

# COMPUTER NETWORKS

## OSI Model

### PHYSICAL LAYER

→ provides functionality to data link layer.

### physical layers

### PBNTSPAN

Network  
Security

Application

### Datalink

### Network

### Transport

### Presentation

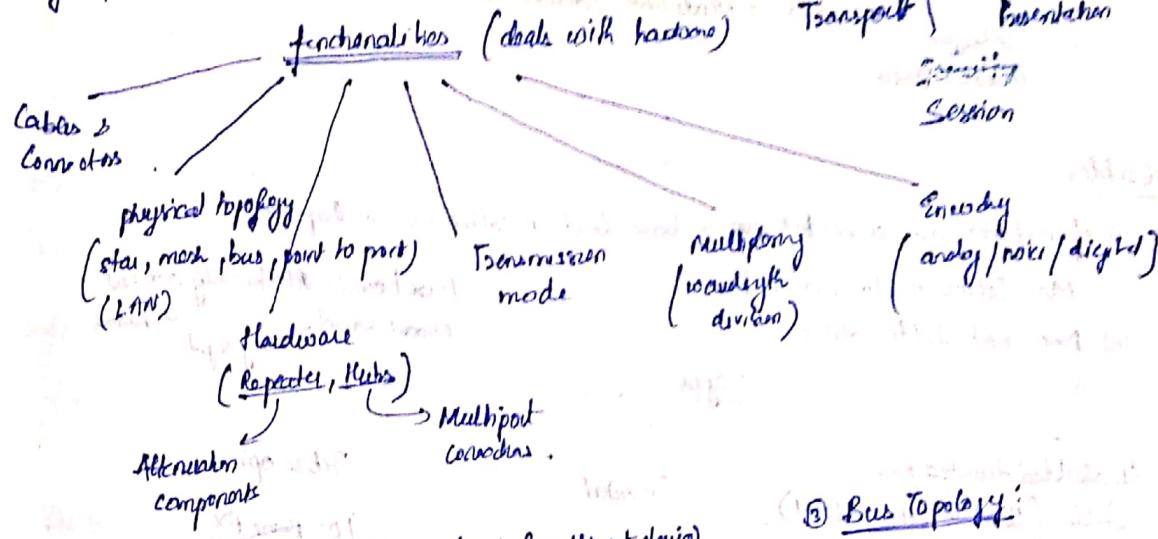
### Session

### Transport

### Session

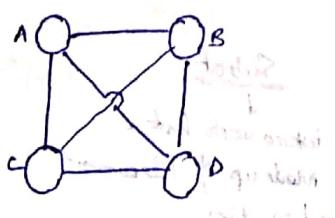
### Transport

### Session

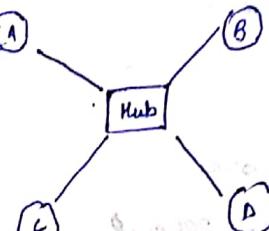


### ① Mesh

→ Security is high

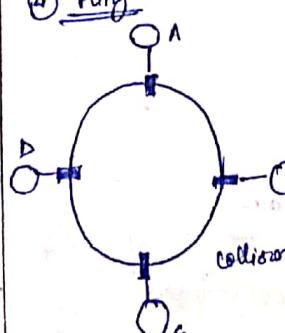


### ② Star/Hub (Multipoint device)



→ No Reliability  
No security.

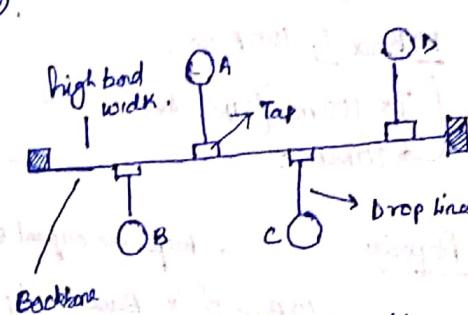
### ③ Ring



→ Reliability and security is low

→ Unidirectional only.

→ Taken to Reduce collisions.



→ No Reliability is checked by  
Single point of failure = 0

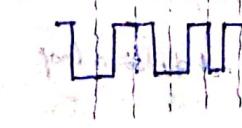
→ Security is less.

(high bandwidth : range of frequencies)  
with a given band, i.e. particular  
that is used for transmission

### Data Encoding

#### ① Manchester Encoding

• 1 → 0 1 0 1 0 1 ...

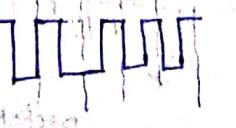


0 → 1 or 1 → 0

### Manchester Encoding

#### ② Differential Manchester Decoding

• 1 0 1 0 1 0 1 ...



0 → 1 or 1 → 0

## Various Networking Devices in Computer Networks

- 1) Cables
  - 2) Hub
  - 3) Switches
  - 4) Gateway
  - 5) Firewall
  - 6) Repeater
  - 7) Bridges
  - 8) Routers
  - 9) IDS
  - 10) Modem.
- Hub : Security.
- Intranet detection services
- Modulator / demodulator.

### Cables

(i) if  $n$  devices are connected over a base band,  $n$  collisions can happen.

Max Collision domain :  $n$  devices connected to a single bus.

(ii) Does not filter data

Base band - At a time only one signal

Broad band - Many signals at a time.

#### Types

Unshielded twisted pair cable (Used in Ethernet LAN).

Coaxial

Fibre optics

10 Base T, 100 Base T

10 Base 2

100 Base FX

Ts 100 mts (attenuation)

10 Base 3

(2 km Attenuation)

10 Mbps.

### Repeater

keeps the original strength source

10 Base 2 : Base band, 10Mbps, 200m, 0.

### Subnet

interconnect host.  
Made up of telecommuni-  
cation lines.

(i) Forwarding

(ii) No filtering

(iii) Collision domain's n (inside the device).

