



UMD DATA605 - Big Data Systems

8.4: Map Reduce Algorithms

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MapReduce: Applications

- **Intuition**

- Break massive jobs into independent map tasks
- Aggregate via shuffle
- Combine in reduce

- **Major classes of applications**

- Text processing and search
 - E.g., tokenization, inverted index, log analysis
- Large data transforms
 - E.g., ETL pipelines, joins, global sort, deduplication across petabytes
- Data mining and machine learning
 - E.g., Counting co occurrences, feature extraction, *k*-means iterations
- Graph and link analysis
 - E.g., PageRank, connected components

- **Typical outputs**, e.g.,

- Counts
- Aggregates
- Reorganized datasets for downstream systems

Cost Measures for Distributed Algorithms

- What matter is the real dollar cost not just $O(\cdot)$
 - **Total cost** \approx CPU + storage + network
 - **Communication cost**
 - Total I/O across all processes, e.g., shuffling 1 TB
 - **Elapsed communication cost**
 - Max I/O along the critical path
 - **Elapsed computation cost**
 - Wall clock with p workers, sensitive to skew and stragglers
- **Dominant term heuristic**
 - *"If one cost dominates, ignore the others for first order reasoning"*
- **Practical note**
 - *"Adding more machines trades \$ for time and may not fix skew"*

Total Cost Model for MapReduce

- **Total cost**

$$C_{total} = C_{compute} + C_{io} + C_{network} + C_{storage}$$

- **Notation**

- $|I|$ input GB, $|S|$ shuffle GB, $|O|$ output GB
- p_m mappers, p_r reducers
- T_m map hours, T_r reduce hours
- $c_{vm} = \text{\$/VM hour}$, $c_{io} = \text{\$/GB I/O}$, $c_{sh} = \text{\$/GB shuffle}$

- **Compute**

$$C_{compute} = c_{vm} (p_m T_m + p_r T_r)$$

where T_m, T_r include skew and stragglers

- Skew effects = a heavy key or hotspot inflates the heaviest task and the critical path

- **I/O**

$$C_{io} = c_{io} (|I| + 2|S| + |O|)$$

- Shuffle often dominates when $\sum_i |S_i| \gg |I|, |O|$

Total Cost Model for MapReduce

- **Notation**

- $|I|$ input GB, $|S|$ shuffle GB, $|O|$ output GB
- $c_{sh} = \$/\text{GB shuffle}$ $c_{eg} = \$/\text{GB egress}$, $c_{st} = \$/\text{GB hour}$
- R HDFS replication

- **Network**

$$C_{network} = c_{sh}|S| + c_{eg}|O|_{eg}$$

where $|O|_{eg}$ is data leaving the provider

- **Storage**

$$C_{stor} = c_{st}R(|I| + |O|)$$

- Plug in variables and unit prices to get C_{total}
- Apply the dominant term heuristic to prioritize optimization
- Tuning levers
 - Use Combiners
 - Compression
 - Better partitioning
 - Early filtering to reduce $\sum_i |S_i|$ and stragglers

Inverted Index using MapReduce

- **Goal:** build a mapping from words to the list of documents they appear in
 - Map phase:
 - Input: (docID, content)
 - Emit: (word, docID) for each word in content
 - Reduce phase:
 - Input: (word, [docID_1, docID_2, ...])
 - Emit: (word, list of unique docIDs)
- Useful in search engines and information retrieval
- Requires tokenization and normalization of content
- Deduplication of document IDs in reducer

Join Operations using MapReduce

- Goal: join two datasets based on a common key
- Types of joins:
 - Inner join, left/right outer join, full outer join
- Map phase:
 - Tag records by source and emit key as join attribute
- Reduce phase:
 - Input: (key, [records from both sources])
 - Perform join logic in reducer
- Example: joining employee records with department data by deptID
- Optimization strategies:
 - Replicated join for small datasets
 - Reduce-side join for large datasets
- Careful data tagging and partitioning required

Sorting and Grouping in MapReduce

- Goal: organize data by keys or values for further analysis
- Map phase:
 - Emit data with key as sort/group criterion
- Shuffle and sort phase automatically sorts by key
- Reduce phase:
 - Receives sorted keys and can perform grouped aggregation
- Total sort:
 - Requires partitioning and sampling for global order
- Partial sort:
 - Sort within each partition
- Example: sort sales data by date or group by product ID
- Often used as a preprocessing step for reporting

Graph Processing with MapReduce

- Many graph algorithms are iterative
- PageRank:
 - Compute importance score of web pages
- Map phase:
 - Emit contributions to neighboring nodes
- Reduce phase:
 - Sum contributions to update PageRank score
- Requires multiple MapReduce rounds for convergence
- Graph is represented as adjacency lists
- Each iteration refines the scores
- Can be implemented using chained MapReduce jobs or systems like Apache Giraph

Statistical Aggregation and Log Analysis

- Statistical Aggregation:
 - Map: emit relevant quantities (e.g., value, 1 for count)
 - Reduce: compute sums, averages, variances
 - Example: (sensorID, temperature) → compute average temperature per sensor
- Log Analysis:
 - Map: parse logs, extract fields (timestamp, IP, status)
 - Emit: (key, value) for desired metrics
 - Reduce: aggregate (e.g., count errors, hits per IP)
 - Example: (status code, 1) → count occurrences
- Useful for monitoring, alerting, and trend analysis
- Can handle large-scale logs from distributed systems