

Dr Greg Wadley



# INFO90002

## Database Systems & Information Modelling

Week 05  
Normalisation

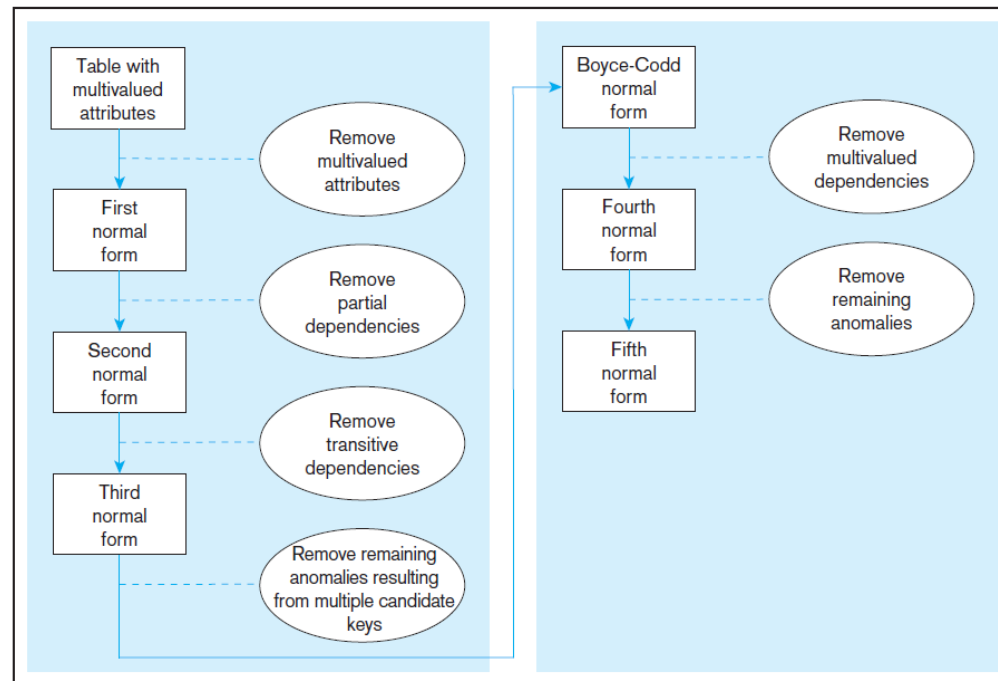


- Normalisation
  - 1<sup>st</sup> Normal Form (1NF)
  - 2<sup>nd</sup> Normal Form (2NF)
  - 3rd Normal Form (3NF)
  - Boyce Codd Normal Form (BCNF)

# Normalisation is...

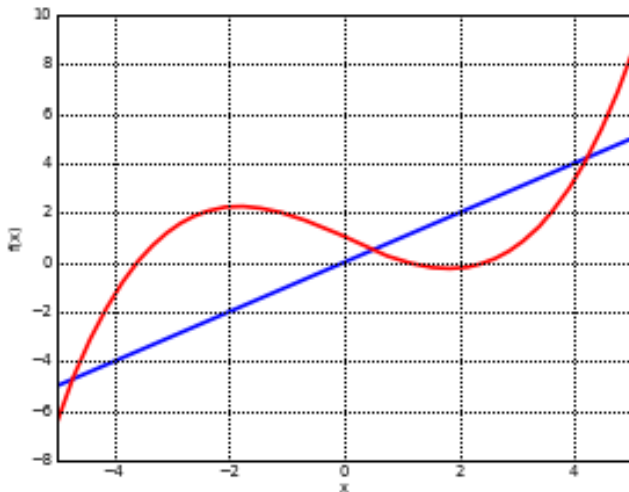
- A theoretical foundation for the relational model
- Application of a series of rules that improve db design
- Last phase of logical design
- Normalisation rules are designed to
  - Prevent data redundancy
  - Prevent update anomalies
- Normalisation is biased towards optimizing *Updates*
  - i.e. it assumes non-key columns are updated frequently
  - thus it can slow down *Selects* because of the need to join tables
  - thus designers sometimes “de-normalise” some tables to improve performance
- Every attribute must be functionally dependent on  
“The key, the whole key, and nothing but the key.”

