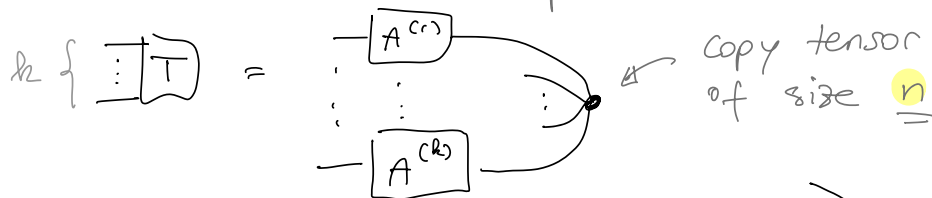


$T \in \mathbb{C}^{d_1} \otimes \dots \otimes \mathbb{C}^{d_k}$ k -partite tensor

$$T = \sum_{i=1}^n \underbrace{a_i^{(1)} \otimes a_i^{(2)} \otimes \dots \otimes a_i^{(k)}}_{\text{simple tensors}}$$



$$\left. \begin{array}{c} | \\ | \\ | \\ | \\ | \end{array} \right\} = \sum_{i=1}^n \underbrace{|i i \dots i\rangle}_{k \text{ times}}$$

The **tensor rank** of T is the **smallest** n for which such a decomposition exists. **$R(T)$**

Example In QIT $|\psi\rangle \in \mathbb{C}^{d_1} \otimes \dots \otimes \mathbb{C}^{d_k}$
a pure quantum state

$$R(|\psi\rangle) = 1 \Leftrightarrow |\psi\rangle = |\psi_1\rangle \otimes \dots \otimes |\psi_k\rangle$$

i.e. ψ separable (or product)

$\Rightarrow R$ is an integer-valued entanglement measure for pure multipartite q. states.

\rightarrow Note that R is invariant under local unitary operations. (LU)

$$R((U_1 \otimes \dots \otimes U_k) \cdot T) = R(T)$$

$$R \left[\begin{array}{c} \boxed{u_1} \\ \boxed{u_2} \\ \vdots \\ \boxed{u_n} \end{array} \rightarrow T \right] = R \left[\begin{array}{c} \text{---} \\ \text{---} \\ \vdots \\ \text{---} \end{array} \rightarrow T \right]$$

Proof if $r = R(T) \Rightarrow T = \sum_{i=1}^r \bigotimes_{s=1}^h a_i^{(s)}$

$$(U_1 \otimes \dots \otimes U_n) T = \sum_{i=1}^r \bigotimes_{s=1}^h (U_s a_i^{(s)})$$

$$\left[\begin{array}{c} \vdots \\ T \end{array} \right] = \left[\begin{array}{c} A^{(0)} \\ \vdots \\ A^{(s)} \end{array} \right] \Rightarrow$$

$$\left[\begin{array}{c} (u_1) \\ \vdots \\ (u_n) \end{array} \right] T = \left[\begin{array}{c} u_1 A^{(1)} \\ \vdots \\ u_n A^{(n)} \end{array} \right] T'$$

$$\Rightarrow R((U_1 \otimes \dots \otimes U_n) T) \leq R(T) = r$$

Also: $R(T = (U_1^{-1} \otimes \dots \otimes U_n^{-1}) \cdot T') \leq R(T')$

Actually $R(X_1 \otimes \dots \otimes X_n T) = R(T)$

if invertible X_s

if $T' = (X_1 \otimes \dots \otimes X_n) T$

for invertible X_s then T' can be

obtained from T by **SLOCC**
 (stochastic local operation and classical comm.)
 with non-zero probability.
 \Rightarrow tensor rank is a **SLOCC invariant**.

