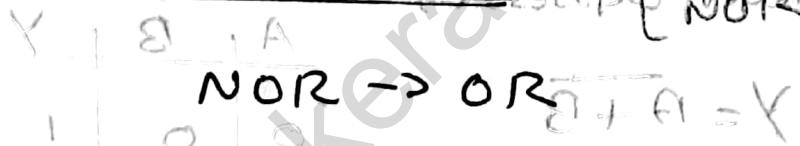
$$Y = \overline{A+B}$$

$$Y = \overline{\overline{X}+X}$$

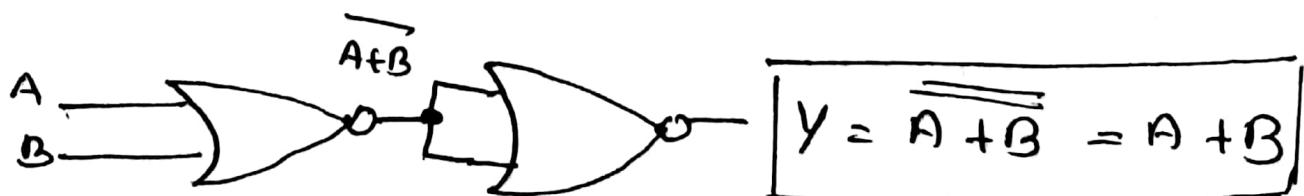
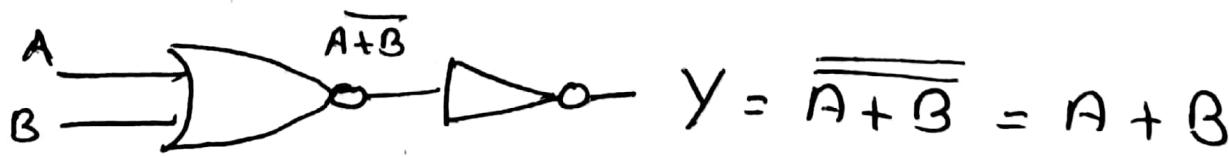
$$Y = \overline{\overline{X}}$$

(ii) OR function

[NOR complement $\rightarrow \text{OR}$]



An OR function is generated by simply inverting the output of NOR gate.



A	B	$A \cdot B$
0	0	0
0	1	0
1	0	0
1	1	1

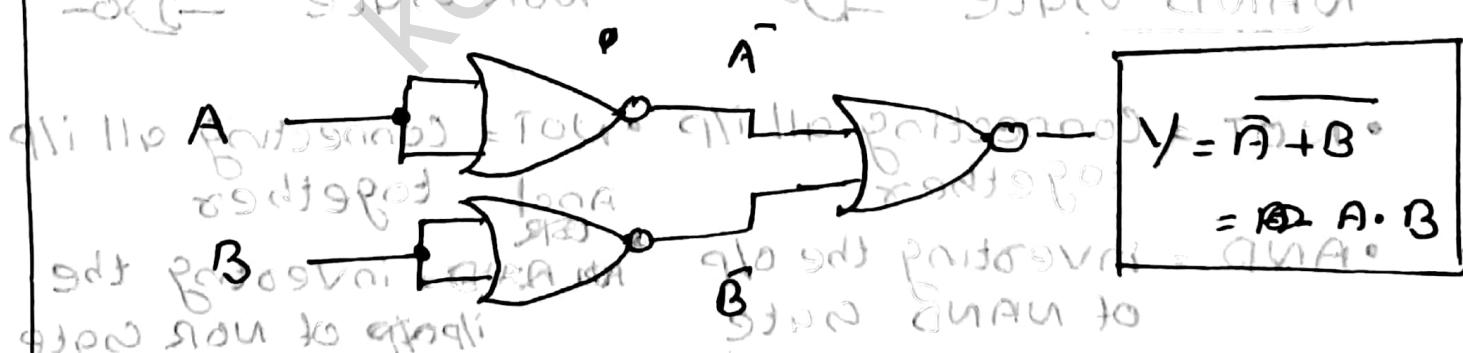
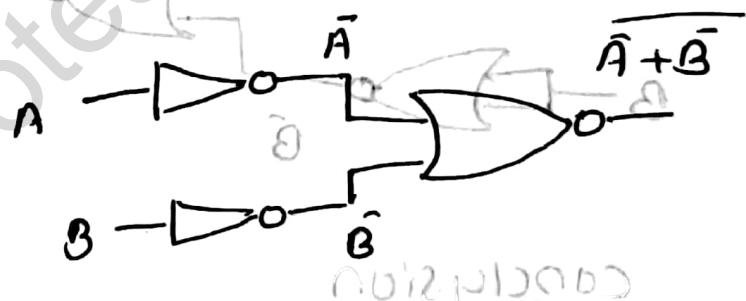
A	B	$\bar{A} + \bar{B}$	$\bar{A} + B$
0	0	1	1
0	1	1	0
1	0	1	0
1	1	0	1

AND function

$\text{NOR} \rightarrow \text{AND}$

An AND function is generated by simply inverting both input of NOR gate

$$Y = \overline{\overline{A} \cdot \overline{B}} = \overline{\overline{A} + \overline{B}} = \overline{(\overline{A} \cdot \overline{B})} = \overline{\overline{A} \cdot \overline{B}} = \overline{A} + \overline{B}$$



$$Y = A + B$$

$$Y = A \cdot B$$

For More Study Materials : <https://www.keralanotes.com/>

(iv)

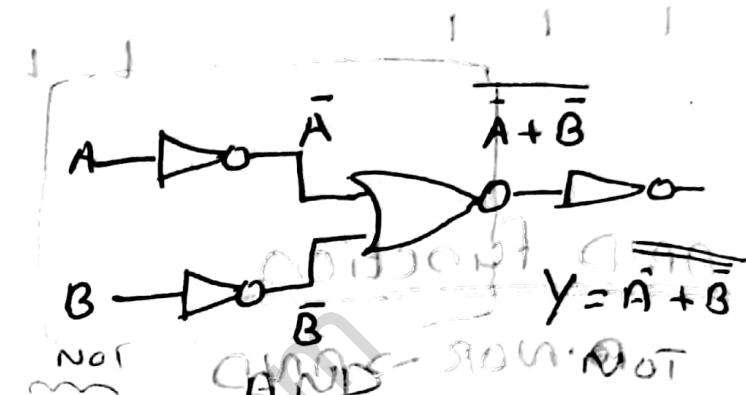
NAND Function

An NAND function is generated by simply inverting both the inputs of NOR gate and invert the output.

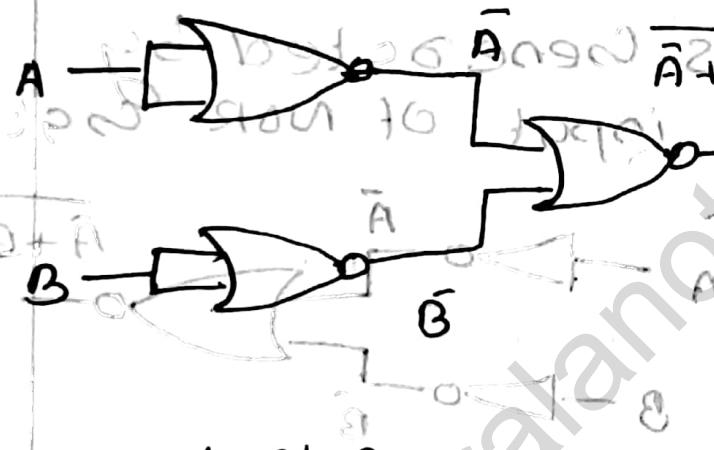
$$Y = \overline{A \cdot B}$$

$$= \overline{\overline{A} + \overline{B}} \quad [\overline{\overline{A}} = A]$$

$$= \overline{\overline{A} + \overline{B}}$$



AND



Conclusion

NAND gate \Rightarrow $\overline{A \cdot B}$

NOR gate \Rightarrow $\overline{A + B}$

- NOT = connecting all i/p together
- AND = inverting the o/p of NAND gate
- OR = " i/p
- NOR = inverting the i/p of NAND gate and invert the output
- NOT = connecting all i/p together
- OR = " o/p
- NAND = inverting all i/p of NOR gate and invert o/p

Karnaugh Map

Given $SOP = \bar{B}\bar{A} + \bar{B}A + BA + B\bar{A}$
We have

Karnaugh map is introduced to simplification of Boolean function in a easy way.

2 variable : $\rightarrow SOP$
 $\rightarrow POS$

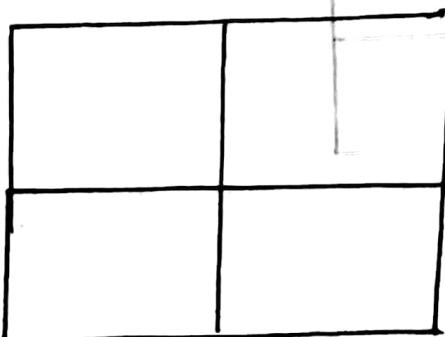
3 Variable : $\rightarrow SOP$
 $\rightarrow POS$

4 Variable : $\rightarrow SOP$
 $\rightarrow POS$

5 Variable : $\rightarrow SOP$
 $\rightarrow POS$

2 variable KMAP

$$\bar{A}\bar{B} + A\bar{B} + AB + \bar{A}B$$



$$2^2 = 4$$

$$\boxed{2^n} = 2^2 = 4$$

3 Variable K map

$$\bar{A}BC + AB\bar{C} + ABC + \bar{A}\bar{B}\bar{C}$$

of 16 cells of 4x4 grid

$$2^3 = 8 \quad \begin{matrix} 100 \\ 101 \\ 110 \\ 111 \end{matrix} : \text{grid size}$$

4 Variable K map

$$\bar{A}BCD + A\bar{B}CD + AB\bar{C}D + AB\bar{C}\bar{D}$$

of 16 cells of 4x4 grid

$2^4 = 16 \text{ cells}$

$$N = S_0 =$$

5 variable k map

		CDE 0 0							
		11 01 10 00							
		00	001	011	010	110	111	101	100
A	B	0 1 0	0 1 1	0 1 0	1 1 0	1 1 1	1 0 1	1 0 0	0 0 0
D	E	0 1 0	0 1 1	0 1 0	1 1 0	1 1 1	1 0 1	1 0 0	0 0 0
C		0 0 1	0 1 0	1 0 0	1 1 0	1 0 0	0 1 0	0 0 0	0 0 0
F		0 1 0	0 1 1	0 1 0	1 0 0	1 0 1	1 0 0	0 1 0	0 1 0
G		0 1 1	1 0 0	1 0 1	1 1 0	1 1 1	1 0 1	1 0 0	1 1 0
H		1 0 0	1 0 1	1 1 0	1 1 1	1 0 0	1 0 1	1 0 0	1 1 0
I		1 0 1	1 1 0	1 1 1	1 1 0	1 0 0	1 0 1	1 0 0	1 1 0
J		1 1 0	1 1 1	1 1 0	1 1 1	1 0 0	1 0 1	1 0 0	1 1 0
K		0 0 0	0 0 1	0 1 0	0 1 1	1 0 0	1 0 1	0 0 0	0 0 0
L		0 0 1	0 1 0	0 1 1	1 0 0	1 0 1	1 0 0	0 0 0	0 0 0
M		0 1 0	0 1 1	1 0 0	1 0 1	1 0 0	1 0 1	0 0 0	0 0 0
N		0 1 1	1 0 0	1 0 1	1 1 0	1 1 1	1 0 1	1 0 0	1 1 0
O		1 0 0	1 0 1	1 1 0	1 1 1	1 0 0	1 0 1	1 0 0	1 1 0
P		1 0 1	1 1 0	1 1 1	1 1 0	1 0 0	1 0 1	1 0 0	1 1 0
Q		1 1 0	1 1 1	1 1 0	1 1 1	1 0 0	1 0 1	1 0 0	1 1 0
R		0 0 0	0 0 1	0 1 0	0 1 1	1 0 0	1 0 1	0 0 0	0 0 0
S		0 0 1	0 1 0	0 1 1	1 0 0	1 0 1	1 0 0	0 0 0	0 0 0
T		0 1 0	0 1 1	1 0 0	1 0 1	1 0 0	1 0 1	0 0 0	0 0 0
U		0 1 1	1 0 0	1 0 1	1 1 0	1 1 1	1 0 1	1 0 0	1 1 0
V		1 0 0	1 0 1	1 1 0	1 1 1	1 0 0	1 0 1	1 0 0	1 1 0
W		1 0 1	1 1 0	1 1 1	1 1 0	1 0 0	1 0 1	1 0 0	1 1 0
X		1 1 0	1 1 1	1 1 0	1 1 1	1 0 0	1 0 1	1 0 0	1 1 0
Y		0 0 0	0 0 1	0 1 0	0 1 1	1 0 0	1 0 1	0 0 0	0 0 0
Z		0 0 1	0 1 0	0 1 1	1 0 0	1 0 1	1 0 0	0 0 0	0 0 0
A'	B'	0 0 0	0 0 1	0 1 0	0 1 1	1 0 0	1 0 1	0 0 0	0 0 0
C'	D'	0 0 0	0 0 1	0 1 0	0 1 1	1 0 0	1 0 1	0 0 0	0 0 0
E'	F'	0 0 0	0 0 1	0 1 0	0 1 1	1 0 0	1 0 1	0 0 0	0 0 0
G'	H'	0 0 0	0 0 1	0 1 0	0 1 1	1 0 0	1 0 1	0 0 0	0 0 0
I'	J'	0 0 0	0 0 1	0 1 0	0 1 1	1 0 0	1 0 1	0 0 0	0 0 0
K'	L'	0 0 0	0 0 1	0 1 0	0 1 1	1 0 0	1 0 1	0 0 0	0 0 0
M'	N'	0 0 0	0 0 1	0 1 0	0 1 1	1 0 0	1 0 1	0 0 0	0 0 0
O'	P'	0 0 0	0 0 1	0 1 0	0 1 1	1 0 0	1 0 1	0 0 0	0 0 0
Q'	R'	0 0 0	0 0 1	0 1 0	0 1 1	1 0 0	1 0 1	0 0 0	0 0 0
S'	T'	0 0 0	0 0 1	0 1 0	0 1 1	1 0 0	1 0 1	0 0 0	0 0 0
U'	V'	0 0 0	0 0 1	0 1 0	0 1 1	1 0 0	1 0 1	0 0 0	0 0 0
W'	X'	0 0 0	0 0 1	0 1 0	0 1 1	1 0 0	1 0 1	0 0 0	0 0 0
Y'	Z'	0 0 0	0 0 1	0 1 0	0 1 1	1 0 0	1 0 1	0 0 0	0 0 0