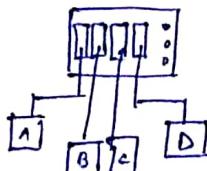
### Hub (source LAN)

(i) Multicast Repeater

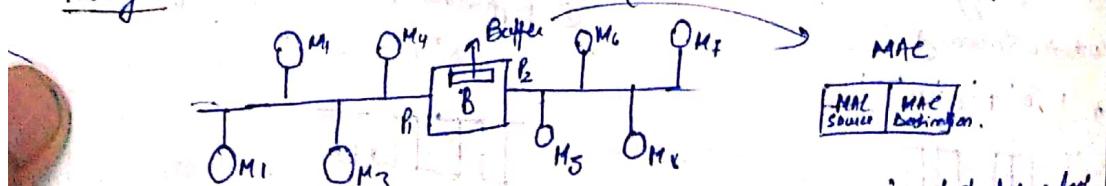
(ii) forwarder and broadcaster

(iii) No filtering

(iv) Collision domain's n.



Bridges : to connect two different LANs. (physical and Data link layer).



(i) Forwarding : if bridge has a packet, it forwards it.

(ii) filtering : filters for two different LAN's.

(iii) Collision domain : collision not occur, because bridge has a store and forward strategy. Bridge maintains buffer.

(iv) Use data Unit protocol : They make a Minimum spanning tree, so that path

## Type of addressing

Static (Entered manually)

MAC	Port
M <sub>1</sub>	P <sub>1</sub>
M <sub>2</sub>	P <sub>1</sub>
M <sub>3</sub>	P <sub>2</sub>
M <sub>4</sub>	P <sub>2</sub>

Dynamic or Transparent

→ learn themselves

→ Store addresses by broadcasting things first time.

→ The machine's have to acknowledge message passing and receiving.

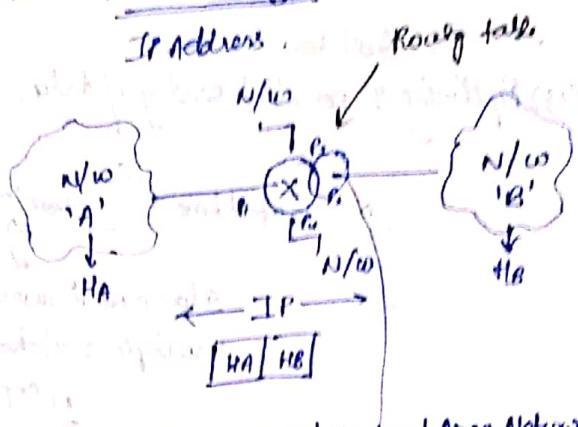
Router (Internet) (connects Physical, Data link, Network layer)  
(WAN) - Wide Area Network.

→ Forwarding: Router using Routing table

→ flooding: To broadcast the package in all directions

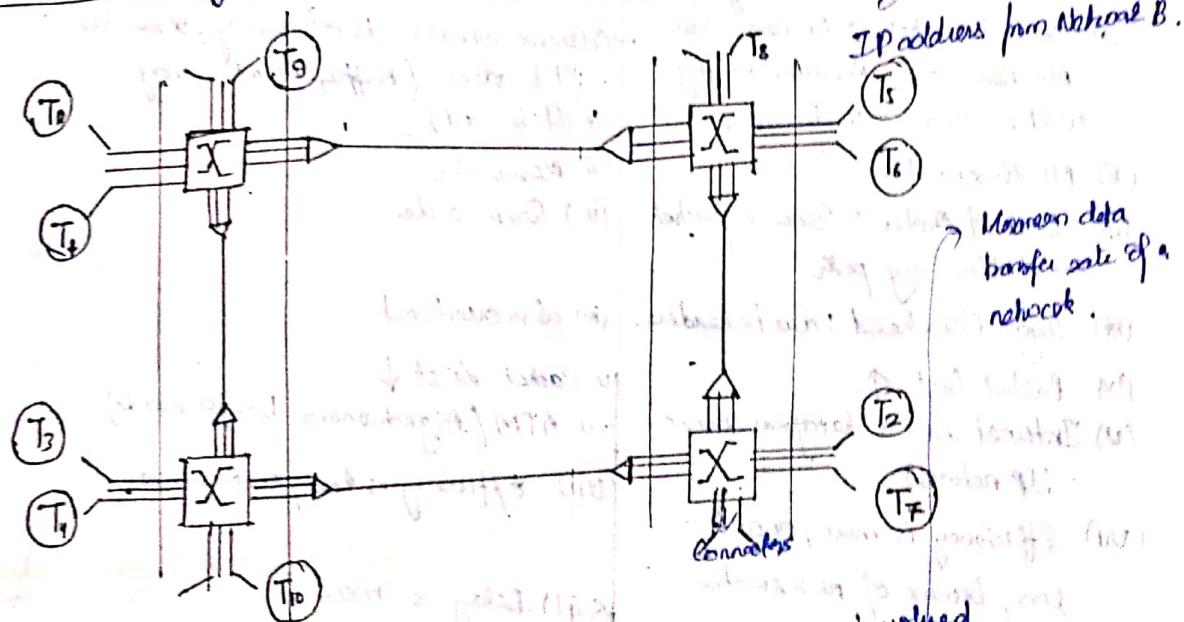
→ Filtering: Using its Routing table.

→ Collision Domain: No, star and forward method is used.



• MAC address is used in Local Area Network  
• ARP: Used to Request MAC address

## Circuit Switching



Total (i) Physical Layer: No other Layer's involved.  
Circuit switching used a dedicated path to communicate.

(ii) Continuous flow: data flows continuously.

(iii) No headers: No headers like MAC address, IP address.

→ Efficiency: Low, due to resource fixing.

