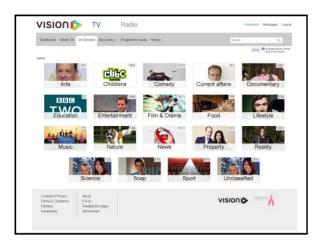


Quick Live Overview of Vision IPTV









Challenge in Video-on-Demand (aka the Problem)

VoD Distribution Efficiency

- VoD requests handled naively independent flow per request
- These are duplicated minutes, hours or days later (by same or different user)
- Identical delivery of media objects through the same network segments
- End-to-end capacity of network infrastructure must grow continuously to match the increasing number of Internet video users
- The increasing popularity of VoD and especially of HD content worsens this
- What is NOT a solution :
 - Multicast : VoD requests are not for the same content at the same time
 - Peer-to-Peer : Limited storage and uplink resources on user devices (peers) –
 cannot guarantee high QoE for the users

Key Characteristics of Video-on-Demand

- · High-throughput end-to-end
 - Not just high egress capacity at origin video servers, but also adequate bandwidth available in all networks in between video source and users
- Distance matters between source VoD server and user
 - (Standard) TCP used for VoD can become bottleneck as it requires ACKs for every window of data packets sent
 - TCP's throughput is inversely related to
 petwork latency or RTT.

Distance (Server to User)	Network RTT	Typical Packet Loss	Throughput	4GB DVD Download Time
Local: <100 mi.	1.6 ms	0.6%	44 Mbps (high quality HDTV)	12 min.
Regional: 500-1,000 mi.	16 ms	0.7%	4 Mbps (basic HDTV)	2.2 hrs.
Cross-continent: ~3,000 mi.	48 ms	1.0%	1 Mbps (SD TV)	8.2 hrs.
Multi-continent: ~6,000 mi.	96 ms	1.4%	0.4 Mbps (poor)	20 hrs

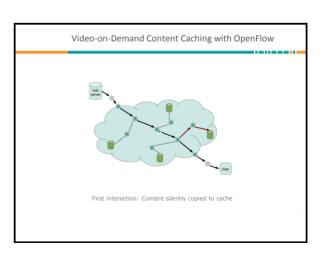
Effect of Distance on Throughput and Download Time [1]

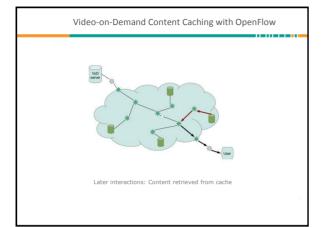
- We need a solution that :
 - Ensures high-throughput end-to-end
 - Minimizes distance between source video content server and user

SDN?

[1] E. Nygren, R. K. Sikaraman, and J. Sun. The Akamai Network: a Platform for High-Performance Internet Applications. SIGOPS Oper. Syst. Rev., 44(3):2–19, 201

Video-on-Demand Content Caching with OpenFlow **CoperFlow switch **R Department of Cache An OpenFlow network with peripheral content caches





Caching Requirements (1)

An OpenFlow-based content caching architecture should satisfy the following functional requirements:

- 1. Should identify cacheable content without any significant impact on the user's request
- 2. Should cache content transparently to the user
- 3. Should deliver content transparently to the user
- 4. Should retain the underlying content delivery mechanism to avoid fundamental changes to the service
- 5. Should be content agnostic
- 6. Should be easily integrated in a production network
- 7. Should be able to use multiple cache instances
- 8. Should be able to add or remove cache instances without service interruption

Worth thinking:
How could we do the
above without SDN?

Caching Requirements (2)

An OpenFlow-based content caching should satisfy the following non-functional requirements:

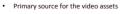
- Should optimize network utilization. For example, it should not unreasonably cache content that is infrequently requested and thus increase the network utilisation unpressarily.
- 2. Should adjust its run-time functionality and improve the users' QoE by maintaining a high level view of the network based on run-time metrics (e.g. buffering times etc.)
- 3. Should support load balancing between carefully and strategically located in-network caches

Worth thinking:
How could we do the
above without SDN?