5 variable karnaugh map (Grey code)

$2^5 = 32$ cells

4 variable k map

		CD ̄D ̄D D				grey code			
		CD	̄D	̄D D	D	̄C D	C D	̄C ̄D	C ̄D
		00	01	11	10	10	01	00	00
A	B	0 0	0 1	1 1	1 0	1 0	0 1	0 0	0 0
D	E	0 0	0 1	1 1	1 0	1 0	0 1	0 0	0 0
C		0 1	1 0	1 1	0 0	0 1	0 0	1 1	1 1
F		1 0	1 1	0 0	0 1	0 0	1 1	1 0	1 0
G		1 1	0 0	0 1	1 0	0 0	1 1	0 0	0 1
H		0 0	1 1	0 0	1 0	1 1	0 0	0 1	1 1
I		0 1	0 0	1 1	0 1	0 0	0 1	1 0	0 0
J		1 0	0 1	0 0	1 1	0 1	1 0	0 0	1 0
K		0 0	0 1	0 0	0 1	0 0	0 1	0 0	0 1
L		0 1	0 0	0 0	0 1	0 0	0 1	0 0	0 1
M		1 0	0 1	0 0	0 1	0 0	0 1	0 0	0 1
N		0 0	1 1	0 0	1 0	0 0	1 1	0 0	0 1
O		0 1	0 0	1 1	0 1	0 0	0 1	1 0	0 0
P		1 0	0 1	0 0	1 1	0 1	0 0	0 1	1 0
Q		0 0	0 1	0 0	0 1	0 0	0 1	0 0	0 1
R		0 1	0 0	0 0	0 1	0 0	0 1	0 0	0 1
S		1 0	0 1	0 0	0 1	0 0	0 1	0 0	0 1
T		0 0	1 1	0 0	1 0	0 0	1 1	0 0	0 1
U		0 1	0 0	1 1	0 1	0 0	0 1	1 0	0 0
V		1 0	0 1	0 0	1 1	0 1	0 0	0 1	1 0
W		0 0	0 1	0 0	0 1	0 0	0 1	0 0	0 1
X		0 1	0 0	0 0	0 1	0 0	0 1	0 0	0 1
Y		1 0	0 1	0 0	0 1	0 0	0 1	0 0	0 1
Z		0 0	1 1	0 0	1 0	0 0	1 1	0 0	0 1
A'	B'	0 0	0 1	1 1	1 0	1 0	0 1	0 0	0 0
C'	D'	0 0	0 1	1 1	1 0	1 0	0 1	0 0	0 0
E'	F'	0 0	0 1	1 1	1 0	1 0	0 1	0 0	0 0
G'	H'	0 0	0 1	1 1	1 0	1 0	0 1	0 0	0 0
I'	J'	0 0	0 1	1 1	1 0	1 0	0 1	0 0	0 0
K'	L'	0 0	0 1	1 1	1 0	1 0	0 1	0 0	0 0
M'	N'	0 0	0 1	1 1	1 0	1 0	0 1	0 0	0 0
O'	P'	0 0	0 1	1 1	1 0	1 0	0 1	0 0	0 0
Q'	R'	0 0	0 1	1 1	1 0	1 0	0 1	0 0	0 0
S'	T'	0 0	0 1	1 1	1 0	1 0	0 1	0 0	0 0
U'	V'	0 0	0 1	1 1	1 0	1 0	0 1	0 0	0 0
W'	X'	0 0	0 1	1 1	1 0	1 0	0 1	0 0	0 0
Y'	Z'	0 0	0 1	1 1	1 0	1 0	0 1	0 0	0 0
A''	B''	0 0	0 1	1 1	1 0	1 0	0 1	0 0	0 0
C''	D''	0 0	0 1	1 1	1 0	1 0	0 1	0 0	0 0
E''	F''	0 0	0 1	1 1	1 0	1 0	0 1	0 0	0 0
G''	H''	0 0	0 1	1 1	1 0	1 0	0 1	0 0	0 0
I''	J''	0 0	0 1	1 1	1 0	1 0	0 1	0 0	0 0
K''	L''	0 0	0 1	1 1	1 0	1 0	0 1	0 0	0 0
M''	N''	0 0	0 1	1 1	1 0	1 0	0 1	0 0	0 0
O''	P''	0 0	0 1	1 1	1 0	1 0	0 1	0 0	0 0
Q''	R''	0 0	0 1	1 1	1 0	1 0	0 1	0 0	0 0
S''	T''	0 0	0 1	1 1	1 0	1 0	0 1	0 0	0 0
U''	V''	0 0	0 1	1 1	1 0	1 0	0 1	0 0	0 0
W''	X''	0 0	0 1	1 1	1 0	1 0	0 1	0 0	0 0
Y''	Z''	0 0	0 1	1 1	1 0	1 0	0 1	0 0	0 0

A, B, C, D

00 01 10 11

$\bar{A}\bar{B}$	m_0 0000 $\bar{A}\bar{B}\bar{C}\bar{D}$	m_1 0001 $\bar{A}\bar{B}\bar{C}D$	m_2 0010 $\bar{A}\bar{B}cD$	m_3 0011 $\bar{A}\bar{B}CD$	m_4 0100 $\bar{A}\bar{B}\bar{C}\bar{D}$	m_5 0101 $\bar{A}\bar{B}\bar{C}D$	m_6 0110 $\bar{A}\bar{B}cD$	m_7 0111 $\bar{A}\bar{B}CD$	m_8 1000 $A\bar{B}\bar{C}\bar{D}$	m_9 (1001 $A\bar{B}\bar{C}D$	m_{10} 1010 $A\bar{B}cD$	m_{11} 1011 $A\bar{B}CD$	m_{12} 1100 $AB\bar{C}\bar{D}$	m_{13} 1101 $AB\bar{C}D$	m_{14} 1110 $ABcD$	m_{15} 1111 $ABC\bar{D}$
$\bar{A}\bar{B}$																
01																
AB																
11																
AB																
10																

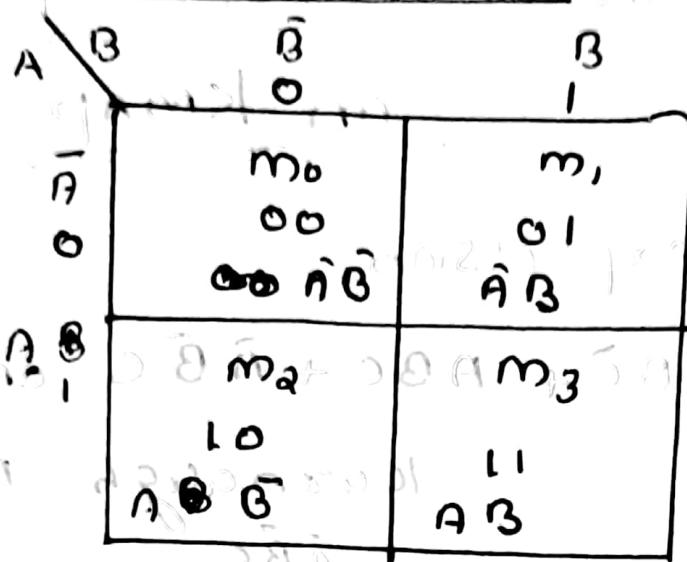
0000	—	0
0001	—	A
0010	—	2
0011	—	3
0100	—	4
0101	—	5
0110	—	6
0111	—	7
1000	—	8
1001	+	9
1010	—	10
1011	—	11
1100	—	12
1101	—	13
1110	—	14
1111	—	15

3 variable K map

qpm A, B, C

\bar{A}	$\bar{B}\bar{C}$	$\bar{B}c$	$B\bar{C}$	Bc
0	m_0 000 $\bar{A}\bar{B}\bar{C}$	m_1 001 $\bar{A}\bar{B}c$	m_3 011 $\bar{A}B\bar{C}$	m_2 010 $\bar{A}B\bar{C}$
1	m_4 100 $A\bar{B}\bar{C}$	m_5 101 $A\bar{B}c$	m_7 111 $AB\bar{C}$	m_6 110 $AB\bar{C}$
0				
1				