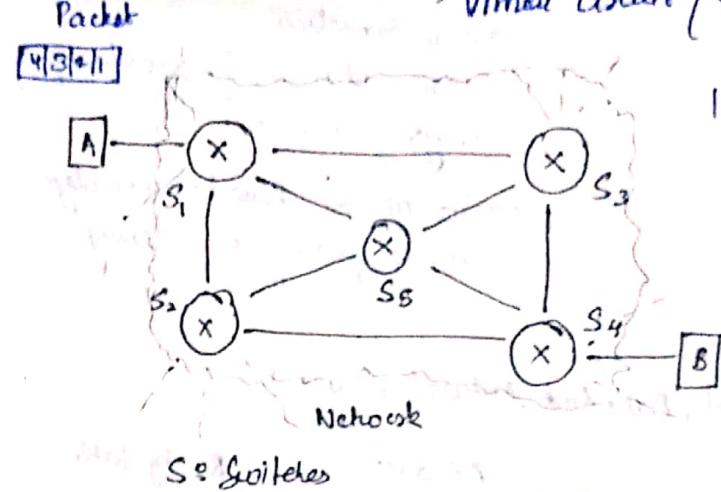
(iv) Delay less: No delay in communication, No star and forward.

$$\text{Total time} = \text{Setup Time} + \text{Transmission Time} + \text{Propagation Delay} + \text{Teardown Time}$$

## Packet Switching

→ Datagram Service (works on Network layer)

→ Virtual Circuit (works in Data link layer)



(i) Pipelining : parallel sending of data.

$$\text{Total time} = \underbrace{(\text{Transmission Time})}_{\text{Store and forward through multiple switches.}} + \frac{\text{PD}}{\text{propagation delay}}$$

Store and forward through multiple switches.

$$\frac{n(T.T)}{2}$$

no of switches

## Datagram Switching Vs. Virtual Switched Circuit.

### Datagram Switching

(i) Connectionless - No Reservation.  
No Resource is Reserved. Everything will be allocated on demand.

(ii) No Reservation

(iii) Out of Order : Since a packet can follow any path

(iv) High Overhead : due to headers.

(v) Packet Lost ↑

(vi) Internet uses a datagram service - IP network.

(vii) Efficiency is more, cost is less, because of no reservation

(viii) Delay is more

### Virtual Circuit

Collection oriented - Before sending, reservation will be done (Buffer, CPU and memory will be used).

(i) Reservation

(ii) Same order.

(iii) less overhead

(iv) Packet Lost ↓

(v) ATM (Asynchronous transfer mode)

(vi) Efficiency is less, cost is more.

(vii) Delay is more

Introduction

→ A computer network is a set of nodes connected by communication links. Mainly used for Resource sharing.

• A node can be a computer, printer or any other device capable of sending / receiving data generated by other nodes in the network.

• A communication link can be wired link or wireless link. Link carries the information.

• End device : They are either the starting point or end point in the communication.

Basic Characteristics

## (i) Fault tolerance

- ability to:

  - Continue working despite failures
  - Ensure no loss of service.

Fault tolerance

Scalability

Security

Basic Characteristics

Quality of Service (QoS)

## (ii) Scalability (Eg: Internet)

- ability to:

  - Grow based on the needs
  - Have good performance after growth

## (iii) Quality of Service (QoS)

- ability to:

  - Set Priorities
  - Manage data traffic to reduce data loss, delay etc.

IV) Security

## ability to provide:

- Unauthorized access
- Misuse
- Forgery

## ability to prevent:

- Confidentiality
- Integrity
- Availability.

Network Protocol & Communications

## (i) SIMPLEX

- Communication is always unidirectional
- One device can transmit and the other device will receive.

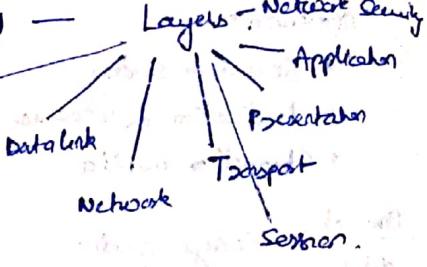
## (ii) HALF DUPLEX

- Communication is in both directions but not at the same time.
- If one device is sending, the other can only receive and vice versa.

## (iii) DUPLEX / FULL DUPLEX

- Communication is in both direction simultaneously.

## (ii) CISCO packet Networks

OSI Model

### (III) PROTOCOLS → Rules governing all methods of communication.

All communication schemes will have following things in common:

- \* Source or sender

- Destination or receiver

- Channel or media.

#### Protocol

define → Message encoding

OR

Message formatting & encapsulation.

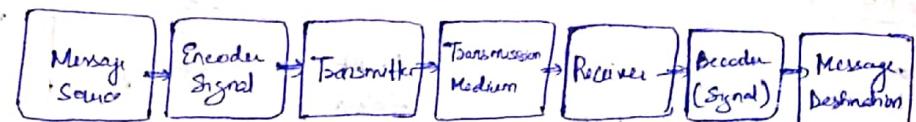
#### Elements of a protocol

Message binary

Message size (breaking of message into packets)

Message delivery options.

#### Message Encoding



(in form of waves)

#### Message formatting & encapsulation

- Agreed format

- Encapsulate the information to identify the sender and the receiver rightly.

#### Message framing

- Flow Control (speed synchronization)

- Response Timeout (synchronization and Acknowledgement)

#### Message delivery options

- Unicast (Sender sending to exactly one destination)

- Multicast (Sending to set of receivers, but not to all)

- Broadcast (to all participants in network).

# IP address is added to the data to be send.

- Every computer in the network is identified by its IP address.

- Data contains both source and destination IP address.

- Acknowledgment, if not done → receiver has to resend the package.

- Message packets are ensured to ensure synchronized receiving and avoid memory loss.

#### Peer to Peer Network

- Decentralized administration

- All peers are equal

- simple sharing applications

- Not Scalable (new devices can't be added to the network)

↳ due to port limitation

#### Client Server Network

- Central administration

- Request - Response model

- Scalable

- Server may be overloaded.

