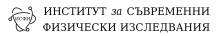
# Лекция 5. Графови невронни мрежи. DSCM023 "Прогнозиране чрез анализ на данни - III"





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План 2

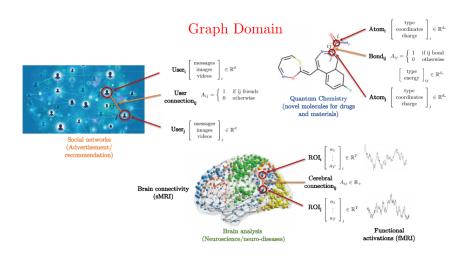
Графи и въпроси върху операцията конволюция

Спекрална теория на графи

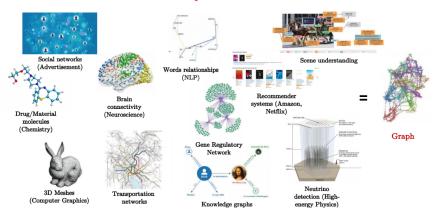
Конволюция на графи in inverse space

# Графи и въпроси върху операцията конволюция

https://youtu.be/Iiv9R6BjxHM



# Graph Domain

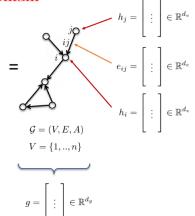


# Graph Domain

- Graphs G are defined by :
  - Vertices V
  - Edges E
  - Adjacency matrix A



- Graph features :
  - Node features : h<sub>i</sub>, h<sub>i</sub> (atom type)
  - Edge features :  $e_{ij}$  (bond type)
  - Graph features : g (molecule energy)



# Операцията "конволюция"

Input Kernel Output

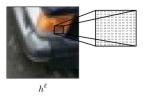
15

### • Convolutional layer (for grids):

$$h^{\ell+1} = w^{\ell} * h^{\ell}$$

$$n_1 \times n_2 \times d \qquad n_1 \times n_2 \times d$$

$$3 \times 3 \times d$$



Image/Hidden features



 $w^{\ell}$ 

Pattern/kernel (learned by backpropagation)



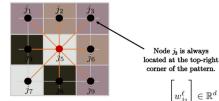
 $h^{\ell+1}$ 

## Convolution

- How to define convolution?
  - Definition #1: Convolution as template matching
  - O(n) by parallelization and for compact support patterns



$$\begin{split} &= \sum_{j \in \mathcal{N}_i} \langle w_j^\ell, h_{ij}^\ell \rangle \\ &= \sum_{j \in \mathcal{N}_i} \langle \left[ w_j^\ell \right], \left[ h_{ij}^\ell \right] \rangle \end{split}$$



All nodes of the template  $w^l$  are always ordered/positioned the same way!

$$\left[w_{j_3}^{\ell}\right] \in \mathbb{R}^d$$

Template features at i.

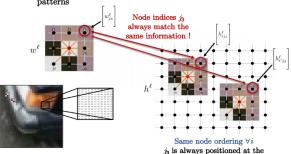
Node j3 is always

corner of the pattern.

### Convolution

top-right corner of the template.

- How to define convolution?
  - Definition #1 : Convolution as template matching
  - O(n) by parallelization for compact support patterns



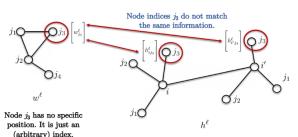
$$\begin{split} h_i^{\ell+1} &= w^{\ell} * h_i^{\ell} \\ &= \sum_{j \in \mathcal{N}_i} \langle w_j^{\ell}, h_{ij}^{\ell} \rangle \\ &= \sum_{j \in \mathcal{N}_i} \langle \left[ w_j^{\ell} \right], \left[ h_{ij}^{\ell} \right] \rangle \\ \\ \end{bmatrix} & \left\langle \left[ w_{j_3}^{\ell} \right], \left[ h_{ij_3}^{\ell} \right] \right\rangle \end{split}$$

These matching scores are always for the top-right corner between the template and the image patches.

 $\langle \left| w_{j_3}^{\ell} \right|, \left| h_{i'j_3}^{\ell} \right| \rangle$ 

# Graph Convolution

- Can we extend template matching for graphs?
  - Main issues :
    - No node ordering: How to match template features with data features when nodes have no given position (index is not a position)?



No node ordering on graphs:

The correspondence of nodes is lost on graphs. There is no up, down, right and left on graphs.

$$\begin{split} & \left\langle \left[ w_{j_3}^{\ell} \right], \left[ h_{ij_3}^{\ell} \right] \right\rangle \\ & \left\langle \left[ w_{j_3}^{\ell} \right], \left[ h_{i'j_3}^{\ell} \right] \end{split}$$

These matching scores have no meaning as they do not compare the same information.