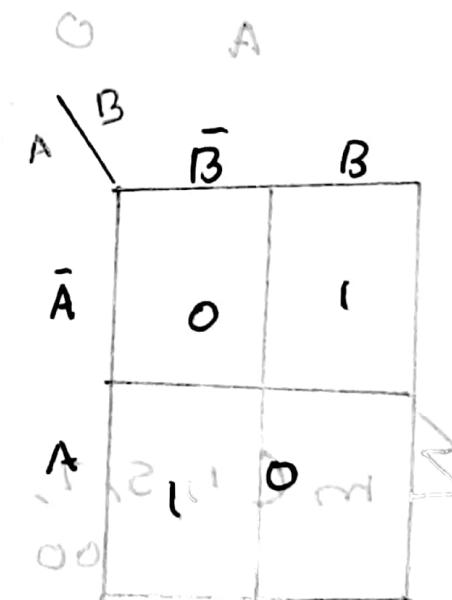
2 Variable KMAP



Representation of Truth Table

on KMAP

A	B	Y
0	0	0
0	1	1
1	0	1
1	1	0



like this we can also do with

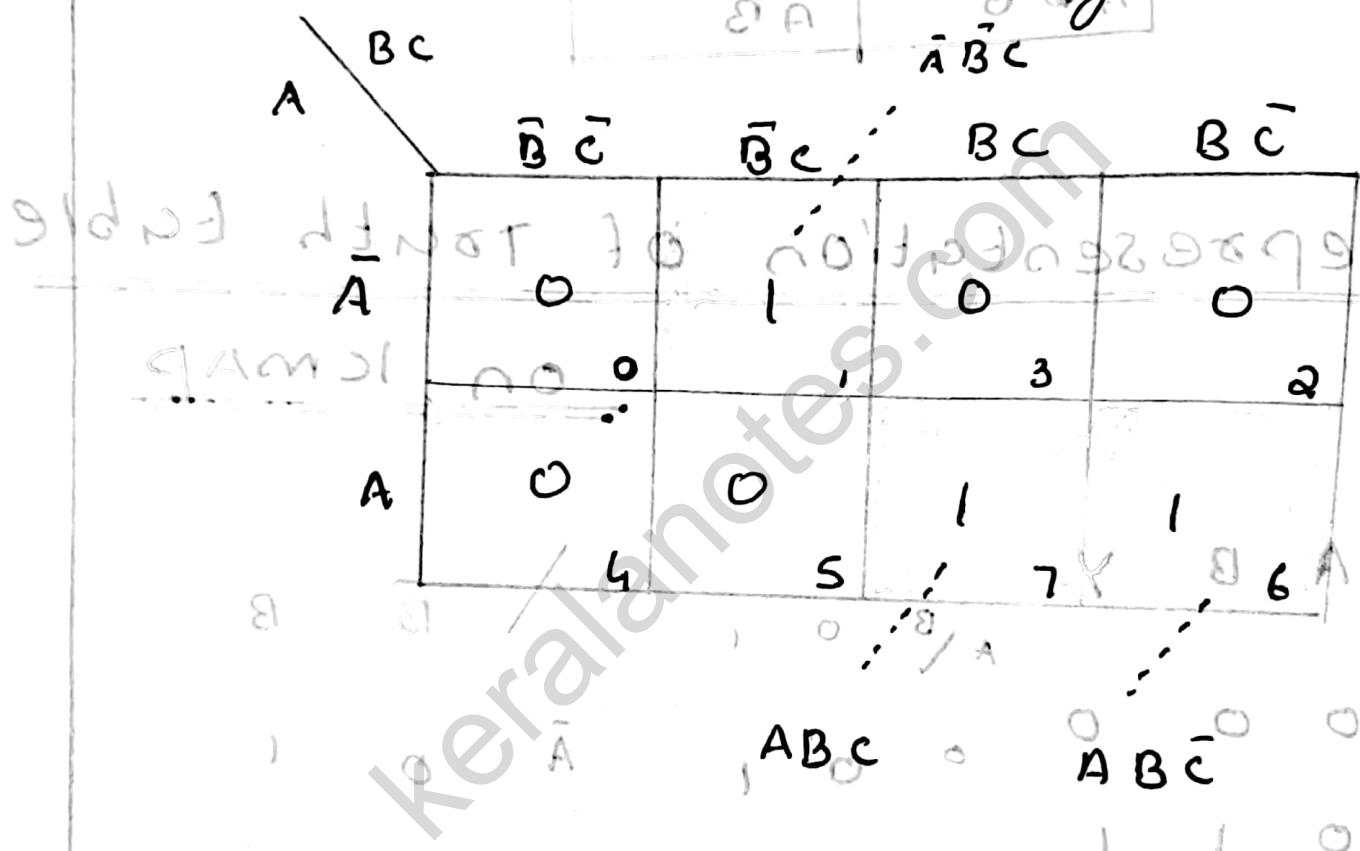
4 variable

11m 21m 31m 41m 11
01m 11m 21m 31m 01

Representation of Boolean Expression on KMAP

→ Plot Boolean expression

$Y = AB\bar{C} + A\bar{B}C + \bar{A}\bar{B}\bar{C}$, on the
Karnaugh map



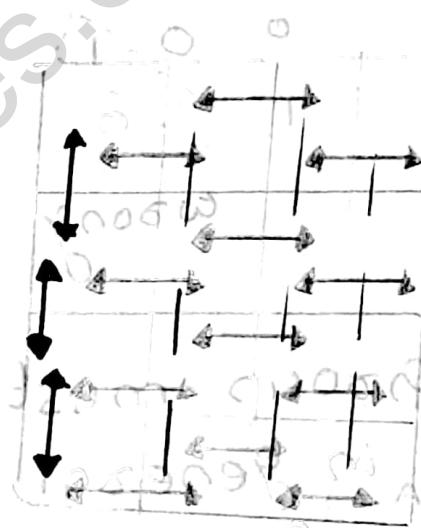
$$\sum m(1, 5, 7, 9, 15)$$

AB	00	01	11	10	11
00	m_0	m_1	m_3	m_2	
01	m_4	m_5	m_7	m_6	
11	m_8	m_{13}	m_{15}	m_{14}	
10	m_9	m_{11}	m_{12}	m_{10}	

1	2	3	4	5	6	7	8
1							
2							
3							
4							
5							
6							
7							
8							

Rules of KMAP

neighboring cells

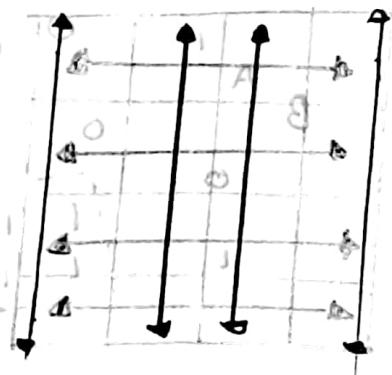


(a) Neighboring cells in the rows are adjacent

(b) Neighboring cells in the columns are adjacent

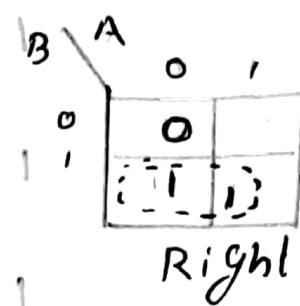
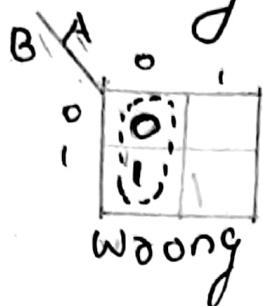
(c) Leftmost and corresponding adjacent rightmost cells are adjacent

(d) Top and corresponding bottom cells are adjacent

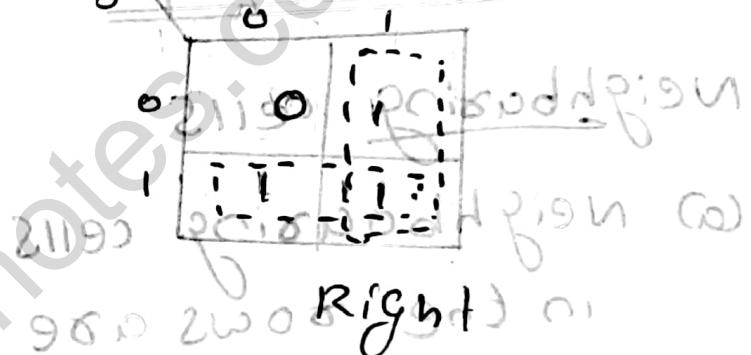


RULES:

1. Group may not include any cells containing a zero



2. Groups may be horizontal or vertical, but not diagonal



3. Group must contain 1, 2, 4, 8, 16 or

In general

~~2, 4, 8, 16 cells~~ [20] region (d)

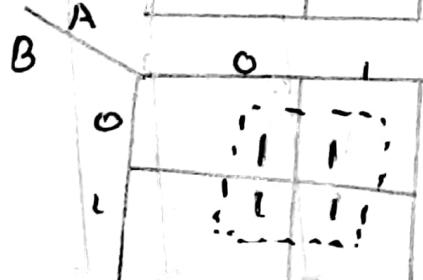


2

Right

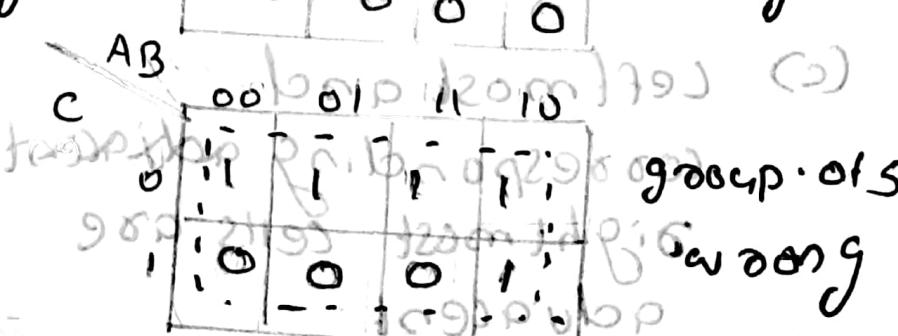


group of 3



Right

group of 4



group of 5

(16)
3
4
2
1

4. Each group should be as large as possible

		AB	C		
		00	01	11	10
0	0	1	1	1	1
	1	0	0	1	1

Right

		AB	C		
		00	01	11	10
0	0	1	1	1	1
	1	0	0	1	1

wrong

5. Each cell containing a one must be in at least one group

		AB	C		
		00	01	11	10
0	0	0	0	11	11
	1	0	0	0	11

Group I
one present
Group II
atleast one
group

6. Groups may overlaps

		AB	C		
		00	01	11	10
0	0	1	1	1	1
	1	0	0	1	1



groups overlaps

Right

		AB	C		
		00	01	11	10
0	0	11	11	11	11
	1	0	0	11	11

Groups
not overlaps

wrong

Some examples: ~~multiple quiesce~~
variable

00	01	11	10
①			
② *	④		
③	0		

00	01	11	10
①			
② *	⑤		
③	0		

AB	CD	00	01	11	10
00	①				
01	*	②		⑤	
11	③		01		

AB	CD	00	01	11	10
00	①				
01	*	②			
11	④				

A	B	C	D	00	01	11	10
0	①	*	③				
1		②			①		

A	B	C	D	00	01	11	10
0	*	②			①		
1		③				②	*

→ Numerical problems: SOP KMAP
[2, 3, 4 variables]

→ DON'T CARE conditions

→ NUMERICAL PROBLEMS

: - Don't care conditions

Grouping of two adjacent ones

Examples:

	$\bar{B}C$	$\bar{B}\bar{C}$	$\bar{B}C$	BC	$B\bar{C}$
	00	01	11	10	
	0	0	1	1	0
\bar{A}	0	0	1	1	0
A	1	0	0	0	0

\bar{A} → common in row

\bar{C} → common in

$\bar{B}C$, BC

$\bar{A} + \bar{B}C$ / row → A common in $\bar{B}\bar{C}$ & $B\bar{C}$
 $\Rightarrow \bar{C}$

	$\bar{B}C$	$\bar{B}\bar{C}$	$\bar{B}C$	BC	$B\bar{C}$
	00	01	11	10	
	0	0	0	0	0
\bar{A}	0	0	0	0	0
A	1	0	0	0	1

→ \bar{A}

A	$B\bar{C}$	$\bar{B}C$	$\bar{B}\bar{C}$	BC	$\bar{B}C$
	0	0	0	1	0
	0	0	0	1	0

Dual CARE (combination
in row no common)

→ \bar{A}

A	$B\bar{C}$	$\bar{B}C$	$\bar{B}\bar{C}$	BC	$\bar{B}C$
	0	0	0	1	0
	0	0	0	1	0

$1 \times (\bar{B}C) \Rightarrow \bar{B}C$ ← Numerical
Group 1 → $\bar{A}C$

Group 2

→ $\bar{B}C$

→ \bar{A}

A	$B\bar{C}$	$\bar{B}C$	$\bar{B}\bar{C}$	BC	$\bar{B}C$
	0	0	0	1	0
	0	0	0	1	0

$\Rightarrow \bar{A}C + BC$

Group 1 → $\bar{A}C$

Group 2

→ $\bar{B}C$

→ $\bar{A}C$ required

→ AB

→ A

A - word
common
3 & 5 &
3 <<

$\Rightarrow \bar{A}C + AB$

Grouping of four adjacent cells

→

A	$\bar{B}C$	$\bar{B}\bar{C}$	$\bar{B}'C$	BC	$B\bar{C}$
\bar{A}	00	01	11	10	00
AB	1	1	1	0	1
	1	1	1	0	1

$\Rightarrow A$

- column not considered because we use all columns so net effect = 1

→

AB	$\bar{B}CD$	$\bar{C}D$	$\bar{C}D$	CD	CD
$\bar{A}\bar{B}$	00	01	11	10	00
$\bar{A}B$	01	0	0	1	0
$A\bar{B}$	11	0	0	1	0
$A\bar{B}$	10	0	0	1	0

CD

→

AB	$\bar{C}D$	$\bar{C}D$	CD	CD	00
$\bar{A}\bar{B}$	00	01	11	10	00
$\bar{A}B$	01	0	0	0	0
$A\bar{B}$	11	0	0	0	0
$A\bar{B}$	10	0	0	0	0

BD

$$Y = BD$$

→

AB	$\bar{C}D$	$\bar{C}D$	CD	CD	00
$\bar{A}\bar{B}$	00	01	11	10	00
$\bar{A}B$	01	0	0	0	0
$A\bar{B}$	11	0	0	0	0
$A\bar{B}$	10	0	0	0	0



AB	$\bar{C}D$	$\bar{C}D$	$\bar{C}D$	CD	CD
$\bar{A}\bar{B}$	00	0	0	0	0
$\bar{A}B$	01	0	0	0	0
$A\bar{B}$	10	1	0	0	1
$A\bar{B}$	11	0	1	1	1

$$Y = AD$$



AB	$\bar{C}D$	$\bar{C}D$	$\bar{C}D$	CD	CD
$\bar{A}\bar{B}$	00	01	10	11	00
$\bar{A}B$	01	0	0	0	0
$A\bar{B}$	10	0	0	0	1
$A\bar{B}$	11	0	0	0	0

$$Y = \bar{B}D$$



AB	$\bar{C}D$	$\bar{C}D$	$\bar{C}D$	CD	CD
$\bar{A}\bar{B}$	00	0	0	0	0
$\bar{A}B$	01	0	0	0	0
$A\bar{B}$	10	0	0	0	1
$A\bar{B}$	11	1	1	1	1

$Y = AB + AD + AC$

Group 1: AB

Group 2: AD

Group 3: AC

Grouping of Eight adjacent cells

→

AB	CD	CD	CD	CD	CD	CD
$\bar{A}\bar{B}$	00	00	01	11	10	01
$\bar{A}B$	00	0	0	0	0	0
$A\bar{B}$	01	1	1	01	01	1
AB	11	1	1	01	01	1
$A\bar{B}$	10	0	0	0	0	0

$y = B$

~~00111011
00000000~~

→

AB	CD	CD	CD	CD	CD	CD
$\bar{A}\bar{B}$	00	00	01	11	10	01
$\bar{A}B$	00	0	1	1	0	0
$A\bar{B}$	01	1	1	01	01	1
AB	11	1	1	01	01	1
$A\bar{B}$	10	0	1	1	0	0

$y = D$

~~00111011
00000000~~

→

AB	CD	CD	CD	CD	CD	CD
$\bar{A}\bar{B}$	00	00	01	11	10	01
$\bar{A}B$	00	1	1	1	1	1
$A\bar{B}$	01	0	0	0	0	0
AB	11	0	0	0	0	0
$A\bar{B}$	10	1	1	1	1	1

$y = \bar{B}$

~~00111011
00000000~~

↓

	$\bar{C}D$	$\bar{C}\bar{D}$	$\bar{C}D$	CD	$C\bar{D}$	
$\bar{A}B$	00	1	0	0	1	0
$\bar{A}B$	01	1	0	0	1	0
$A\bar{B}$	11	1	0	0	1	1
$A\bar{B}$	10	1	0	0	1	1
$\bar{B} = Y$						\bar{D}