## Components of a Computer Network

- Components
  - Nodes
  - Media
  - Services

Nodes → End nodes (End devices)

Intermediate nodes

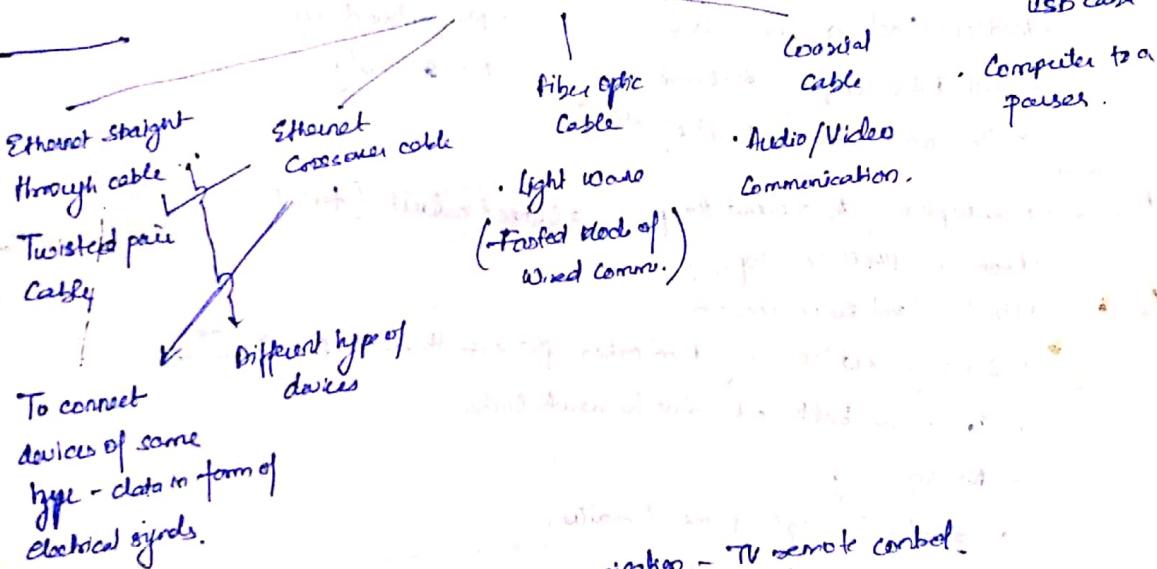
End nodes → Computer, Notebook printers, VoIP phones, Telepresence endpoints, Security cameras, Mobile handheld devices

Intermediate nodes → Switches, WAP/wireless Access point, Router, Security Devices (firewall), Bridges

Nodes → Hub, Repeater, Cell tower.

Media → Wired Medium (Guided Medium) / Unguided - Wireless medium

Wired Medium



Wireless → Infrared (e.g. short range communication - TV remote control).

Media → Radio (e.g. Bluetooth, Wi-Fi)

RFILUVXG.

Microwaves (e.g. Cellular system)

Satellite (e.g. GPS (Long Range communication))

Services → e-mail, storage services, Online game, Voice over IP, file sharing, Video streaming, Instant messaging, World Wide Web.

\* Voice over IP is an intermediary Node.

Classification of Computer Networks - LAN, MAN, WAN.

LAN → Interconnects with a Limited Area

Centrally managed switch.

May be Wired (Hub, Switch)

Ethernet.

Wireless

(Wi-Fi)

MAN → Interconnects users with computer resources in a geographic region of the size of a city

Devices → Switches / Hub → for LAN

→ Router / Bridges → for MAN to connect two switches or hubs.

WAN → Extends over large geographical Area. - Telecommunication Network

- Devices : End devices and Intermediary devices

Internets → WideWAN.

Cloud Computing → On-demand availability of computer system resources, especially data storage and computing power, without direct active management by the user.

Network Topology : Arrangement of nodes of a complete network

Physical Topology

- placement of various nodes,

→ Bus, Any, Star, Mesh, Hybrid.

Bus →

- Common transmission medium

- Bidirectional communication

- Suited for temporary Network

- One node does not affect others

Logical Topology

- Deals with the data flow in the networks.

- No fault tolerant

- No redundancy

- No security

Air → Bus topology in a closed loop

• Closed circuit (for 2 node - 2 cable)

- P2P - to - P2P LAN topology

- Unidirectional communication

- Sending and receiving data takes place with the help of a token

- Can cause bottleneck due to weak links

- No security

- Widest single point of failure

- High load, low performance

Star →

- Centralized Management - hub or switch

- All nodes pass through hub or switch

- Scalable

- Single point of network

point of congestion

- Bottlenecks due to overloaded switch / Hub

- Repeaters connect two or more star topology

Mesh →

- 100% fault tolerance and reliable

- Each node is interconnected

- Issues with broadcast message

- Expensive

# No traffic problem in mesh topology truly.

## Basics of IP Addressing

IP Address → Internet protocol - Every node in the computer network is identified with the help of IP address.

IPv4

- logical address - Can be changed based on logic or location of devices
- Assigned by manually or dynamically
- Represented in decimal and it has 4 octets ( $x.x.x.x$ ) // 32 bits

0.0.0.0 to  $\frac{255.255.255.255}{1\text{byte } 1\text{byte } 1\text{byte } 1\text{byte}}$

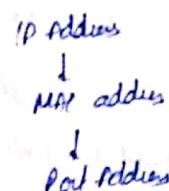
# ipconfig gives IP address of the computer.

## Basics of MAC Addressing

MAC address

- Media Access Control - Used to identify node in LAN network.
- location = location of a person  
may change
- MAC address = Name of a person  
will not change
- Routers need IP address (used by MAN/WAN)
- MAC address is used by switches. (used by LAN)
- physical address - unique, and cannot be changed. - Assigned by the manufacturer.
- represented in hexa-decimal.
- Example: 48 bits ( $\frac{70-30-84-09-60-FC}{1\text{byte } (-) \text{ separator } 1\text{byte } 1\text{byte } 1\text{byte } 1\text{byte } 1\text{byte}})$   
 $6 \times 8 \text{ bits} = 48 \text{ bits}$   
dash. dot colon

# ipconfig /all → to see MAC address details.



## Basics of Port Addressing

→ Port - Communication endpoint.

- Unique for every process.
- fixed port number and dynamic port number ( $0 - 65535$ )  
 $65535 = 2^{16} \text{ or } (2^3)^5$   
 $= 2^8 \times (2^3)^2$   
 $= 2^8 \times 8^2 \text{ bits}$

# Use Resource Monitor to check port number.

Switching Techniques - Path for multiple connections.

Circuit Switching

Circuit Switching

- dedicated path between sender & receiver
- before data transfer, connection will be established first
- Eg: Telephone network
- phases
  - Connection establishment
  - Data transfer
  - Connection disconnection.