OpenFlow-assisted In-network Caching Architecture

Entities (1)

- Any hardware or software OpenFlow Switch
 - Must be able to communicate with the VoD server, the CSS(s) and the OpenFlow controller, but not necessarily directly



Could be located anywhere on the Internet (reachable by IP)



- Any kind of OpenFlow Controller (e.g. Floodlight, NOX, POX)
 - Should be reachable by the OpenFlow Switch
 - Runs L2 learning switch : allows the switch to forward on MAC-to-Port pairing
 - Exposes a JSON-RPC Flow Pusher interface to OCI

Entities (2): OpenFlow Caching Intelligence (OCI)

- Orchestrator of in-network caching functionality
 - Provides a JSON-RPC interface to retrieve requests for content to be cached in a highly flexible and configurable fashion (Cache as a Service)
 - Used by network administrator or even content providers (via SLAs)
 - Supports regular expressions to fine tune requests for content (e.g. particular video, all videos from a domain, a type of video from any domain (n.b. with later versions of OpenFlow))

METHOD	PARAMETERS	RESULT
start-expr	{ "expr" : <expr> }</expr>	<boolean></boolean>
stop-expr	{ "expr" : <expr> }</expr>	<boolean></boolean>
list-expr-all	None	[{'expr': <expr>, 'port': <port>},]</port></expr>

- 2. Implements the caching logic : what should be cached where at each point in time $\frac{1}{2}$
 - Enhanced to support resource monitoring and load-balancing

Entities (2): OpenFlow Caching Intelligence (OCI)

Orchestrator of in-network caching functionality



METHOD	PARAMETERS	RESULT	
hello	{ "host" : <host>, "port" : <port> }</port></host>	<node-id></node-id>	
Keep-alive	{ "node-id" : <node-id> }</node-id>	<boolean></boolean>	
goodbye	{ "node-id" : <node-id> }</node-id>	<boolean></boolean>	

- 4. Manages the OpenFlow switches of the network via the Controller
 - Adding/removing flows to switches via the Flow Pusher API of the controller so that users' requests are served appropriately

URI	DESCRIPTION ARGUMENTS	
/wm/staticflowentrypusher/json	Add/Delete static flow	HTTP POST data (add flow), HTTP DELETE (for deletion)

Entities (3)

- Key-Value Store to maintain a list of :
 - All names of videos that have been requested for caching



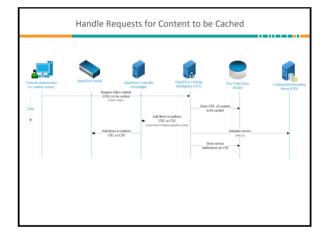
- Videos that have been cached and where
- Status of CSS (online/offline, reachable etc.), their location and resources
- Caching and Streaming Server (CSS)

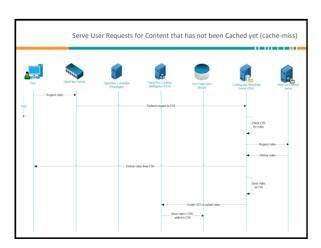


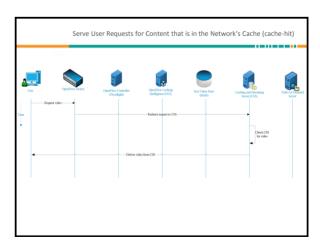
- Multiple CSS instances in the network, possibly connected directly to the switch and consequently to the user: Lower latency and faster response times (high QoE)
- Three operations :
 - 1. Communicate its status to the OCI
 - 2. Caching content that is requested from the user
 - 3. Stream content that is being already cached

Three Essential Operations

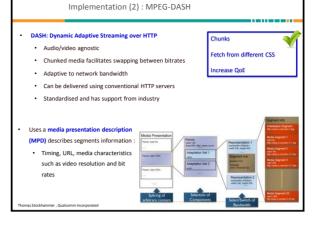
- 1. Handle requests for content to be cached (Cache as a Service)
 - From network admins/content providers
- 2. Serve user requests for content that has not been cached yet
 - Serve user and cache content for future use
- 3. Serve user requests for content that is in a network's cache