$y = \bar{D}$

Simplify Boolean Expression

1. $Y = \underbrace{\bar{A}\bar{B}\bar{D}}_{\text{not Standard SOP}} + \underbrace{A\bar{B}\bar{C}\bar{D}}_{\text{not Standard SOP}} + \underbrace{\bar{A}\bar{B}D}_{\text{not Standard SOP}} + A\bar{B}C\bar{D}$

not Standard SOP so,

$$\begin{array}{c} C \\ \bar{C} \\ A = Y \end{array}$$

$$Y = \underbrace{\bar{A}\bar{B}\bar{C}\bar{D}}_{AB} + \underbrace{\bar{A}\bar{B}\bar{C}\bar{D}}_{CD} + \underbrace{A\bar{B}\bar{C}\bar{D}}_{\bar{C}D} + \underbrace{01\bar{B}A}_{C\bar{D}}$$

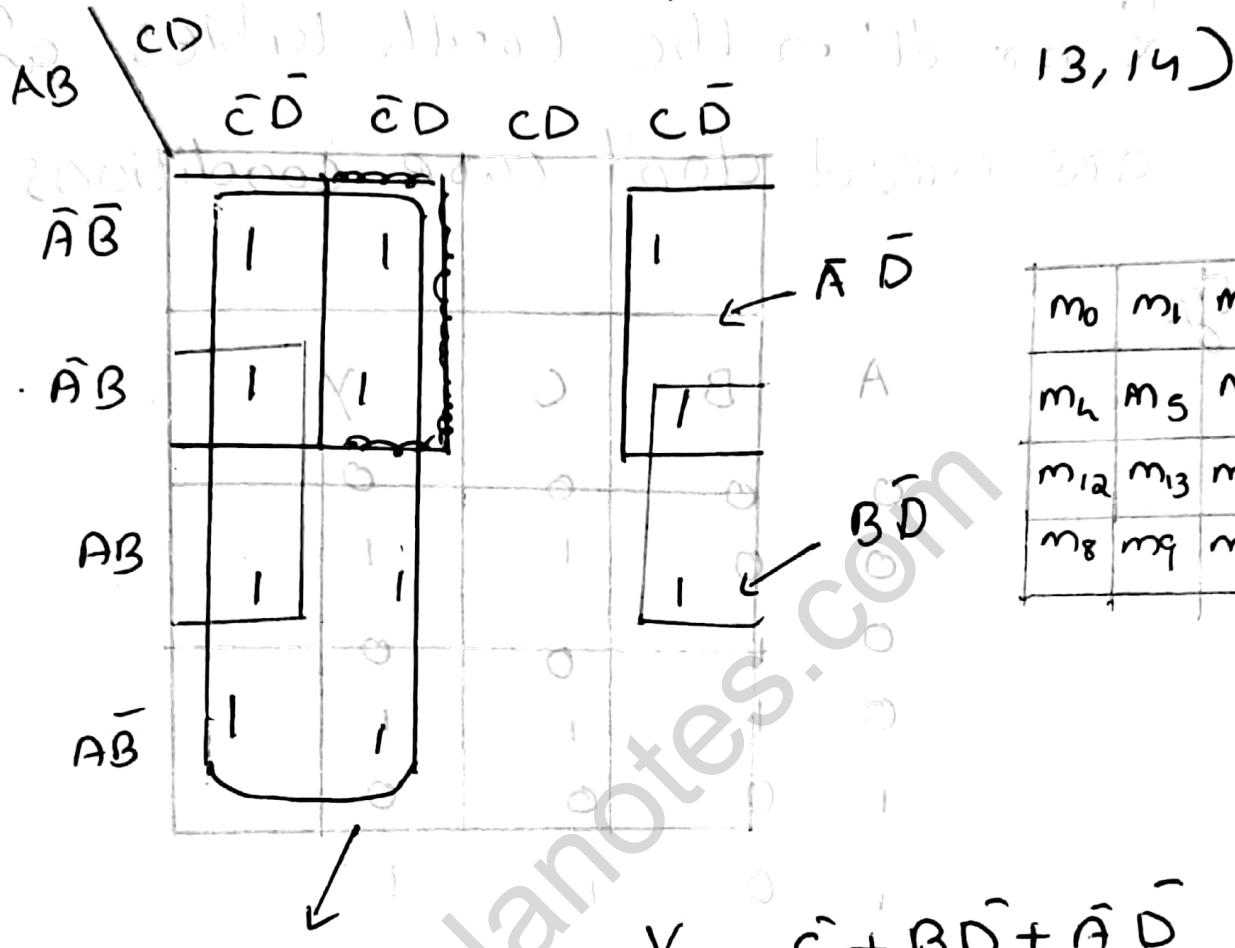
$$\cancel{\bar{A}\bar{B}\bar{C}\bar{D}} + \cancel{\bar{A}\bar{B}\bar{C}\bar{D}} + \cancel{A\bar{B}\bar{C}\bar{D}} + \cancel{01\bar{B}A}$$

	$\bar{C}D$	$\bar{C}\bar{D}$	$\bar{C}D$	CD	$C\bar{D}$	
$\bar{A}\bar{B}$	00	1	0	1	1	0
$\bar{A}B$	01	0	1	0	0	0
$A\bar{B}$	11	0	0	1	0	0
$A\bar{B}$	10	1	1	0	0	1
$\bar{B} = Y$						$\bar{A}\bar{D}$

	$\bar{C}D$	$\bar{C}\bar{D}$	$\bar{C}D$	CD	$C\bar{D}$	
$\bar{A}\bar{B}$	00	1	0	1	1	0
$\bar{A}B$	01	0	1	0	0	0
$A\bar{B}$	11	0	0	1	0	0
$A\bar{B}$	10	1	1	0	0	1
$\bar{B} = Y$						$A\bar{B}$

$$Y = \bar{A}B\bar{D} + \bar{A}\bar{D}$$

Q. $F(A, B, C, D) = \sum m(0, 1, 2, 4, 5, 6, 8, 9, 12, 13, 14)$



m_0	m_1	m_3	m_2
m_4	m_5	m_7	m_6
m_{12}	m_{13}	m_{15}	m_{14}
m_8	m_9	m_{11}	m_{10}

$$Y = \underline{\bar{C} + BD + \bar{A}\bar{D}}$$

Don't Care Terms

(Incompletely specified functions)

- In some logic circuits, certain input conditions never occur. Therefore the corresponding output never appears.

- In such case output level is defined, it can be either HIGH or LOW
- These output level are indicated by 'x' or 'd' in the truth tables and are called don't care conditions.

Eg:

A	B	C	y
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	x
1	1	1	x

} Don't care condition

- when forming groups of cell, treat the don't care cell as either a 1 or 0 or ignore the don't cares.
- This is helpful if it allows us to form a larger groups than would otherwise be possible without the don't cares. There is no requirement

to group all or any of the
only use them in group if it
simplifies the logic

A	$\bar{B}C$	$\bar{B}\bar{C}$	$\bar{B}C$	BC	BC
\bar{A}	0	0	0	0	1
A	1	0	(1)	x	x

$$\text{out} = A \bar{B} C$$

B^C	$\bar{B}\bar{C}$	$\hat{B}C$	BC	$\bar{B}C$
A	00	01	11	10
\bar{A}	0	0	0	0
$A\bar{I}$	0	1	\times	\times
I				

out: AC

Find the Reduced SOP using KMAP

$$F(A, B, C, D) = \sum_m (1, 3, 7, 11, 15) +$$

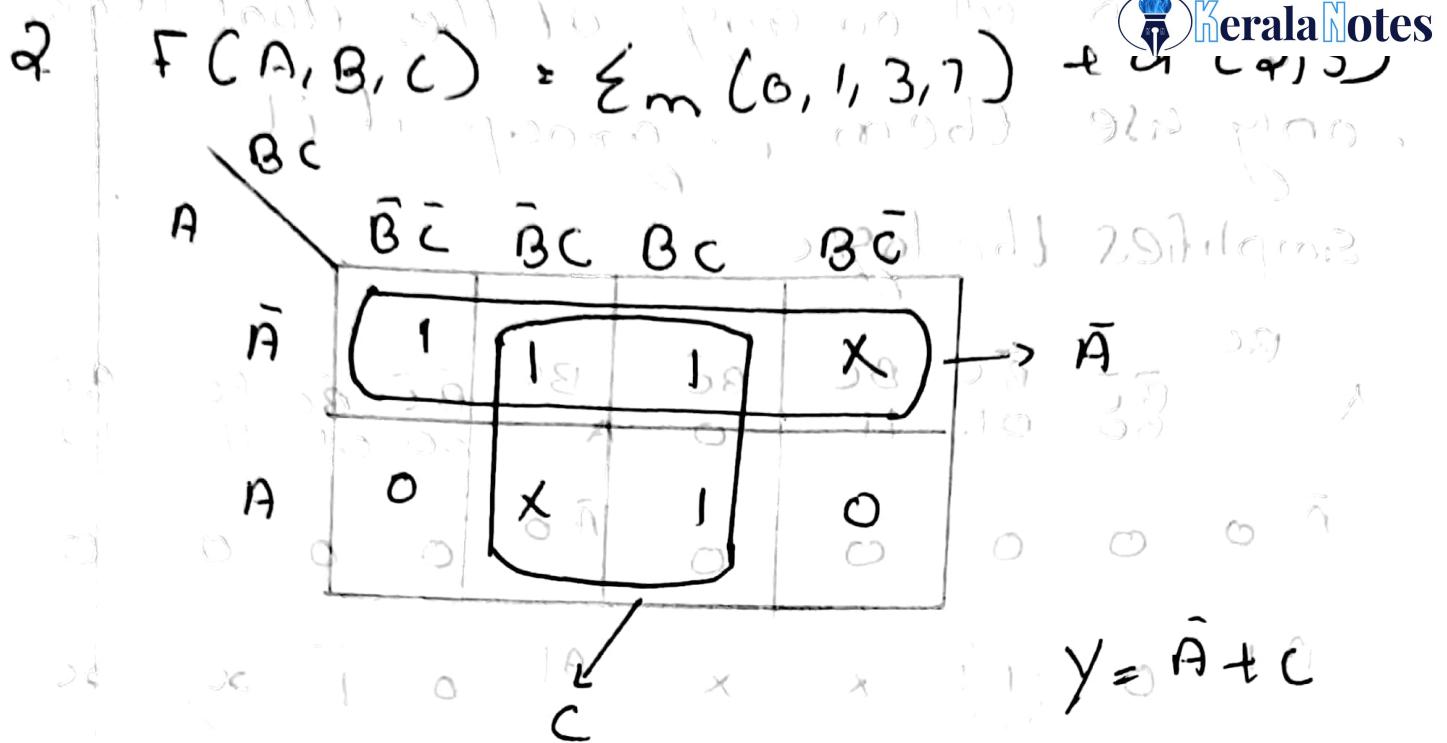
~~(5x20 + 10x10) = CP CD CD = CD~~ don't care

$\bar{A}\bar{B}$	x	1	1	x
$\bar{A}B$	x	0	1	0
AB	0	0	1	0
$A\bar{B}$	0	0	1	0

$$Y = \frac{C\cancel{A}}{\bar{A}\bar{B}} + C\bar{D}$$

$$+ \underset{CD}{\cancel{\sigma \chi \bar{\chi} w}} + \bar{\rho} \bar{\phi} c \bar{w} + \rho \bar{\phi} c \bar{w} \leftarrow$$

$$\cancel{\sigma \bar{\phi} w} + \cancel{\sigma \rho \bar{\phi}}$$



problems

1. Using K-map Simplify Boolean function F as sum of products using the don't care condition of

$$F(w, x, y, z) = w' (x'y\bar{z} + x'yz + xy\bar{z}) + (x + \bar{x})z'(y + w)$$

$$d(w, x, y, z) = w'y\bar{z}(y'z + yz) + wyz$$

$$\Rightarrow F(w, x, y, z) = \overline{w}\overline{x}\overline{y}\overline{z}$$

$$\text{CD} = X \quad \cancel{\overline{x}\overline{y}\overline{z}} + \cancel{\overline{x}\overline{y}\overline{z}} + \cancel{\overline{x}\overline{y}z} + \cancel{xy\overline{z}}$$

$$\Rightarrow \overline{w}\overline{x}y + \overline{w}\overline{x}\overline{y} + \overline{w}\overline{x}yz + \overline{x}\overline{y}z + w\overline{x}\overline{z}$$

$$\bar{\omega} \bar{x} \bar{y} + \bar{\omega} \bar{x} \bar{y} \bar{z} + \bar{\omega} x \bar{y} z + \bar{\omega} x \bar{y} \bar{z}$$

$$\frac{z}{\bar{z}} \quad \frac{z}{\bar{z}} \quad \text{Simplifying} \quad \frac{\omega}{\bar{\omega}} \quad \frac{y}{\bar{y}}$$

$$F(\bar{\omega}, x, y, z) :$$

Simplifying X

$$\Rightarrow \bar{\omega} \bar{x} \bar{y} z + \bar{\omega} \bar{x} \bar{y} \bar{z} + \bar{\omega} x \bar{y} z +$$

$$\bar{\omega} x \bar{y} \bar{z} + \bar{\omega} x y z + \bar{\omega} x \bar{y} \bar{z} + \cancel{\bar{\omega} \bar{x} y \bar{z} + \bar{\omega} \bar{x} \bar{y} \bar{z}}$$

$$\cancel{\bar{\omega} \bar{x} y \bar{z} + \bar{\omega} \bar{x} \bar{y} \bar{z}} + \cancel{\bar{\omega} x y \bar{z}}$$

$$d(\bar{\omega}, x, y, z) \Rightarrow$$

$$\bar{\omega} x \bar{y} z + \bar{\omega} x y z + w y z$$

$$\begin{matrix} x \\ \bar{x} \end{matrix}$$

$$\Rightarrow \bar{\omega} x \bar{y} z + \bar{\omega} x y z + w x y z + w \bar{x} y z$$

$$\cancel{w x} \quad \cancel{y z} \quad \cancel{y z} \quad \cancel{y z} \quad \cancel{y z} \quad \cancel{x z}$$