Message Switching

Datagram approach

Packet Switching

Virtual Circuit approach

- Message Switching
- Share and forward mechanism.
  - Intermediary nodes is first storage and then forwarded.
  - Not suited for real time application.
- Packet Switching
- Each packet is sent individually. Each packet has source and destination IP address with sequence number.
  - Sequence number → Route the packets
    - | → Detect missing packets
    - | → Send Acknowledgement.
- (connection oriented)
- Datagram Approach (Connectionless switching)
- Each independent entity is called as datagram.
  - Path is not fixed.
  - Intermediary nodes take routing decision.
- Virtual Circuit Switching
- A preplanned route is established.
  - Call request and call accept packets are used to establish the connection between sender and receiver.
  - Path is fixed for the duration of a logical connection.

### Laying in Computer Networks

Layers → Decomposing problem into more manageable components (layers).

- Advantages
  - Modular Design
  - Easy to troubleshoot.

Protocols → Protocols in each layer govern the communication in that layer.

- OSI Model
- OSI : Open Systems Interconnection
  - Model for understanding and designing a network architecture that is flexible, robust, and interoperable.
  - OSI is only a guidance, only a reference model.
  - Purpose is to allow how to facilitate communication between different systems without requiring changes to the logic of the underlying hardware and software.

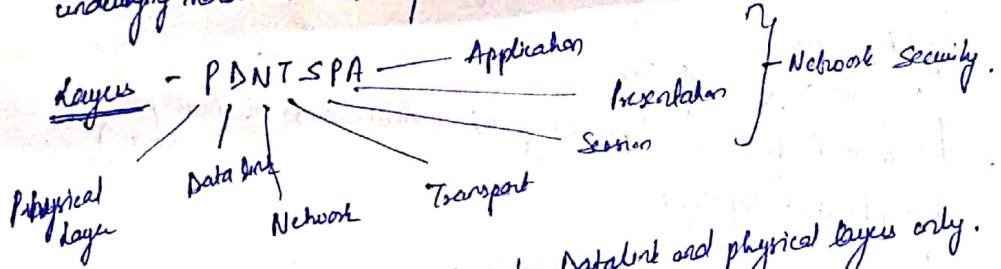
- TCP/IP Model
- TCP : Transmission Control Protocol
  - IP : Internet Protocol
  - Hierarchical protocol making of interactive modules, each of which provides specific functionality.

Every layer  
functions of  
one addressing node.

## The OSI Model

Objectives → Layering means decomposing the problem of building a network into more manageable layers.

Purpose → Communicate between different systems without exposing changes to the logic of underlying hardware and software.



# Intermediate Nodes operate on Network, Datalink and physical layers only.

Application Layer → Enables user to access the network resources

- ↳ File transfer and Access Management (FTAM)
- ↳ Mail Services
- ↳ Directory Services

Presentation Layer → Concerned with the syntax and semantics of the information exchanged between two systems.

↳ Translation: Converting data into a common format

↳ Encryption: Protect data from unauthorized access

↳ Compression: to send multimedia (reduce network traffic).

Session Layer → Maintains, synchronizes the interaction among communicating devices.

↳ Dialog Control: Determines data flow - simplex, duplex, half duplex

↳ Synchronization: Checkpoint synchronization

Transport Layer → Responsible for process-to-process delivery of the entire message.

↳ Port Addressing: handling data to right process.

↳ Segmentation and Reassembly

↳ Connection Control: Establishes connection between nodes

↳ End-to-end flow control: between end-to-end devices

↳ Error control.

Network Layer → Responsible for delivery of data from the original source to the destination network

↳ Logical Addressing (IP Address)

↳ Routing: finding best routes.

↳ Flow control, Error control

↳ Access Control: If two or more devices are connected to the same link, to determine which device has access

↳ Flow control, Error control

↳ Physical Addressing (MAC addressing)

↳ Flow control, Error control

↳ Access Control: If two or more devices are connected to the same link, to determine which device has access

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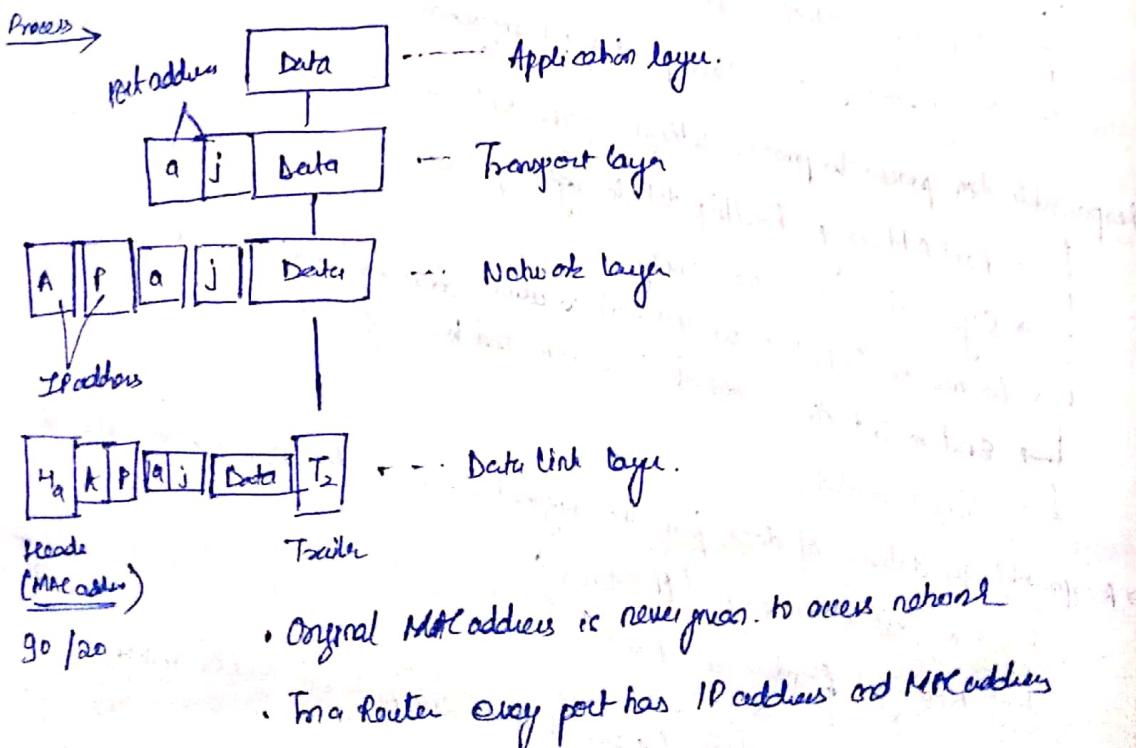
- Physical layer
- Responsible for transmitting bits over a medium. provides electrical and mechanical specifications.
    - ↳ wired or wireless
    - ↳ signal conversion.
  - Representation of bits : encoding (how 0 and 1's are converted to signals)
  - Data rate
  - Synchronization of bits
  - Line Configuration (point to point communication or multipoint layer).
  - Physical topology.
  - Transmission mode / Data flow.