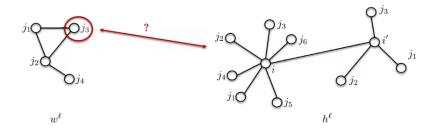
$$h_i^{\ell+1} = w^{\ell} *_{\mathcal{G}} h_i^{\ell}$$

$$= \sum_{j \in \mathcal{N}_i} \langle w^{\ell}, h_{ij}^{\ell} \rangle$$

r Bresson

# Graph Convolution

- Can we extend template matching for graphs?
  - Main issues :
    - No node ordering: How to match template features with data features?
    - Heterogeneous neighborhood: How to deal with different neighborhood sizes?



# Graph Convolution

- How to define convolution?
  - Definition #1 : Template matching
  - Definition #2 : Convolution theorem
    - Fourier transform of the convolution of two functions is the pointwise product of their Fourier transforms

$$\mathcal{F}(w*h) = \mathcal{F}(w) \odot \mathcal{F}(h) \quad \Rightarrow \quad w*h = \mathcal{F}^{-1}(\mathcal{F}(w) \odot \mathcal{F}(h))$$

- Generic Fourier transform has O(n²) complexity, but if the domain is a grid then complexity can be reduced to O(nlogn) with FFT<sup>[1]</sup>.
- Can we extend the Convolution theorem to graphs?
  - How to define Fourier transform for graphs?
  - How to compute fast spectral convolutions in O(n) time for compact kernels?

$$w *_{\mathcal{G}} h \stackrel{?}{=} \mathcal{F}_{\mathcal{G}}^{-1}(\mathcal{F}_{\mathcal{G}}(w) \odot \mathcal{F}_{\mathcal{G}}(h))$$

Права и обратна трансформация:

$$H(f) = \int_{-\infty}^{\infty} h(t)e^{2\pi i f t} dt,$$
  $h(t) = \int_{-\infty}^{\infty} H(f)e^{-2\pi i f t} df$ 

then
$H(-f) = [H(f)]^*$
$H(-f) = -[H(f)]^*$
H(-f) = H(f) [i.e., $H(f)$ is even]
H(-f) = -H(f) [i.e., $H(f)$ is odd
H(f) is real and even
H(f) is imaginary and odd
H(f) is imaginary and even
H(f) is real and odd

$$F(g * f) = F(g).F(f),$$
  $g * f = F^{-1}(F(g).F(f))$ 

Права и обратна трансформация:

$$H(f) = \int_{-\infty}^{\infty} h(t)e^{2\pi i f t} dt,$$
  $h(t) = \int_{-\infty}^{\infty} H(f)e^{-2\pi i f t} df$ 

If ... then ... 
$$h(t) \text{ is real} \qquad H(-f) = [H(f)]^*$$

$$h(t) \text{ is imaginary} \qquad H(-f) = -[H(f)]^*$$

$$h(t) \text{ is even} \qquad H(-f) = -[H(f)]^*$$

$$h(t) \text{ is even} \qquad H(-f) = H(f) \quad [\text{i.e.}, H(f) \text{ is even}]$$

$$h(t) \text{ is real and even}$$

$$h(t) \text{ is real and even}$$

$$h(t) \text{ is imaginary and even}$$

$$h(t) \text{ is real and odd}$$

$$H(f) \text{ is real and odd}$$

$$H(f) \text{ is real and odd}$$

$$F(q * f) = F(q).F(f),$$
  $q * f = F^{-1}(F(q).F(f))$ 

За дискретни данни  $h_k$ :

$$H_n = \sum_{k=0}^{N-1} h_k e^{2\pi i k n/N}, h_k = \frac{1}{N} \sum_{n=0}^{N-1} H_n e^{-2\pi i k n/N}, \sum_{k=0}^{N-1} |h_k|^2 = \frac{1}{N} \sum_{n=0}^{N-1} |H_n|^2 = \frac{1}{N} \sum_{$$

Колко опрации са нужни за дискретното преобразувание на Фурие  $H_n = \sum_{k=0}^{N-1} W^{nk} h_k$ , където  $W \equiv e^{2\pi i/N}$ ? Може би  $\sim N^2$ ...

Колко опрации са нужни за дискретното преобразувание на Фурие  $H_n = \sum_{k=0}^{N-1} W^{nk} h_k$ , където  $W \equiv e^{2\pi i/N}$ ? Може би  $\sim N^2$ ... Оказва се, че това може да стане с $\sim N \log_2 N$  операции. Danielson и Lanczos (1942) показват, че дискретното преобразование на Фурие с дължина N може да се запише като **сума от две** дискретни преобразования на Фурие, всяко с дължина N/2. Едното е преобразуванието на стойностите, който стоят на четни позиции, а другото - на нечетните:

$$F_k = \sum_{i=0}^{N-1} W^{ij} f_j = \dots = F_k^e + W^k F_k^o$$

# Спекрална теория на графи

ENCYCLOPEDIA OF MATHEMATICS AND ITS APPLICATIONS

### Eigenspaces of graphs

D. Cvetković University of Belgrade

P. Rowlinson University of Stirling

S. Simić University of Belgrade Conference Board of the Mathematical Sciences

# CBMS

Regional Conference Series in Mathematics

Number 92

Spectral Graph Theory

Fan R. K. Chung

Published for the Conference Board of the Mathematical Sciences



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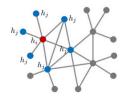
# Graph Laplacian