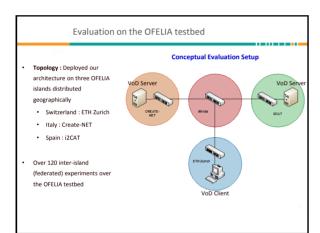


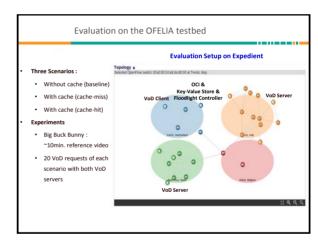


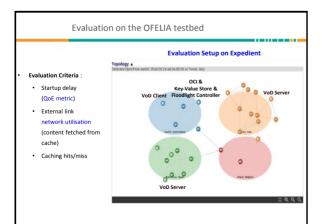












Results						
	CREATE-NET (Italy)			i2CAT (Spain)		
	Without Cache	Cache-miss	Cache-hit	Without Cache	Cache-miss	Cache-hit
Average Startup Delay (s)	2.484	2.088	1.639	2.212	1.982	1.441
Improvement over Baseline (%)	-	16.02	34.02	-	10.40	34.85
Standard Deviation (σ)	0.208	0.225	0.226	0.145	0.138	0.109
External Link Usage (Bytes)	105,734,144	105,827,872	0	105,734,144	105,827,872	0

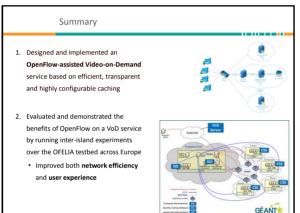
• Key results :

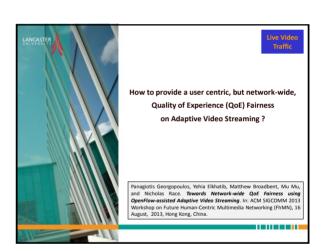
- In tests over both islands we reduced the startup delay up to 35% -> increased QoE for end-user
- External link utilisation reduced to virtually zero (only background traffic remained)
 - Indicatively, the full streaming of our ~10min video saved ~100MBytes for just one client session

Average Startup Delay External Link Usage ORATE-NET Without Cache Cache-miss Cache-hit - 35% Improvement even in a bandwidth reach environment (OFELIA testbed) Reinforced by relatively low standard deviation values Greater improvements would be possible on next generation OpenFlow switches where packet processing will take place on the hardware path

Advantages of our VoD In-network Caching Architecture

- 1. Provides an interface for cacheable content in an "open", highly-configurable, controllable and flexible manner
- 2. Centrally controlled caching: efficient load balancing, allows pre-caching of frequent content
- Easily deployable in a production network: the underlying delivery video mechanism will remain the same in an OpenFlow network (existing hardware and software can be retained, no fundamental changes in service)
- Fully transparent to the user: no need to install any extra software or have to sacrifice any of his
 local network or storage to be able to stream HD content with high efficiency.
- 5. Caching very close to the user :
 - Reduces network utilisation as requests are served locally: minimize the amount of packets
 that are required to traverse the network from the source media provider to the user
 - The video QoE of the end-user will improve, as the user will experience lower latency, smaller buffering times and higher video quality as content is now located locally





Adaptive Video Streaming (e.g. MPEG-DASH) aims to increase QoE and maximise connection utilisation (supporting chunks encoded at different bitrates) Many implementations are bursty and unstable in nature and naively estimate available bandwidth from a one-sided client perspective No account of other devices in the network Results in unfairness; video streams fight over link's capacity which causes network congestion (video quality degradation, frame freezing etc.) and potentially lowers QoE for all clients Counter productive!

The Problem

Potential Solution

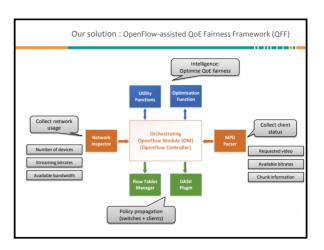
Split available bandwidth to current users on the network?

But naïve network resource fairness (equal split) is unfair:

You could easily satisfy a user watching a video on his smartphone, but it is much harder for an HD TV

xkbps

**xkbp



Poptimisation Function finds the optimum bitrate for each streaming video device in the network that results in equivalent QoE levels for all devices

But the utility functions are not continuous, i.e. we don't have available encodings for all possible bitrates

Implemented branch and bound optimisation algorithm that downgrade all clients to the maximum feasible bitrate (max-min fairness)

Very modest computational overhead < 0.3 sec for optimising 100 Utility Functions with 10 different bitrates each