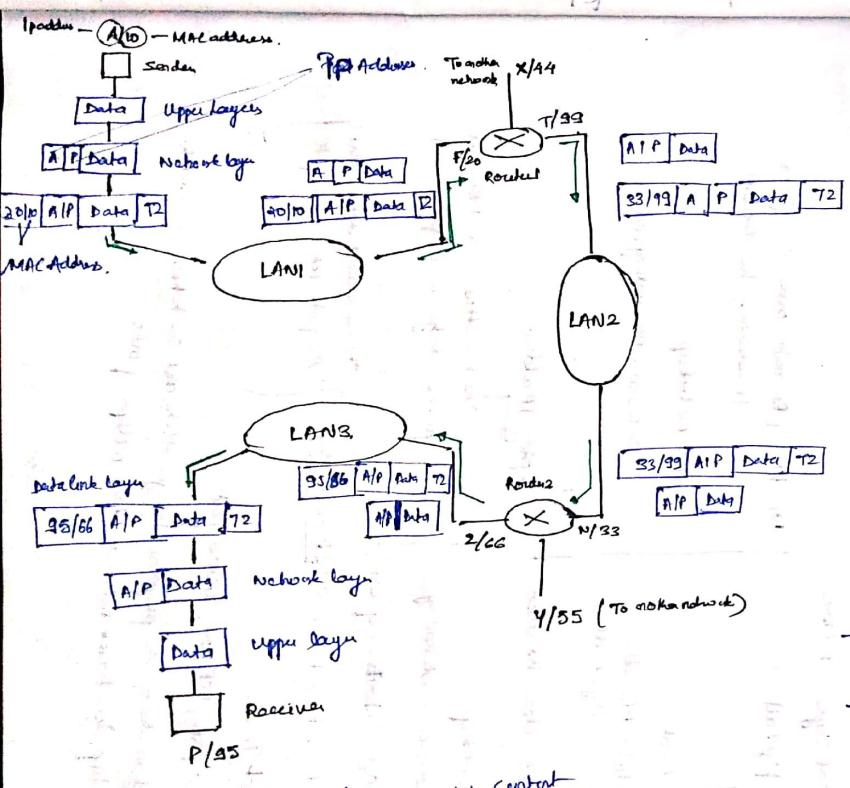
### Addressing in Network

- Port Addressing • Port addressing handled by Transport layer
- IP addressing • IP addressing by network layer.

- General
- Source and destination port no. : 16 bits
  - IP address IPv4 : 32 bits
  - IPv6 : 128 bits

- IP Addressing
- Mac Addresses are of 48 bits.
  - MAC address
  - Routers take decision based on IP addresses.
  - Every node has a Address (MAC or IP).





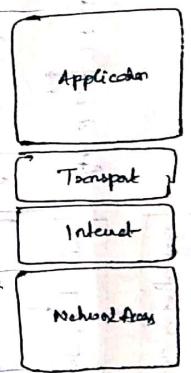
→ Intermediary devices can access Physical layer, Data link layer, Network layer.

### TCP/IP protocol

OSI Ref.

TCP/IP

- functionality of SPA layer
- Protocol: HTTP, DNS, DHCP, PTP



- functionality of T layer
- Protocol: TCP, UDP
- functionality of I layer
- Protocol: IPv4, IPv6, ICMPv4, ICMPv6
- functionality of N layer
- Protocol: PPP, Frame Relay, Ethernet

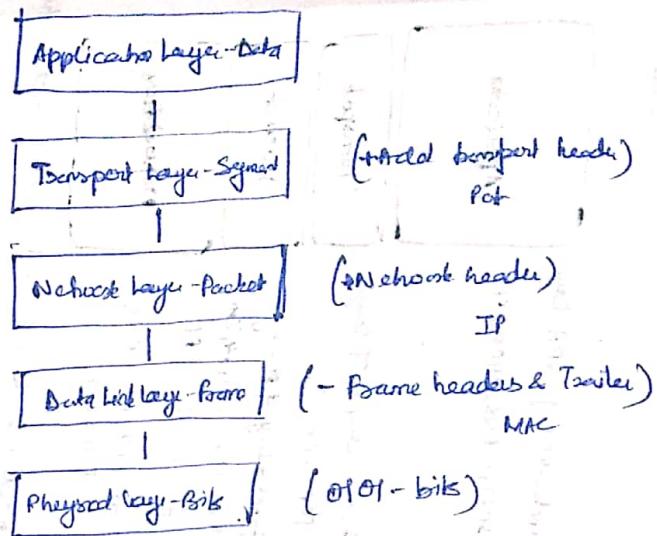
→ Network Access: controls hardware devices and media that make up the Network

→ Internet: Determines best path through the network

→ Transport: Supports communication between devices across direct networks

→ Application: Represents data to the user, plus encoding and decoding control.

Proto col Data Unit :- Data, Segment, packet, frame and bits.



### Basic Networking Commands

(i) ipconfig :- gives us IP configuration details.

- Every IP address is complemented by a subnet mask
- Default gateway is Router's IP address

Layer 3 information

ie Network layer - IP Address .

(ii) ipconfig /all :-

Layer 2 information

ie Data link - MAC address  
(Physical address)

(iii) clscs clean the screen.

