
MODULE *Channel*

Taken from “Specifying systems”

EXTENDS *Naturals*

CONSTANT *Data*

VARIABLE *chan*

$TypeInvariant \triangleq chan \in [val : Data, ready : \{0, 1\}, ack : \{0, 1\}]$

In the begining, the *ack* and *ready* flags are the same.

$Init \triangleq \wedge TypeInvariant$
 $\wedge chan.ack = chan.ready$

When the flags are the same, you can send the message.

$Send(d) \triangleq \wedge chan.ready = chan.ack$
 $\wedge chan' = [chan \text{ EXCEPT } !.val = d, !.ready = 1 - chan.ready]$

When the flags are not the same, you can confirm the message.

$Receive \triangleq \wedge chan.ready \neq chan.ack$
 $\wedge chan' = [chan \text{ EXCEPT } !.ready = 1 - chan.ready]$

$Next \triangleq (\exists d \in Data : Send(d)) \wedge Receive$

$Spec \triangleq Init \wedge \square [Next]_{chan}$

THEOREM $Spec \Rightarrow \square TypeInvariant$