- Normal Form
  - the state of a relation resulting from applying rules about the functional dependency of some attributes upon others
- Design progresses in stages from 1NF via 3NF, possibly as far as 6NF
  - in this course we only cover 1NF to 3.5NF (Boyce-Codd NF)
  - this is often sufficient in practice



flowchart from  
Hoffer et al.  
*Modern Database  
Management*

- Functional dependency
  - a set of attributes  $X$  *determines* another set of attributes  $Y$  iff each value of  $X$  is associated with only one value of  $Y$
  - written  $X \rightarrow Y$  ( similar to  $y = f(x)$  )
- Determinants
  - the attribute(s) on the *left hand side* of the arrow
- Key and Non-Key attributes
  - each attribute is either part of the primary key or it is not
- Partial functional dependency
  - a functional dependency of one or more non-key attributes upon *part (but not all)* of the primary key
- Transitive dependency
  - a functional dependency between 2 (or more) non-key attributes



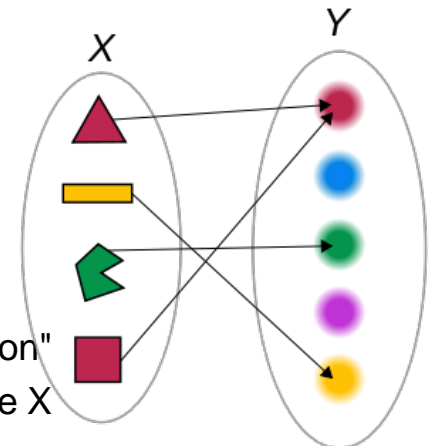
$$y = f(x)$$

for each  $x$  there is no more than one  $y$

CustId	Firstname	Lastname
1	Fred	Brown
2	Fred	Smith
3	Julia	Brown
4	Tony	Abbott
5	Julia	Gillard



CustId determines Firstname and Lastname  
neither Firstname nor Lastname determine anything



the "colour-of-the-shape function"  
 $X$  determines  $Y$ , but  $Y$  does not determine  $X$

- A relationship between the attributes of one relation
  - one or more attributes determine the value of another attribute
- A primary key must functionally determine all non-key attributes of an entity
  - stock exchange company code  $\rightarrow$  company name and details
    - If we know the company code, we know the value of company name
    - (<http://www.asx.com.au/prices/understanding-asx-ticker-codes.htm>)
  - studentId, subjectCode  $\rightarrow$  dateCompleted, resultObtained
    - for any combination of studentId and subjectCode, there is only one dateCompleted and one resultObtained
- Determinant
  - only primary keys should be determinants

- 1NF
  - Any *multivalued* attributes and repeating groups have been removed. There is a SINGLE value (possibly null) at the intersection of each row and column of the table.
- 2NF
  - Any *partial* functional dependencies have been removed (ie all non-key attributes are identified by the WHOLE key)
- 3NF
  - Any *transitive* dependencies have been removed (ie non-key attributes are identified by ONLY the PRIMARY key)
- BCNF
  - Any remaining anomalies that result from functional dependencies have been removed (ie because there was more than one candidate primary key for the same non-keys)



- To begin:
  - Write in relational notation  
List all the attributes that must be recorded, as one big relation
  - Don't include anything that you can derive (e.g. total sales)
  - Use parentheses “()” to indicate repeating groups
  - Underline the primary key
  - R1 (InvoiceNo, Date, CustomerNo, CustomerName, CustomerAddress, ClerkNo, ClerkName, (ProductNo, ProductDesc, UnitPrice, Qty))