(iv) >nstartip

>~~www~~.resacademy

> NC (about my operator)

> baceert (base root) IP address

(v) DNS (Domain name server gives IP address for the domain)

cmd > ns lookup

> www.resacademy.org

(vi) ping

> ping IP address

Check whether its reachable or not

(receives acknowledgments of packets sent)

Shows several details about the path taken by the packets

### Cisco packet Tracer

i) Serial ports, USB port, Ethernet port

ii) Hub works in the physical layer. Hubs and switches are used to setup local area network. Has multiple ports

## Hub / Network Hub

- Used to Set up LAN
- Has multiple ports
- follows star topology.
- Cable: Straight Through, crossover cable.

100BaseT

- cheaper than switches
- works good for smaller networks

cons

- Issues with band cost
- Nonroutable

## Switch

- Switch has memory. Can store MAC Address in memory. (In conjunction with interface / port).
- It's a layer 2 device.

## Router

- A Networking device that forwards packets between computer networks.
- Connected to at least two networks, commonly two LANs or WANs or a LAN and its ISP's network.
- Operates in Network Layer. Stores routing table in memory. The interfaces of a router are part of a LAN.
- Default gateway information is stored with all devices. This is the IP address of the interface in Router.

## Repeater

- physical layer device
- Repeater just Regenerates the signal.
- A hub port device

## Bridge

- Bridge = Repeater + functionality of reading MAC address.
- It is a layer 2 device.
- Used for interconnecting two LANs on the same protocol.
- A two port device.

## Transparent Bridges

## Source Routing Bridges

- Routing is performed by source station and the path specifies which route to follow.

- These are the bridges in which no stations are completely unaware of the bridge's existence.

- Recomfiguration of the stations is unnecessary even if bridge is added or removed from network.

## Multi-layer Switch

- Can act as a switch and a router

## Firewall

- Security Device
- filters incoming and outgoing traffic.

## Bridge

- Act as a layer 2 bridge in a router

## Modem

- Modulates and demodulates information



## → MAC Sublayer

• Implemented by hardware, typically in the computer NIC.

• Two primary responsibilities:

### → Data Encapsulation

- Frame assembly and frame disassembly.
- Adds header and trailer to the network layer PDU.

• Functions:

- Framing
- Physical addressing or MAC addressing
- Error control

### → Media Access Control

• Responsible for the placement of frames on the media and the removal of frames from the media.

## FRAMING

i). frames are packets of zeros and ones.

• The start and end of the frame is set by a protocol. (boundary bits)

• frame = Header + Network Layer PDU + Tailer.

Protocol Data Unit

• Types of framing

### fixed size framing

- Size of frame is fixed and so the frame length acts as delimiter of the frame.
- Consequently, does not require additional boundary bits to identify the start & end of the frame.

• Size of each frame to be transmitted may be different.

### Variable-size framing

## Framing Approaches

(i)

### Bit Oriented

• Frame is viewed as a collection of bits.

• Protocol: HDLC (High Level Data Link Control)

### Character Oriented

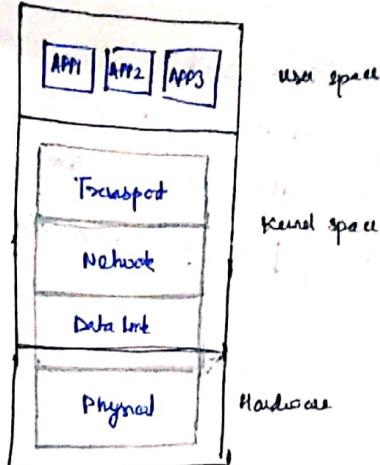
• E.g. SONET - Synchronous Optical Network

### Byte Oriented

• Each frame viewed as a collection of bytes (characters) rather than bits.

- Protocol:
  - BSCP: binary synchronous communication protocol
  - DDCMP: digital data communication protocol
  - PPP: point-to-point protocol.

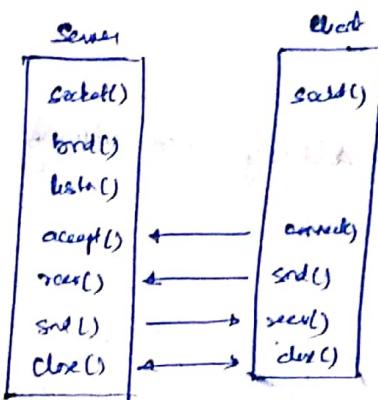
## Socket Programming



Use spell

→ logical zipper between two objects is called socket. Form end-to-end connection in TCP, and end-to-end connections in case of UDP.

→ A set of system calls to get the services from TCP/IP protocol stack (not available in Kernels)



→ i) Socket Programming framework/API's

~~server() announces a port and client needs to make a connection to that particular port.~~

- This will create the server side opening of the logical pipe and it will bind to socket with your TCP/IP protocol stack.

→ Bond( ) function : with the port number you are specifying it will find that port number with that socket.

→ listen() function: help you to make the server go to the listening state,

→ connect(): to initiate a connect call; Client sends a SYN packet, Server returns an ACK packet.

→ `send()`: send to data

`scelene()`: to select data.

### (i) Socket Type

• TCP is Reliable whereas UDP is Unreliable

• Types of sockets: — Stream Socket (SOCK\_STREAM)

Datagram Socket (SOCK\_DGRAM)

✓ Unzulässige, korrigierte Form

• Low Socket - helps in directly interacting with IP layer.

### (iii) Socket API

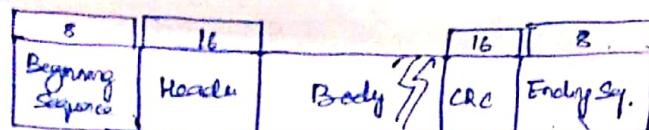
- int s=socket(domain, type, protocol); - Create a socket
  - domain : Communication domain, typically used AF\_INET (IP4 Protocol)
  - type : Type of the socket - SOCK\_STREAM or SOCK\_DGRAM
  - Protocol : Specifies protocols - usually set to 0 - Explained !
- int status = bind(sockfd, &addrport, size); - Reserves a port for the socket
  - sockfd : Socket identifier
  - addrport : struct sockaddr\_in - the (IP) address and port of the machine (address usually set to INADDR\_ANY chooses a local address)
  - size : Size of the sockaddr structure .