Core operator in Spectral Graph Theory

$$\mathcal{G} = (V, E, A) \longrightarrow \underset{n \times n}{\Delta} = I - D^{-1/2} A D^{-1/2} \qquad \text{Normalized Laplacian}$$
 where 
$$\underset{n \times n}{D} = \operatorname{diag}(\sum_{j \neq i} A_{ij})$$

- Interpretation:
  - Measure of smoothness: Difference between local value  $h_i$  and its neighborhood average values  $h_i$ .

$$(\Delta h)_i = h_i - \sum_{j \in \mathcal{N}_i} \frac{1}{\sqrt{d_i d_j}} A_{ij} h_j$$



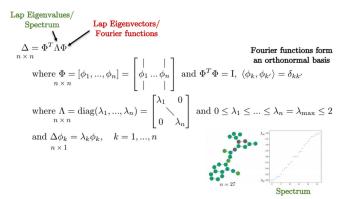
$$\Delta \equiv \nabla^2 = \sum_i \partial_{x_i}^2$$

$$\Delta \equiv \nabla^2 = \sum_i \partial_{x_i}^2$$

- Фурие трансформация на граф се нарича разлагането на матрицата на Лаплас за този граф по собствени вектори (създаващи, т.нар. Фурие базис на граф).

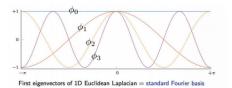
### Fourier Functions

### • Eigen-decomposition of graph Laplacian:

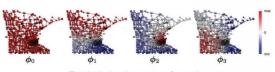


# Fourier Functions

• Grid/Euclidean domain:



Graph domain :



First Laplacian eigenvectors of a graph

Fourier functions related to graph geometry (s.a. communities, hubs, etc) Spectral graph clustering<sup>[1]</sup>

## Fourier Transform

• Fourier series: Decompose function h with Fourier functions<sup>[1]</sup>:

$$\begin{split} h &= \sum_{k=1}^{n} \underbrace{\langle \phi_{k}, h \rangle}_{\mathcal{F}(h)_{k} = \hat{h}_{k} = \phi_{k}^{T} h} \underbrace{\overset{q}{\wedge}_{h} \times 1}_{n \times 1} \\ &= \sum_{k=1}^{n} \hat{h}_{k} \phi_{k} \\ &= \underbrace{\Phi \hat{h}}_{\mathcal{F}^{-1}(\hat{h})} \end{split}$$

Fourier transforms are one line of code (linear operations) 
$$\begin{cases} & \mathcal{F}(h) = \Phi^T h & \text{Fourier Transform/} \\ & \text{coefficients of Fourier Series} \\ & & \\ & \mathcal{F}^{-1}(\hat{h}) = \Phi \hat{h} & \text{Inverse Fourier Transform} \\ & &$$

Invertible transformation

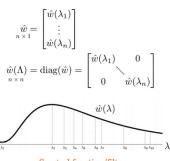
# Конволюция на графи in inverse space

## Convolution Theorem

Fourier transform of the convolution of two functions is the pointwise product of their Fourier transforms:

$$\begin{array}{l} w * h \\ n \times 1 \end{array} = \underbrace{\mathcal{F}^{-1}}_{\Phi} \left( \underbrace{\mathcal{F}(w)}_{\Phi^T w = \hat{w}} \underbrace{\mathcal{F}(h)}_{\Phi^T h} \right) \\ = \underbrace{\Phi}_{n \times n} \left( \underbrace{\hat{w}}_{n \times 1} \odot \underbrace{\Phi^T h}_{n \times 1} \right) \\ = \underbrace{\Phi}_{n \times n} \left( \underbrace{\hat{w}}_{n \times n} A \right) \underbrace{\Phi^T h}_{n \times 1} \\ = \underbrace{\Phi}_{n \times n} \left( \underbrace{\Phi \Lambda \Phi^T}_{\Delta} A \right) h \\ = \underbrace{\hat{w}}_{n \times n} \left( \underbrace{\Phi \Lambda \Phi^T}_{n \times 1} A \right) h \\ = \underbrace{\hat{w}}_{n \times 1} \left( \underbrace{\Phi \Lambda \Phi^T}_{n \times 1} A \right) h \end{array}$$

Expensive computation  $O(n^2)$ No FFT



Spectral function/filter

Notebook 24

```
https:
```

//github.com/xbresson/spectral\_graph\_convnets/blob/
master/02\_graph\_convnet\_lenet5\_mnist\_pytorch.ipynb