- Separate repeating groups into new tables

- if an attribute is multi-valued,
- or a group of attributes is repeated several times for one entity,
- create a new table containing the repeating data

example: paper invoices typically include a repeating group which represents each product sold in that sale

**XinCube Inc**  
380 Francisco St  
San Francisco  
CA 94133 US  
Tel: (415) 989-1188 Fax: (415) 989-2288  
Email: admin@xincube.com  
Website: www.xincube.com

Invoice No. INV09080012

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**Invoice**

**Bill To**  
John  
Synex Inc  
128 AA Juanita Ave  
Glendora  
CA 91740 US

**Ship To**  
John  
Synex Inc  
128 AA Juanita Ave  
Glendora  
CA 91740 US

Date	14-Aug-2009	Order No		Sales Person	Charles Wooten
Shipping Date	13-Aug-2009	Shipping Terms		Terms	COD

ID	SKU / Description	Unit Price (USD)	Qty	Amount (USD)
PS.V880.005	AMD Athlon X2DC-7450, 2.4GHz/1GB/160GB/SMP-DVD/VB	580.00	6.00	3,480.00
PS.V880.037	PDC-E5300 - 2.6GHz/1GB/320GB/SMP-DVD/FDD/VB	645.00	4.00	2,580.00
LC.V890.002	LG 18.5" WLCD	230.00	10.00	2,300.00
HP.Q754.071	HP LaserJet 5200	1,103.00	1.00	1,103.00

Note  
All Payments must be made only in the form of a crossed cheque or cash payable to  
Xin Cube Inc

Sub Total (USD):	9,463.00
Discount (USD):	0.00
Sales Tax (USD):	780.70
Shipping (USD):	0.00
Total (USD):	10,243.70
Deposit (USD):	0.00

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Amount Due (USD): 10,243.70



## Our example (1NF)

- R1 (InvoiceNo, Date, CustomerNo, CustomerName, CustomerAddress, ClerkNo, ClerkName, (ProductNo, ProductDesc, UnitPrice, Qty))
- Repeating group is (ProductNo, ProductDesc, UnitPrice, Qty)
- The new table is as follows:
  - R2 (*InvoiceNo*, ProductNo, ProductDesc, UnitPrice, Qty)  
(foreign keys get italics or dotted underline)
- The repeating fields will be removed from the original relation, leaving the following.
  - R1 (InvoiceNo, Date, CustomerNo, CustomerName, CustomerAddress, ClerkNo, ClerkName)
- These two relations together comprise a database in 1NF

- Relevant only for relations whose primary key is a COMPOSITE key
- Remove partial functional dependencies.
  - an attribute is dependent on only *part of* the primary key
  - i.e. dependent on *one* of the columns in a composite primary key
- Create a separate table, containing the functionally dependent data, and the part of the key on which it depends.

## Our example (2NF)

- R1 (InvoiceNo, Date, CustomerNo, CustomerName, CustomerAddress, ClerkNo, ClerkName)
  - R2 (InvoiceNo, ProductNo, ProductDesc, UnitPrice, Qty)
- 
- R1 doesn't have a composite key, so it is already in 2NF
  - R2 is not in 2NF, because ProductDesc and UnitPrice are determined by ProductNo, which is *part of* the primary key. We need to create a new table from these fields
- 
- The new table will contain the following fields:
    - R3 (ProductNo, ProductDesc, UnitPrice)

## Our example (2NF)

- All of these fields except the primary key will be removed from the original table. The primary key will be left in the original table to allow linking of data:
  - R2 (InvoiceNo, ProductNo, Qty)
- Our database is now in second normal form:
  - R1 (InvoiceNo, Date, CustomerNo, CustomerName, CustomerAddress, ClerkNo, ClerkName)
  - R2 (InvoiceNo, ProductNo, Qty)
  - R3 (ProductNo, ProductDesc, UnitPrice)

- Remove *transitive dependencies*
  - a functional dependency where one non-key attribute is functionally dependent on another non-key attribute
  - thus its value is only indirectly determined by the primary key.
- Create a separate table containing the determinant attribute and the fields that are functionally dependent on it.
  - keep a copy of the key in the original table.

