

## EECS 212

Northwestern 2019

### Homework 8

Due at 3:59pm, June 4th (Tuesday), 2019

Congratulations! Finally you are here!

*Scoring:* There are total 20 points in this homework. Problem 1 and Problem 4 carries 3 points each. Problem 2 carries 8 points. Problem 3 carries 6 points.

Please write down sufficient steps how you get the result. Simply writing down a final answer will not give your full credit. You are responsible for any mistake resulted from illegible submissions.

#### Problem 1

For any two positive integers  $a, b$ , define  $k(a, b)$  to be the largest  $k$  such that  $a^k \mid b$  but  $a^{k+1} \nmid b$ . Given two positive integers  $x, y$ , show that

(a)  $k(a, \gcd(x, y)) = \min\{k(a, x), k(a, y)\}$  for any positive integer  $a$ .

(b)  $k(a, \text{lcm}(x, y)) = \max\{k(a, x), k(a, y)\}$  for any positive integer  $a$ .

*Hint:* Think of the prime factorization of the numbers.

#### Problem 2

Consider any regular simple graph  $G$  on  $n$  vertices of degree  $d$  and let  $A$  be its adjacency matrix. Consider the matrix  $L = dI - A$  where  $I$  is the identity matrix of size  $n \times n$  (it has 1s on the diagonal and 0 elsewhere). For any set  $S \subseteq \{1, 2, \dots, n\}$ , let  $1_S \in \mathbb{R}^n$  be a vector such that the  $j$ th co-ordinate is  $1_S(j) = 1$  if  $j \in S$  and  $1_S(j) = 0$  if  $j \notin S$ .

(a) Show that for any vector  $x \in \mathbb{R}^n$  show that

$$x^T L x = \sum_{(i,j) \in E} (x_i - x_j)^2.$$

Note that in the right hand side of the above formula, there are  $nd/2$  terms; one term for each edge in  $E$ . Hence both  $(i, j)$  and  $(j, i)$  are both captured by one term that corresponds to the edge between  $i, j$ .

(b) Show that  $(1_S^T L 1_S)$  is the number of edges between the set of vertices  $S$  and the rest of graph (i.e. number of edges  $(u, v)$  such that  $u \in S$  and  $v \in \{1, 2, \dots, n\} \setminus S$ ).

- (c) Show that the matrix  $A$  has an eigenvalue of  $d$ .
- (d) Show that if the graph is disconnected, there are at least two eigenvalues equal to  $d$  (note that you need to exhibit two eigenvectors of this graph, whose eigenvalue is  $d$ ).

### Problem 3

A general would like to count his soldiers before a battle, but they are too numerous for him to do so by counting them one by one. Instead, he asks his soliders to get into a number of different formations consisting of rows of particular lengths. He can then determine the total number of soldier by counting the number of soldiers remaining who do not form a complete row in each of these formations. In this problem, we will find out how he can do so.

Let  $m_1, m_2, \dots, m_k$  be  $k$  prime numbers. For some unknown positive integer  $x$ , suppose following relations holds:

$$\begin{aligned} x &\equiv a_1 \pmod{m_1} \\ x &\equiv a_2 \pmod{m_2} \\ &\dots \\ x &\equiv a_k \pmod{m_k} \end{aligned} \tag{1}$$

We want to find the smallest  $x$  that satisfies the above congruence relations. Let  $M = m_1 m_2 \dots m_k$  and  $M_i = M/m_i$ .

- (a) Show that there exists some integer  $M_i^{-1}$  such that

$$\begin{aligned} a_i M_i M_i^{-1} &\equiv a_i \pmod{m_i} \\ a_i M_i M_i^{-1} &\equiv 0 \pmod{m_j} \text{ for } (i \neq j). \end{aligned}$$

- (b) Show that

$$x_0 \equiv (a_1 M_1 M_1^{-1} + a_2 M_2 M_2^{-1} + \dots + a_k M_k M_k^{-1}) \pmod{M}$$

satisfies the above relations.

- (c) Show that if  $x_1, x_2$  both satisfy the above relations given by (1), then  $x_1 - x_2$  is a multiple of  $M = m_1 m_2 \dots m_k$ . Use this to conclude that  $x_0$  is indeed the smallest positive integer satisfying the above relations.

*Hint:* Use the fact that if  $\gcd(m, p) = 1$  then there is some  $m'$  (which we will denote by  $m^{-1}$ ) such that  $mm' \equiv m'm \pmod{p} \equiv 1 \pmod{p}$ .

#### Problem 4

In class we know that, for two positive integers  $a, b$ , there exists two positive integers  $s, t$  such that

$$sa + tb = \gcd(a, b).$$

In this problem, we will modify the Euclids Algorithm to find the  $s$  and  $t$ .

**gcd'**( $a, b$ )

1. If  $a < b$ , swap  $a, b$
2.  $r \leftarrow \text{remainder}(a, b)$
3. If  $r = 0$ , return  $(0, 1)$ ;
4.  $(s', t') \leftarrow \text{gcd}'(b, r)$
5. return  $(t', s' - (a - r)t'/b)$

We will figure out why the above algorithm works in part (a) and (b).

Let  $r = \text{remainder}(a, b)$  and assume

$$s'b + t'r = \gcd(b, r)$$

when  $r > 0$ .

(a) Show that if  $r = 0$ , then  $s = 0$  and  $t = 1$ .

(b) Show that if  $r > 0$ ,

$$t'a + (s' - \frac{(a - r)t'}{b})b = \gcd(a, b)$$

*Hint:* Represent  $r$  using  $r = a - (a - r)/b \cdot b$ .

#### Problem 5 (Optional: not for grade)

Suppose we have some data points  $(m_1, a_1), \dots, (m_k, a_k) \in \mathbb{R}^2$ , we want to find a polynomial  $p(x)$  of degree  $k - 1$  such that  $p(m_i) = a_i$  for  $i = 1, \dots, k$ .

It turns out that we can actually use the idea of Problem 3 to find this  $p(x)$ . In fact, we can view this problem in the same format as that of Problem 4:

Find  $p(x)$  with degree less than  $k - 1$ , satisfying

$$p(x) \equiv a_1 \pmod{(x - m_1)}$$

$$p(x) \equiv a_2 \pmod{(x - m_2)}$$

$\dots$

$$p(x) \equiv a_k \pmod{(x - m_k)}$$

- (a) Find an analogy for the  $M_i$  and  $M_i^{-1}$  respectively in this problem.
- (b) Construct a  $p(x)$  using the same manner as part (b) of Problem 3. Verify that your  $p(x)$  indeed satisfy  $p(m_i) = a_i$  for  $i = 1, \dots, k$ .