Example for matrices ($k=2$)

$$R(A) = \text{rk}(A) = \# \text{ of non-zero singular values.}$$

$$A = \sum_{i=1}^{\hat{n}} |a_i\rangle\langle b_i| \Leftrightarrow \sum_{i=1}^{\hat{n}} |a_i\rangle \otimes |b_i\rangle$$

Examples

$k=2$

- $|00\rangle \in \mathbb{C}^2 \otimes \mathbb{C}^2$

$$R(|00\rangle) = 1 \quad |0\rangle \otimes |0\rangle$$

$$\downarrow$$

$$\text{rank} \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix} = 1 \quad \leadsto \text{separable}$$

- $|S\rangle = \frac{1}{\sqrt{2}} (|00\rangle + |11\rangle)$

$$R(|S\rangle) = 2 \quad \Rightarrow \text{entangled}$$

$$\downarrow$$

$$\text{rank} \left(\frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \right) = \text{rank}(I_2) = 2$$

$k=3$

$$|\varphi\rangle \in \mathbb{C}^2 \otimes \mathbb{C}^2 \otimes \mathbb{C}^2 \quad 3 \text{ qubits.}$$

[Dür, Vidal, Cirac '2000] can be entangled in 2 different ways

- $|000\rangle = |0\rangle \otimes |0\rangle \otimes |0\rangle$

$$R(|000\rangle) = 1 \Rightarrow \text{separable}$$

- $\frac{1}{\sqrt{2}}(|000\rangle + |011\rangle) = |0\rangle_A \otimes |\Sigma\rangle_{BC}$
 $R(\uparrow) = 2$ entangled but bi-separable $A|BC$

- $|W\rangle = \frac{1}{\sqrt{3}}(|100\rangle + |010\rangle + |001\rangle)$

- $|GHZ\rangle = \frac{1}{\sqrt{2}}(|000\rangle + |111\rangle)$

$$R(|W\rangle) = 3 \text{ and } R(|GHZ\rangle) = 2$$

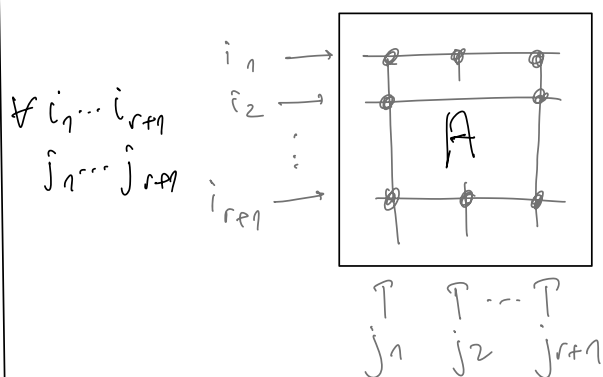
But : $|W\rangle$ and $|GHZ\rangle$ are entangled "differently"

Interlude For matrices, the sets

$$\text{Mat}_r := \{ A \in M_d(\mathbb{C}) : \text{rk}(A) \leq r \}$$

are closed

$$A \in \text{Mat}_r \Leftrightarrow \text{all of its } (r+1) \text{ minors are } 0$$



$$\det \left(A \begin{matrix} i_1 \dots i_{r+1} \\ j_1 \dots j_{r+1} \end{matrix} \right) = 0$$

\uparrow
 polynomial in the entries of A .

$A \in \text{Mat}_r \quad (\Rightarrow) \quad \left\{ \binom{d}{r+1}^2 \text{ polynomials are zero at } A \right\}$

$\Rightarrow \text{Mat}_r$ is closed

For tensors, this is not true!

Example $R(|W\rangle) = 3$, but it is the limit

$$|W\rangle = \lim_{n \rightarrow \infty} T_n \quad \text{with} \quad R(T_n) = 2$$

$(\Rightarrow) \{T \in (\mathbb{C}^2)^{\otimes 3} : R(T) \leq 2\}$ is not closed

$$|100\rangle + |010\rangle + |001\rangle = \lim_{\varepsilon \rightarrow 0} \frac{(|0\rangle + \varepsilon|1\rangle)^{\otimes 3} - |000\rangle}{\varepsilon}$$