## Our example 3NF

- R1 (InvoiceNo, Date, CustomerNo, CustomerName, CustomerAddress, ClerkNo, ClerkName)
- R2 (InvoiceNo, ProductNo, Qty)
- R3 (ProductNo, ProductDesc, UnitPrice)
- R2 and R3 are in 3NF as there are not any transitive dependencies
  - All the attributes only depend on the key
- R1 is not in 3NF
  - Transitive dependencies exist
    - CustomerName and CustomerAddress depend on CustomerNo
    - ClerkName depends on ClerkNo

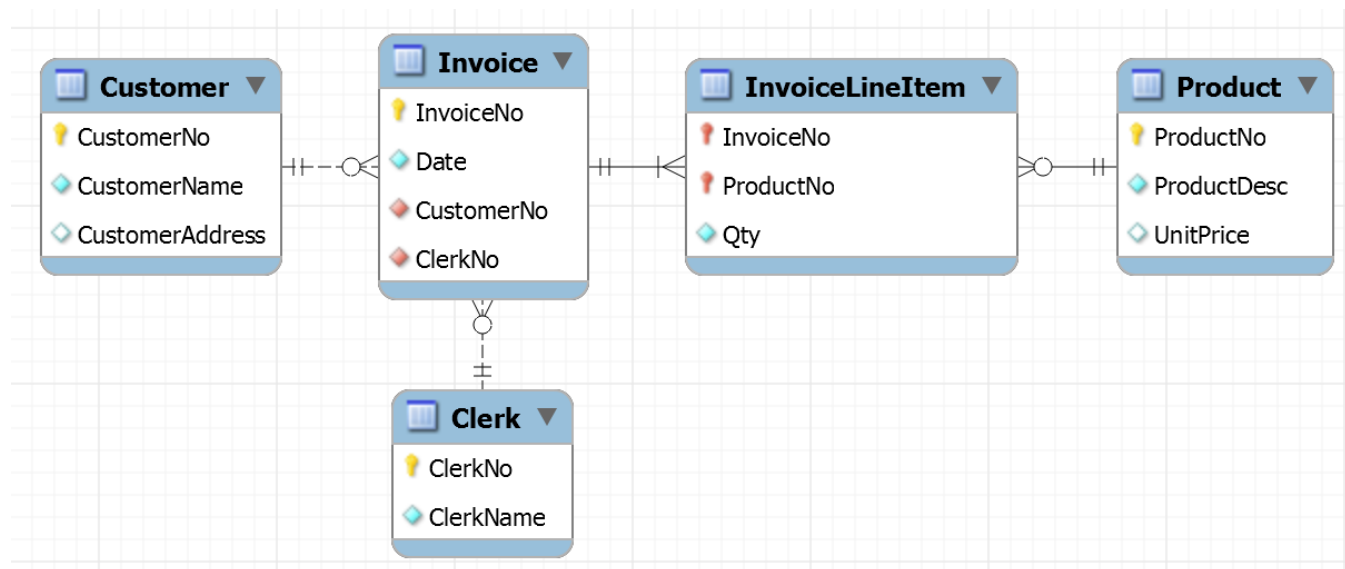


- The new tables will be:
  - R4 (CustomerNo, CustomerName, CustomerAddress)
  - R5 (ClerkNo, ClerkName)
- All of these fields except the primary keys will be removed from the original table. The PKs will be left in the original table (as FKs) to allow linking of data, as follows:
  - R1 (InvoiceNo, Date, *CustomerNo*, *ClerkNo*)
- Together with the unchanged tables below, these tables make up the database in third normal form.
  - R2 (InvoiceNo, ProductNo, Qty)
  - R3 (ProductNo, ProductDesc, UnitPrice)



- Repetition of data
  - Individual facts are stored many times.
- Update anomalies
  - To change a fact, we must change it many times.
- Deletion anomalies
  - Information about entity A is stored inside entity B.  
If we delete all B rows, we lose our record of A.
- Insertion anomalies
  - To record a new A, we must insert a B.

