

BRFS Specification

Bruno Filesystem (formerly BOOT-ROOT)

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Abstract

This specification document describes the BRFS filesystem structure used to store data on storage devices. This provides a standard common description of the filesystem for developers to implement freely.

Revision History

Revision	Date	Author(s)	Description
0.1		bruneo32	Created
0.2		bruneo32	Unknown
0.3		bruneo32, arf20	This document

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1 Introduction

1.1 Scope

This document defines the Bruno Filesystem. As a filesystem it provides a way of structuring data in a block-based (i.e. LBA) storage device. It is meant for embedded systems where a complex filesystem is not needed, this is not a replacement for any modern desktop filesystem such as ext4, because it lacks basic features of journaling. Although BRFS is able to address large volumes, it is not recommended.

1.2 Definitions

Key words will be referred to with a `monospace font`.

- block: Minimum filesystem unit of data
- unspecified: May be implementation dependent

1.3 Advantages and disadvantages

Advantages	Disadvantages
TODO when defined	

1.4 Volume layout

Superblock
Root directory
<other files>
Free space

2 Superblock

The superblock records properties of the enclosed filesystem, such as the block size, pointer size and attribute size. It is 1 block in size. The remaining block will be padded with zeroes.

2.1 Superblock layout

Size (bytes)	Field (Unsigned integer)	Value
4	Magic number	"BRFS" 0x42524653
1	Block size in LBAs	0-255 (+1)
1	Pointer size in bytes	2, 4, or 8
Pointer	Free space pointer	
Pointer	First free block	
	Padding...	0x00

2.2 Theoretical limits

Property	Limit
Block size	256
Pointer size	64
Attribute size	256
Addressable blocks	2^{64}
Addressable LBAs	$256 \cdot 2^{64}$
Absolute maximum capacity (512-byte LBA)	$512 \cdot 256 \cdot 2^{64} \approx 2 \text{ YiB}$

The maximum capacity of the filesystem is calculated as follows

$$C = L \cdot B \cdot 2^p \quad (1)$$

Where p is pointer size, B is block size and L is LBA size.

Some examples of reasonable configurations (assuming 512-byte LBA) are $p = 32$, $B = 8$, which gives 16 TiB capacity; or for more efficient storage, $p = 64$, $B = 1$: 8 ZiB; for embedded systems perhaps only a $p = 16$ $B = 1$ is needed, for 32 MiB.

3 File

BRFS is a file based filesystem. Regular files and directories are both files.

Data
[Padding]
Next block pointer

3.1 Next block pointer

In the end of each file's block, lies a pointer to the next block of the file. This pointer is a linear offset of blocks.

Pointer number 0 is reserved to denote EOF, and pointer 1 refers to next block. Pointer space starts with 2, the block after the first block of the root directory, which coincides with the global block offset.

	Block
0	Superblock
1	Root directory
2	First file
3	...

3.2 Directory

A directory is a special file that holds file entries one after another. There is no limit on the number of entries. The directory ends with a null character 0x00, following the string terminator of the last entry's file name.

file0
...
filen
null terminator 0x00
[Padding]
Next block pointer

3.2.1 Directory entry

Describes file entry on directory

Type (Size)	Field
8	File size
22	Attributes
Pointer	First block
unspecified	File name (null-terminated string)

3.2.2 Attributes

Following the POSIX.1-2017 standard, and inspired in some linux ext4 attributes. It takes 22 bytes. See linux kernel and man7.

Type (size)	Name	Description
uint16	mode	File type and mode
uint32	uid	User ID of owner
uint32	gid	Group ID of owner
uint32	ctime	Creation time
uint32	atime	Last access time
uint32	mtime	Last modification time