$\bar{\omega} \bar{x}$	1	1	1	1
$\bar{\omega} x$	1	1	1	1
$w \bar{x}$	1	1	1	1
$\bar{x} z$	1	1	1	1

$$\bar{\omega} x$$

$$(w x)$$

$$w \bar{x}$$

$$\bar{x} z$$

X

X

1

$$\bar{\omega} z$$

$$Y = \overline{AB} + \overline{C}\overline{D} + \overline{B}\overline{D}$$

~~$$Y = \overline{AB} + \overline{C}\overline{D}$$~~

: (s, p, x, w) 7

$$Y = \overline{B}\overline{D} + \overline{C}\overline{D} + \overline{B}\overline{C}$$

$$+ \overline{B}\overline{C}\overline{W} + \overline{B}\overline{C}\overline{W} + \overline{B}\overline{C}\overline{W} + \overline{B}\overline{C}\overline{W}$$

~~$$\overline{B}\overline{C}\overline{W} + \overline{B}\overline{C}\overline{W}$$~~

$= (s, p, x, w) 10$

$$\overline{B}\overline{W} + \overline{B}\overline{C}\overline{W} + \overline{B}\overline{C}\overline{W}$$

\overline{x}
 \overline{w}

2. Reduce the following expression using k-map and implement the real minimal expression in universal logic

1. $F(A, B, C, D) = \Sigma_m(0, 1, 2, 3, 5, 7, 8, 9, 10, 12, 13)$

1 x

1 $\overline{x}w$

\overline{s}

$\bar{A}B$	$\bar{C}D$	$\bar{C}\bar{D}$	CD	$C\bar{D}$	$\bar{C}\bar{D}$
$\bar{A}\bar{B}$	$\bar{x} + \bar{w}$	$x + w$	$x + \bar{w}$	$\bar{x} + \bar{w}$	$\bar{x} + w$
$\bar{A}B$	$\bar{x} + w$	$x + \bar{w}$	$x + w$	$\bar{x} + \bar{w}$	$\bar{x} + w$
$A\bar{B}$	$x + \bar{w}$	$\bar{x} + w$	$\bar{x} + \bar{w}$	$x + w$	$x + \bar{w}$
AB	$x + w$	$\bar{x} + \bar{w}$	$\bar{x} + w$	$x + \bar{w}$	$\bar{x} + \bar{w}$

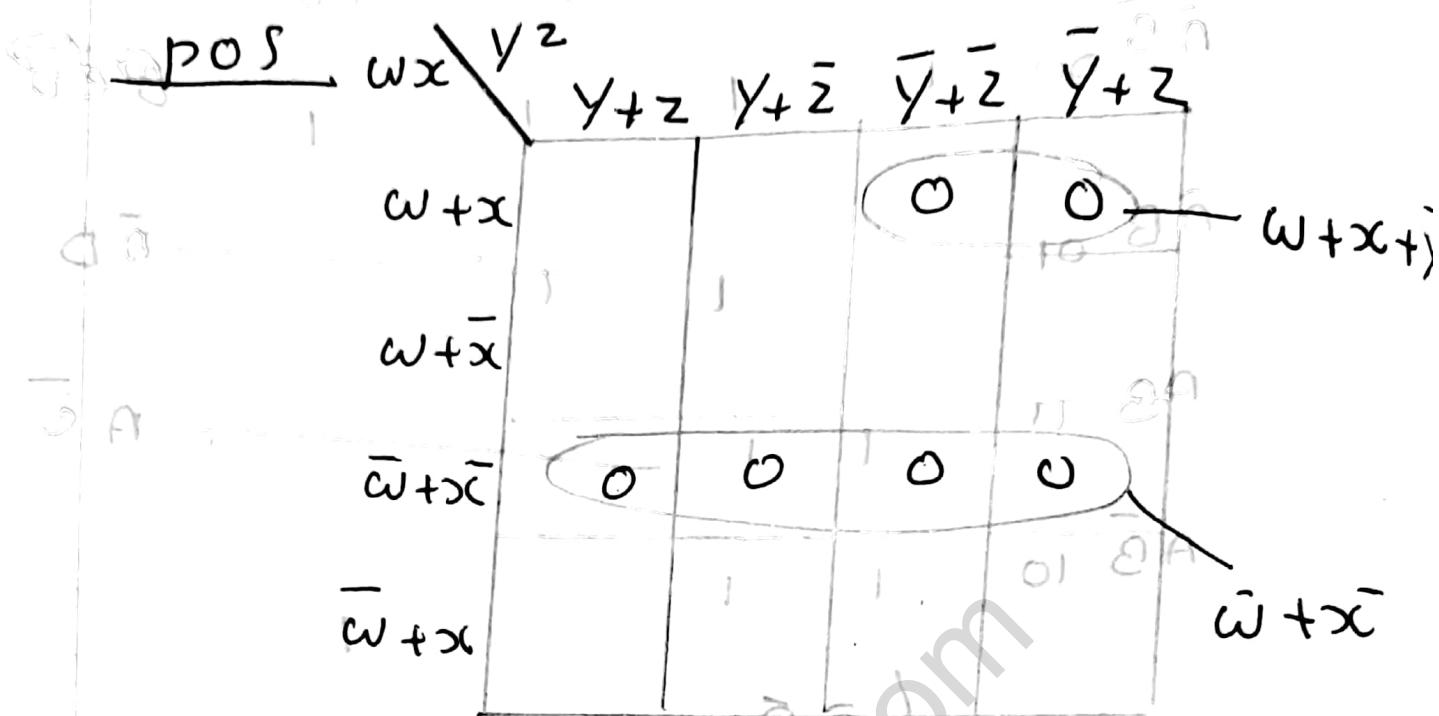
$$(\bar{x} + w) \cdot (\bar{x} + \bar{w}) = \bar{x}$$

$$Y = \bar{x} + w + \bar{A}D + \bar{A}\bar{C} + \bar{B}\bar{C}D + \bar{B}\bar{D}$$

3. Simplify the given Boolean function f
 $f(w, x, y, z) = \Sigma(2, 3, 12, 13, 14, 15)$
(i) sum of product (ii) product of sum

wx	y_2	\bar{y}_2	\bar{y}_2	y_2	\bar{y}_2	$\bar{x}w$
$\bar{w}\bar{x}$	\bar{y}_2	\bar{y}_2	\bar{y}_2	y_2	\bar{y}_2	$\bar{w}\bar{x}y$
$\bar{w}x$	\bar{y}_2	\bar{y}_2	y_2	\bar{y}_2	\bar{y}_2	xw
wx	\bar{y}_2	\bar{y}_2	\bar{y}_2	\bar{y}_2	\bar{y}_2	$\bar{x}w$
$\bar{x}w$	\bar{y}_2	\bar{y}_2	\bar{y}_2	\bar{y}_2	\bar{y}_2	\bar{y}_2

$$\text{SOP } y = \bar{w_1} \bar{x_1} y + w_2 x_2$$



$$\text{pos } \dot{\bar{y}} = (\omega + x + \bar{y}) \cdot (\bar{\omega} + \bar{x})$$

Simplify the Boolean function

$$F(w, x, y, z) = \sum_{m=0}^8 (0, 5, 7, 9, 10, 11, 14, 15)$$

~~Diagram 5 - (5, 14, x, y, w, o) 7~~

$$\omega_x \quad \begin{array}{|c|c|} \hline 1 & 1 \\ \hline \end{array} \quad \omega_x \quad \overbrace{\quad \quad \quad}^{\omega_{xz}} \quad \omega_x$$

$$wx - wy = \cancel{wx}$$

$\omega \bar{x}$

$$Y = \bar{x}\bar{y}\bar{z} + \bar{w}xz + wy + w\bar{z}$$

Q. Reduce the following function using Karnaugh map technique and implement using basic gates.

$$F(A, B, C, D) = \bar{A}\bar{B}D + A\bar{B}\bar{C}D + \bar{A}BD + ABC\bar{D}$$

$$\begin{array}{c} C \\ \bar{C} \end{array} \quad \begin{array}{ccccc} A & B & C & D & F \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 & 1 \\ 0 & 1 & 1 & 0 & 1 \\ 0 & 1 & 1 & 1 & 1 \end{array} \quad \begin{array}{c} C \\ \bar{C} \end{array}$$

$$\Rightarrow \bar{A}\bar{B}CD + \bar{A}\bar{B}\bar{C}D + A\bar{B}\bar{C}\bar{D} + \bar{A}BCD + \bar{A}B\bar{C}D + A\bar{B}C\bar{D}$$

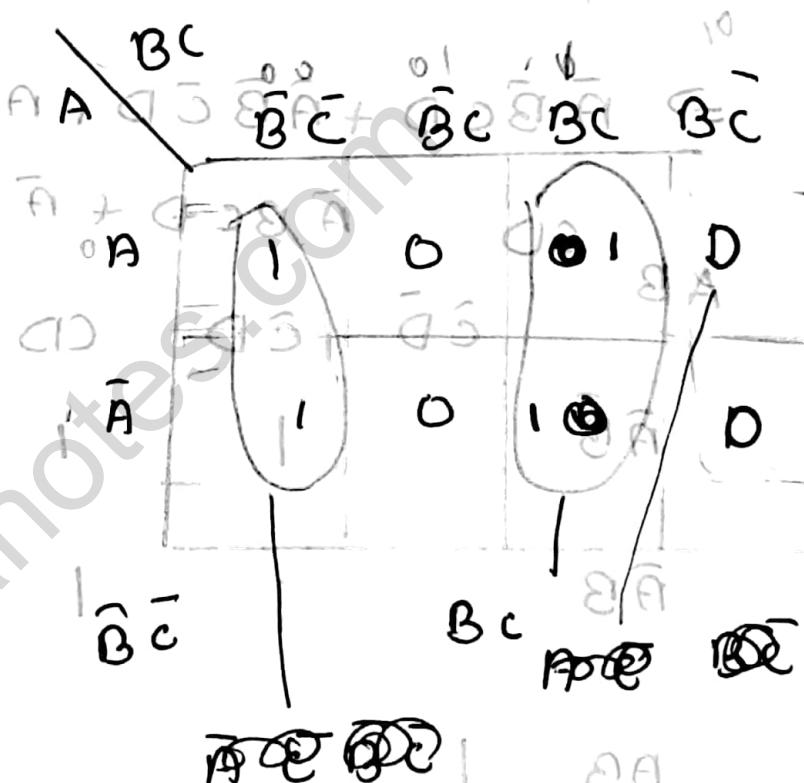
	$\bar{A}\bar{B}$	$\bar{A}B$	AB	$A\bar{B}$	
$\bar{C}\bar{D}$	0	1	1	0	$\bar{A}D$
$\bar{C}D$	1	0	0	0	
CD	0	1	1	1	
$C\bar{D}$	1	1	0	0	
A	0	0	0	1	
B	0	1	1	1	

$$Y = \bar{A}D + A\bar{B}\bar{D}$$

$$+ \bar{A}C\bar{B}\bar{D} + \bar{A}B\bar{C}\bar{D} + \bar{A}B\bar{C}D + \bar{A}B\bar{D} + A\bar{B}D = Y$$

6. Simplify the logic function by the truth table using the Karnaugh map method. y is the output variable and A, B , and C are the input variables.