- We can name the relations now...
  - Customer (CustomerNo, CustomerName, CustomerAddress)
  - Clerk (ClerkNo, ClerkName)
  - Product (ProductNo, ProductDesc, UnitPrice)
  - Invoice (InvoiceNo, Date, *CustomerNo*, *ClerkNo*)
  - InvoiceLineItem (*InvoiceNo*, *ProductNo*, Qty)

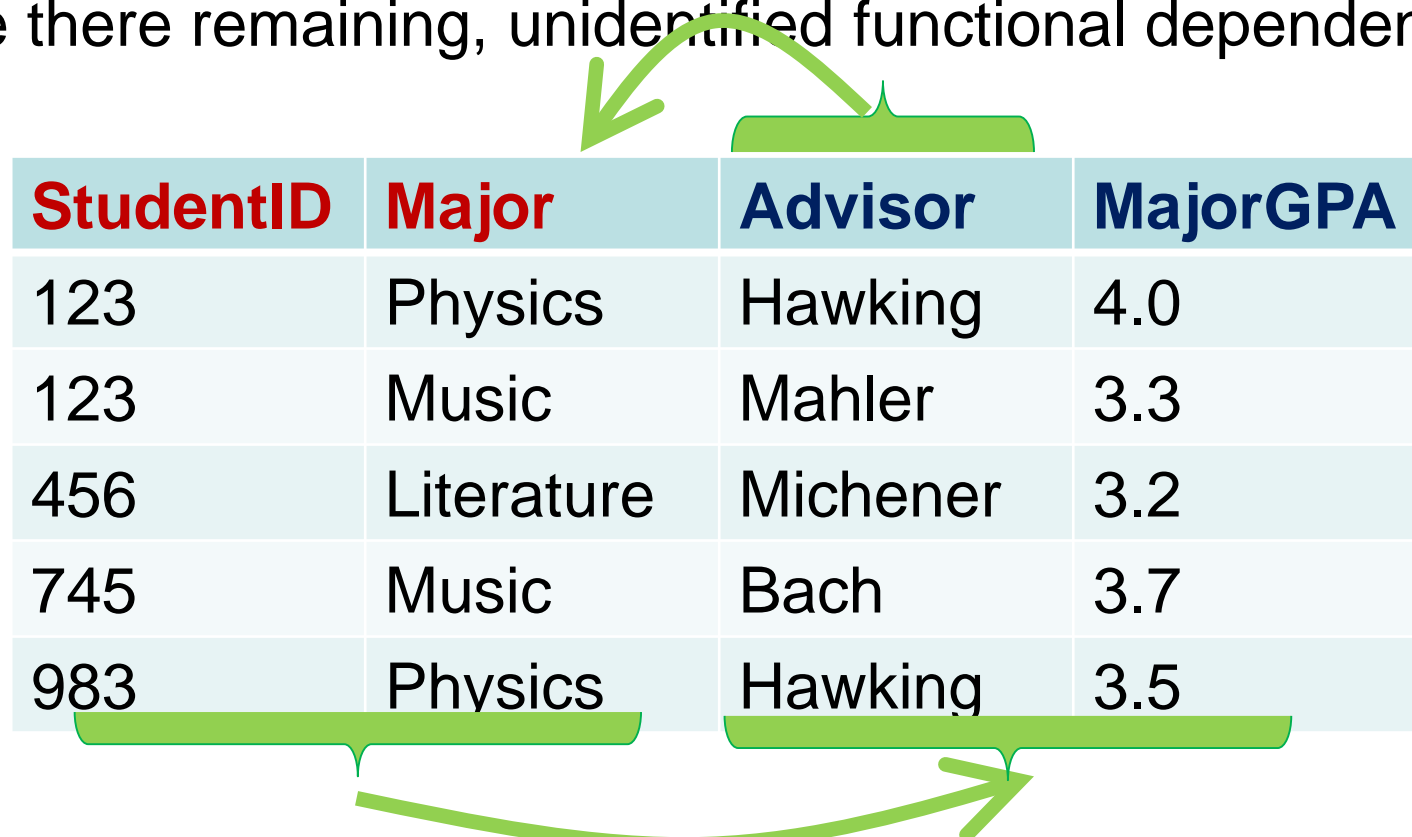




- It is often sufficient to stop normalizing at 3NF
- BCNF, 4NF, 5NF and even 6NF have been defined
  - we'll touch on BCNF
  - it is suggested that you read about the others, see
  - [http://en.wikipedia.org/wiki/Database\\_normalization#Normal\\_forms](http://en.wikipedia.org/wiki/Database_normalization#Normal_forms)
  - but we will not assess beyond 3NF

- Can arise when ...
  - not every determinant in the relation is a candidate key
- A simple test is to check whether ...
  - "Every non-key attribute must provide a fact about the key, the whole key, and nothing but the key."

- Consider the following relation
  - Student (StudentID, Major, Advisor, MajorGPA)
- It is in 3NF ... but ...
- Are there remaining, unidentified functional dependencies?



StudentID	Major	Advisor	MajorGPA
123	Physics	Hawking	4.0
123	Music	Mahler	3.3
456	Literature	Michener	3.2
745	Music	Bach	3.7
983	Physics	Hawking	3.5

- Update anomaly
  - Hawking is replaced by Einstein
- Insertion anomaly
  - Babbage advises comp. sci., but has no students
- Deletion anomaly
  - e.g. delete the only student with a particular advisor (456)

StudentID	Major	Advisor	MajorGPA
123	Physics	Hawking	4.0
123	Music	Mahler	3.3