## High-level Data Link Control (HDLC)

### I) BIT ORIENTED APPROACH.

- Simply view the frame as a collection of bits.
- HDLC is bit level Data Link Control.

### (ii) NALC Frame format.



• This sequence is also beneficial during any time that the link is idle so that the sender and receiver can keep their clock synchronized.

- Beg and ending seq : 0111110.

- Header Address and Control field.

Body : Payload (variable size)

CRC : Cyclic Redundancy check - Error detection.

(iii)

Types of HDLC Field.

(in control field)

I-frame

Information frame

• 1st bit is 0

i.e., this frame is carrying information

S-frame

Supervisory frame

• 1st two bits is 10

for flow & error control mechanism

0-frame

• Un-numbered frame.

• 1st two bits is 11.

Using link management and miscellaneous activities

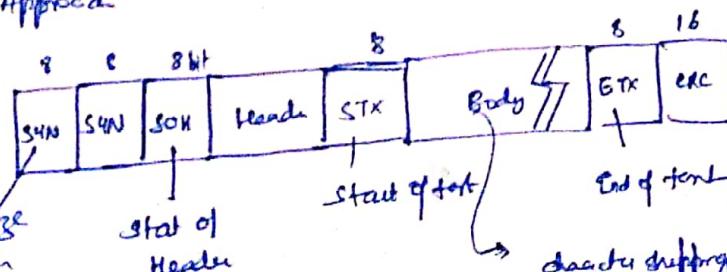
## Bit Stuffing

- A Bit is stuff to properly distinguish frame sequence.

## BSC : Binary Synchronous communications protocol (BSC)

### A Sentinel Approach

body + guard



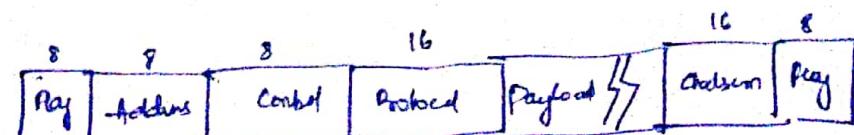
character stuffing: whenever a flag or escape character appears in the body.

## PPP Protocol

- A won protocol over Internet Links used in broadcast communications having heavy loads and high speeds

- PPP

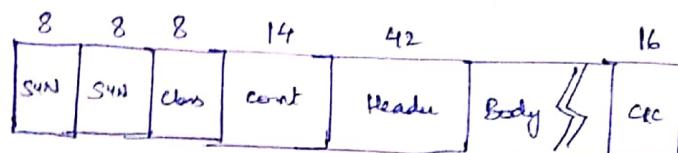
- character stuffing is done using DLE: Data Link Escape protocol.



- Flag is 0111110 : marks the beginning and the end of the frame.
- Address is set to 1111111 in case of broadcast.
- Control field : Control set at 1100,0000
- Protocol: 1 or 2 bytes that define the type of data contained in the payload field.
- Payload : This carries the data from the network layer.. Maximum length of the payload field is 1500 bytes . this may be negotiated between the endpoints of communication.
- Characters stuffing : An extra byte of information is added.

DDCMP : Digital Data Communication Message Protocol

- Uses a byte counting approach.
- Count field determines how many bytes are contained in the frame body.



- Danger : if transmission error could corrupt the count field , then the end of the frame would not be correctly detected by the receiver.

## ERROR DETECTION

- Data can be corrupted during transmission.
- For reliable communication, errors must be detected and corrected.
- Error detection and correction are implemented either at the data link layer or the transport layer of OSI model.

### Types of Error.

#### BIT ERROR

- only 1 bit in the data unit has been changed.

- Error detection is done by the Receiver. To detect or correct error, need to send some extra bits with the data. These bits are called a redundant bit.

A Generating function is used for redundancy check.

A Checking function is used for checking consistency.

- Error Correction
  - (i) Receiver can have the code retransmit entire data if (ii) use an error-correcting code, automatically correct certain errors.

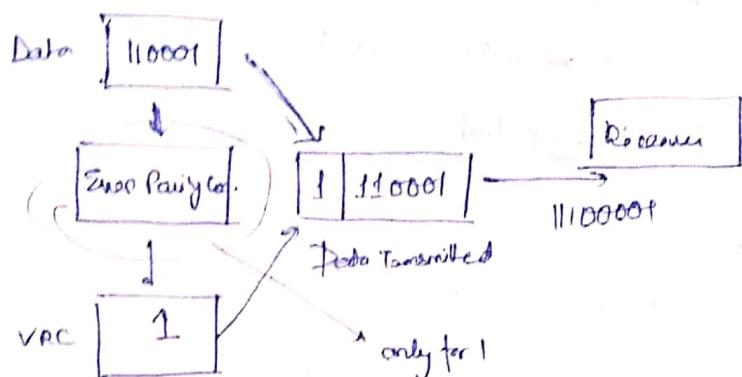
#### BURST ERROR

- 2 or more bits in data unit have changed.

## Parity & Detection Techniques

- 1. VRC: Vertical Redundancy Check
- 2. LRC: Longitudinal Redundancy check (LRC)
- 3. Check sum
- 4. CRC: Cyclic Redundancy check (CRC)

### Vertical Redundancy check (parity check)



Performance

- Can detect single bit error
- Can detect burst errors only if the no. of errors is odd.

### Longitudinal Redundancy check (Two dimension parity)

- In LRC, a block of bits is organized in rows and columns.
- parity bit is calculated for each column and sent along with the data.
- Block of parity act as redundant bit.
- At last LRC is added.

### Checksum

$$\text{Checksum} = \text{check} + \text{sum}$$

(a) Checksum: Creating : Sender side

(b) Checksum Validation : Receiver side

- Break the original message into 'k' number of blocks with 'n' bits in each block.
- Sum all the 'k' data blocks.

98  
47  
51

- Add the carry to the sum, if any
- Do 1's complement to the sum + Checksum (n bits)

$$\begin{array}{r}
 150 \\
 +33 \\
 +07 \\
 \hline
 90 = 3011
 \end{array}$$

for receiver : The result of ALL blocks including parity block, sum must be 1111.

→ Communication Protocols