A	B	C	y
0	0	0	1
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

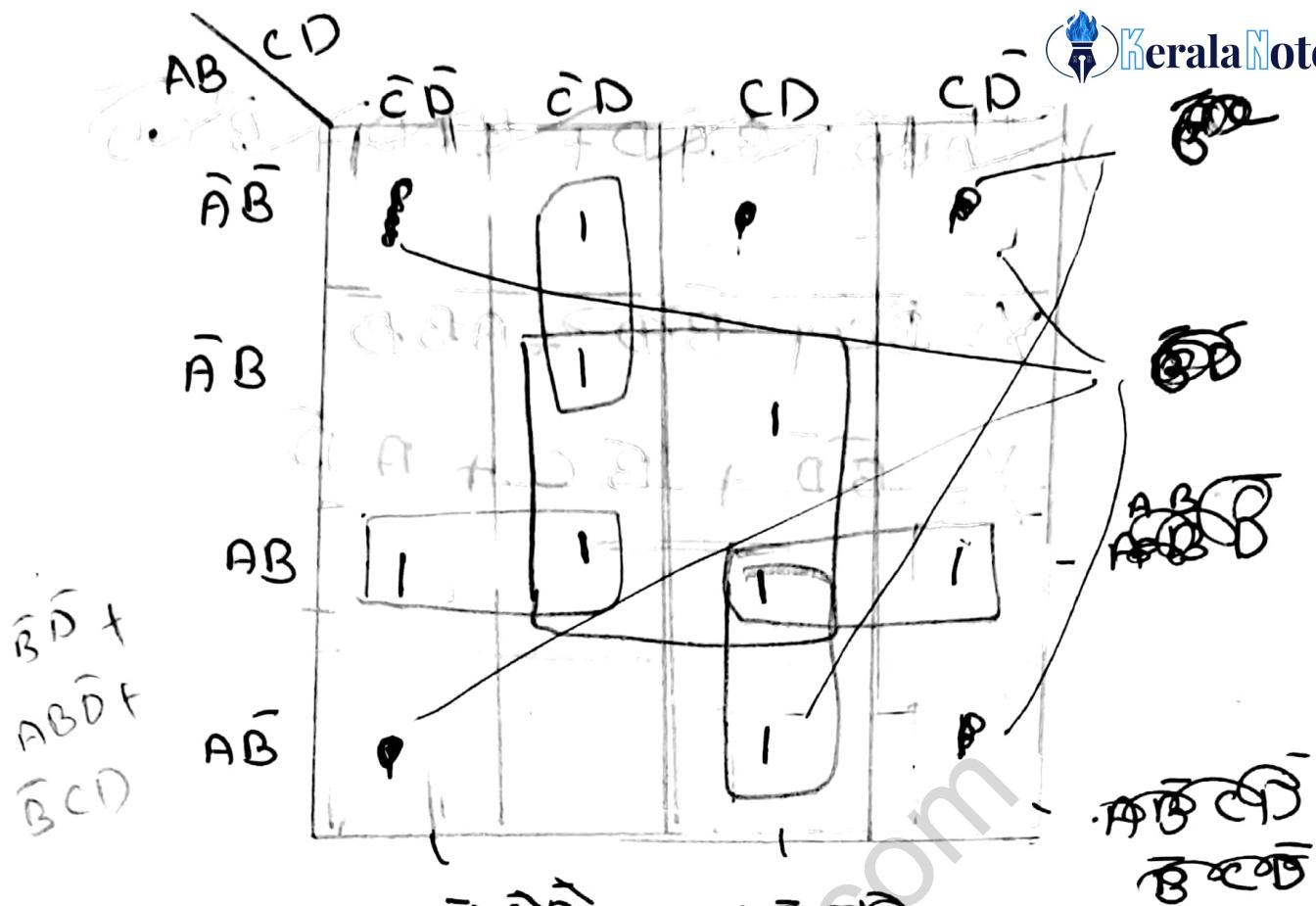


$$Y = \cancel{ABC} + \cancel{ACB} \quad \cancel{BC} + \cancel{B'C} \quad Y = \cancel{C}$$

$$Y = \underline{\bar{B}'\bar{C}} + \underline{B'C}$$

7. Reduce the following function to its minimum sum of product form:

$$Y = \bar{A}\bar{B}\bar{C}\bar{D} + \bar{A}B\bar{C}\bar{D} + \bar{A}\bar{B}C\bar{D} + \cancel{A}\bar{B}C\bar{D} + A\bar{B}\bar{C}\bar{D} + A\bar{B}\bar{C}\bar{D} + \cancel{A}\bar{B}CD + \cancel{A}\bar{B}\bar{C}D$$



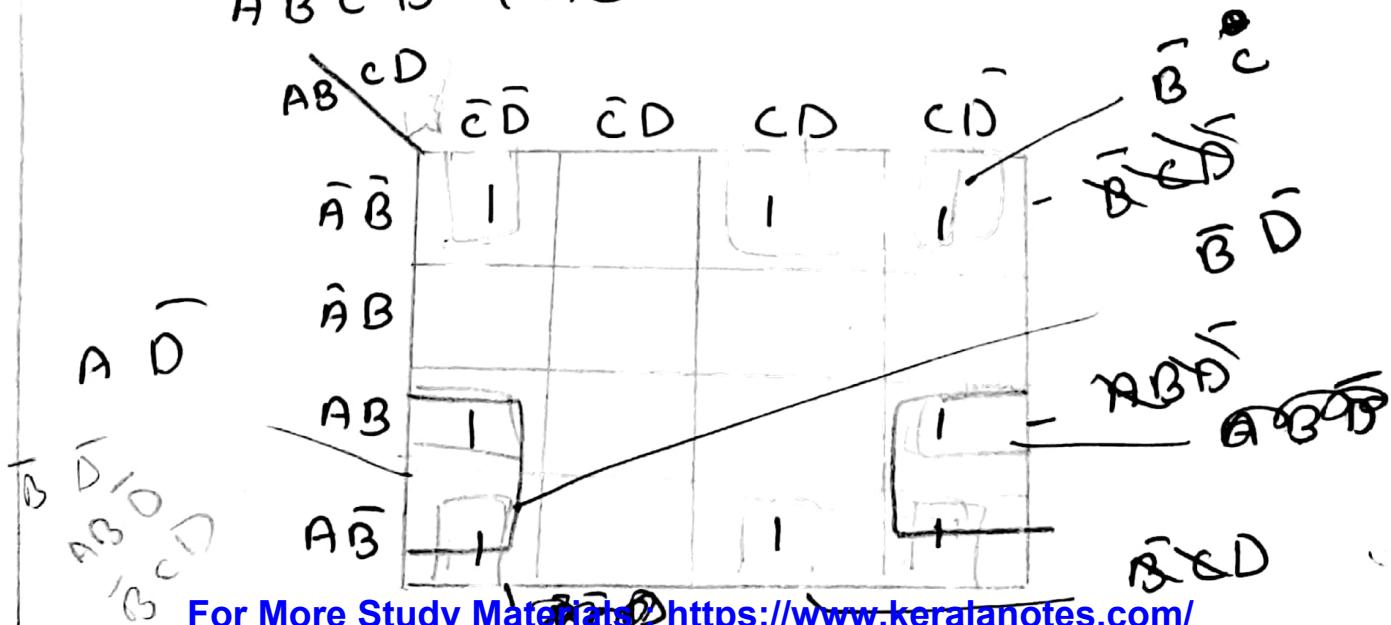
$$Y = \bar{A}\bar{C}D + A\bar{B}\bar{C} + \bar{A}B\bar{C}\bar{D}$$

$$ACD + \bar{A}BC$$

$$Y = A\bar{B}D + \bar{A}\bar{C}D + \bar{B}\bar{C}D + \bar{B}\bar{C}\bar{D}$$

Reduce the following four variable function to the minimum sum of product form

$$Y: \bar{A}\bar{B}\bar{C}\bar{D} + A\bar{B}CD + A\bar{B}C\bar{D} + A\bar{B}CD + A\bar{B}\bar{C}\bar{D} + AB\bar{C}\bar{D} + \bar{A}BCD + \bar{A}\bar{B}\bar{C}\bar{D}$$



$$Y = \cancel{AB\bar{D}} + \cancel{B\bar{C}\bar{D}} + \cancel{\bar{B}CD} + \cancel{A\bar{B}\bar{D}}$$

$$Y = \cancel{B\bar{C}D} + \cancel{\bar{B}CD} + \cancel{A\bar{B}\bar{D}}$$

$$Y = \cancel{\bar{B}\bar{D}} + \underline{BC} + \cancel{AD}$$

പ്രാഥിക

സംഖ്യ

പ്രാഥിക

$\bar{A}\bar{B}\bar{C} + \bar{C}\bar{B}\bar{A} + \bar{B}\bar{C}\bar{A}$

ഒരു

$\bar{A}\bar{B}\bar{C} + \bar{B}\bar{C}\bar{A} + \bar{C}\bar{A}\bar{B}$

$\bar{A}\bar{B}\bar{C} + \bar{A}\bar{B}\bar{C} + \bar{A}\bar{B}\bar{C} + \bar{A}\bar{B}\bar{C}$

$\bar{A}\bar{B}\bar{C} + \bar{A}\bar{B}\bar{C}$

ഇലാസ്റ്റിനൈറ്റിനു സൂത്രം ഫോളോഓം ചെയ്യാം

ഈ സൂത്രം നിലനിൽക്കുന്നതു എന്ന് പറയാം

സൂത്രം ലഭിച്ചു

$+ \bar{A}\bar{C}\bar{B} + \bar{A}\bar{B}\bar{C} + \bar{A}\bar{B}\bar{C} + \bar{A}\bar{B}\bar{C} + \bar{A}\bar{B}\bar{C} : Y$

$\bar{A}\bar{B}\bar{C} + \bar{A}\bar{B}\bar{C} + \bar{A}\bar{B}\bar{C} + \bar{A}\bar{B}\bar{C}$

\bar{A}	\bar{B}	\bar{C}	\bar{D}	$\bar{A}\bar{B}$	$\bar{A}\bar{C}$	$\bar{A}\bar{D}$	$\bar{B}\bar{C}$	$\bar{B}\bar{D}$	$\bar{C}\bar{D}$	$\bar{A}\bar{B}\bar{C}$	$\bar{A}\bar{B}\bar{D}$	$\bar{A}\bar{C}\bar{D}$	$\bar{B}\bar{C}\bar{D}$	$\bar{A}\bar{B}\bar{C}\bar{D}$
1	1	1	1	1	1	1	1	1	1	1	1	1	1	1

സൂത്രം

$\bar{A}\bar{B}\bar{C}$

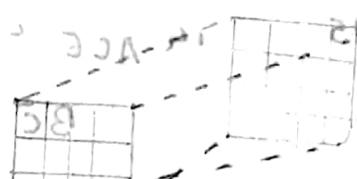
part - 3

5 Variable K-map

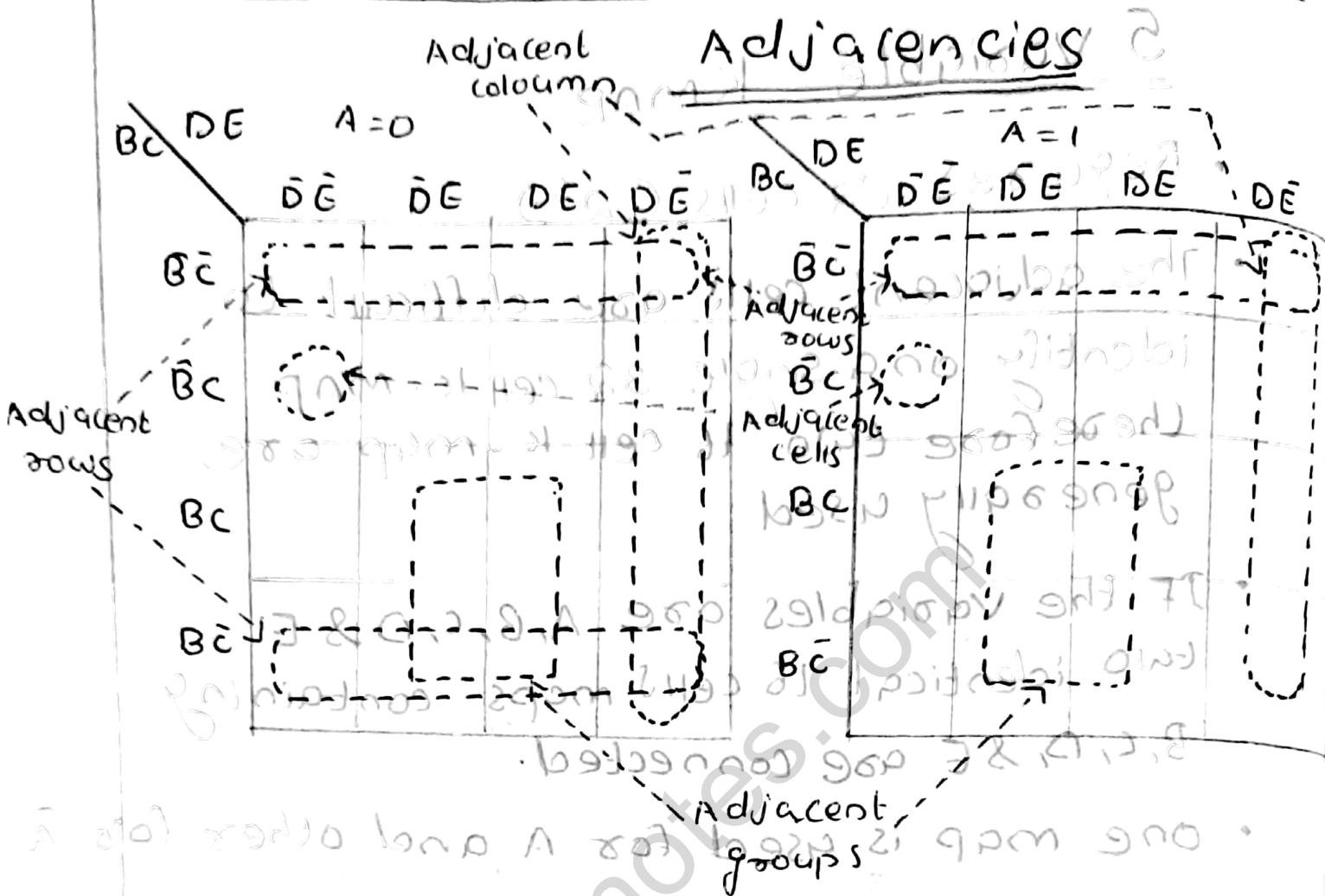
- Requires 32 cells (2^5)
 - The adjacent cells are difficult to identify on a single 32 cell IC-map. Therefore two 16 cell IC-maps are generally used.
 - If the variables are A, B, C, D & E two identical 16 cell maps containing B, C, D, & E are connected.
 - one map is used for A and other for \bar{A}