$\underbrace{\hspace{10em}}_{R(\cdot) = 2}$

$$(|0\rangle + \varepsilon|1\rangle)^{\otimes 3} = |000\rangle + \varepsilon \cdot |W\rangle + \varepsilon^2(\dots) + \varepsilon^3|111\rangle$$

Border rank

$$\underline{R} = \inf \{ r : \exists T_n \rightarrow T \text{ with } R(T_n) \leq r \}$$

$$\underline{R}(|W\rangle) = 2 \quad \text{vs} \quad R(|W\rangle) = 3$$

Theorem $\{T : \underline{R}(T) \leq r\}$ is the zero set of some finite family of polynomials

In $\mathbb{C}^2 \otimes \mathbb{C}^2 \otimes \mathbb{C}^2$ there is only one such poly.

3-tangle $\tau_3 = 4 \cdot \text{Cayley hyperdet.}$
(1850's)

$\tau_3(14\rangle)$ is a poly of degree 4 in
the 8 variables.

- $\tau_3(W) = 0$.
- $\tau_3(\text{GHZ}) > 0$.
- τ_3 is a SLOCC invariant.

In particular:

$$\begin{array}{ccc} \text{GHZ} & \xrightarrow{\text{SLOCC}} & W \\ W & \xrightarrow{\text{SLOCC}} & \text{GHZ} \end{array}$$

In some sense, GHZ is more
entangled than W.

Fact deciding $R(T) \leq r$ is NP-hard
(for $k \geq 3$)

3SAT: $\varphi = (x_1 \text{ OR } \bar{x}_2 \text{ OR } x_5) \text{ AND } (x_2 \text{ OR } x_3 \text{ OR } \bar{x}_1)$
 $\text{AND } (\dots)$

Given such a statement, is it satisfiable?

To any 3SAT problem $\varphi \mapsto T_\varphi \in \mathbb{C}^{\otimes \dots \otimes \mathbb{C}}$
s.t. φ satisfiable $\Leftrightarrow R(T_\varphi) \leq \dots$
where all the \dots are polynomial in size of φ

The matrix multiplication tensor

$$\overset{m \times n}{(A, B)} \overset{n \times p}{\longrightarrow} \overset{m \times p}{C = A \cdot B}$$

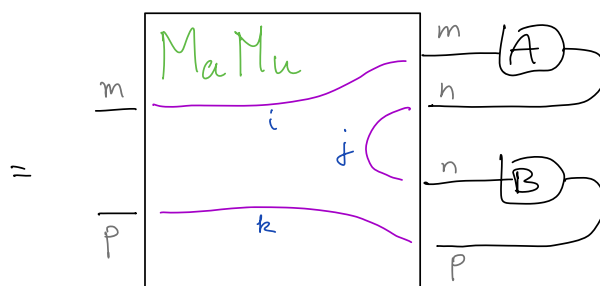
is a bilinear operation

$$C_{ik} = \sum_{j=1}^n A_{ij} B_{jk} \quad \begin{matrix} i \in [m] \\ k \in [p] \end{matrix}$$

$$\text{---} \boxed{C} \text{---} = \text{---} \boxed{A} \text{---} \boxed{B} \text{---}$$

In vectorized form:

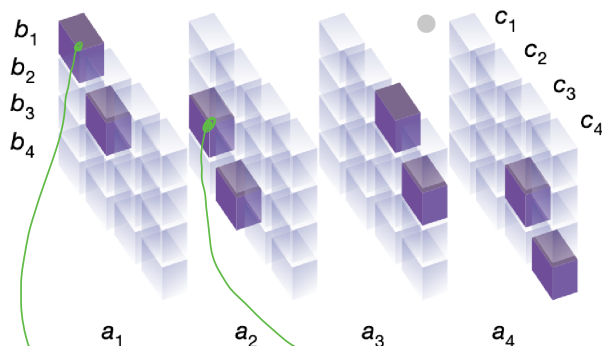
$$\text{---} \boxed{C} \text{---} = \text{---} \boxed{A} \text{---} \boxed{B} \text{---}$$



