

# BRFS Specification

## Bruno Filesystem (formerly BOOT-ROOT)

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October 17, 2023

### Abstract

This specification document describes the BRFS filesystem structure used to store data on storage devices. This provides a standard common description of the filesystem for developers to implement freely.

## Revision History

Revision	Date	Author(s)	Description
0.1		bruneo32	Created
0.2		bruneo32	Unknown
0.3		bruneo32, arf20	This document

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# 1 Introduction

## 1.1 Scope

This document defines the Bruno Filesystem. As a filesystem it provides a way of structuring data in a block-based (i.e. LBA) storage device. It is meant for embedded systems where a complex filesystem is not needed, this is not a replacement for any modern desktop filesystem such as ext4, because it lacks basic features of journaling. Although BRFS is able to address large volumes, it is not recommended.

## 1.2 Definitions

Key words will be referred to with a `monospace font`.

- block: Minimum filesystem unit of data
- unspecified: May be implementation dependent

## 1.3 Advantages and disadvantages

Advantages	Disadvantages
TODO when defined	

## 1.4 Volume layout

Superblock
Root directory
<other files>
Free space

# 2 Superblock

The superblock records properties of the enclosed filesystem, such as the block size, pointer size and attribute size. It is 1 block in size. The remaining block will be padded with zeroes.

## 2.1 Superblock layout

Size (bytes)	Field (Unsigned integer)	Value
4	Magic number	"BRFS" 0x42524653
1	Block size in LBAs	0-255 (+1)
1	Pointer size in bytes	2, 4, or 8
Pointer	Available space (RENAME)	
Pointer	First free block	
b-6	Padding...	0x00

## 2.2 Theoretical limits

Property	Limit
Block size	256
Pointer size	64
Attribute size	256
Addressable blocks	$2^{64}$
Addressable LBAs	$256 \cdot 2^{64}$
Absolute maximum capacity (512-byte LBA)	$512 \cdot 256 \cdot 2^{64} \approx 2 \text{ YiB}$

The maximum capacity of the filesystem is calculated as follows

$$C = L \cdot B \cdot 2^p \quad (1)$$

Where  $p$  is pointer size,  $B$  is block size and  $L$  is LBA size.

Some examples of reasonable configurations (assuming 512-byte LBA) are  $p = 32$ ,  $B = 8$ , which gives 16 TiB capacity; or for more efficient storage,  $p = 64$ ,  $B = 1$ : 8 ZiB; for embedded systems perhaps only a  $p = 16$   $B = 1$  is needed, for 32 MiB.

### 3 File

BRFS is a file based filesystem. Regular files and directories are both files.

Data
[Padding]
Next block pointer

#### 3.1 Next block pointer

In the end of each file's block, lies a pointer to the next block of the file. This pointer is a linear offset of blocks.

Pointer number 0 is reserved to denote EOF, and pointer 1 refers to next block. Pointer space starts with 2, the block after the first block of the root directory, which coincides with the global block offset.

	Block
0	Superblock
1	Root directory
2	First file
3	...

#### 3.2 Directory

A directory is a special file that holds file entries. There is no limit on the number of entries.

file0
...
filen
[Padding]
Next block pointer

##### 3.2.1 Directory entry

Describes file entry on directory

Type (Size)	Field
8	file size
22	Attributes
Pointer	First block
unspecified	Filename (c-string)

##### 3.2.2 Attributes

Following the POSIX.1-2017 standard, and inspired in some linux ext4 attributes. It takes 22 bytes.

Value	Name
uint16	mode
uint32	uid
uint32	gid
uint32	ctime
uint32	atime
uint32	mtime