456	Literature	Michener	3.2
745	Music	Bach	3.7
983	Physics	Hawking	3.5

- Modify the relation so that
  - the determinant in the relation that is not part of the key becomes a component of the primary key of the revised relation
  - the attribute that is functionally dependant on that determinant becomes a non key attribute
- The relation becomes
  - Student (StudentID, Advisor, Major, MajorGPA)
- Can this be normalized further?
  - Its not in 2NF (partial functional dependency) – so normalise it
  - Student (StudentID, Advisor, MajorGPA)
  - Advisor (Advisor, Major)



- Student (StudentID, Advisor, MajorGPA)
- Advisor (Advisor, Major)

STUDENT

<u>SID</u>	<u>Advisor</u>	MajGPA
123	Hawking	4.0
123	Mahler	3.3
456	Michener	3.2
789	Bach	3.7
678	Hawking	3.5

ADVISOR

<u>Advisor</u>	Major
Hawking	Physics
Mahler	Music
Michener	Literature
Bach	Music



# Normalisation Exercise

Now it is your turn!

The Melbourne Multi-Campus School

**STUDENT RECORD**

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Student Id: 123456

Name: Joe Bloggs

Date of Birth: 1<sup>st</sup> April 1990

Campus: Fitzroy

Campus Address: 123 Smith St, Fitzroy 3065

Subject Code	Subject Name	YearTaken	Result
MAT	Maths	2015	
PHY	Physics	2015	
ART	Art	2014	60
BIO	Biology	2014	90

Graduation Date: \_\_\_\_\_

- The *Melbourne Multi-Campus School* has been using a paper-based information system.
- Student data is recorded on these sheets of paper.
- Now they are going to computerize.
- How shall we store the data in a relational database?