Distributed Computing

16. GraphX

Graph-Based Analytics

- •When MapReduce showed up, several engines for efficient graph-based analytics were introduced
- -Many machine learning & network analysis applications
- -Exploit the same cluster, but have ad-hoc optimizations
- •GraphX brings the same concepts to Spark, in a quite small package
- -First version: 2,500 lines of code
- •Reference: paper at USENIX OSDI '14

Example: PageRank

- Evaluate how important a node is in a graph (centrality)
- Can we see which page is more important in the Web graph (pages and links)?
- Idea: a random walker starts from a random page and
 - With probability d=0.85, clicks a link at random
 - With probability (1-d) 0.15, goes to another random page
- The score of a page is the probability the random walker gets there

Single-Machine PageRank

- •rank = [1/n, 1/n, ...] # array of length n
- For m iterations:
 - $new_rank = [(1-d)/n, (1-d)/n, ...] # length n$
 - For each node x with t neighbors
 - For each neighbor y of x
 - •new_rank[y] += d * rank[x] / t
 - rank = new_rank

Pregel: Think Like a Vertex

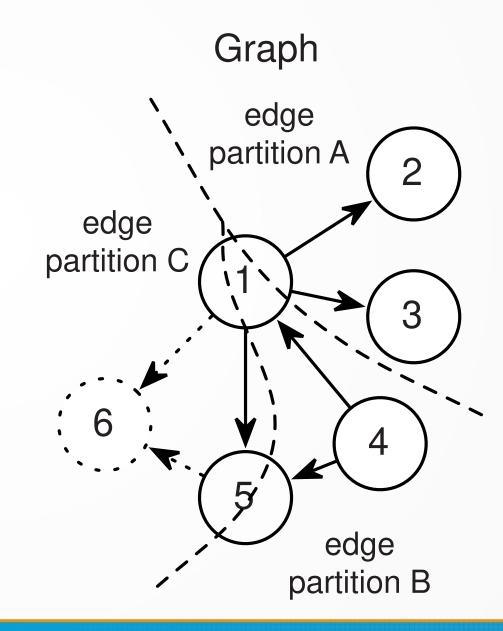
```
def PageRank(v: Id, msgs: List[Double]) {
  // Compute the message sum
  var msqSum = 0
  for (m <- msgs) { msgSum += m }
  // Update the PageRank
  PR(v) = 0.15 + 0.85 * msgSum
  // Broadcast messages with new PR
  for (j <- OutNbrs(v)) {
    msg = PR(v) / NumLinks(v)
    send_msg(to=j, msg)
  // Check for termination
  if (converged(PR(v))) voteToHalt(v)
```

GAS: Gather, Apply, Scatter

```
It's a message combiner! We can gather locally
                before sending informations around:)
def Gather(a: Double, b: Double) = a + b
def Apply(v, msgSum) {
  PR(v) = 0.15 + 0.85 * msqSum
  if (converged(PR(v))) voteToHalt(v)
def Scatter(v, j) = PR(v) / NumLinks(v)
```

Partitioning

- *The largest thing in a graph is the edge list
- *Hence, divide by edges
- -Some nodes will be "replicated" in different partitions (mirror nodes)
- Partition to minimize the number of them
- Copying a node's data (e.g., their PageRank value) is cheaper than sending messages along edges in different machines



Inactive Nodes

- •In many iterative algorithms some nodes "turn off"
- •E.g., a node already reached convergence
- •We can mark that to stop sending updates to them and accelerate iterations