A = 0 / Å			
BC	DE	DE	DE
BC	DE	DE	DE
BC 00	0	1	3
BC 01	4	5	7
BC 11	12	13	15
BC 10	8	9	11

$\bar{B}\bar{C}$	00	01	11	10
$\bar{B}\bar{C} = 00$	(3, 4, 5, 6, A)	7	8	9
	16	17	19	18
$\bar{B}\bar{C} = 01$	20	21	23	22
$\bar{B}\bar{C} = 11$	28	29	31	30
$\bar{B}\bar{C} = 10$	24	25	27	26



5 variable KMAP with ex



problems

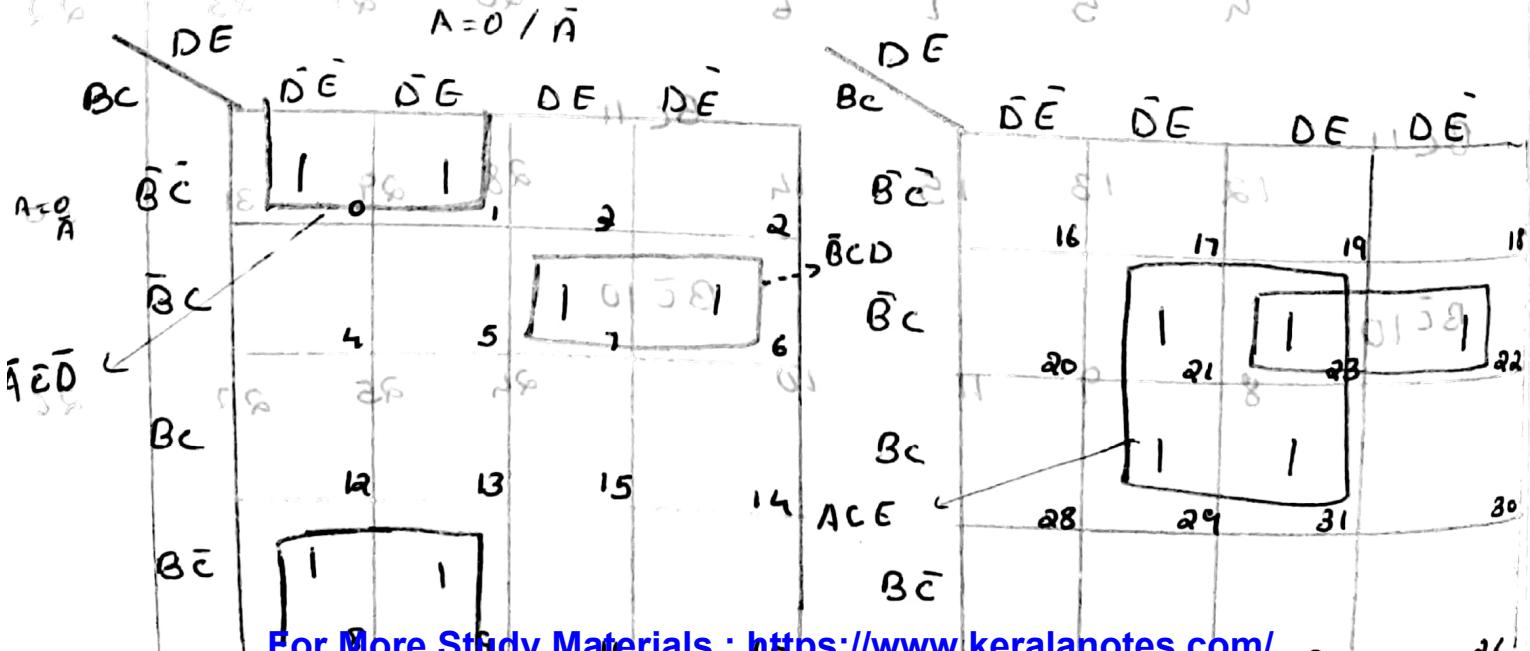
$$1. F(A, B, C, D, E) = \sum m(0, 1, 6, 7, 8, 9, 21, 22, 23, 29, 31)$$

0-7 → 3 variable

0-15 → 4 variable

0-31 → 5 variable

A=0 / A

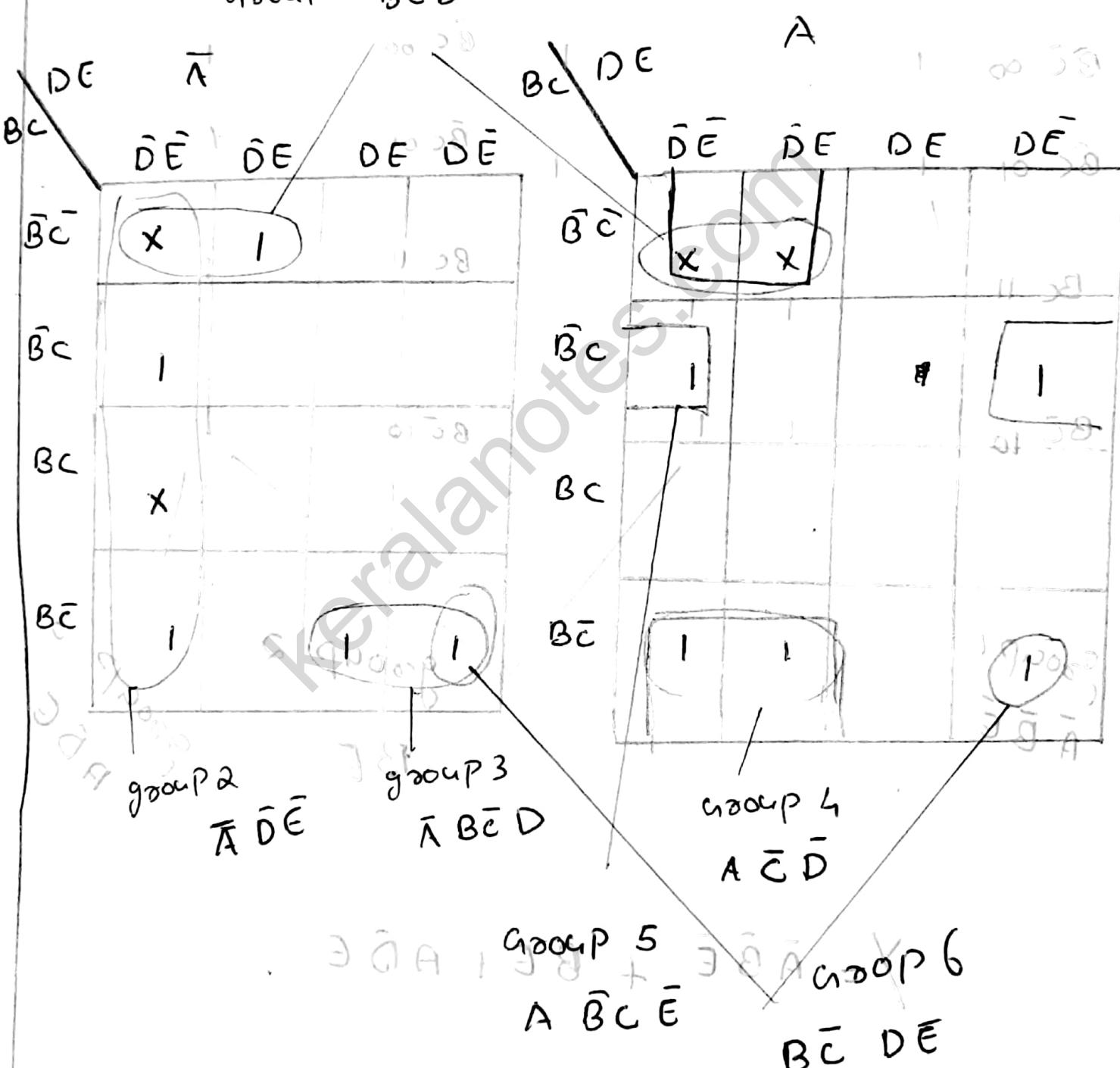


$$Y = \bar{A}\bar{C}\bar{D} + \bar{B}\bar{C}D + A\bar{C}\bar{E}$$

$$f(A, B, C, D, E) = \sum m(1, 6, 8, 10, 11, 20, 23) \\ (18, 26, 27, 28, 24, 25, 26)$$

$$\bar{d}(0, 12, 16, 17)$$

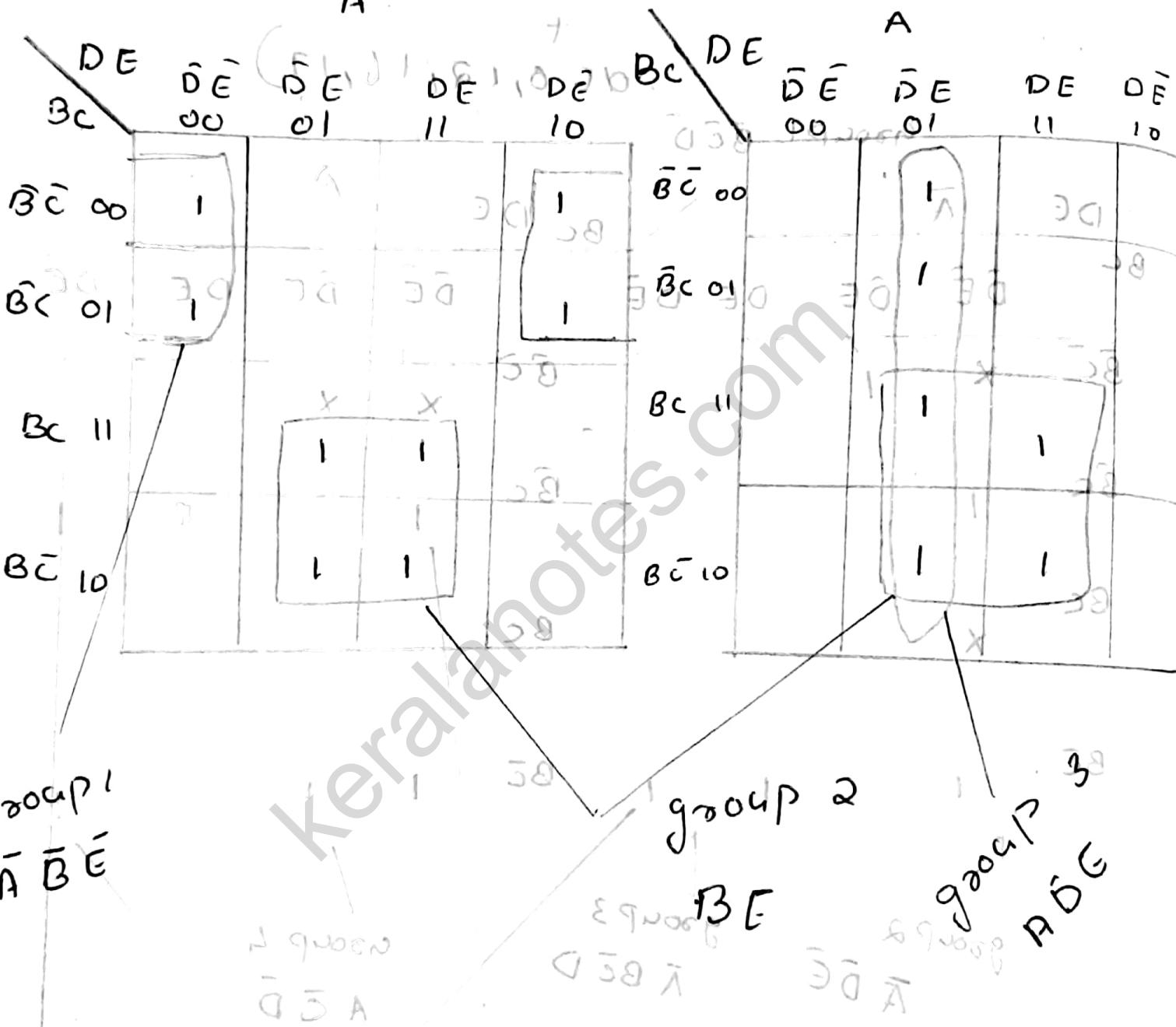
group 1 $\bar{B}\bar{C}\bar{D}$



$$Y = \bar{B}\bar{C}\bar{D} + \bar{A}\bar{D}\bar{E} + \bar{A}B\bar{C}D + A\bar{C}\bar{D} + A\bar{B}C\bar{E}$$