$$MaMu_{m,n,p} = \sum_{i=1}^m \sum_{j=1}^n \sum_{k=1}^p \underbrace{|j\rangle\langle i|}_{\text{"A"}} \otimes \underbrace{|k\rangle\langle j|}_{\text{"B"}} \otimes \underbrace{|i\rangle\langle k|}_{\text{"C"}}$$



$$\begin{pmatrix} c_1 & c_2 \\ c_3 & c_4 \end{pmatrix} = \begin{pmatrix} a_1 & a_2 \\ a_3 & a_4 \end{pmatrix} \cdot \begin{pmatrix} b_1 & b_2 \\ b_3 & b_4 \end{pmatrix}$$



$$c_1 = a_1 b_1 + a_2 b_3$$

Claim

$R(\text{MatMu}) = \text{minimal number of multiplications needed to compute the matrix product.}$

$$\Rightarrow R(\text{MatMu}_{2,2,2}) \leq 8$$

In general

$$\text{MatMu}_{2,2,2} = \sum_{i=1}^r \underbrace{\alpha_i}_{\mathbb{C}^4} \otimes \underbrace{\beta_i}_{\mathbb{C}^4} \otimes \underbrace{\gamma_i}_{\mathbb{C}^4}$$

$$\text{MatMu}(i,j,k) = 1 \Rightarrow \boxed{\text{shaded}} \\ = 0 \Rightarrow \boxed{\text{empty}}$$

$$= \text{MatMu}_{2,2,2}$$

$$\uparrow \\ \mathcal{M}_2 \otimes \mathcal{A}_2 \otimes \mathcal{A}_2 \\ \text{"A" } \quad \text{"B" } \quad \text{"C" }$$

$$\downarrow \\ \mathbb{C}^4 \otimes \mathbb{C}^4 \otimes \mathbb{C}^4$$

Claim : $R(\text{MatMu}_{2,2,2}) = 7$

Strassen '69

$$m_1 = (a_1 + a_4)(b_1 + b_4)$$

$$m_2 = (a_3 + a_4)b_1$$

$$m_3 = a_1(b_2 - b_4)$$

$$m_4 = a_4(b_3 - b_1)$$

$$m_5 = (a_1 + a_2)b_4$$

$$m_6 = (a_3 - a_1)(b_1 + b_2)$$

$$m_7 = (a_2 - a_4)(b_3 + b_4)$$

$$c_1 = m_1 + m_4 - m_5 + m_7$$

$$c_2 = m_3 + m_5$$

$$c_3 = m_2 + m_4$$

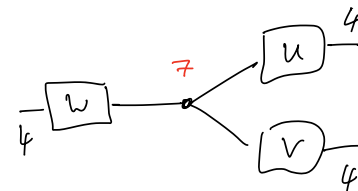
$$c_4 = m_1 - m_2 + m_3 + m_6$$

$$u = \begin{pmatrix} 1 & 0 & 1 & 0 & 1 & -1 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 & 1 & 0 \\ 1 & 1 & 0 & 1 & 0 & 0 & -1 \end{pmatrix}$$

$$v = \begin{pmatrix} 1 & 1 & 0 & -1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 1 \\ 1 & 0 & -1 & 0 & 1 & 0 & 1 \end{pmatrix}$$

$$w = \begin{pmatrix} 1 & 0 & 0 & 1 & -1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 & 0 & 0 \\ 1 & -1 & 1 & 0 & 0 & 1 & 0 \end{pmatrix}$$

$$\begin{matrix} 4 & & 4 \\ \text{---} & \boxed{\text{MatMu}} & \text{---} \\ & 4 & \end{matrix} =$$



Using this, the cost of multiplying matrices
of size d goes from $d^3 \longrightarrow d^{\log_2 7} \approx d^{2.808}$

Open problem Can this be $d^{2 + o(1)}$