3 Simplify the boolean function

$$f(A, B, C, D, E) = \sum m(0, 2, 4, 6, 9, 11, 13, 15, 17, \\ 18, 20, 21, 25, 27, 29, 31)$$



$$Y = \bar{A}\bar{B}\bar{E} + B\bar{E} + A\bar{D}\bar{E}$$

$$\bar{B}\bar{C}\bar{A} + \bar{B}\bar{C}A + B\bar{C}\bar{A} + \bar{B}CA + \bar{B}\bar{C}\bar{A} = Y$$

$$\bar{B}\bar{C}\bar{A} + \bar{B}\bar{C}A + B\bar{C}\bar{A} + \bar{B}CA + \bar{B}\bar{C}\bar{A} = Y$$

$$F(A, B, C, D, E) = \sum m(0, 5, 6, 8, 17, 18, 21, 16, 20, 24, 25, 26, 27, 29, 31)$$

	\bar{A}	$D\bar{E}$	$D\bar{E}$	$D\bar{E}$	$B\bar{C}$	$D\bar{E}$	A	$D\bar{E}$	$D\bar{E}$
$\bar{B}\bar{C} 00$	1					1			
$\bar{B}\bar{C} 01$		1				1			
$B\bar{C} 11$					1				
$B\bar{C} 10$	1	1	1	1			1	1	1

group 1

group 5

group 4

group 2

group 6

$A\bar{B}\bar{D}\bar{E}$

group 2

$\bar{A}\bar{B}\bar{C}\bar{D}\bar{E}$

group 3

$\bar{A}\bar{B}CD\bar{E}$

$(\bar{C}\bar{D}\bar{E})$

\bar{B}

\bar{B}

\bar{B}

Simplification of 3 POS Expression

(MAX TERMS KMAP)

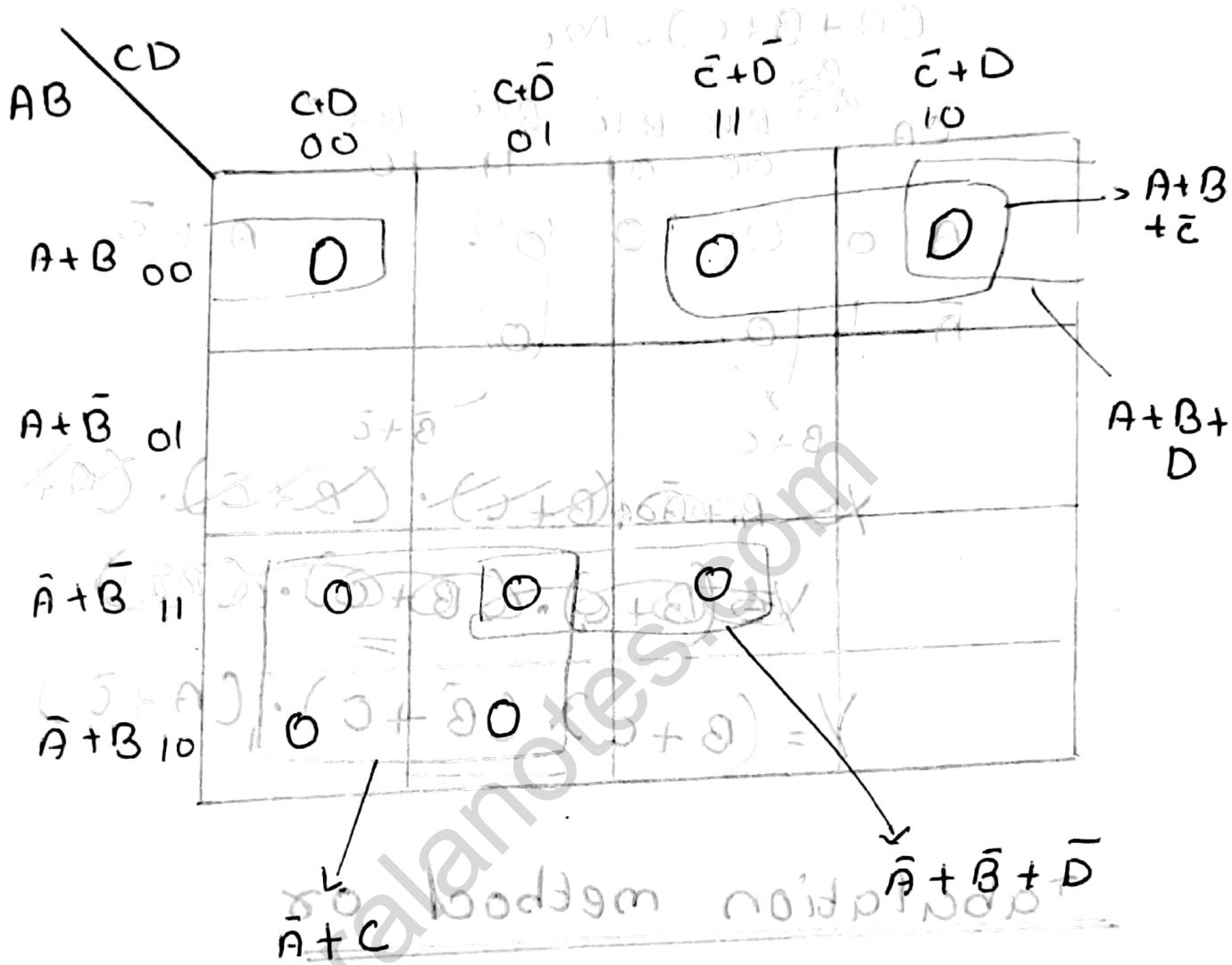
steps to solve

1. plot the k-map and place 1's in those cells corresponding to the 1's in the TT-0s Maxterms in the product of sum Expressions
2. check the kmap for adjacent 1's and encircle those 1's
3. Form the simplified pos Expression by taking pos terms of all the groups

		CD	C+D	C+D	$\bar{C}+\bar{D}$	$\bar{C}+D$
		AB	$A \oplus B$	$A \oplus B$	$\bar{A} \oplus B$	$\bar{A} \oplus B$
		00	01	11	10	(J43)
$A+B$	00	$A+B+C+$ 0	$A+B+C+$ 0	M_0	M_1	M_3
$A+\bar{B}$	01	M_4	M_5	M_7	M_6	
$\bar{A}+\bar{B}$	11	M_{12}	M_{13}	$\bar{A}+\bar{B}+\bar{C}+\bar{D}$ M_{15}	M_{14}	
$\bar{A}+B$	10	M_8	M_9	M_{11}	M_{10}	

problems

1. $F(A, B, C, D) = \sum m(0, 2, 3, 8, 9, 12, 13, 15)$



~~$$F = (A + C)(A + B + D)(A + B + C)$$~~

minimise the expression

2.

$$Y = (A + B + C)(A + B + C̄)(A + B + C̄)$$

$(A + B + C)(A + B + C̄) = 1$ (bottom right corner)

Ans:

$$(A+B+C) = m_0, (A+\bar{B}+\bar{C}) = m_3,$$

$$(\bar{A}+B+\bar{C}) = m_7, (A+\bar{B}+C) = m_4,$$

$$(A+B+C) = M_0$$

	$\bar{B}\bar{C}$	$B\bar{C}$	$B\bar{C}$	$\bar{B}C$	$\bar{B}C$
\bar{A}	00	01	11	10	00
A	0	0	0	0	1
	$B+C$	$\bar{B}+C$	$\bar{B}+C$	$\bar{B}+C$	$\bar{B}+C$

$$A + \bar{C}$$

$$X = \cancel{(A+\bar{B}+0)} \cdot \cancel{(B+C)} \cdot \cancel{(A+\bar{C})}$$

$$\cancel{Y = (B+0)} \cdot \cancel{(B+C)} \cdot \cancel{(A+\bar{C})}$$

$$Y = \underline{(B+C)} \cdot \underline{(B+C)} \cdot \underline{(A+\bar{C})}$$

Tabulation method or

Quine McCluskey method

- Based on the concept of prime implicants
- Prime implicant is a product of sum term

$$(S+B+A) (S+\bar{B}+A) (S+B+A) = Y$$

- Q Simplify Boolean expression using Quine McCluskey method (Tabular method)

$$F(A, B, C, D) = \sum m(0, 1, 3, 7, 8, 9, 11, 13)$$

Step 0: Convert the decimal to Binary number

Decimal Number	Equivalent Binary Number	minterms
0	0000	m_0
1	0001	m_1
2	0011	m_2
3	0111	m_3
4	1000	m_4
5	1001	m_5
6	1011	m_6
7	1101	m_7
8		
9		
10		
11		
12		
13		
14		
15		

Step 1: Arrange the minterm according to Number's of 1's

Group	minterm no.	In Binary			
		A	B	C	D
0	0	0	0	0	0
1	1	0	0	0	1
2	(8, 15)	1	0	0	0
3	(3, 11)	0	0	1	1
4	9	1	0	0	1
5	7	0	1	1	1
6	11	1	1	0	1
7	15	1	1	1	1

Step 2: • Compare each minterm in Group n with each minterm in Group $(n+1)$ and identify the matched pairs

- A matched pair is a pair of minterms which differ only in one variable (gray code)

S m.	M in te rm	In Binary				S m.	M in te rm	In Binary					
		A B C D						A B C D					
		A	B	C	D			A	B	C	D		
0	0	0	0	0	0	0	0	0	0	0	-		
1	-	0	1	0	0	0	0	1	0	0	0		
2	3	0	1	0	1	1	1	0	0	-	1		
3	7	0	1	1	1	2	1	1	0	0	-		
4	15	0	1	1	0	0	0	0	1	1	1		
5	0	0	0	0	0	1	0	0	0	1	1		
6	8	1	0	0	0	1	0	0	0	0	0		
7	12	1	1	0	0	2	1	0	0	0	1		
8	11	1	1	1	0	3	1	1	0	1	1		
9	13	1	1	0	1	4	1	0	1	0	-		
10	14	1	1	1	1	5	1	0	0	1	1		
11	16	1	0	1	1	6	0	1	0	0	-		
12	17	1	0	1	0	7	0	1	1	0	1		
13	18	1	1	1	0	8	0	1	0	1	1		
14	19	1	1	0	1	9	0	0	1	0	1		
15	20	1	1	1	1	10	0	0	0	1	1		
16	21	1	0	1	1	11	0	0	1	0	1		
17	22	1	0	1	0	12	0	1	0	0	1		
18	23	1	1	1	0	13	0	1	0	1	1		
19	24	1	1	0	1	14	0	0	1	0	1		
20	25	1	1	1	1	15	0	0	0	1	1		

Table 2

3. (7, 15)

(11, 15)

Table 3

Step 3: Now compare all the minterms of table 3 with those in the adjacent groups.

As shown in table 4

Group P minutiae	In Binary				P Job 5	E minutiae	In Binary			
	A	B	C	D			A	B	C	D
0 (0,1)	0	0	0	0	10	0,1,8,9	-	0	0	-
(0,8)	-	0	0	0	10	0,8,1,9	-	0	0	-
1 (1,3)	0	0	1	1	12	1,3,9,11	-	0	-	1
(1,9)	5	0	0	1	12	1,4,3,11	-	0	-	1
(8,9)	1	0	0	-	13	3,7,11,15	-	0	-	1
-	0	0	-	-	13	3,11,7,15	-	0	-	1
2 (3,7)	0	0	1	1	13	3,11,7,15	-	0	-	1
(3,11)	-	0	1	1	13	3,11,7,15	-	0	-	1
(9,11)	+	0	0	1	13	3,11,7,15	-	0	-	1
-	0	-	-	-	13	3,11,7,15	-	0	-	1
3 (7,15)	-	1	1	1	13	3,11,7,15	-	0	-	1
(11,15)	!	-	1	1	13	3,11,7,15	-	0	-	1

Table 3

Table 3

Step 4: Repeat the procedure till grouping if can group the quads of minterms in the adjacent groups of table to obtain groups of eight minterms. These are no such matching.

- check again for matched pairs if not exist stop
- o o - P, E, S

Now prepare prime implicant table as shown in table 5.

Group	Minterms	Binary SOP [$I - A \bar{B} \bar{C}$]			
		A	B	C (P, I)	D
-	- II, P, E, I	-	0	0	-
1	$m_0 - m_1 - m_8 - m_9$	-	0	0	-
	$m_0 - m_8 - m_1 - m_9$	- 1	0	- 0 0	(P, E)
2	$m_1 - m_3 - m_9 - m_{11}$	- 1	0	0 -	(I, E)
	$m_1 - m_9 - m_3 - m_{11}$	-	0	-	1
3.	$m_3 - m_7 - m_{11} - m_{15}$	- 1	- 1	- 1 -	(E, I)
	$m_3 - m_{11} - m_7 - m_{15}$	- 1	- 1	- 1 -	(E, I)

Table 5

$$\bar{B}\bar{C} + \bar{B}D + CD - \text{prime implicant}$$

Now we have to find essential

Step 5: Find essential prime implicants.

PI prime implicant	Minterms Groups and Boolean Representation	Given minterms
		0 1 3 7 8 9 11 15
✓	$\bar{B}\bar{C}$ 0, 1, 8, 9	X X X X
✗	$\bar{B}D$ 1, 3, 9, 11	X X X X
✓	CD 3, 7, 11, 15	X X X X

Rounding which have only one value (x) Table 6 Here $\bar{B}D$ does not included in rounding so it is not PI

From table 6 prime implicants are

$\bar{B}\bar{C}$ and CD

$\bar{B}D$ is

not

Required output

$$Y = \bar{B}\bar{C} + CD$$

Implicant

no logical error in this.

Final output

given logic is same as required to draw.

Final output

most important part of this logic is this.

Final output

Final output